

Title: A new perspective on holographic entanglement

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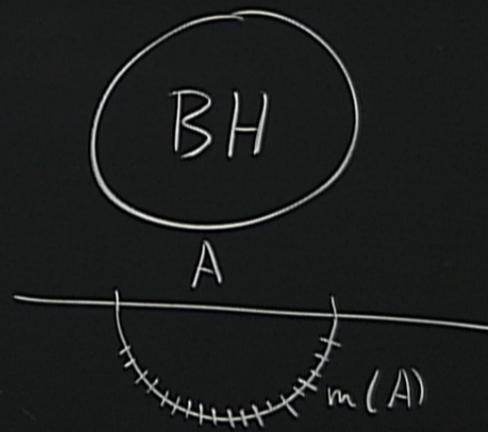
Abstract: We will present a reformulation of the Ryu-Takayanagi holographic entanglement entropy formula which does not involve the areas of surfaces. The reformulation leads to a picture of entanglement entropy of boundary regions being carried by Planck-thickness "bit threads" in the bulk. We will argue that this picture resolves a number of conceptual difficulties surrounding the RT formula.

A new perspective on holographic entanglement
(w/ M. Freedman, to appear)

In gravity, area = entropy

$$S_{\text{BH}} = \frac{1}{4G_N} \text{area}(\text{horizon})$$

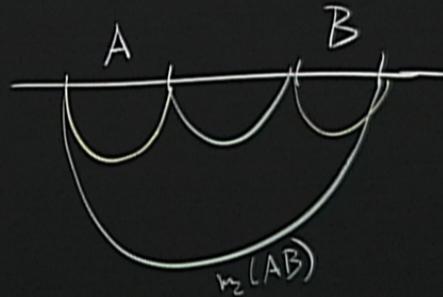
$$S(A) = \frac{1}{4G_N} \text{area}(m(A))$$



Do the bits "live on" the surface?

Puzzles:

1) $m(A)$ can jump



$S(AB)$

$$m_1(AB) = m(A) \cup m(B)$$

$$m_2(AB) \neq \quad "$$

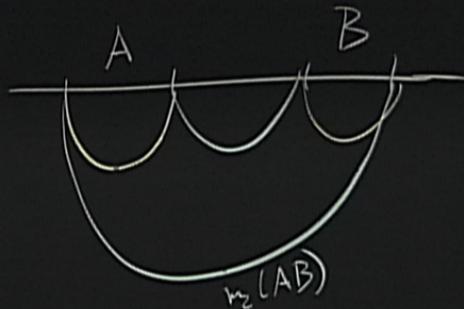
2)

CAUTION
For safety do not touch the electrical wiring
inside panels in this building or any other
building.
If you experience an electrical problem
please contact the building manager.
THANK YOU FOR YOUR COOPERATION.

Do the bits "live on" the surface?

Puzzles:

