

Title: Optimal Thresholds for Fracton Codes and Random Spin Models with Subsystem Symmetry

Speakers: Hao Song

Series: Perimeter Institute Quantum Discussions

Date: March 02, 2022 - 11:00 AM

URL: <https://pirsa.org/22030033>

Abstract: Fracton models provide examples of novel gapped quantum phases of matter that host intrinsically immobile excitations and therefore lie beyond the conventional notion of topological order. Here, we calculate optimal error thresholds for quantum error correcting codes based on fracton models. By mapping the error-correction process for bit-flip and phase-flip noises into novel statistical models with Ising variables and random multi-body couplings, we obtain models that exhibit an unconventional subsystem symmetry instead of a more usual global symmetry. We perform large-scale parallel tempering Monte Carlo simulations to obtain disorder-temperature phase diagrams, which are then used to predict optimal error thresholds for the corresponding fracton code. Remarkably, we found that the X-cube fracton code displays a minimum error threshold (7.5%) that is much higher than 3D topological codes such as the toric code (3.3%), or the color code (1.9%). This result, together with the predicted absence of glass order at the Nishimori line, shows great potential for fracton phases to be used as quantum memory platforms. If time allows, I will also present some of our more recent progress on fractons.

Reference: arXiv:2112.05122.

Zoom Link: <https://pitp.zoom.us/j/97053396111?pwd=Ny9tK295dGVacENJMzg0aHRObjZlZz09>

Optimal Thresholds for Fracton Codes and Random Spin Models with Subsystem Symmetry

HS, J Schönmeier-Kromer, K Liu, O Viyuela, L Pollet, MA Martin-Delgado, arXiv:2112.05122.

Hao Song

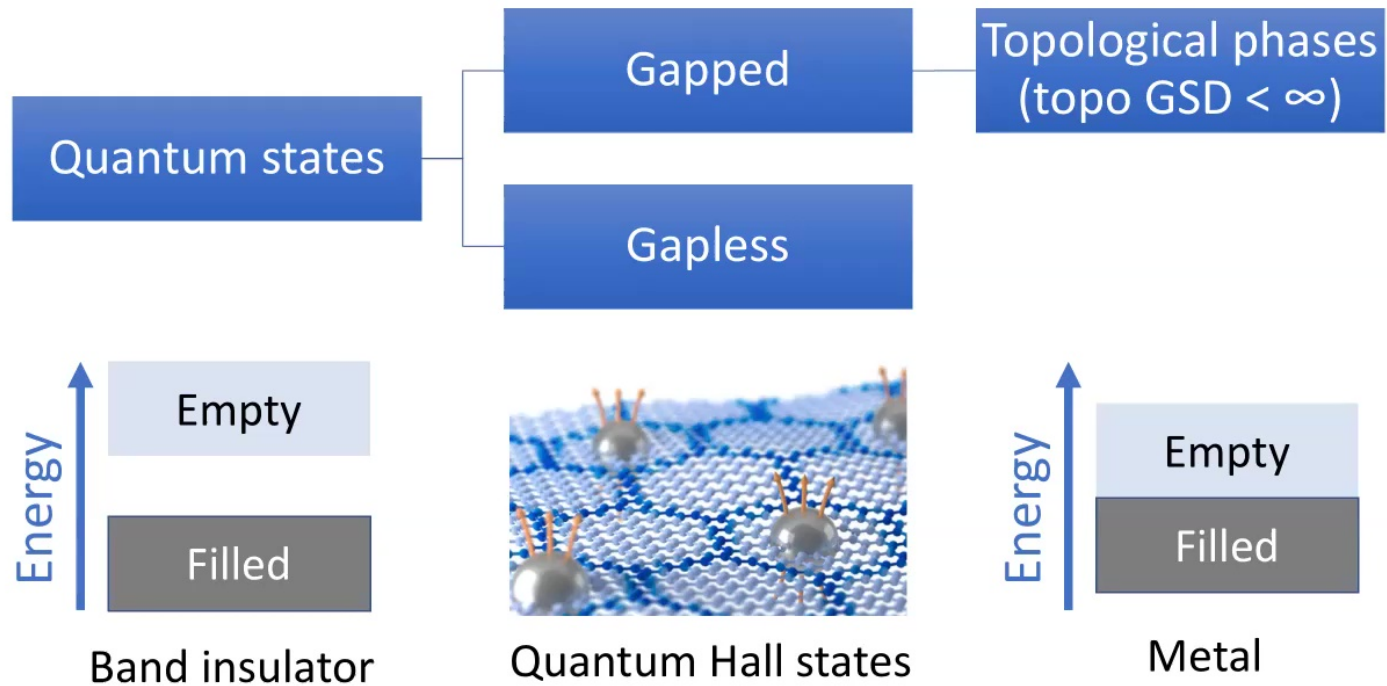
McMaster University

March 2, 2022 @ Perimeter Institute

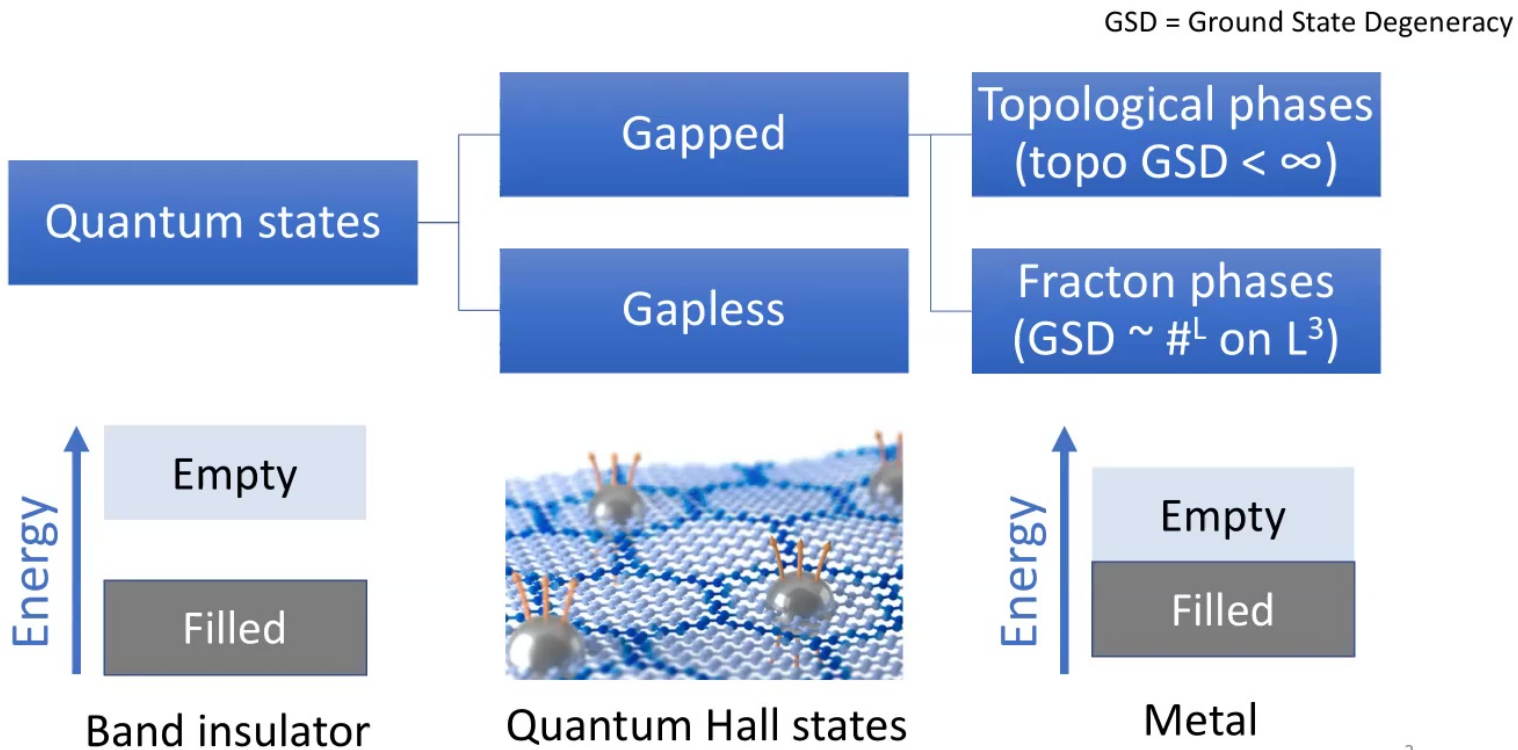
Email: haosongphys@gmail.com



Gapped and gapless states



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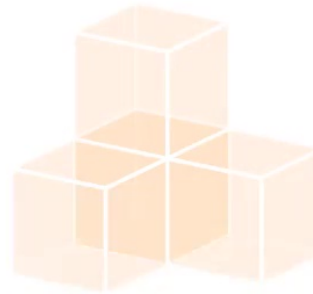
What is a fracton?

