

Title: PSI 17/18 - Foundations of Quantum Mechanics - Lecture 6

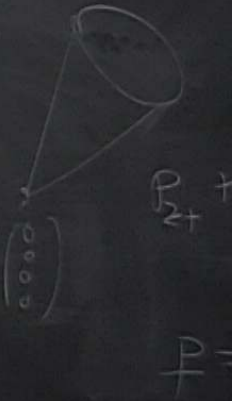
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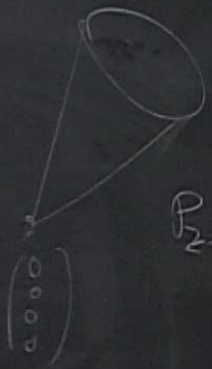
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Abstract:

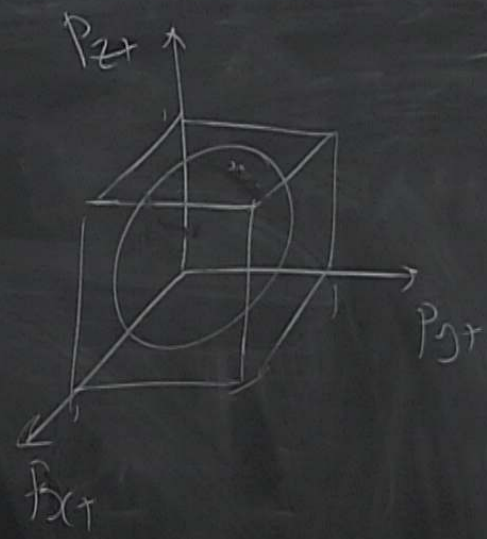
The shape of Quantum state space .

$$\rho = \begin{pmatrix} p_{z+} & a \\ a^* & p_{z-} \end{pmatrix} \Leftrightarrow \mathbb{F} = \begin{pmatrix} p_{z+} \\ p_{z-} \\ p_{x+} \\ p_{x-} \end{pmatrix}$$

