

Title: Lecture - Quantum Matter, PHYS 777

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Collection/Series: Quantum Matter (Elective), PHYS 777, March 31 - May 2, 2025

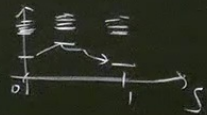
Subject: Condensed Matter

Date: April 10, 2025 - 10:15 AM

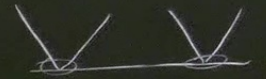
URL: <https://pirsa.org/25040005>

Recap ① Lieb-Robinson bound: $\|[A_i(t), B_j]\| \leq \text{const} \cdot e^{-c|i-j|}$ for $|i-j| \geq v_L t$

② Two gapped $|\psi_1\rangle, |\psi_2\rangle \iff |\psi_2\rangle = U_{\text{eff}} |\psi_1\rangle = \mathcal{T} e^{-i \int_0^t \tilde{H}_{\text{eff}} ds} |\psi_1\rangle$
 $H(s), s \in [0,1]$

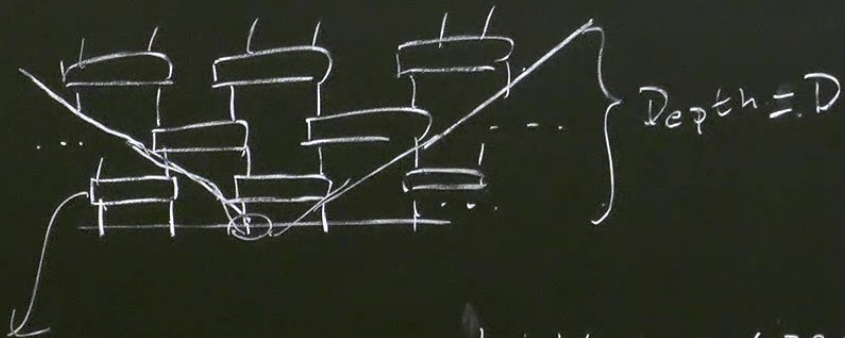


③ U_{eff} can be efficiently simulated using Q.C. with low depth



Quantum Circuit:

Setup:



Unitary, width $\sim l$

Light cone size $\sim D \cdot l$

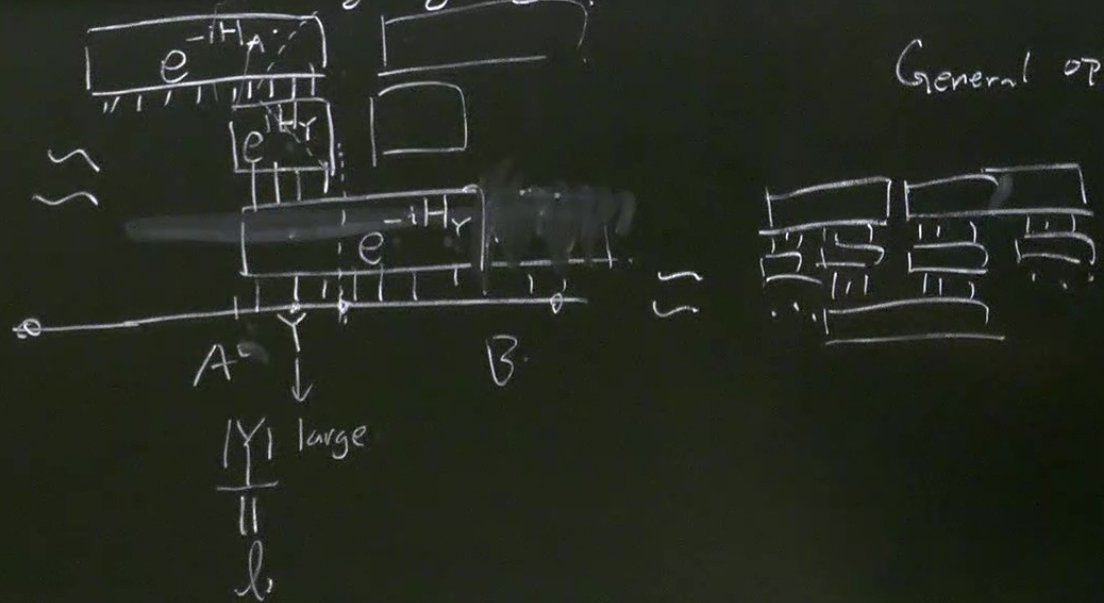
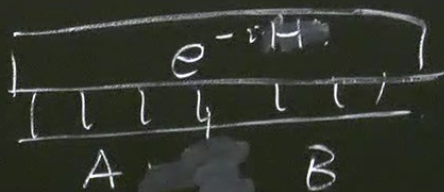
Setup: local $H(s) = \sum_z H_z(s)$, $\|H_z\| = 0$ if $|z| \geq \text{Const.}$, $U = T e^{-i \int \tilde{H}(s) ds}$.
 Thm: U can be simulated in Depth $O(\text{poly} \log \frac{L^{\text{cl}}}{\epsilon})$, $\|V - U\| \leq \epsilon$.

