

Title: Holographic complexity by design

Date: Jul 20, 2016 05:00 PM

URL: <http://pirsa.org/16070023>

Abstract:

Holographic complexity

(by design)

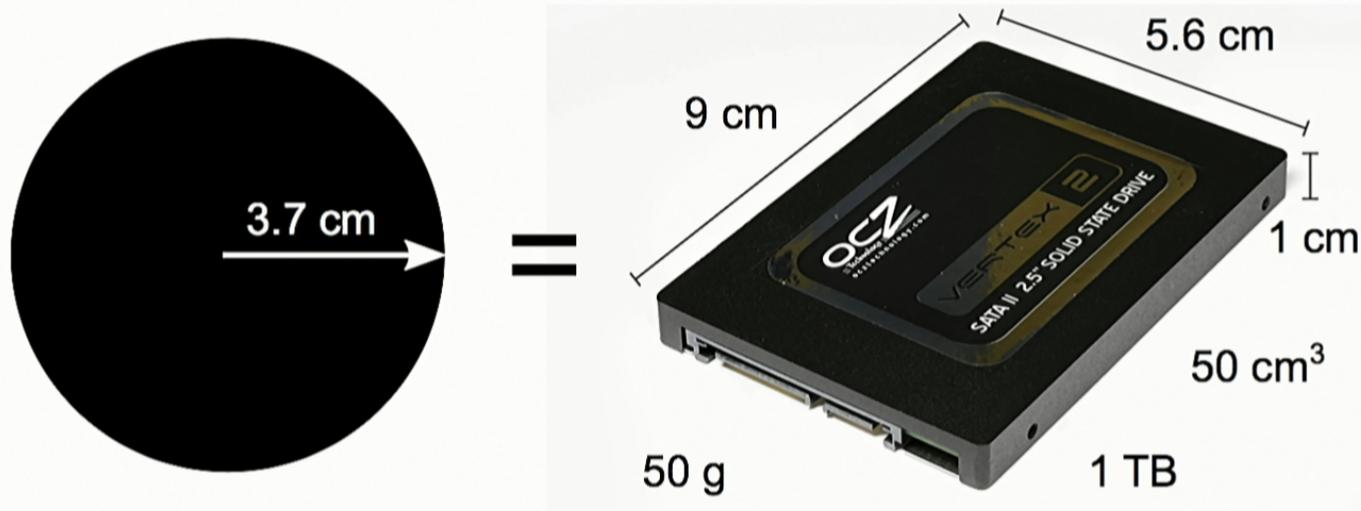
Dan Roberts

The Holmes Building, 632 Massachusetts Ave, Cambridge, MA 02139

July 20, 2016

[arXiv:1509.07876](#) with Adam Brown, Brian Swingle, Leonard Susskind, and Ying Zhao.
[arXiv:1512.04993](#) with Adam Brown, Brian Swingle, Leonard Susskind, and Ying Zhao.
[quant-ph/1608.?????](#) with Beni Yoshida.

