

Title: Gauge color codes

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URL: <http://pirsa.org/14070012>

Abstract: I will describe a new class of topological quantum error correcting codes with surprising features. The construction is based on color codes: it preserves their unusual transversality properties but removes important drawbacks. In 3D, the new codes allow the effective transversal implementation of a universal set of gates by gauge fixing, while error-detecting measurements involve only 4 or 6 qubits. Furthermore, they do not require multiple rounds of error detection to achieve fault-tolerance.



# Outline

- Introduction
- Transversal gates in color codes
- Gauge color codes
- Single-shot fault-tolerant error correction



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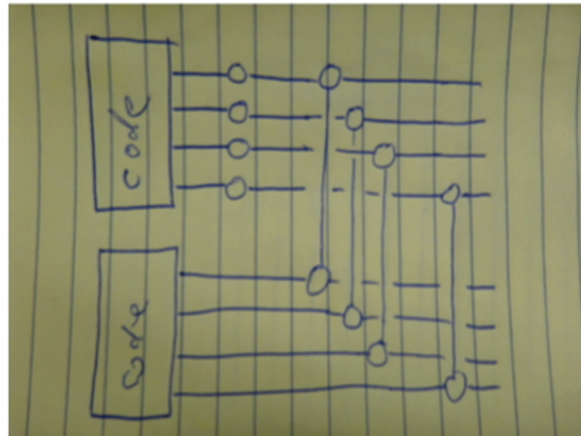
Local codes



Local computation

# Transversal gates

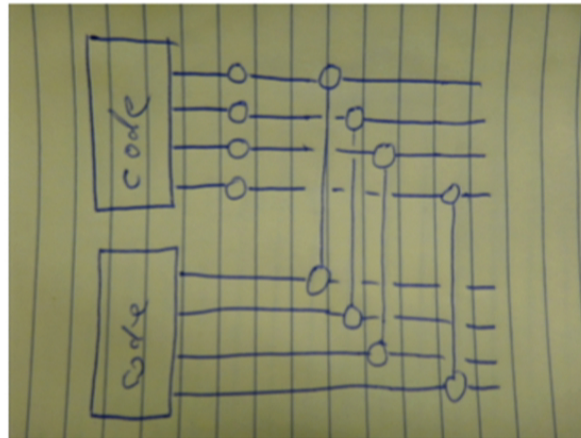
- Aim: do not spread errors
- How: act on subsystems separately





# Transversal gates

- Aim: do not spread errors
- How: act on subsystems separately



# Transversal gates

No code admits a universal transversal set of gates

Eastin & Knill '09

Alternatives:

- Magic state distillation Bravyi & Kitaev '03
- Gauge fixing Paetznick & Reichardt '13
- Concatenation Jochym-O'Connor & Laflamme '14

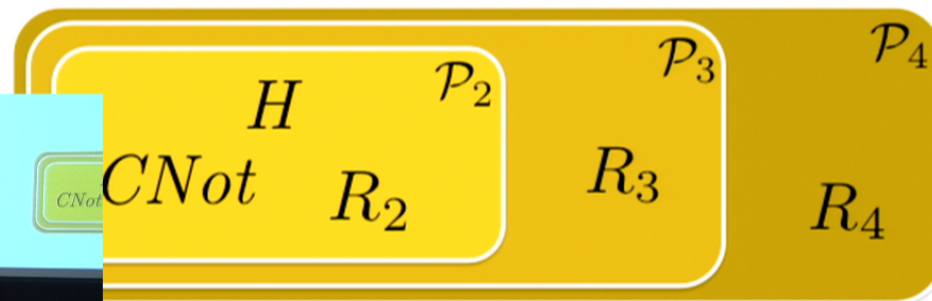


# Dimensional restrictions

Bravyi & Koenig '13

$$\mathcal{P}_D := \{U \mid U\mathcal{P}U^\dagger \subseteq \mathcal{P}_{D-1}\}, \quad \mathcal{P}_1 := \mathcal{P}$$

Gottesman & Chuang '99



$$R_D := \begin{pmatrix} 1 & 0 \\ 0 & e^{2\pi i/2^D} \end{pmatrix}$$



# Effectively local operations

- Local quantum op + global classical computation
- Caution: can yield non-local errors!