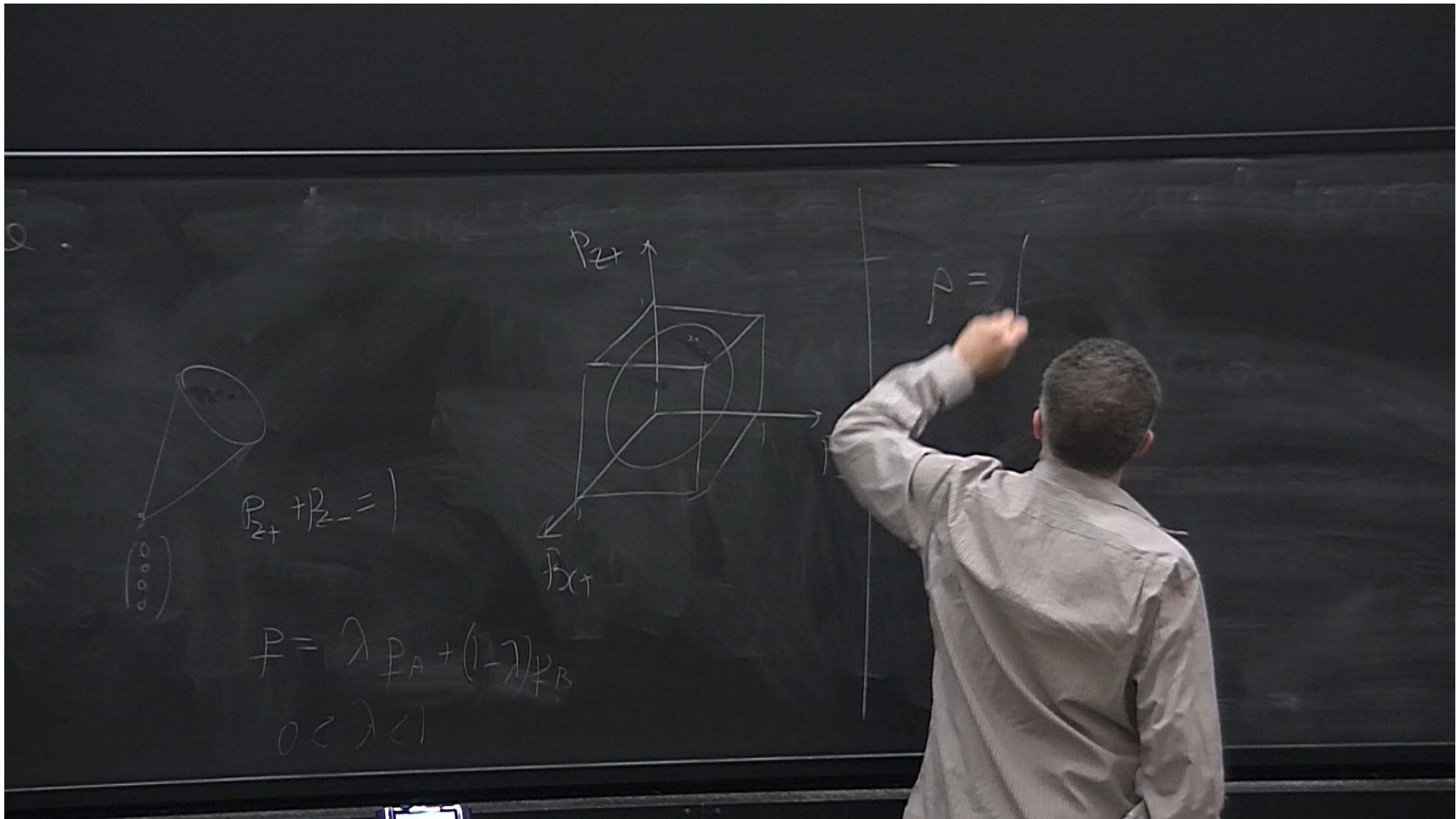


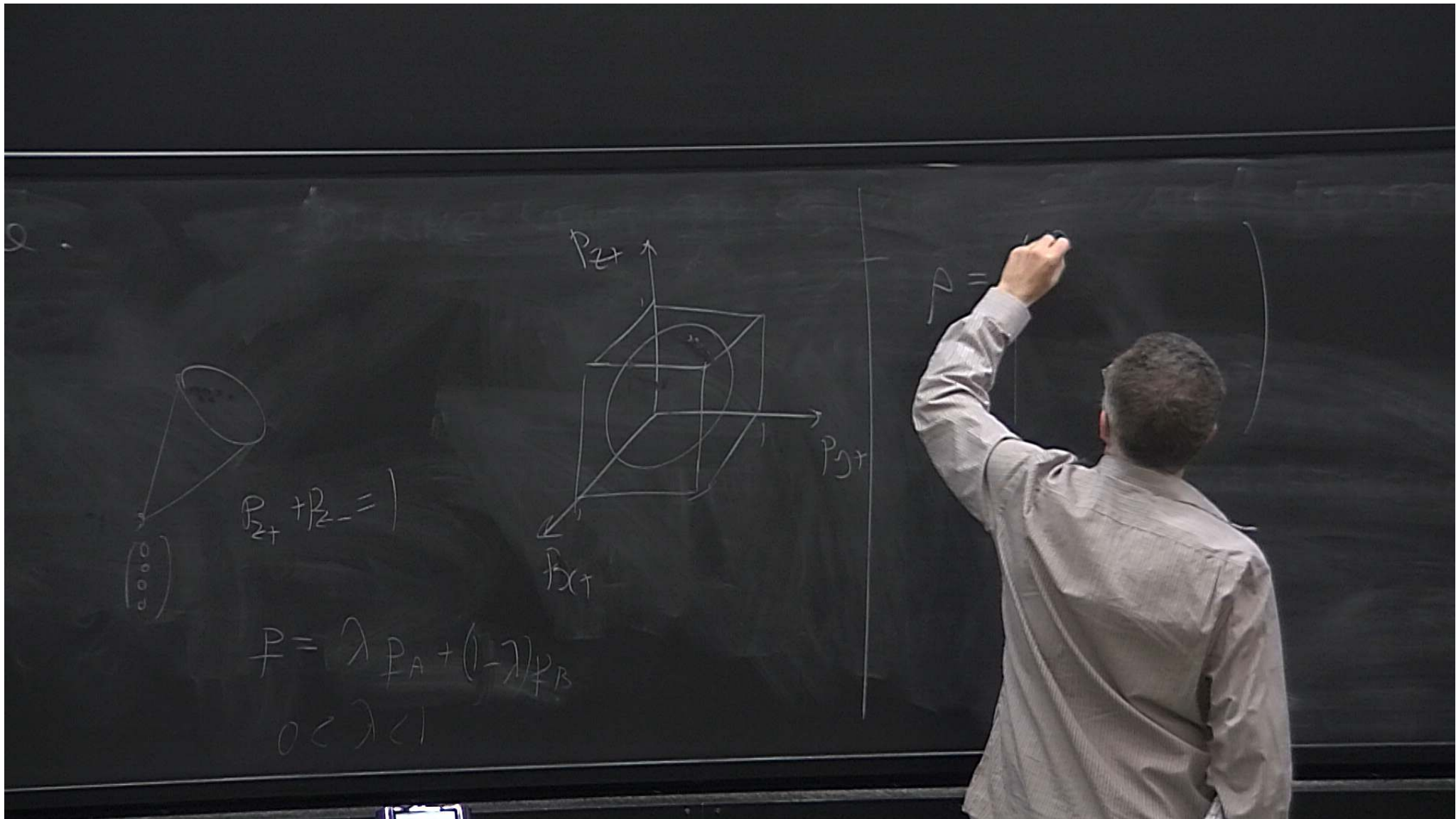


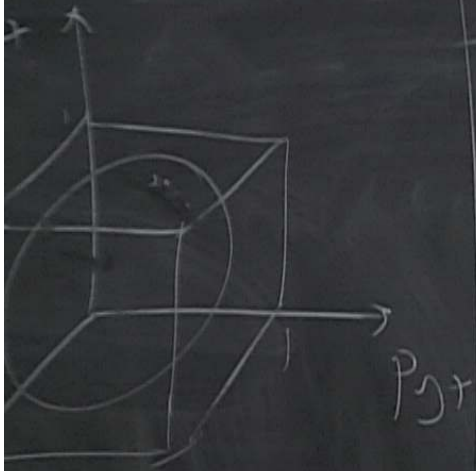
$$P_{z+} + P_{z-} = 1$$



$$P = \lambda P_A + (1-\lambda) P_B$$
$$0 < \lambda < 1$$







$$P = \begin{pmatrix} P_1 & & & \\ & a_{11} & & \\ & & \ddots & \\ & & & a_{nn} \\ & & & & P_n \end{pmatrix} \iff P = \begin{pmatrix} P_1 \\ \vdots \end{pmatrix}$$