Example: X-cube model

- Exactly solvable stabilizer Hamiltonian on cubic lattice (3D).

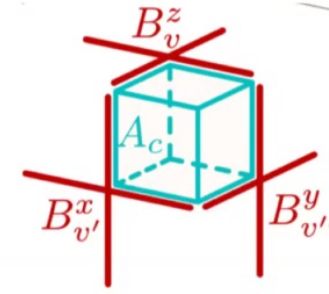
$$H_{\text{Xcube}} = - \sum_c A_c - \sum_v (B_v^x + B_v^y + B_v^z)$$

- Ground states $A_c = B_v^\mu = 1$.
- GSD on 3-torus of size $L \times L \times L$:
 - $\log_2 \text{GSD} = 6L - 3$ (sub-extensive)
 - Gapped
 - Locally indistinguishable
- Similar to topological order but not fully topological



a qubit per edge ℓ
on cubic lattice

$$A_c = \prod_{\ell \in c} X_\ell$$



— = Pauli Z — = Pauli X

$$B_v^\mu = \prod_{\ell \sim v: \ell \perp \mu} Z_\ell$$

3



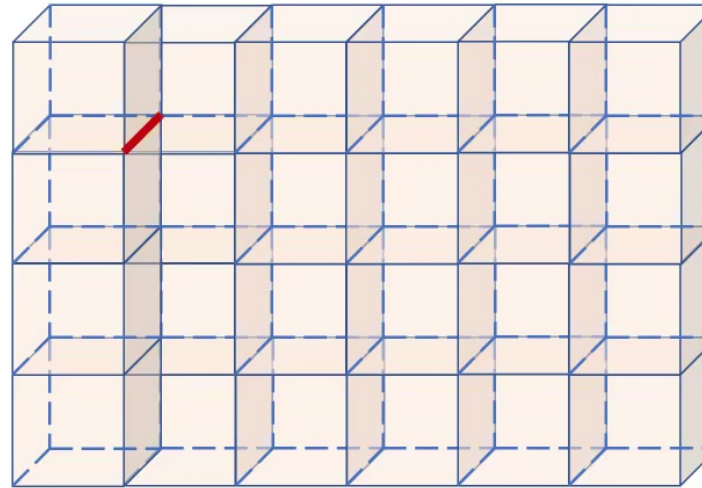
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- cyan cube: $A_c = -1$.
- Isolated A-excitations are created at corners of a membrane operator



$$\text{red line} = Z$$

A fracton is an emergent quasiparticle which fractionalizes into pieces while moving.
It is immobile in the conventional sense!

4



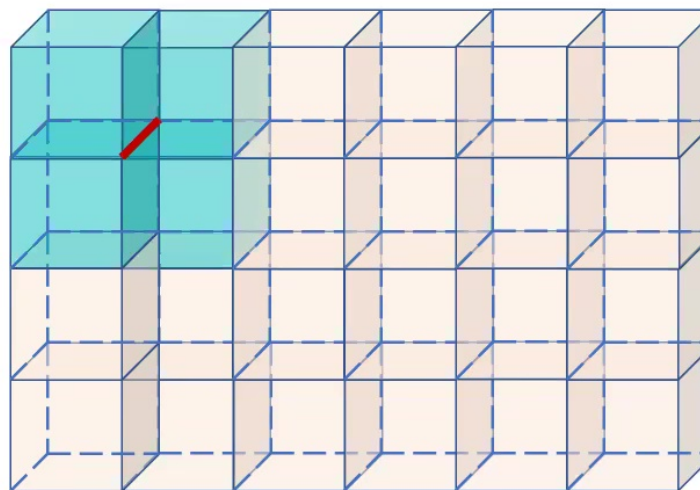
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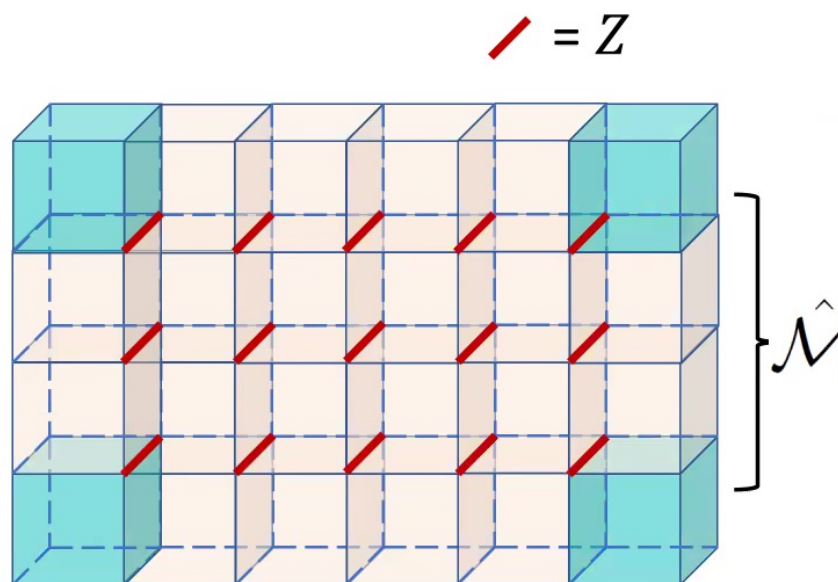
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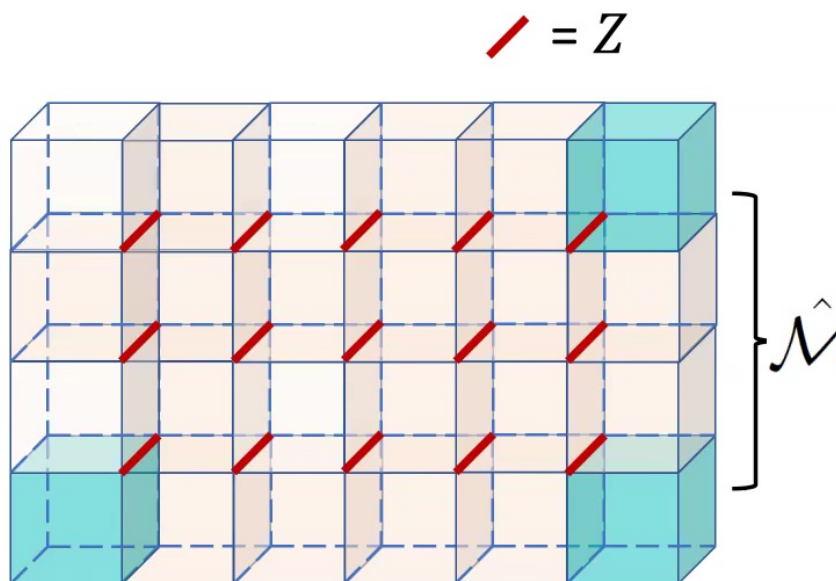
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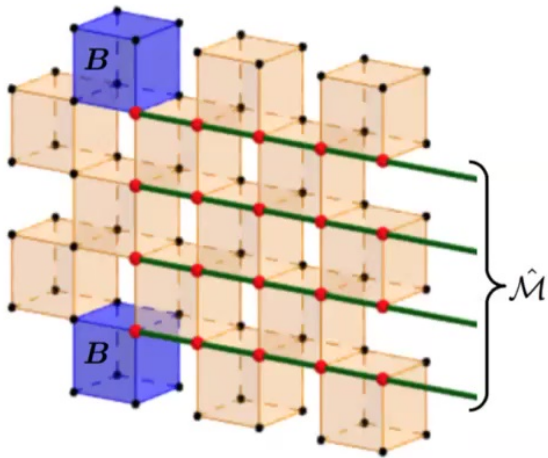
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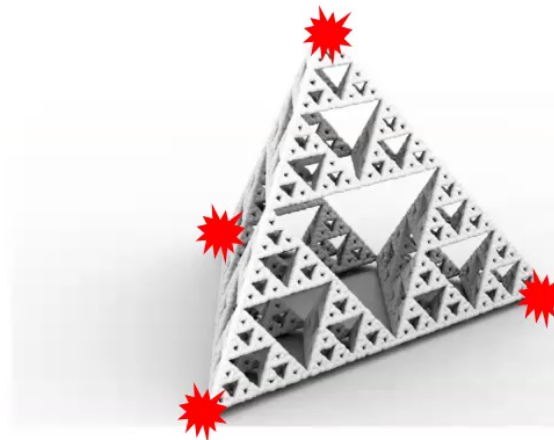