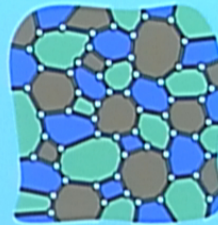


# Effectively local operations

- Defines a class of codes with error correction/ initialization local in this sense.
- Contains abelian anyons, but not non-abelian ones!



## Color codes

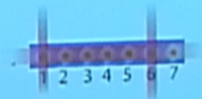


### Topology + transversality !

- 2D: Clifford group  $\rightarrow$  Distillation
- $D > 2$ : Cnot +  $R_D := \begin{pmatrix} 1 & 0 \\ 0 & e^{2\pi i/2^D} \end{pmatrix}$
- 3D  $\rightarrow$  Hadamard via ancilla  $|+\rangle$

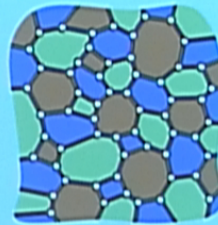
Demonstrated with 7 trapped ions

Nigg et al '14





## Color codes



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# Color codes

## Difficulties:

- Transversal gates conditional on the existence of suitable lattices
- 3D: many-body measurements of 20+ qubits
- Ancilla encoded qubit for  $H$  gate



## Gauge Color codes

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**NOT**

- Transversal gates conditional on the existence of suitable lattices
- 3D: many-body measurements of ~~20+~~ **just 6** qubits
- **NO** Ancilla encoded qubit for  $H$  gate
- Effectively local (**single-shot**) FT QEC
- Constant time overhead



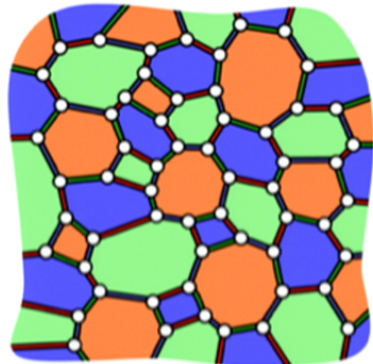
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## 2D color codes

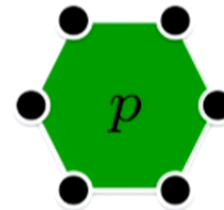


- 2-colexes
- 1 qubit per vertex
- 2 stabilizer generators per plaquette

$$\mathcal{S} = \langle X_p, Z_p \rangle_p$$

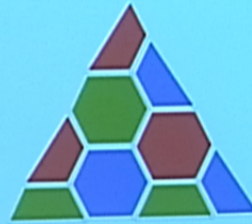
$$X_p = \bigotimes_{i \in p} X_i$$

$$Z_p = \bigotimes_{i \in p} Z_i$$





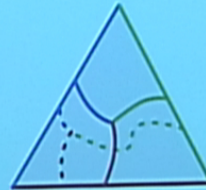
## 2D color codes



Triangular geometry:

- 1 encoded qubit
- Logical operators:  $\hat{X}, \hat{Z}$

$$\hat{O} := O^{\otimes \text{all}}$$



String-like logical operators

Must connect the 3 triangle edges



## 2D color codes

### Trasversal gates

- $H$  and CNot work: self-dual CSS codes
- $R_2$  is more tricky

$$R_2^\dagger Z R_2 = Z \quad R_2^\dagger X R_2 = i Z X$$

$$\hat{R}_2^\dagger \hat{Z} \hat{R}_2 = \hat{Z} \quad \hat{R}_2^\dagger \hat{X} \hat{R}_2 = \pm i \hat{Z} \hat{X} \quad \checkmark$$

$$\hat{R}_2^\dagger Z_p \hat{R}_2 = Z_p \quad \hat{R}_2^\dagger X_p \hat{R}_2 = (-1)^{q/2} Z_p X_p \quad \times$$

$q = \text{\#qubits in plaquette } p$

## 2D color codes



Instead, inverse gate on some qubits!

