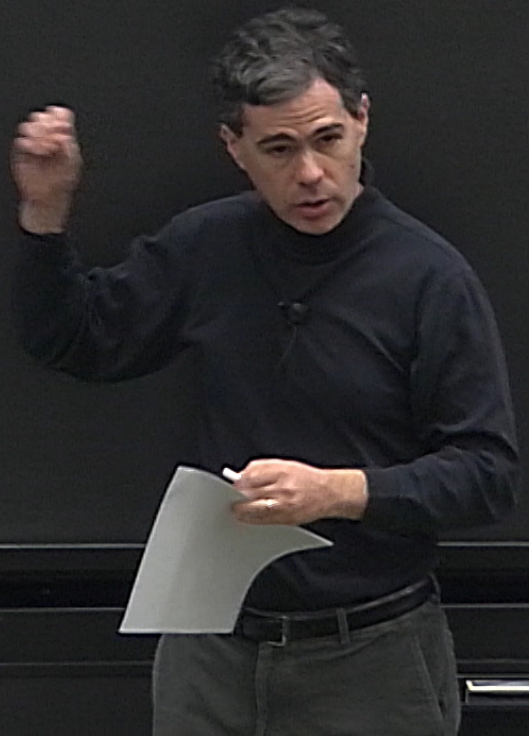
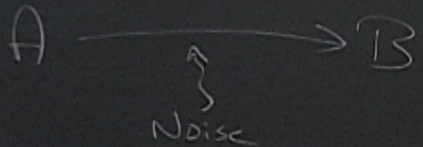


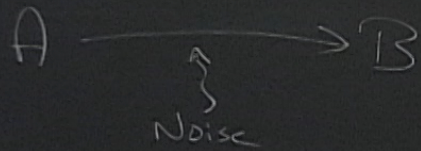
Title: PSI 2017/2018 - Quantum Information - Lecture 10

Date: Mar 05, 2018 11:30 AM

URL: <http://pirsa.org/18030023>

Abstract:





Repetition code:

0 → 000 ← 010
1 → 111

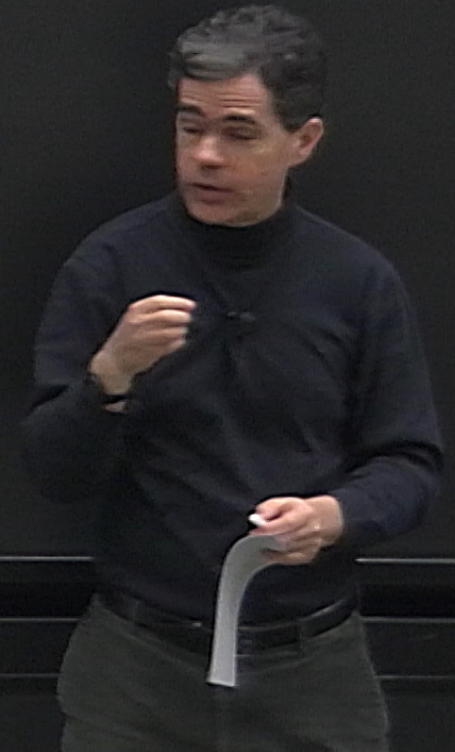
Prob. p of error per bit.

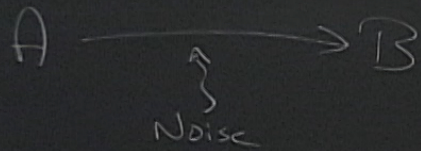
Prob. no errors $(1-p)^3$

Prob. 1 error $3p(1-p)^2$

Prob. 2 errors $3p^2(1-p)$

Prob. 3 errors p^3



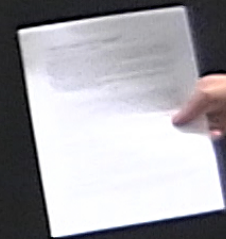


Repetition code:

$0 \rightarrow 000 \leftarrow 010$
 $1 \rightarrow 111$

Prob. p of error per bit.

Prob. no errors	$(1-p)^3$	}	$O(p^2)$
Prob. 1 error	$3p(1-p)^2$		
Prob. 2 errors	$3p^2(1-p)$		
Prob. 3 errors	p^3		

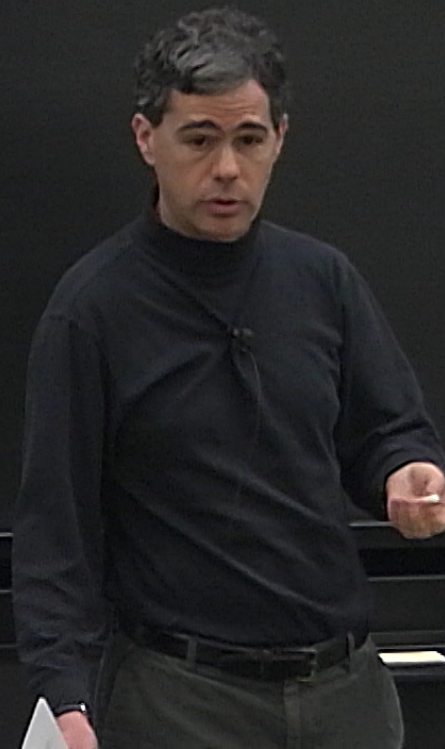


Problems w/ QECCs

1. No-cloning theorem
2. Measurement (e.g. to identify errors) can destroy superpositions.

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Problems w/ QECCs

1. No-cloning theorem
2. Measurement (e.g. to identify errors) can destroy superpositions.
3. Must deal with phase errors as well as bit flip errors.
4. Deal with continuous rotations, decoherence, ...

