

Title: Topological orders and topological quantum computing - Lecture 20230309

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Collection: Mini introductory course on topological orders and topological quantum computing

Date: March 09, 2023 - 10:00 AM

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Abstract: In this mini course, I shall introduce the basic concepts in 2D topological orders by studying simple models of topological orders and then introduce topological quantum computing based on Fibonacci anyons. Here is the (not perfectly ordered) syllabus.

Overview of topological phases of matter

Z₂ toric code model: the simplest model of 2D topological orders

Quick generalization to the quantum double model

Anyons, topological entanglement entropy, S and T matrices

Fusion and braiding of anyons: quantum dimensions, pentagon and hexagon identities

Fibonacci anyons

Topological quantum computing.

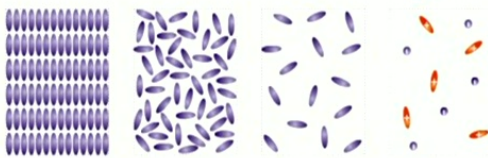
An Over brief introduction to topological phases

TOPOLOGICALLY ORDERED MATTER PHASES

Landau-Ginzburg theory of symmetry breaking
order



temperature



Solid

Liquid

Gas

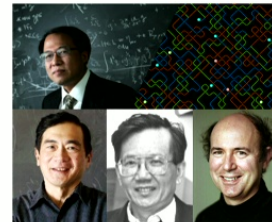
Plasma



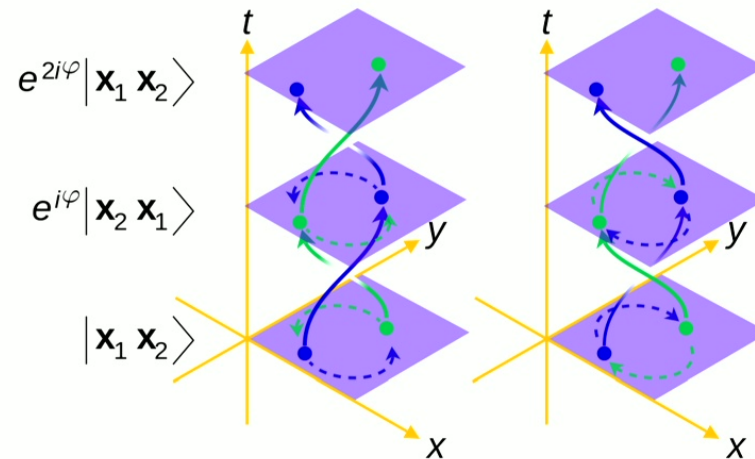
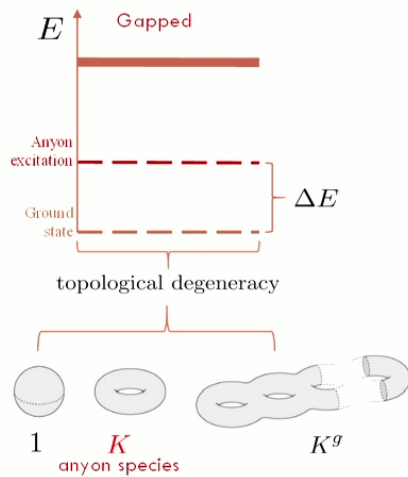
Distinct fractional quantum Hall states share the same symmetry

Landau-Ginzburg symmetry breaking **fails**

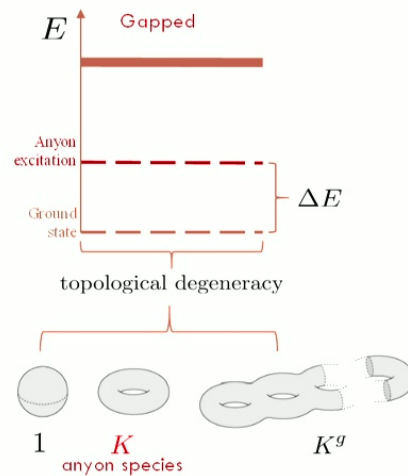
Topological orders



TOPOLOGICAL ORDER IN 2D



TOPOLOGICAL ORDER IN 2D



$\nu = \frac{1}{m}$ FQHS Abelian topological order

- Landau levels.
- $K = m$.
- Anyons species: $\{1, a_1, \dots, a_{m-1}\}$
- Self-Statistics:

$$a_j a_j |\Psi\rangle = e^{i\pi \frac{2}{m}} |\Psi\rangle$$

Topological spin

- Braiding (Aharonov-Bohm):

$$a_j a_k |\Psi\rangle = e^{i\pi \frac{jk}{m}} |\Psi\rangle$$

- Fusion

$$a_j \times a_k = a_{j+k} \mod m$$

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anyon species

$$a_j \times a_k = a_{j+k} \pmod m$$

FIBONACCI TOPOLOGICAL ORDER

I



GSD = 2

Ground states as noncontractible loops

Non-Abelian topological order

Anyon a	1	τ
Quantum dimension d_a	1	$\frac{1+\sqrt{5}}{2}$
Topological spins h_a	0	$2/5$

Fusion : $\tau \times 1 = \tau$,

 $\tau \times \tau = 1 + \tau$ 

$\tau \times \tau \times \tau = (1 + \tau) \times \tau = \tau + \tau + 1$ 

$\dim \mathcal{H}(n \tau\text{'s}) = \# \text{ of outcomes in fusing } n \tau\text{'s}$

$$\dim \mathcal{H} = \begin{pmatrix} 0 & \tau & 2\tau & 3\tau & 4\tau & \cdots & n\tau \\ 1 & 1 & 2 & 3 & 5 & \cdots & d_\tau^{n-2} \end{pmatrix}_{n \rightarrow \infty}$$

Non-local Hilbert space

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Non-local Hilbert space

COLLOQUIUM @ IQC

A diagram showing two vertical arrows pointing down to the labels $|0\rangle$ and $|1\rangle$. A horizontal line with a small vertical tick in the center is positioned below these labels, and the word "qubit" is written in red below the line.

