

Title: Infinite composite systems and cellular automata in operational probabilistic theories

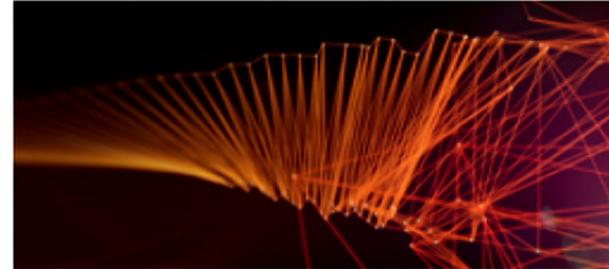
Date: Jul 30, 2018 11:30 AM

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Abstract:

Cellular automata are a central notion for the formulation of physical laws in an abstract information-theoretical scenario, and lead in recent years to the reconstruction of free relativistic quantum field theory. In this talk we extend the notion of a Quantum Cellular Automaton to general Operational Probabilistic Theories. For this purpose, we construct infinite composite systems, illustrating the main features of their states, effects and transformations. We discuss the generalization of the concepts of homogeneity and locality, in a framework where space-time is not a primitive object. We show that homogeneity leads to a Cayley graph structure of the memory array, thus proving the universality of the connection between homogeneity and discrete groups. We conclude illustrating the special case of Fermionic cellular automata, discussing three relevant examples: Weyl and Dirac quantum walks, the Thirring automaton and the simplest families of automata on finite graphs.

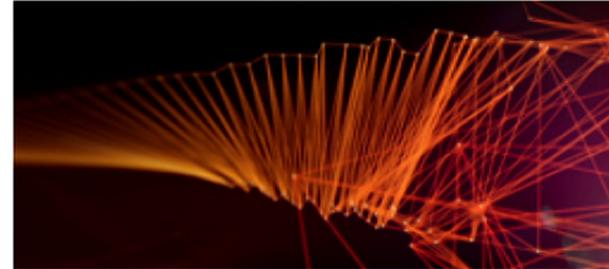
FOUNDATIONS OF QUANTUM MECHANICS
Perimeter Institute, July 30 2018



Infinite composite systems and cellular automata in operational probabilistic theories

Paolo Perinotti
Dipartimento di Fisica,
Università di Pavia



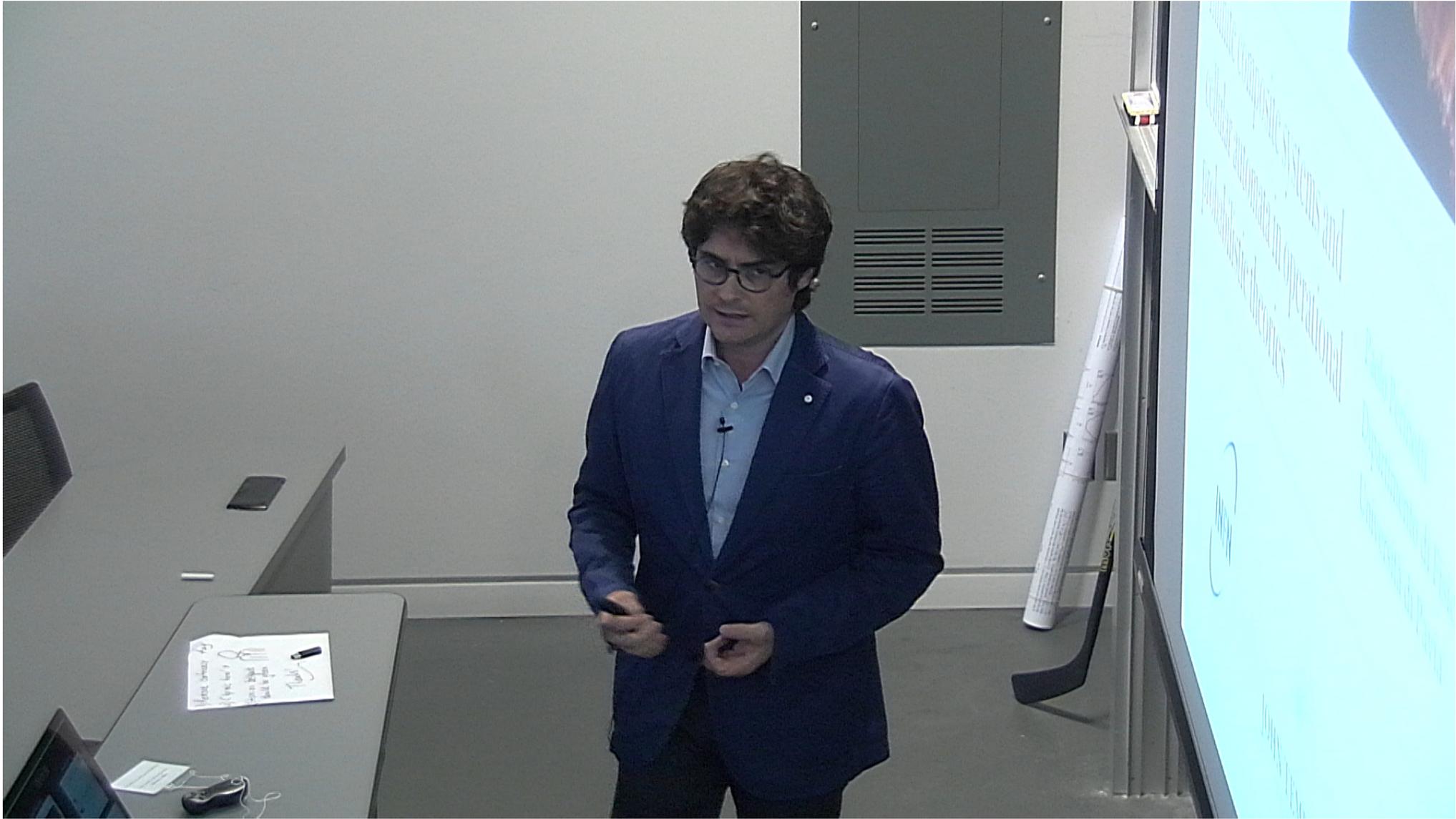


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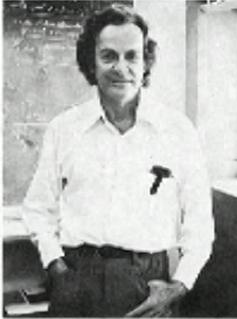
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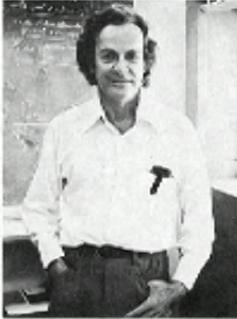
Digital Universe



“I want to talk about the possibility that there is to be an exact simulation, that the computer will do exactly the same as nature.”

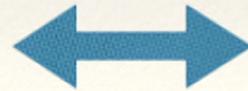
R. Feynman, *Int. J. Theo. Ph.* **21**, 467 (1982)

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Physical law

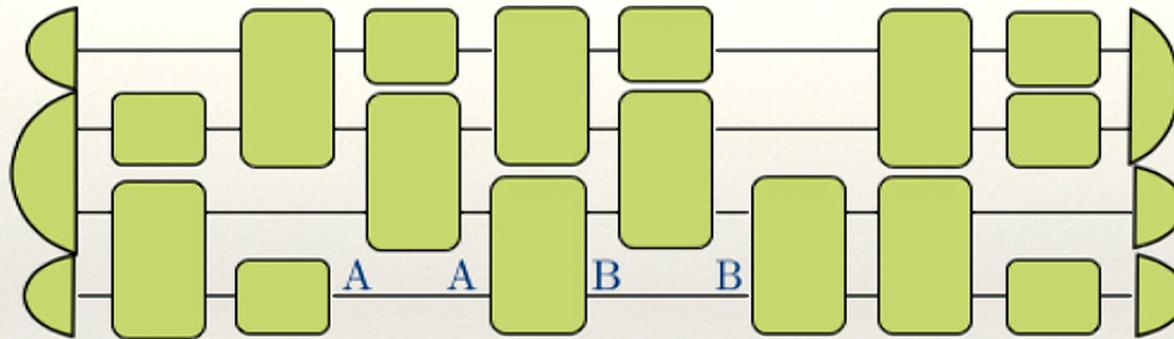


Algorithm

R. Feynman, *Int. J. Theo. Ph.* **21**, 467 (1982)

Operational Language

- ❖ Operational theory: tests with composition rules

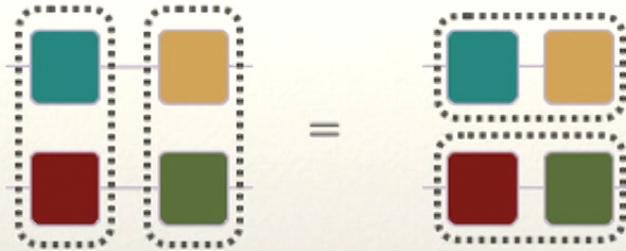


$$\text{---} \square \text{---} = \text{---} \{\mathcal{F}_i\}_{i \in J} \text{---}$$

