

Title: No-free-information principle in general probabilistic theories

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Abstract: <p>In quantum theory, the no-information-without-disturbance and no-free-information principles express that those observables that do not disturb the measurement of another observable and those that can be measured jointly with any other observable must be trivial, i.e., coin tossing observables. We show that in the framework of general probabilistic theories these principles do not hold in general. In this way, we obtain characterizations of the probabilistic theories where these principles hold and we show that the two principles are not equivalent. </p>

No-free-information principle in general probabilistic theories

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The three principles

1. No broadcasting

You can't make two coffes out of one.



The three principles

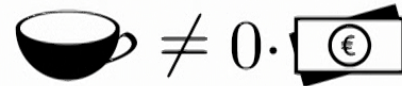
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2. No information without disturbance

You can't have the coffee for free.



The three principles

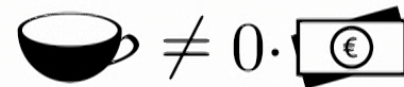
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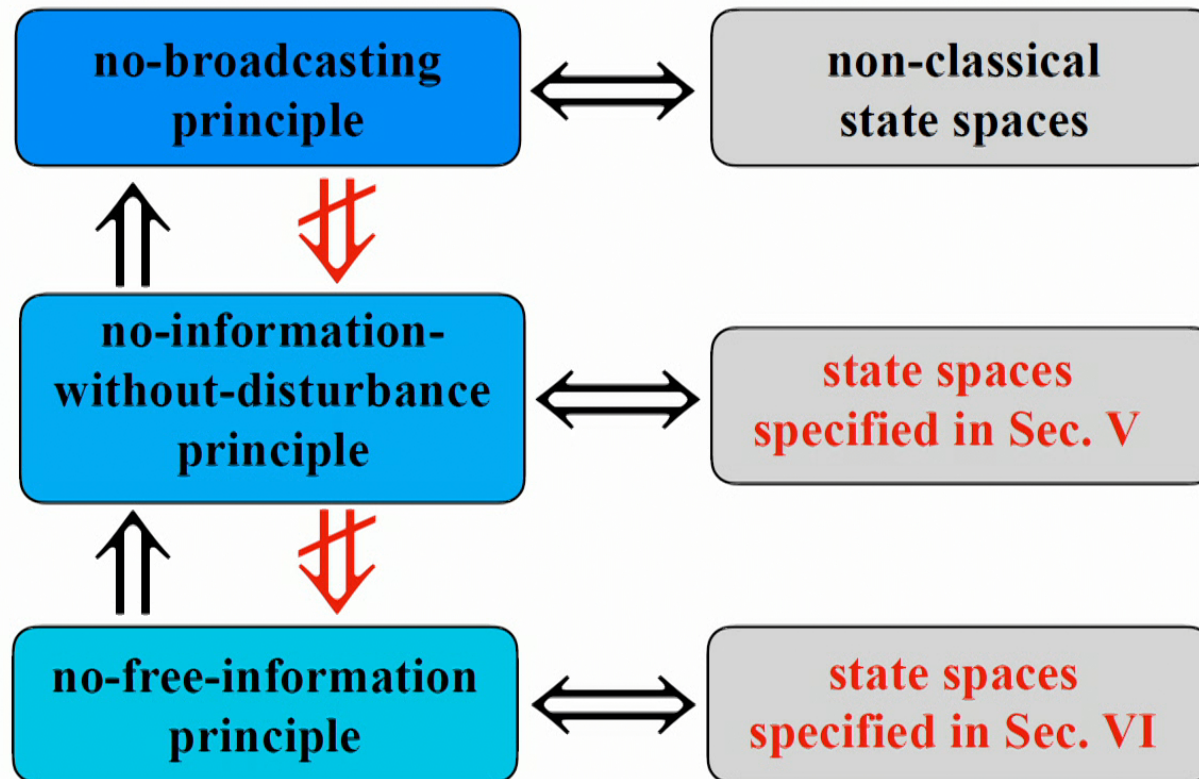


3. No free information

You can't have the coffee for free even if you pay for the lunch.