Diagram illustrating a full 1-qubit operation using a quantum circuit with two qubits, σ_1 and σ_2 . The circuit consists of two stages. Stage 1: A CNOT gate with σ_1 as control and σ_2 as target, followed by a rotation gate on σ_1 . Stage 2: A CNOT gate with σ_2 as control and σ_1 as target, followed by a rotation gate on σ_2 . The final state is shown as a product of two single-qubit states.

$$\begin{pmatrix} e^{-i4\pi/5} & 0 \\ 0 & e^{i3\pi/5} \end{pmatrix} \left(\begin{array}{cc} -\frac{e^{-i\pi/5}}{d_\tau} & -\frac{ie^{-i\pi/10}}{\sqrt{d_\tau}} \\ -\frac{ie^{-i\pi/10}}{\sqrt{d_\tau}} & -\frac{1}{d_\tau} \end{array} \right)$$

Full 1-qubit operation

 $su(2)_3/\mathbb{Z}_2$ MTC

$su(2)_3/\mathbb{Z}_2$ MTC

COLLOQUIUM @ IQC

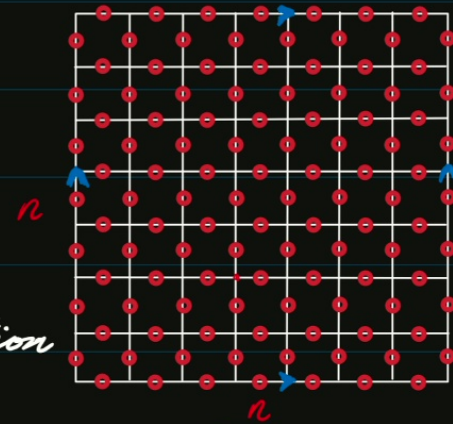
The simplest Lattice Model of 2D Topological Phases

$\mathbb{Z}_2$ toric code model

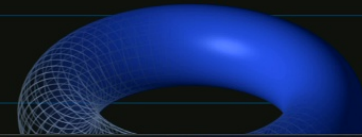
- The set-up

- The playground

a square lattice with periodic boundary condition



- The input data (degrees of freedom)



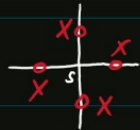
- The output data

doubled $\mathbb{Z}_2$ topological phase (to be explained)

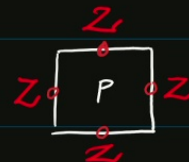
- The Hamiltonian

$$H = -\sum_s A_s - \sum_p B_p$$

$$A_s = \bigotimes_{r \in S} X_r$$



$$B_p = \bigotimes_{r \in P} Z_r$$



$$Z|0\rangle = |0\rangle, \quad Z|1\rangle = -|1\rangle$$

$$X|0\rangle = |1\rangle, \quad X|1\rangle = |0\rangle$$

$\Rightarrow A_s$ & B_p have eigenvalues $1, -1$

$$\text{Clearly, } [A_s, A_{s'}] = [B_p, B_{p'}] = [A_s, B_p] = 0$$



What are these operators
 A_s & B_p ?

operators A_s & B_p :

$$A_s |\psi\rangle = B_p |\psi\rangle = |\psi\rangle, \quad \forall s, p$$

Do you remember how to construct such ground states?

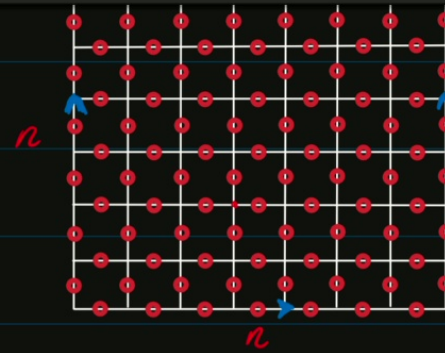
Use projectors! We can however be a bit smarter than naively applying all possible projectors.

- Consider the trivial product state $|0\rangle^{\otimes N}$ $\leftarrow$ # of spins $2n^2$

Obviously,

$$B_p |0\rangle^{\otimes N} = |0\rangle^{\otimes N} \quad \forall p \quad \text{already a } +1 \text{ eigenstate of all } B_p.$$

▷ total # d.o.f.?



$2n^2$ spins. Are they all independent?

▷ Total # constraints on ground states

$$A_s |g.s.\rangle = |g.s.\rangle \quad \forall s \Rightarrow n^2 \text{ constraints}$$

$$+ \quad B_p |g.s.\rangle = |g.s.\rangle \quad \forall p \Rightarrow n^2 \text{ constraints}$$

$2n^2$ spins. Are they all independent?

