

Title: Boundary States and Entanglement Spectrum from Strong Subadditivity

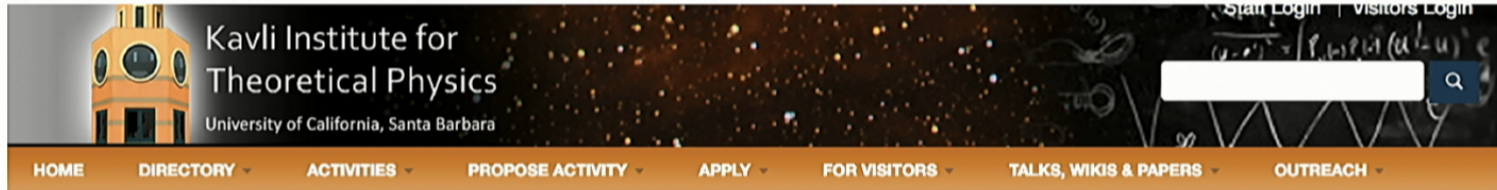
Date: Jul 21, 2016 05:00 PM

URL: <http://pirsa.org/16070061>

Abstract: In this talk I will consider quantum states satisfying an area law for entanglement (e.g. as found in quantum field theory or in condensed matter systems at sufficiently low temperature). I will show that both the boundary state and the entanglement spectrum admit a local description whenever there is no topological order. The proof is based on strong subadditivity of the von Neumann entropy. For topological systems, in turn, I'll show that the topological entanglement entropy quantifies exactly how many extra bits are needed in order to have a local description for the boundary state. This latter result is based on a recent strengthening of strong subadditivity.

Based on joint work with Kohtaro Kato (University of Washington)

# #qinfo17



## Quantum Physics of Information

**Coordinators:** Fernando Brandao, Veronika Hubeny, Stephen Jordan, Renato Renner

**Scientific Advisors:** Daniel Harlow, Patrick Hayden, John Preskill

Quantum information science began about 30 years ago as the union of two of the biggest scientific developments of the last century, quantum mechanics and computer science. The original motivation was to understand the new possibilities offered by quantum mechanics to information processing and computation. In recent years, it has emerged that the field also offers a new perspective for the study of physics, from condensed matter and thermodynamics to quantum gravity.

Recent breakthroughs in the application of quantum information in other areas of physics have served to increase the interest of researchers in high energy, gravity, and condensed matter, to work at the intersection of these fields. This program has the objective of cultivating these growing interdisciplinary discussions. It is the hope that the program will be a turning point in the exciting, rapidly developing dialogue between quantum information and physics as a whole.

Some of the topics covered in the program are quantum computational complexity in physics and the simulation of physical systems, advances in concepts and methods of quantum information of relevance to many body physics and thermodynamics, as well as the role of quantum information in quantum gravity, field theory, and the foundations of quantum theory.



### DATES

Sep 18, 2017 - Dec 15, 2017

### INFORMATION

[Apply](#)

**Application deadline is:**  
Oct 2, 2016.

Deadline: October 2



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1 0 1 0 1  
0 1 Q Stephen Jordan

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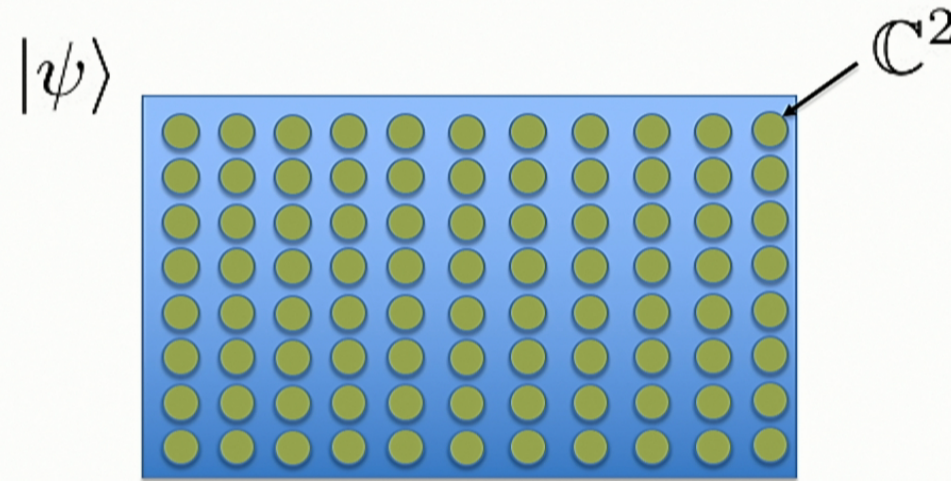
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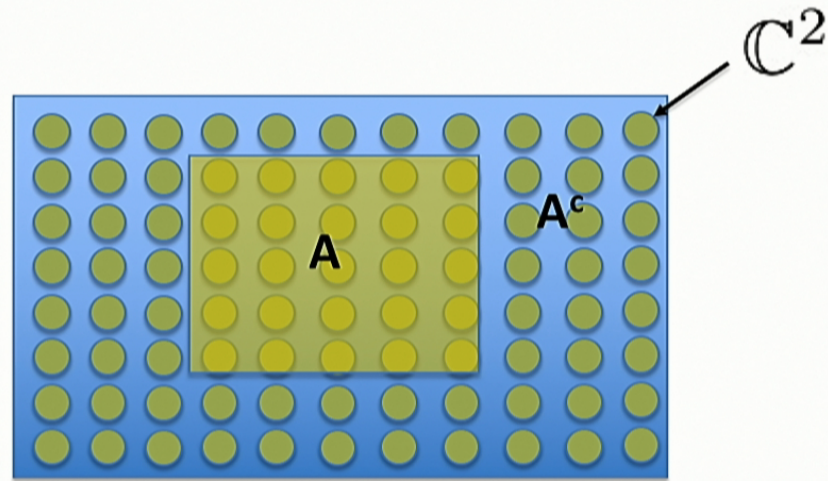
# Entanglement in Many-Body Quantum States





