

Title: The holographic map as a conditional expectation

Speakers: Thomas Faulkner

Series: Quantum Fields and Strings

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Abstract: I will study quantum error correcting codes that model aspects of the AdS/CFT correspondence. In an algebraic approach I will demonstrate the existence of a consistent assignment, to each boundary region, of conditional expectations that preserve the code subspace. This allows us to give simple derivations of well known results for these holographic code, and also to derive a few new results.&nbsp;

I will also make a connection to the theory of QFT super-selection sectors.



①

# The Holographic Map as a Conditional Expectation

2008.04810

Main take away:

## Concrete:

- Quantum Error correcting codes that model AdS/CFT naturally give rise to conditional expectations: well studied mathematical structure in operator algebras

- Strengthens connections to Superselection



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- Quantum Error Correcting codes that model AdS/CFT naturally give rise to conditional expectations: well studied mathematical structure in operator algebras

- Strengthens connections to Superselection sectors elucidated by Casini, Huerta, Mugan, Pontello

Speculative:



AdS/CFT naturally give rise to conditional expectations: well studied mathematical structure in operator algebras

- Strengthens connections to superselection sectors elucidated by Casini, Huerta, Mugan, Pontello

**Speculative:**

- Study multiple boundary regions
- Exact recovery (studied here) has shortcomings that are (hopefully) not fatal



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- give a picture of entanglement wedge phase transitions.

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A model of entanglement wedge reconstruction (EWR)

Will call "complementary recovery"

Dong, Harlow, Wall; Harlow; Kang-Kolchmeyer  
( $\infty$ -dims)



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$V: \mathcal{X} \rightarrow \mathcal{K}$  is an isometric embedding

$$V^\dagger V = \mathbb{I}_{\mathcal{X}} \quad VV^\dagger = e$$

$e$ : code subspace projector. on  $\mathcal{K}$ .

③

EWR: how does this holographic map restrict



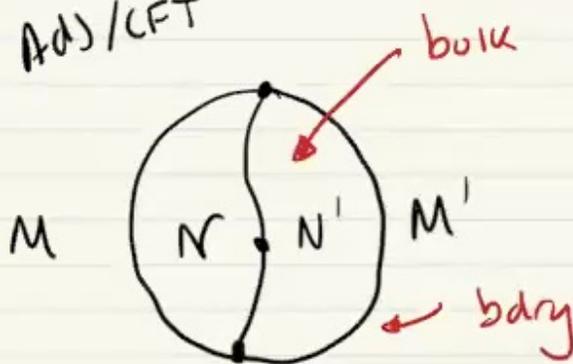
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Picture:

AdS/CFT

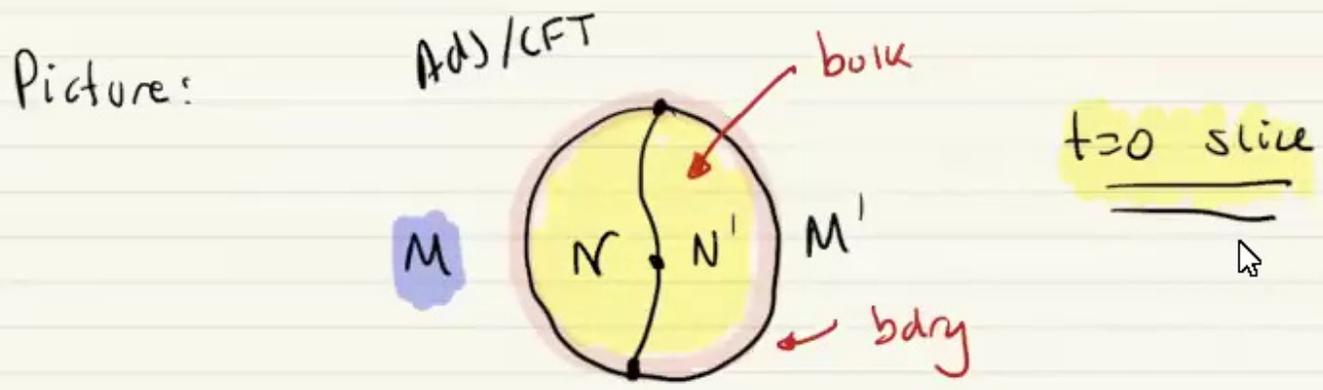


$t=0$  slice



③

EWR: how does this holographic map restrict to local algebras on  $\mathcal{X}$  &  $\mathcal{K}$ ?



- $M$ : bdry algebra say  assume

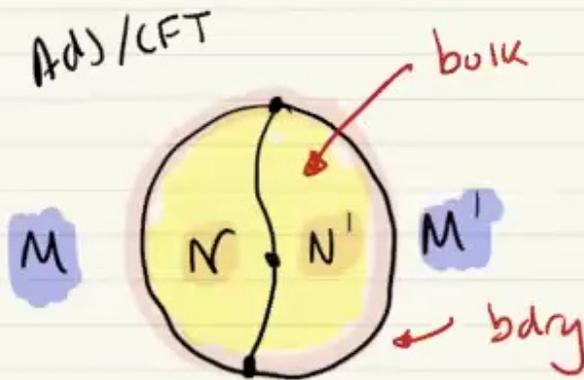


3

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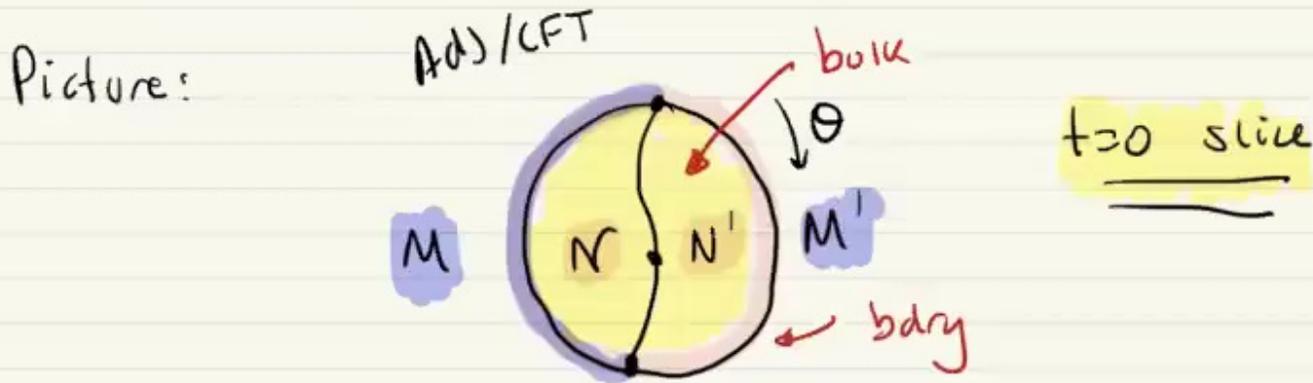
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to local algebras on  $\mathcal{H}$  &  $\mathcal{K}$ :

