

Title: Conditional entanglement of purification

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Abstract: <p>We study the conjectured holographic duality between entanglement of purification and the entanglement wedge cross-section. We generalize both quantities and prove several information theoretic inequalities involving them. These include upper bounds on conditional mutual information and tripartite information, as well as a lower bound for tripartite information. These inequalities are proven both holographically and for general quantum states. In addition, we use the cyclic entropy inequalities to derive a new holographic inequality for the entanglement wedge cross-section, and provide numerical evidence that the corresponding inequality for the entanglement of purification may be true in general. Finally, we use intuition from bit threads to extend the conjecture to holographic duals of suboptimal purifications.</p>

## Conditional Entanglement of Purification

1710.07643, 1701.07424, 1708.09393

Entanglement: the quantum or classical correlation between two subsystems.

For a pure state, the entanglement entropy quantifies the amt of ent. between a subsystem of the pure state and the rest of the pure state

$$\rho_A, \rho_A = \text{tr}_B |\Psi\rangle\langle\Psi|_{AB}, \quad S(\rho_A) = -\text{tr} \rho_A \log \rho_A$$

$$\rho_A, \rho_A = \text{tr}_B \rho_{AB}, \quad I(A:B) = S(\rho_A) + S(\rho_B) - S(\rho_{AB})$$

Ent. entropy is known to be inv. under LOCC.

For all quantum states

$$S(\rho) \geq 0$$

$$S(\rho_A) + S(\rho_B) - S(\rho_{AB}) \geq 0$$

$$S(\rho_{AB}) + S(\rho_{BC}) - S(\rho_B) - S(\rho_{ABC}) \geq 0$$

Purification



## Purification

$\rho_{AB} \rightarrow |\Psi\rangle_{ABC}$ , as long as  $\dim \mathcal{H}_C \geq 2^{S(\rho_{AB})}$

$S(A)$  is independent of your choice of purification!

What is the "minimal" purification of  $\rho_{AB}$ ?

Ent. of purification (Terhal et al. 2002)

$$E_p(A:B) = \inf_{A', B'} S(AA'), \text{ when } \rho_{AB} = \text{tr}_{AB'} |\Psi\rangle_{AA'BB'}, \quad \mathcal{H}_{AA'} \otimes \mathcal{H}_{BB'}$$



## observations

For pure  $\mathcal{P}_{AB}$ ,  $E_p(A:B) = S(A) = S(B)$ .

For product  $\mathcal{P}_{AB} = \mathcal{P}_A \otimes \mathcal{P}_B$ ,  $E_p(A:B) = 0$ .

For  $\mathcal{P}_{AB} = \sum p_i |X_i\rangle_A \otimes |X_i\rangle_B$ ,  $E_p(A:B) = -\sum p_i \log p_i \equiv$  Shannon entropy

$$1. 0 \leq E_p(A:B) \leq \min(S(A), S(B))$$

$$2. E_p(A:BC) \geq E_p(A:B)$$

$$3. E_p(A:B) \geq \frac{I(A:B)}{2}$$

$$4. E_p(A:BC) \geq \frac{I(A:B)}{2} + \frac{I(A:C)}{2}$$

$$5. \text{For pure } \mathcal{P}_{ABC}, E_p(A:BC) \geq E_p(A:B) + E_p(A:C) \text{ "polygamous"}$$



3. For pure  $\rho_{ABC}$ ,  $I_P(A:B) = I_P(A|C)$

1. standard purification  $\Rightarrow S(AA') = S(A) + S(B)$ .

2. For pure  $\rho_{ABC}$   $E_P(A:B) = S(A) \geq E_P(A|B) \checkmark$

mixed " " constrained opt. on the LHS, less cons. opt. on the RHS.

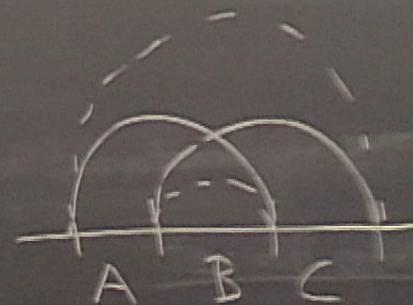
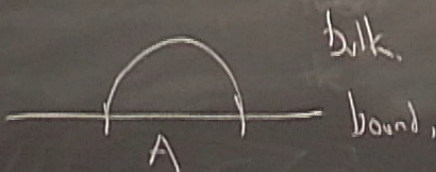
3.  $S(AA'|AB') \leq S(A'|A) + S(B'|B)$ .

$$\underbrace{S(AA'BB')}_{\uparrow 0} - S(AB) \leq \underbrace{S(AA') - S(A)}_{\uparrow E_P} + \underbrace{S(BB') - S(B)}_{\uparrow E_{P'}}$$



# Holography

$$S(A) = \frac{\text{Area}(A)}{4G}$$



Converts hard lin. alg. problems into geometric questions,

holog. states:

$$S(AB) + S(AC) + S(BC) \geq S(A) + S(B) + S(C) + S(ABC)$$

$$S(ABC) + S(BCD) + \dots + S(EAB) \geq S(AB) + S(BC) + \dots + S(EA) + S(AB-DE)$$

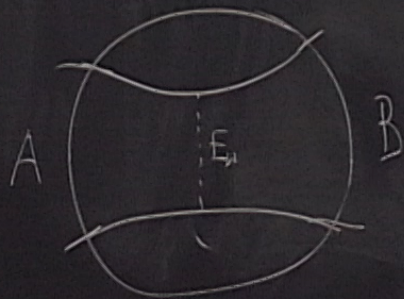




HPPS, constraints on theories, entropy const. constraints.