$$\begin{aligned} |0\rangle &\rightarrow |000\rangle \\ |1\rangle &\rightarrow |111\rangle \\ \alpha|0\rangle + \beta|1\rangle &\rightarrow \alpha|000\rangle + \beta|111\rangle \neq (\alpha|0\rangle + \beta|1\rangle)^{\otimes 3} \end{aligned}$$

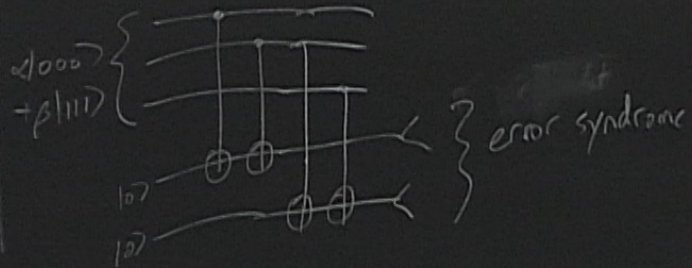
errors)

$$|0\rangle \rightarrow |000\rangle$$

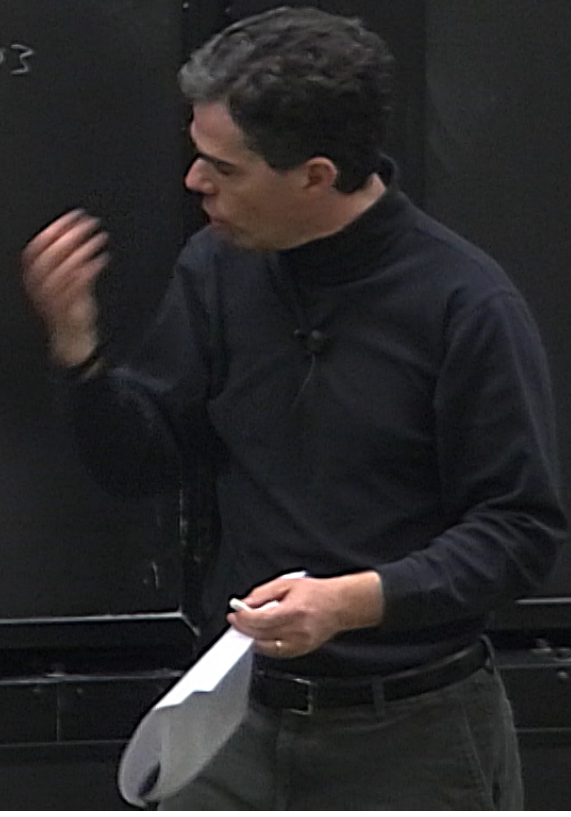
$$|1\rangle \rightarrow |111\rangle$$

$$\alpha|0\rangle + \beta|1\rangle \rightarrow \alpha|000\rangle + \beta|111\rangle \neq (\alpha|0\rangle + \beta|1\rangle)^{\otimes 3}$$

$$\text{Bit flip } X_2 \rightarrow \alpha|010\rangle + \beta|101\rangle$$



00	no error
10	1st qubit
01	3rd qubit
11	2nd qubit



$$\begin{aligned}
 |0\rangle &\stackrel{H}{\longleftrightarrow} |+\rangle = |0\rangle + |1\rangle \\
 |1\rangle &\longleftrightarrow |-\rangle = |0\rangle - |1\rangle
 \end{aligned}$$

Z in $|+\rangle, |-\rangle$ basis acts
like X .

$$\begin{aligned}
 |0\rangle &\rightarrow |+++ \rangle \\
 |1\rangle &\rightarrow |-- \rangle \\
 \alpha|0\rangle + \beta|1\rangle &\rightarrow \alpha|+++ \rangle + \beta|-- \rangle
 \end{aligned}$$

9-qubit code:

$$|0\rangle \rightarrow (|000\rangle + |111\rangle) \otimes (|000\rangle + |111\rangle) \otimes (|000\rangle + |111\rangle)$$

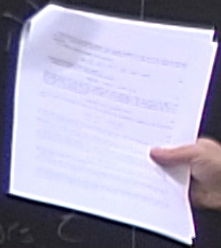
$$|1\rangle \rightarrow (|000\rangle - |111\rangle) \otimes (|000\rangle - |111\rangle) \otimes (|000\rangle - |111\rangle)$$

Corrects bit flip^X errors w/in each set of 3

Compare phases between sets of 3 to correct phase errors Z

Also correct $Y = -iXZ = \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix}$.

