

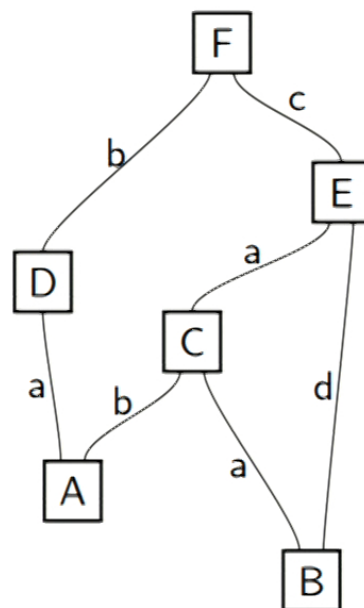
Title: PSI 17/18 - Foundations of Quantum Mechanics - Lecture 9

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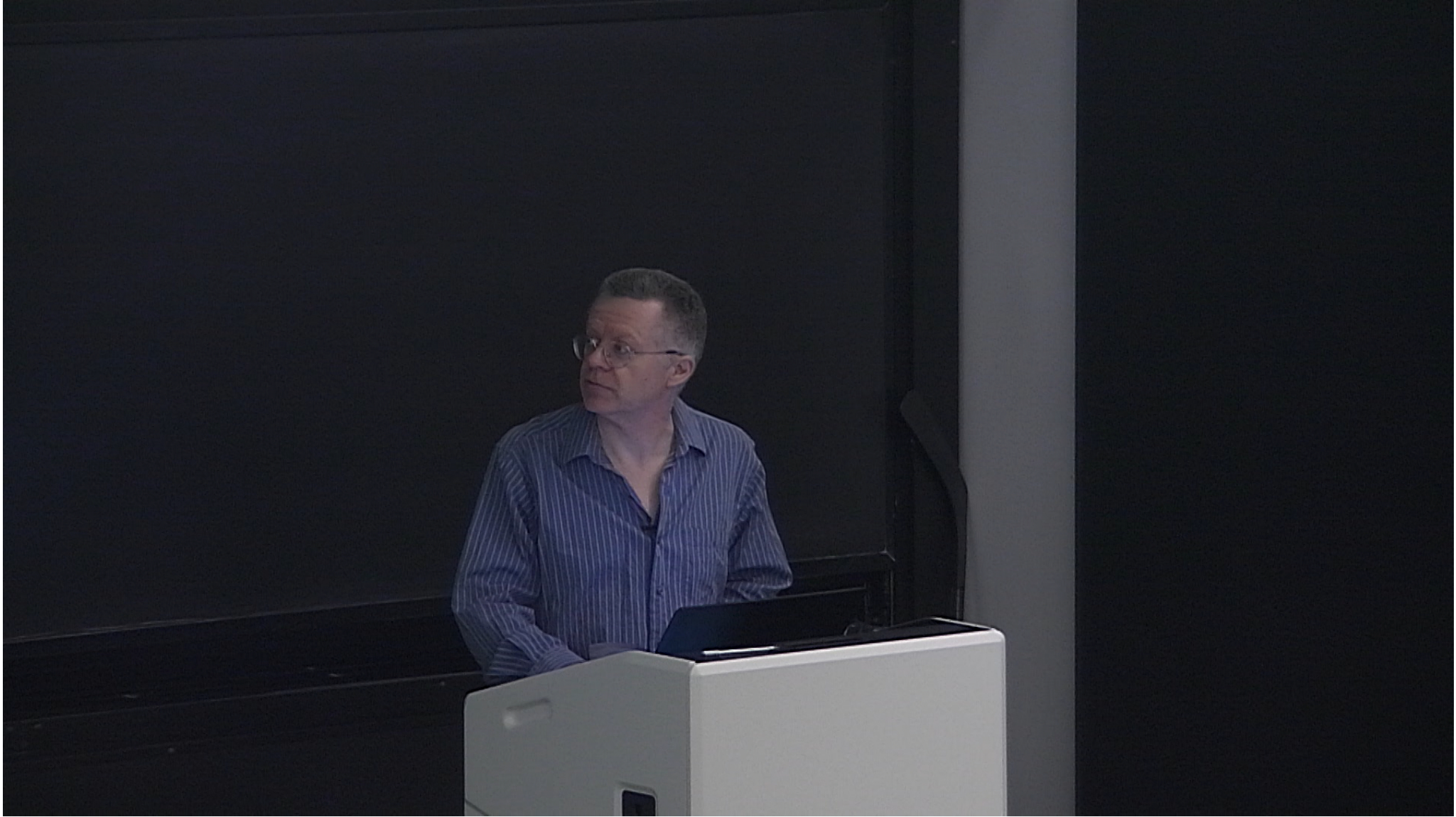
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Abstract:

Consider a circuit

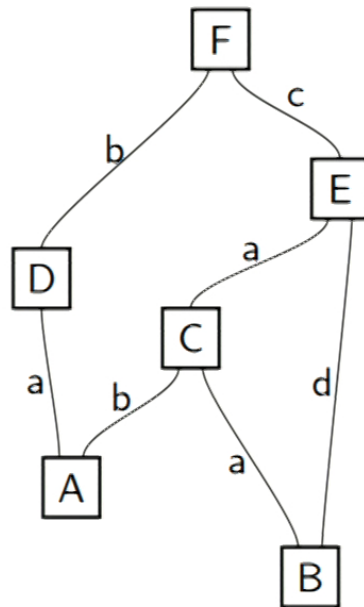


How do we calculate the probability for this circuit in standard framework of QT?