A recent conjecture by Tak. et al., Swingle, Zaldar, et al

$\exists$  A holog. description of  $E_P$ ,  $E_W$ , the entanglement wedge cross section.



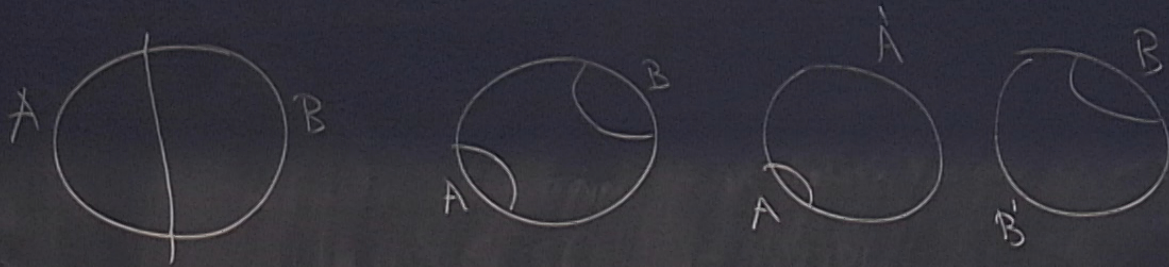
$E_W = \inf_{\text{partitions}}$  that separate A from B in the ent. wedge.

$E_W$  obeys all the same obs. and reqs that  $E_P$  is known to obey.

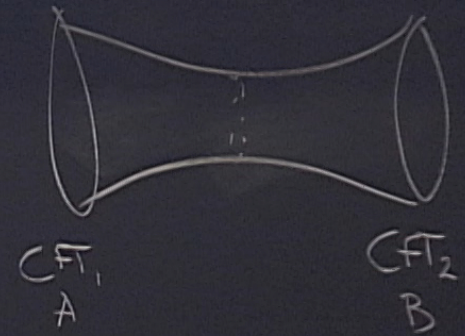


$$4. \mathbb{E}_P(A:BC) \geq \frac{I(A:B)}{2} + \frac{I(A:C)}{2}$$

5. For pure  $\rho_{ABC}$ ,  $\mathbb{E}_P(A:BC) \geq \mathbb{E}_P(A:B) + \mathbb{E}_P(A:C)$  "polygamous."



$$\mathbb{E}_W(A:BC) \geq \frac{1}{2} I(A:BC) \geq \frac{1}{2} I(A:B) + \frac{1}{2} I(A:C) \geq 0$$





$E_w(A:B|C)$  is the min partitioning of the Ent Wedge of ABC / Ent wedge of C. into regions adjacent to A or B.

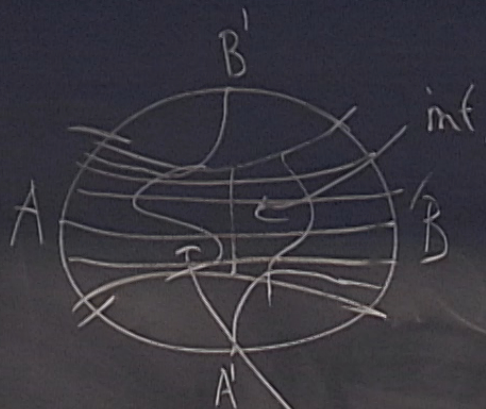
$$f_1(E_w(A:B|C)) \geq I(A:B|C) \geq 0$$

$$f_2(E_w(A:B|C)) \geq C_K(\cdot) \geq 0$$



$$\begin{aligned}
 E_p(A:B|C) &= \inf_{A'B'C^A} S(AA'C^A) \text{ where } |\psi\rangle_{A'B'B'C^A} \text{ is a purif of } \rho_{ABC} \text{ with } \left( \begin{array}{c} A \\ A' \end{array} \right) \left( \begin{array}{c} B \\ B' \end{array} \right) \left( \begin{array}{c} C \\ C^A \end{array} \right) \\
 &\geq \underline{E_p(A:B|C) + S(ABC) + S(C)} = S(AA'C^A) + S(BB'C^B) + S(ABC) + S(C) \text{ for some } C^A, C^B, A', B' \\
 &\geq S(AA'BC) + S(BB'C^B) + S(AC^A) + S(C^A, C^B) \\
 &\geq S(AA'BC) + S(BB'C^B) + S(AC) + S(C^A) \\
 &\geq S(BC^B) + S(AC) + S(C^A) \geq S(BC) + S(AC)
 \end{aligned}$$





$$E_p = \inf ( \frac{S(AA')}{\dots} )$$

$S(AA')$