Title: Toward a Quantum-Safe Future

Speakers: Michele Mosca

Series: Colloquium

Date: April 08, 2020 - 2:00 PM

URL: http://pirsa.org/20040005

Abstract: There has been tremendous progress in the many layers needed to realize large-scale quantum computing, from the hardware layers to the high level software. There has also been vastly increased exploration into the potentially useful applications of quantum computers, which will drive the desire to build quantum computers and make them available to users. I will describe some of my research in quantum algorithmics and quantum compiling.

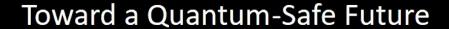
The knowledge and tools developed for these positive applications give us insight into the cost of implementing quantum cryptanalysis of today's cryptographic algorithms, which is a key factor in estimating when quantum computers will be cryptographically relevant (the "collapse time"). In addition to my own estimates, I will summarize the estimates of 22 other thought leaders in quantum computing.

What quantum cryptanalysis means to an organization or a sector depends not only on the collapse time, but also on the time to migrate to quantum-safe algorithms as well as the shelf-life of information assets being protected. In recent years, we have gained increasing insight into the challenges of a wide-scale migration of existing systems. We must also be proactive as we deploy new systems. Open-source platforms, like OpenQuantumSafe and OpenQKDNetwork, are valuable resources in helping meet many of these challenges.

While awareness of the challenges and the path forward has increased immensely, there is still a long road ahead as we work together with additional stakeholders not only to prepare our digital economy to be resilient to quantum attacks, but also to make us more resilient to other threats that emerge.

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Perimeter Institute Colloquium

8 April 2020

Michele Mosca















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Security and readiness are too often an afterthought

https://www.wsj.com/articles/michael-hayden-says-u-s-is-easy-prey-for-hackers-1434924058

June 21, 2015 11:19 p.m. ET

THE WALL STREET JOURNAL.

Michael Hayden Says U.S. Is Easy Prey for Hackers

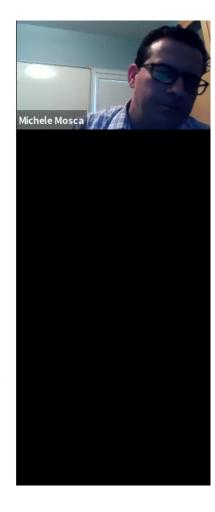
Former CIA and NSA chief says 'shame on us' for not protecting critical information better

GEN. HAYDEN:

Then, all of us just fell in love with the ease and convenience and scale, so we decided to take things we used to keep if not in a safe, at least in our desk drawer, and put it up here, where it's by definition more vulnerable.

•••

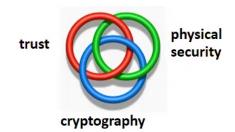
the next sound you hear will not be a digital bugle signaling the arrival of the digital cavalry to come save the day. The government ain't coming. You're not quite on your own, but you are more on your own up here [in cyberspace] than you in your lifetime have ever experienced being on your own down here.



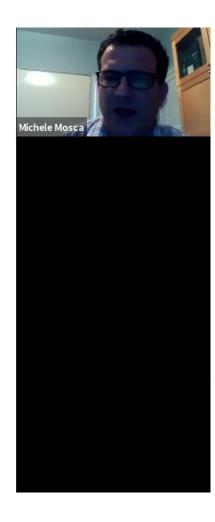
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Cryptography is a foundational pillar of cybersecurity

- Cryptography allows us to achieve information security while using untrusted communication systems.
- e.g. Do you update your software and anti-virus daily? Why do you trust the source?
- N.B. Cryptography is susceptible to "record now, decrypt later".



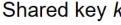




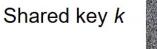
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Message Mto be encrypted













CipherText(k,M) =



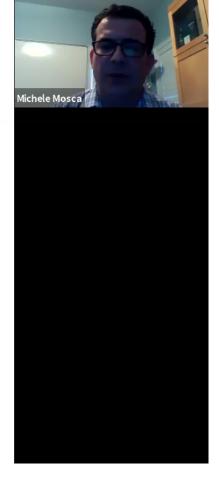




Shared key k



Decrypted Message M



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- Symmetric encryption
- Key Establishment
- Authentication

Message *M* to be encrypted



Shared key k





CipherText(k,M) =





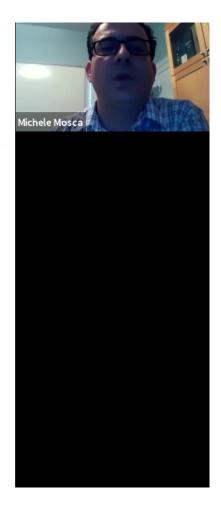


Shared key k



Decrypted Message *M*

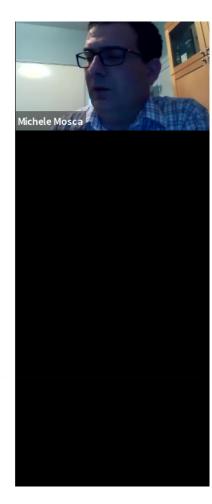




So many different vulnerabilities

- · Fundamentally vulnerable cryptography
- · Cryptography implementation errors
- User errors
- · Platform implementation errors
- Platform design errors
- Admin errors

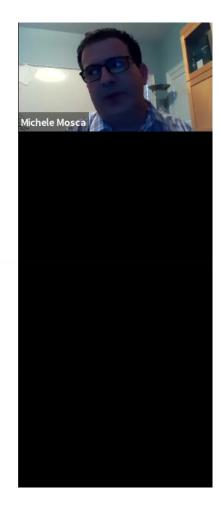
- Corrupt users
- Corrupt admin



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Vulnerabilities, ranked, from bad to worse?

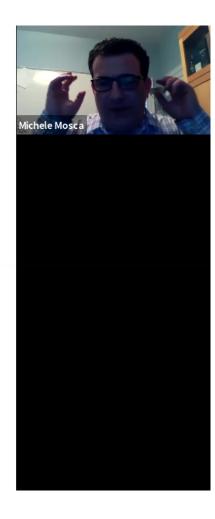
- User errors
- · Corrupt users
- Admin errors
- Corrupt admin
- Platform implementation errors
 - Platform design errors
- Crypto implementation errors



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Vulnerabilities, ranked, from bad to worse?

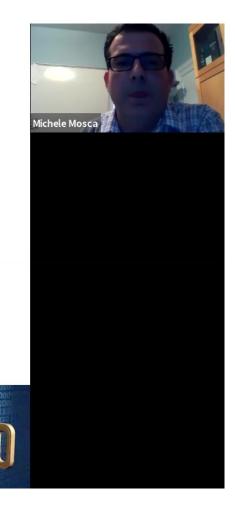
- User errors
- Corrupt users
- Admin errors
- Corrupt admin
- Platform implementation errors
 - Platform design errors
- Crypto implementation errors
- Fundamentally vulnerable cryptography



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Vulnerabilities, ranked, from bad to worse?

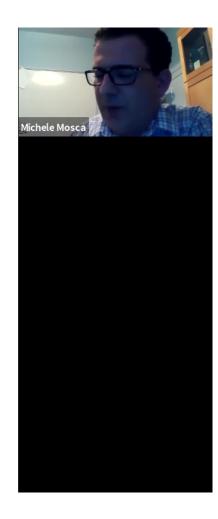
- User error
- Corrupt users
- Admin errors
- Corrupt admin
- Platform implementation errors
 - Platform design errors
- Crypto implementation errors
- Fundamentally vulnerable cryptography



Unpredictable new vulnerabilities

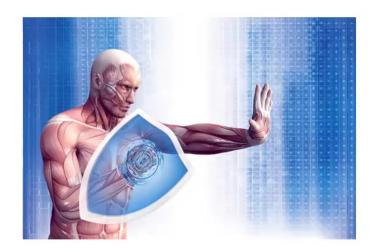


"As is the norm, an unexpected problem occurred today."