▷ Total # constraints on ground states

$$A_s |g.s.\rangle = |g.s.\rangle \quad \forall s \Rightarrow n^2 \text{ constraints}$$

$$+ \underline{B_p |g.s.\rangle = |g.s.\rangle \quad \forall p \Rightarrow n^2 \text{ constraints}}$$

$$= 2n^2 \quad \text{Are we done?}$$

▷ Total # independent constraints

$$\prod_s A_s \equiv 1 \quad \prod_p B_p \equiv 1 \quad \because \text{torus}$$

$\Rightarrow 2n^2 - 2$ independent constraints

⇒ $2n^2 - 2$ independent constraints

▷ Total # independent d.o.f

$$2n^2 - (2n^2 - 2) = 2$$

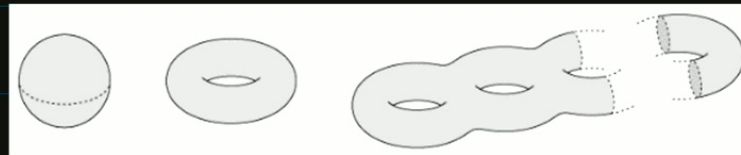
▷ $GSD_{\text{torus}} \equiv 2^2 = 4$ — regardless of the # of spins

That is, as a stabilizer code, 2 logical qubits can be encoded in the ground-state Hilbert space of this model on the torus. This is why the model is named $\mathbb{Z}_2$ toric code.

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Exercise: What is the GSD on the sphere?

▷ In fact



GSD = 1

4

4^g

Exercise: Try $g=2$

the system's size L . [Bravyi & Hastings, J. Math. Phys. 51, 093512 (2010)]

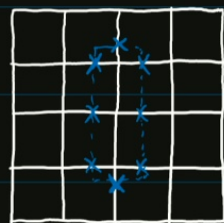
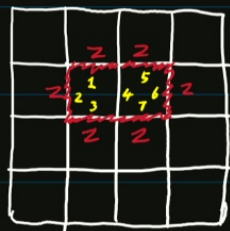


The 2 logical qubits encoded in the ground states of the $\mathbb{Z}_2$ toric code on the torus are **robust** against local perturbations and may be used for **quantum memory** (Kitaev).

But how can we operate on these logical qubits? The answer is related to the other way of counting GSD, i.e.
What's the other way of counting GSD?

◦ Spin- $\frac{1}{2}$ system $\Rightarrow$ $\mathbb{Q}$ must be made of X, Z (Y?)

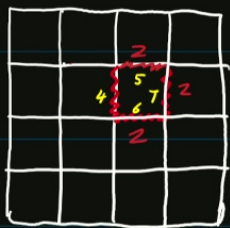
◦ Trivial symmetry generators



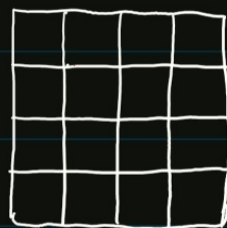
$\therefore$ They don't change a ground state at all.

Why?