# Entanglement in Many-Body Quantum States

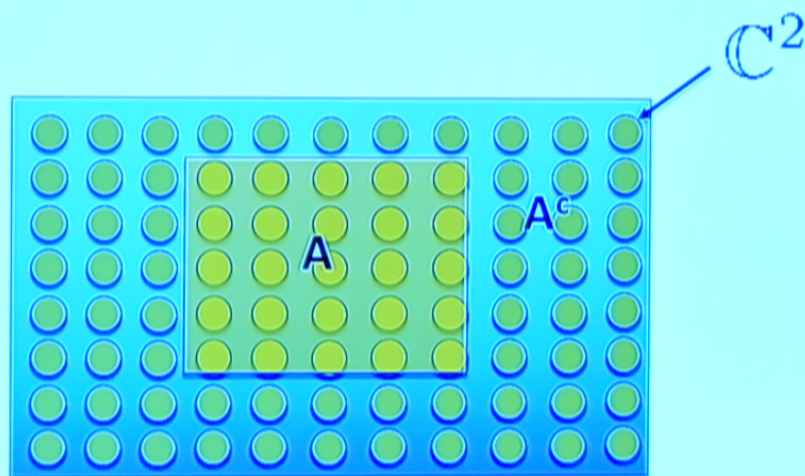
$$|\psi\rangle_{AA^c}$$



$$\text{Entanglement Entropy: } S(A) = -\text{tr}(\rho_A \log \rho_A)$$

# Entanglement in Many-Body Quantum States

$$|\psi\rangle_{AA^c}$$



Entanglement Entropy:  $S(A) = -\text{tr}(\rho_A \log \rho_A)$

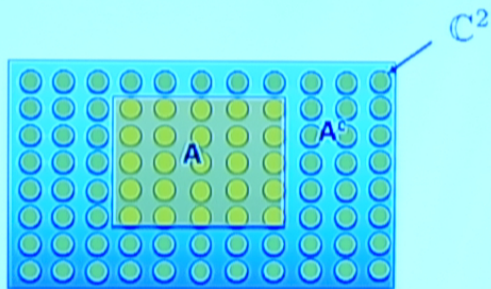
For generic quantum states:  $S(X) \approx \text{vol}(X)$  (Page '93)

What's the behavior of EE for interesting states of matter?



# Area Law

$|\psi\rangle_{AA^c}$



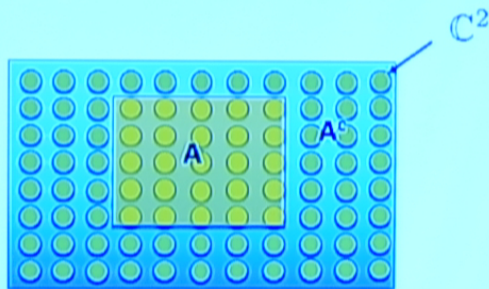
Entanglement is “localized”,  
concentrated around the boundary

For every region  $X$ :  $S(X) = \alpha|\partial X| - \gamma + \dots$

e.g. gapped models, 2+1 CFT (from RT formula)

# Area Law

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For every region  $X$ :  $S(X) = \alpha|\partial X| - \gamma + \dots$

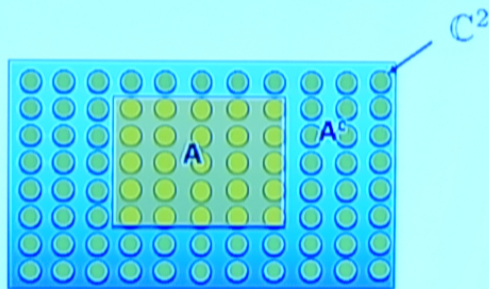
$\gamma$ : Topological EE  
(signature topological order)

$$\gamma = \log \mathcal{D}, \quad \mathcal{D} = \sqrt{\sum_a d_a^2} \quad \mathcal{D}: \text{Quantum dimension}$$



# Area Law

$|\psi\rangle_{AA^c}$



Entanglement is “localized”,  
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For every region  $X$ :  $S(X) = \alpha|\partial X| - \gamma + \dots$

- Topological EE quantifies “non-local entanglement”

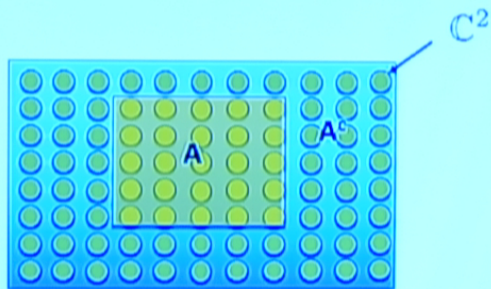
(Kitaev '12)  $\gamma = 0$  : state is adiabatically connected to trivial phase

(Kim '13)  $\log(N) \leq 2\gamma$   $N :=$  number topologically protected states

⋮

# Area Law

$|\psi\rangle_{AA^c}$

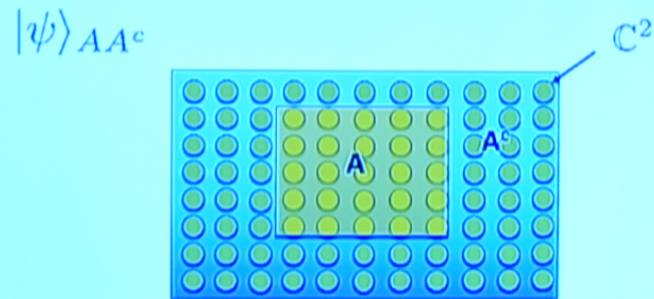


Entanglement is “localized”,  
concentrated around the boundary

What are the consequences of an area law?  
What’s the influence of TEE on the boundary?