Fracton order (3+1D)

- Unconventional topological order
 - Interesting gapped quantum phases beyond topological quantum field theories? Yes!
 - Self-correcting quantum memories in $d=3$? Not fully realized yet!



Type-I: Chamon, Bravyi,
Leemhuis, Terhal (2005, 2010)
Vijay, Haah, Fu (2015, 2016)



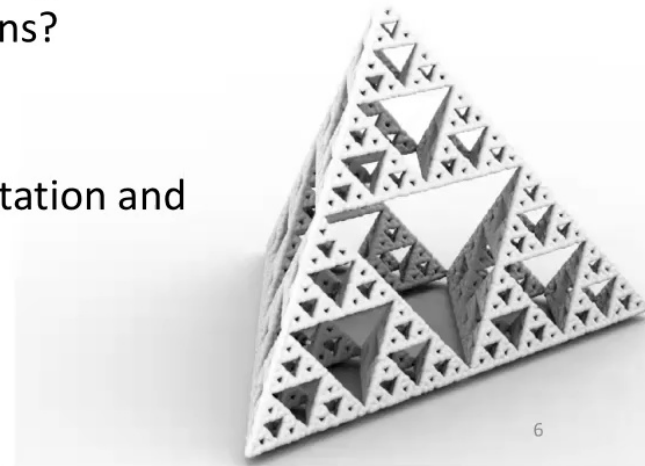
Type-II: Haah (2011)
Yoshida (2013)

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Open questions on fracton Order (3+1D)

- How to properly define such phases in the continuum, RG fixed point?
- A systematic way to distinguish or characterizing fracton phases, like algebraic theory of fractons?
- Its potential applications for quantum computation and error correction?

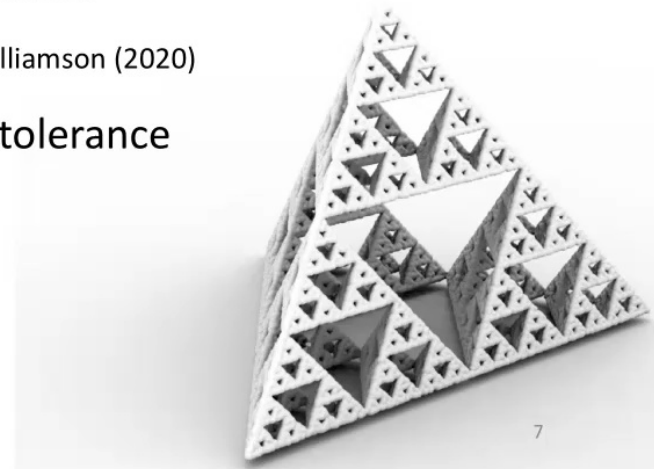


Self-error correction vs active error correction

- Self-error correction: errors are correctable by thermal bath. “Decoder = thermal bath.”
- Various practical active decoders, e.g., RG decoder.
 - Fracton order may allow more efficient decoders

Brown and Williamson (2020)

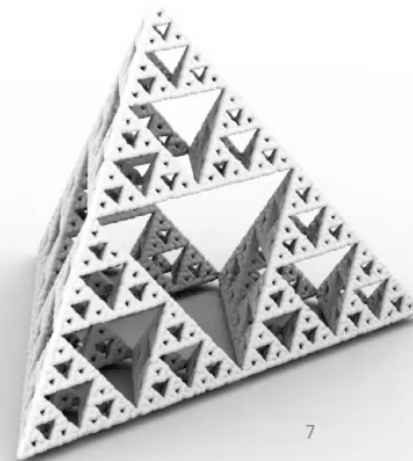
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- Practical decoders may sacrifice some fault tolerance for efficiency
- The theoretical limit of fault tolerance of fracton codes (in dependent of the choice of decoders)?

Brown and Williamson (2020)



Outline

1. Review of quantum error correction in toric code

[Dennis, Kitaev, Landahl, Preskill, 2001]

2. Optimal Thresholds for the X-cube Code

[arXiv:2112.05122]

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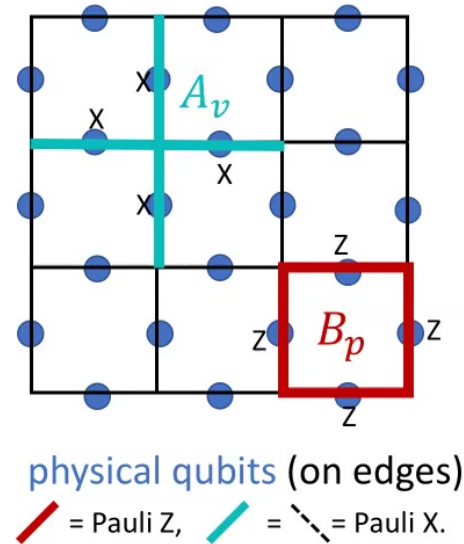
Toric code as a topological quantum memory

- Toric code is a Calderbank-Shor-Steane (CSS) code allows us to treat Z and X errors separately

- Stabilizer generators

$$A_v = X^{\otimes 4}, \quad B_p = Z^{\otimes 4}.$$

- Code space C is selected by $A_v = B_p = 1$.
 - $\dim C = 4 = 2^2$ on torus (i.e. periodic boundary condition).