$Z_1 Z_2 Z_3 Z_4$

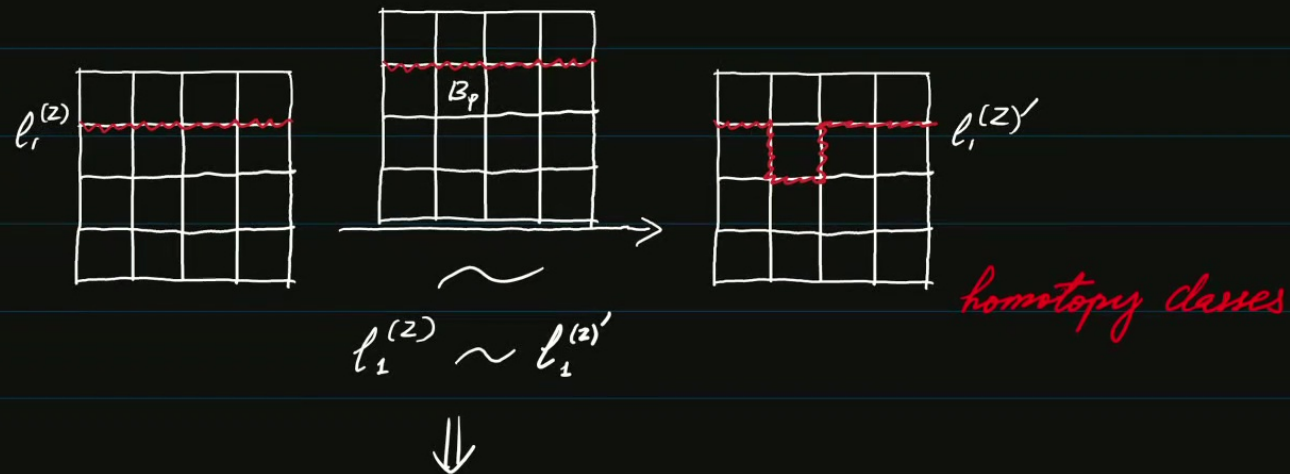


$Z_4 Z_5 Z_6 Z_7$



contractible loops

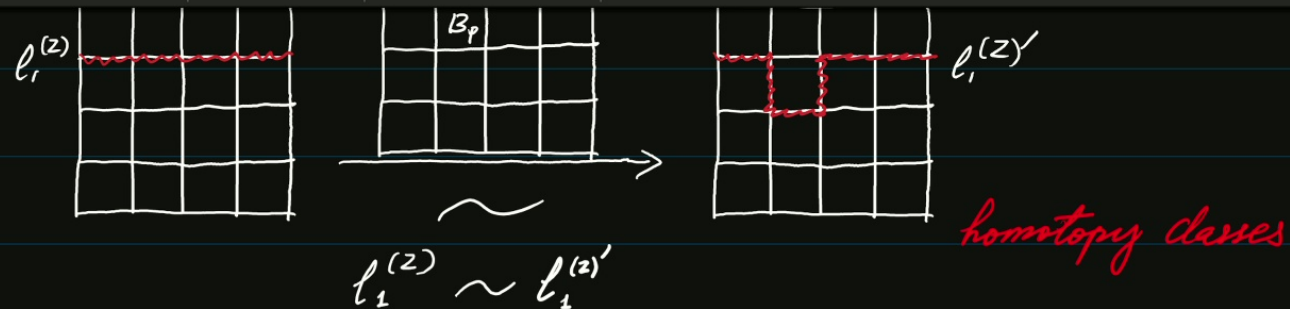
◦ Equivalent loop operators



◦ Need only to consider the four equivalent classes of loops

$l_0^{(z)}, l_1^{(z)}, l_0^{(x)}, l_1^{(x)}$

Which are precisely the symmetry generators (linearly independent)



◦ Need only to consider the four equivalent classes of loops

$$l_0^{(z)}, l_1^{(z)}, l_0^{(x)}, l_1^{(x)}$$

Which are precisely the symmetry generators (linearly independent) of the symmetry of the $\mathbb{Z}_2$ toric code model.

◦ Do they all commute?

Maximally, only two of them commute

e.g. $[l_0^{(z)}, l_1^{(x)}] \neq 0$, $[l_0^{(z)}, l_0^{(x)}] = 0$

There are four commuting set:

(1) $l_0^{(z)}, l_1^{(z)}$

Can choose any set to specify a ground-state basis.

(2) $l_0^{(z)}, l_0^{(z)}$

(3) $l_0^{(x)}, l_1^{(z)}$

These operators all have eigenvalues ± 1

(4) $l_1^{(z)}, l_1^{(x)}$



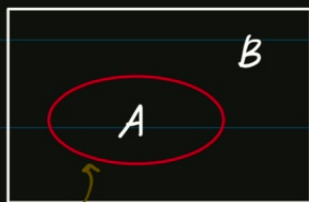
$$GSD = 2^2 = 4$$

- | | | |
|-----|------------------------------|------------------------------|
| (1) | $\ell_0^{(2)}, \ell_1^{(2)}$ | $\ell_1^{(x)}, \ell_0^{(x)}$ |
| (2) | $\ell_0^{(2)}, \ell_0^{(2)}$ | $\ell_1^{(x)}, \ell_1^{(x)}$ |
| (3) | $\ell_0^{(x)}, \ell_1^{(x)}$ | $\ell_1^{(2)}, \ell_0^{(2)}$ |
| (4) | $\ell_1^{(x)}, \ell_1^{(x)}$ | $\ell_0^{(2)}, \ell_0^{(2)}$ |

• Topological entanglement entropy (TEE)

I

- Recall EE:



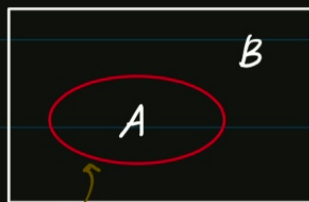
entanglement cut (boundary)

$$\rho_A = \text{Tr}_B \rho_{AB} \quad \text{partial trace}$$

$$S_A = -\text{Tr}(\rho_A \ln \rho_A) \quad \text{von Neumann entropy}$$

measures A's ignorance of B

- Recall EE:



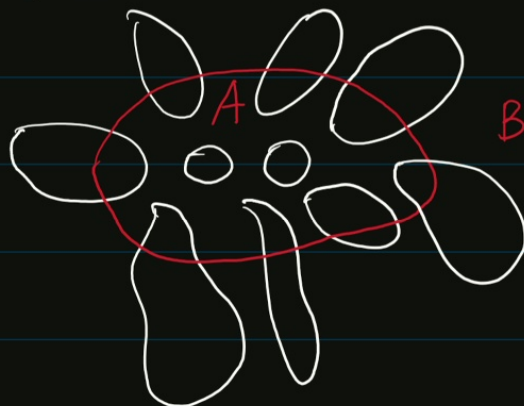
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- TEE



$$S_A = \alpha L - \ln^2 \quad \text{length of the entanglement boundary}$$

determined by anyons

constant γ : TEE

$\gamma = \ln 2$ in $\mathbb{Z}_2$ toric code

$\gamma = \ln D$ in general

total quantum dimension

- If we cut a loop operator open, we create a pair of excitations.