G. Chiribella, G. M. D'Ariano, and PP, Phys. Rev. A **81**, 062348 (2010)

Main property

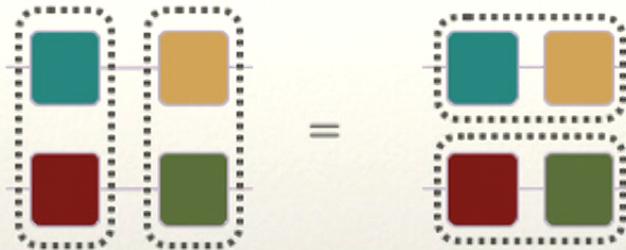
❖ Most important rule:



$$(A \otimes B)(C \otimes D) = (AC) \otimes (BD)$$

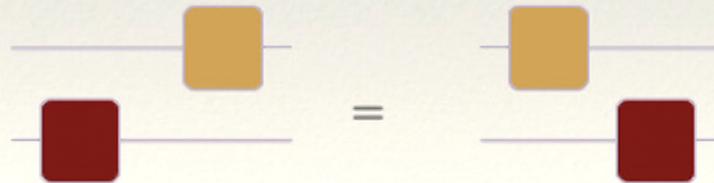
Main property

- ❖ Most important rule:



$$(A \otimes B)(C \otimes D) = (AC) \otimes (BD)$$

- ❖ Consequence:



Probabilistic theories

Every test of type $I \rightarrow I$ is a probability distribution

$$\rho_i \text{---} a_j = \Pr(a_j, \rho_i)$$

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Probabilistic theories

Every test of type $I \rightarrow I$ is a probability distribution $\rho_i \text{---} a_j = \Pr(a_j, \rho_i)$

States are functionals on effects and vice-versa

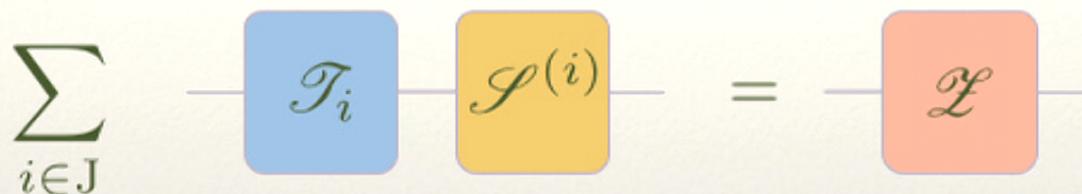


Real vector spaces $\text{St}_{\mathbb{R}}(A), \text{Eff}_{\mathbb{R}}(A)$

G. Chiribella, G. M. D'Ariano, and PP, Phys. Rev. A **81**, 062348 (2010)

Causal theories

- ❖ Possibility of arbitrary conditional tests



- ❖ Causality implies no backward signalling

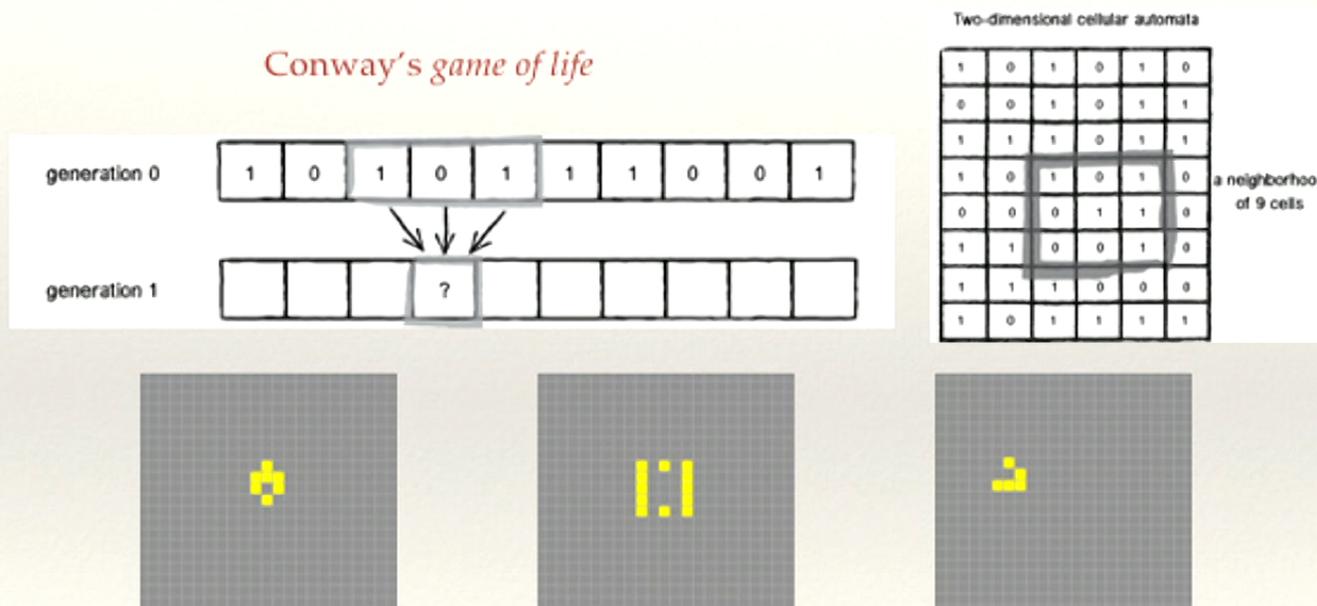
$$p_a(\rho_i) := \sum_j \left(\rho_i \begin{array}{c} \text{A} \\ \hline a_j \end{array} \right) = p(\rho_i)$$

$$\sum_j \begin{array}{c} \text{A} \\ \hline a_j \end{array} = \begin{array}{c} \text{A} \\ \hline e \end{array}$$

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Cellular Automata

J. Von Neumann and A. W. Burks, "Theory of self-reproducing automata" 1966

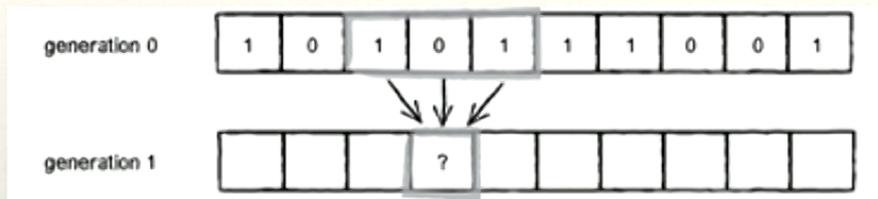


<http://web.stanford.edu/~cdebs/GameOfLife/>

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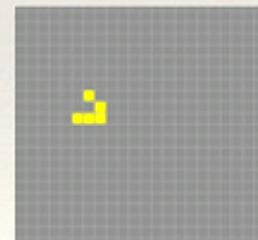
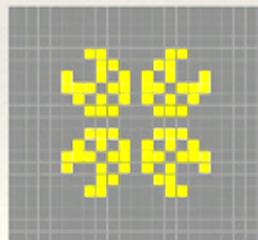
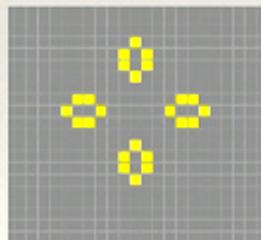
Conway's game of life



Two-dimensional cellular automata

1	0	1	0	1	0
0	0	1	0	1	1
1	1	1	0	1	1
1	0	1	0	1	0
0	0	0	1	1	0
1	1	0	0	1	0
1	1	1	0	0	0
1	0	1	1	1	1

a neighborhood of 9 cells

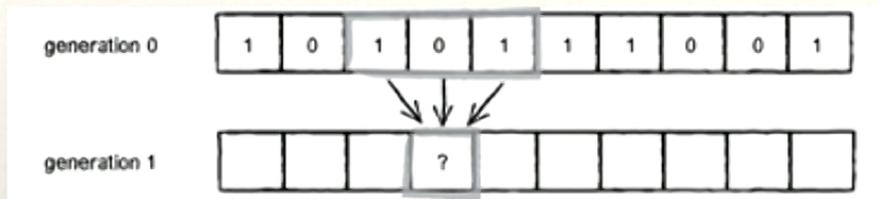


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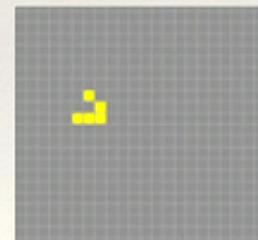
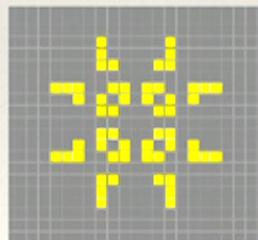
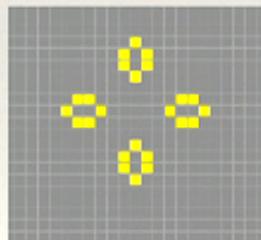
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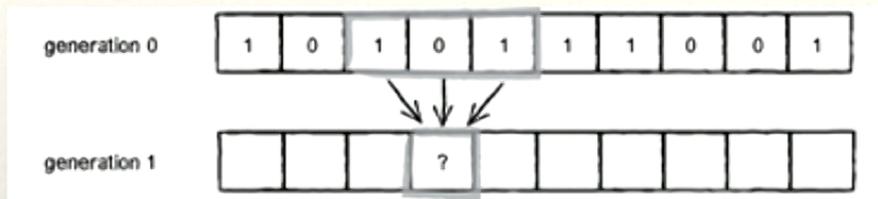