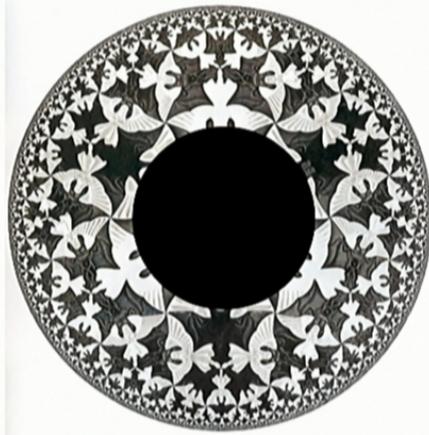
Conference summary: 1975



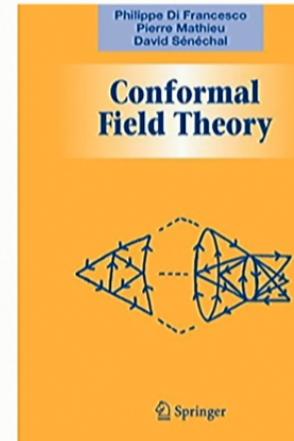
$$R = 2G_N M$$

$$S_{BH} = \frac{A}{4G_N}$$

Conference summary: 1997

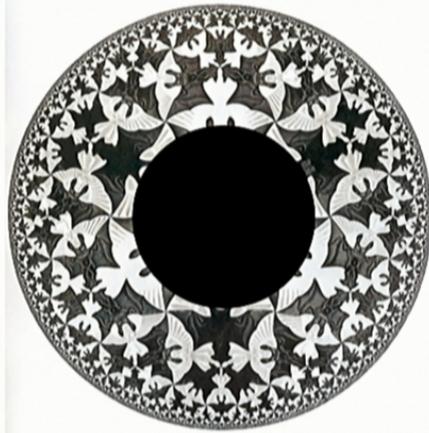


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$$AdS = CFT$$

Conference summary: now

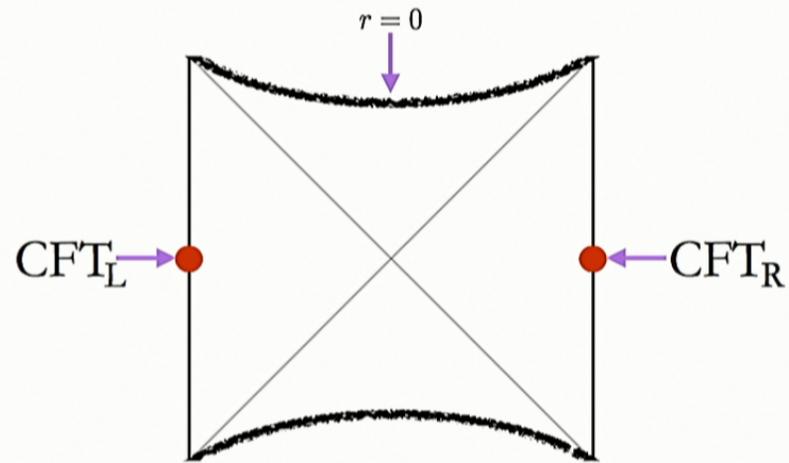


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$$S_X = \min_x \frac{A(x)}{4G_N} \Big|_{\partial X = \partial x}$$

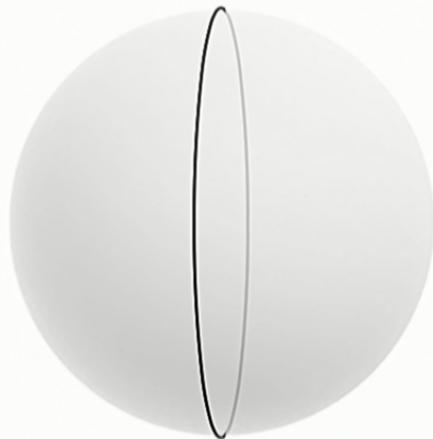
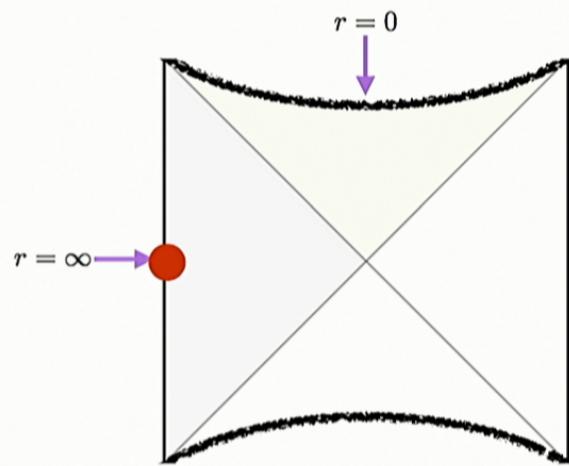
$$S_X = -\text{tr} \{ \rho_x \log \rho_x \}$$

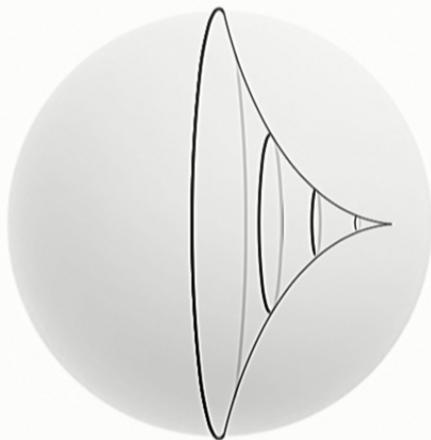
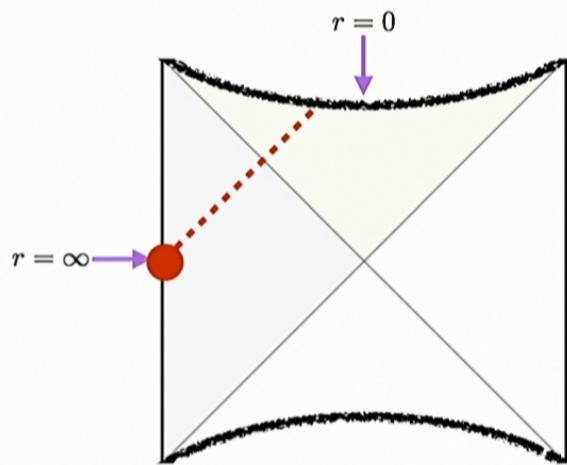


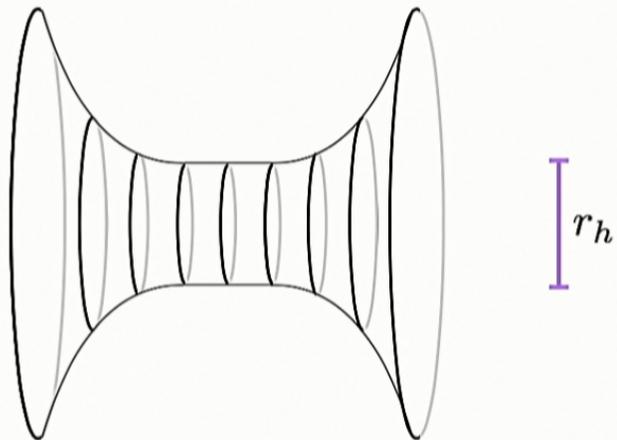
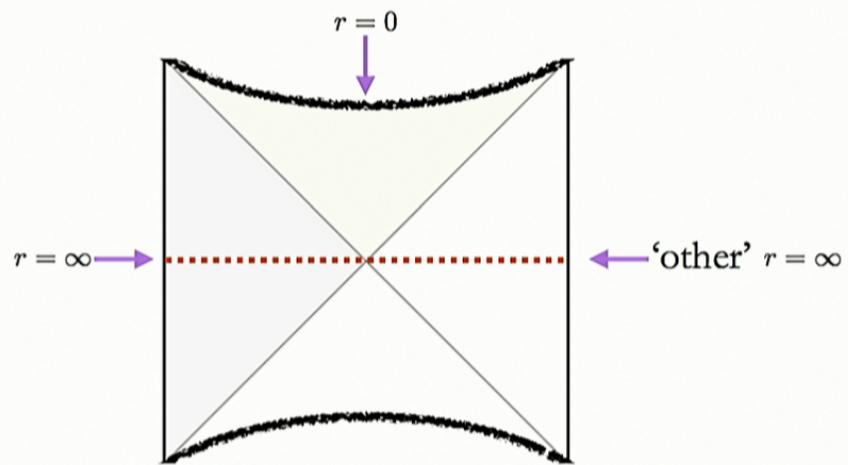
$$|\text{TFD}\rangle = \sum_i e^{-\beta E_i/2} |E_i\rangle_L |E_i\rangle_R$$