Consider again the circuit



$$A^{a_1 b_2} B^{a_3 d_4} C^{a_5}_{b_2 a_3} D^{b_6}_{a_1} E^{c_7}_{a_5 d_4} F_{b_6 c_7}$$

In the operator tensor formulation

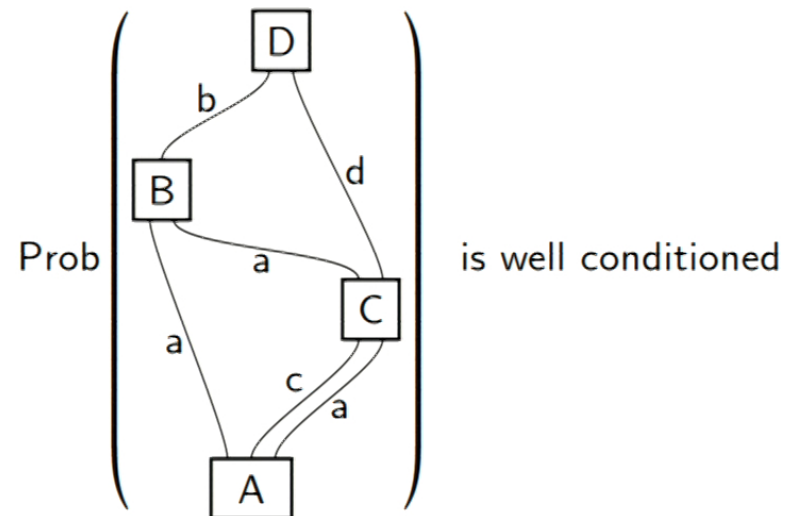
$$\text{Prob}(A^{a_1 b_2} B^{a_3 d_4} C^{a_5}_{b_2 a_3} D^{b_6}_{a_1} E^{c_7}_{a_5 d_4} F_{b_6 c_7}) = \hat{A}^{a_1 b_2} \hat{B}^{a_3 d_4} \hat{C}^{a_5}_{b_2 a_3} \hat{D}^{b_6}_{a_1} \hat{E}^{c_7}_{a_5 d_4} \hat{F}_{b_6 c_7}$$

We will explain what the RHS means later.



## Introducing probabilities

**Assump 1** *We can associate a probability with any given circuit (the probability that the circuit “happens”), and this probability depends only on the specification of the given circuit (the knob settings and outcome sets at the operations, and the wiring).*



## The $p(\cdot)$ function

We define the function  $p(\cdot)$  as follows

$$p(\alpha A + \beta B + \dots) := \alpha \text{Prob}(A) + \beta \text{Prob}(B) + \dots$$

for *circuits*  $A, B, \dots$  and real numbers  $\alpha, \beta, \dots$  (these can be negative).

Will use this to define a notion of equivalence.

## Example of equivalence

Have

$$\alpha A^{a_1} + \beta B^{a_1} \equiv \gamma C^{a_1} + \delta D^{a_1}$$

if

$$p([\alpha A^{a_1} + \beta B^{a_1}]E_{a_1}) = p([\gamma C^{a_1} + \delta D^{a_1}]E_{a_1}) \quad \text{for all } E_{a_1}$$



## General definition of equivalence

We consider expressions like

$$\text{expression} = \alpha + \beta A + \gamma B + \dots$$

where  $A, B, \dots$  are fragments.

**Equivalence:** *We write*

$$\text{expression}_1 \equiv \text{expression}_2$$

*if*

$$p(\text{expression}_1 E) \equiv p(\text{expression}_2 E)$$

*for any fragment  $E$  that makes the contents of the argument on both sides of this equation into a linear sum of circuits.*

Equivalence is a weaker notion than equality.

## Another example of equivalence

In general, we have

$$A \equiv \text{Prob}(A) \quad \text{for any circuit } A$$

Proof: For any circuit  $E$

$$p(AE) = p(A)p(E) = p(\text{Prob}(A)E)$$

This example illustrates how equivalence is a weaker notion than equality.

# Fiducial preparations

Fiducial preparations

$$\begin{array}{c} a \\ | \\ \bullet \\ | \\ \nabla \end{array} \iff a_1 X^{a_1} \text{ where } a_1 = 1 \text{ to } K_a$$

For any preparation  $A^{a_1}$  (summation over  $a_1$  implicit below)

$$A^{a_1} \equiv a_1 A \quad \iff \quad \begin{array}{c} a \\ | \\ \boxed{A} \end{array} \equiv \boxed{A} \text{---} \begin{array}{c} a \\ | \\ \bullet \\ | \\ \nabla \end{array}$$



# Fiducial results

Fiducial results

$$\begin{array}{c} \triangle \\ | \\ a \end{array} \bullet^a \iff X_{a_1}^{a_1} \text{ where } a_1 = 1 \text{ to } K_a$$

For any result for a system of type a

$$B_{a_1} \equiv B_{a_1} X_{a_1}^{a_1} \iff \begin{array}{c} \boxed{B} \\ | \\ a \end{array} \equiv \begin{array}{c} \triangle \\ | \\ a \end{array} \bullet^a \circ \boxed{B}$$

## The hopping metric

We define *the hopping metric*

$$\bullet \xrightarrow{a} \bullet := p \left( \begin{array}{c} \triangle \bullet a \\ a \\ a \bullet \nabla \end{array} \right) \Leftrightarrow \begin{array}{c} \triangle \bullet a \\ a \\ a \bullet \nabla \end{array} \equiv \bullet \xrightarrow{a} \bullet$$

## Black and white dots

We define

$$\boxed{A} \text{---} \bullet := \boxed{A} \text{---} \circ \text{---} \bullet \quad \bullet \text{---} \boxed{B} := \bullet \text{---} \bullet \text{---} \circ \text{---} \boxed{B}$$

Hence

$$\boxed{A} \text{---} \circ \text{---} \bullet \xrightarrow{a} \circ \text{---} \boxed{B} = \boxed{A} \text{---} \circ \text{---} \bullet \xrightarrow{a} \bullet \text{---} \boxed{B} = \boxed{A} \text{---} \bullet \xrightarrow{a} \circ \text{---} \boxed{B} := \boxed{A} \xrightarrow{a} \boxed{B}$$