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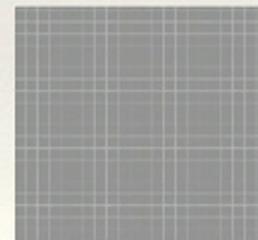
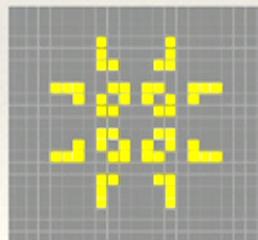
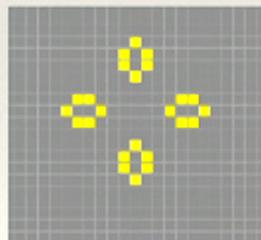
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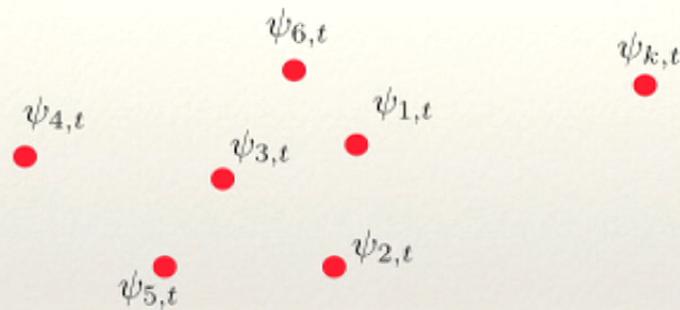
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Linear Fermionic Cellular Automata

$$\{\varphi_i^\dagger, \varphi_j\} = \delta_{ij} I, \quad \{\varphi_i, \varphi_j\} = 0$$

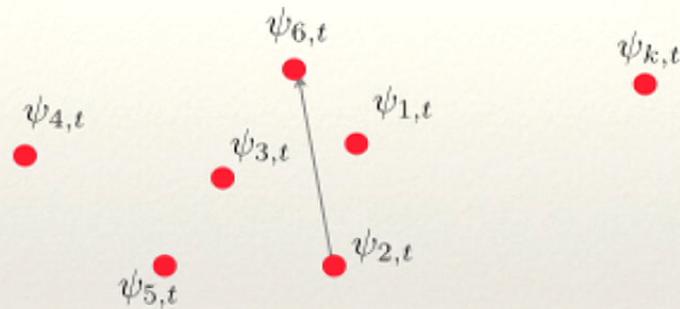


$$\psi_{i,t+1} = \sum_{j \in N_i} A_{i,j} \psi_{j,t} \quad A_{ij} \in M_{s_i \times s_j}$$

G. M. D'Ariano and PP, Phys. Rev. A **90**, 062106 (2014).

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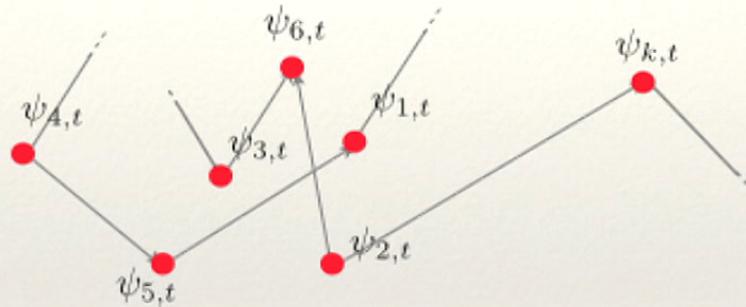


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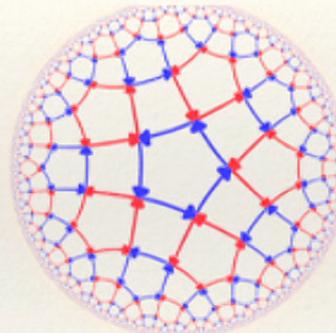
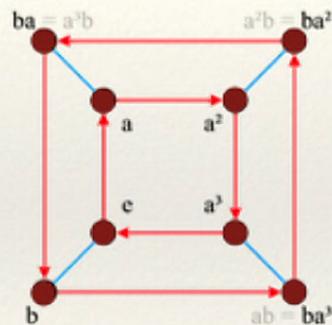
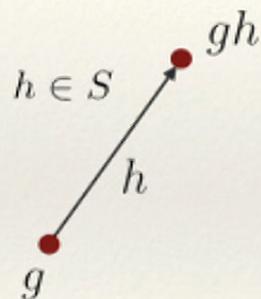


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Homogeneity and Cayley Graphs

- ◇ Homogeneity: the memory array is structured as a Cayley graph



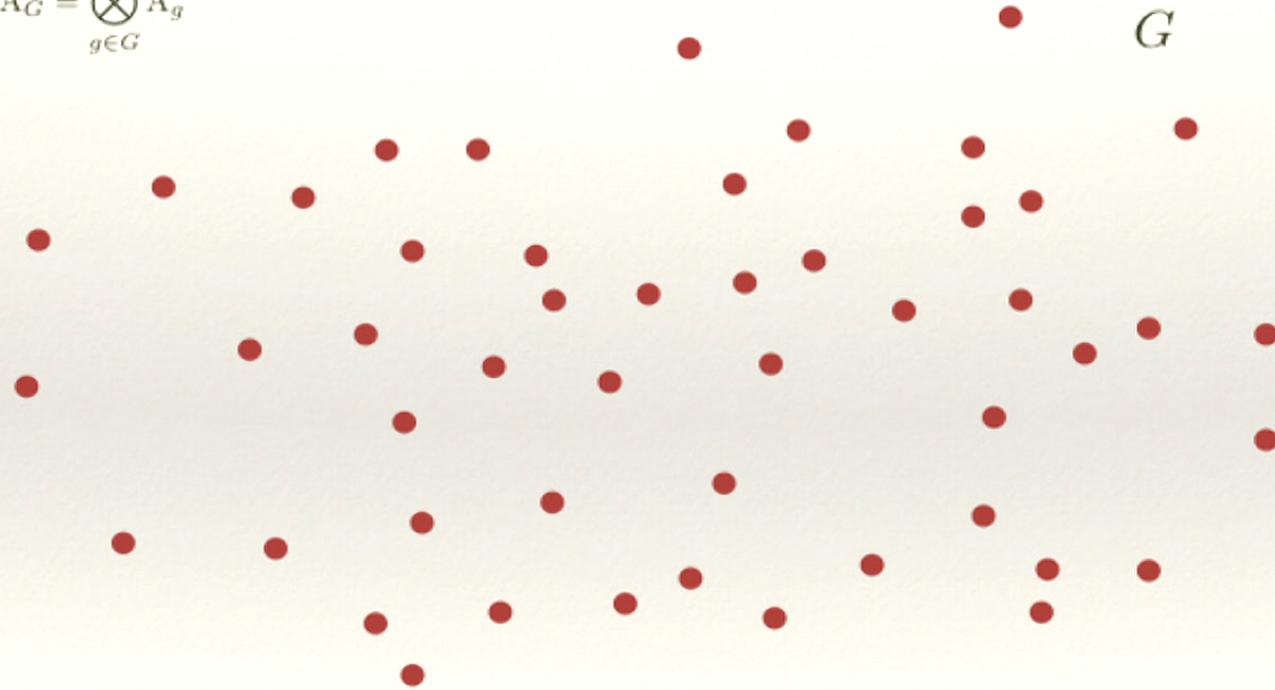
G. M. D'Ariano and PP, Phys. Rev. A **90**, 062106 (2014)
G. M. D'Ariano and PP, Front. Phys. **12**(1), 120301 (2017).

In general OPTs

- ❖ We do not have Hilbert spaces - only state / effect spaces
- ❖ No field operators - only transformations
 - ❖ There is no notion of a linear automaton
- ❖ Can we generalise the above construction?
- ❖ Are our results generic or specific of FQT?