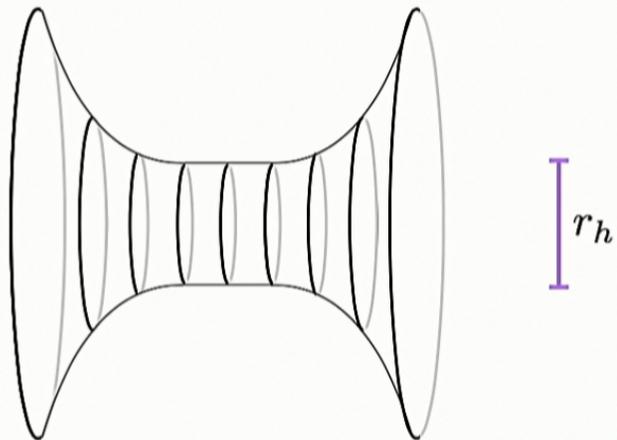
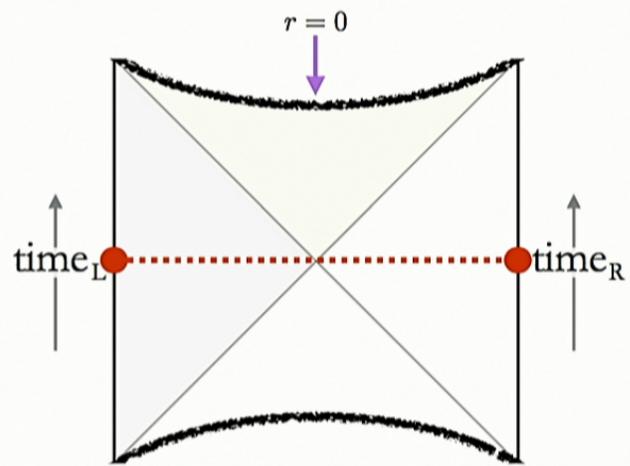
Question:

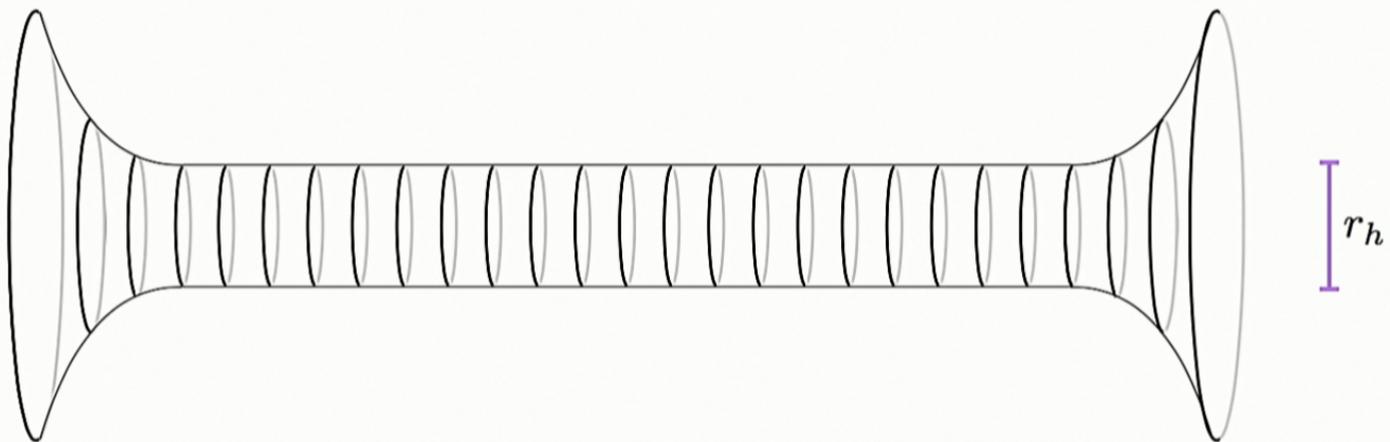
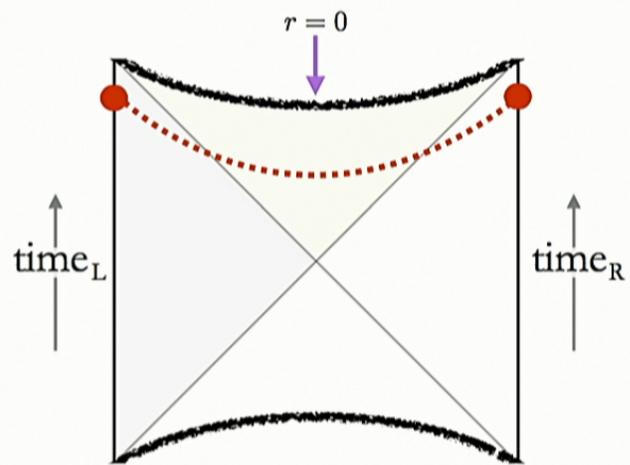
- ▶ What clock ticks for the longest time in AdS?