The results



GPT refresher

- K : state space, compact convex subset of \mathbb{R}^n
- $A(K)$: linear space of affine functions $f : K \rightarrow \mathbb{R}$
- $A(K)^+$: cone of positive affine functions, $f(x) \geq 0, \forall x \in K$
- $E(K)$: effect algebra of affine functions $f : K \rightarrow [0, 1]$
- $A(K)^*$: dual vector space to $A(K)$
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In quantum theory:

- K : density matrices over finite dimensional Hilbert space
- $A(K)$: self-adjoint operators, $A(\rho) = \text{Tr}(\rho A)$
- $A(K)^+$: cone of positive semi-definite matrices, $A \geq 0$
- $E(K)$: set of effects, $0 \leq A \leq \mathbb{1}$
- $A(K)^*$: self-adjoint operators
- $A(K)^{*+}$: cone of positive semi-definite matrices, $A \geq 0$

Measurements

Ω : sample space, finite measurable set of outcomes

$\mathcal{P}(\Omega)$: set of probability measures on Ω , simplex

Definition

Measurement is an affine map

$$m : K \rightarrow \mathcal{P}(\Omega).$$

Let $A \subset \Omega$, measurable

$m(x; A)$ = measure of the set A with respect to measure $m(x)$

Compatibility of measurements

Ω_1, Ω_2 sample spaces

Definition

Measurements

$$m_1 : K \rightarrow \mathcal{P}(\Omega_1)$$

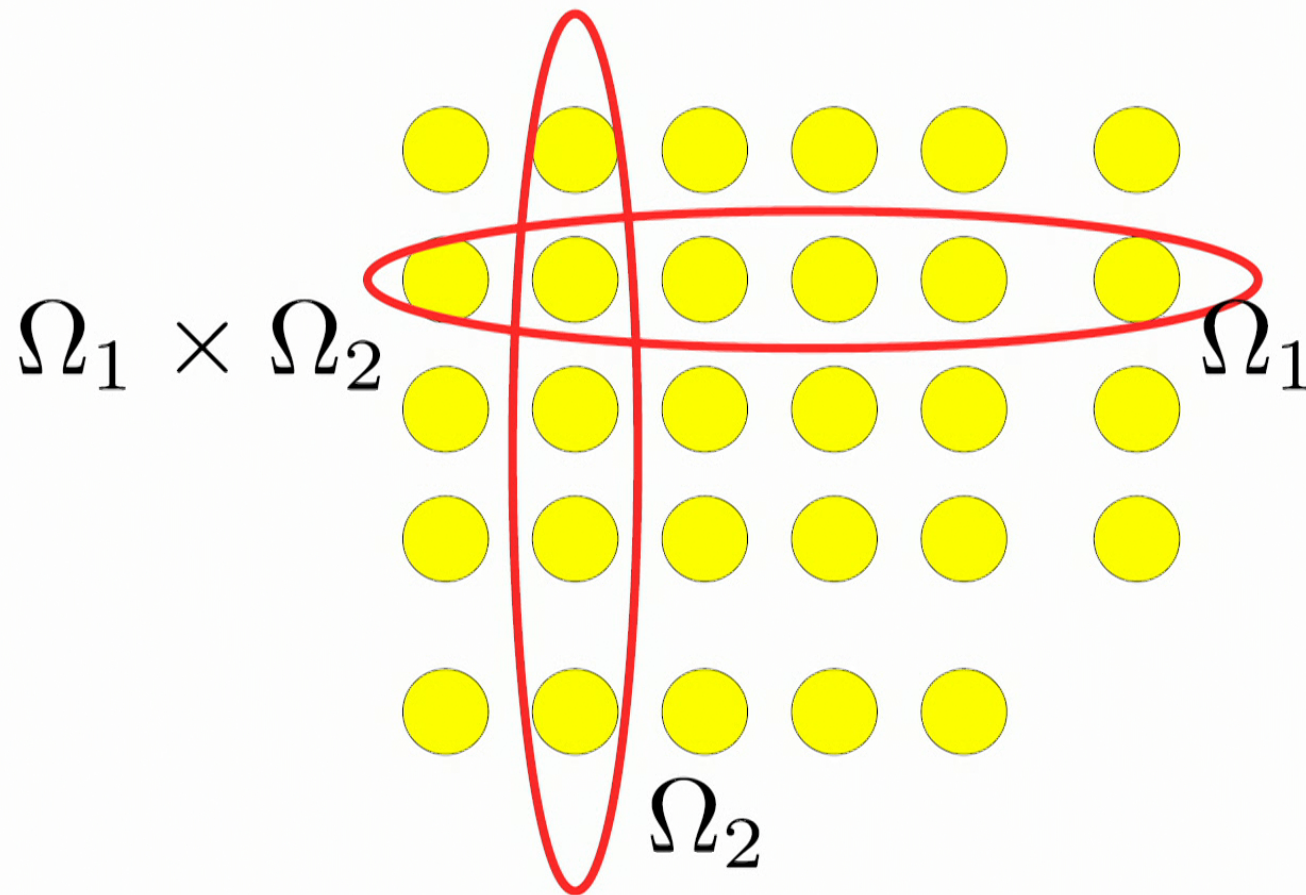
$$m_2 : K \rightarrow \mathcal{P}(\Omega_2)$$