slightly weaker version: $D \sim O(1)$, $l \sim \text{Poly Log}(\frac{L^d}{\epsilon})$

For any local operator

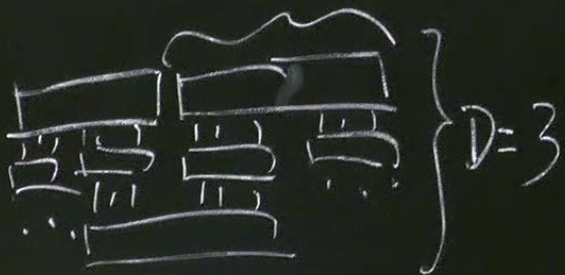
General operator:

Key idea:



For any local operator O_i , effect: e^{-iH} to $e^{-L \cdot C}$.

General operator: $\sum \lambda_a O_1^{(a)} O_2^{(a)} O_3^{(a)} \dots O_L$ error $\underline{L e^{-L \cdot C}}$.



$$U O_1 O_L U^\dagger = U O_1 U^\dagger U O_L U^\dagger$$

$D=3$, Gate size l .

$$\text{Error } L e^{-L \cdot C} \leq \epsilon \Rightarrow l \sim \text{Poly} \log \frac{L}{\epsilon}$$

$$|\psi_1\rangle \simeq |\psi_2\rangle \quad \text{iff} \quad \exists U_{FD} : |\psi_1\rangle = U_{FD} |\psi_2\rangle.$$

- All product states belong to some phase ("trivial phase")

$$|prod\rangle = |a_1\rangle \otimes |a_2\rangle \otimes \dots \otimes |a_L\rangle.$$

$$|0\rangle = \begin{array}{c} \downarrow \quad \downarrow \quad \downarrow \\ |-\rangle \otimes |-\rangle \otimes \dots \otimes |-\rangle \end{array}$$

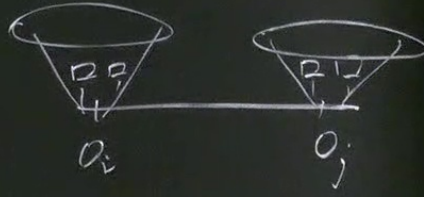
- If $|\psi\rangle = U_{\text{FD}} |0\rangle$ for some U_{FD} , $|\psi\rangle$ is "Short-range entangled" (SRE)

Immediate consequence:

$$\langle O_i O_j \rangle_c = \langle (O_i - \langle O_i \rangle) (O_j - \langle O_j \rangle) \rangle \sim e^{-|i-j|/\xi} \quad (0 \text{ for strictly local gates})$$

for $|i-j| \gg d \cdot D$.

Proof: $\langle \psi | O_i O_j | \psi \rangle = \langle 0 | \underbrace{U_{\text{FD}}^\dagger O_i U_{\text{FD}}}_{\tilde{O}_i} \underbrace{U_{\text{FD}}^\dagger O_j U_{\text{FD}}}_{\tilde{O}_j} | 0 \rangle$



if $|i-j| \gg d \cdot D$

$$= \langle 0 | \tilde{O}_i | 0 \rangle \langle 0 | \tilde{O}_j | 0 \rangle = \langle O_i \rangle_\psi \langle O_j \rangle_\psi$$

If $\exists U_{FP}$ s.t. $|\psi\rangle = U_{FP} |0\rangle$: $|\psi\rangle$ is Long-range entangled (LRE)