# Area Law



Entanglement is “localized”,  
concentrated around the boundary

What are the consequences of an area law?

What’s the influence of TEE on the boundary? This talk:

Area Law



TEE determines locality of  
i) Boundary State  
ii) Entanglement Spectrum

by strong subadditivity and  
stronger subadditivity

# Quantum Information 1.01: Fidelity

... it's a measure of distinguishability between two quantum states.

Given two quantum states their fidelity is given by

$$F(\rho, \sigma) := \text{tr}((\rho^{1/2} \sigma \rho^{1/2})^{1/2})$$

It tells how distinguishable they are by any quantum Measurement

Ex 1:  $F=1$ : same state

Ex 2:  $F=0$  : perfectly distinguishable states



# Quantum Information 1.01: Relative Entropy

... it's another measure of distinguishability  
between two quantum states.

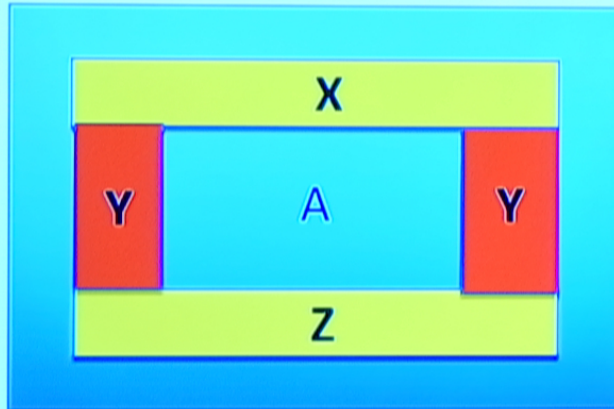
$$\text{Def: } S(\rho\|\sigma) := \text{tr}(\rho(\log(\rho) - \log(\sigma)))$$

Gives optimal exponent for distinguishing the two states

$$\text{Pinsker's inequality: } S(\rho\|\sigma) \geq -\frac{1}{2} \log F(\rho, \sigma)$$

$$S(\rho\|\sigma) \approx 0 \implies \rho \approx \sigma$$

# Topological EE and Locality of Boundary States

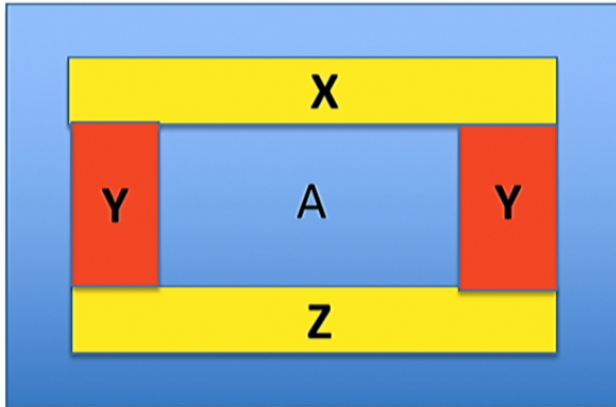


$\rho_{XYZ}$ : reduced state on XYZ

XYZ Boundary of A



# Topological EE and Locality of Boundary States



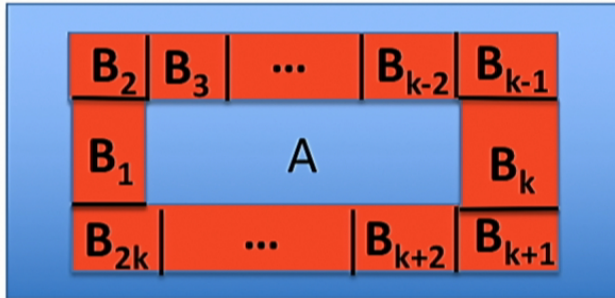
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**Result 1.** If  $S(X) = \alpha|\partial X| - \gamma + \dots$  :

$$\gamma \approx \min_{H_{XY}, H_{YZ}} S(\rho_{XYZ} \| \exp(H_{XY} + H_{YZ}) / \text{tr}(\dots))$$

# Topological EE and Locality of Boundary States



$\rho_{XYZ}$ : reduced state on XYZ

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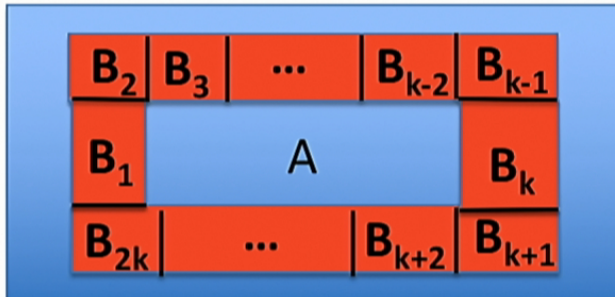
**Result 1.** If  $S(X) = \alpha|\partial X| - \gamma + \dots$  :

$$\gamma \approx \min_{H_{XY}, H_{YZ}} S(\rho_{XYZ} \| \exp(H_{XY} + H_{YZ}) / \text{tr}(\dots))$$

$$\approx \min_{H_{B_1 B_2}, \dots, H_{B_{2k-1} B_{2k}}} S(\rho_{B_1 \dots B_{2k}} \| \exp(H_{B_1 B_2} + \dots + H_{B_{2k-1} B_{2k}}) / \text{tr}(\dots))$$