are compatible if there exists $m : K \rightarrow \mathcal{P}(\Omega_1 \times \Omega_2)$ such that for all $A_1 \subset \Omega_1, A_2 \subset \Omega_2$ and $\forall x \in K$

$$m_1(x; A_1) = m(x; A_1 \times \Omega_2),$$

$$m_2(x; A_2) = m(x; \Omega_1 \times A_2).$$

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$$\left. \begin{matrix} E_i \\ F_j \end{matrix} \right\} \text{POVMs}$$

$$H_{ij}$$

$$E_i = \sum_j H_{ij}$$

$$F_j = \sum_i H_{ij}$$

Direct convex hull

Definition

Let K_A, K_B be state spaces, then $K_A \oplus K_B$ is a state space of ordered and weighted pairs of states, i.e.

$$K_A \oplus K_B = \{(\lambda x, (1 - \lambda)y) : x \in K_A, y \in K_B, \lambda \in [0, 1]\}.$$

$$\begin{matrix} E_i \\ F_j \end{matrix} \} \text{povms}$$

$$E_i = \sum_j H_{ij}$$

$$F_j = \sum_i H_{ij}$$

$$\left(\begin{array}{c|c} \lambda s_1 & 0 \\ \hline 0 & (1-\lambda) s_2 \end{array} \right)$$

No information without disturbance

Theorem

There exists nontrivial observable compatible with the identity channel if and only

$$K = \bigoplus_{i=1}^n K_i.$$

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Let τ be a coin-toss measurement, then m is compatible with the identity channel if and only if $\lambda m + (1 - \lambda)\tau$ is compatible with the identity channel for all $\lambda \in [1, 0)$.

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Noise does not help!

Some classification

Proposition

Let $\dim(K) = 2$, then $K = K_1 \oplus K_2$ if and only if K is a triangle.

Proof.

By counting dimensions. \square

$$\begin{aligned} E_i & \text{ Divergences} \\ F_j & \\ H_{ij} & \\ E_i &= \sum_j H_{ij} \\ F_j &= \sum_i H_{ij} \end{aligned}$$

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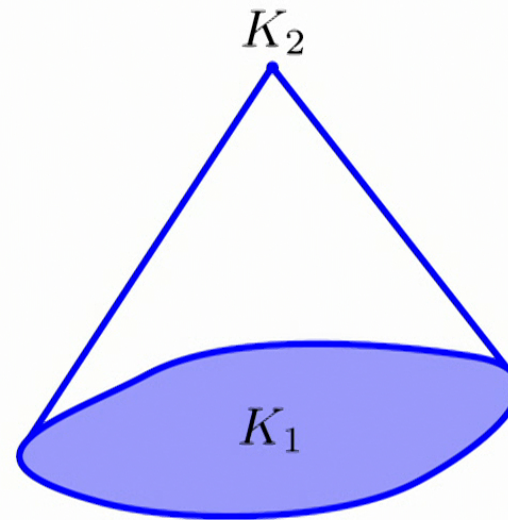
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Proposition

*Let $\dim(K) = 3$, then
 $K = K_1 \oplus K_2$ if and only if
 $K_2 = \{x\}$ and K is
pyramid-shaped.*



Post-processings and convex combinations

Let m_1, m_2 be measurements

$$m_1 : K \rightarrow \mathcal{P}(\Omega_1)$$

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Post-processing is a map $\nu : \mathcal{P}(\Omega_2) \rightarrow \mathcal{P}(\Omega_1)$ that allows us to construct measurement

$$m'_2 = \nu \circ m_2 : K \rightarrow \mathcal{P}(\Omega_1)$$

Post-processing is a form of order that gives rise to an equivalence.

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Definition

For $\lambda \in [0, 1]$ we define convex combination of measurements as the map

$$(\lambda m_1 + (1 - \lambda) m'_2) : K \rightarrow \mathcal{P}(\Omega_1)$$

Simulability of measurements

Definition

Measurements m_1, \dots, m_k simulate a measurement m if there are

- post-processing ν_1, \dots, ν_k
- numbers $\lambda_1, \dots, \lambda_k$, s.t. $\lambda_i \in [0, 1]$ for $i = 1, \dots, k$,
 $\sum_{i=1}^k \lambda_i = 1$

such that

$$m = \sum_{i=1}^k \lambda_i \nu_i \circ m_i$$

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Definition

We say that a measurement m is simulation irreducible if it can be simulated only by post-processing equivalent measurements.