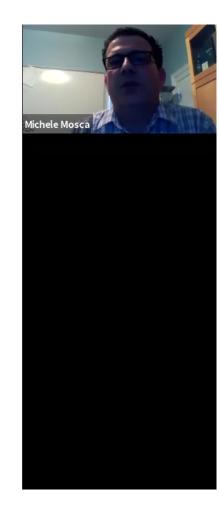


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Do we have a strong cyber immune system?

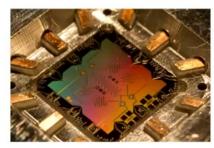


As an example, let's consider one emerging threat.

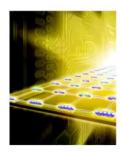


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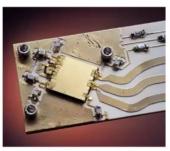
A new paradigm for computation: quantum computation



E. Lucero, D. Mariantoni, and M. Mariantoni



© Harald Ritsch



Y. Colombe/NIST



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Strong desire to implement quantum technology



Designing new materials, drugs, etc.



Optimizing



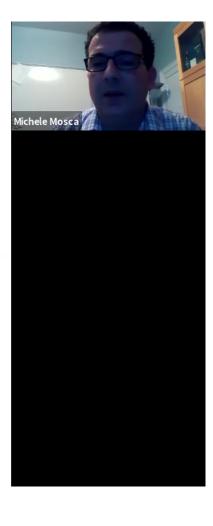
Sensing and measuring



Secure communication



What else???

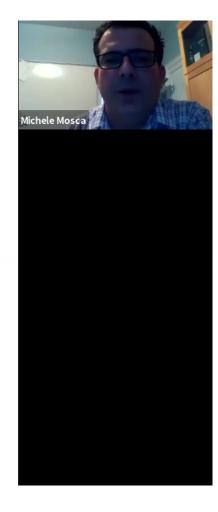


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What sorts of practical applications?

Possibilities include:

- Optimizing the design of new materials
 - For example, next generation materials could allow more efficient energy capture or transport or storage.
- Simulating chemical reactions at the quantum level
 - Potential applications include more efficient yields for chemical processes like the production of fertilizers.
- Optimization of designs or allocation or resources
 - For example, optimizing in the insertion of dampers in buildings to protect against earthquakes.



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IBM's new 53-qubit quantum computer is its biggest yet

The system will go online in October.



Forbes

IBM Doubles Its Quantum **Computing Power Again**





August 10, 2018

Google moves toward quantum supremacy with 72-qubit computer



physicsworld

POLICY AND FUNDING | NEWS

Quantum Circuits bags \$18m in first-round financing



Intel brings Quantum computing a step closer to reality

intel is betting on its fabrication expertise to push quantum computing into the mainstream

2 in 0



quantum computing real. Google, IBM, Microsoft among other prominent big names in the Industry are already working on quantum machines that can work outside the confines of academia. Intel is betting on its

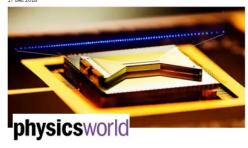
Technology

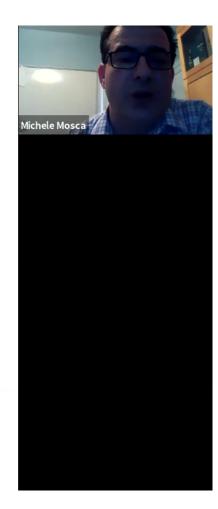
Microsoft Edges Closer to Quantum **Computer Based on Elusive Particle**

Researchers make Majorana fermions, but now must try to control them

By Jerenny Kahn March 28, 2018, 11-48 AM EDT. Corrected March 28, 2018, 2:09 PM EDT.

Ion-based commercial quantum computer is a first 17 Dec 2018





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China is opening a new quantum research supercenter

The country wants to build a quantum computer with a million times the computing power presently in the world.

By Jeffrey Lin and P.W. Singer October 10, 2017





NATIONAL LABORATORY FOR QUANTUM INFORMATION SCIENCES

The \$10 billion National Laboratory for Quantum Information Sciences in Hefei will be the center of China's attempt to take the global lead in quantum computing and sensing.

Alibaba puts 11-qubits quantum power on public cloud

Together with Chinese Academy of Sciences, Alibaba Cloud has unleashed superconducting quantum computing services on its public cloud, running on a processor with 11 quantum bits of power.

By Elleen Yu for By The Way | March 1, 2018 -- 1411 GMT (0611 PST) | Topic: Cloud



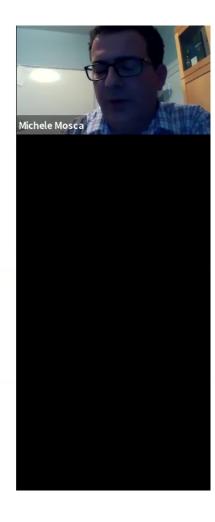
China's race for the mother of all supercomputers just got more crowded

Baidu, Alibaba and Tencent jockey for position in the development of quantum computing, which delivers a faster and more efficient approach to processing information than today's fastest computers

PUBLISHED : Monday, 12 Merch, 2018, 9:03am UPDATED : Monday, 12 March, 2018, 9:02am



Baidu has entered the race to build quantum computers



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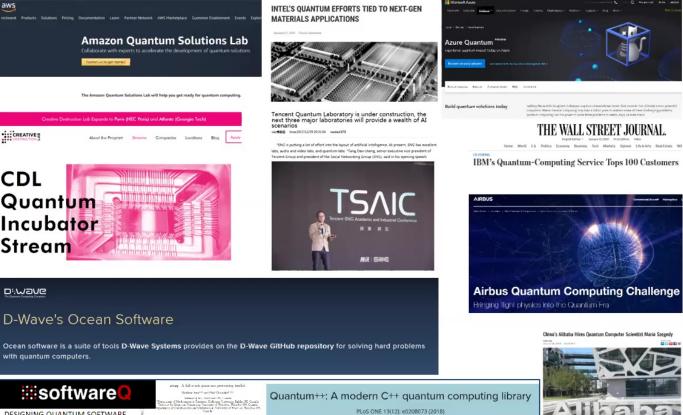


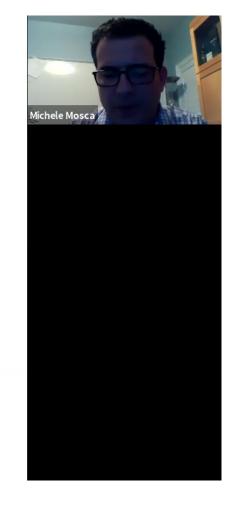
with quantum computers.

software

DESIGNING QUANTUM SOFTWARE







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What can you do **now** with a quantum computer?

- Currently, not enough qubits for a realistic speed-up (in circuit model). Proofof-principle quantum programs.
- Main prize: understand NOW what it takes to use one and identify potential business cases, so your business will benefit ASAP once the tech is out there
- Priority: finding applications and optimizing algorithms
- Resource counts continue to decline as we improve algorithms/software; make sure you are prepared
- Plausible that in next 2-4 years (e.g. IARPA LogiQ) that a fault-tolerant qubit is demonstrated, followed by intense scaling effort
- Possible that some quantum advantage is attainable with sufficiently large NISQ devices or quantum annealers.



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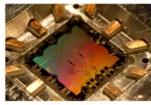
One serious problem for public-key cryptography

Algorithms for Quantum Computation: Discrete Logarithms and Factoring

Peter W. Shor
AT&T Bell Labs
Room 2D-149
600 Mountain Ave.
Murray Hill, NJ 07974 USA

In Proceedings, 35th Annual Symposium on Foundations of Computer Science, Santa Fe, NM, November 20 -22, 1994, IEEE Computer Society Press, pp. 124-134.