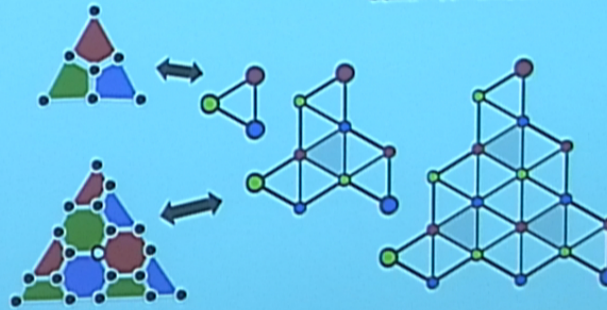
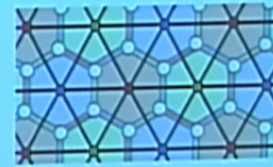
$$R_2 \quad R_2^\dagger = Z R_2$$

● ○



## 2D color codes

Dual picture:  
simplicial lattice





## Color codes

- Simplicial  $(D+1)$ -colored lattice
- Geometry = colored  $D$ -simplex
- Qubits =  $D$ -simplices
- Logical ops =  $\hat{X}, \hat{Z}$

Parameter:

$$d = 1, \dots, D - 1$$

$$\bar{d} := D - d$$

**Dimensions**

	$x$	$z$
$\mathcal{S}$	$d - 1$	$\bar{d} - 1$
errors / logical ops	$\bar{d}$	$d$

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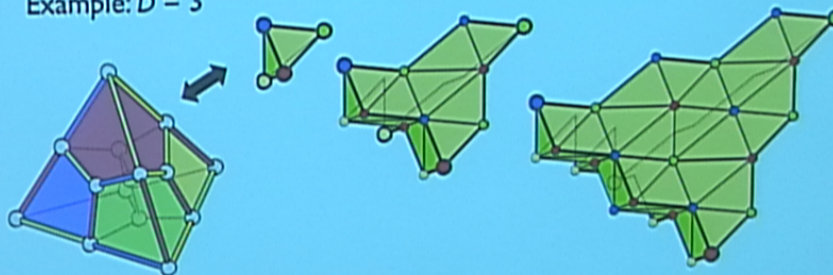
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$\mathcal{S}$	$d - 1$	$\bar{d} - 1$
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# Color codes

Example:  $D = 3$

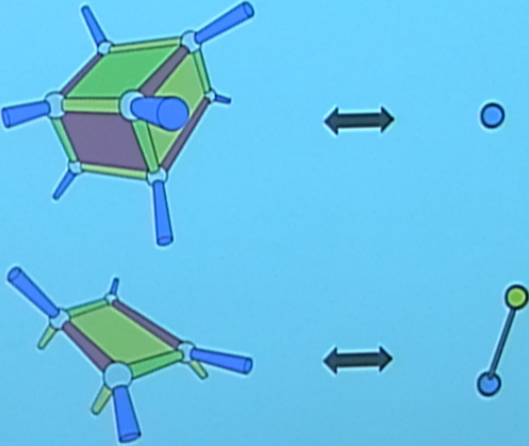


	$\mathcal{S}_X$	$\mathcal{S}_Z$	
$d = 1$			→ 4 or 6 qubits
$d = 2$			→ up to 24 qubits



# Color codes

Example:  $D = 3$





## Transversality trick

CSS code with logical operators  $\hat{X}, \hat{Z}$

$R_n$  via transversal  $R_n^k$  for some  $k$  if

$$\begin{aligned} |S_i| &\equiv 0 \pmod{2^n} \\ |S_i \cap S_j| &\equiv 0 \pmod{2^{n-1}} \\ |S_i \cap S_j \cap S_k| &\equiv 0 \pmod{2^{n-2}} \\ &\dots \end{aligned}$$