1) $m(A)$ can jump



$$S(AB)$$

$$m_1(AB) = m(A) \cup m(B)$$

$$m_2(AB) \neq \text{" "}$$

2) MI: $I(A:B) = S(A) + S(B) - S(AB)$

Info-theory meaning of MI $S(A)$ $H(B|A) = S(AB) - S(A)$

classical: A  B

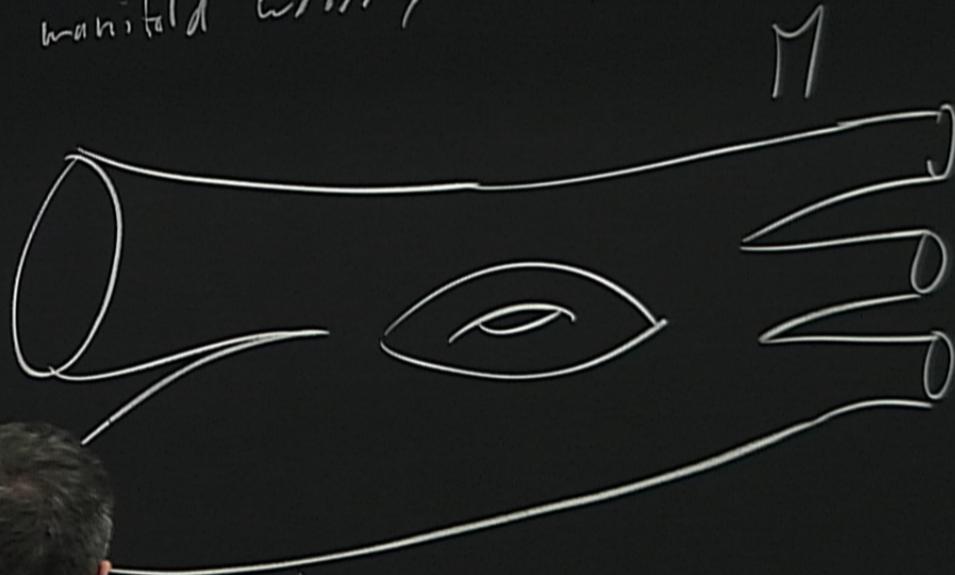
$I(A:B)$

$$H(A|B) = S(A) - S(B)$$

cond. entropy $S(B)$

Min-cut / max-flow

Riemannian manifold w/ h/ky M

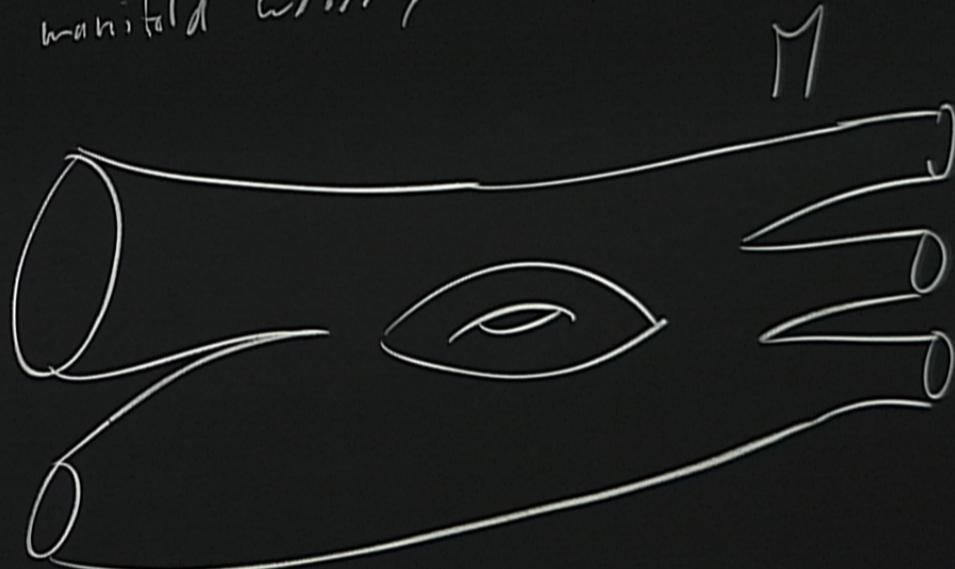


Flow: vector field v $|v| \leq 1$

Min-cut / max-flow

(HPY 15)

Riemannian manifold w/ hky M



Flow: vector field v :
1) $|v| \leq 1$
2) $D_x v^i = 0$

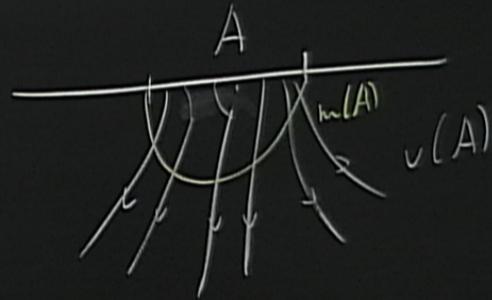
Thm: $\max_v \int_A v = \min_{m \sim A} \text{area}(m)$