9



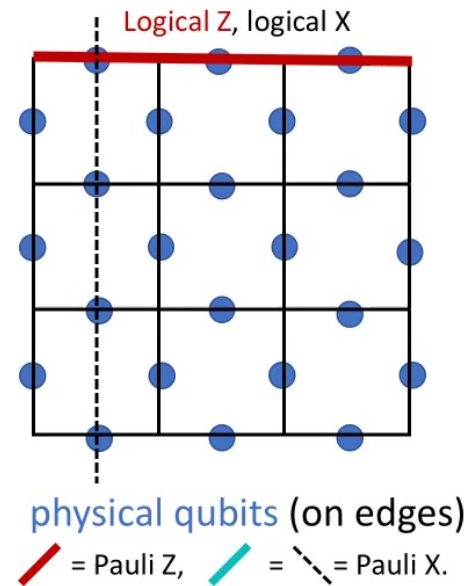
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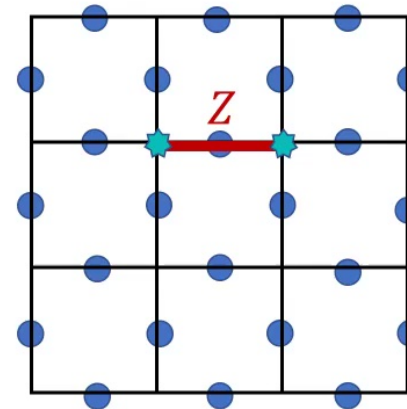
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 - An Z error \rightarrow two flipped A_v (“**A-syndrome**” = the set of flipped A-operators.)



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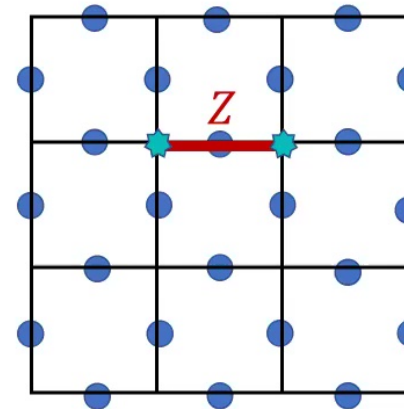
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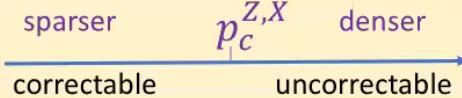
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- **Dense** errors \rightarrow uncorrectable logical errors.



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Suppose Z (or X) errors occur at each qubit independently with probability p .

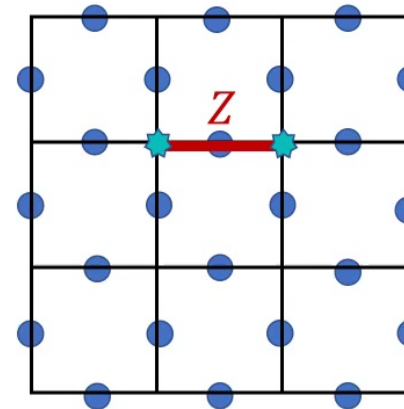




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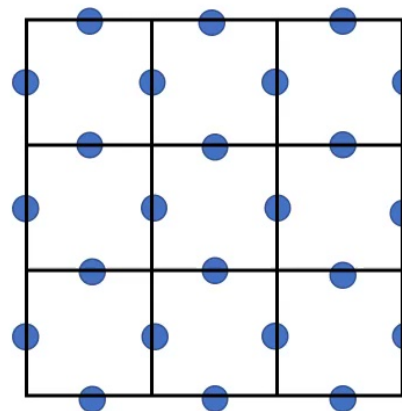
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sparser $p_c^{Z,X}$ denser
— correctable — uncorrectable



Toric code in chain complex terminology

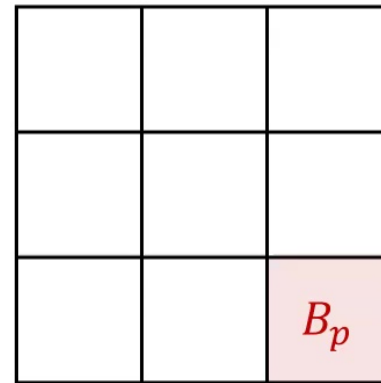


Q = label set of **qubits** (edges)

10



Toric code in chain complex terminology



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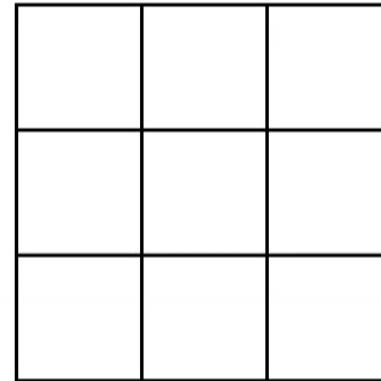
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- Let $\mathbb{Z}_2^{\mathcal{S}} = \text{Fun}(\mathcal{S}, \mathbb{Z}_2 = \{0,1\})$ be the set of \mathbb{Z}_2 -valued functions (i.e., **indicator functions**) on set \mathcal{S} .
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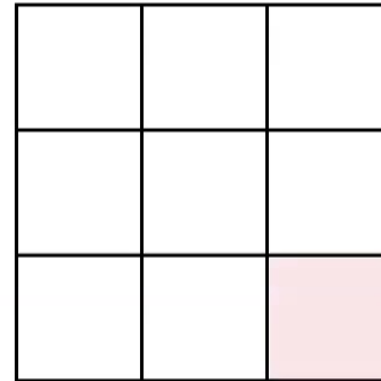


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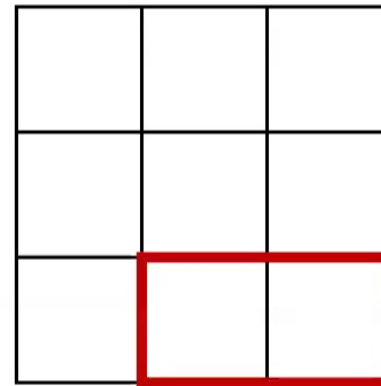



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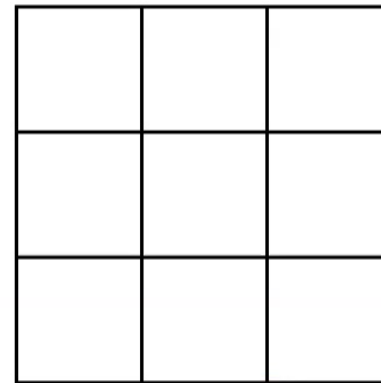



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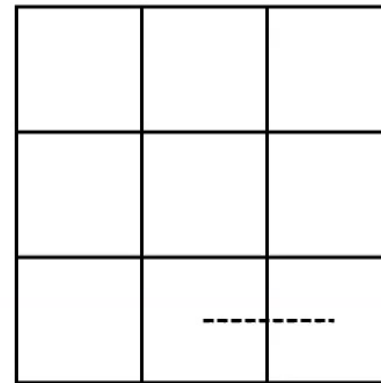


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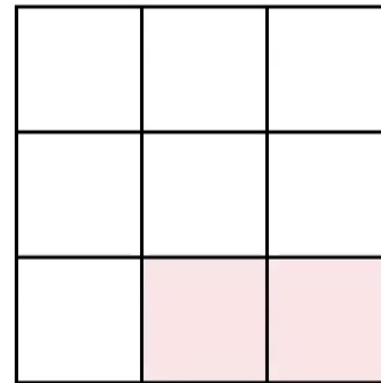


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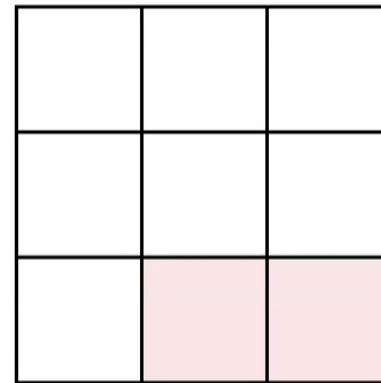


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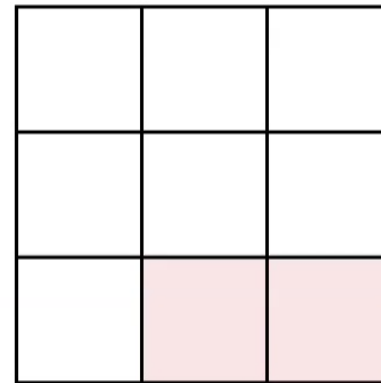


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- Logical Z-operators $\in \ker \partial_A^\dagger$, but $\notin \text{im} \partial_B$.