We have

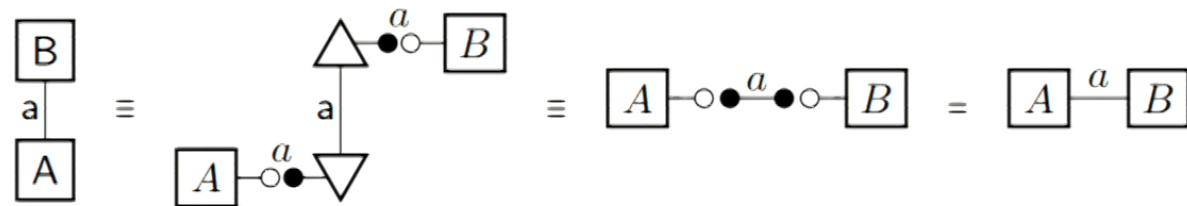
$$\text{---} \circ \bullet \text{---} = \text{---} = \text{---} \bullet \circ \text{---}$$

Hence, we can insert and delete pairs of black and white dots as we like.  
Consistency requires

- ▶  $\circ \text{---} \circ$  to be the inverse of  $\bullet \text{---} \bullet$
- ▶  $\circ \text{---} \bullet$  to be equal to the identity
- ▶  $\bullet \text{---} \circ$  to be equal to the identity



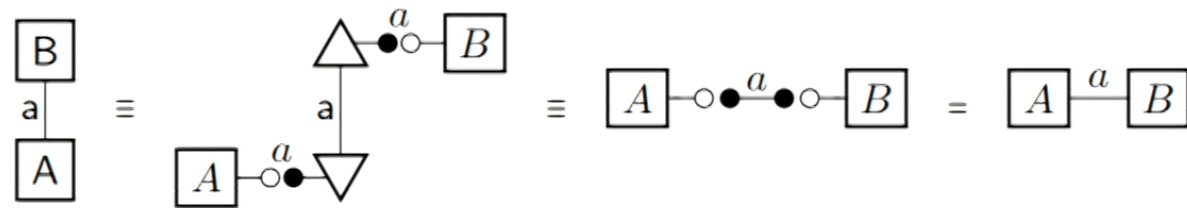
## The steps for a simple circuit



Hence

$$\text{Prob} \left( \begin{array}{c} B \\ a \\ A \end{array} \right) = A \text{---} a \text{---} B$$

## The steps for a simple circuit

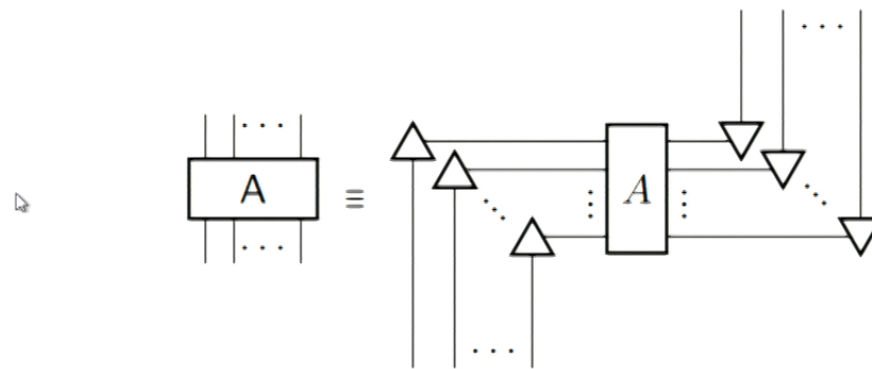


Hence

$$\text{Prob} \left( \begin{array}{c} \boxed{B} \\ \text{a} \\ \boxed{A} \end{array} \right) = \boxed{A} \text{---}^a \boxed{B}$$

## Assumption 2

**Assumption 2: Operations are fully decomposable.** We assume that any operation can be written as



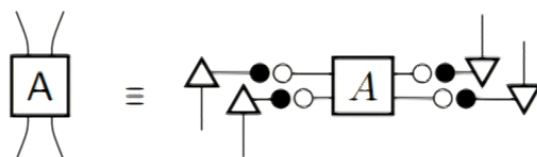
In words we will say that any operation is equivalent to a linear combination of operations each of which consists of an result for each input and a preparation for each output.

This is equivalent to tomographic locality.

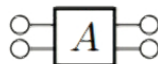


## Duotensor with all white dots

Inserting black and white dots (with black next to the fiducial elements)



Therefore



(with all white dots) provides the weights in the sum over fiducial elements.

This is an example of a duotensor.

## What are duotensors?

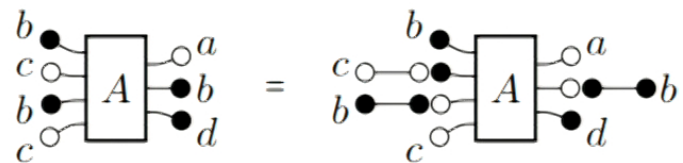
- ▶ Like tensors except that each index is associated with two bases.
- ▶ They transform like tensors but with respect to two bases.
- ▶