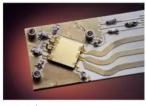






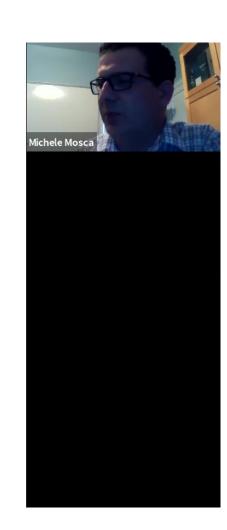


Christian Lagerek/Alamy



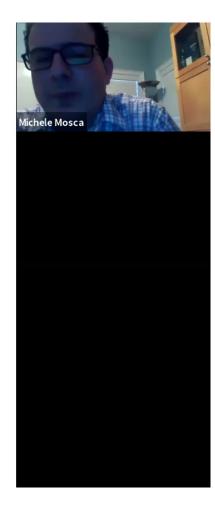
Y. Colombe/NIST

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How secure will our current crypto algorithms be?

Algorithm	Key Length	Security level (Conventional Computer)	
RSA-1024	1024 bits	80 bits	~0 bits
RSA-2048	2048 bits	112 bits	~0 bits
ECC-256	256 bits	128 bits	~ <mark>0</mark> bits
ECC-384	384 bits	192 bits	∼0 bits
AES-128	128 bits	128 bits	~64 bits
AES-256	256 bits	256 bits	~128 bits



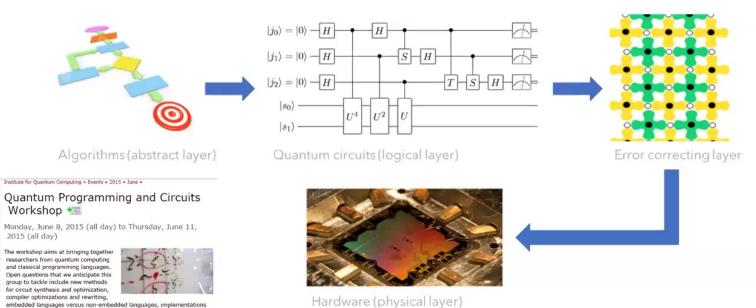
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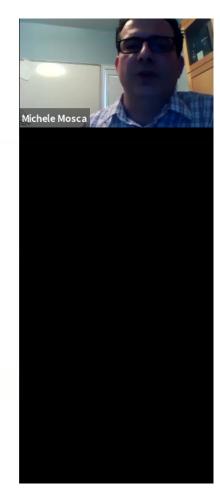
^{*}For symmetric crypto, known attacks are "only" quadratic, so doubling key or hash lengths suffices (against known attacks).

How large of a quantum computer is needed?

https://qsoft.iqc.uwaterloo.ca/

embedded languages versus non-embedded languages, implementations of type systems and error reporting for quantum languages, techniques for verifying the correctness of quantum programs, and new techniques for compiling efficient circuits and protocols for fault-tolerant questions and

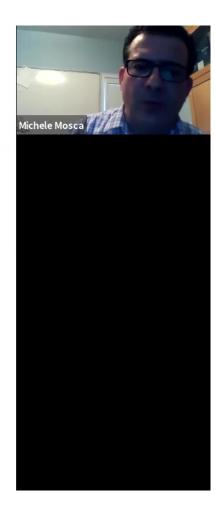




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Some useful circuit synthesis and optimization tools

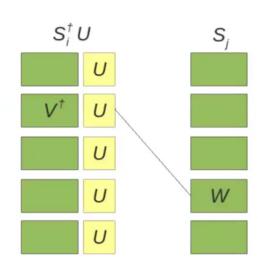
- Brute force exhaustive synthesis of multi-qubit unitaries
- Parallel collision-finding algorithms applied to circuit synthesis
- Optimal T-depth synthesis of one-qubit unitaries
- Optimization of T-depth via matroid partitioning
- Optimizing phase polynomials via Reed-Muller decoding

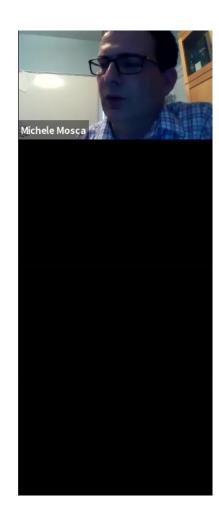


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Synthesis tools

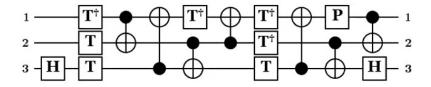
General multi-qubit MITM [AMMR13, TCAD]
Applied parallel collision-finding method
[Di Matteo, M 16, QST]



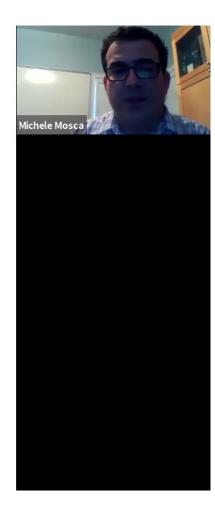


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Efficient Toffoli circuits



• T-depth 3 Toffoli (40% improvement)



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T-gate parallelization

 Selinger showed that the T-depth of the Toffoli could be reduced to 1 with the addition of enough ancilla qubits.

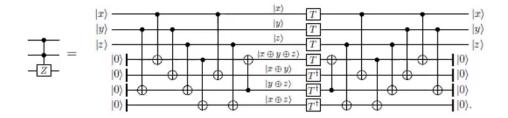
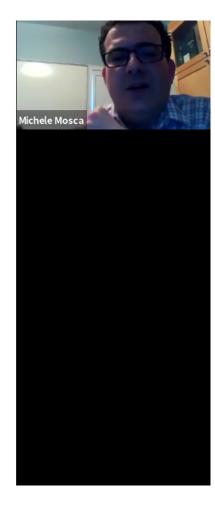


Figure 1: T-depth 1 representation of the Toffoli gate

arXiv:1210.0974v2 [quant-ph] 3 Apr 2013

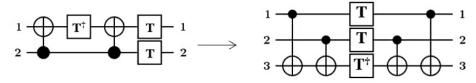


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A framework for T-gate parallelization (AMM13)

arXiv:1303.2042v1 [quant-ph] 8 Mar 2013

Our result: Can write {CNOT, T} circuits in T-depth 1 using unbounded ancillae



e.g.
$$U: |a,b\rangle \to e^{\frac{i\pi}{4}(-(a\oplus b)+a+b)}|a,b\rangle$$

Idea: prepare qubits in states that are XORs of inputs then perform all T gates at once

$$|ab0\rangle \rightarrow |a,b,a \oplus b\rangle \rightarrow e^{\frac{i\pi}{4}(-(a\oplus b)+a+b)}|a,b,a \oplus b\rangle \rightarrow e^{\frac{i\pi}{4}(-(a\oplus b)+a+b)}|a,b,0\rangle$$

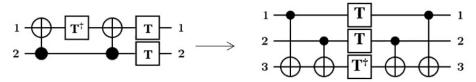


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A framework for T-gate parallelization (AMM13)

arXiv:1303.2042v1 [quant-ph] 8 Mar 2013

Our result: Can write {CNOT, T} circuits in T-depth 1 using unbounded ancillae



$$U: |a,b,c...\rangle \rightarrow e^{\frac{i\pi}{4}\sum_{i=1}^{k} f_i(a,b,c,...)} |g(a,b,c,...\rangle$$

e.g.
$$U: |a,b\rangle \to e^{\frac{i\pi}{4}(-(a\oplus b)+a+b)}|a,b\rangle$$

Idea: prepare qubits in states that are XORs of inputs then perform all T gates at once

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A framework for T-gate parallelization (AMM13)

arXiv:1303.2042v1 [quant-ph] 8 Mar 2013

Also considered any given number of ancilla qubits.

e.g. if we have 80 terms in 6 variables and 2 ancilla qubits (so 8 qubits total), can we partition the variables into 10 parts, and have T-depth 8?

It depends: each part should be able to reconstruct the input bits.