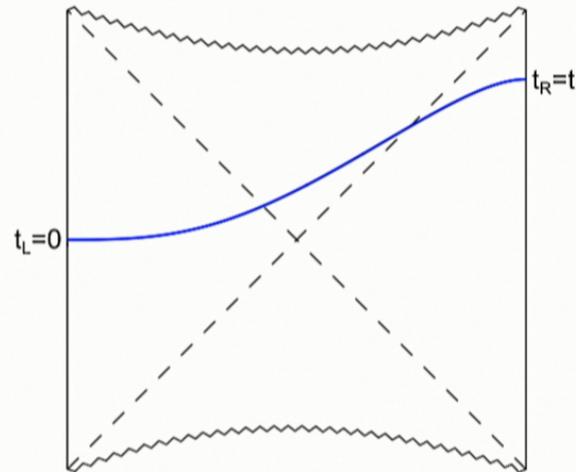




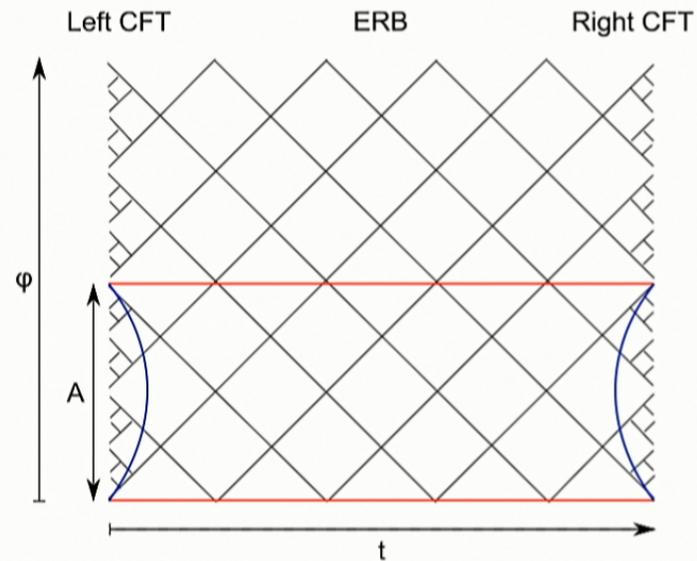
wormhole length $\sim t_L + t_R$

Entanglement entropy?

[Hartman/Maldacena] showed that the entanglement entropy of a subregion of linear size L of the thermofield double state grows linearly with time until saturating at time $O(L)$.



$$\begin{aligned} S_A &\sim t, & t < L, \\ S_A &\sim L, & t > L. \end{aligned}$$



Question:

- ▶ What clock ticks for the longest time in a CFT?

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COMPLEXITY?

computational gate complexity of a quantum state

$$\text{e.g. } |\psi(t_L, t_R)\rangle = \sum_i e^{-\beta E_i/2 + iE_i(t_L + t_R)} |E_i\rangle_L |E_i\rangle_R$$

DEFINITION? starting in a reference state

$$\text{e.g. } |\text{TFD}\rangle = \sum_i e^{-\beta E_i/2} |E_i\rangle_L |E_i\rangle_R$$

how many fundamental gates

e.g. unitaries each of which
act only on two-qubits

are needed to make the state

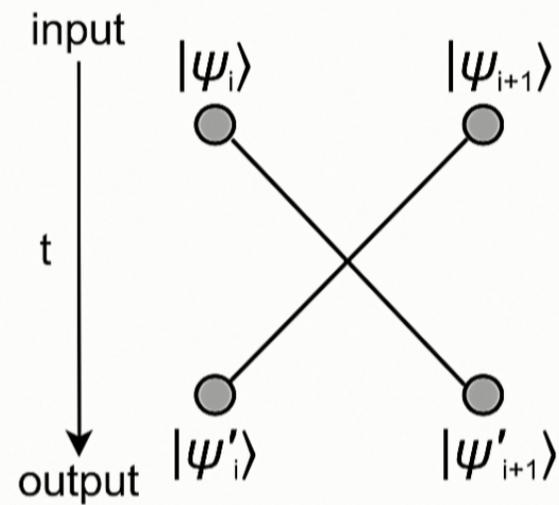
e.g. to within an accuracy ϵ

- can be exponentially large (Hilbert Space is huge)
- expected to grow linearly (at early times)

Complexity of time evolution

Discretize $U(t) = e^{-iHt}$ with a set of 2-body unitaries, or gates

$$U^{(i,i+1)} = \sum_{jk} u_{jk}^{(i,i+1)} |j\rangle\langle k| \approx e^{-ih_{i,i+1}\delta}$$