$$K = N + 2 \cdot \frac{N(N-1)}{2} = N^2$$

normalized states  $|N^2 - 1|$

boundary of set of states  $N^2 - 2$

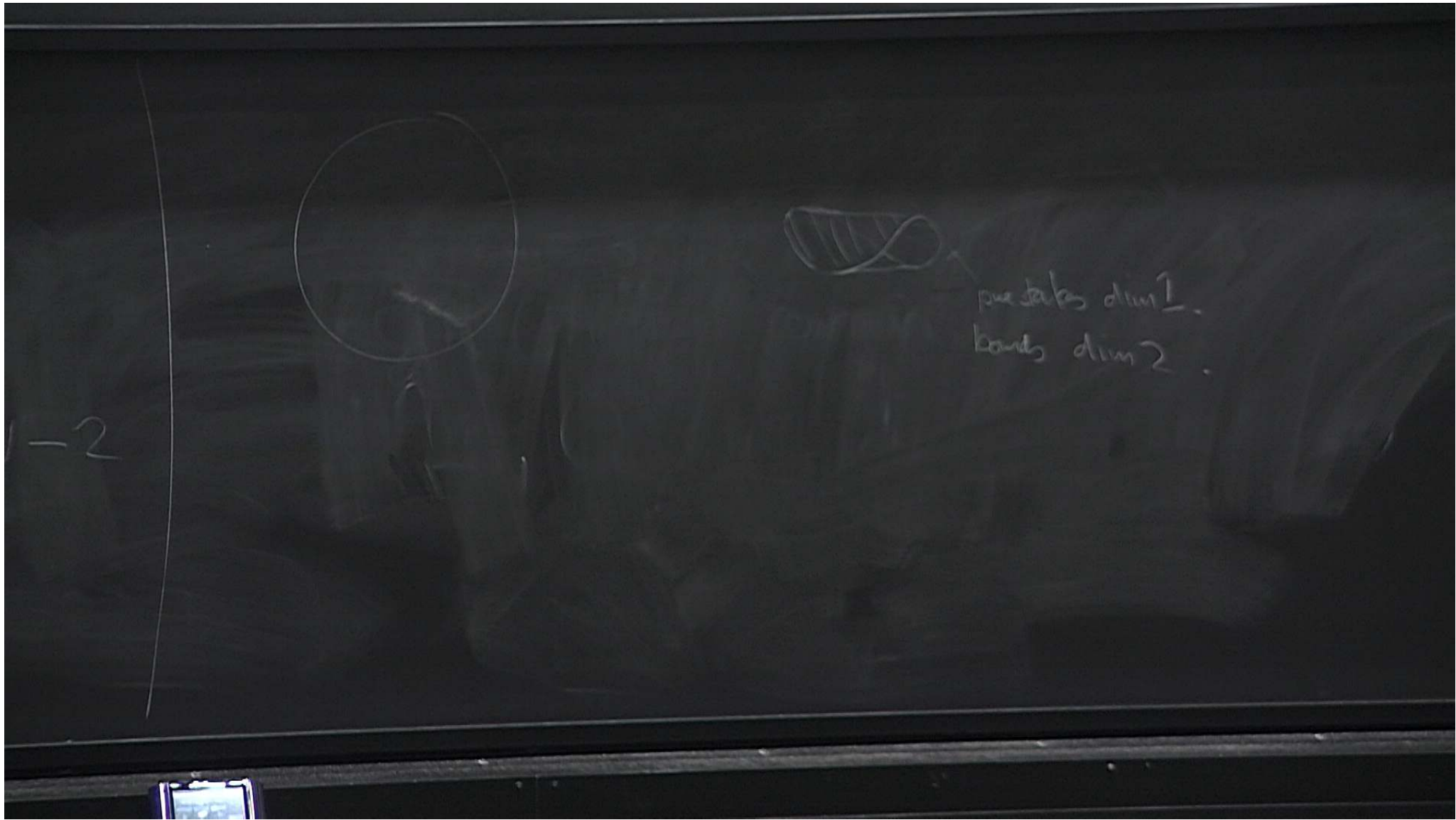
dimension of set of pure states  $= L =$

normalized states  $N-1$

boundary of set of states  $N^2-2$

dimension of set of pure states  $= L = 2N-2$

$$|\psi\rangle = \sum_{n=1}^N a_n |n\rangle$$





pentakis dim 1.  
kardis dim 2.

-2

$$\rho = \rho^2$$

$N^2$  equations  
quadratic in probs.



pure states dim 1.  
bands dim 2.

2004

$$P = P^2$$

$N^2$  equations  
quadratic in probs.



pure states dim 1.  
bands dim 2.