▷ charge (e-type) excitations

e-type/charge



Let $|\psi_0\rangle \in \mathcal{H}_0$

Then, e.g., $Z_1 Z_2 |\psi_0\rangle \notin \mathcal{H}_0$

$$\therefore A_s(Z_1 Z_2 |\psi_0\rangle) \neq Z_1 Z_2 |\psi_0\rangle$$

$$A_{s'}(Z_1 Z_2 |\psi_0\rangle) \neq Z_1 Z_2 |\psi_0\rangle$$

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In fact, $A_s(Z_1 Z_2 |\psi_0\rangle) = -Z_1 Z_2 |\psi_0\rangle$ $\therefore \{A_s, Z_1\} = \{A_{s'}, Z_2\} = 0$

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▷ charge (e-type) excitations

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$A_{s'}(Z_1 Z_2 |\psi_0\rangle) = -Z_1 Z_2 |\psi_0\rangle$

$\Downarrow$

a pair of *charge (e-type) excitations* are created at s & s' .

Why charge?

$$\therefore \{A_s, Z_1\} = \{A_{s'}, Z_2\} = 0$$

$$A_{s'}(Z_1 Z_2 |\psi_0\rangle) = -Z_1 Z_2 |\psi_0\rangle$$



a pair of *charge (e-type) excitations* are created at s & s' .

Why charge?

$Z_1 Z_2$ is an example of *e-string operators*.

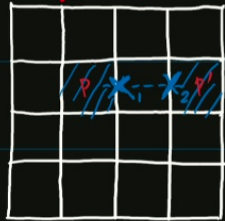
Generically, given any path L along certain edges of the lattice.
an *e-string operator* is defined by

$$W_e = \bigotimes_{r \in L} Z_r$$

which creates a pair of charge excitations at the two ends of L .

▷ Flux (m-type) excitations

m-type/flux



Let $|\psi_0\rangle \in \mathcal{H}_0$

Then, e.g., $X_1 X_2 |\psi_0\rangle \notin \mathcal{H}_0$

$$\therefore B_p(X_1 X_2 |\psi_0\rangle) \neq X_1 X_2 |\psi_0\rangle$$

$$B_{p'}(X_1 X_2 |\psi_0\rangle) \neq X_1 X_2 |\psi_0\rangle$$

In fact, $B_p(X_1 X_2 |\psi_0\rangle) = -X_1 X_2 |\psi_0\rangle$ $\therefore \{B_p, X_1\} = \{B_{p'}, X_2\} = 0$

$$B_{p'}(X_2 X_2 |\psi_0\rangle) = -X_1 X_2 |\psi_0\rangle$$



A pair of flux (m-type) excitations are created at p & p' .

Why flux?

X_1, X_2 is an example of m -type string operators.

Generically, given a path L crossing certain edges of the lattice, an m -string operator is defined by

$$W_m = \bigotimes_{r \in L} X_r$$

- All types of anyons

4 types:

1,	e ,	m ,	$e := em$
<i>vacuum</i> <i>trivial</i>	<i>charge</i>	<i>flux</i>	<i>dyon</i>

$\overset{4}{\# \text{ anyon types}} = \overset{4}{\text{GSD terms}}$

To be proven later

- All types of anyons

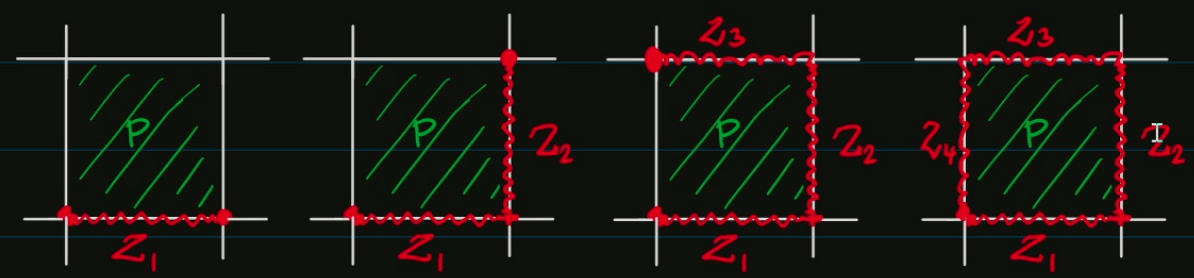
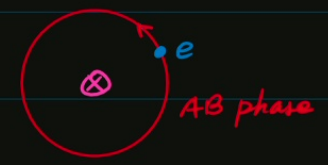
4 types:

1. e , m , $\epsilon := em$
vacuum charge flux dyon
trivial

4 4
anyon types = GSD_{torus}

To be proven later

The names "charge" & "flux" can be justified by
the Aharonov-Bohm effect



by certain m -string operator:

You then create a pair of charges

at s_1 & s_2 by a Z operator:



Show that if you move the charge at s_2 by Z operators around the plaquette as in the diagrams above, until the two charges annihilate, the state acquires a minus sign.

- Fusion

▷ e is created by e -string operator W_e made of Z_i , but $Z^2 = 1$.

$$\therefore e \times e = 1$$

▷ m is created by m -string operator W_m made of X_i , but $X^2 = 1$.

$$m \times m = 1$$

around the plaquette as in the diagrams above, until the two charges annihilate, the state acquires a minus sign.

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$$\therefore m \times m = 1$$

▷ e is the found state of e & m , so $e \times m = e$

▷ Thus, $e \times e = 1$, $e \times m = m$, $m \times e = e$

around the plaquette as in the diagrams above, until the two charges annihilate, the state acquires a minus sign.