Complexity of time evolution: $e^{-0iH\delta}$

input



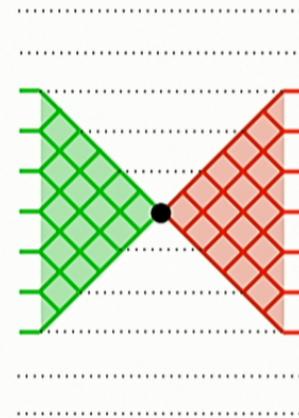
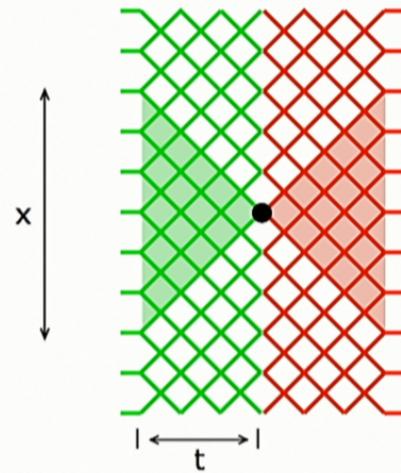
output

complexity \sim size of wormhole?

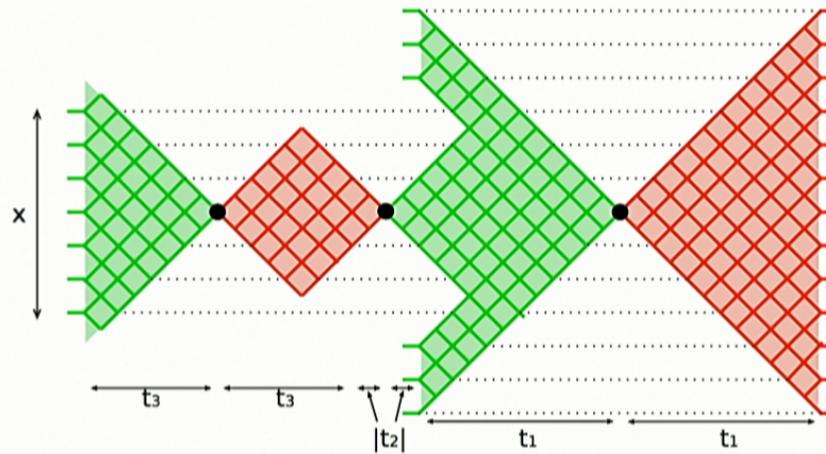
- EVIDENCE:
- both expected to grow linearly (at early times)
 - can be exponentially large (max out at same time)

Complexity means minimal

$$W(t) = e^{-Ht} W e^{Ht}$$



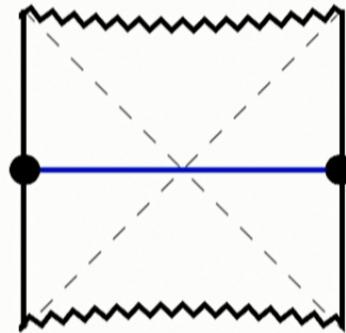
Complexity means minimal



$$W(t_3)W(t_2)W(t_1)$$

Black holes and the butterfly effect

Start with the thermofield double state.

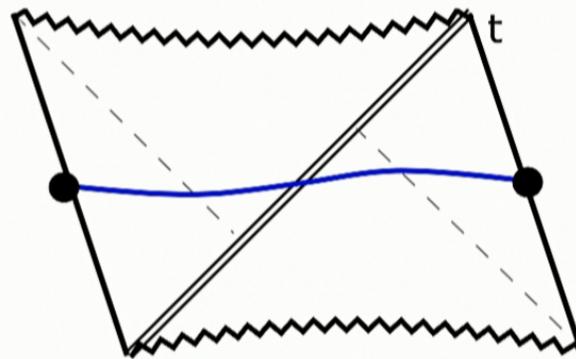


$|TFD\rangle$

[Shenker/Stanford]¹⁰⁰⁰

Black holes and the butterfly effect

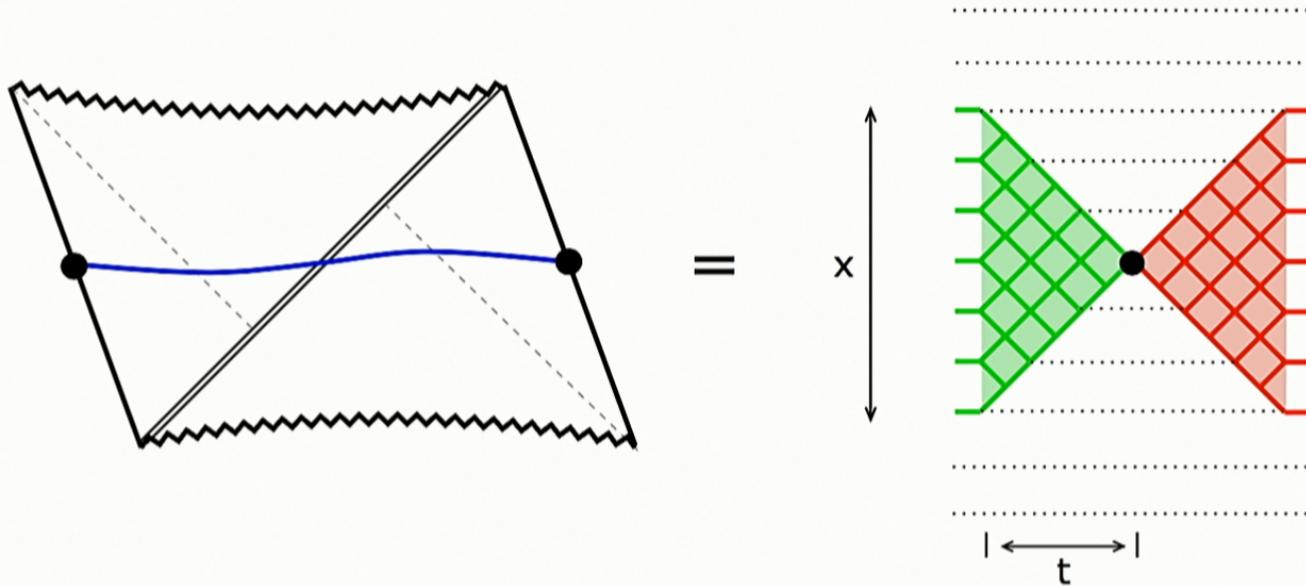
Start with the thermofield double state. Evolve to time t on the right boundary, act with an operator W at position x . Finally, evolve back to time $t_R = 0$.