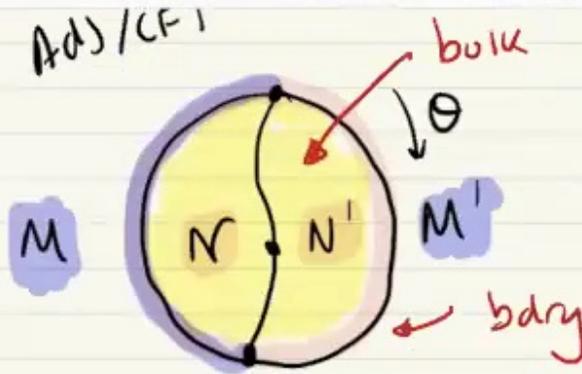


- $M$ : bdry algebra say  $\begin{matrix} \uparrow \\ \theta \\ \rightarrow \end{matrix}$  assume it is a vn algebra.  $M \subset B(\mathcal{K})$

- $N$ : bulk algebra associated to entanglement wedge. Also assume it is a vn algebra



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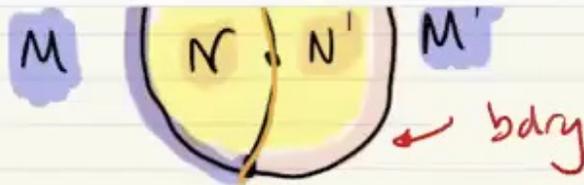


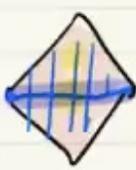
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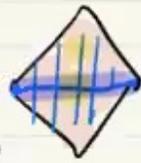
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$$M' = \{ m' \in B(\mathcal{K}) : [m', M] = 0 \}$$



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(4)

- New developments in AT approach to AdS/CFT



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- Now developments in QI approach to AdS/CFT



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- Now developments in QI approach to AdS/CFT

RT  $\rightarrow$  FLM  $\rightarrow$  JLMS allowed for a

precise characterization of how  $N$  &  $M$

are related



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Def:

- DHW condition:  $\forall \psi_1, \psi_2 \in \mathcal{H}$  then:

$$\omega_{\psi_1} |_{N'} = \omega_{\psi_2} |_{N'} \Rightarrow \omega_{\psi_1} |_{M'} = \omega_{\psi_2} |_{M'}$$

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Will say "N is reconstructable from M"

Comments: (i)  $\omega_{\psi}(\cdot) \equiv \langle \psi | \cdot | \psi \rangle$  is a

positive linear functional  $\sim$  "state"

algebraic version of density matrix

$\omega_{\psi}$  is global defined on  $B(\mathcal{H}) \rightarrow$  all

correlation functions for  $|\psi\rangle$



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$\omega_{\psi}|_{N'}$  means restrict to  $N' \subset B(\mathcal{H})$

Algebraic version of a reduced dens. matrix



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(5)

ii) **physically**: information about  $N$

(via differences in the state  $w_{t_1|N} \neq w_{t_2|N}$ )

cannot be detected on  $M'$ . **A version of**

**privacy**

iii) often stated as:

$$S_{rel}(\psi_1 | \psi_2; N') = 0 \Rightarrow S_{rel}(V\psi_1 | V\psi_2; M') = 0$$

derived in ADS/CFT from JLMS



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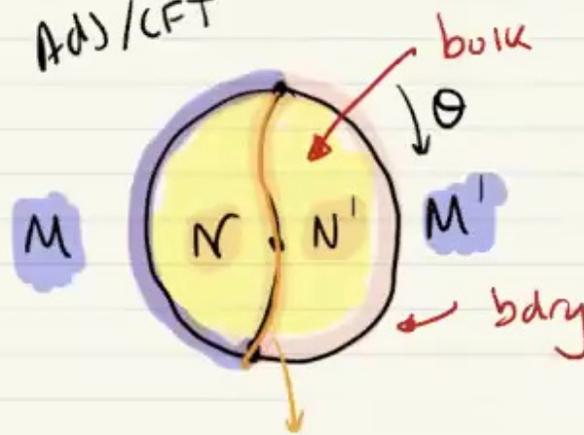


(3)

EWR: how does this holographic map restrict to local algebras on  $\mathcal{X}$  &  $\mathcal{K}$ ?