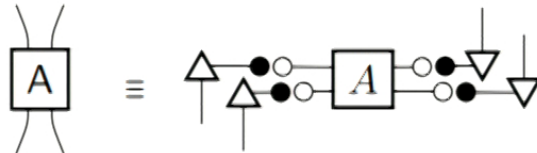
- ▶ Have map

$$\begin{array}{cc} \circ & \bullet \\ & A \\ \bullet & \circ \end{array}$$

- ▶ Can change colours of dots using ●—● and ○—○



All white dots gives coefficients in sum over fiducials

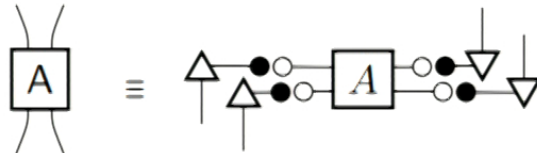


All black dots gives fiducial probabilities

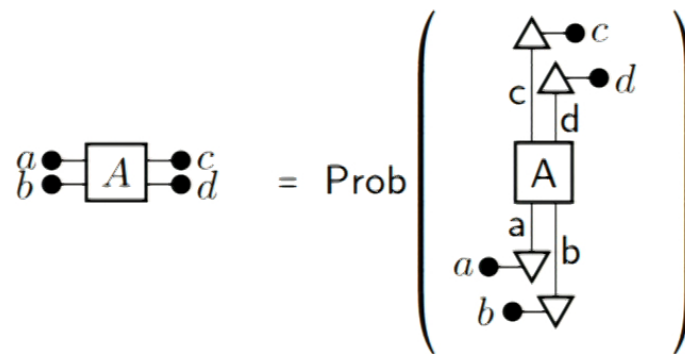
$$\begin{matrix} a & \bullet & \\ b & \bullet & \end{matrix} \begin{matrix} \bullet \\ \bullet \end{matrix} \begin{matrix} \text{---} \\ \text{---} \end{matrix} \boxed{A} \begin{matrix} \bullet & \bullet \\ \bullet & \bullet \end{matrix} \begin{matrix} c \\ d \end{matrix} = \text{Prob} \left( \begin{matrix} \triangle \bullet c \\ c \triangle \bullet d \\ \triangle \bullet d \\ \triangle \bullet c \\ \triangle \bullet c \\ \triangle \bullet d \\ \triangle \bullet d \\ \triangle \bullet c \end{matrix} \right)$$



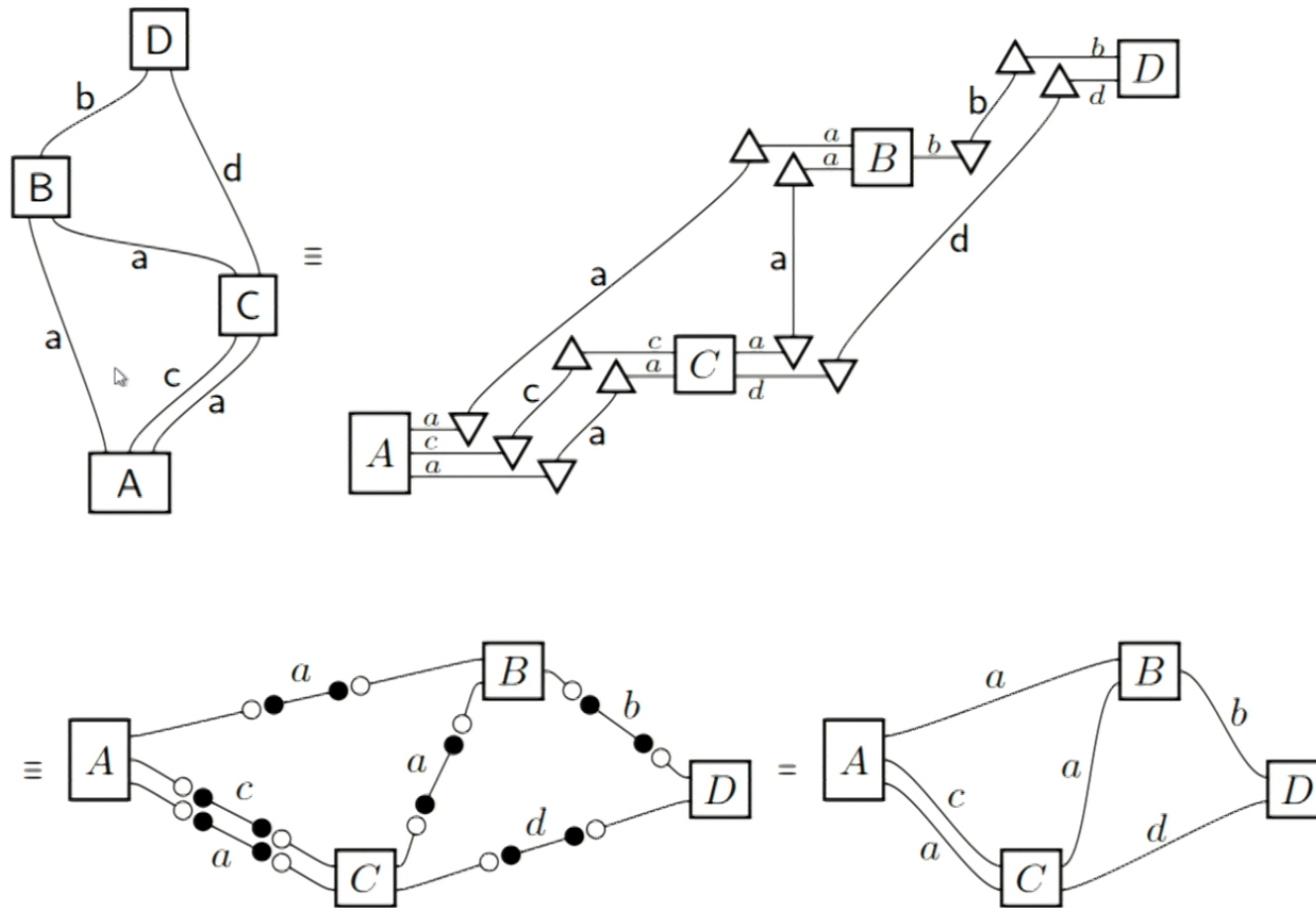
All white dots gives coefficients in sum over fiducials



All black dots gives fiducial probabilities



## General circuits



## Hilbert space notation

- ▶ We define  $\mathcal{H}_{a_1}, \mathcal{H}_{b_2}, \dots$  having dimensions  $N_a, N_b, \dots$
- ▶ We define  $\mathcal{H}^{a_1}, \mathcal{H}^{b_2}, \dots$  having dimensions  $N_a, N_b, \dots$
- ▶ We define

$$\mathcal{H}_{a_1 b_2 \dots c_3}^{d_4 e_5 \dots f_6} := \mathcal{H}_{a_1} \otimes \mathcal{H}_{b_2} \otimes \dots \otimes \mathcal{H}_{c_3} \otimes \mathcal{H}^{d_4} \otimes \mathcal{H}^{e_5} \otimes \dots \otimes \mathcal{H}^{f_6}$$

These are all taken to be complex Hilbert spaces.