$$R_{2\theta} = \begin{pmatrix} 1 & 0 \\ 0 & e^{i2\theta} \end{pmatrix} = e^{i\theta} \begin{pmatrix} e^{-i\theta} & 0 \\ 0 & e^{i\theta} \end{pmatrix} = e^{i\theta} (\cos \theta I - i \sin \theta Z)$$

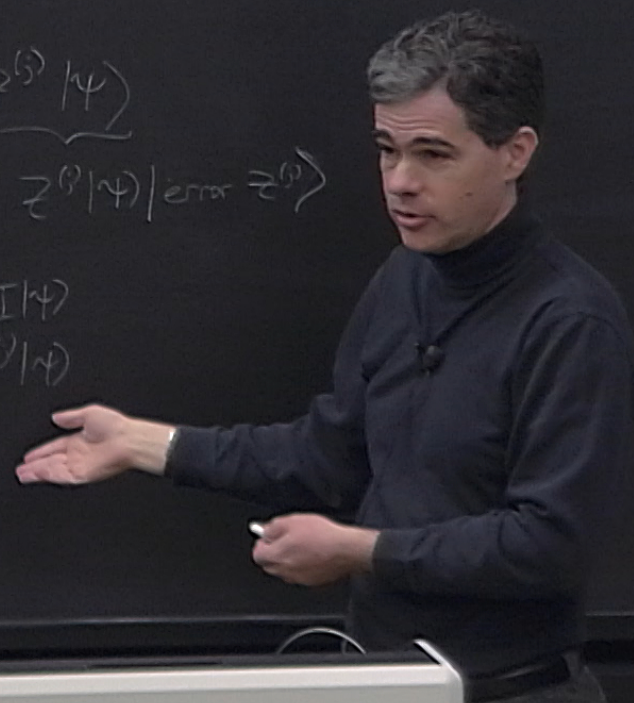


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$$R_{z\theta}^{(j)} |\psi\rangle = \cos\theta I |\psi\rangle - i\sin\theta Z^{(j)} |\psi\rangle$$

Add ancilla
 Unitaries $\rightarrow \cos\theta I |\psi\rangle | \text{no error} \rangle - i\sin\theta Z^{(j)} |\psi\rangle | \text{error } z^{(j)} \rangle$

Measure ancilla:
 "No error": Prob. $\cos^2\theta$, $I|\psi\rangle$
 "error $z^{(j)}$ ": Prob. $\sin^2\theta$, $Z^{(j)}|\psi\rangle$



Thm. If a QECC corrects errors A, B ,
then it also corrects $\alpha A + \beta B$. In particular,
if it corrects I, X, Y, Z on single ^(or +) qubits,
it corrects any operation on single ^(or +) qubits.

$C = 0101$
 $D = 1111$

$\sqrt{1}$
 $\sqrt{2}$
 $\sqrt{3}$
 4

Thm. If a QECC corrects errors A & B ,
 then it also corrects $\alpha A + \beta B$. In particular,
 if it corrects I, X, Y, Z on single qubits,
 it corrects any operation on single qubits.

$$\text{CPTP map } \rho \rightarrow \sum_k A_k \rho A_k^\dagger$$

$$|\psi\rangle\langle\psi| \rightarrow \sum_k A_k |\psi\rangle\langle\psi| A_k^\dagger$$

1.
 ✓ 1.
 ✓ 2.
 ✓ 3.
 4.

$$U_{\Sigma}^{(n)} = I + \varepsilon(U^{(1)} + U^{(2)} + \dots + U^{(n)}) + O(\varepsilon^2)$$

$$U_{\Sigma} = I + \varepsilon U'$$

Pauli group: on n qubits

$\mathcal{P}_n = \{ \text{tensor products of } I, X, Y, Z \text{ on } n \text{ qubits w/ phase } \pm 1, \pm i \}.$

- Any element squares to $+I$ or $-I$.
- Any two $P, Q \in \mathcal{P}_n$ either commute $[P, Q] = 0$
or anticommute $\{P, Q\} = 0$.

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Def. Weight of $P \in \mathcal{P}_n$ is # qubits where P is not I .

$z \otimes z$

$z \otimes z$

$z \otimes z$

$z \otimes z$

$z \otimes z$

$z \otimes z$

$X \otimes X \otimes X \otimes X \otimes X \otimes X$

$X \otimes X \otimes X \otimes X \otimes X \otimes X$

Any product has +1 eigenvalue for valid codewords.

$$S = \{ M \mid M \in \mathcal{P}_n, M|\psi\rangle = |\psi\rangle \forall |\psi\rangle \in \text{code} \}$$

Has following properties:

- Does not contain $-I$.
 - Is a group: $M, N \in S \Rightarrow MN|\psi\rangle = M|\psi\rangle = |\psi\rangle \checkmark$
 - Is an Abelian group: $M, N \in S \Rightarrow [M, N]|\psi\rangle = (MN - NM)|\psi\rangle = 0 \checkmark$
- Paulis commute or anticommute $\Rightarrow [M, N] = 0$.

code words.