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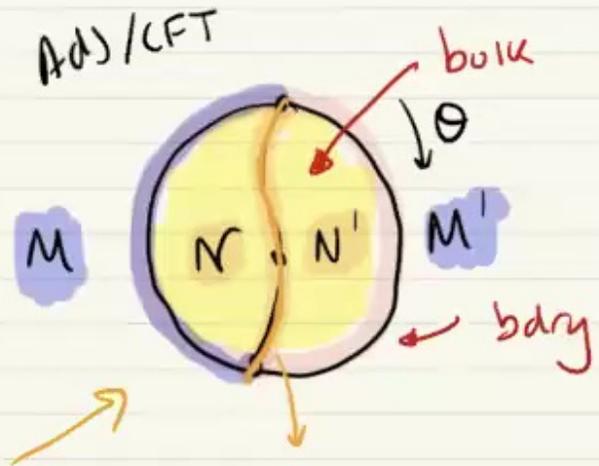
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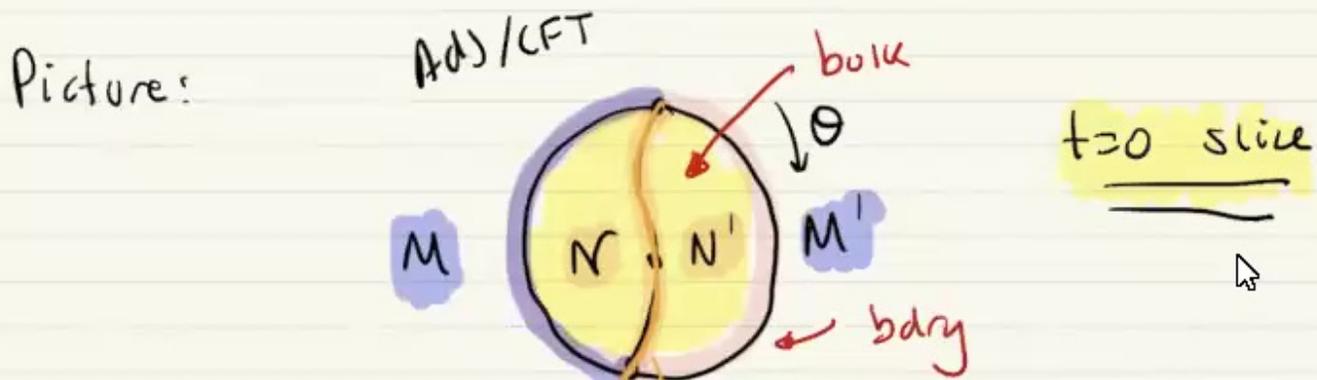
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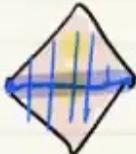
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$S_M = \frac{\text{Area}}{4\pi G} + S_N$

JLMS  $S_{\text{ren}}(\psi_i)$

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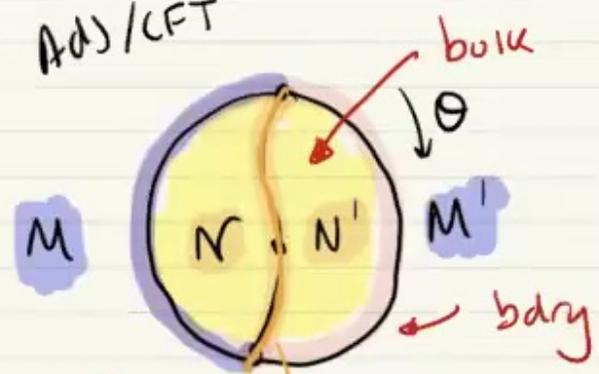


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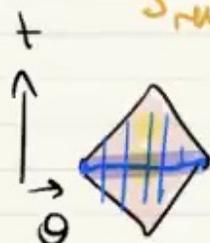
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- (6)

## Complementary Recovery:

Def:

" $N$  is  $c$ -reconstructable from  $M$  if"

- a)  $N$  is reconstructable from  $M$
- b)  $N'$  is reconstructable from  $M'$




---

To understand this structure we give a well known theorem (Accardi, Cecchini; <sup>DHW</sup> Harlow; Karg-Kolchmeyer)

The following are equivalent:



$$S_{rel}(\psi_1, \psi_2; N') = 0 \Rightarrow S_{rel}(\psi_1, \psi_2; M') = 0$$

derived in AdS/CFT from JLMS

(iv) correct up to non-perturbative errors in

$\ln N$ , although only for "reconstructable wedge" & not "entanglement wedge":

these are only different for very large  
code subspaces or near EW transitions

We will assume EXACT



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derived in ADS/CFT from JLMS

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$b_N$ , although only for ("reconstructable wedge" & not "entanglement wedge":)

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**EXACT**

b)  $N'$  is reconstructable from  $M'$

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Main  
Theorem

The following are equivalent:

(i) DHW  $(w_{\psi_1} - w_{\psi_2})|_{N'} = 0 \Rightarrow (w_{V\psi_1} - w_{V\psi_2})|_{M'} = 0$

$\forall \psi_1, \psi_2 \in \mathcal{X}$

(ii)  $\left( \begin{array}{l} \alpha_G : B(\mathcal{K}) \rightarrow B(\mathcal{X}) \\ \text{(body)} \quad \text{(bulk)} \end{array} \quad \alpha_G \equiv V^\dagger(\cdot)V \right)$

restrict to:



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 DHW†

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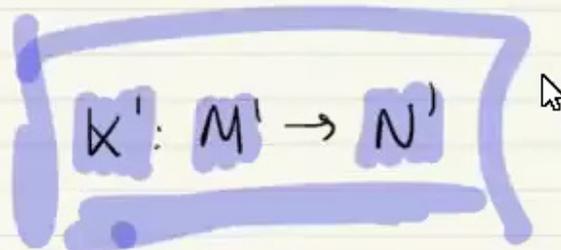
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restricts to:

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( $\alpha'$  is faithful)

(injective)

Ⓣ

(iii)  $\exists \beta : N \rightarrow M$  a  $1^*$ -homomorphism

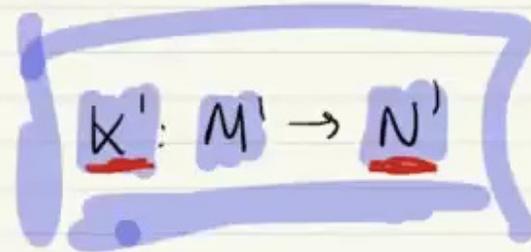
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$$(ii) \left( \alpha_G : B(K) \rightarrow B(N) \quad \alpha_G \equiv V^+(M) V \right)$$

(obj)                      (bulk)

restricts to:

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$$(iii) \exists \beta : N \rightarrow M \quad \text{a } \lambda\text{-homomorphism}$$

(injective)      ⑦

$$\beta(n_1 n_2) = \beta(n_1) \beta(n_2)$$

(unital  $\beta(1) = 1$ ) & such that



(iii)  $\exists$   $(\beta) \underline{N \rightarrow M}$  a  $\lambda$   $*$ -homomorphism

$$\beta(\underline{n_1 n_2}) = \underline{\beta(n_1)} \underline{\beta(n_2)}$$

(unital  $\beta(1) = 1$ ) & such that

$$\beta(n) V = V_n$$

Comments:

- $\alpha \leftrightarrow \beta$  called the dual map. it is a version of the Petz map (without  $J$ 's)
- $\alpha$  &  $\beta$  are quantum channels in the



(unital  $\beta(1) = 1$ ) & such that

$$\beta(V) = V$$

Comments:

- $\alpha \leftrightarrow \beta$  called the dual map. it is a version of the Petz map (without  $J$ 's)
- $\alpha$  &  $\beta$  are quantum channels in the Heisenberg picture (completely positive, normal)
- Explains the name "reconstructable"



S.t  $\beta(n)|v\rangle = |v\rangle$  so acts

correctly on the code subspace. (8)

- we have stated a version which requires a single cyclic & sep. vector  $|v\rangle$  for  $M$

- $\beta(N) \cong N^B$  is a  $\forall N$  subalgebra of  $M$  b/c of the homomorphism property.

