

Title: PSI 2016/2017 Quantum Information (Review) - Lecture 6 (Daniel Gottesman)

Date: Feb 28, 2017 09:00 AM

URL: <http://pirsa.org/17020117>

Abstract:

$$|0\rangle = (|000\rangle + |111\rangle)^{\otimes 3}$$

$$|1\rangle = (|000\rangle - |111\rangle)^{\otimes 3}$$

$$X = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$$

$$Y = \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix} = iXZ$$

$$Z = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}$$

$$Z \otimes Z \otimes I \otimes I \dots$$

$$Z \otimes Z$$

$$Z \otimes Z$$

$$Z \otimes Z$$

$$Z \otimes Z$$

$$Z \otimes Z$$

$$S = \{M \in \mathcal{M}\}$$

Def: Pauli g
with overall

$$\begin{array}{cccccccc} X & \otimes & X & \otimes & X & \otimes & X & \otimes & X \\ X & \otimes & X & \otimes & X & \otimes & X & \otimes & X \end{array}$$

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$$X \otimes X \otimes I \otimes I$$

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$$S = \{ M \in \mathcal{P}_n \mid M|H\rangle = |H\rangle \forall |H\rangle \in \text{code} \}$$

Stabilizer

Def: Pauli group $\mathcal{P}_n$ = n-qubit tensor products of I, X, Y, Z , with overall phase $\pm 1, \pm i$

$$P, Q \in \mathcal{P}_n \Rightarrow \text{either } [P, Q] = 0 \text{ or } \{P, Q\} = 0 \quad \text{e.g. } (X \otimes X, Z \otimes Z) = 0$$

$$P \in \mathcal{P}_n \Rightarrow P^2 = \pm I$$

S is a group

$$M, N \in S \Rightarrow MN \in S:$$

$$(MN)|\psi\rangle = M|N\psi\rangle = |N\psi\rangle$$

$$-I \notin S$$

S is Abelian

$$M, N \in S \quad MN|\psi\rangle = |N\psi\rangle$$

$$NM|\psi\rangle = |N\psi\rangle$$

$$[M, N]|\psi\rangle = 0$$

$$\Rightarrow [M, N] = 0$$

Given S st. S is an Abelian group and $-I \notin S$, then

$$T(S) = \{ |\psi\rangle \mid M|\psi\rangle = |\psi\rangle \forall M \in S \}$$

Def: For a stabilizer code $C \subseteq \mathcal{H}$,

$$T(S(C)) = C.$$

Prop: If $|S| = 2^r$, S has r generators

$$\text{and } \dim T(S) = 2^{n-r} \quad (n \text{ physical qubits})$$

(i.e. $k = n - r$ logical qubits)

S is Abelian

$$\begin{aligned} M, N \in S \quad & MN|\psi\rangle = N|\psi\rangle \\ & NM|\psi\rangle = N|\psi\rangle \\ & [M, N]|\psi\rangle = 0 \\ & \Rightarrow [M, N] = 0 \end{aligned}$$

Def. For a stabilizer code $C \subseteq \mathcal{H}$,
 $T(S(C)) = C$.

$$M \in S, E \in \mathcal{P}_n, \{E, M\} = 0$$

$$M(E|\psi\rangle) = -E M|\psi\rangle = -(E|\psi\rangle)$$

$$[E, M] = 0$$

$$M(E|\psi\rangle) = E M|\psi\rangle = E|\psi\rangle$$

$$N(S) = \{ N \in \mathcal{P}_n \mid [N, M] = 0 \ \forall M \in S \}$$

$$S \subseteq N(S) \quad |N(S)| = 2^{n+k} \cdot 4$$

Undetectable errors = $N(S) \setminus S$.

E, F have same error syndrome

$\Leftrightarrow E, F$ commute with same elements of S

$$\Leftrightarrow [E^{\dagger} F, M] = 0 \ \forall M \in S$$

$$\Leftrightarrow E^{\dagger} F \in N(S)$$

$$\begin{aligned} E^{\dagger} F \in S \quad & E^{\dagger} F|\psi\rangle = |\psi\rangle \quad \forall |\psi\rangle \in T(S) \\ \Rightarrow F|\psi\rangle &= E|\psi\rangle \end{aligned}$$

Thm. S corrects a set C of errors
iff $E^+F \notin N(S) \setminus S \quad \forall E, F \in C$.

Def. Distance of S is minimum weight
of $N(S) \setminus S$. Code is degenerate if S has operators
with weight $< \text{dist}$.
A code correct t -qubit errors if $\text{dist} \geq 2t+1$

X	Z	Z	X	I
I	X	Z	Z	X
X	I	X	Z	Z
Z	X	I	X	Z

$\{ \# \text{ logical qubits} \}$
 $[[5, 1, 3]]$
 $\uparrow$ $\# \text{ physical qubits}$ $\uparrow$ distance

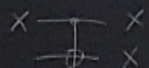
X	X	X	X	I	I	I
X	X	I	I	X	X	I
X	I	X	I	X	I	X
Z	Z	Z	Z	I	I	I
Z	Z	I	I	Z	Z	I
Z	I	Z	I	Z	I	Z

CSS code

$[[7, 1, 3]]$

Fault Tolerance:

Control error propagation:



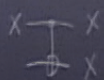
- FT error correction
- FT state preparation
- FT measurement
- FT universal set of gates

Def: A universal set of gates G is a set that can be multiplied together to get (close to) any unitary $U \in U(2^n)$

E.g. $\{\text{single-qubit gates}, \text{CNOT}\}$ exact
 $\{H, R_{\pi/8}, \text{CNOT}\}$ approximate
(ϵ gate)

Fault Tolerance

Control error propagation



- FT error correction
- FT state preparation
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Def. A universal set of gates G is a set that can be multiplied together to get (close to) any unitary $U \in SU(2^n)$

E.g. $\{\text{single-qubit gates, CNOT}\}$ exact
 $\{H, R_{\pi/8}, \text{CNOT}\}$ approximate
 $\begin{pmatrix} e^{-i\pi/8} & 0 \\ 0 & e^{i\pi/8} \end{pmatrix}$