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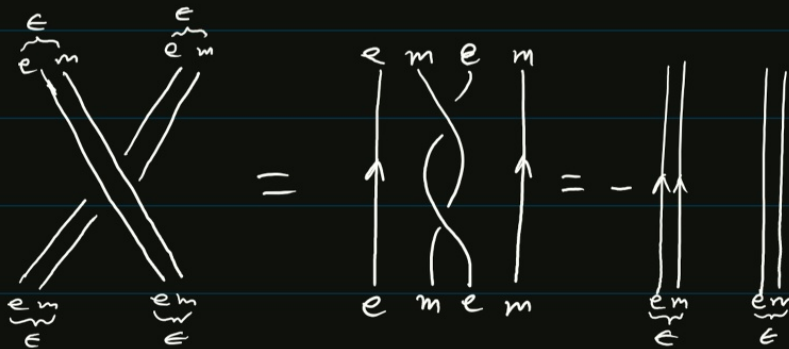
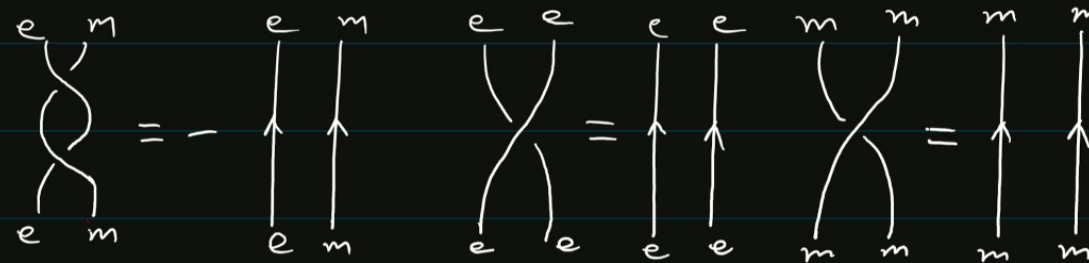
▷ e is the bound state of e & m , so $e \times m = e$

▷ Thus, $e \times e = 1$, $e \times e = m$, $m \times e = e$

▷ Thus, $e \times e = 1$, $e \times m = m$, $m \times e = e$

- Braiding statistics

You may also find:



e, m are self-bosons
but mutually semions
 e is a self-fermion.

These anyons are Abelian anyons, as however you braid them, the state of the system acquires merely a phase (sign) factor.

- The mathematical structure

$$(1, e) \times (1, m)$$

$$D[\mathbb{Z}_2] = \mathbb{Z}_2 \times \mathbb{Z}_2 \quad \text{the quantum double of } \mathbb{Z}_2$$

Replace
 $\mathbb{Z}_2$ by any
finite group G

The Kitaev $\mathbb{Z}_2$ toric code model is
also called the $\mathbb{Z}_2$ quantum double model

- Generalization: The quantum double (QD) model

$\mathbb{Z}_2$ by any
finite group G

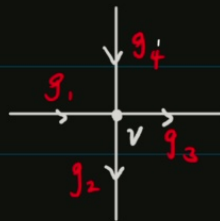
also called the $\mathbb{Z}_2$ quantum double model

• Generalization: The quantum double (QD) model

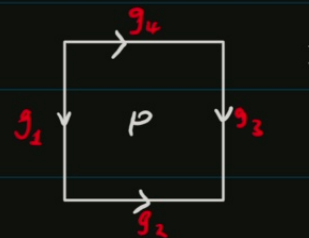
- The Hamiltonian

Let Γ be a square lattice (not necessarily square but for simplicity)

$$H = -\sum_{v \in \Gamma} A_v - \sum_{p \in \Gamma} B_p$$



$g_i \in G$



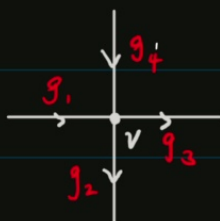
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identity element of G

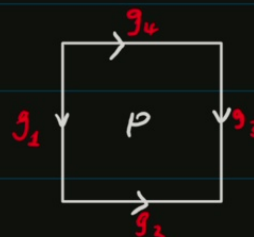
- The Hamiltonian

Let Γ be a square lattice (not necessarily square but for simplicity)

$$H = - \sum_{v \in \Gamma} A_v - \sum_{p \in \Gamma} B_p$$



$$g_i \in G$$



$$A_v = \frac{1}{|G|} \sum_{g \in G} A_v^g$$

$$B_p = \delta_{g_1 g_2 g_3 g_4, e} \quad \text{identity element of } G$$

$$A_v^g |g_1, g_2, g_3, g_4\rangle$$

$$= |g_1 g^{-1}, g g_2, g g_3, g g_4\rangle$$

$$g g_i g^{-1}$$

flatness condition / no-flux

in $\mathbb{Z}_2$ toric code model, this $\mathbb{1}$ is omitted because it only shifts the energy by a constant

E.g. $G = \mathbb{Z}_2$.

$$A_v = \frac{1}{2} (1 + X_1 X_2 X_3 X_4)$$

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in the $\mathbb{Z}_2$ toric code model, this $\mathbb{1}$ is omitted because it only shifts the energy by a constant

- GSD on the torus

For G abelian, $GSD_{\text{torus}} = |G|^2$

For G non-abelian, complicated [PRB 87, 125114 (2013)]

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$\uparrow$ conjugacy classes in G $\uparrow$ centralizer of conjugacy class A .

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$$\forall h, h' \in A, Z^A \cong Z^{h'} \text{ (exercise)}$$

$\Rightarrow$ can label a centralizer by Z^A .