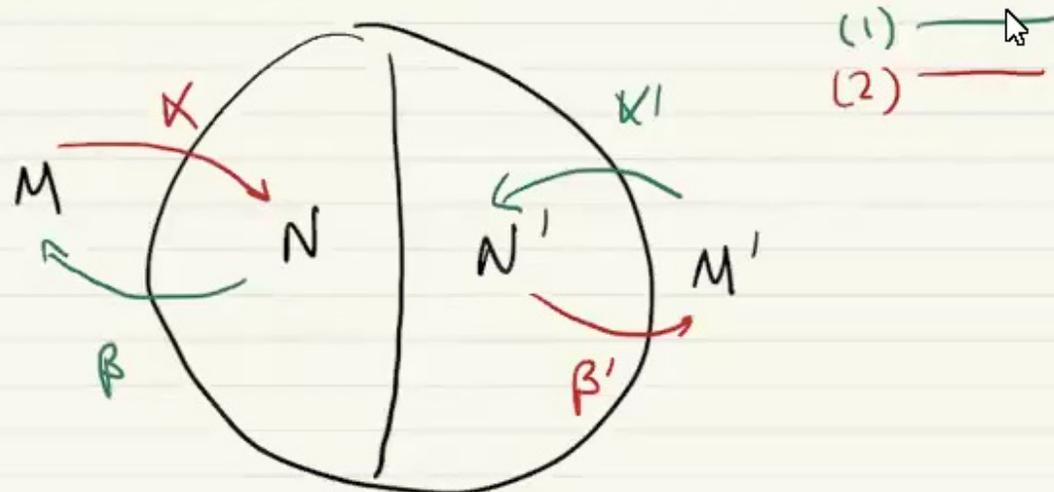
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a single cyclic & sep. vector  $\overline{|v\rangle}$  for  $M$

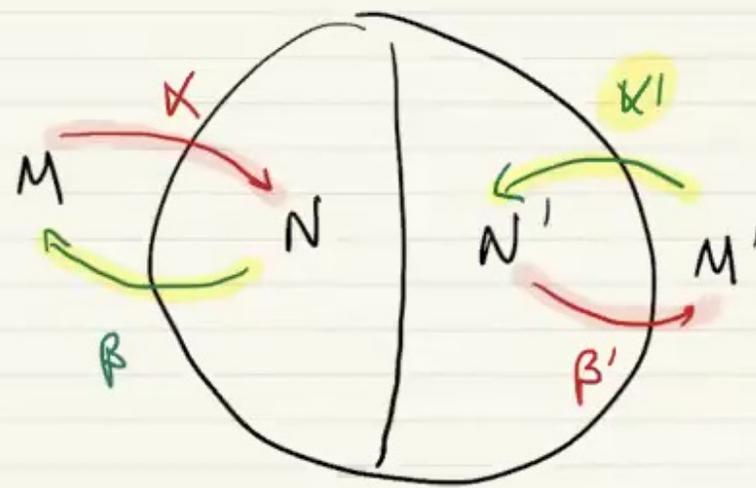
- $\beta(N) \cong N^{\beta}$  is a  $\sqrt{N}$  subalgebra of  $M$   
b/c of the homomorphism property.

Complementary recovery: apply theorem twice!





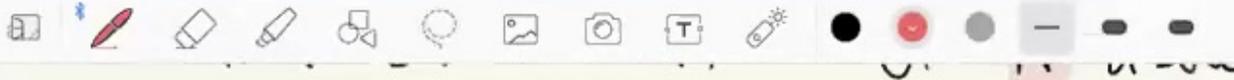
Complementary recovery: apply theorem twice!



(1) — green  
(2) — red

" $N$  is  $\epsilon$ -reconstructable from  $M$ "

\* You should think of  $N$  above  
as defining Entanglement wedge



as defining Entanglement Wedge

\*  $\mathcal{C}$ -reconstructability is a strong condition.

\* reconstructable is weaker

-  $\tilde{\mathcal{N}} \subset \mathcal{N}$  any subalgebra is reconstructable  
 [will show iff statement later]



-  $\tilde{N}_a, \tilde{N}_b$  reconstructable  $\tilde{N}_a \vee \tilde{N}_b$  is  
reconstructable (from  $M$ )

However  $N$  is unique [if it exists]

Claim:  $N$   $C$ -reconstructable from  $M$  then:

$$E = (\beta, \kappa) : M \rightarrow N^{\beta} C M$$

is a faithful conditional expectation



$\text{Id}$  preserves the whole subspace.

$$\underline{W_{\psi \neq 0} E = W_{\psi} \quad \forall \psi \in \mathcal{H}}$$

Conditional Expectation:

- $E: \underline{A \rightarrow B}$   $B \subset A$  subalgebra  $E(b) = b$

s.t.  $E(b_2 a b_1) = b_2 E(a) b_1$  ( $E(b) = b$ )

$b_1, b_2 \in B \quad a \in A \quad \uparrow$  "bimodule property"

- Existence is non-trivial for general  $\forall N$  algebras



algebras

- Non-commutative version of conditioning on some random variable  $B$

$$B|X_A \supset B(X_B) \otimes \mathbb{I}_C$$

- E.g.:  $X_A = \overline{X_B \otimes X_C}$

$$E(a) = \frac{\mathbb{I}_C}{\dim X_C} \otimes \text{tr}_C(a)$$

(10)

OR  $\int dg V(g)^{\dagger} a V(g)$  Haar measure

$A \rightarrow B = \text{fixed pt algebra of } V(g)$



OR  $\int dg V(g)^{\dagger} a V(g)$  Haar measure (10)