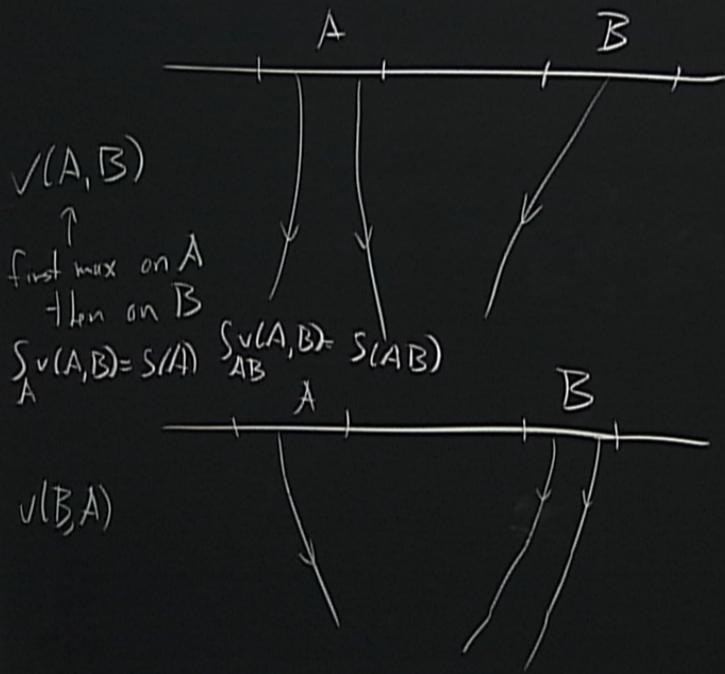
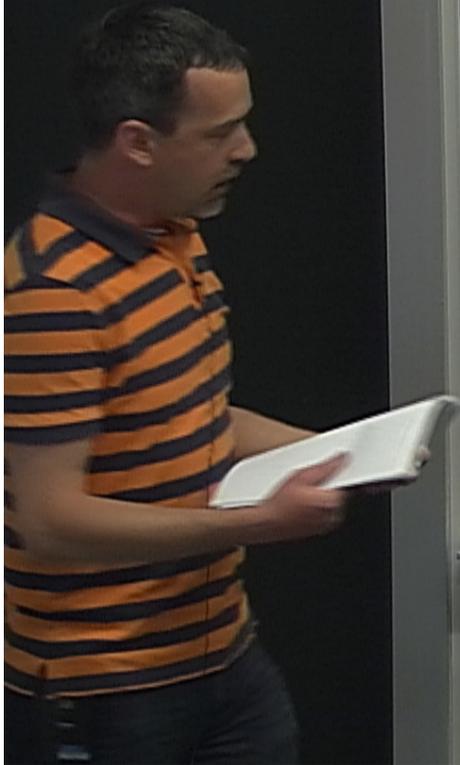
On $m(A)$, $v(A) = n/m(A)$
 \uparrow max flow

elsewhere, freedom in choosing $v(A)$

RT 2.0: $S(A) = \frac{1}{4G_N} \max_v \int_A v$

Threads have cross-section $4G_N$ threads leaving A





$V(A, B)$

↑
first max on A
first min on B

$$\int_A V(A, B) = S(A)$$

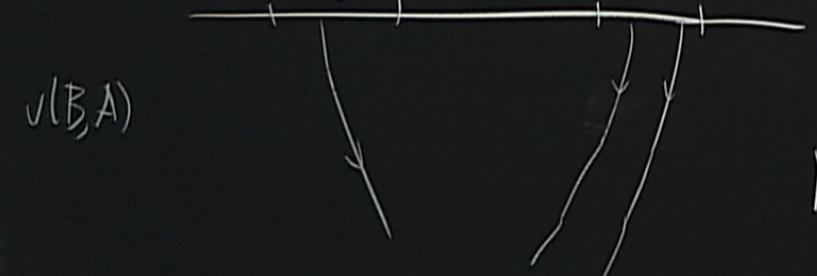
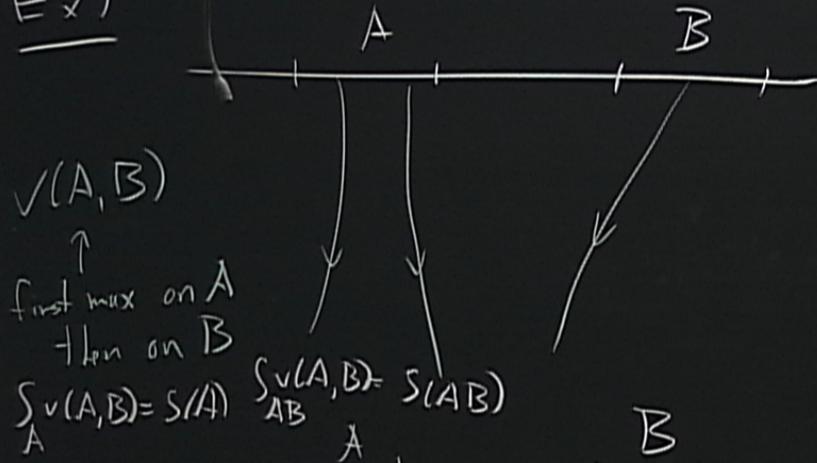
$$\int_{AB} V(A, B) = S(AB)$$