- Anyon types

For G Abelian, $|G|$ G -charges

$|G|$ G -fluxes

$$\Rightarrow |G|^2 \text{ anyon types} = GSD_{\text{torus}}$$

For G non-Abelian, an anyon is labeled by a pair

I

flux $\rightarrow$ charge

(A, μ)

a dyon in general

a conjugacy class

an irrep of Z^A

Why a flux is labeled by a conjugacy class?

Why charge is an irrep of Z^A ?

In a gauge theory, a **gauge flux** is defined (measured) by the **holonomy** of the gauge field around the flux (**A-B phase** in electrodynamics)

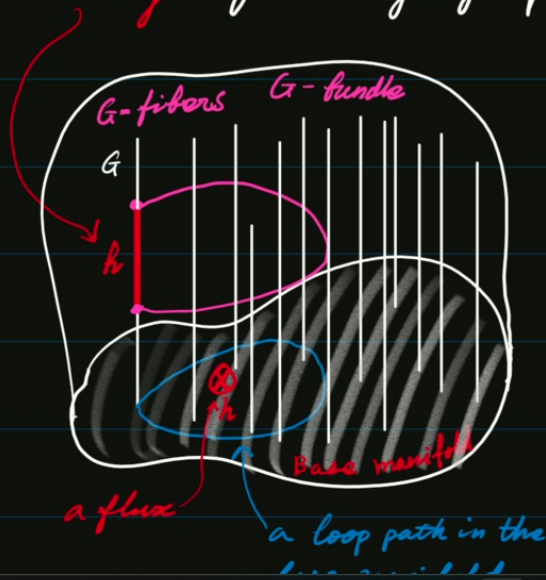
flux $h = \text{holonomy } h \in G$

a physical quantity should be gauge invariant,
i.e., invariant under the gauge transformation

$$h \rightarrow g h g^{-1}$$



Fluxes should take value in **conjugacy**



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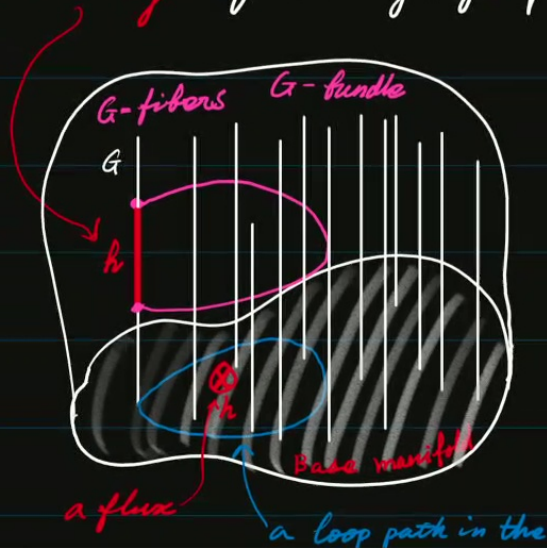
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$\triangleright (e, \mu)$ pure charges

(A, μ) dyon
 ↑ nontrivial

$(A, 1)$ pure fluxes
 ↑ nontrivial conjugacy class
 ↖ trivial rep of Z^A

$\triangleright \# \text{anyon types} = \sum_A \# \text{irrep}(Z^A) = \text{GSD}_{\text{torus}}$

- Mathematical structure

$$\mathcal{D}[G] = \mathcal{C}[G] \times \mathcal{F}[G]$$

group algebra functions of G

take every group
element as a basis vector

$\mathcal{F}[G] = \{f \mid f: G \rightarrow \mathbb{C}\}$ is a vector space

What is the basis of this space?

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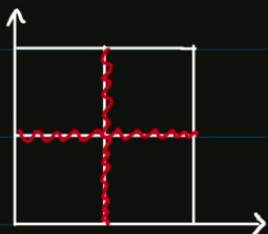
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 What is the basis of this space?

modular transformations



The goal is to turn a basis given by a set of such loops into a basis given by another set of loops.

⇒ need to transform the lattice

In 2D, all such transformations form a group:

$$SL(2, \mathbb{Z}) := (S, T \mid S^4, (ST)^3 S^{-2})$$

integer matrices relations: $S^4 = e, (ST)^3 = S^2$

S & T are the two generators of the group, and in the defining representation are

$$S = \begin{pmatrix} 0 & -1 \\ 1 & 0 \end{pmatrix} \quad T = \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix}$$

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$T:$

Shear

Shear twist

Unit represent the S, T transformation on the ground state

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$\ell_0^{(2)}$

Exercise: Act S & T transformations on the $(\ell_0^{(z)}, \ell_1^{(z)})$ basis to find the 4-dimensional representation of S & T over the ground-state Hilbert space of $\mathbb{Z}_2$ toric code on the torus, i.e., the two 4×4 unitary matrices representing S & T . The answer is:

in the anyon-basis

$$S_{\text{standard}} = \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & 1 & -1 & -1 \\ 1 & -1 & 1 & -1 \\ 1 & -1 & -1 & 1 \end{pmatrix} \begin{matrix} 1 \\ e \\ m \\ \epsilon \end{matrix}, \quad T_{\text{standard}} = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & -1 \end{pmatrix} \begin{matrix} 1 \\ e \\ m \\ \epsilon \end{matrix}$$