$A \rightarrow B =$  fixed pt algebra of  $V(g)$

---

Back to claim: check it fixes sub-algebra

- $E(\beta(n)) = \beta_0 \times \beta(n) = \beta(n)$

$$V^{\dagger} \beta(n) V = V^{\dagger} V n = n$$

reconstruction  $\curvearrowright$



$$\bullet \quad \underline{E(\beta(n))} = \beta_0 \times \beta(n) = \underline{\beta(n)}$$

$$\underline{\underline{X_0 \beta = Id}}$$

$$\underline{V^+ \beta(n) V} = V^+ V n = n$$

reconstruction

Q Info

$\psi = \text{Noise}$

$\beta = \text{Decoding}$

$$X_0 \beta = Id$$

$$\Rightarrow D_0 N = Id$$

$$\bullet \quad \underline{W_{V^+} \circ E(m)} = \langle \psi | V^+ E(m) V | \psi \rangle$$

$$= \langle \psi | V^+ \beta(X(m)) V | \psi \rangle$$

$$= \langle \psi | X(m) | \psi \rangle = \underline{W_{V^+}(m)}$$

Consequences: • Equality of Spect. through

the code (JLMS) [Petz]



- $E(\underline{\beta(n)}) = \beta_0 \times \beta(n) = \underline{\beta(n)}$

$X_0 \beta = Id$

$\underline{V^+ \beta(n) V} = V^+ V n = n$   
reconstruction

Q/Info  
 $\mathcal{N} = \text{Noise}$   
 $\beta = \text{Decoding}$   
 $X_0 \beta = Id$   
 $\Rightarrow \underline{D_0 \mathcal{N} = Id}$

- $\underline{W_{V^+} \circ E(m)} = \langle \psi | V^+ E(m) V | \psi \rangle$   
 $= \langle \psi | V^+ \beta(X(m)) V | \psi \rangle$   
 $= \langle \psi | X(m) | \psi \rangle = \underline{W_{V^+}(m)}$

Consequences: • Equality of Spect. through

the code (JLMS) [Petz]



$\beta = \text{Decoding}$

$\alpha \circ \beta = \text{Id}$

$\Rightarrow \underline{D \circ N = \text{Id}}$

$$= \langle \psi | V^\dagger \beta(\alpha(m)) | \psi \rangle$$

$$= \langle \psi | \alpha(m) | \psi \rangle = \underline{W_V^\dagger(m)}$$

Consequences: • Equality of Srel. through

the code (JLMS) [Petz]

- Equality of bulk & bdy mod

flow [ well known Takesaki ] (11)

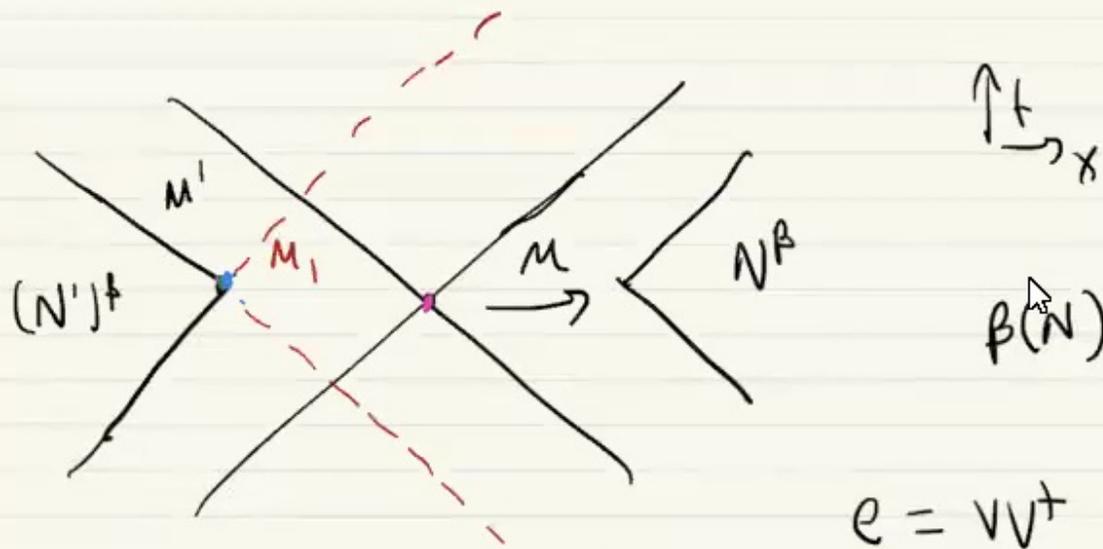
Also see: Besteuer & Karg for the converse



Also see: *Besteau & Kary* for the converse



A schematic (some what misleading) picture:



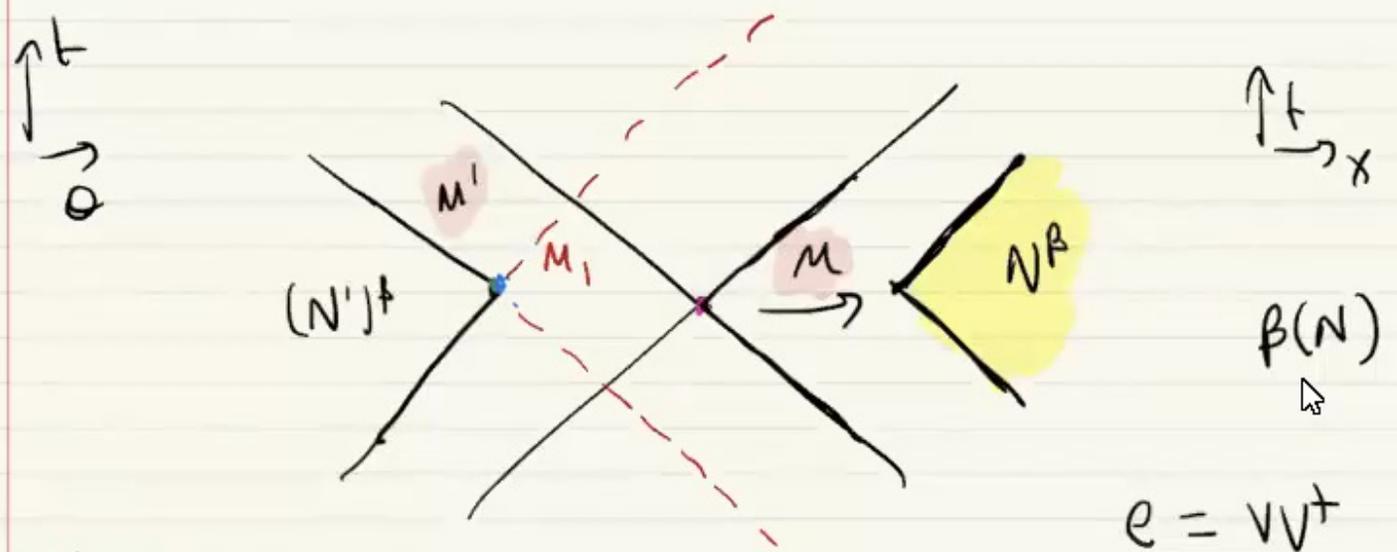
$$(N')^R \equiv \beta'(N')$$

- One can show:  $|(N')^t|' = (M v e)$

Handwritten notes at the top of the page, partially obscured by the toolbar.



A schematic (some what misleading) picture!



$(N')^B \equiv \beta'(N')$

- One can show:  $[(N')^A]^T = (M v e)$

 $p(N)$ 

$$(N')^{\beta} \equiv \beta'(N')$$

$$e = vv^{\dagger}$$

- One can show:  $[(N')^{\beta}]' = (Mve)$

$$\text{where } e = vv^{\dagger}$$

gen by these

- $Mve \equiv M_1$  is called the Jones extension, plays important role in index theory for subfactors.



extension, plays important role in  
Index theory for subfactors.

(12)

- Relative commutant plays an important role:

$$N^c \equiv M \cap (N^B)'$$

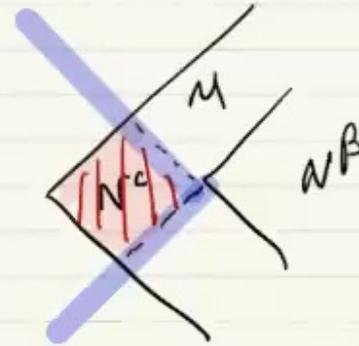


- Ryu Takayanagi Area operator: [Hofstadter]



- Relative commutant plays an important role.

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- Ryu Takayanagi Area operator: [Hulow]  
Naturally associated to  $E$ : [Takasaki]

Factorization  $M \cong \overbrace{(M \cap (N^B)')^{N^c}} \otimes N^B$

(almost: requires factors & normalcy)

Naturally associated to  $E$ : Takasaki

$$\text{Factorization } M \cong \overbrace{(M \wedge (N^A)^A)^{N^c}} \otimes N^B$$

(almost requires factors & normalcy)

States fixed by  $E$ :  $P \circ E = P$   $P$  on  $M$

$$P = P|_{N^c} \otimes P|_{N^B}$$

$$S_M(P) = S_{N^c}(P) + \overbrace{S_{N^A}(P)}^{P = \omega_{N^A}} = S_N(\omega_{N^A})$$

Bulk Entropy.



$N^c$

$N^c$

Bulk Entropy

(13)

State on  $N^c$ :  $\rho(n^c) = \rho_0 E(n^c)$

but  $E(n^c) \in N^B$  and commutes!

$$E(n^c) n^B = E(n^c n^B) = E(n^B n^c) = n^B E(n^c)$$

thus  $E(n^c) \in Z(N^B)$  the center!

$$\Rightarrow E(n^c) = \sum \pi_a \gamma_a(n^c)$$

Central projections

$$E(n^c) n^{\beta} = E(n^c n^{\beta}) = E(n^{\beta} n^c) = n^{\beta} E(n^c)$$

thus  $\underline{E(n^c)} \in Z(N^{\beta})$  the center!

$$\Rightarrow E(n^c) = \sum_a \pi_a \left( \gamma_a(n^c) \right) \quad \text{minimal central projections.}$$

$$S_M(p) = \sum_a p(\pi_a) S_{N^c}(\gamma_a) + S_{N^{\beta}}(p)$$

$$\downarrow p = w_{\psi}:$$

$$S_M(w_{\psi}) = w_{+}(\hat{\mathcal{L}}_N) + S_N(w_{\psi})$$

$$\Rightarrow E(n^c) = \sum_a \pi_a(\gamma_a(n^c)) \quad h$$

linear functional on  $N^c$

$$S_M(p) = \sum_a p(\pi_a) S_{N^c}(\gamma_a) + S_{N^p}(p)$$

$$\downarrow p = w_{V\psi}:$$

$$S_M(w_{V\psi}) = w_+ (\hat{d}_N) + \underline{S_N(w_{V\psi})}$$

$$\hat{d}_N = \sum_a \pi_a(S_{N^c}(\gamma_a))$$

central projections for  $N$



$$\gamma_a = W_f(\Pi_a)$$

only depends on  $\epsilon$

This is exactly the "RT formula" with associated area operator  $\hat{\mathcal{L}}_N$  as discussed in Harlow

Note: really type-I since we have assumed

$S_M$  etc. are finite. In type-III

setting [QFT] we can still define  $\hat{\mathcal{L}}_N$

---

A minimization formula; , Plays role of area law



Note: Really **type-I** since we have assumed

$S_M$  etc. are finite. In type-III

setting [QFT] we can still define  $\hat{L}_N$

A minimization formula: ↓ Plays role of area term.

$$S_M(\omega_{V4}) \leq \inf ( \overbrace{S_{N^c}(\omega_{V4})} + S_N(\omega_{V4}) )$$

Thm:

$\{ N: \text{reconstructable from } M \}$

↑  $\sim S_{gen}?$

$$N^c = M \wedge (\beta(N))'$$



Equality is achieved iff  $\exists$  an  $N$  <sup>(15)</sup>

s.t.  $N$  is  $C$ -reconstructible from  $M$

Very suggestive, but some issues with analogy  
to the RT formula minimization:

(1) not all subalgebras of EW have a  
larger  $S_{gen}$  ( $S_{gen}$  decreases along

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↘ Exact Recovery



(16)

Generalities: How far can we push this?