$$e^{iHt} W e^{-iHt} |TFD\rangle = W(t) |TFD\rangle$$

[Shenker/Stanford]¹⁰⁰⁰

Cancellations in the ERB

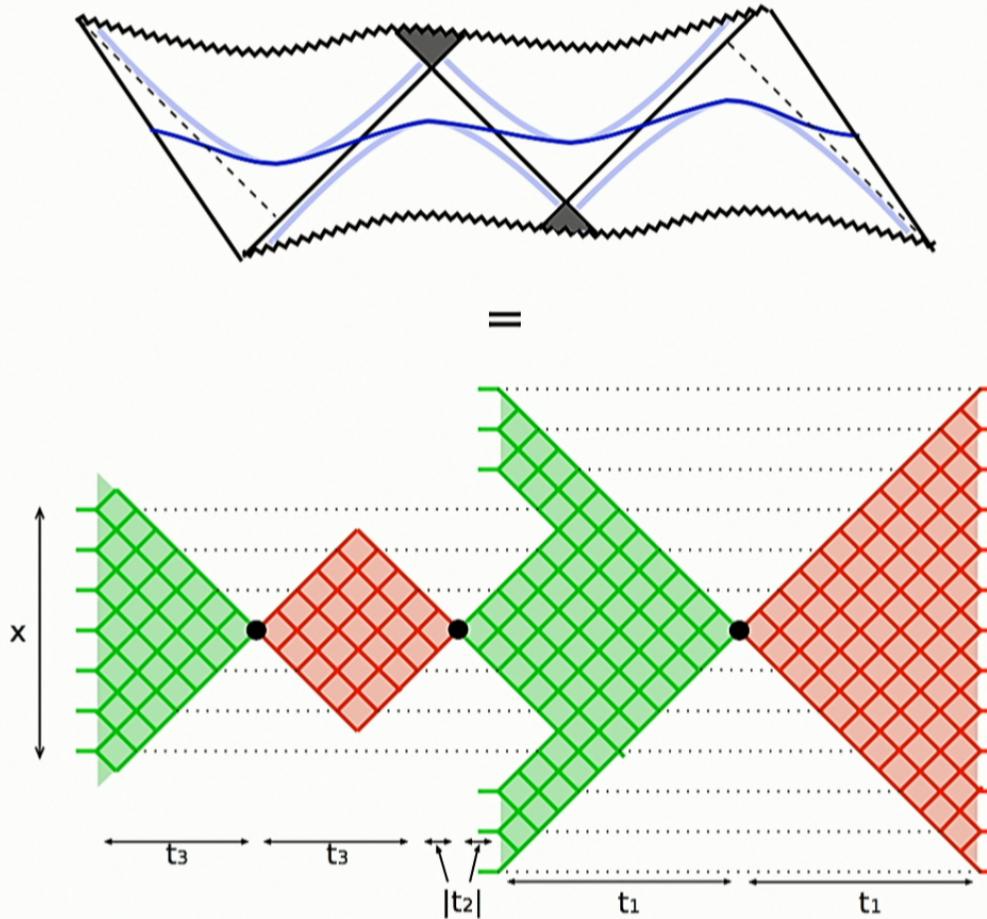


[Stanford/Susskind and DR/Stanford/Susskind]

complexity \sim size of wormhole?

- EVIDENCE:
- both expected to grow linearly (at early times)
 - can be exponentially large (max out at same time)
 - perturb black hole with boundary operator
 - in CFT, leads to increase in complexity (chaos)
 - in wormhole, leads to shockwave, increases size
- the two increases match
including delicate cancellations!
- single/multiple/localized perturbations

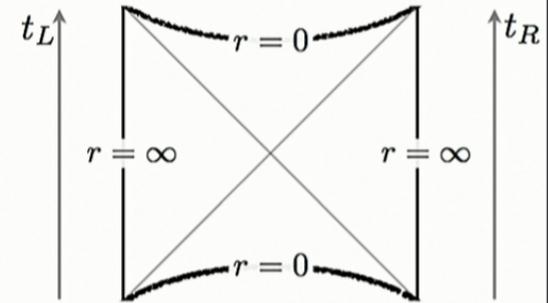
Multiple cancellations in the ERB



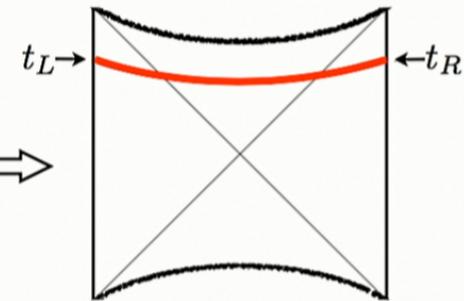
how to characterize size of wormhole?