No free information

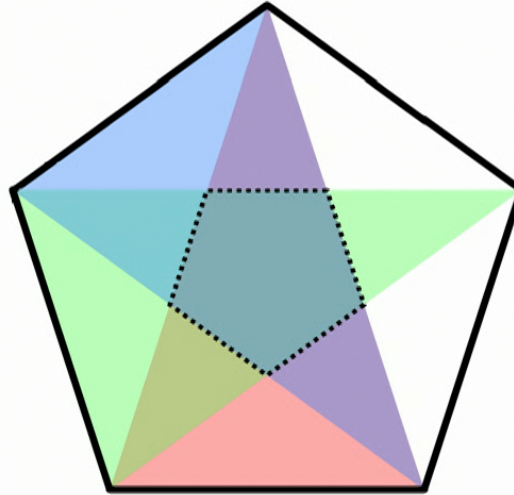
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A measurement m is compatible with every other measurement on K if and only if it is compatible with every simulation irreducible measurement on K .

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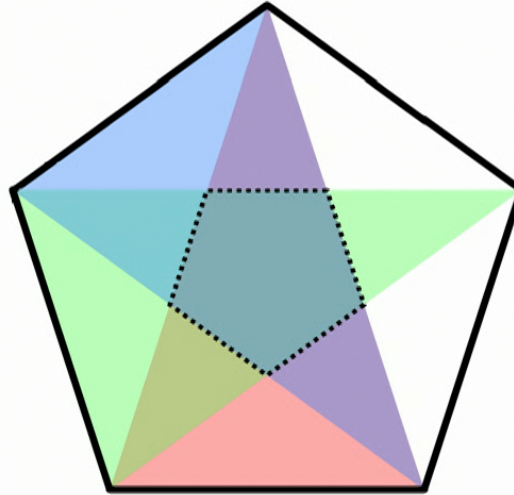
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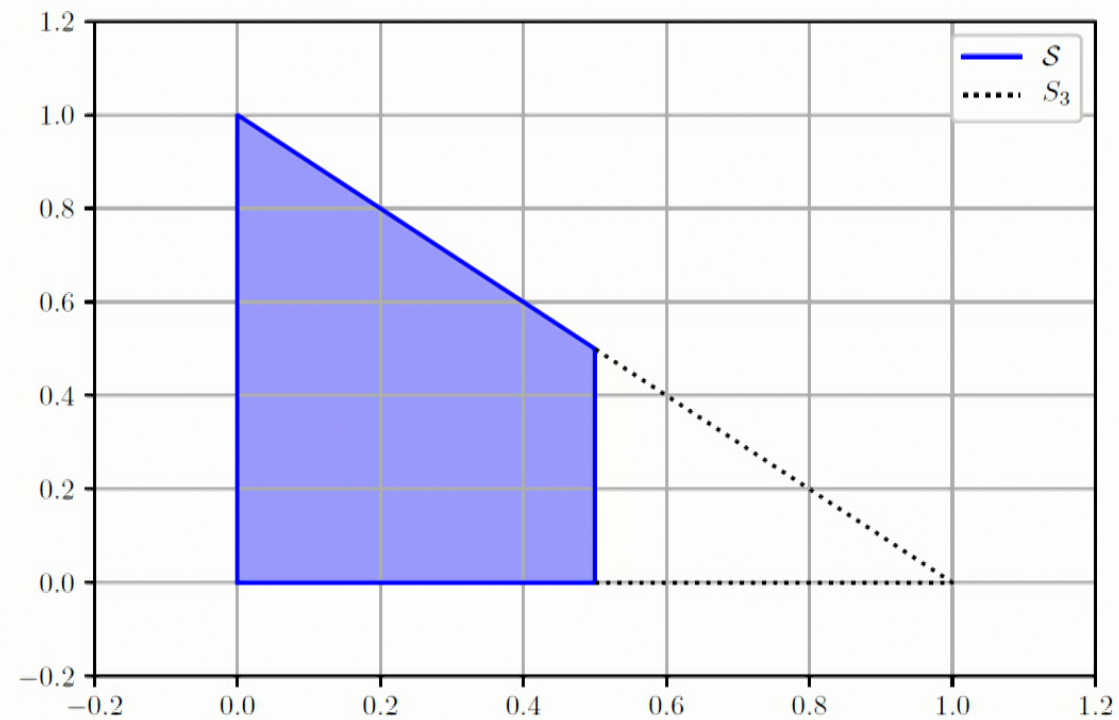
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Noise helps! (sometimes)