# Topological EE and Locality of Boundary States



$\rho_{XYZ}$ : reduced state on XYZ

XYZ Boundary of A

Obs 1:  $\gamma = 0$

$$\implies \rho_B \approx \exp(H_{B_1 B_2} + \dots H_{B_{2k-1} B_{2k}} / \text{tr}(\dots))$$

Obs 2: Thermal states has same on-site symmetries as original state

Obs 3: Thermal state is max entropy state consistent with local constraints

# TEE gives number of non-local bits

Interpretation relative entropy (Anshu *et al* '14)

Alice



knows  $\rho$

Bob

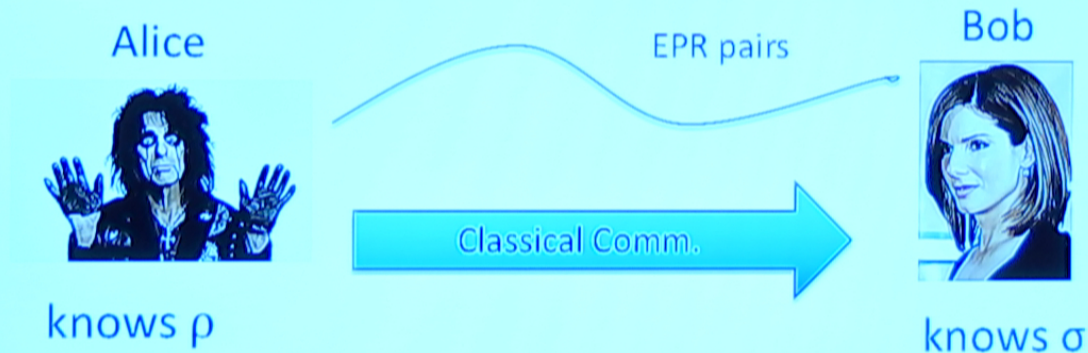


knows  $\sigma$



# TEE gives number of non-local bits

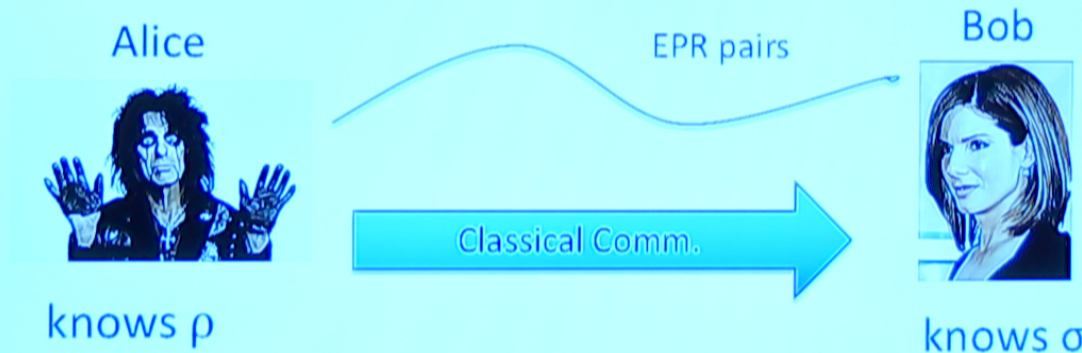
Interpretation relative entropy (Anshu *et al* '14)



What's the minimum classical comm. required for Bob to learn  $\rho$ ?  
(i.e. to be able to prepare a copy of  $\rho$ )

# TEE gives number of non-local bits

Interpretation relative entropy (Anshu *et al* '14)



$\approx S(\rho \parallel \sigma)$  necessary and sufficient for Bob to prepare a copy of  $\rho$

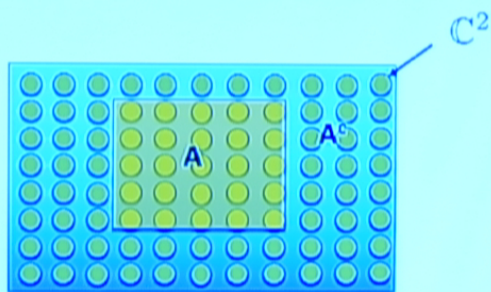
$$\gamma \approx \min_{\sigma \in \text{Local Gibbs State}} S(\rho_{B_1 \dots B_{2k}} \parallel \sigma) \text{ gives number of non-local bits of } \rho$$

obs: Consistent with  $\gamma = \log(\text{quantum dimension})$



# Entanglement Spectrum

$|\psi\rangle_{AA^c}$



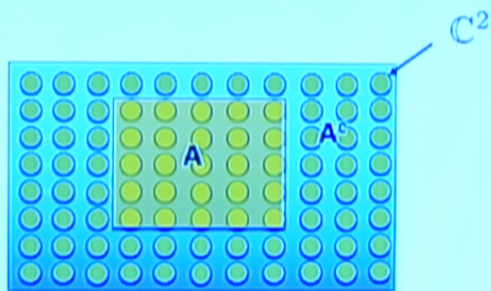
$\lambda(\rho_A)$  : eigenvalues of  $\rho_A$   
Entanglement Spectrum

Area law statement about  $-\sum_i \lambda_i \log \lambda_i$

What can we say about the whole spectrum?

# Entanglement Spectrum

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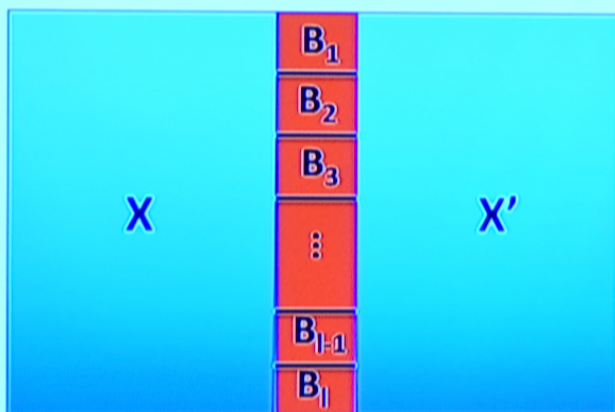
(Haldane, Li '08, Cirac, Poiblanç, Schuch, Verstraete '11, ...)