Example. Ferromagnet/Cat/GHZ state

$$|\uparrow\uparrow\uparrow\dots\uparrow\rangle + |\downarrow\downarrow\downarrow\dots\downarrow\rangle$$

$$\langle Z_i Z_j \rangle = 1, \quad \langle Z_i \rangle = \langle Z_j \rangle = 0, \quad \langle Z_i Z_j \rangle_c \rightarrow \text{const} \neq 0$$

$$\text{Critical point, } \langle Z_i Z_j \rangle_c \rightarrow \frac{1}{|i-j|^\alpha}$$

long-range entangled (LRE)

Next, weak: Examples of LRE
not from $\langle O_i O_j \rangle_c$.

"Topological order".

$\langle O_i O_j \rangle_c \rightarrow \text{const} \neq 0$

Back to SRE, can be interesting if $\exists$ symmetry

If $|\psi\rangle = U_{FD} |0\rangle$, symmetry G . It may be the case that $\Rightarrow$
individual gates in U_{FD} has to break G .

$$G|\psi\rangle = |\psi\rangle$$

$$G|0\rangle = |0\rangle$$

$$G U_{FD} G^\dagger = U_{FD}$$

$$U_{FD} = \text{[Circuit Diagram]} \rightarrow G u_i G^\dagger \neq u_i$$

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$$U_{FD} = \left(\begin{array}{c} \text{---} \\ \text{---} \\ \text{---} \\ \text{---} \end{array} \right) G u_i G^\dagger \neq u_i$$

$$\downarrow$$

$$|\psi\rangle = p e^{-i}$$

If $|\psi_1\rangle = U_{FD} |\psi_2\rangle$, gates in U_{FD} are symmetric

$\Rightarrow |\psi_1\rangle, |\psi_2\rangle$ belong to same SPT phase.

case that $\Rightarrow |\psi\rangle$ is a (nontrivial) symmetry-protected topological (SPT) state.
 has to break G .

$\downarrow$
 $|\psi\rangle = \rho e^{-i \int_0^1 H(s) ds} |0\rangle$, but $H(s) = \sum_s \underbrace{H_s(s)}_{\text{has to break } G \text{ for some } s}$

Example: Cluster chain:

1d qubit chain, $L = \text{even}$

$\mathbb{Z}_2 \times \mathbb{Z}_2$ symmetry: $X_{\text{even}} = \prod_{i \text{ even}} X_i$ $X_{\text{odd}} = \prod_{i \text{ odd}} X_i$

$$H_{cl} = - \sum_i Z_i X_{i+1} Z_{i+2}$$

vs. trivial chain $H_0 = - \sum_i X_i$

From Tutorial $H_{cl} = U_{FD}^\dagger H_0 U_{FD}$,

$\Rightarrow$ G.S. of H_{cl} is unique & SRE

$\mathbb{Z}_2 \times \mathbb{Z}_2$ symmetric.

$$U_{FD} = \prod_i \underbrace{CZ_{i,i+1}}_{\text{break } \mathbb{Z}_2 \times \mathbb{Z}_2}$$

• If 1d

Imm

Proof

$\Rightarrow$ G.S. of H_{cl} is unique & SRE

break $Z_2 \times Z_2$

$$| \rightarrow \rangle = | + \rangle, | \leftarrow \rangle = | - \rangle$$

$$| \text{Cluster} \rangle = | \uparrow + \uparrow - \downarrow + \downarrow \dots \rangle + | \uparrow - \downarrow + \downarrow - \uparrow \dots \rangle + \dots$$

$$= \sum_{\substack{\{S_j\} \\ j=1 \dots L/2}} \bigotimes_{2j-1} | Z_{2j-1} = S_j \rangle \bigotimes_{2j} | X_{2j} = S_{2j-1} S_{2j+1} \rangle$$

$\rightarrow$ In comparison: $| + + + \dots \rangle = | \uparrow + \uparrow + \downarrow + \downarrow \dots \rangle + | \uparrow + \downarrow + \uparrow + \downarrow \dots \rangle$

Symmetric state is a condensate of domain walls. (KW duality: $\langle \tilde{Z}_i \tilde{Z}_j \rangle \neq 0$)
 What's nontrivial about $| \text{Cluster} \rangle$: domain walls are "decorated" with charge of X_{even} of Z_{odd}