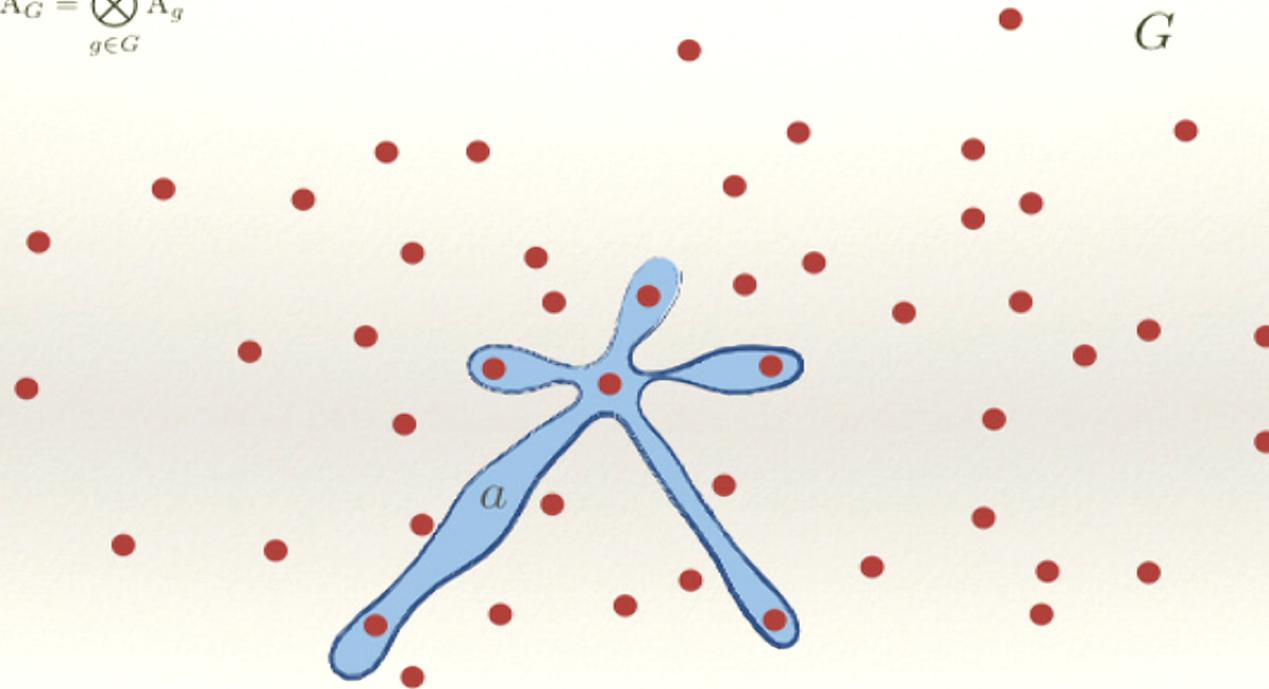
(Quasi-)Local effects

$$A_G = \bigotimes_{g \in G} A_g$$



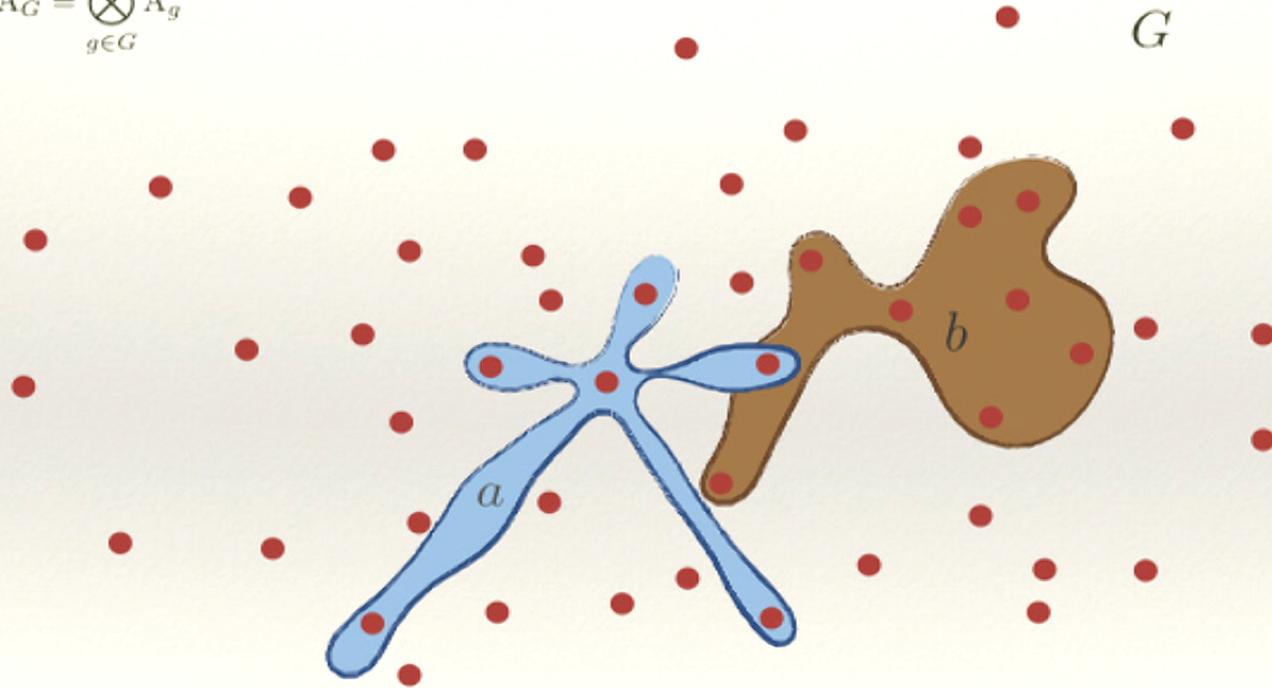
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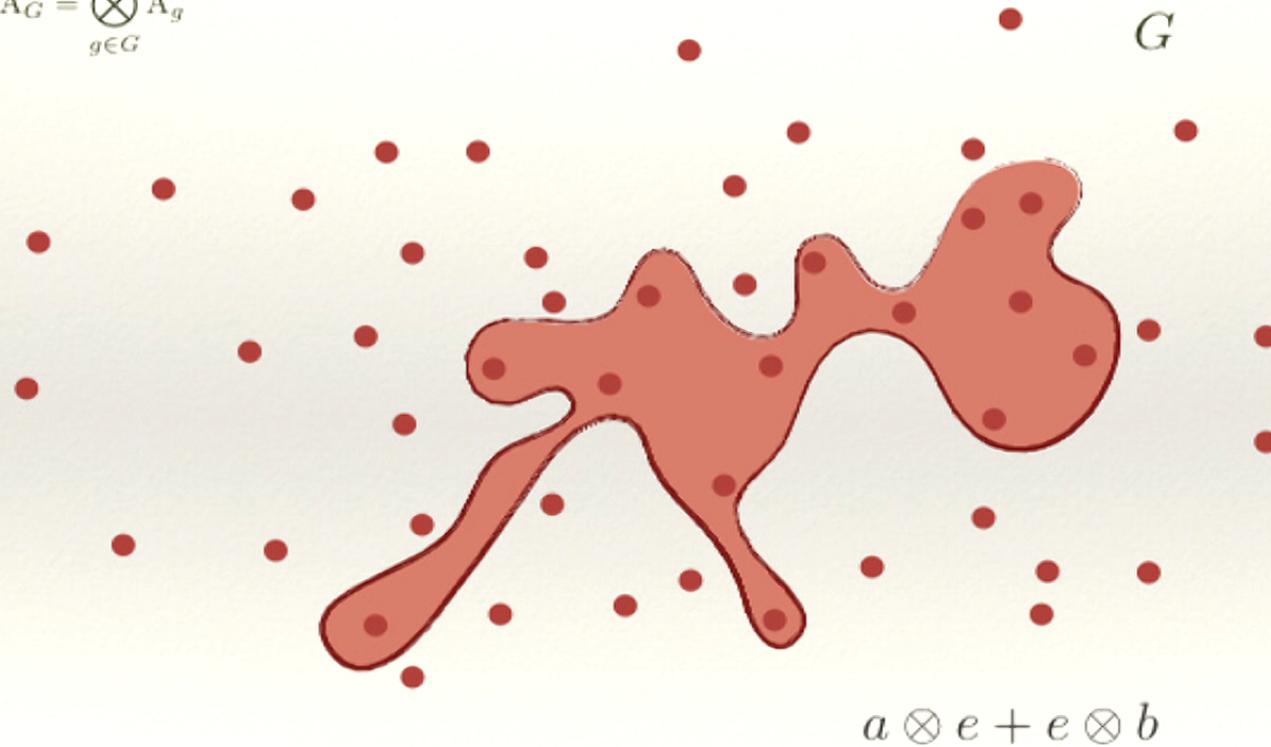
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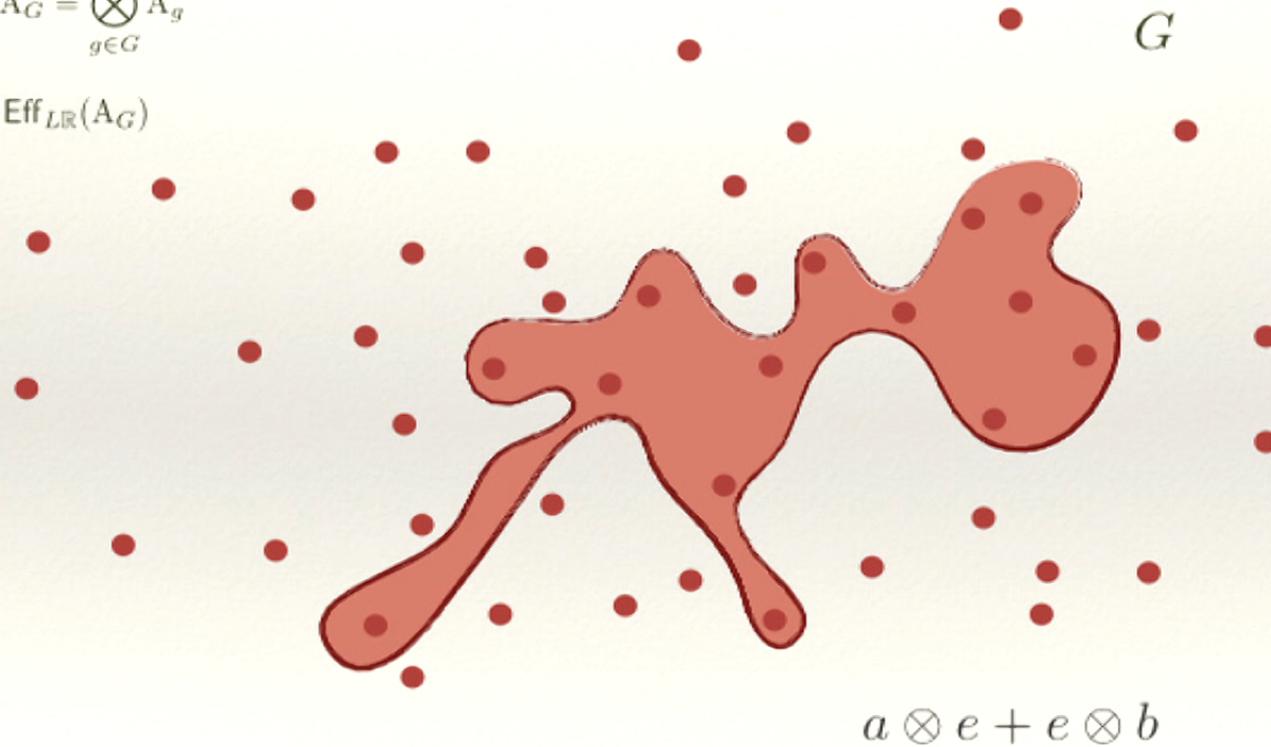
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$$\text{Eff}_{LR}(A_G)$$