2004 Jones Linden

fuchs

$A$  (Hermitian),  $\neq 1$

$$\text{tr}(\rho^2) = \text{tr}(\rho^3) = 1$$



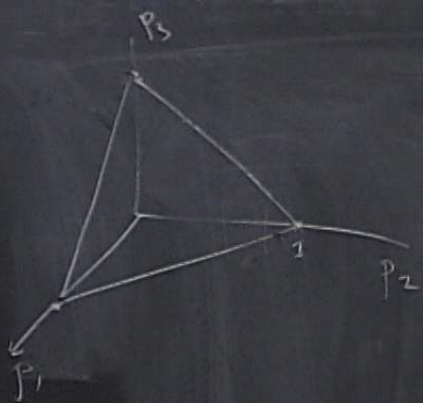
pure states dim 1  
bands dim 2

2004 Jones Linden  
Fuchs

1 cubic eqn  
1 quadratic eqn

$\rho$  (Hermitian)  $\rho^\dagger$

$$\text{tr}(\rho^2) = \text{tr}(\rho^3) = 1$$



$$\sum_k p_k = 1$$

$$\sum_k p_k^2 = 1$$

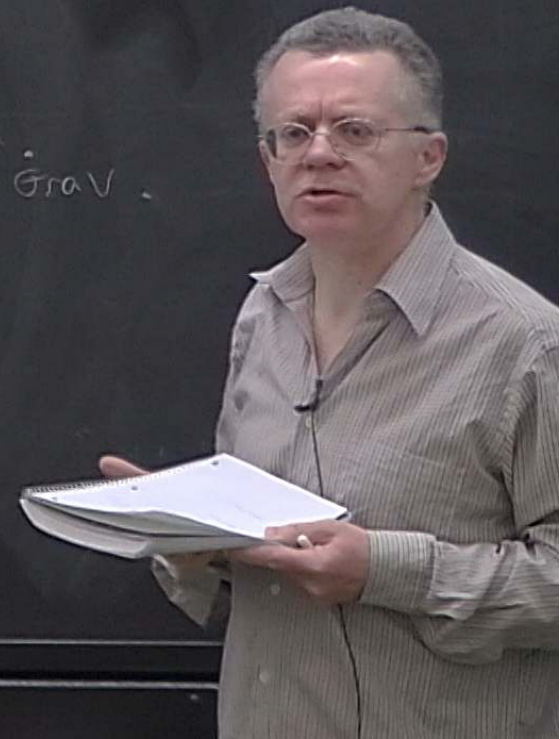
Why QT?

Why Kepler's laws?

Why QT?

Why Kepler's laws?

Newton  
3 laws of motion  
+ Law of Grav.



Why QT?

Why Kepler's laws?

Why Lorentz trans.

Newton

3 laws of mech

+ Law of Grav.

Why QT?

Why Kepler's laws?

Why Lorentz trans.

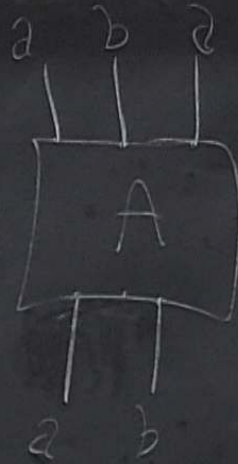
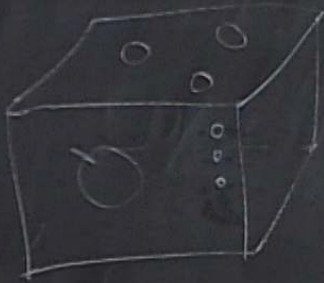
Newton.

3 laws of mech  
+ Law of Grav.

Einstein's SR.

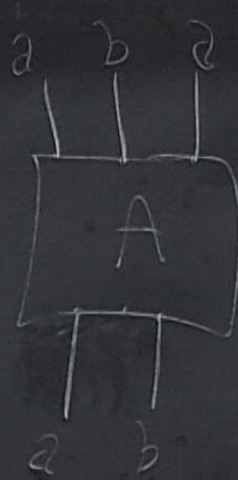
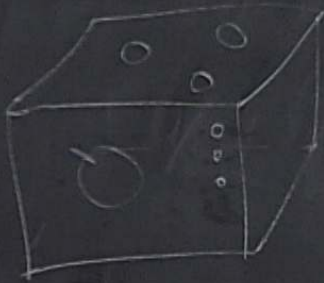


apparatus



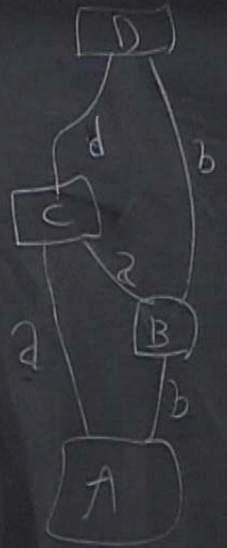
(setting  
+ outcome  
given)