Matroid partitioning

 poly-time algorithm for finding minimal partitions of matroids (a set with an independence relation) Michele Mosca

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T-parallelization algorithm

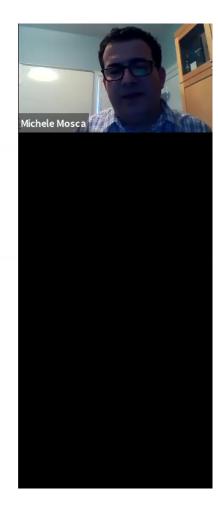
- First write circuit over {CNOT, T, H}
- For each {CNOT, T} subcircuit:
 - Find f_1, \dots, f_k, g s.t.

$$U: |abc \dots\rangle \to e^{\frac{i\pi}{4}\sum_{i=1}^k f_i(a,b,c,\dots)} |g(a,b,c,\dots\rangle$$

- Find minimal partition P of $f_1, ..., f_k$
- For each $p \in P$, synthesize T-depth 1 circuit for

$$V: |abc \dots\rangle \to e^{\frac{i\pi}{4}\sum_{f \in p} f(a,b,c,\dots)} |abc \dots\rangle$$

- Benchmark results (18 arithmetic circuits):
 - 32.1% avg T-count reduction
 - 53.5% avg T-depth reduction (no ancillas)
 - 82.4% avg T-depth reduction (as many ancillas as inputs)



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One-qubit gate synthesis (thanks to Vadym Kliuchnikov for slides)

▶ Algorithm building block: Quantum Fourier Transform

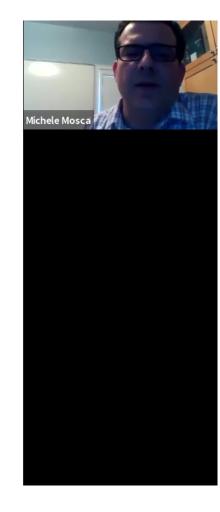
▶ Uses
$$R_Z(\phi) = \begin{pmatrix} e^{-i\phi/2} & 0 \\ 0 & e^{i\phi/2} \end{pmatrix}, \phi = \frac{\pi}{2^n}$$

- ► Available gates in **Clifford+T** gate library:
 - Single qubit:

$$T = \begin{pmatrix} 1 & 0 \\ 0 & \omega \end{pmatrix}, \omega = e^{i\pi/4} \qquad H = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix}$$

$$P = \begin{pmatrix} 1 & 0 \\ 0 & \omega^2 \end{pmatrix}, \text{ Pauli: } X, Y, Z$$

$$Two qubit: CNOT = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix}$$



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Known approaches to unitary approximation

[Before number-theoretic methods]

Name	Gates	Classical	Number	Resource
		run-time	of ancillae	state
[1] Solovay-Kitaev	$O\left(\log^{3.97}(1/\varepsilon)\right)$	$O\left(\log^{2.71}\left(1/arepsilon ight) ight)$	0	no
[2] Brute force	$O(\log(1/arepsilon))$	$2^{O(\log^c(1/\varepsilon))}$	0	no
[3] Phase kickback	$O(\log(1/arepsilon))$	$O(\log(1/arepsilon))$	$O(\log(1/\varepsilon))$	yes
[4] PAR	O(1)	$O(\log(1/arepsilon))$	O(1)	yes
[5]States ladder	$O\left(\log^{>1.12}(1/arepsilon) ight)$			yes

- [1] C. M. Dawson and M. A. Nielsen, "The Solovay-Kitaev algorithm" [2005], [1995,1997][Any universal gate set]
- [2] A. G. Fowler, "Constructing arbitrary Steane code single logical qubit fault-tolerant gates" [2004][Clifford+T]
- [3] A. Y. Kitaev, A. H. Shen, and M. N. Vyalyi, Classical and Quantum Computation [1999][Clifford+T]
- [4] N. C. Jones et al. "Simulating chemistry efficiently on fault-tolerant quantum computers" [2012][Clifford+T]
- [5] G. Duclos-Cianci, K. M. Svore "A State Distillation Protocol to Implement Arbitrary Single-qubit Rotations" [2012][Clifford+T]





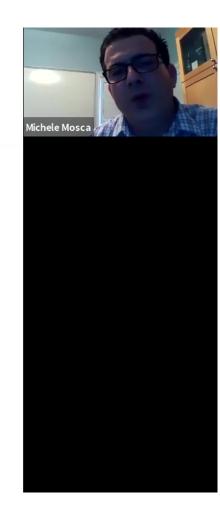
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Non-constructive results

- ▶ Exists gate set saturating bound $O(\log(1/\varepsilon))$ for SU(2):
 - A. Lubotzky, R. Phillips and P. Sarnak "Hecke operators and distributing points on the sphere I", [1986]
 - A. Lubotzky, R. Phillips and P. Sarnak "Hecke operators and distributing points on the sphere II", [1987]
- ▶ Exists gate set saturating bound $O(\log(1/\varepsilon))$ for SU(d):
 - A. W. Harrow, B. Recht, I. L. Chuang "Efficient Discrete Approximations of Quantum Gates", [2001]
- Any universal gate set with algebraic entries saturates bound $O(\log(1/\varepsilon))$ in SU(2)
 - ▶ J. Bourgain, A. Gamburd, "On the spectral gap for finitely generated subgroups of SU(2)", [2008]
- Any universal gate set with algebraic entries saturates bound $O(\log(1/\varepsilon))$ in SU(d)
 - ► J. Bourgain, A. Gamburd, "A Spectral Gap Theorem in SU(d)" [2011]

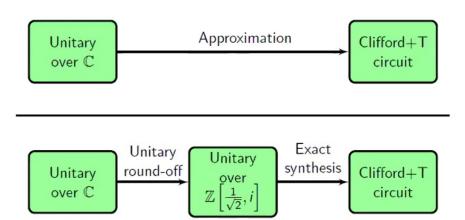
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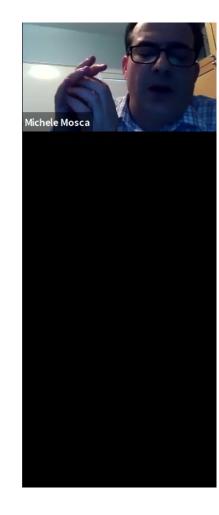
Pirsa: 20040005 Page 33/68

New approach to unitary approximation



$$\mathbb{Z}\left[\frac{1}{\sqrt{2}},i\right] := \left\{\left.\frac{1}{\sqrt{2}^n}\left(a + b\omega + c\omega^2 + d\omega^3\right)\right| a,b,c,d,n \in \mathbb{Z}\right\}, \omega = e^{i\pi/4}$$

[KMM13] V. Kliuchnikov, D. Maslov, and M. Mosca, "Fast and efficient exact synthesis of single qubit unitaries generated by Clifford and T gates", in Quantum Information and Computation. [arXiv:1206.5236]



Pirsa: 20040005 Page 34/68

Single qubit exact synthesis: main result

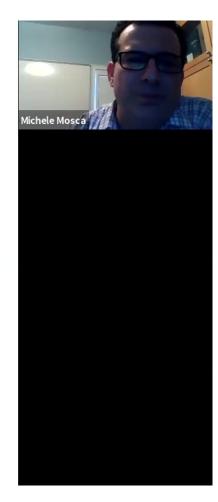
► Gate set
$$\mathscr{G}$$
: $H = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix}$, $T = \begin{pmatrix} 1 & 0 \\ 0 & e^{i\pi/4} \end{pmatrix}$
 $P = \begin{pmatrix} 1 & 0 \\ 0 & i \end{pmatrix}$, Pauli X, Y, Z

Theorem

Single qubit unitary U can be implemented exactly using $\mathscr G$ if and only if it has entries of the form:

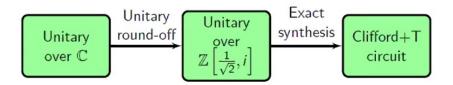
$$\frac{1}{\sqrt{2}^n} (a + b\omega + c\omega^2 + d\omega^3) a, b, c, d, n \in \mathbb{Z}$$
$$\omega = e^{i\pi/4}.$$

There exists an efficient algorithm that produces circuit that implements unitary with a minimal number of H and T gates.