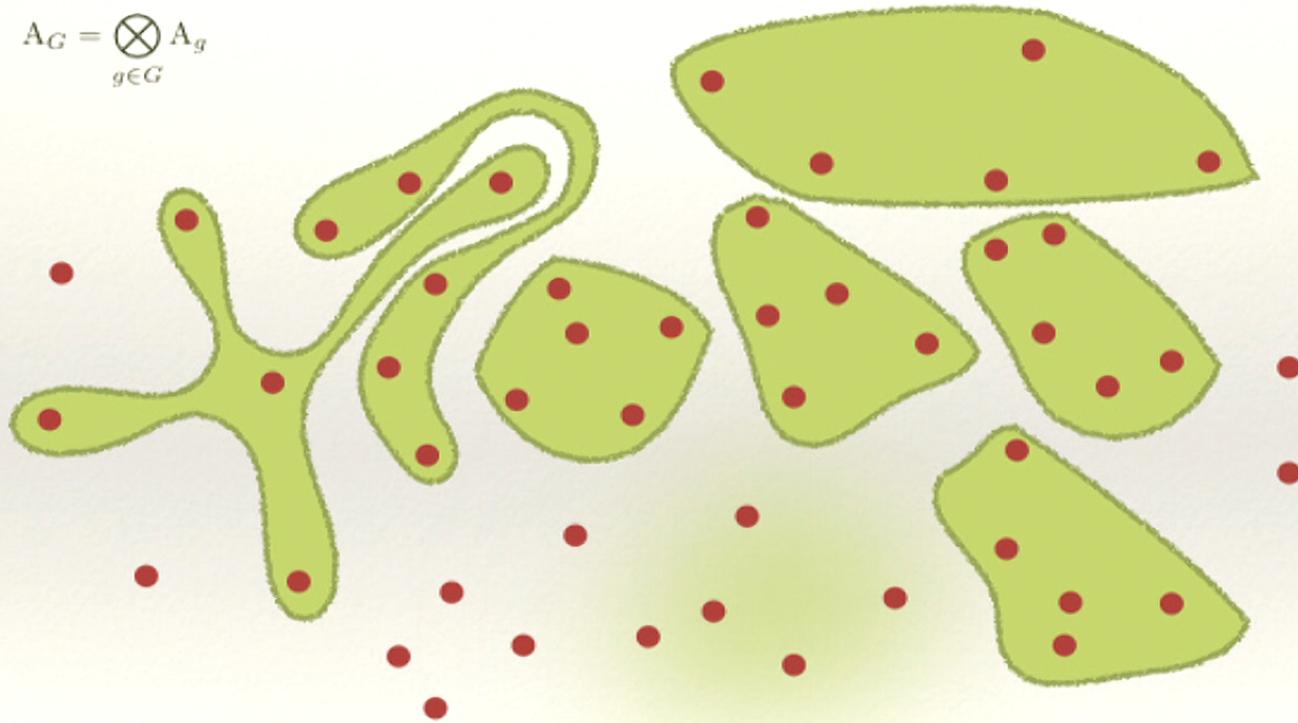


Remarks

- ❖ Why starting from effects?
 - ❖ Causality: unique deterministic effect
 - ❖ State space: functionals on the space of effects
 - ❖ Much “larger” than the space of quasi-local states
- ❖ States: functionals that locally behave as states

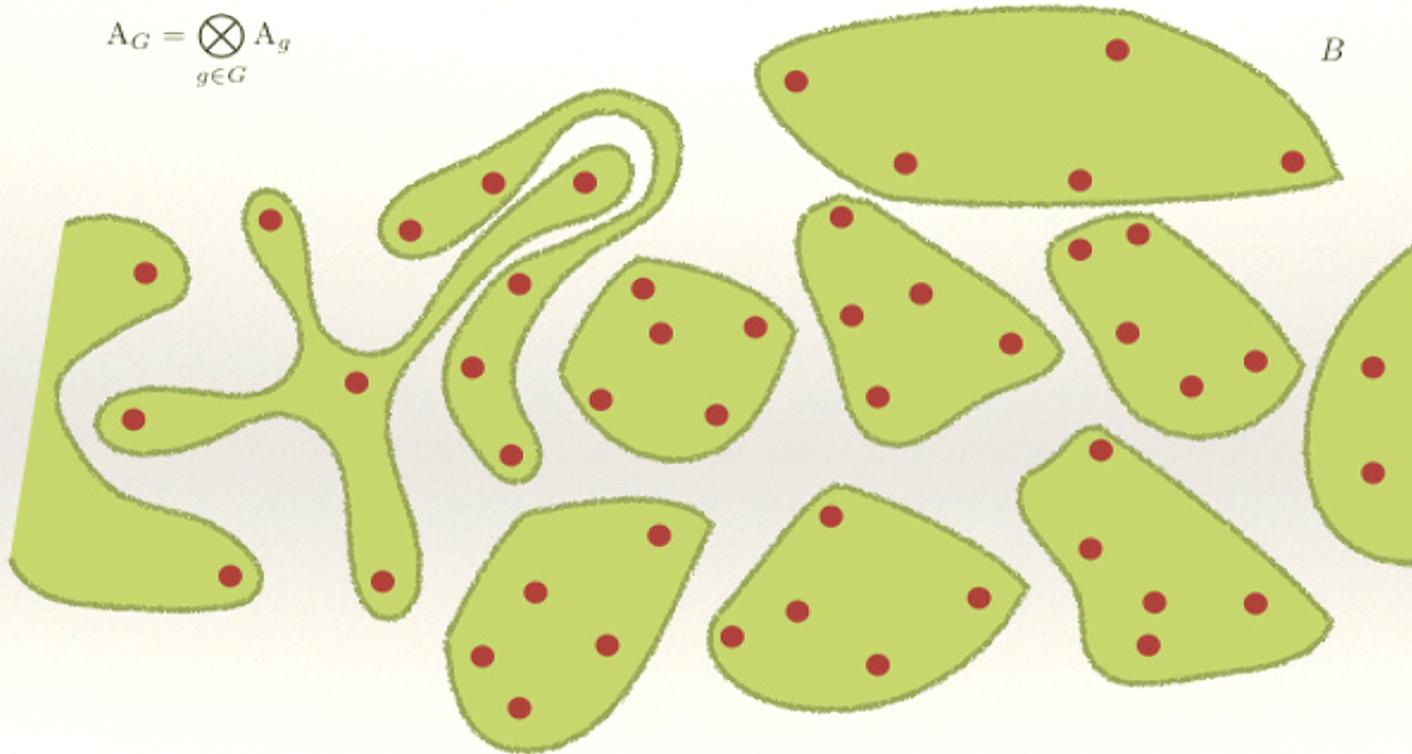
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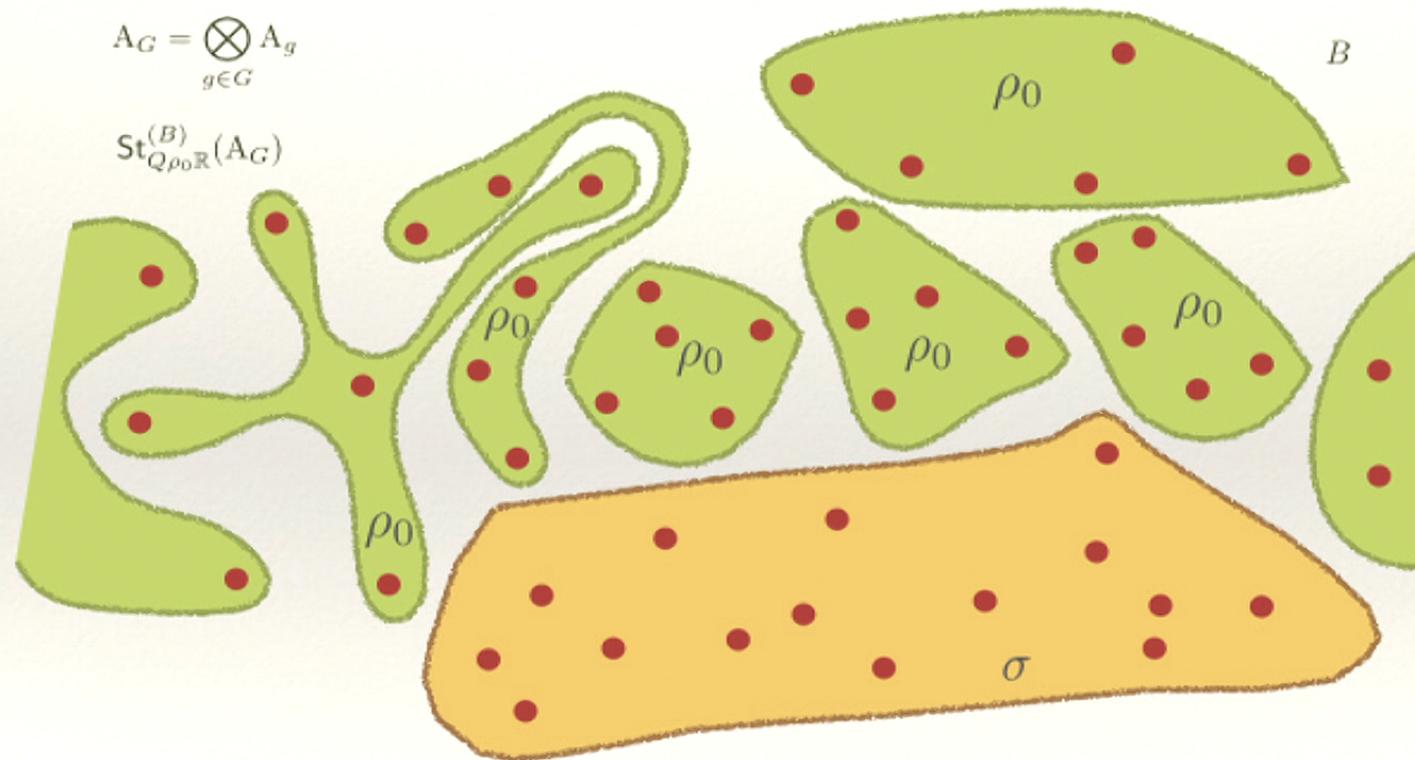