- Multiple boundary alg.  $M_1, M_2$

① Entanglement wedge Nesting:  $\mathcal{E}_1, \mathcal{E}_2 \sim$  Entanglement Wedges.

$\mathcal{E}_1$  is  $\mathcal{C}$ -reconstructable from  $M_1$

$\mathcal{E}_2$  is  $\mathcal{C}$ -reconstructable from  $M_2$

Then:

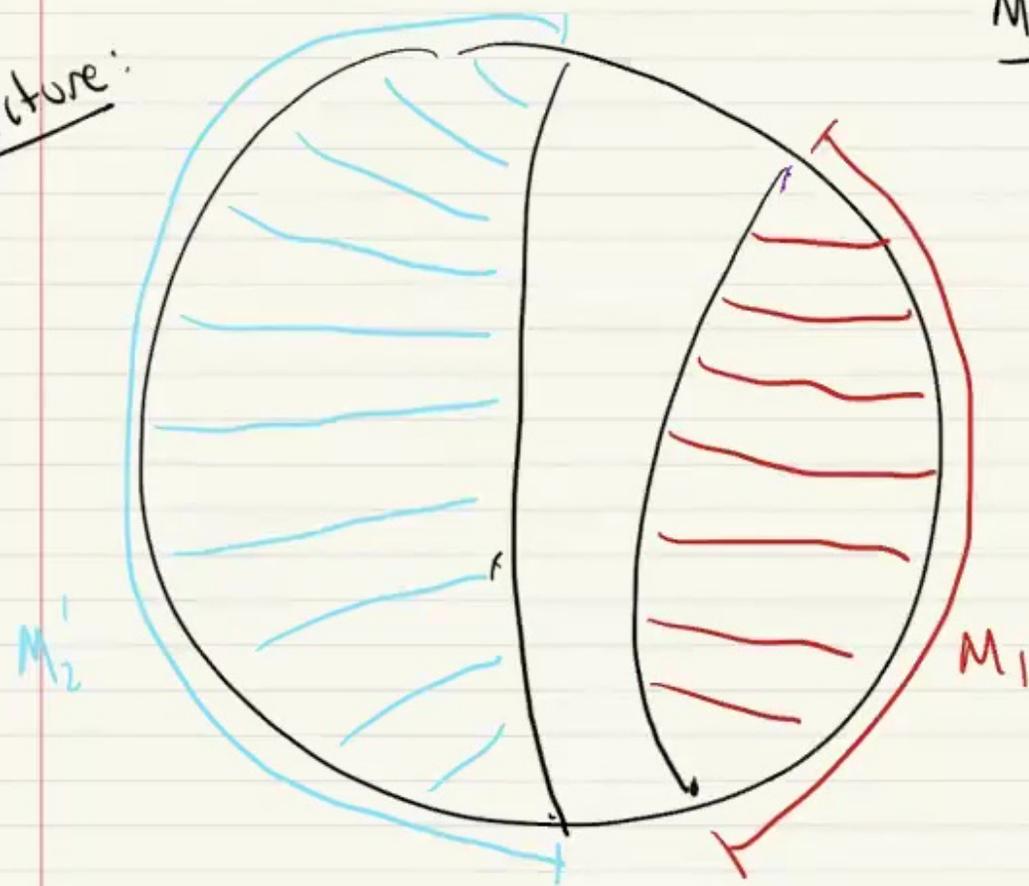
$$M_1 \subset M_2 \Rightarrow \mathcal{E}_1 \subset \mathcal{E}_2$$



Then:

$$M_1 \subset M_2 \Rightarrow \mathcal{E}_1 \subset \mathcal{E}_2$$

Picture:



$$\frac{M_1 \subset M_2}{\rightarrow}$$

$$[M_1, M_2'] = 0$$



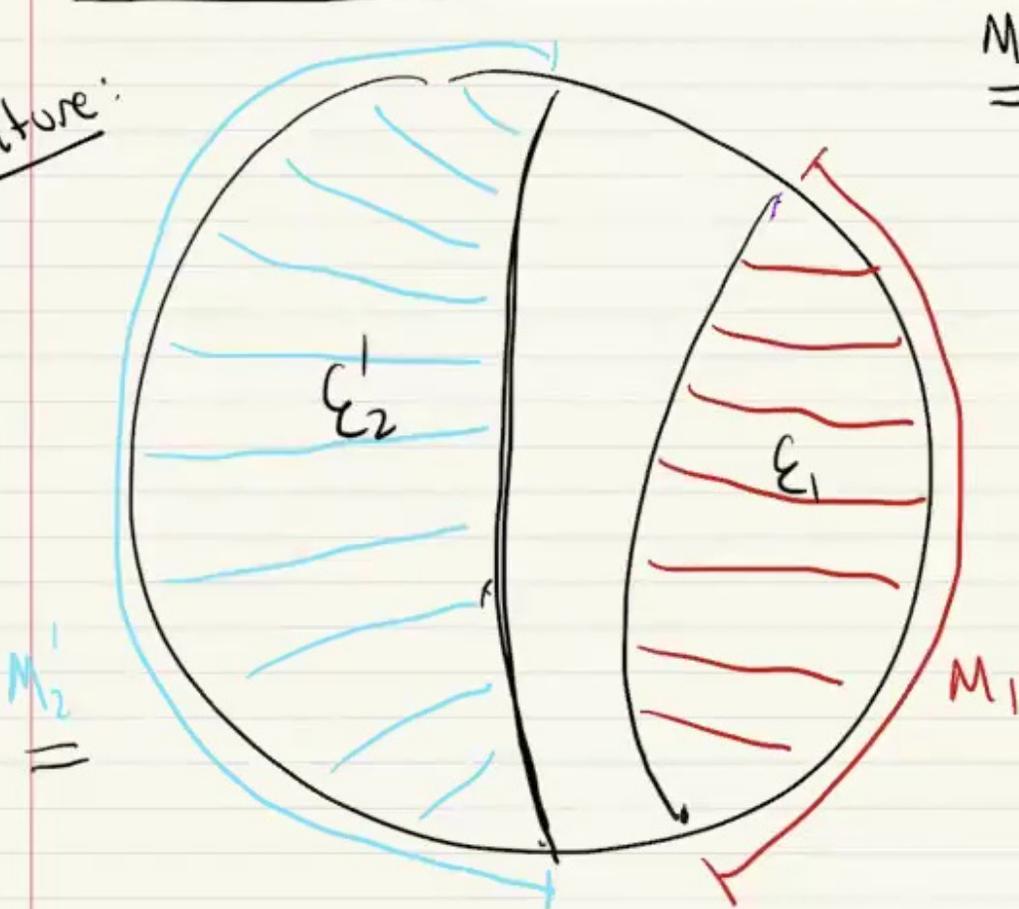
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Theri.