## Space of operators

We define

$$\mathcal{V}_{a_1 b_2 \dots c_3}^{d_4 e_5 \dots f_6}$$

as the space of Hermitian operators acting on  $\mathcal{H}_{a_1 b_2 \dots c_3}^{d_4 e_5 \dots f_6}$ .

We write

$$\hat{A}_{a_1 b_2 \dots c_3}^{d_4 e_5 \dots f_6} \Leftrightarrow \begin{array}{c} \begin{array}{ccc} d & e & f \\ | & | & | \\ \dots & & \end{array} \\ \boxed{\hat{A}} \\ \begin{array}{ccc} | & | & | \\ a & b & c \end{array} \end{array}$$

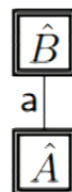
for elements of  $\mathcal{V}_{a_1 b_2 \dots c_3}^{d_4 e_5 \dots f_6}$ .

## Notation

- ▶ We write  $\hat{A}^{a_1} \hat{B}^{b_2}$  instead of  $\hat{A}^{a_1} \otimes \hat{B}^{b_2}$  (in  $\mathcal{V}^{a_1 b_2} = \mathcal{V}^{a_1} \otimes \mathcal{V}^{b_2}$ ).
- ▶ We write  $\hat{A}_{a_1} \hat{C}^{b_2}$  instead of  $\hat{A}_{a_1} \otimes \hat{C}^{b_2}$ .
- ▶ We write  $\hat{A}_{a_1} \hat{D}^{a_2}$  instead of  $\hat{A}_{a_1} \otimes \hat{D}^{a_2}$ .
- ▶ Order not important (information in the integers).  $\hat{A}_{a_1} \hat{C}^{b_2} = \hat{C}^{b_2} \hat{A}_{a_1}$

When we have a repeated integer:

- ▶ We write  $\hat{A}^{a_1} \hat{B}_{a_1}$  for  $\text{trace}(\hat{A}^{a_1} \hat{B}_{a_1})$ , denoted graphically by



A wire or a repeated integer means we take the trace of the product of the given operators.

- ▶ When have a more complicated example like

$$\hat{A}_{a_1 b_2}^{b_3 c_4} \hat{B}_{a_5 b_3}^{d_6 c_7}$$

we take product in appropriate ( $b_3$ ) subspace and then take *partial trace* in that subspace. This is accomplished very naturally using full decomposability of operators.



## Fiducial operators

We introduce a fiducial (spanning) set of operators for  $\mathcal{V}^a$

$${}_{a_1} \hat{X}^{a_1} \iff a \bullet \nabla^a \quad \text{where } a_1 = 1 \text{ to } K_a$$

Similarly, we introduce a fiducial (spanning) set of operators for the space  $\mathcal{V}_{a_1}$

$$\hat{X}_{a_1}^{a_1} \iff \triangle^a \bullet a \quad \text{where } a_1 = 1 \text{ to } K_a$$

## The hopping metric

In the context of operator tensors the hopping metric is given by

$$\bullet \overset{a}{\text{---}} \bullet := \begin{array}{c} \triangle \bullet a \\ | \\ a \\ | \\ \bullet \nabla \end{array}$$

and its inverse is represented by  $\circ \text{---} \circ$ .

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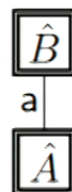
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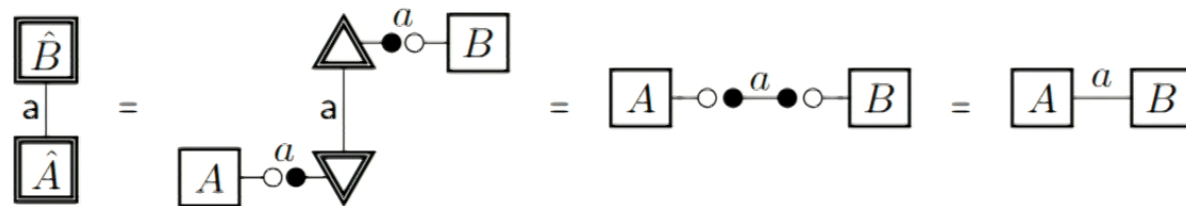
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## The steps for a simple operator circuit





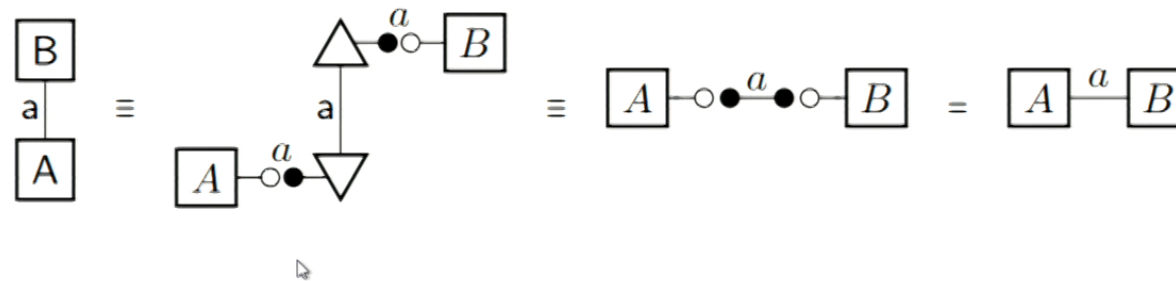
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and its inverse is represented by  $\circ \text{---} \circ$ .

## Compare with a simple (operation) circuit



Since Hilbert spaces are complex ....

$$\mathcal{V}_{a_1 b_2 \dots c_3}^{d_4 e_5 \dots f_6} = \mathcal{V}_{a_1} \otimes \mathcal{V}_{b_2} \otimes \dots \otimes \mathcal{V}_{c_3} \otimes \mathcal{V}^{d_4} \otimes \mathcal{V}^{e_5} \otimes \dots \otimes \mathcal{V}^{f_6}$$

we have ....