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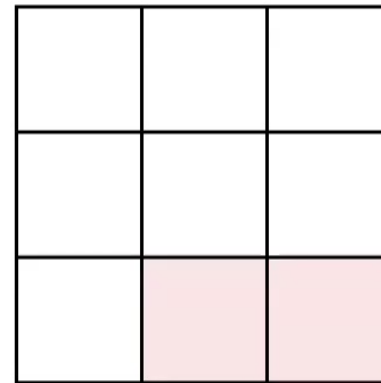


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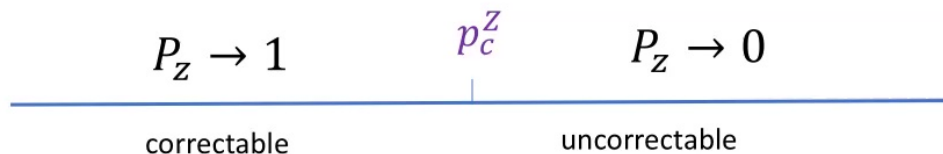
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Maximal success probability in Z-error correction

- All Z-error configurations η, η' acts trivially on the code space, iff $\eta - \eta' \in \text{im} \partial_B$.
Z-error equivalence class $[\eta] = \eta + \text{im} \partial_B \in \mathbb{Z}_2^Q / \text{im} \partial_B$.
- For each possible A-syndrome $\sigma \in \mathbb{Z}_2^A$, there are four compatible Z-error equivalence classes $[\eta_\sigma]$, $[\eta_\sigma] + \mathfrak{z}_1$, $[\eta_\sigma] + \mathfrak{z}_2$, and $[\eta_\sigma] + \mathfrak{z}_1 + \mathfrak{z}_2$.
- Choose η_σ such that $[\eta_\sigma]$ has the largest probability among the four.
- maximal success primality in Z-error correction $P_Z = \sum_\sigma [\eta_\sigma]$.
- Suppose Z errors occur at each qubit independently with probability p .
By analogy to statistical-mechanical models, as system size $\rightarrow \infty$,



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Mapping to statistical-mechanical models

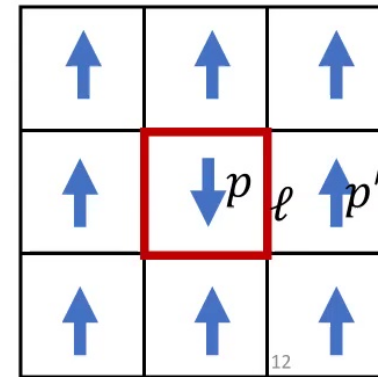
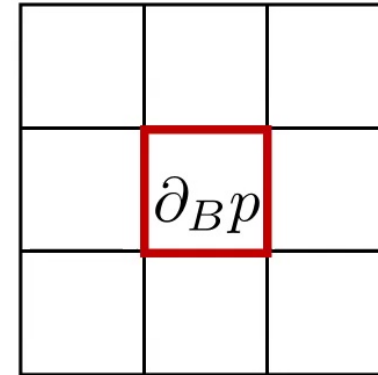
- For Z-error configuration η with $\eta(\ell) = 0$ or 1 on edges with or without an error,

$$\text{pr}(\eta; p) = \prod_{\ell \in \mathcal{Q}} p^{\eta(\ell)} (1-p)^{1-\eta(\ell)} \propto \left(\frac{p}{1-p} \right)^{\sum_{\ell} \eta(\ell)}$$

- Introduce auxiliary temperature $e^{-\frac{2}{T}} = \frac{p}{1-p}$
- The relation is called the **Nishimori line** in the p-T plane.
- The total probability of a Z-error equivalence class $[\eta] = \eta + \text{im}\partial_B \rightarrow$ partition function of random Ising model at T

$$H_{\eta}^{\mathcal{B}} = - \sum_{\ell \in \mathcal{Q}} (-1)^{\eta(\ell)} \prod_{p \in \partial_B^+ \ell} S_p = - \sum_{\langle pp' \rangle} J_{pp'} S_p S_{p'}$$

$J_{pp'} = \pm 1$ with probability $1-p$ and p .



Mapping to statistical-mechanical models

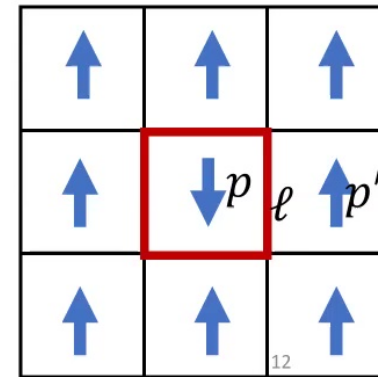
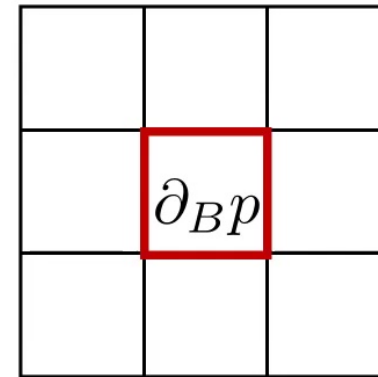
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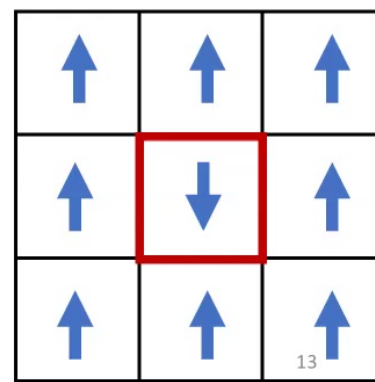
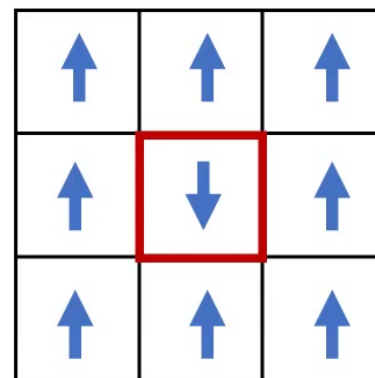
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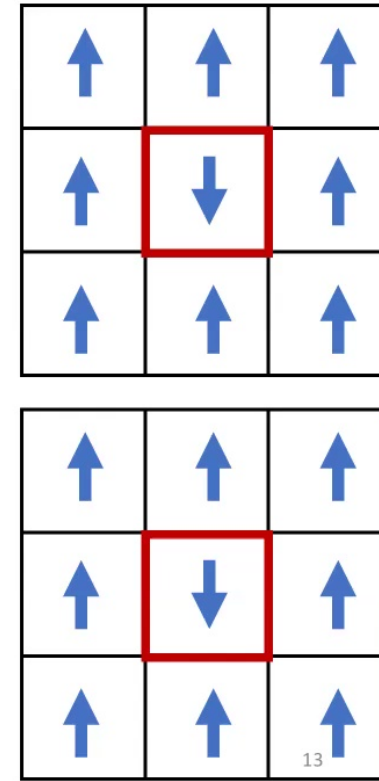
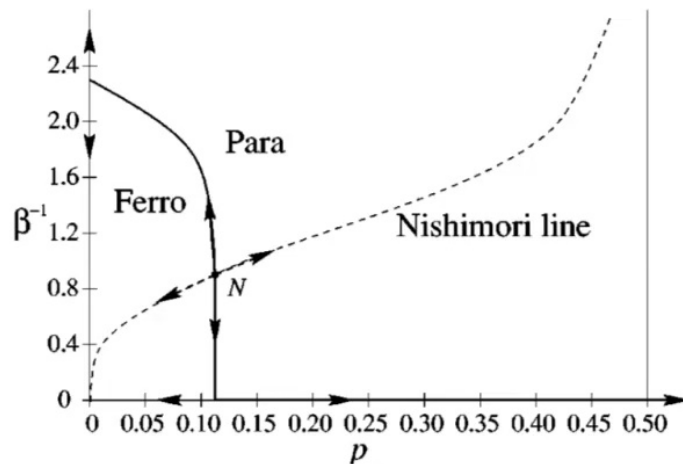
Wisdom from statistical-mechanical models