$\gamma=0$ : matches spectrum thermal state local model

$\gamma \neq 0$ : matches spectrum thermal state local model  
after projecting into topological superselection sector



# Entanglement Spectrum



We assume translation invariance  
s.t.  $\rho_X = \rho_{X'}$

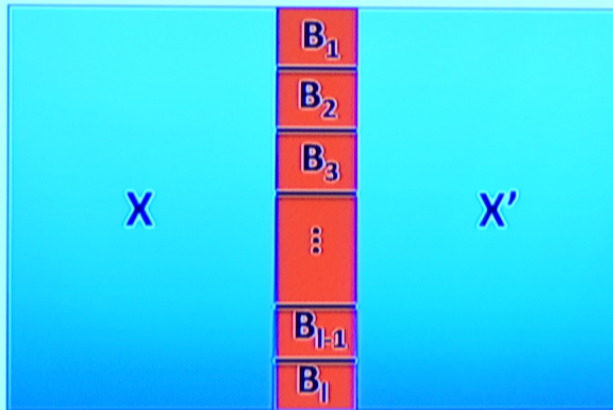
Result 2: If  $S(X) = \alpha|\partial X| - \gamma + \dots$  :

$$\gamma = 0 \implies \lambda(\rho_X)^{\otimes 2} \approx \lambda(e^{\sum_k H_{B_k, B_{k+1}}})$$

$$\gamma \neq 0 \implies \lambda(\rho_X)^{\otimes 2} \approx \lambda(\sigma),$$

$$\text{tr}_{B_1}(\sigma) = e^{\sum_{k>1} H_{B_k, B_{k+1}}}$$

# Result 2 from 1



From area law assumption:  
(more later)

$$\rho_{XX'} \approx \rho_X \otimes \rho_{X'}$$

$$\lambda(\rho_{XX'}) = \lambda(\rho_B) \implies \lambda(\rho_X) \otimes \lambda(\rho_{X'}) \approx \lambda(\rho_B)$$

Uhlmann's theorem There is an isometry  $U : B \rightarrow B_X B_{X'}$  s.t.

$$U|\psi\rangle_{XB_{X'}} \approx |\phi\rangle_{XB_X} \otimes |\phi'\rangle_{XB_{X'}} \quad \rho_X = \text{tr}_{B_{X'}}(|\phi\rangle\langle\phi|_{XB_X})$$

U maps degrees of freedom of X and X' into B



# Why does it hold?

We want to show:

$$\begin{aligned}\gamma &\approx \min_{H_{XY}, H_{YZ}} S(\rho_{XYZ} \| \exp(H_{XY} + H_{YZ}) / \text{tr}(\dots)) \\ &\approx \min_{H_{B_1 B_2}, \dots, H_{B_{2k-1} B_{2k}}} S(\rho_{B_1 \dots B_{2k}} \| \exp(H_{B_1 B_2} + \dots + H_{B_{2k-1} B_{2k}}) / \text{tr}(\dots))\end{aligned}$$

$\chi = 0$  : follow from *strong subadditivity* (SSA) (Lieb, Ruskai '73)

$$S(AB) + S(BC) \leq S(ABC) + S(B)$$

$\chi \neq 0$  : follows from a *strengthening* of SSA (Fawzi and Renner '14)

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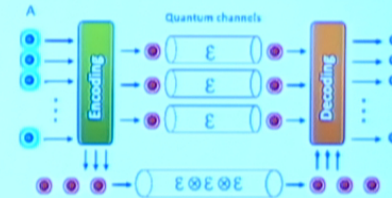
$X \neq \emptyset$  : follows from a *strengthening* of SSA (Fawzi and Renner '14)



# Applications of SSA

Used to prove optimal rates for nearly every quantum information protocol.

- Channel capacities (classical, quantum, private)
- Distillable Entanglement
- ....

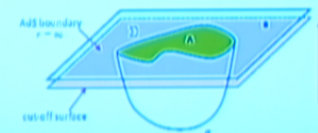


(Casini, Huerta, Myers ...) SSA + Lorentz Invariance:

- Entropic proof of the  $c$ -theorem (irreversibility of renormalization flow)
- Proof of Bekenstein's and Bousso's bound



(Ryu-Takayanagi, Headrick, ...) Test for holographic proposals of entropy

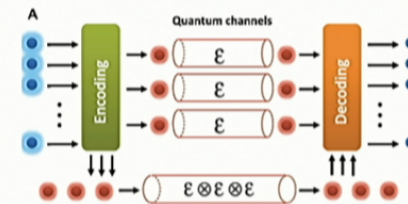


Many others...

# Applications of SSA

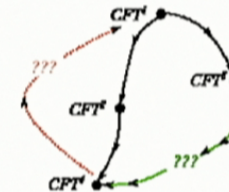
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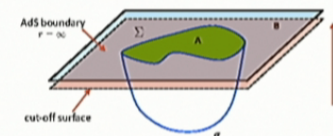


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Many others...