$$\frac{1}{G\ell_{\text{AdS}}} \text{volume}$$

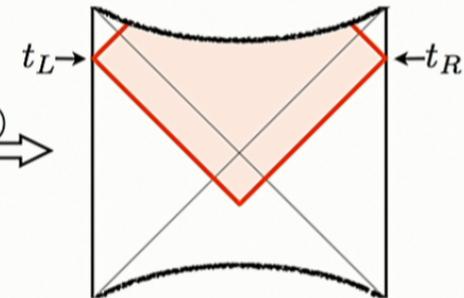
$$\frac{1}{\hbar} \text{action}$$



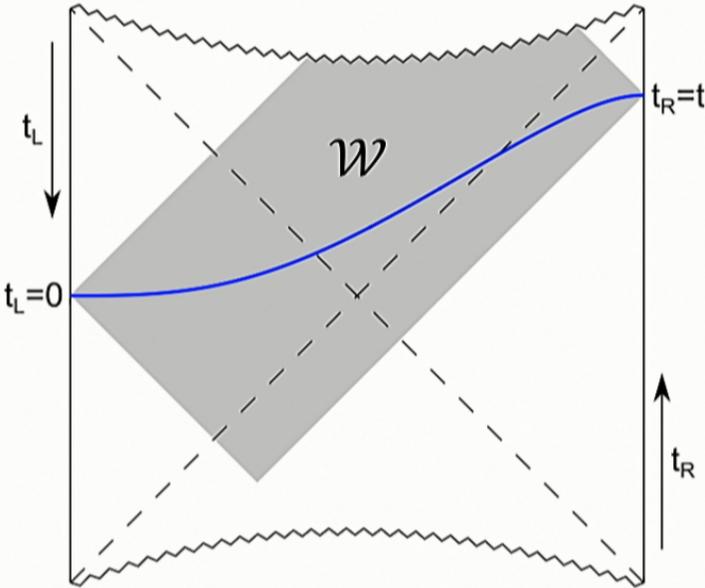
(old)
maximal slice \Rightarrow



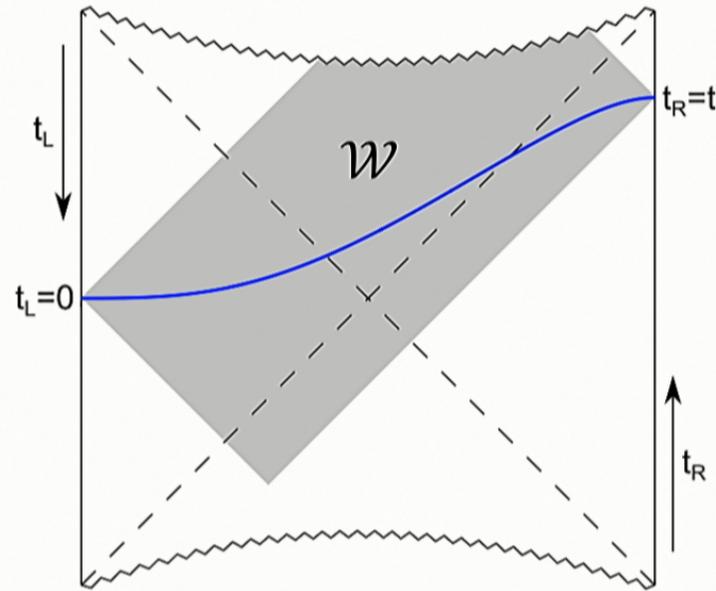
(new/improved)
WdW patch \Rightarrow



Wheeler-DeWitt patch

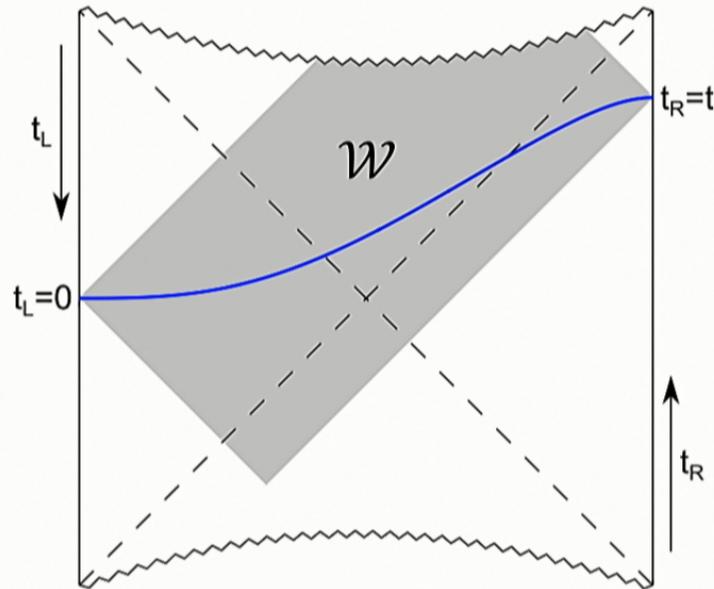


Conjecture



$$\text{Complexity} = \frac{\text{Action}(\mathcal{W})}{\pi \hbar}$$

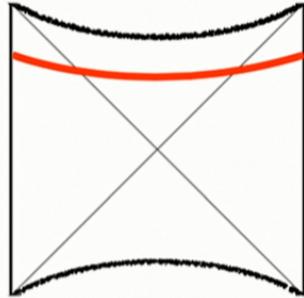
(Another) conjecture



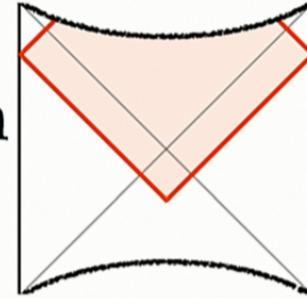
$$\frac{d\text{Action}(\mathcal{W})}{dt} \leq 2M$$

(Neutral) black holes are fastest computers in nature.

$\frac{1}{G l_{\text{AdS}}}$ **volume**



vs. $\frac{1}{\hbar}$ **action**



prefer action because:

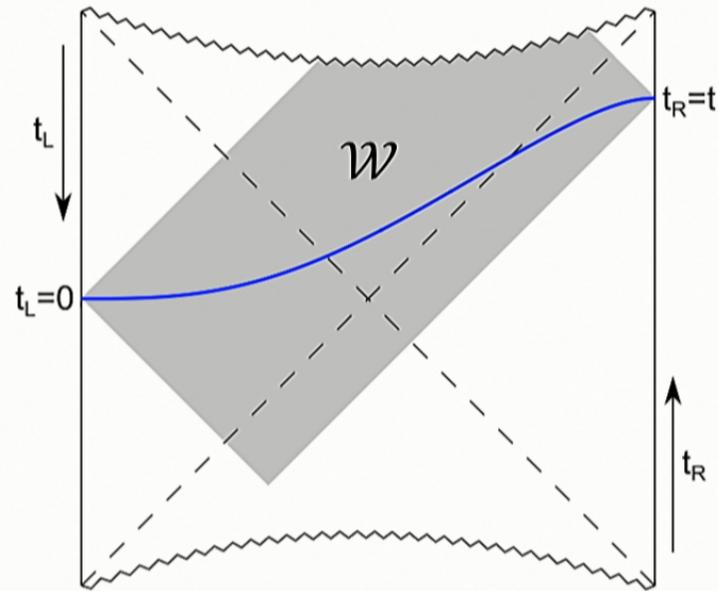
- naturally “dimensionless” (no arbitrary length scale)
- $d(\text{action})/dt = 2M$ (independent of size and dimension)

connect to $\frac{d\text{Complexity}}{dt} \leq \frac{2M}{\pi\hbar}$

black holes saturate limit implies

$$\text{Complexity} = \frac{\text{Action}}{\pi\hbar}$$

Wheeler-DeWitt patch



$$\text{Action}(\mathcal{W}) = \frac{1}{16\pi G_N} \int_{\mathcal{W}} \sqrt{-g} (\mathcal{R} - 2\Lambda) + \frac{1}{8\pi G_N} \int_{\partial\mathcal{W}} \sqrt{|h|} K,$$

Complexity Equals Action ?

FURTHER WORK:

- **precise definition of complexity?**
- precise definition of action?
- relate imprecision in two definitions?
- reference state? (“complexity of formation”)
- classical proof that black holes maximize action?
- more general black holes?
- higher-derivative theories and singularities?
- principle of least computation?