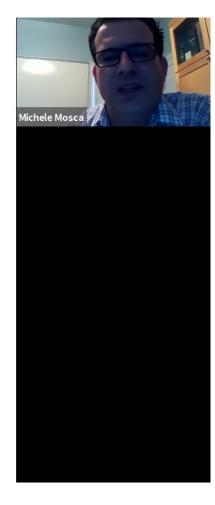
Pirsa: 20040005 Page 35/68

Rigorous efficient single qubit exact synthesis (with two ancilla qubits)



$$\begin{array}{ll} \text{Goal Approximate } \Lambda \left(e^{i\phi} \right) : \alpha \left| 0 \right\rangle + \beta \left| 1 \right\rangle \rightarrow \alpha \left| 0 \right\rangle + e^{i\phi} \beta \left| 1 \right\rangle \\ \text{Solution Approximately implement} \\ \alpha \left| 000 \right\rangle + \beta \left| 100 \right\rangle \rightarrow \alpha \left| 000 \right\rangle + e^{i\phi} \beta \left| 100 \right\rangle \\ \end{array}$$

- ▶ Find a unitary U that approximately prepares $e^{i\phi}|00\rangle$ starting from $|00\rangle$
- lacktriangle Execute U controlled on the first qubit



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Rounding off state $e^{i\phi} |00\rangle$

$$(\sin(\phi) + i\cos(\phi), 0, 0, 0)$$

$$\downarrow \downarrow$$

$$|v\rangle = \frac{1}{2^n} (\lfloor 2^n \sin(\phi) \rfloor + i \lfloor 2^n \cos(\phi) \rfloor, a + ib, c + id, 0)$$

Solving constraint $||v\rangle| = 1$

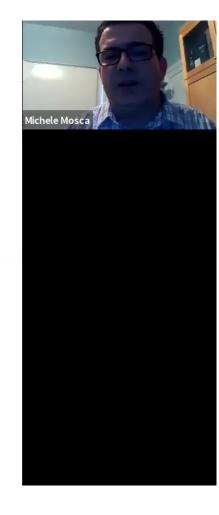
*:
$$a^2 + b^2 + c^2 + d^2 = 4^n - [2^n \sin(\phi)]^2 - [2^n \cos(\phi)]^2 = M$$

Complexity of finding a solution to *

$$O(\log^2(M)\log\log(M))$$

[*] M. O. Rabin and J. O. Shallit, "Randomized algorithms in number theory," Communications on Pure and Applied Mathematics, vol. 39, no. S1, pp. S239–S256, 1986.

[KMM13b] V. Kliuchnikov, D. Maslov, and M. Mosca, "Asymptotically optimal approximation of single qubit unitaries by Clifford and T circuits using a constant number of ancillary qubits", Physical Review Letters, **110**, 190502 (2013). [arXiv:1212.0822]



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Post-Quantum Cryptography

Volume 9606 of the series Lecture Notes in Computer Science pp 29-43.

Date: 04 February 2016

Applying Grover's Algorithm to AES: Quantum Resource Estimates

Markus Grassl, Brandon Langenberg, Martin Roetteler 🖾 , Rainer Steinwandt

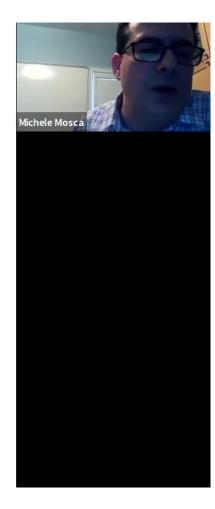
Our AES analysis, e.g. 192-bit AES: 5.9x10⁶ qubits, 2¹²¹ surface code cycles, 2^{137.5} total cost

Estimating the cost of generic quantum pre-image attacks on SHA-2 and SHA-3

Matthew Amy 1,4 , Olivia Di Matteo 2,4 , Vlad Gheorghiu 3,4 , Michele Mosca 3,4,5,6 , Alex Parent 2,4 , and John Schanck 3,4

		SHA-256	SHA 3-256
Grover	T-count	1.27×10^{44}	2.71×10^{44}
	T-depth	3.76×10^{43}	2.31×10^{41}
	Logical qubits	2402	3200
	Surface code distance	43	44
	Physical qubits	1.39×10^{7}	1.94×10^{7}
Distilleries	Logical qubits per distillery	3600	3600
	Number of distilleries	1	294
	Surface code distances	$\{33, 13, 7\}$	${33, 13, 7}$
	Physical qubits	5.54×10^5	1.63×10^{8}
Total	Logical qubits	$2^{12.6}$	2^{20}
	Surface code cycles	$2^{153.8}$	$2^{146.5}$
	Total cost	$2^{166.4}$	$2^{166.5}$

Table 3. Fault-tolerant resource counts for Grover search of SHA-256 and SHA3-256.



Bottom line for RSA-2048

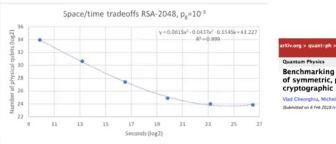


Fig. 4. RSA-2048 space/time tradeoffs with physical error rate per gate $p_g=10^{-3}$. The scale is logarithmic (base 2). Approximately $y(16.3987)\approx 1.72\times 10^8$ physical qubits are required to break the scheme in one day (24 hours). The number of T gates

quotes are required to oreast the scheme in one day (22 none). The manner of s_2 and in the circuit is 2.41×10^{12} , the corresponding number of logical qubits is 4098, and the total number of surface code cycles is 4.69×10^{14} . The classical security parameter



RSA-2048 - GIDNEY & EKERĂ UPDATES

Quantum Physics

How to factor 2048 bit RSA integers in 8 hours using 20 million noisy qubits

Craig Gidney, Martin Ekerå Gubmitted on 23 May 2019)

Guberniste on x seq xv xv x We significanly receive the cost of facialising integers and computing discrete legarithms over finite fields on a countries component by combining schoneques from Criffino-Hou 1996, Zelia 2006, receive 2012, Servi-Falest 2017, Servi 2019, Collary-Folder 2013, Gellery 2015 Wincountries of the component of the component of the component of the component of the contribution of the component of the contribution of the component of the contribution of the contribution of the contribution of the contribution of the component of the contribution of the component of the contribution of the contribution of the component of the contribution of the contribution of the component of the contribution of the component of the contribution of the con

no. of physical gubits reduced by ~10^2

is approximately 112 bits.

	Physical asymptote			Approach		Estimated costs			
Historical cost			Braction time		That Old you	Executive	Physical qubits	Espected runtime	Expected welcom
evolution at n - 2048		(microsconds)	(microscouds)	summetivity	stdategy	-transgr	(millions)	(days)	[megaqubindays]
Fowler et al. 2012 [6]	0.1%	4	0.1	planer	12081 T	magla threaded	21667	1.1	1100
O'Gorman et al. 2017 [18]	0.1%	to	1	arbitrary	block CCZ		230	3.7	850
Cheerghin et al. 2019 [10]	0.1%	0.2	6.1	planar	1100 T	stegle threaded	170	1	170
(ours) 2019 (1 factors)	0.135	1	10	ploner	L CC2	serial distillation	36	6	96
(irur) 2019 (1 thread)	0.1%	1	10	pluser	14 CCZ	single threafol	19	01.30	6.6
(cores) 2010 (porrallell)	0.150	4	10	mlamor	26 CVY	about the extended	20	0.31	K. 0-

TABLE II. Historical estimates of the expected costs of factoring n=2048 bit RSA integers, and the assumptions they used. Our spacetime volumes can be directly compared to the volume from Fowfer et al. (we achieve a 165x improvement), because Fowfer et al's continuate is dominated by distillation and changing the reaction time doesn't affect this volume. It is unclear how to compare O'Gorman et al.'s volume to ours, because of the difference in connectivity. Multiplying the volume from Gheorghiu et al. by 5, to account for the difference in cycle time, allows comparison to our volume (we achieve a 140x improvement). See Appendix B for details on each entry in this table.