# Conditional Mutual Information

Given  $\rho_{ABC}$ ,

$$\begin{aligned} I(A : C|B) &:= S(AB) + S(BC) - S(ABC) - S(B) \\ &= S(\rho_{ABC} \| \exp(\log(\rho_{AB}) + \log(\rho_{BC}) - \log(\rho_B))) \end{aligned}$$

**Strong subadditivity:**  $I(A : C|B) \geq 0$

# Conditional Mutual Information

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**Stronger subadditivity** (Fawzi-Renner '14):

$$I(A : C|B) \geq \frac{1}{2} \min_{\Lambda: B \rightarrow BC} -\log(F(\rho_{ABC}, \Lambda(\rho_{AB})))$$

$$I(A : C|B) \approx 0 \implies I_A \otimes \Lambda^{B \rightarrow BC}(\rho_{BC}) \approx \rho_{ABC}$$

 quantum channel



# Conditional Mutual Information

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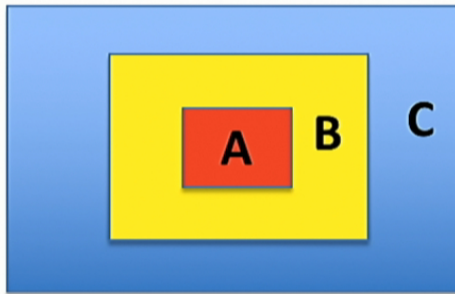
$$I(A : C|B) \geq \frac{1}{2} \min_{\Lambda: B \rightarrow BC} -\log(F(\rho_{ABC}, \Lambda(\rho_{AB})))$$

Can reconstruct the state ABC from reduction on AB by acting on B only



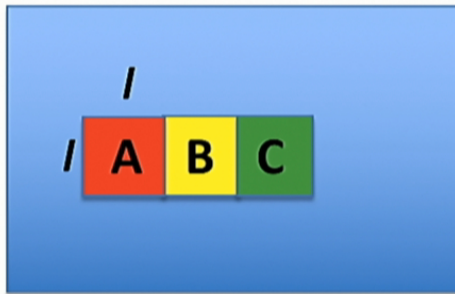


# Consequence of Area Law: State Reconstruction



For every ABC with trivial topology:

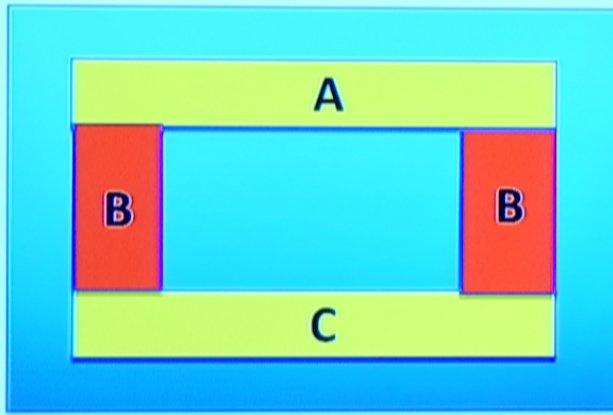
$$I(A : C|B) \approx 0$$



$$\begin{aligned} I(A : C|B) &= S(AB) + S(BC) - S(ABC) - S(B) \\ &= \alpha(|\partial(AB)| + |\partial(BC)| - |\partial(ABC)| - |\partial(B)|) + \dots \\ &= \alpha(6l + 6l - 8l - 4l) + \dots \end{aligned}$$

# TEE as Conditional Mutual Info

(Kitaev, Preskill '05, Levin, Wen '05)



$$\gamma = I(A : C|B) + \dots$$

$$\begin{aligned} & I(A : C|B) \\ &= S(AB) + S(BC) - S(ABC) - S(B) \\ &= \alpha(\partial(AB) + |\partial(BC)| - |\partial(ABC)| - |\partial(B)|) - \gamma - \gamma + \gamma + 2\gamma + \dots \\ &= \gamma + \dots \end{aligned}$$

Non zero TEE gives an obstruction to reconstruct  $\rho_{ABC}$  from  $\rho_{AB}$  by acting on B



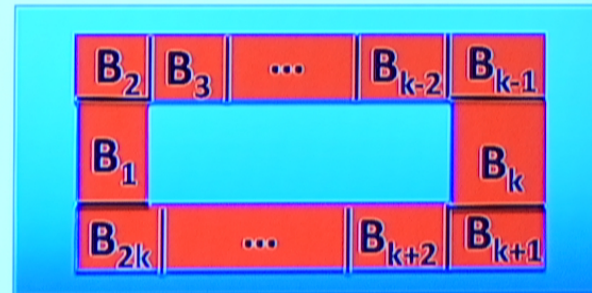
# Why does it work?

We want to show:

$$\begin{aligned} \gamma &\approx \min_{H_{XY}, H_{YZ}} S(\rho_{XYZ} \parallel \exp(H_{XY} + H_{YZ}) / \text{tr}(\dots)) \\ &\approx \min_{H_{B_1 B_2}, \dots, H_{B_{2k-1} B_{2k}}} S(\rho_{B_1 \dots B_{2k}} \parallel \exp(H_{B_1 B_2} + \dots + H_{B_{2k-1} B_{2k}}) / \text{tr}(\dots)) \end{aligned}$$

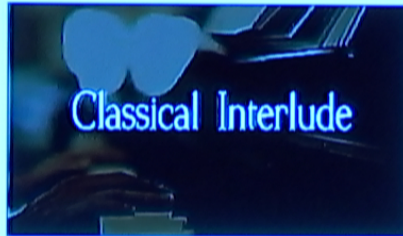
Let's start with the case  $\gamma=0$ .