$V(B, A)$

$$S(A) = S(B) = 2$$

$$S(AB) = 3$$

Ex 1



$$S(A) = S(B) = 2$$
$$S(AB) = 3$$

threads free to move between A, B
 $= I(A:B)$

threads stuck on A = $H(A|B)$

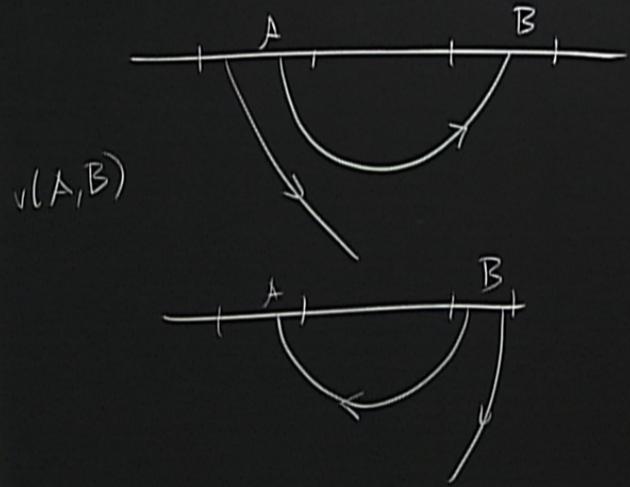
$$H(A|B) = S(AB) - S(B)$$
$$= \int_{AB} v(B,A) - \int_B v(B,A)$$
$$= \int_A v(B,A)$$

A

Info-theory meaning of $M(L) S(A)$ $H(B|A) = S(AB) - S(A)$



Ex 2. $S(A) = S(B) = 2$ $S(AB) = 1$



between A, B

$H(A|B)$

$v(B, A)$

Flow v

Thm

RT 20

Th

CAUTION

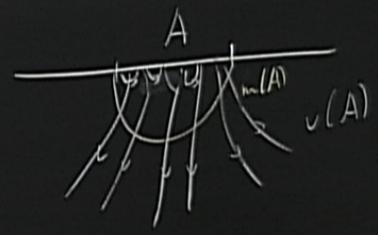
CAUTION

Flow vector field v . 1) $|v| \leq 1$
 2) $D_x v^x = 0$

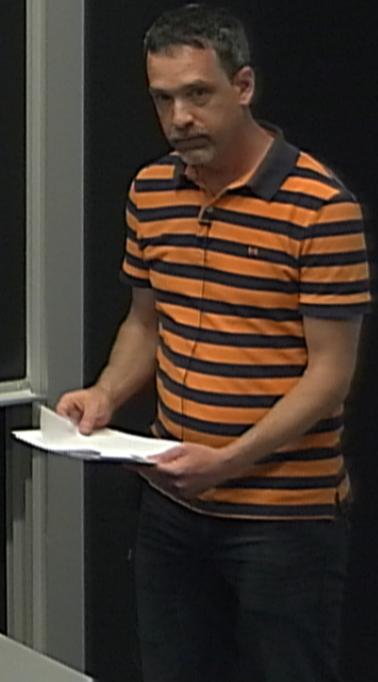
$$\int_A v = \int_{m \sim A} v$$

Thm: $\max_v \int_A v = \min_{m \sim A} \text{area}(m)$

On $m(A)$, $v(A) = n|_{m(A)}$
 ↑ max flow
 elsewhere, freedom in choosing $v(A)$

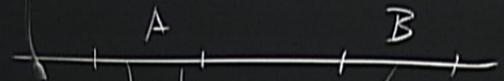


RT 2.0: $S(A) = \frac{1}{4G_N} \max_v \int_A v$
 = max # threads leaving A
 Threads have cross-section $4G_N$



CAUTION

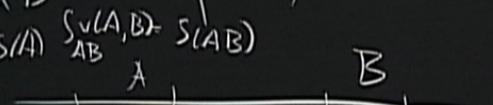
Ex 1



$$S(A) = S(B) = 2$$

$$S(AB) = 3$$

$v(A, B)$
 ↑
 first max on A
 then on B
 $\int_A v(A, B) = S(A)$



$v(B, A)$

threads free to move between A, B
 $= I(A, B)$

threads stuck on A = $H(A|B)$

$$H(A|B) = S(AB) - S(B)$$

$$= \int_{AB} v(B, A) - \int_B v(B, A)$$

$$= \int_A v(B, A)$$



≥ 0

Conditional MI

$$I(A; B | C) = S(A|C) + S(B|C) - S(C) - S(ABC)$$
$$= \int_A v(C, A, B) - \int_A v(C, B, A)$$

≥ 0

≥ 0

Conditional MI

$$I(A;B|C) = S(A|C) + S(B|C) - S(C) - S(ABC)$$
$$= \int_A v(C,A,B) - \int_A v(C,B,A)$$

≥ 0

CAUTION

$v(B, A)$

$$I(A:B) = H(A) - H(A|B) \\ = \int_A v(A, B) - \int_A v(B, A) \geq 0$$



between A, B
 $H(A|B)$
 $v(B, A)$

Condition
 $I(A:B)$

Questions

1) Tripartite info $I_3(A:B:C)$
Superadditivity
Monogamy of MI

$$I_3 \leq 0$$

2) What kinds of entanglement + correlation can be represented by threads?

3) Higher-curvature corrections?

4) Quantum "