- Probabilities for the four classes $[\eta]$, $[\eta] + \mathfrak{z}_1$, $[\eta] + \mathfrak{z}_2$, and $[\eta] + \mathfrak{z}_1 + \mathfrak{z}_2$ with the same syndrome differs by a contractible domain wall (on random Ising model side).
 - In the ordered phase, one dominates \rightarrow correctable
 - In the disordered phase, they are comparable \rightarrow uncorrectable



Wisdom from statistical-mechanical models

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 - In the ordered phase, one dominates \rightarrow correctable
 - In the disordered phase, they are comparable \rightarrow uncorrectable



Outline

1. Review of quantum error correction in toric code

[Dennis, Kitaev, Landahl, Preskill, 2001]

2. Optimal Thresholds for the X-cube Code

[arXiv:2112.05122]

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Logical operators in X-cube code

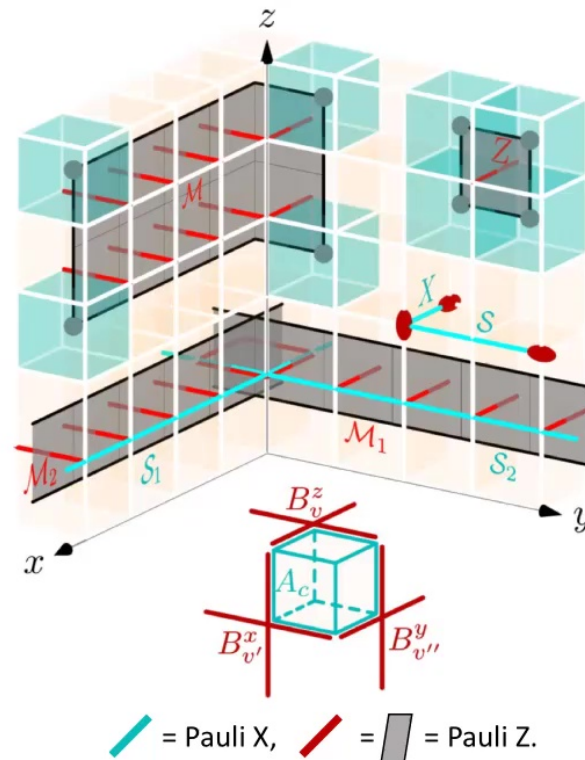
- Toric code is a Calderbank-Shor-Steane (CSS) code allows us to treat Z and X errors separately

- Stabilizer generators

$$A_v = X^{\otimes 12}, \quad B_p = Z^{\otimes 4}.$$

- Code space C is selected by $A_v = B_p = 1$.

- $\dim C = 2^{6L-3}$ on $L \times L \times L$ torus (i.e. periodic boundary condition).
- $6L-3$ logical qubits
- Logical Z operators: $\mathcal{M}_1, \mathcal{M}_2, \dots$
- Logical X operators: $\mathcal{S}_1, \mathcal{S}_2, \dots$



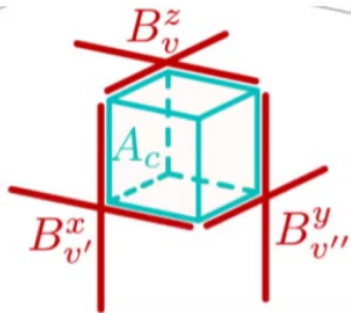
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Mapping to statistical-mechanical models



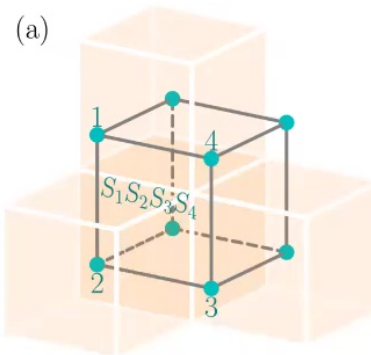
Hao Song



X-cube fracton code
(on a cubic lattice with a qubit per edge)

Suppose X (and Z) errors occur independently at each qubit with probability p .

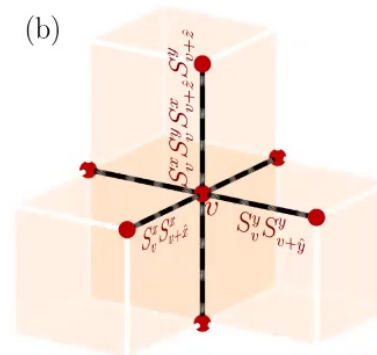
$$\text{pr}(\eta; p) = \prod_{\ell \in Q} p^{\eta(\ell)} (1-p)^{1-\eta(\ell)} \propto \left(\frac{p}{1-p} \right)^{\sum_{\ell} \eta(\ell)}$$



Random Plaquette Ising model

$$H_{\eta}^A = - \sum_{\ell \in Q} (-1)^{\eta(\ell)} \prod_{c \in \partial_A^{\dagger} \ell} S_c$$

For X-error equivalence class
 $[\eta]_X = \eta + \text{im} \partial_A$



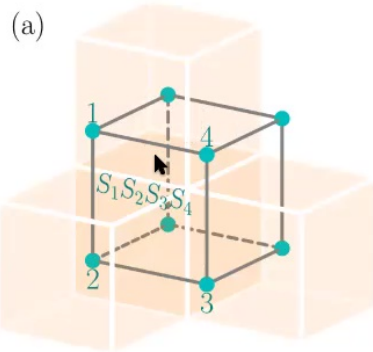
Random Anisotropically
Coupled Ashkin-Teller model

$$H_{\eta}^B(\{S_v^x, S_v^y\}_v) = - \sum_v \sum_{\mu=x,y,z} J_v^{\mu} S_v^{\mu} S_{v+\hat{\mu}}^{\mu},$$

$$S_v^z \equiv S_v^x S_v^y, \quad J_v^{\mu} \equiv (-1)^{\eta((v, v+\hat{\mu}))}$$

For Z-error equivalence class
 $[\eta]_Z = \eta + \text{im} \partial_B$

Planar subsystem symmetry and order parameters



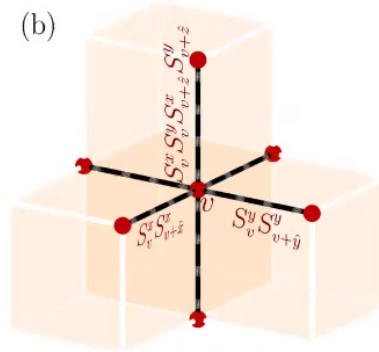
Random Plaquette Ising model

Symmetry: $S_c \rightarrow -S_c$ for c on any xy -, yz -, or zx -plane

$$G^A(\mathbf{r}) = [\langle S_c S_{c+\hat{z}} S_{c+\mathbf{r}} S_{c+\mathbf{r}+\hat{z}} \rangle]$$

$$Q^A := \frac{1}{L^3} \sum_{z=0}^{L-1} \left[\left\langle \left| \sum_{x,y=0}^{L-1} S_{c(x,y,z)} S_{c(x,y,z+1)} \right| \right\rangle \right]$$

planewise dipole moment



Random Anisotropically Coupled Ashkin-Teller model

Symmetry: $S_v^\mu \rightarrow -S_v^\mu$, $S_v^\nu \rightarrow -S_v^\nu$ for v on any μ ν -plane