Need to show  $\rho_{B_1 \dots B_{2k}}$  is close to thermal assuming all conditional mutual information are small, i.e. **approximately independence**



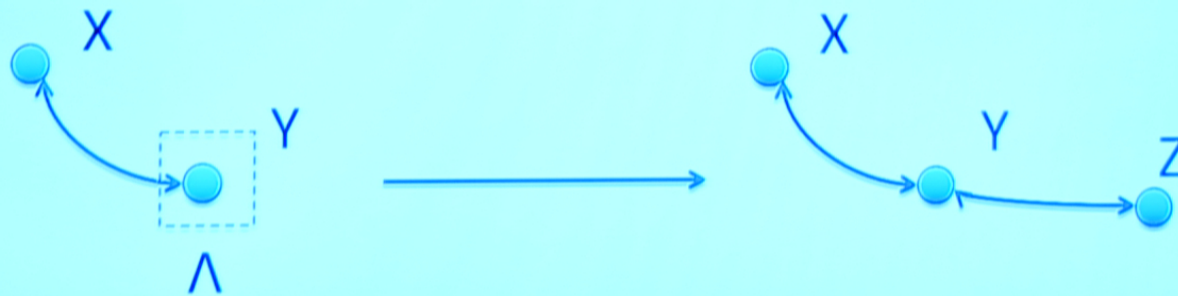
$$I(B_1 \dots B_{j-1} : B_{j+1} \dots B_{2k-1} | B_j B_{2k}) \approx 0$$

# Markov Chain



$X, Y, Z$  with distribution  $p(x, y, z)$

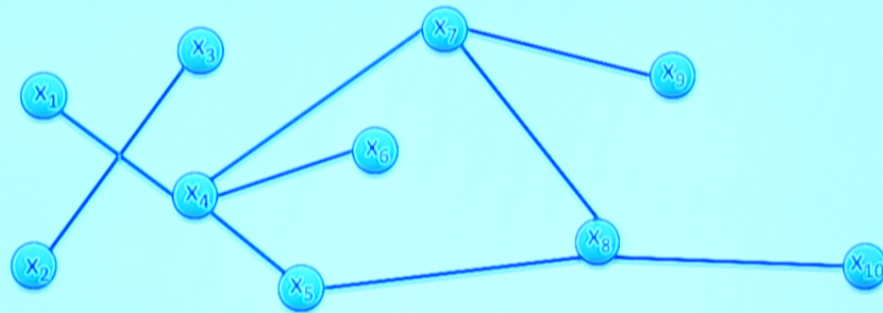
- i)  $X$ - $Y$ - $Z$  Markov if  $X$  and  $Z$  are independent conditioned on  $Y$
- ii)  $X$ - $Y$ - $Z$  Markov if there is a channel  $\Lambda : Y \rightarrow YZ$  s.t.  $\Lambda(p_{XY}) = p_{XYZ}$



$$\text{iii) } I(X : Y|Z)_p = \mathbb{E}_{z \sim p(z)} I(X : Y)_{p(x,y|z=z')}$$



# Markov Networks

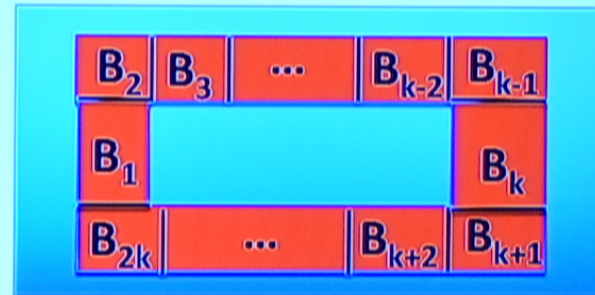


We say  $X_1, \dots, X_n$  on a graph  $G$  form a Markov Network if  $X_i$  is independent of all other  $X$ 's conditioned on its neighbors

Ex: Markov chains 

# Going Back

Need to show  $\rho_{B_1 \dots B_{2k}}$  is close to thermal assuming all conditional mutual information are small (approximately independence)



$$I(B_1 \dots B_{j-1} : B_{j+1} \dots B_{2k-1} | B_j B_{2k}) \approx 0$$

We want a **quantum** and **approximate** version of **Hammersley-Clifford**, but only for 1D chains



# Quantum Markov Chain

Classical:  $X, Y, Z$  with distribution  $p(x, y, z)$

- i)  $X$ - $Y$ - $Z$  Markov if  $X$  and  $Z$  are independent conditioned on  $Y$
- ii)  $X$ - $Y$ - $Z$  Markov if there is a channel  $\Lambda : Y \rightarrow YZ$  s.t.  $\Lambda(p_{XY}) = p_{XYZ}$

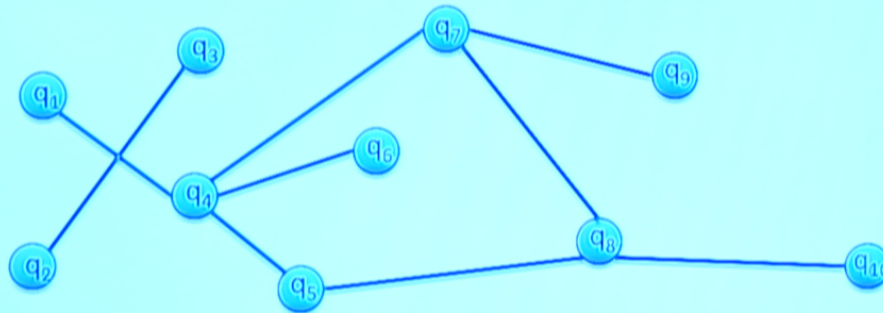
Quantum:

(Hayden, Jozsa, Petz, Winter '03)