apparatus

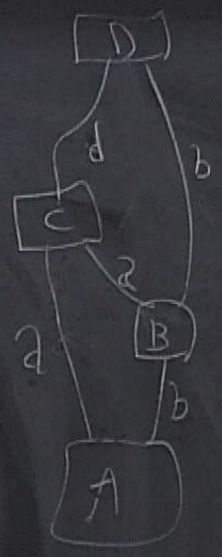
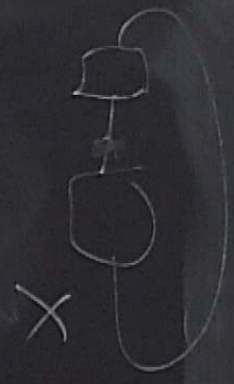


operation

(setting  
+ outcome  
given)



ation.  
ething  
outcome  
given)



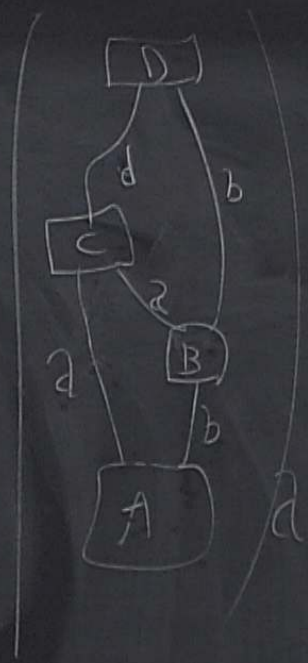
DAG

$$a = \text{mathsf{a}}$$



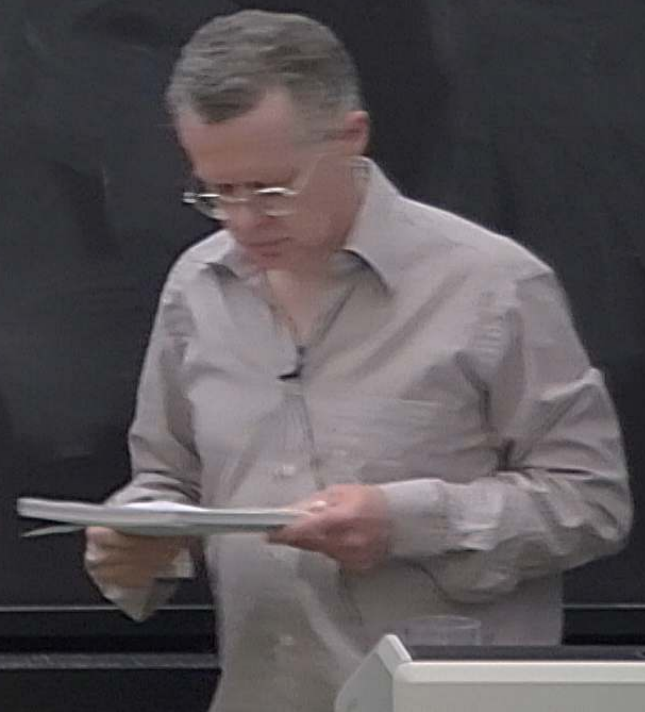
ation.  
ething  
outcome  
given)

prob

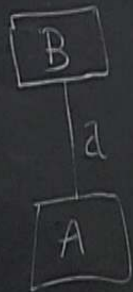


DAG

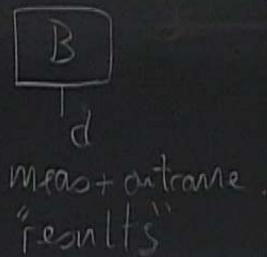
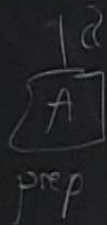
$$a = \text{mathsf{a}}$$