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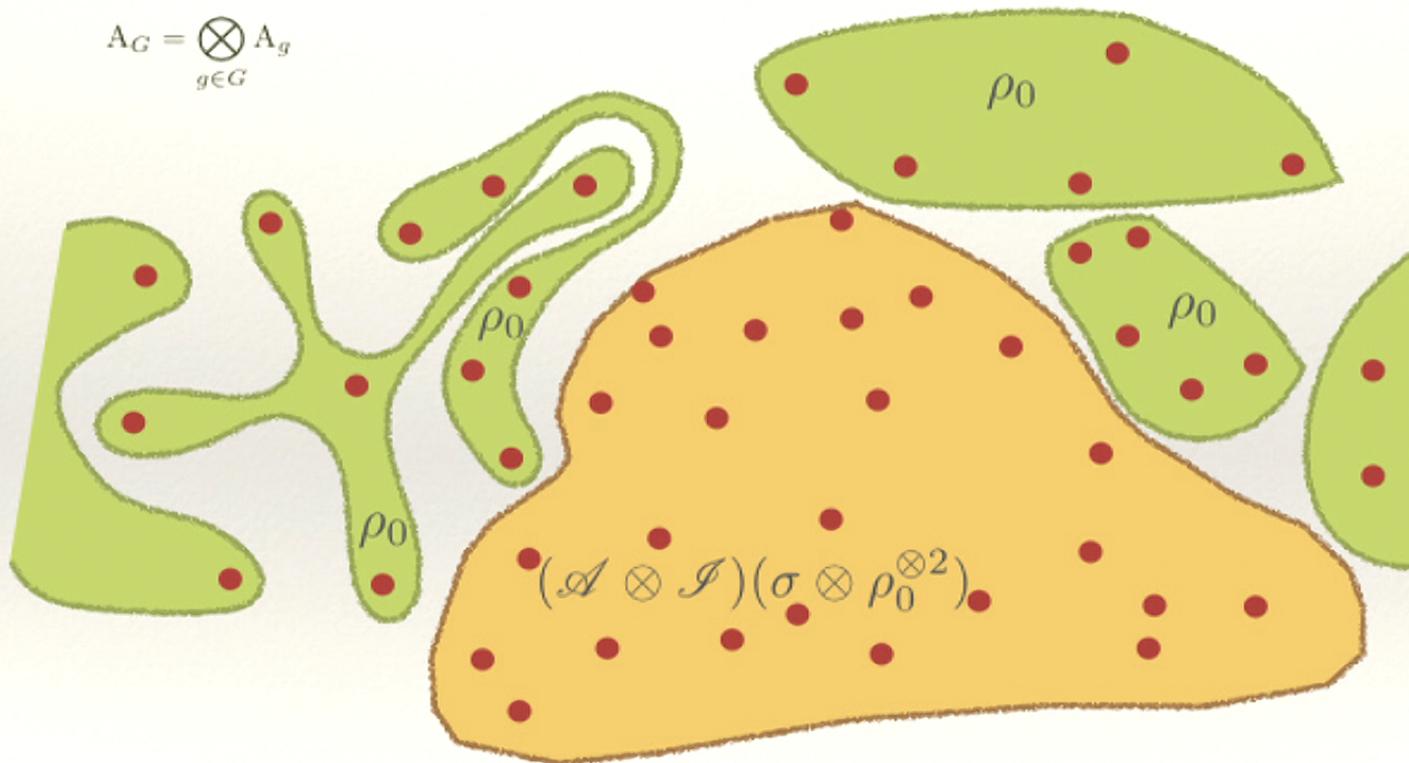


(Quasi-)Local states



(Quasi-)Local transformations

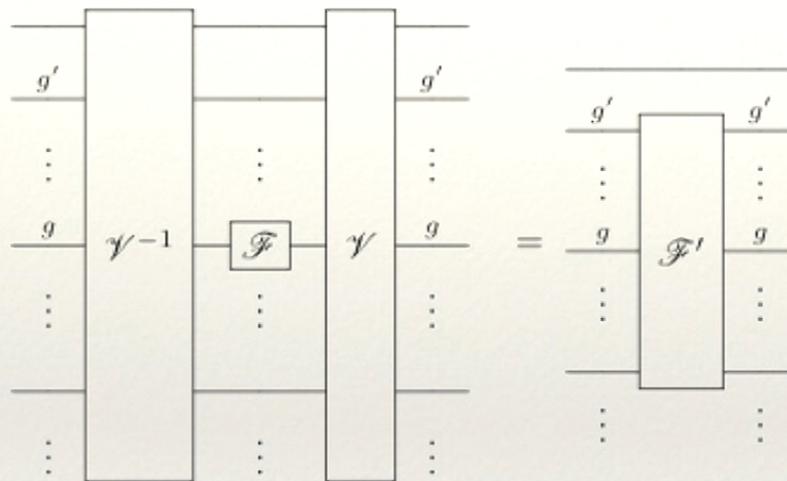
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Cellular Automata

- ❖ Maps $\mathcal{V} : \text{Eff}_{Q\mathbb{R}}(A_G) \rightarrow \text{Eff}_{Q\mathbb{R}}(A_G)$
- ❖ Admissible $(\mathcal{V} \otimes \mathcal{I}_C)\text{Eff}_Q(A_G C)^* \subseteq \text{Eff}_Q(A_G C)^*$
- ❖ Reversible $\mathcal{V}\mathcal{V}^{-1} = \mathcal{V}^{-1}\mathcal{V} = \mathcal{I}$
with \mathcal{V}^{-1} admissible
- ❖ The map $\mathcal{V} \cdot \mathcal{V}^{-1}$ preserves local transformations

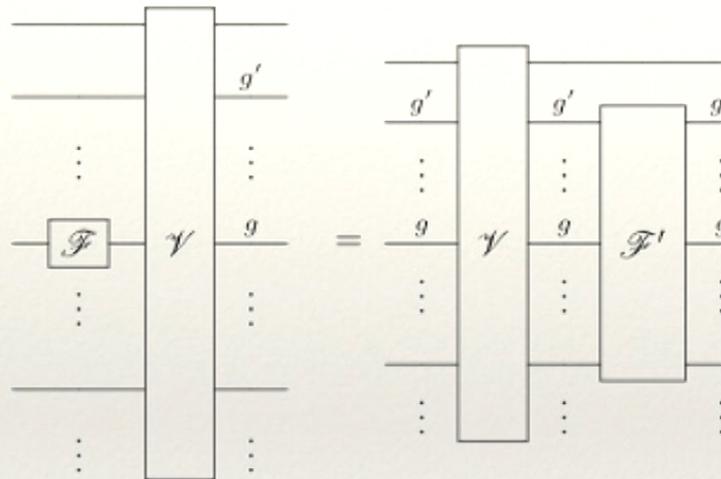
Causal influence



Then $g \rightsquigarrow_g g'$

g directly influences g' through \mathcal{V}

Causal influence



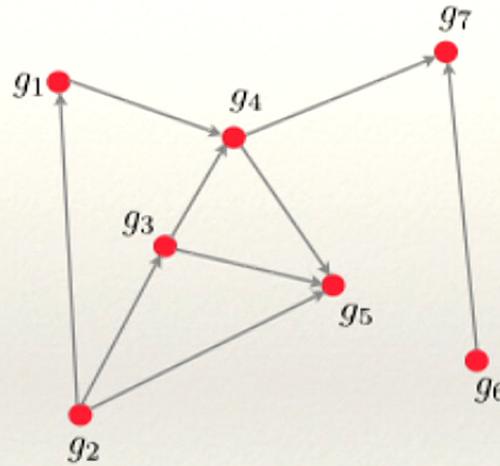
$$\mathcal{V} \mathcal{F} \mathcal{V}^{-1} = \mathcal{F}'$$

$$\mathcal{V} \mathcal{F} = \mathcal{F}' \mathcal{V}$$

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g directly influences g' through \mathcal{V}