- complexity and horizon transparency?
- lots of puzzles!

Complexity lower bounds?

How many circuits can we make with g gates acting on n qubits with Δ steps?

There are $g \binom{n}{2}$ choices to make at each step, therefore we can make at most

$$\# \text{ circuits} = (gn^2)^\Delta.$$

If we have a collection of different circuits we want to make \mathcal{E} , then if we want to ensure we can make all of them, we can determine the minimal number of steps

$$\Delta \geq \frac{\log |\mathcal{E}|}{\log(gn^2)}.$$

How random?

For an ensemble \mathcal{E} , we can characterize how generic it is by using the frame potential

$$F_{\mathcal{E}}^{(k)} = \frac{1}{|\mathcal{E}|^2} \sum_{U, V \in \mathcal{E}} |\text{tr}\{V^\dagger U\}|^{2k}.$$

The frame potential is maximized by $\mathcal{E} = I$ and minimized by the Haar ensemble

$$F_{Haar}^{(k)} < F_{\mathcal{E}}^{(k)} < F_I^{(k)}.$$

If for an ensemble $F_{Haar}^{(k)} < F_{\mathcal{E}}^{(k)}$, then that ensemble is a k -design. It will reproduce k moments of the Haar ensemble.

Complexity and chaos

For any ensemble \mathcal{E} of unitary operators, with $\tilde{B} = U^\dagger B U$

$$\frac{1}{2^{4kn}} \sum_{A_1, \dots, B_1, \dots} \left| \langle A_1 \tilde{B}_1 \cdots A_k \tilde{B}_k \rangle_{\mathcal{E}} \right|^2 = \frac{1}{2^{2(k+1)n}} \cdot F_{\mathcal{E}}^{(k)}$$

where summations are over all possible Pauli operators.

If we assume H is generic, then our ensemble could be of Hamiltonians. A disorder average of out-of-time-order correlators would give the frame potential. If the system self-averages, then we might use this to lower-bound the complexity for a particular system.