The  $S_i$  are the supports of X stabilizer generators



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## Transversality trick

$$\begin{array}{cc} R_n^k & R_n^{-k} \\ \bullet & \circ \end{array}$$

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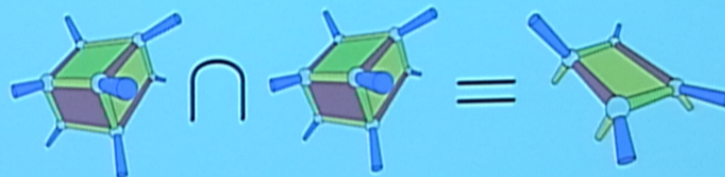
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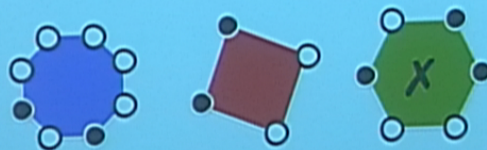


## Transversality trick

For color codes intersections are constrained by coloring



Choose set as support of error for  
'bad plaquette' syndrome

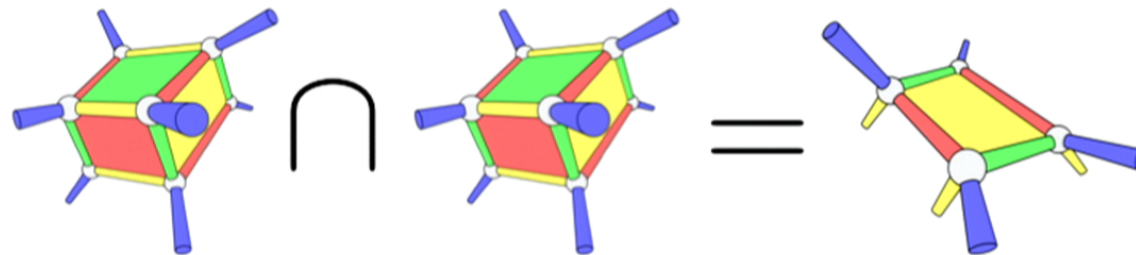


**Only constraint:**

$$D \geq nd$$

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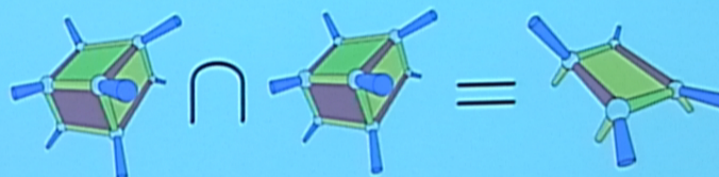
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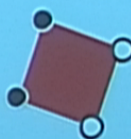


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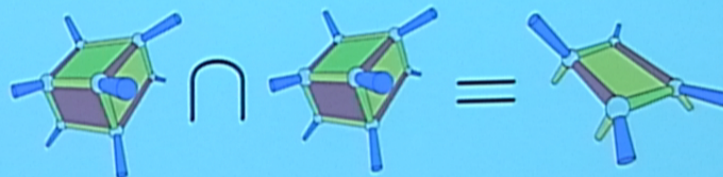
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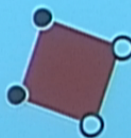


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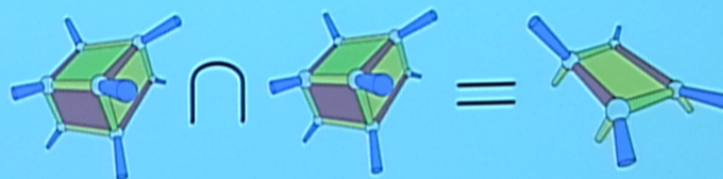
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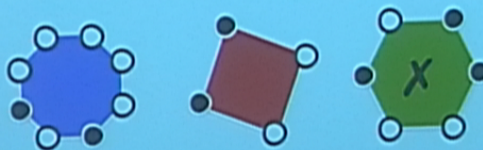