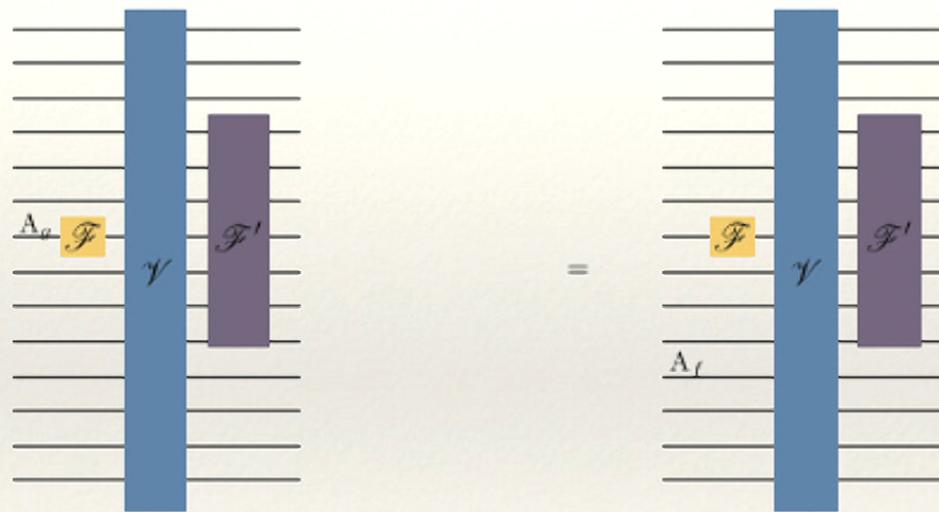
Graph of influence relations



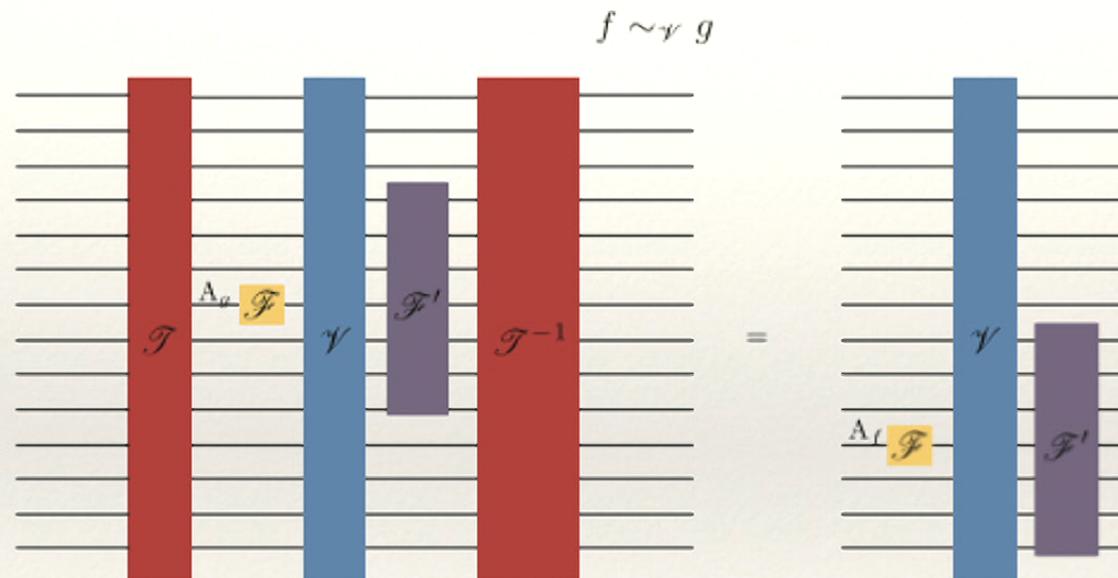
$g_1 \rightsquigarrow g_4$ $g_2 \rightsquigarrow g_1$ $g_3 \rightsquigarrow g_4$ $g_4 \rightsquigarrow g_5$ $g_6 \rightsquigarrow g_7$
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 $g_2 \rightsquigarrow g_5$

Homologous regions

$$f \sim_{\gamma} g$$



Homologous regions

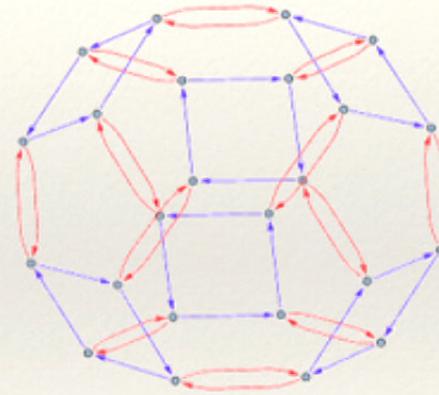
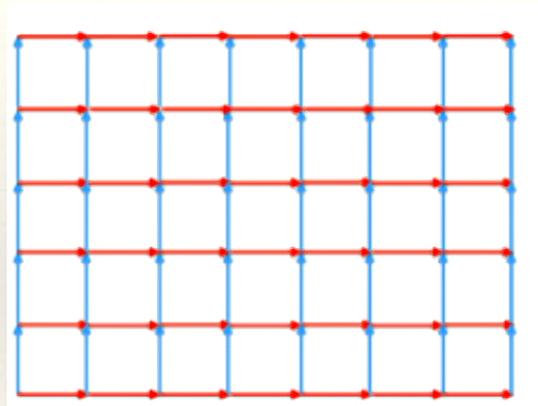


Homogeneity

“The evolution of every two different systems can be distinguished only with reference to a third system”

- ❖ For every f, g , f and g are homologous
- ❖ Two pairs (e, f) and (e, g) are never homologous

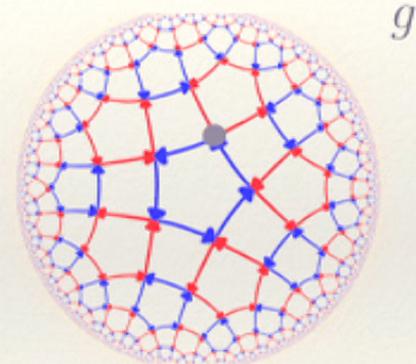
Homogeneity



The influence graph is a Cayley graph: it corresponds to a group

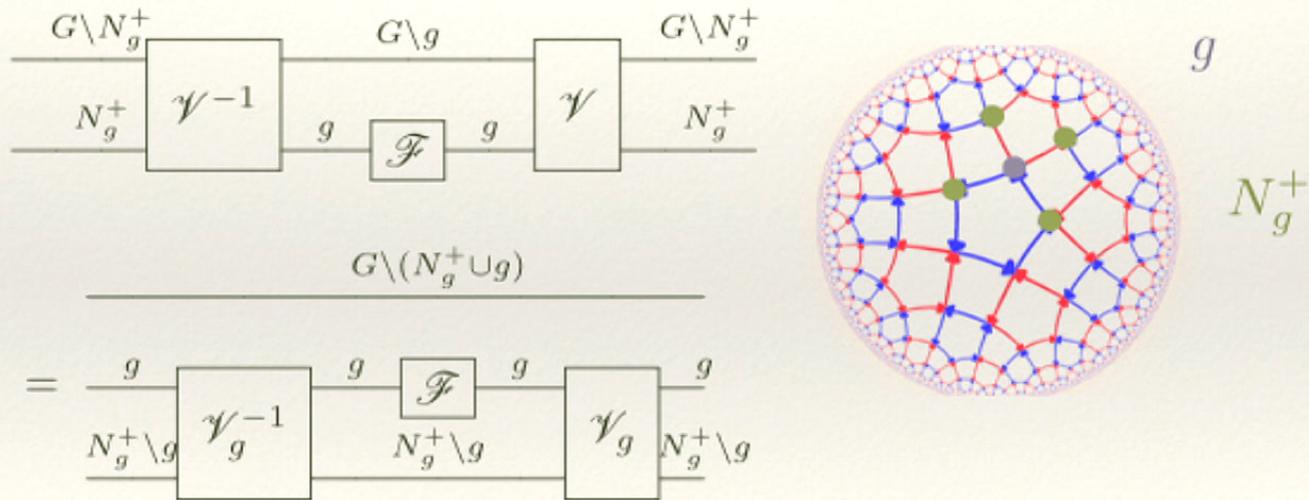
Locality

“The global evolution rule can be reduced to the evolution of bounded regions”



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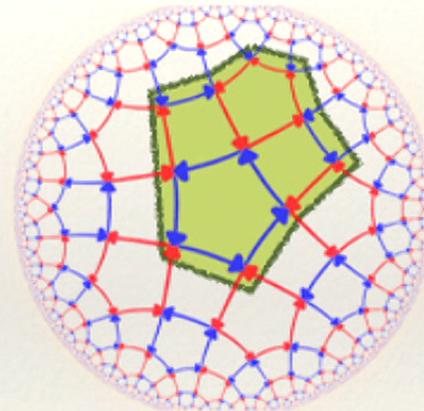


Locality

“The evolution rule can be reduced to the evolution of bounded regions”

i. Every vertex has finitely many arrows

ii. The elementary closed paths are bounded



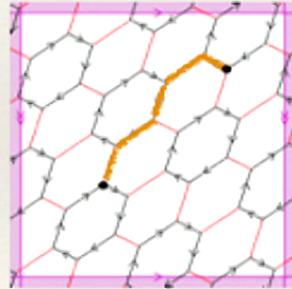
The signalling graph is the Cayley graph of a finitely presented group

Emergent space

Homogeneity and locality of evolution



Essentially unique geometry



$$a + bd(\mathbf{E}(g), \mathbf{E}(g'))_R \leq d(g, g')_G \leq a + \frac{1}{b}d(\mathbf{E}(g), \mathbf{E}(g'))_R$$

Example: Fermionic QCA

- ❖ OPT: Fermionic theory
- ❖ Automata: can be expressed as

$$\mathcal{T}(\varphi_i) = \varphi'_i$$

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- ❖ **Wrapping lemma**: the local rule for \mathcal{T} defines automata also on finite groups



B. Schumacher and R. F. Werner, arXiv:quant-ph/0405174 (2004)
PP and L. Poggiali, arXiv:1807.08695

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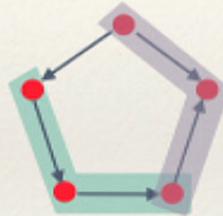
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- ❖ Automata: can be expressed as $\mathcal{T}(\varphi_i) = \varphi'_i$
- ❖ **Wrapping lemma:** the local rule for \mathcal{T} defines automata also on finite groups