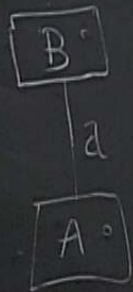
start with simple circuits



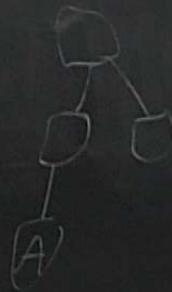
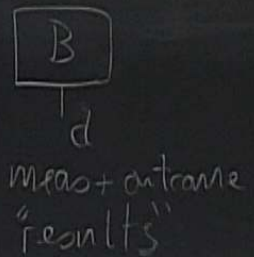
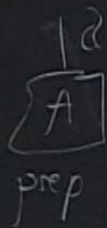
prep  $\Rightarrow$  state  $\Leftrightarrow$  math. object  
that we can use to calculate  
the prob for any circuit  
with this prep.

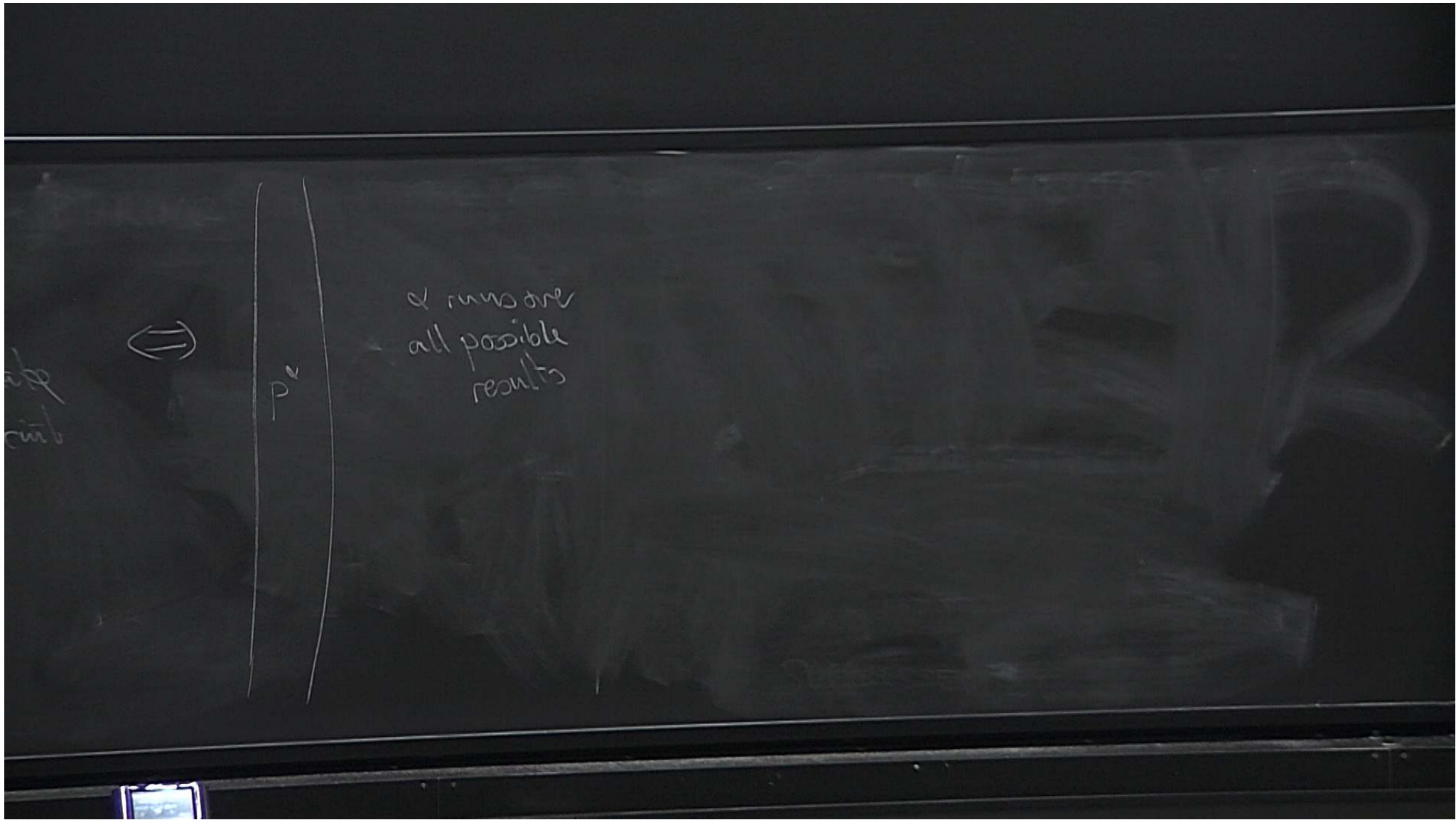


start with simple circuits



prep  $\Rightarrow$  state  $\Leftrightarrow$  math. object  
 that we can use to calculate  
 the prob for any circuit  
 with this prep.





$\Leftrightarrow$

$P_e$

$\alpha$  runs over  
all possible  
results

$$P = \begin{pmatrix} \vdots \\ P^\alpha \\ \vdots \end{pmatrix} \quad \alpha \in F$$

s.t. that we can write

$$\text{prob} = \underline{1} \cdot P$$

$\Leftrightarrow$

$P_e$

$\alpha$  runs over  
all possible  
results

$$P = \begin{pmatrix} \vdots \\ P^\alpha \\ \vdots \end{pmatrix} \quad \alpha \in F$$

s.t. that we can write

$$\text{prob} = \sum_B P \cdot P \xrightarrow{A}$$

$\Leftrightarrow$

$P_e$

$\alpha$  runs over  
all possible  
results

$$P = \begin{pmatrix} \vdots \\ P^\alpha \\ \vdots \end{pmatrix} \quad \alpha \in F$$

s.t. that we can write

$$\text{prob} = \sum_B P \cdot P \rightarrow A$$

$\Rightarrow$  should have  
most possible elements

$$= \begin{pmatrix} \vdots \\ p_\alpha \\ \vdots \end{pmatrix}$$

$$\alpha \in F$$

$$p = \begin{pmatrix} p_{z+} \\ p_{z-} \\ p_{x+} \\ p_{y+} \end{pmatrix}$$

$$F = (z+, z-, x+, y+)$$

that we can write