$M_1 \subset M_2 \Rightarrow \mathcal{E}_1 \subset \mathcal{E}_2$

Picture:



$\frac{M_1 \subset M_2 \rightarrow}{\underline{\underline{[M_1, M_2']}}} = 0$

$\rightarrow [E_1, E_2'] = 0$





$$\text{take } M_2 = M_1, N_2 = \mathcal{E}_1 \Rightarrow N_1 \subset \mathcal{E}_1$$

$\Rightarrow$  Reconstructable iff  $N_1 \subset \mathcal{E}_1$

② Consistency of  $E$ :  $M_1 \subset M_2$

with  $\mathcal{C}$ -reconstructable algebras  $N_1 \subset N_2$  then

$$E_2: M_2 \rightarrow (N_2)^{\mathcal{C}} \quad \text{restricts:}$$

$$E_2|_{M_1} = E_1: M_1 \rightarrow (N_1)^{\mathcal{C}}$$

This consistency proves that exact complementary

recovery is the same structure as a



consistency of  $E_1, E_2$

With  $C^*$ -reconstructible algebras  $N_1 \subset N_2$  then

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Largo Rehren



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Longo Rehren



Casini, Huerta, Magán, Pontello

Here  $\mathcal{N}^{\beta}$   $\mathcal{C}_M$  is the sub-algebra of neutral operators under some global charge [orbifolded theory] of the local operators in  $M$

$E$ : projects to this neutral sector

$$\frac{1}{|G|} \sum_{g \in G} V(g)^{\dagger} M V(g) \leftarrow$$



g<sub>UV</sub>

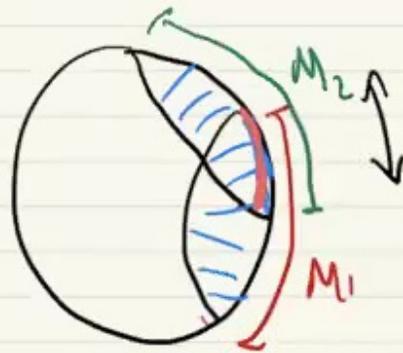
In AdS/CFT: E projects to the low-energy  
gravity degrees of freedom.

I have no specific proposal for E

→ Q. Info arguments tell us one  
should exist



③ There is an important issue: [W. Kelly]



Overlapping regions

Thm:  $N_i$  is  $C$ -reconstructible from  $M_i$

if  $\exists \gamma$  s.t.  $\forall \gamma$  is cyclic & sep for

$M_1 \wedge M_2$  then:

$N_1 \vee N_2$  is  $C$ -reconstructible from  $M_1 \vee M_2$



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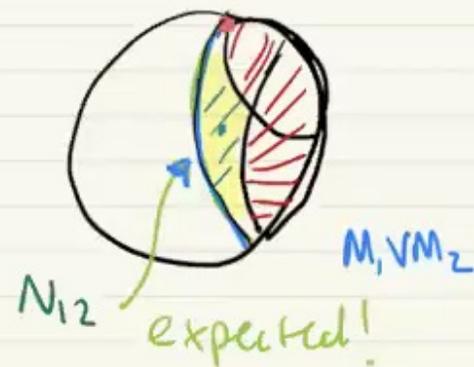
$N_1 \vee N_2$  is C-reconstructable from  $M_1 \vee M_2$

(19)

i.e. the bulk algebras are "additive" for



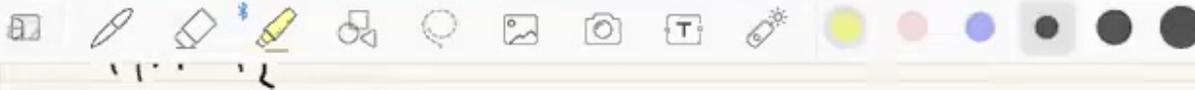
But: not true in Ads/CFT



Uses in an essential way the cyclicity for  $M_1, M_2$

i.e.  $\{ \alpha \vee \gamma ; \alpha + M_1, M_2 \}$  is  
a dense subset of  $K$ .

- Fix: (Kelly) is to have approximate recovery



ie.  $\{x \forall t> ; x + M_1 A M_2\}$  is  
a dense subset of  $K$ .

- Fix: (Kelly) is to have approximate recovery

then cyclicity loses its power!

Cyclicity still holds but:

$$x \forall t> \rightarrow \{s> \quad \|x\| \rightarrow \infty$$



$\&$

$$N_{12} \stackrel{c.rec}{\sim} M_1 VM_2$$

$$\Rightarrow ( N_{12}' \stackrel{c.rec}{\sim} (M_1 VM_2)' )$$

But what do we pick for  $N_{12}$ ?

Holography suggests two options:

$$N_{12} = \underline{N_1 VN_2}$$

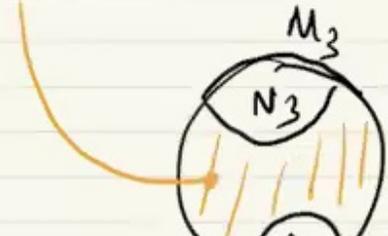
or

$$N_{12} = (\underline{N_3 VN_4})'$$

(21)

• We call such a code

"Dual Addition"





## Future work:

### \* Beyond Exact Recovery

-  $\kappa$ -bit Story (Hayden Penington)

- breakdown of complementary recovery (Akers, Leichenauer, Levine

& Akers, Penington)

\* Lofty goal: axiomatic



- breakdown of complementary  
newery (Akers, Leichenauer, Levine  
& Akers, Penington)

\* Lofty goal: axiomatic  
characterization of quantum gravity  
(in low energy / EFT limit) in AdS