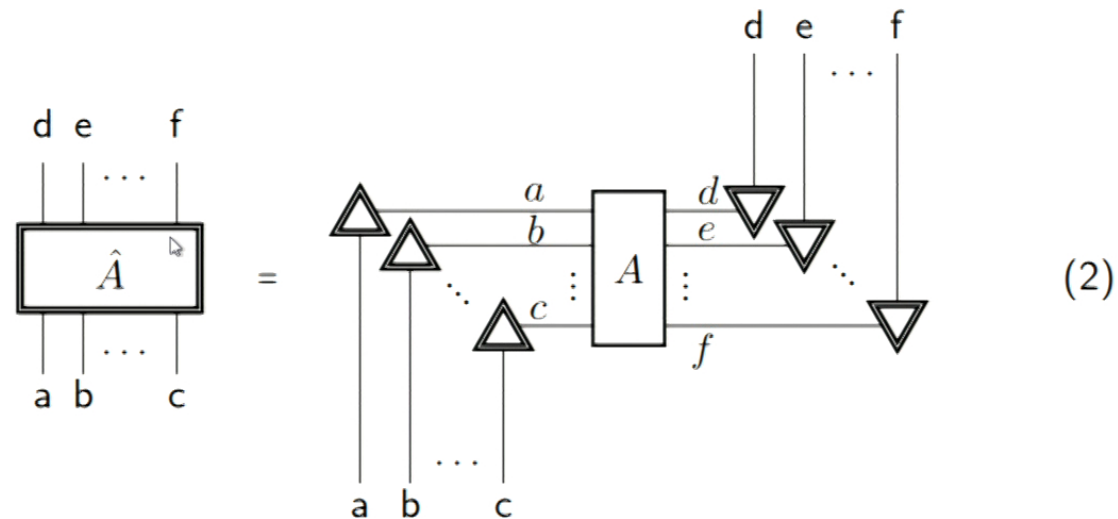


## Full decomposability of operators

We can write any operator as a linear sum over fiducial operators for the inputs and outputs.

$$\hat{A}_{a_1 b_2 \dots c_3}^{d_4 e_5 \dots f_6} = d_4 e_5 \dots f_6 A_{a_1 b_2 \dots c_3} \hat{X}_{a_1}^{a_1} \hat{X}_{b_2}^{b_2} \dots \hat{X}_{c_3}^{c_3} d_4 \hat{X}^{d_4} e_5 \hat{X}^{e_5} \dots f_6 \hat{X}^{f_6} \quad (1)$$

in symbolic notation, or

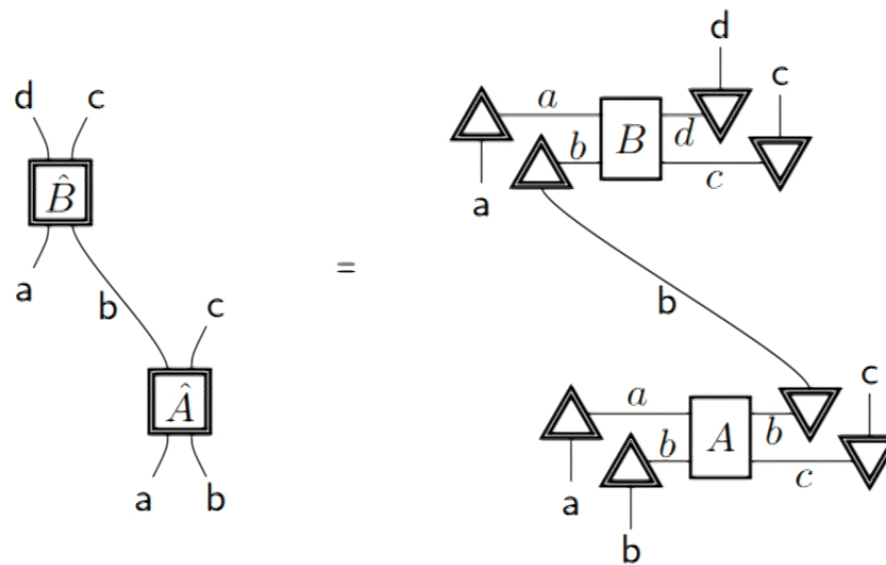


in diagrammatic notation.

## Using full decomposability

$$\hat{A}_{a_1 b_2}^{b_3 c_4} \hat{B}_{a_5 b_3}^{d_6 c_7} = A_{a_1 b_2}^{b_3 c_4} \hat{X}_{a_1}^{a_1} \hat{X}_{b_2}^{b_2} b_3 \hat{X}^{b_3}_{c_4} \hat{X}^{c_4}_{a_5} B_{a_5 b_3}^{d_6 c_7} \hat{X}_{a_5}^{a_5} \hat{X}_{b_3}^{b_3'} d_6 \hat{X}^{d_6}_{c_7} \hat{X}^{c_7}_{c_7}$$