Resource requirements keep going down.

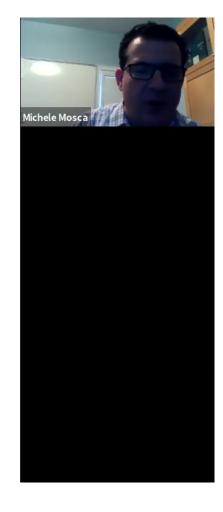




Cloud computing, payment systems, internet, IoT, eHealth, etc...

Secure web browsing, Auto-updates, VPN, Secure email, Blockchain, etc...

Cryptography:RSA, DSA, DH, ECDH, ECDSA,..., SHA, AES



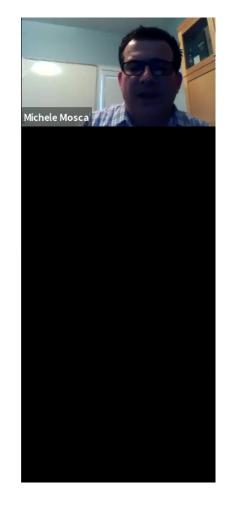
Pirsa: 20040005 Page 40/68



Cloud computing, payment systems, internet, IoT, eHealth, etc...

Secure web browsing, Auto-updates, VPN, Secure email, Blockchain, etc...

Cryptography:RS, SHA, AES



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Non-fault-tolerant quantum devices

i.e. quantum annealers, quantum simulators, NISQ

Not a known threat to cryptography

- Can they capture some of the power of quantum computation?
- Can they simulate themselves or similar systems faster/cheaper than conventional computers?
- Can they solve useful problems better than conventional devices?
- Can the same platforms be leveraged for fault-tolerant quantum computing?

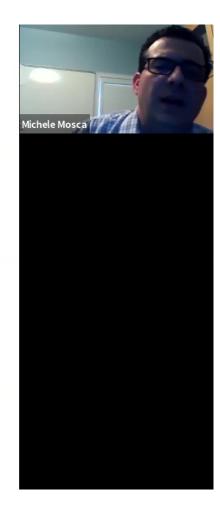
 "Similarly eithers is no proof today that it

"Similarly, although there is no proof today that imperfect quantum machines can compute fast enough to solve practical problems, that may change."



Commercialize early quantum technologies

Masoud Mohseni, Peter Read, Hartmut Neven and colleagues at Google's Quantum Al Laboratory set out investment opportunities on the road to the ultimate quantum machines.



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On some alternative quantum factoring approaches...

Factoring semi-primes with (quantum) SAT-solvers

Michele Mosca*1 and Sebastian R. Verschoor†2

and question the practical effectiveness of this approach for factoring large numbers. We find no evidence that this is a viable path toward factoring large numbers, even for scalable fault-tolerant quantum computers, as well as for various quantum annealing or other special purpose quantum hardware.

https://arxiv.org/pdf/1902.01448.pdf

On speeding up factoring with quantum SAT solvers

Michele Mosca¹, João Marcos Vensi Basso², and Sebastian R. Verschoor³

We present a SAT circuit that can be given to quantum SAT solvers such as annealers in order to perform this step of factoring. If quantum SAT solvers achieve any speedup over classical brute-force search, then our factoring algorithm is faster than the classical NFS.

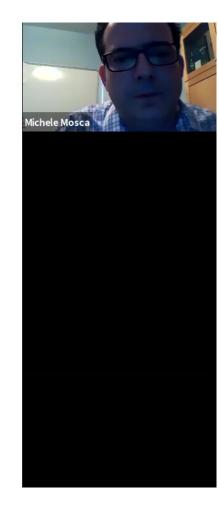
https://arxiv.org/pdf/1910.09592.pdf

A low-resource quantum factoring algorithm

Daniel J. Bernstein^{1,2}, Jean-François Biasse³, and Michele Mosca^{4,5,6}

heuristics, is $L^{p+o(1)}$ where p>1.9. The new time complexity is asymptotically worse than Shor's algorithm, but the qubit requirements are asymptotically better, so it may be possible to physically implement it sooner.

https://eprint.iacr.org/2017/352

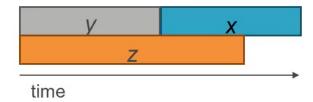


Do we need to worry *now*?

Depends on*:

- security shelf-life (x years)
- migration time (y years)
- collapse time (z years)

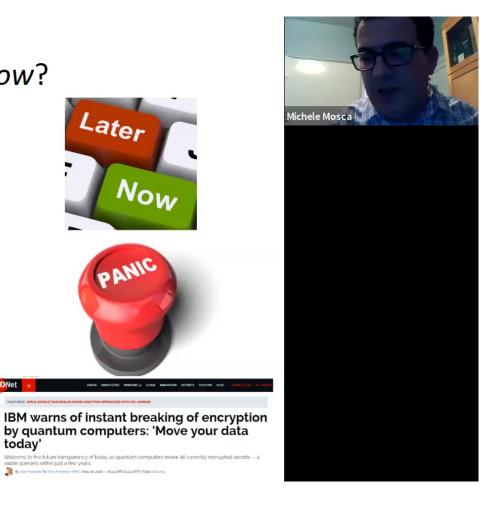
"Theorem": If x + y > z, then worry.



*M. Mosca: e-Proceedings of 1st ETSI Quantum-Safe Cryptography Workshop, 2013. Also http://eprint.iacr.org/2015/1075



by quantum computers: 'Move your data



today' scome to the future transparency of today as quantum co ble scenario within just a few years.

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A very recent milestone

"Quantum supremacy"

nature

Article | Published: 23 October 2019

Quantum supremacy using a programmable superconducting processor

Frank Arute, Kunal Arya, [...] John M. Martinis

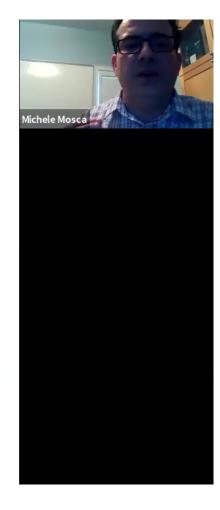
Nature **574**, 505–510 (2019) | Download Citation <u>¥</u>

2139 Altmetric | Metrics >>

Google claims it has achieved 'quantum supremacy' - but IBM disagrees

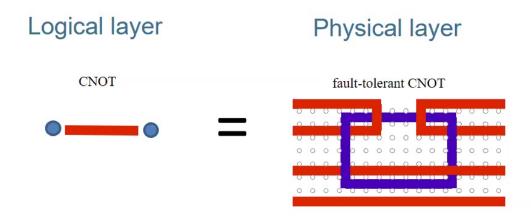
Task that would take most powerful supercomputer 10,000 years 'completed by quantum machine in minutes'

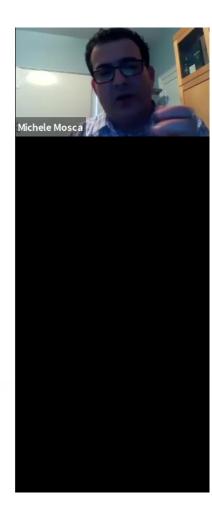




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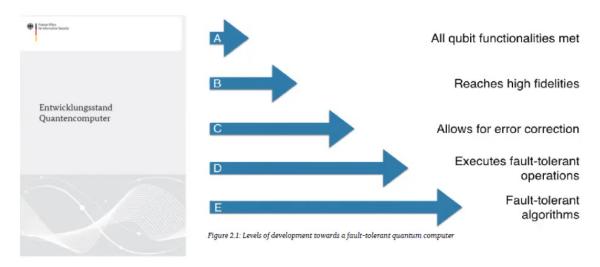
Critical Future Milestone: Scalable fault-tolerant logical qubits



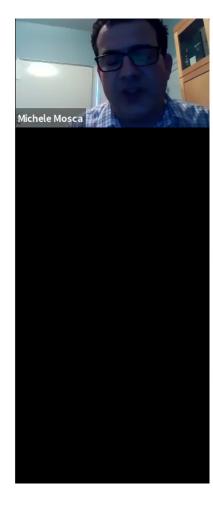


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Estimating 'z'?



https://www.bsi.bund.de/DE/Publikationen/Studien/Quantencomputer/quantencomputer.html (first draft in 2018; updated version 1.1 in 2019)



What is 'z'?