$$G^B(r) = [\langle S_v^x S_{c+r\hat{x}}^x \rangle]$$

$$Q^B := \frac{1}{L^3} \sum_{x,y=0}^{L-1} \left[\left\langle \left| \sum_{z=0}^{L-1} S_{v(x,y,z)}^z \right| \right\rangle \right]$$

linewise magnetization



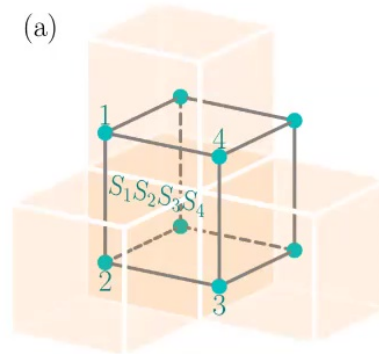
Hao Song



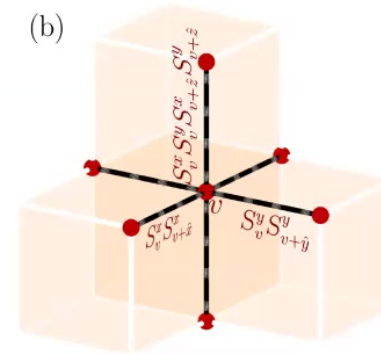
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Phase diagrams (by Monte Carlo simulations)

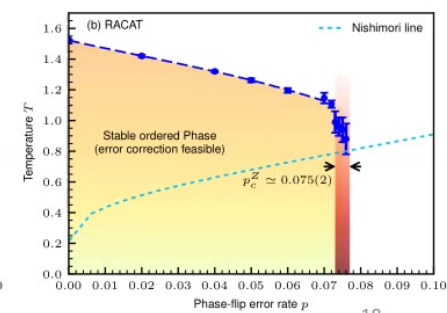
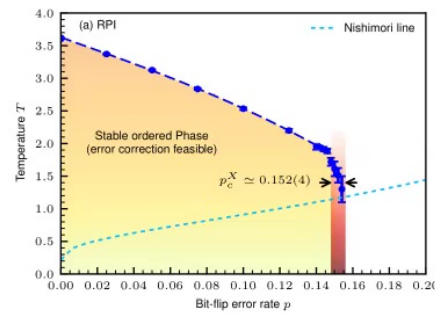
- First-order phase transition for both models at small p . (revealed by energy histogram)
- For larger p , the phase transitions are softened to continuous ones.



Random Plaquette Ising model



Random Anisotropically Coupled Ashkin-Teller model



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Phase diagrams (by N

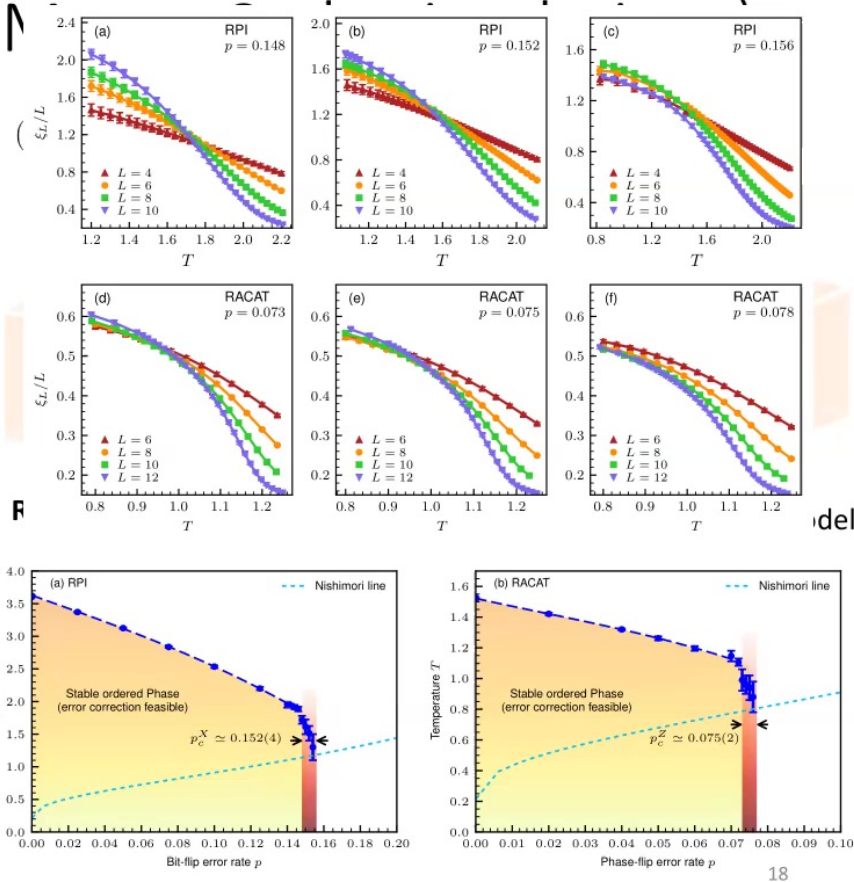
- First-order phase transition for both models at small p . (revealed by energy histogram)

- For larger p , the phase transitions are softened to continuous ones.

- To precisely identify the location of the transition, we study the second-moment correlation length

$$\xi_L := \frac{1}{2 \sin(|\mathbf{k}_{\min}|/2)} \left(\frac{\tilde{G}(\mathbf{0})}{\tilde{G}(\mathbf{k}_{\min})} - 1 \right)^{1/2}$$

- Scaling $\frac{\xi_L}{L} = g \left(L^{-\frac{1}{\nu}} (T - T_c) \right)$
 \rightarrow crossing at $(T_c, g(0))$ if there is a continuous phase transition.



Hao Song

Optimal thresholds of CSS code

- X-cube code (3D)
 - $p_c^X = 15.2\%$, $p_c^Z = 7.5\%$
- 2D toric code & color code
 - $p_c^X = p_c^Z = 10.9\%$

Honecker, Picco, Pujol, 2001

Katzgraber, Bombin, Martin-Delgado, 2009

- 3D toric code
 - $p_c^X = 23.5\%$, $p_c^Z = 3.3\%$

Ozeki and Ito, 1998

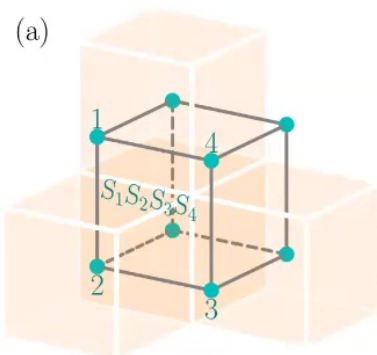
Ohno, Arakawa, Ichinose, Matsui, 2004

- 3D color code
 - $p_c^X = 27.6\%$, $p_c^Z = 1.9\%$

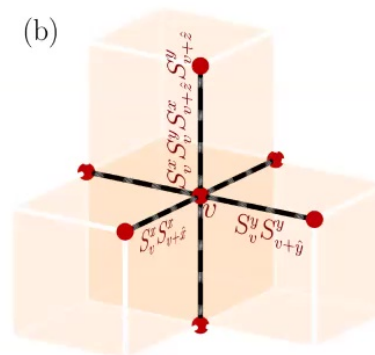
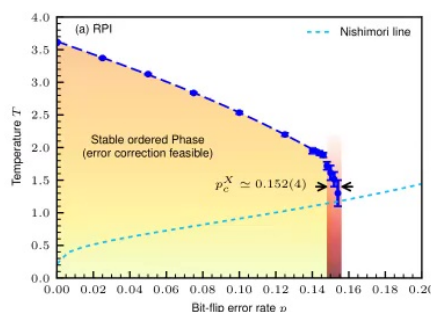
Kubica, Beverland, Brandão, Preskill, Svore, 2018

- Approximate duality (Nishimori 2007) works for the X-cube code:

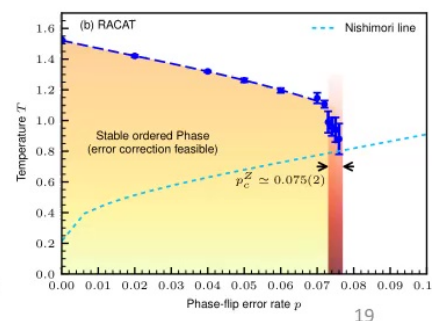
$H(p_c^X) + H(p_c^Z) \approx 1$, where $H(p) = -p \log_2(p) - (1-p) \log_2(1-p)$ is the Shannon entropy.