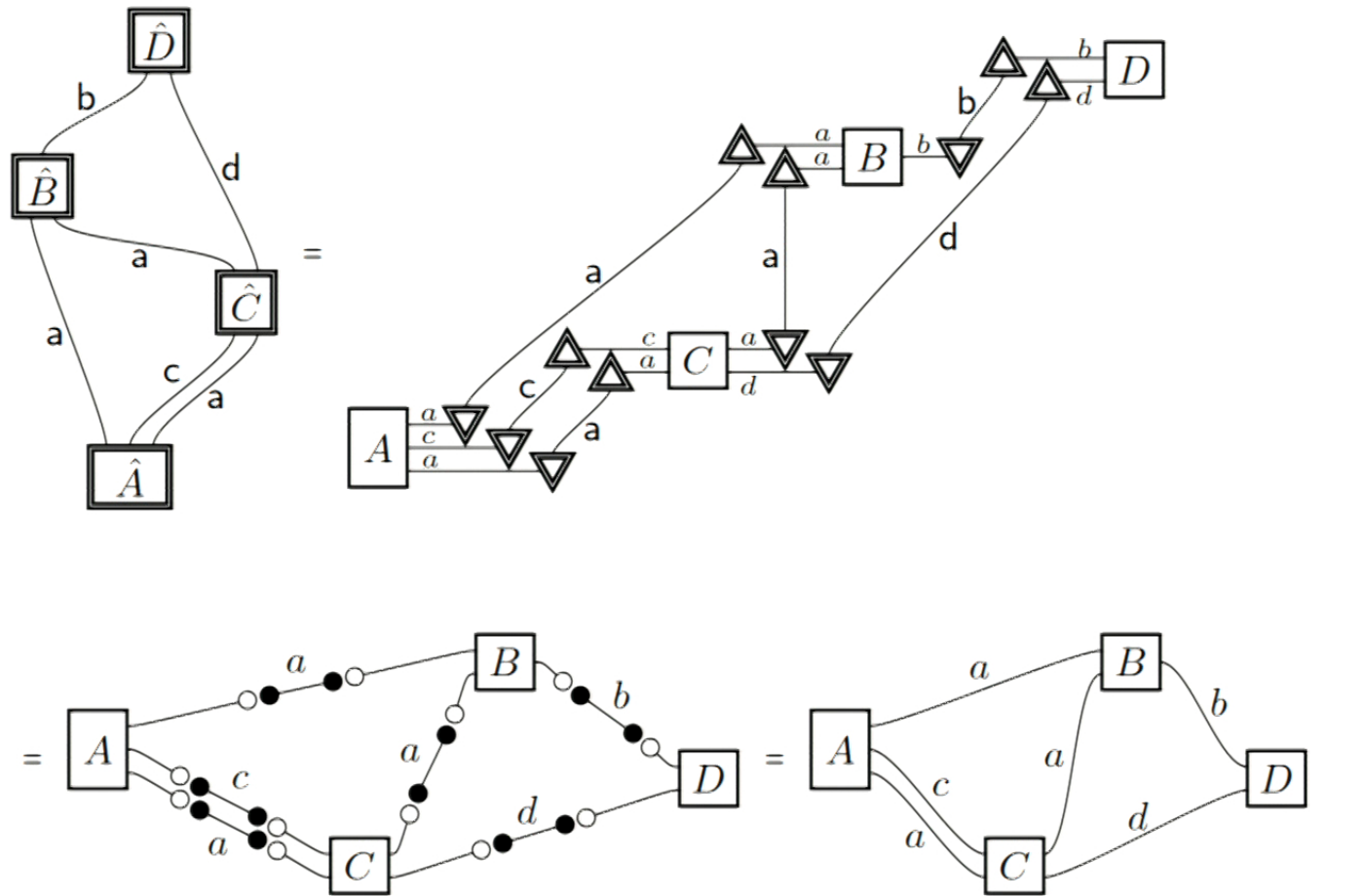
or, in diagrammatic notation,



This operator is in the space

$$\mathcal{V}_{a_1 b_2 a_5}^{c_4 d_6 c_7} = \mathcal{V}_{a_1} \otimes \mathcal{V}_{b_2} \otimes \mathcal{V}_{a_5} \otimes \mathcal{V}^{c_4} \otimes \mathcal{V}^{d_6} \otimes \mathcal{V}^{c_7}$$

## General operator circuit





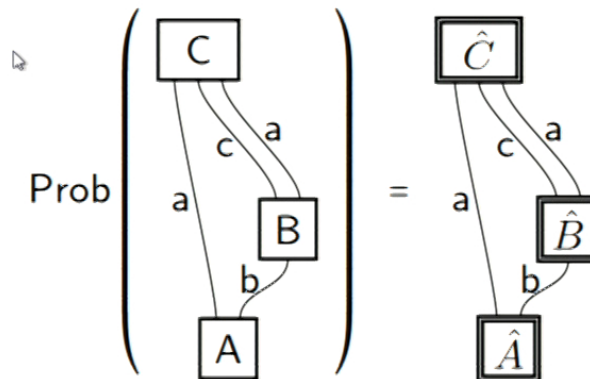
## Operation-operator correspondence.

We will say that operations correspond to operators if there is a mapping from operations,  $A_{a_1 b_2 \dots c_3}^{d_4 e_5 \dots f_6}$ , to operators,  $\hat{A}_{a_1 b_2 \dots c_3}^{d_4 e_5 \dots f_6}$ , such that the probability for *any* circuit comprised of operations is equal to the trace of the circuit operator obtained under this mapping.

If operations correspond to operators then, for example

$$\text{Prob}(A_{a_1 b_2}^{a_1 b_2} B_{b_2 c_3}^{c_3 a_4} C_{a_1 c_3 a_4}) = \hat{A}_{a_1 b_2}^{a_1 b_2} \hat{B}_{b_2 c_3}^{c_3 a_4} \hat{C}_{a_1 c_3 a_4}$$