- Michele Mosca [Oxford, 1996]: "20 qubits in 20 years"
- Microsoft Research [October 2015]: "Recent improvements in control of quantum systems make it seem feasible to finally build a quantum computer within a decade".
- Michele Mosca ([NIST, April 2015], [ISACA, September 2015]): "1/7 chance of breaking RSA-2048 by 2026, ½ chance by 2031"
- Michele Mosca [London, September 2017]: "1/6 chance within 10 years"
- Simon Benjamin [London, September 2017]: Speculates that if someone is willing to "go Manhattan project" then "maybe 6-12 years"
- Michele Mosca [Seattle, November 2019]: 1/5 chance within 10 years

See also: https://globalriskinstitute.org/publications/quantum-threat-timeline/

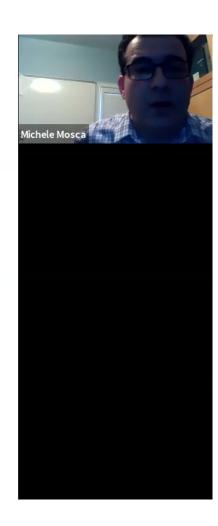
Michele Mosca

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Name	Institution
Scott Aaronson	University of Texas at Austin
Dorit Aharonov	The Hebrew University of Jerusalem
Dave Bacon	Google
Simon Benjamin	University of Oxford
Alexandre Blais	Université de Sherbrooke
Ignacio Cirac	Max Planck Institute of Quantum Optics
Bill Coish	McGill University
David DiVincenzo	Forschungszentrum Jülich
Runyao Duan	Institute for Quantum Computing, Baidu Research
Martin Ekerå	KTH Royal Institute of Technology and Swedish NCSA
Artur Ekert	University of Oxford and National University of Singapore

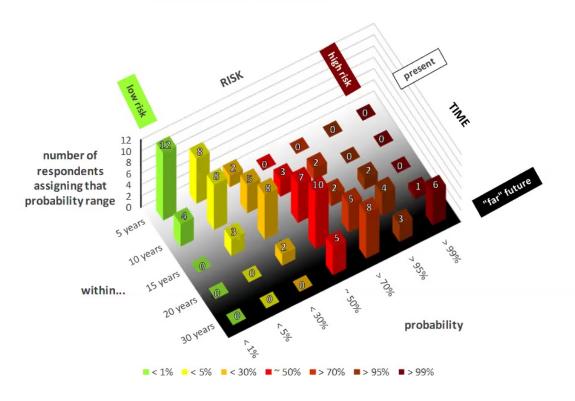
Name	Institution
Daniel Gottesman	Perimeter Institute for Theoretical Physics and Quantum Benchmark Inc
Jungsang Kim	Duke University
Ashley Montanaro	University of Bristol
Andrea Morello	UNSW Sydney
Yasunobu Nakamura	The University of Tokyo
Tracy Northup	University of Innsbruck
Peter Shor	Massachusetts Institute of Technology
Stephanie Simmons	Simon Fraser University
Krysta Svore	Microsoft
Frank Wilhelm- Mauch	Saarland University
Shengyu Zhang	Tencent

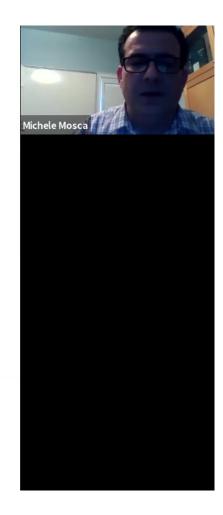




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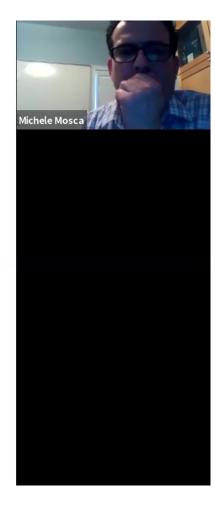
Likelihood of a digital quantum computer able to break RSA-2048 in 24 hours as function of time





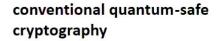
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Quantum-safe cryptographic tool-chest



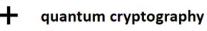
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Quantum-safe cryptographic tool-chest



a.k.a. Post-Quantum Cryptography or Quantum Resistant Algorithms



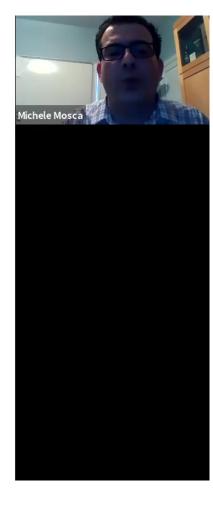






http://www.idquantique.com/p counting/clavis3-qkd-platform/

Both sets of cryptographic tools can work very well together in quantum-safe cryptographic ecosystem



Terminology

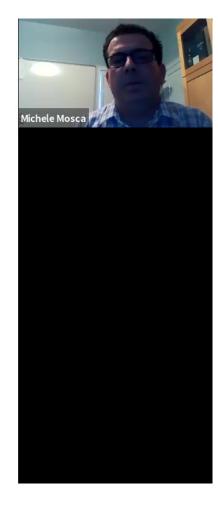
"Quantum-safe"

=

"designed to be safe in the era with large-scale quantum computers"

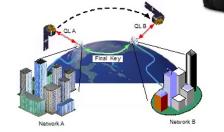
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conventional "post-quantum"/ "quantumresistant" cryptography + quantum cryptography

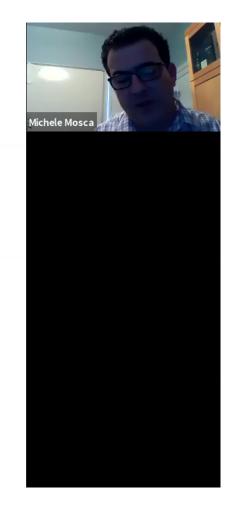


The ultimate key-establishment tool

Quantum physics guarantees the security of the cryptographic key



- •A quantum satellite in LEO can interconnect ground networks located anywhere on Earth.
- •Together with ground-based repeaters, we will eventually have a "quantum internet".



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Quantum key agreement

"Quantum key distribution" (QKD):

Over time, QKD evolves from:

Point-to-Point

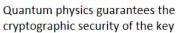
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Trusted Repeater Networks

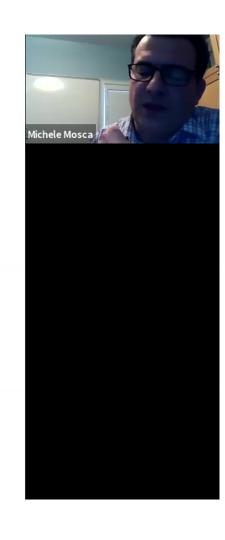
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Untrusted Repeater Networks



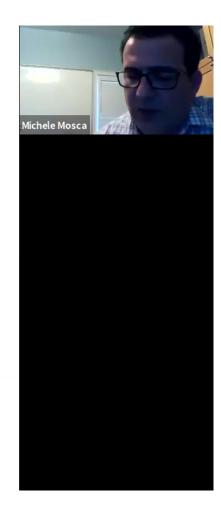






Go "where the puck is going, not where it has been"