$\alpha$  runs over  
all possible  
results

$$p = \begin{pmatrix} \vdots \\ p_\alpha \\ \vdots \end{pmatrix} \quad \alpha \in F$$

$$p = \begin{pmatrix} p_{z+} \\ p_{z-} \\ p_{x+} \\ p_{y+} \end{pmatrix}$$

$$F = \{z+, z-, x+, y+\}$$

s.t. that we can write

$$\text{prob} = \sum_{\substack{A \\ B}} p$$

$p$  should have  
fewest possible elements

prep  
"results"

linear compression is optimal if you allow arbitrary matrices

$$F = \sum_{i=1}^K \lambda_i \phi_i \quad K = |F|$$

choose  $\phi_i$  to be lin indep. we know there are  $K$  such states.  
(because compression is optimal)

prep  
"results"

linear compression is optimal if you allow arbitrary mixtures

$$p = \sum_{i=1}^K \lambda_i p_i + \lambda_0 \underline{0} \quad K = |F|$$
$$0 \leq \lambda_i \quad \sum_{i=1}^K \lambda_i \leq 1$$

choose  $p_i$  to be lin indep. we know there are  $K$  such states.  
(because compression is optimal)

prep  
results

linear compression is optimal if you allow arbitrary mixtures

$$\rho = \sum_{i=1}^K \lambda_i \rho_i + \lambda_0 \rho_0 \quad K = |F|$$
$$0 \leq \lambda_i \quad \sum_{i=1}^K \lambda_i \leq 1$$

choose  $\rho_i$  to be lin indep. we know there are  $K$  such states.  
(because compression is optimal)

to characterize this set of mixed states  
we need  $K$   $\lambda_i$ .

results

linear compression is optimal if you allow arbitrary mixtures

$$p = \sum_{i=1}^K \lambda_i p_i + \lambda_0 \underline{0} \quad K = |F|$$

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