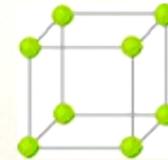
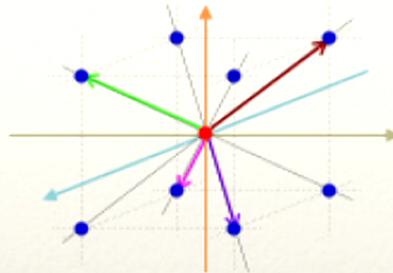


B. Schumacher and R. F. Werner, arXiv:quant-ph/0405174 (2004)
PP and L. Poggiali, arXiv:1807.08695

Linear Fermionic QCA

$$G = \mathbb{Z}^3$$

$$s = 2$$

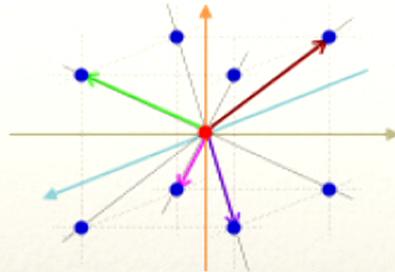


G. M. D'Ariano and PP, Phys. Rev. A **90**, 062106 (2014)

Linear Fermionic QCA

$$G = \mathbb{Z}^3$$

$$s = 2$$



$$\psi_{\mathbf{k},t+1} = W_{\mathbf{k}}^{\pm} \psi_{\mathbf{k},t}$$

G. M. D'Ariano and PP, Phys. Rev. A **90**, 062106 (2014)

Free Quantum Field theory

❖ Composing Weyl automata

$$Z_{\mathbf{k}}^{\pm} = \begin{pmatrix} nW_{\mathbf{k}}^{\pm} & imI \\ imI & nW_{\mathbf{k}}^{\pm\dagger} \end{pmatrix}$$

❖ Dirac's equation

$$i\partial_t\psi_{\mathbf{k},t} = (s\boldsymbol{\alpha} \cdot \mathbf{k} + c\beta)\psi_{\mathbf{k},t}$$

❖ Special class of entangled states

$$M_{\mathbf{k}}^{\pm} = W_{\mathbf{k}}^{\pm} \otimes W_{\mathbf{k}}^{\pm*}$$

❖ Maxwell's equations

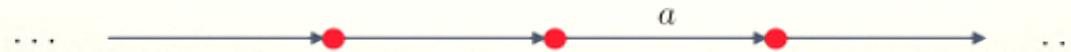
$$\partial_t \operatorname{Re} \mathbf{H}(\mathbf{x}, t) = \nabla \times \operatorname{Im} \mathbf{H}(\mathbf{x}, t),$$

$$\partial_t \operatorname{Im} \mathbf{H}(\mathbf{x}, t) = -\nabla \times \operatorname{Re} \mathbf{H}(\mathbf{x}, t),$$

G. M. D'Ariano and PP, Phys. Rev. A **90**, 062106 (2014),
A. Bisio, G. M. D'Ariano and PP, Ann. Phys. **368**, 177 (2016).

Non-linear FQCA

- $G = \mathbb{Z}$



- $A = A_D V(\chi)$

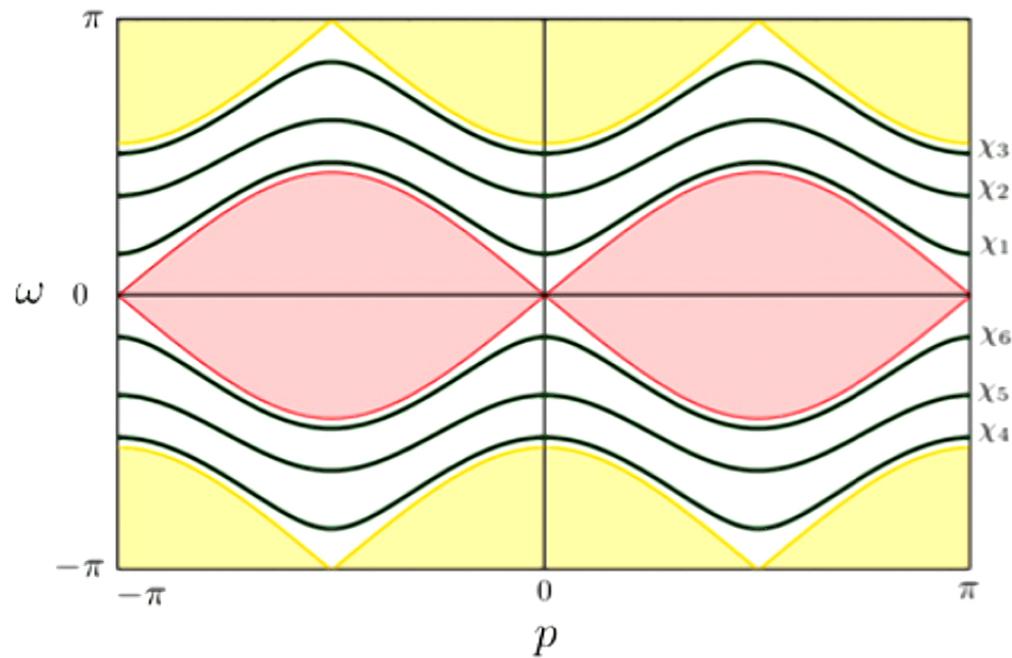
$$A_D \text{ Dirac automaton: } A_D \psi_j(x) A_D^\dagger = W_{Djk}(x, y) \psi_k(y)$$

$$V(\chi) := \exp[i\chi \psi_r^\dagger(x) \psi_r(x) \psi_l^\dagger(x) \psi_l(x)]$$



A. Bisio, G. M. D'Ariano, PP, and A. Tosini, Phys. Rev. A, **97**, 032132, (2018)

Non-linear FQCA

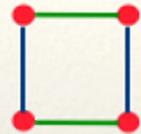


A. Bisio, G. M. D'Ariano, PP, and A. Tosini, Phys. Rev. A, **97**, 032132, (2018)

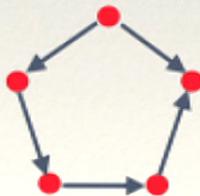
Non-linear FQCA: Case studies

$$\mathcal{T}(\varphi_i) = \varphi'_i$$

$$\mathbb{Z}_2 \times \mathbb{Z}_2$$



$$\mathbb{Z}_5 \quad (\mathbb{Z})$$



$\alpha_x \neq 0$ $\forall x = a, b, c$	$\alpha_x \neq 0$ with $x = a, b, c$ $\alpha_x, \beta_{cx}, \beta_{jx}, \gamma_{jx} \neq 0$	$\alpha_x \neq 0$ and $\alpha_y \neq 0$ with $x = a, b, c$ $\alpha_x, \alpha_y, \beta_{yx}, \beta_{xy} \neq 0$
No admitted solutions	$\alpha_x \alpha_x^* = 1; \gamma_{jx} = \beta_{cx} \beta_{jx}$	$\alpha_x \alpha_x^* + \alpha_y \alpha_y^* = 1$
	$-2 \leq \text{Re}(\beta_{cx}) \leq 0$	
	$\text{Im}(\beta_{ix}) = -\sqrt{-\text{Re}(\beta_{ix})(\text{Re}(\beta_{ix}) + 2)}$	$2 \alpha_x \cos(\theta_i - \theta_{j1}) = - \beta_{j1} $
	$-2 \leq \text{Re}(\beta_{jx}) \leq 0$	$\theta_x - \theta_{yx} = \theta_y - \theta_{xy}$
	$\text{Im}(\beta_{jx}) = -\sqrt{-\text{Re}(\beta_{jx})(\text{Re}(\beta_{jx}) + 2)}$	

$\mu_{410} = \nu_{410} = 0$	$\mu_{410} = 0; \nu_{410} \neq 0$ $\gamma_0, \eta_{10}, \eta_{40} \neq 0$	$\nu_{410} = 0; \mu_{410} \neq 0$ $\alpha_0, \beta_{10}, \beta_{40} \neq 0$
Linear case	$\eta_{10} = \eta_{40}$ $\gamma_0 \gamma_0^* = 1; \nu_{410} \gamma_0 = \eta^2$	$\beta_{10} = \beta_{40}$ $\alpha_0 \alpha_0^* = 1; \mu_{410} \alpha_0 = \beta^2$
	$2\cos(\theta_{\eta_{10}} - \theta_{\alpha_0}) = - \eta_{j1} $	$2\cos(\theta_{\beta_{10}} - \theta_{\alpha_0}) = - \beta_{j1} $

PP and L. Poggiali, arXiv:1807.08695



L. Poggiali



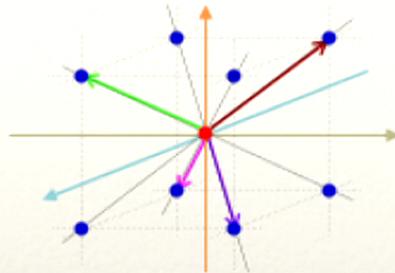
Summary

- ❖ Paradigm of a physical law as an algorithm
- ❖ Emergence of space manifolds from Cayley graphs
 - ❖ Universal feature of automata in OPTs
- ❖ Fermionic cellular automata
 - ❖ Linear case
 - ❖ Finite groups and \mathbb{Z}

Linear Fermionic QCA

$$G = \mathbb{Z}^3$$

$$s = 2$$



$$\psi_{\mathbf{k},t+1} = W_{\mathbf{k}}^{\pm} \psi_{\mathbf{k},t}$$



$$i\partial_t \psi_{\mathbf{k},t} = \mathbf{k} \cdot \boldsymbol{\sigma}^{\pm} \psi_{\mathbf{k},t}$$

Weyl's equations



G. M. D'Ariano and PP, Phys. Rev. A **90**, 062106 (2014)