Records the mutual statistics of the anyons

Records the self statistics of the anyons

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Hopf link ←

$S_{ab} =$

M_{ab}

$\uparrow t$

records the mutual statistics of anyons

$S_{aa} = T_{aa}^2$

Generically,

$$S = \frac{1}{D} \begin{pmatrix} d_1 & d_2 & d_3 & \dots & d_n \\ d_2 & & & & \\ d_3 & & & & \\ \vdots & & & & \\ d_n & & & & \end{pmatrix} \quad \mathbb{I}$$

Total quantum dimension $D = \sqrt{\sum d_i^2}$ *d_i - quantum dimension of anyon i*

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$$S_{aa} = T_{aa}$$

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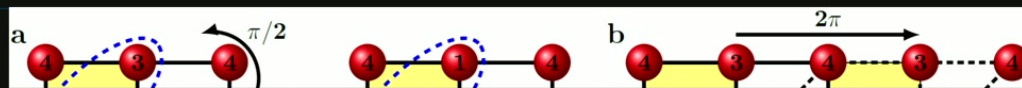
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$d_i \equiv 1$ in the $\mathbb{Z}_2$ -toric code.

S & T matrices are the **finger prints** of topological phases.

I

We can implement S & T transformations on a lattice of spins:

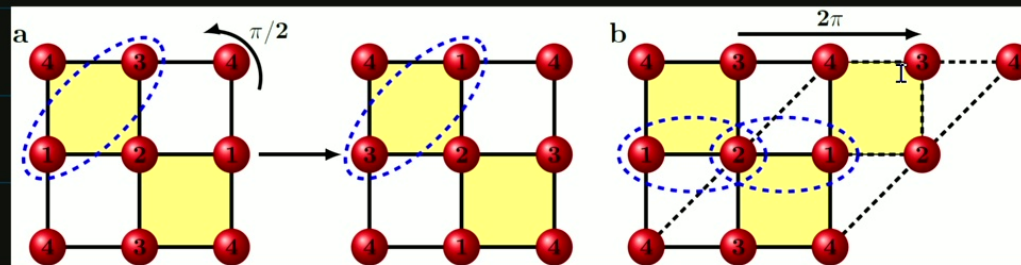


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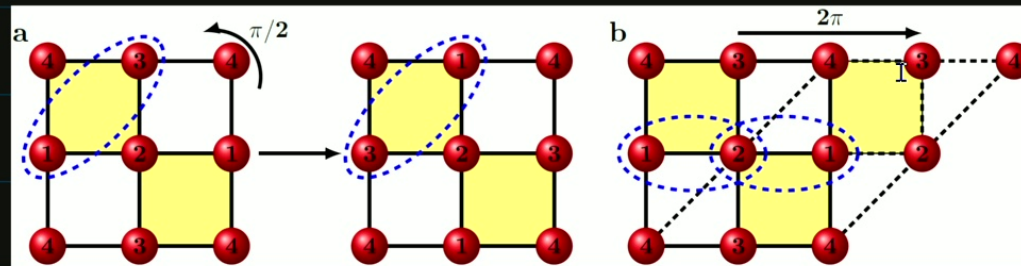
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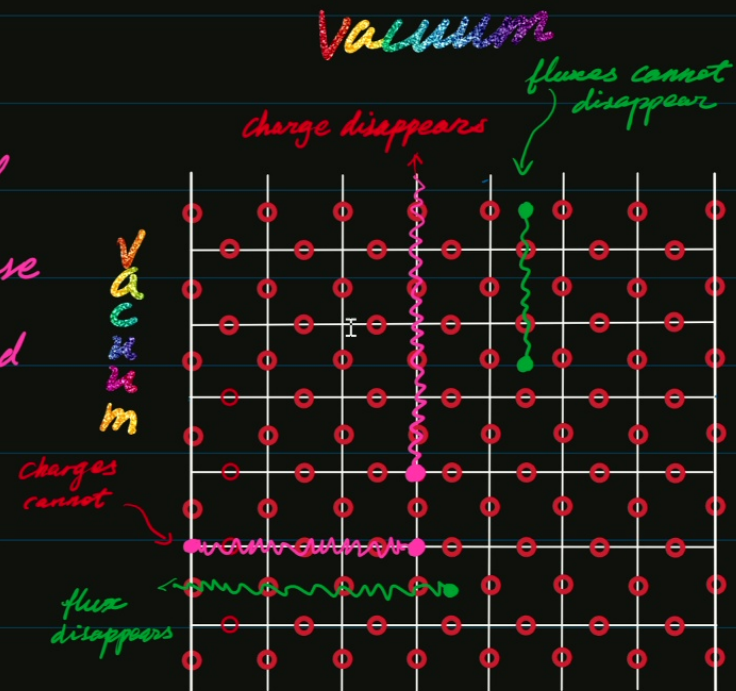
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because charges condense at the boundary & disappear into the vacuum, while fluxes cannot go beyond the boundary.

A smooth boundary is also called a charge- or an e -boundary because fluxes condense at the boundary and disappear into the vacuum, while charges cannot go beyond the boundary.



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beyond the boundary.

A smooth boundary is also called a charge- or an e-boundary because fluxes condense at the boundary and disappear into the vacuum, while charges cannot go beyond the boundary.

Vacuum

charge disappears

fluxes cannot disappear

Vacuum

charges cannot

flux disappears

You see: the trick to creat a logical qubit on the plane is to allow alternate boundaries.

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