Example: measurement compatible with all measurements

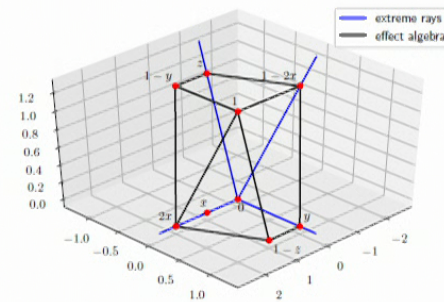
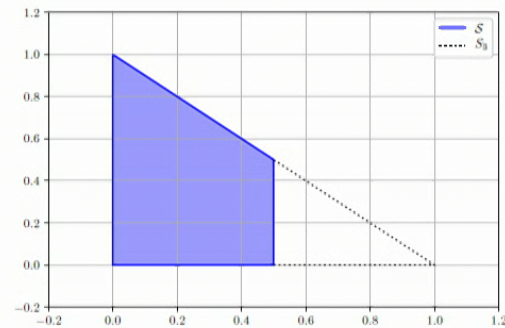


Example: measurement compatible with all measurements

Simulation irreducible measurements:

$$m_2 = 2x \otimes \delta_1 + (1 - 2x) \otimes \delta_2$$

$$m_3 = x \otimes \delta_1 + y \otimes \delta_2 + (1 - x - y) \otimes \delta_3$$



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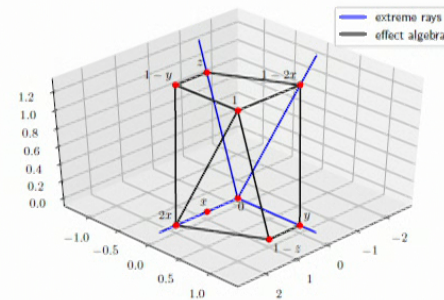
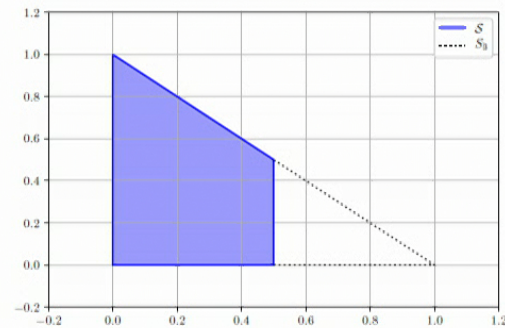
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Measurement compatible with every other measurement

$$m = x \otimes \delta_1 + (1 - x) \otimes \delta_2$$



Example: measurement compatible with all measurements

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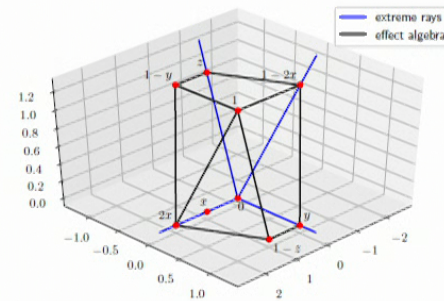
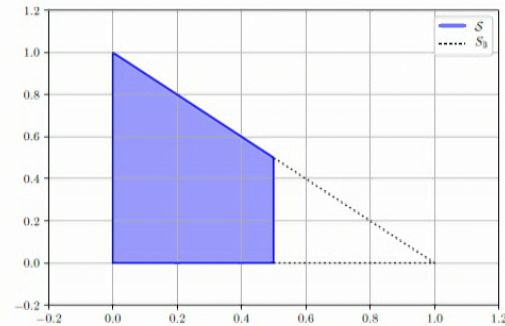
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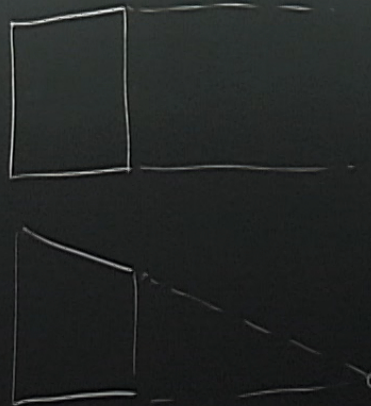
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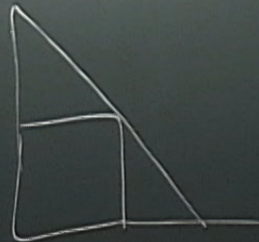
The joint measurement

$$m'_2 = x \otimes \delta_1 + x \otimes \delta_2 + (1 - x) \otimes \delta_3$$



$$\left(\begin{array}{c|c} \lambda s_1 & 0 \\ \hline 0 & (1-\lambda) s_2 \end{array} \right)$$





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Thank you for your attention!

Questions?
Comments?