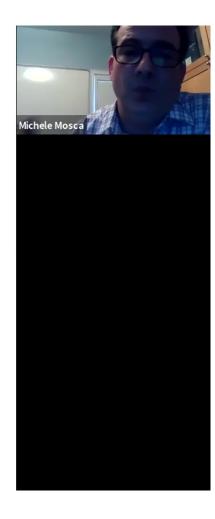




Quantum Internet – the Long Term Vision



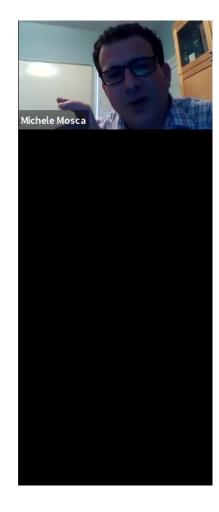
(Thanks to Thomas Jennewein)



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A historical fluke/opportunity

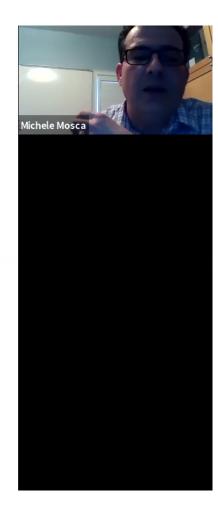
- Our current crypto infrastructure is not nearly as good as it could be.
- In practice it is nearly impossible to replace something "good enough" with something better.



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A historical fluke/opportunity

- Our current crypto infrastructure is not nearly as good as it could be.
- In practice it is nearly impossible to replace something "good enough" with something better.
- Given that we have no choice but to replace fundamental cryptography tools with something quantum-safe, the "toolbox" must be opened.



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Other comments: hybrid hybrid hybrid

Hybrid QKD key agreement: combine QKD key with classical public-key derived key, so breaking derived key requires breaking *both* QKD implementation and public-key derived key.

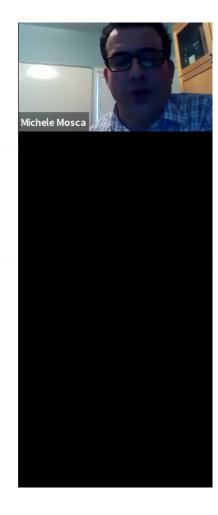
It must be done carefully, but done right it means adding QKD only *increases* your security.

(e.g.

maybe QKD isn't resilient to certain parameters being leaked

maybe QKD equipment or software is open to some sort of sidechannel or other compromise that the post-quantum implementation isn't

certification/validation issues (like FIPS 140-2))

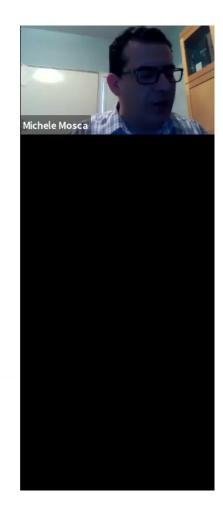


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How easy is it to evolve from one cryptographic algorithm to a quantum-secure one?

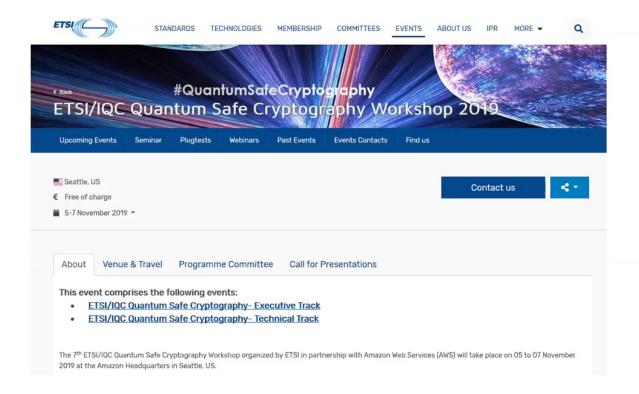
Are the standards and practices ready?

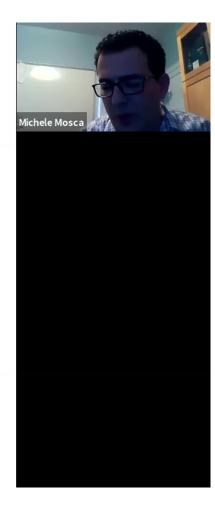




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https://www.etsi.org/events/1607-etsi-iqc-quantum-safe-cryptography-workshop-2019





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So what do we do about it now?

"Execution is 90% planning and 10% doing"







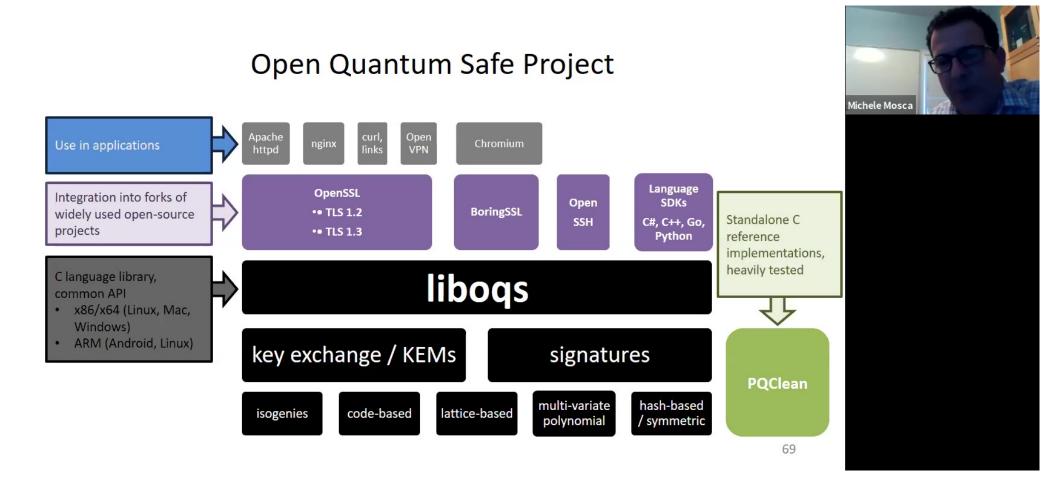
We don't get to call a "time-out" if we're not ready!



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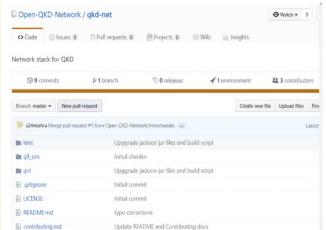
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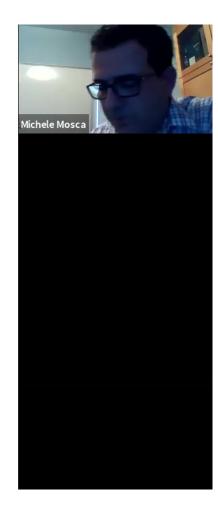
Pirsa: 20040005 Page 65/68

Open QKD Net: The Open Source Project

KMS/QNL/QLL-Sim Software Hosted at Github.com, managed by IQC team Source code license: MIT License QKD technology teams can use the system to issue keys to users with little effort Program language: Java Sample applications written in C





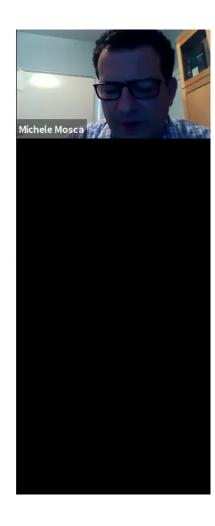


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Historic opportunity







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Thank you!

Comments, questions and feedback are very welcome.

2

Michele Mosca
Associate Faculty, Perimeter Institute for
Theoretical Physics
mmosca@perimeterinstitute.ca
Professor, Faculty of Mathematics
Co-Founder, Institute for Quantum Computing,
University of Waterloo www.iqc.ca/~mmosca
mmosca@uwaterloo.ca

CEO, evolutionQ Inc. @evolutionQinc michele.mosca@evolutionq.com

Co-founder, softwareQ Inc. softwareq.ca



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