Same example in diagrammatic notation

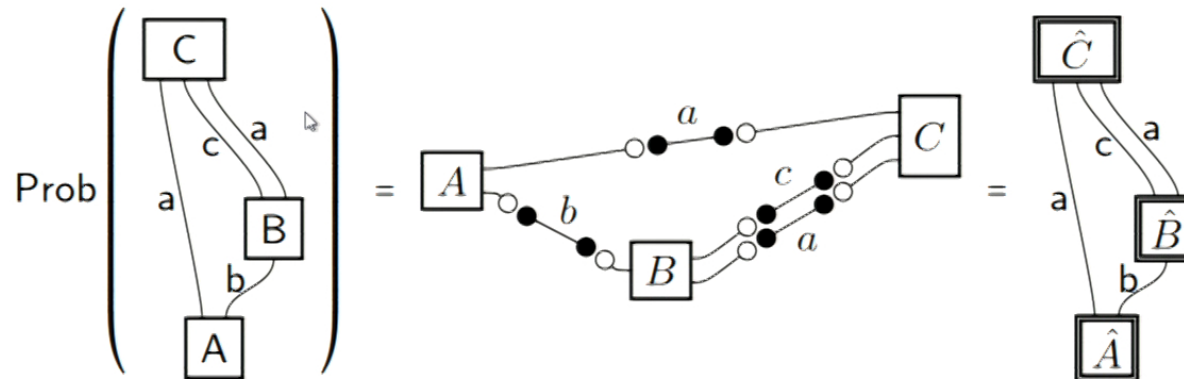


## Condition for correspondence - KEY IDEA!

Have correspondence if can find fiducial sets such that have equal hopping metrics

$$p \left( \begin{array}{c} \triangle \bullet a \\ a \\ \bullet a \triangle \end{array} \right) = \bullet \overset{a}{\text{---}} \bullet = \begin{array}{c} \triangle \bullet a \\ a \\ \bullet a \triangle \end{array}$$

since, for example,

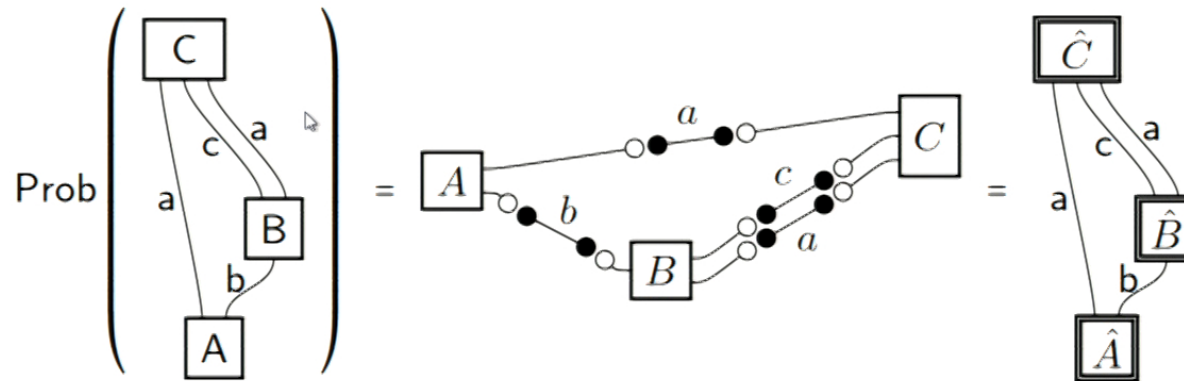


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since, for example,





## Determining operator corresponding to operation

Do local process tomography to get duotensor with all black dots:

$$\begin{array}{c} a \\ b \end{array} \bullet \bullet \boxed{A} \bullet \bullet \begin{array}{c} c \\ d \end{array} = \text{Prob} \left( \begin{array}{c} \triangle \bullet c \\ c \triangle \bullet d \\ \boxed{A} \\ a \triangle \bullet \\ b \triangle \bullet \end{array} \right)$$

Convert to duotensor with all white dots (using  $\circ\text{---}\circ$ ) and then operator given by

The diagram shows an equality between two expressions. On the left is a box labeled  $\hat{A}$  with two vertical lines passing through it. On the right is a more complex expression: a box labeled  $A$  with four inputs (two black dots on top, two white dots on bottom) and two outputs (one white dot on top, one black dot on bottom). To the left of the  $A$  box are two triangles (one pointing up, one pointing down) connected to the top and bottom inputs of  $A$ . To the right of the  $A$  box are two triangles (one pointing down, one pointing up) connected to the top and bottom outputs of  $A$ .

# Physicality theorem

Assume operator circuits must take values between 0 and 1 and we allow at least

1. all rank one projectors,  $\hat{A}^{a_1}$  (for all systems a),
2. all rank one projectors,  $\hat{C}_{b_2}$  (for all systems b),
3. the identity effect,  $\hat{I}_{b_2}$  (for all systems b),

then all operators,  $\hat{B}_{c_3}^{d_4}$ , we must have

Positive input transpose

$$\hat{B}_{c_3}^{d_4}$$

Output trace less than identity

$$\hat{B}_{c_3}^{d_4} \hat{I}_{d_4} \leq \hat{I}_{c_3}$$

We call operators satisfying these two conditions *physical*.

## Complete sets

**A complete set of operations**,  $\{B_{a_1}^{b_2}[l] : l = 1 \text{ to } L\}$  is a set of operations corresponding to the same apparatus use with disjoint outcome sets whose union is the set of all possible outcomes for this apparatus.

**A complete set of physical operators**,  $\{\hat{B}_{a_1}^{b_2}[l] : l = 1 \text{ to } L\}$ , is a set for which every operator positive input transpose and

$$\sum_{l=1}^L B_{a_1}^{b_2}[l] \hat{I}_{b_2} = \hat{I}_{a_1} \quad (5)$$

Elements of a complete set of physical operators satisfy

$$B_{a_1}^{b_2}[l] \hat{I}_{b_2} \leq \hat{I}_{a_1}$$

so are physical.



# Mathematical axioms

**Axiom 1** Operations correspond to operators.

**Axiom 2** All complete sets of physical operators correspond to complete sets of operations.

Or, more glibly,

*All complete sets of operations correspond to complete sets of physical operators and vice versa.*

Equivalent to usual formulation of quantum theory.

## End comments

- ▶ unification of all objects - states, effects, transformations, fragments in general all described by positive operators.
- ▶ simple way to combine them.
- ▶ Don't need to foliate.
- ▶ Curious time asymmetry in the physicality condition.
- ▶ formalism locality.
- ▶ played a crucial role in my recent reconstruction.
- ▶ relation with q-combs, multi-time, general boundary, and Leifer Spekkens work.
- ▶ embraces the quantum pictorialism revolution!
- ▶ quantum field theory, quantum gravity, ....

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