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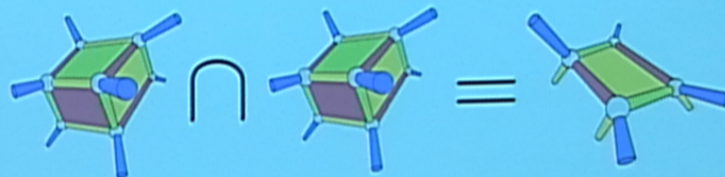
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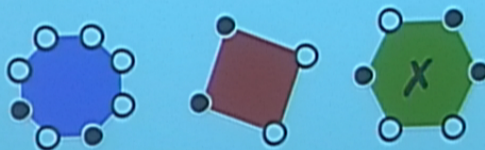


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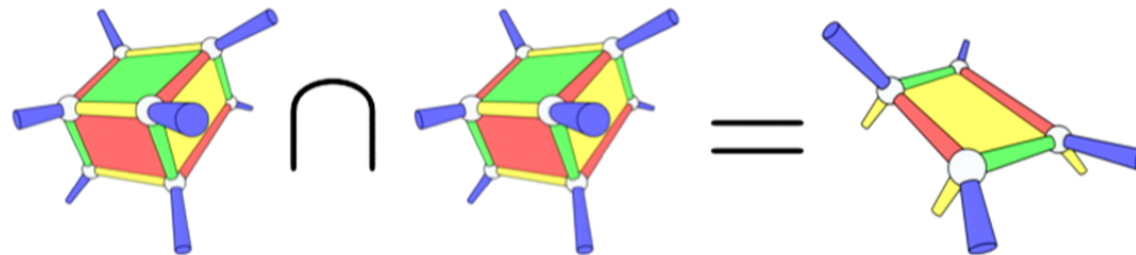


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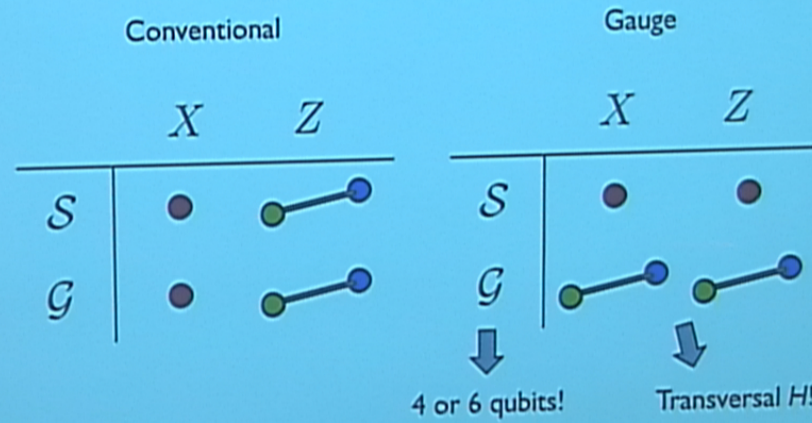
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# Gauge color codes

Example:  $D = 3$

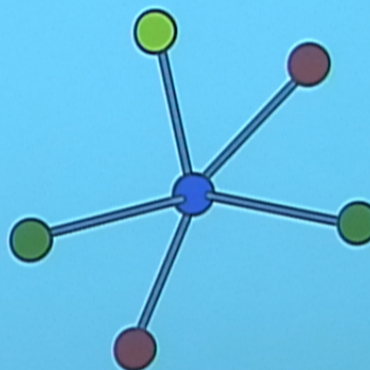




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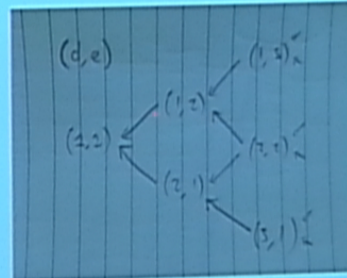
Geometrical relationship between 'syndromes'





## Gauge color codes

- Partial order:  $d_1 \leq d_2, e_1 \leq e_2 \quad \mathcal{S}_1 \subseteq \mathcal{S}_2 \quad \mathcal{G}_1 \supseteq \mathcal{G}_2$



- To move with the arrows, do nothing
- To move opposite, fix the gauge: measure  $\mathcal{S}_2$ , apply  $\mathcal{G}_1$
- Requires a non-local classical computation!