- i)  $\rho_{ABC}$  Markov quantum state if  $A$  and  $C$  are "independent conditioned" on  $B$ , i.e.  $H_B \simeq \bigoplus_k H_{B_{L,k}} \otimes H_{B_{R,k}}$  and

$$\rho_{ABC} = \bigoplus_k p_k \rho_{AB_{L,k}} \otimes \rho_{B_{R,k}C}$$

# Quantum Hammersley-Clifford Theorem



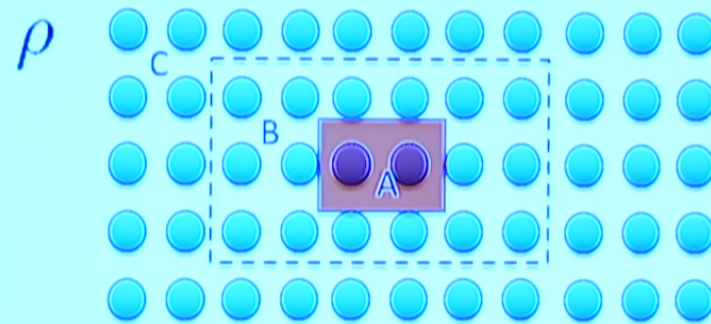
(Leifer, Poulin '08, Brown, Poulin '12) Analogous result holds replacing classical Hamiltonians by *commuting* quantum Hamiltonians

(obs: quantum version more fragile; only works for graphs with no 3-cliques)

Only Gibbs states of commuting Hamiltonians appear. Is there a fully quantum formulation?



# Q. Approximate Markov States



$\rho$  quantum approximate Markov if for every  $A, B, C$

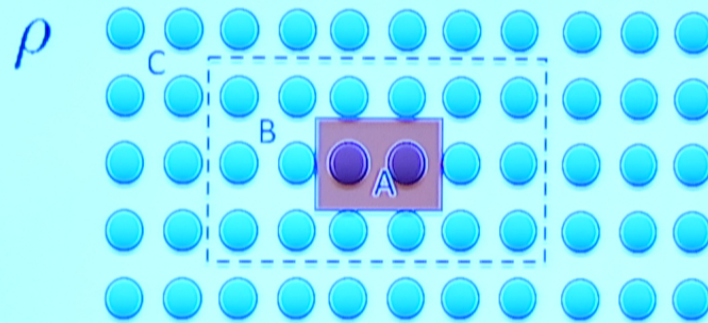
$$I(A : C|B) \rightarrow 0 \text{ when } \text{dist}(A, C) \rightarrow \infty$$

Conjecture

Quantum Approximate Markov  $\iff$  Gibbs state local Hamiltonian

$$\rho = e^{\sum_k H_k}$$

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# Approximate Quantum Markov Chains are Thermal



thm

1. Let  $H$  be a local Hamiltonian on  $n$  qubits. Then

$$I(A : C|B)_{\rho_T} \leq e^{-c'} \sqrt{|B|} + e^{c/T}$$

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2. Let  $\rho_{1\dots n}$  be a state on  $n$  qubits s.t. for every split ABC with  $|B| > m$ ,  $I(A : C|B) \leq \varepsilon$ . Then

$$\min_{H \in \mathcal{H}_{2m}} S(\rho || e^H) \leq \varepsilon \frac{n}{m}$$

$$\mathcal{H}_{2m} := \left\{ H : H = \sum_k H_{k,k+1}, \forall k \text{ supp}(H_{k,k+1}) \leq 2m \right\}$$



## Proof Part 2



Let  $\sigma_{X_1 \dots X_{\frac{n}{m}}}$  be the maximum entropy state s.t.

$$\sigma_{X_i, X_{i+1}} = \rho_{X_i, X_{i+1}} \quad \forall i \in [n/m]$$

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Fact 1 (Jaynes '57):  $\sigma = e^{\sum_k H_{X_k, X_{k+1}}}$

“maximum entropy state given linear constraints is thermal”

$$\operatorname{argmax} (S(\sigma) \text{ s.t. } \operatorname{tr}(\sigma M_i) = c_i) = \exp \left( \sum_i \lambda_i M_i \right)$$



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Fact 1 (Jaynes '57):  $\sigma = e^{\sum_k H_{X_k, X_{k+1}}}$

Fact 2  $\min_{H \in \mathcal{H}_{2m}} S(\rho \| e^H / Z) \leq -S(\rho) - \text{tr}(\rho \log \sigma)$   
 $= S(\sigma) - S(\rho)$

Let's show it's small  $\nearrow$

## Proof Part 2



$$\begin{aligned}
 & S(X_1 \dots X_{n/m})_\sigma \\
 \leq & S(X_1 X_2)_\sigma - S(X_2)_\sigma + S(X_2 \dots X_{n/m})_\sigma \\
 \leq & S(X_1 X_2)_\sigma - S(X_2)_\sigma + S(X_2 X_3)_\sigma - S(X_3)_\sigma + S(X_3 \dots X_{n/m})_\sigma \\
 \leq & \sum_i S(X_i X_{i+1})_\sigma - S(X_{i+1})_\sigma \\
 = & \sum_i S(X_i X_{i+1})_\rho - S(X_{i+1})_\rho
 \end{aligned}$$

Since  $\sigma_{X_i, X_{i+1}} = \rho_{X_i, X_{i+1}} \quad \forall i \in [n/m]$



# Proof Part 1

**Recap:** Let  $H$  be a local Hamiltonian on  $n$  qubits. Then

$$I(A : C|B)_{\rho_T} \leq e^{-c' \sqrt{|B|}} + e^{c/T}$$

We show there is a recovery channel from  $B$  to  $BC$  reconstructing the state on  $ABC$  from its reduction on  $AB$ .

More technical. Uses **Quantum Belief Propagation** equations of Hastings.

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- Locality of EE (area law) implies locality of boundary states and entanglement spectrum
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## Open Questions:

- Applications to high energy/holography?
- Are two copies of entanglement spectrum needed?
- Is the conjecture about approximate Markov chains true?
- Thermal state has same symmetries as original state. Mapping from 2D (zero temperature) to 1D (thermal). Is it useful for classification of (symmetry-protected) phases?