Random Plaquette Ising model



Random Anisotropically Coupled Ashkin-Teller model



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Hao Song

Absence of glass order along the Nishimori line

- Normal and spin glass order correlation functions

$$G^{\mathcal{A}}(\mathbf{r}) = [\langle S_c S_{c+\hat{z}} S_{c+\mathbf{r}} S_{c+\mathbf{r}+\hat{z}} \rangle]$$

$$G_{\text{SG}}^{\mathcal{A}}(\mathbf{r}) = [\langle S_c S_{c+\hat{z}} S_{c+\mathbf{r}} S_{c+\mathbf{r}+\hat{z}} \rangle^2]$$

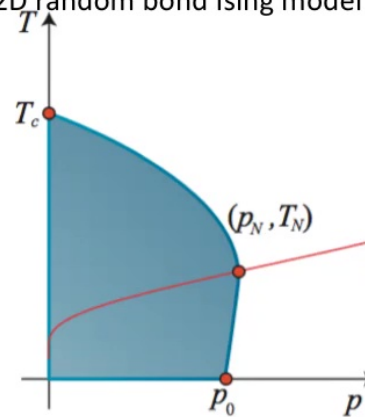
$$G^{\mathcal{B}}(r) = [\langle S_v^x S_{c+r\hat{x}}^x \rangle]$$

$$G_{\text{SG}}^{\mathcal{B}}(r) = [\langle S_v^x S_{c+r\hat{x}}^x \rangle^2]$$

- Glass order

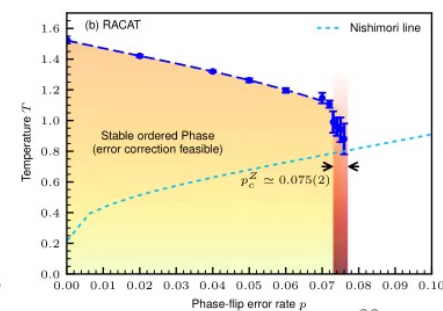
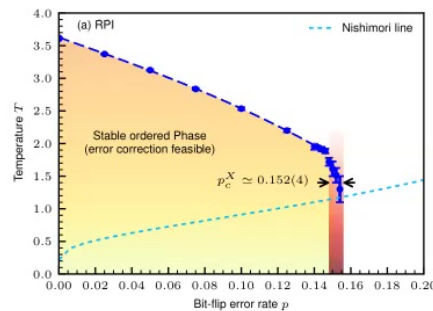
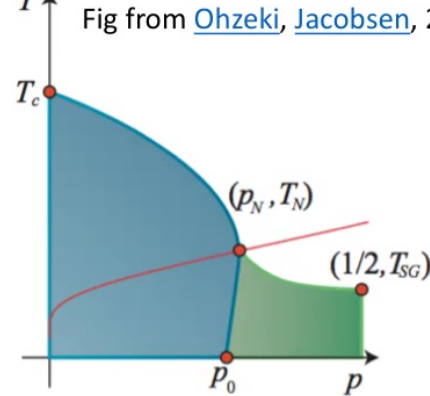
$$\lim_{r \rightarrow \infty} G(r) = 0, \lim_{r \rightarrow \infty} G_{\text{SG}}(r) \neq 0$$

2D random bond Ising model



3D random bond Ising model

Fig from [Ohzeki](#), [Jacobsen](#), 2014



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Absence of glass order along the Nishimori line

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- Glass order

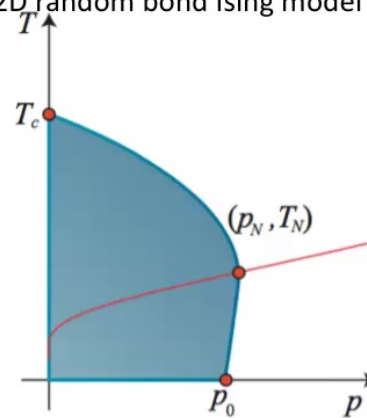
$$\lim_{r \rightarrow \infty} G(r) = 0, \lim_{r \rightarrow \infty} G_{SG}(r) \neq 0$$

- Along the Nishimori line \rightarrow no glass order

- We double-checked Nishimori's argument and showed in general

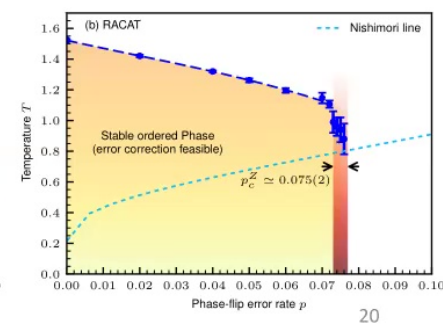
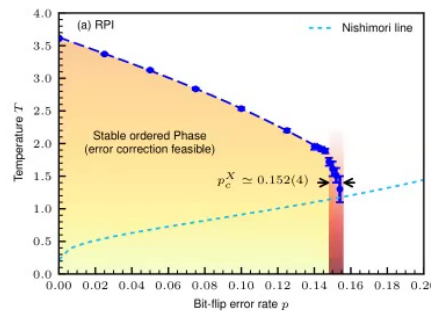
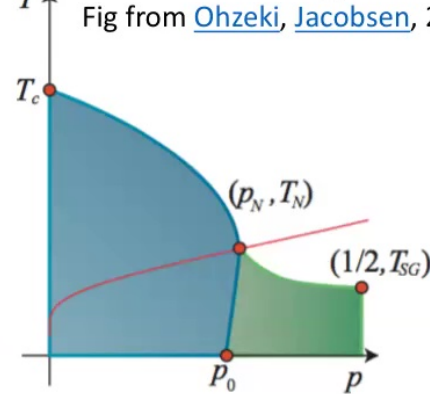
$$G(r) = G_{SG}(r)$$

2D random bond Ising model



3D random bond Ising model

Fig from [Ohzeki](#), [Jacobsen](#), 2014



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Summary

- We make a first study of the optimal thresholds of fracton codes
- For the X-cube model, $p_c^X = 15.2\%$ and $p_c^Z = 7.5\%$ ---higher minimum error threshold (7.5%) than 3D topological codes.
- Random spin models with subsystem symmetry.
- Analytically show no glass order along Nishimori line.
- Numeric suggests no glass order (even below the Nishimori line).
- Approximate duality relation between p_c^X and p_c^Z .

Outlook

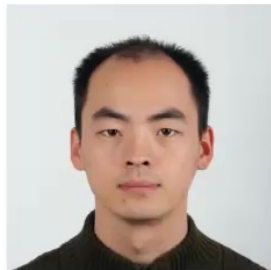
- Haah's code and checkerboard model $\rightarrow p_c^X = p_c^Z \rightarrow 11\%$?
- Measurement errors in fracton codes & high-rank tensor gauge theory in 4D?
- Factonic states for universal quantum computing?



Acknowledgments



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(Munich)



Lode Pollet
(Munich)



Oscar Viyuela
MIT & Harvard



M. A. Martin-
Delgado (Madrid)

HS, J Schönmeier-Kromer, K Liu, O Viyuela, L Pollet, MA Martin-Delgado, arXiv:2112.05122.

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Thank you for your attention!

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