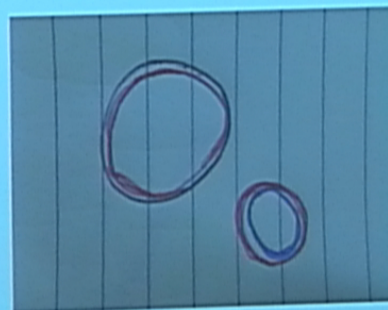


## Gauge color codes

Example:  $D = 3$

For an encoded state we get a gauge syndrome with loops only

Plaquette operators suffice to bring this to a 'vacuum' state





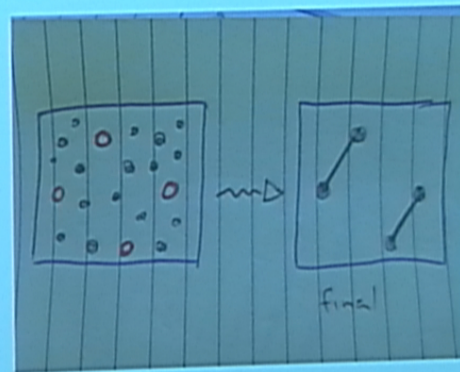
## Setting

- The applied correction is 'inmaterial' up to stabilizers & logical Z operators
- Assume *only* syndrome readout fails, with probability  $p$  independent of reading.
- We can attach *any* Z error with that syndrome to that probability  $p$ , and get an effective error model for the whole process
- *Locality*:  $n$  given failures happen with probability  $< \lambda^n$
- Will the effective model preserve locality?



## Surface codes

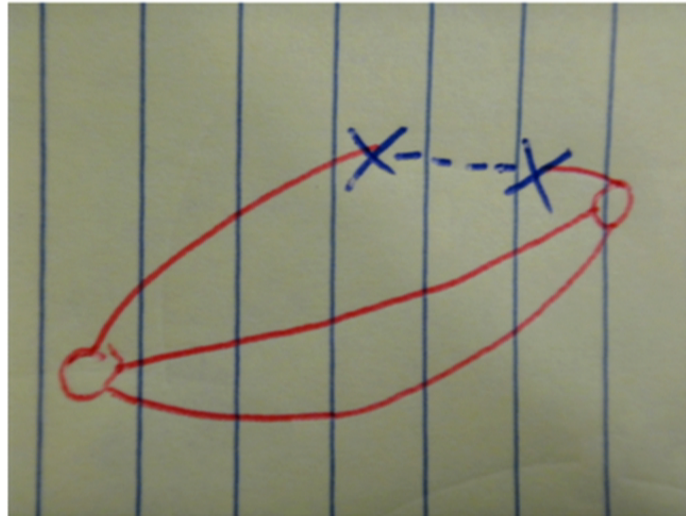
- Large errors induced by a few ghosts!





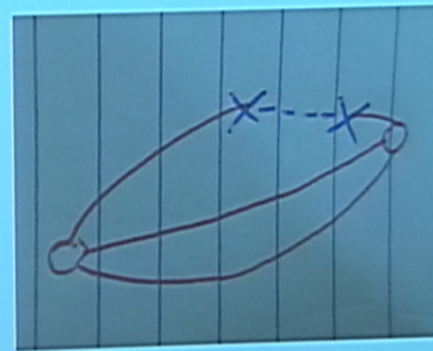
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- Dealing with non-valid measurements is easy



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## Summary & discussion

- 3D gauge color codes have **many** interesting practical properties
- Error thresholds?
- What are the limitations in 2D?
- What about non-geometrical locality?
- Could there be related 3D self-correcting systems?

$$H = - \sum_{g \in \mathcal{G}_0} J_g g$$