

Title: Exact Classical Simulation of the Quantum-Mechanical GHZ Distribution

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Abstract: John Bell has shown that the correlations entailed by quantum mechanics cannot be reproduced by a classical process involving non-communicating parties. But can they be simulated with the help of bounded communication? This problem has been studied for more than twenty years and it is now well understood in the case of bipartite entanglement. However, the issue was still widely open for multipartite entanglement, even for the simplest case, which is the tripartite Greenberger-Horne-Zeilinger (GHZ) state. We give an exact simulation of arbitrary independent von Neumann measurements on general n -partite GHZ states. Our protocol requires $O(n^2)$ bits of expected communication between the parties, and $O(n \log n)$ expected time is sufficient to carry it out in parallel. Furthermore, we need only an expectation of $O(n)$ independent unbiased random bits, with no need for the generation of continuous real random variables nor prior shared random variables. In the case of equatorial measurements, we improve earlier results with a protocol that needs only $O(n \log n)$ bits of communication and $O(\log^2 n)$ parallel time. At the cost of a slight increase in the number of bits communicated, these tasks can be accomplished with a constant expected number of rounds.

Exact simulation of the GHZ distribution with finite expected communication

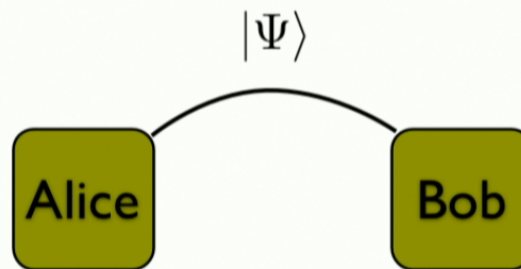
Gilles Brassard
Université de Montréal

Luc Devroye
McGill University

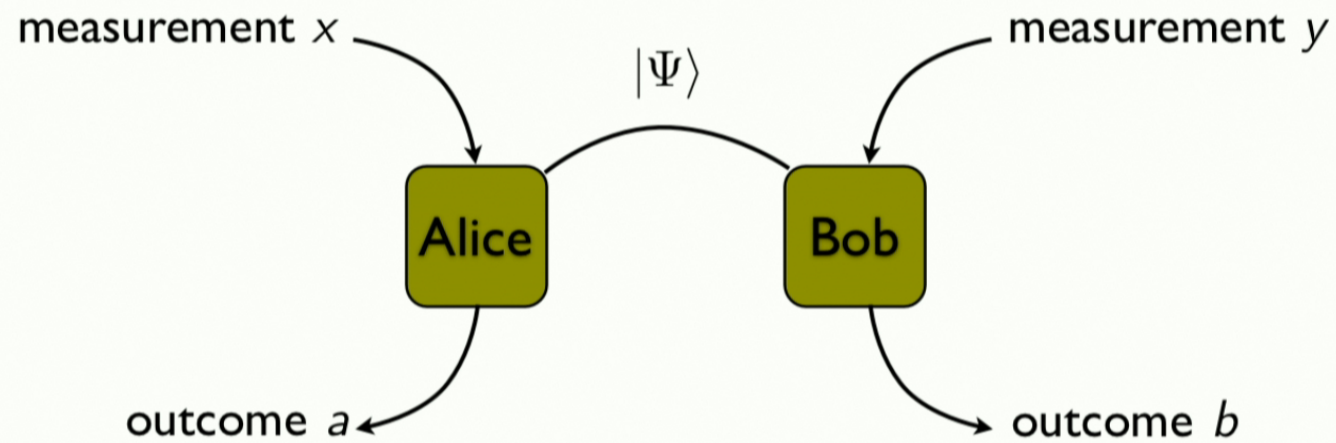
Claude Gravel
Université de Montréal



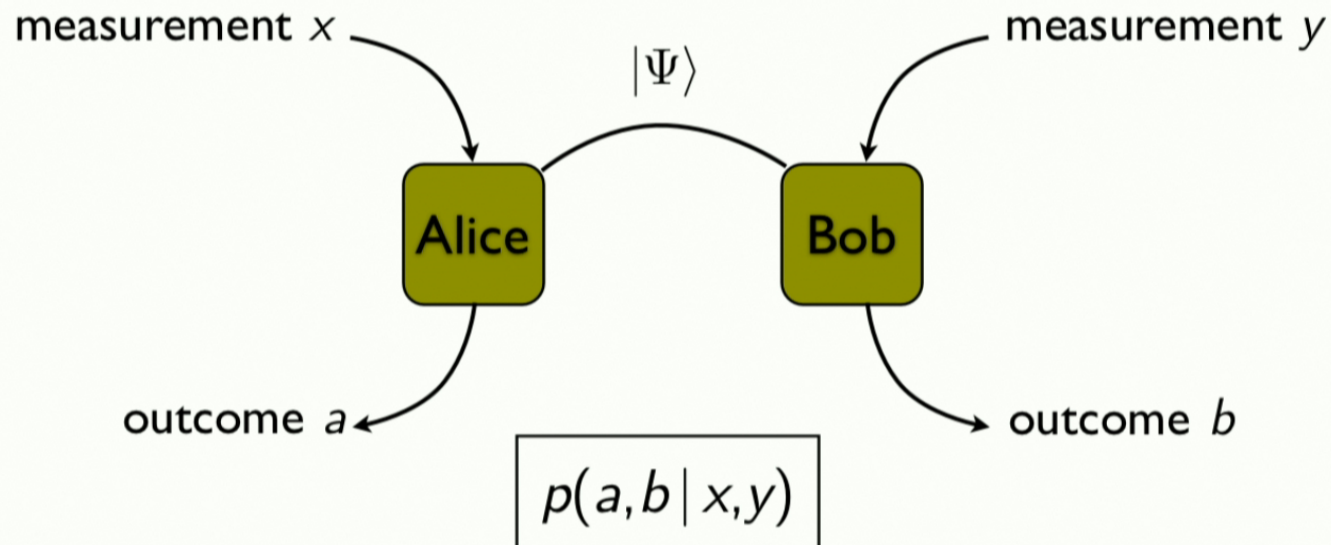
Quantum nonlocality



Quantum nonlocality



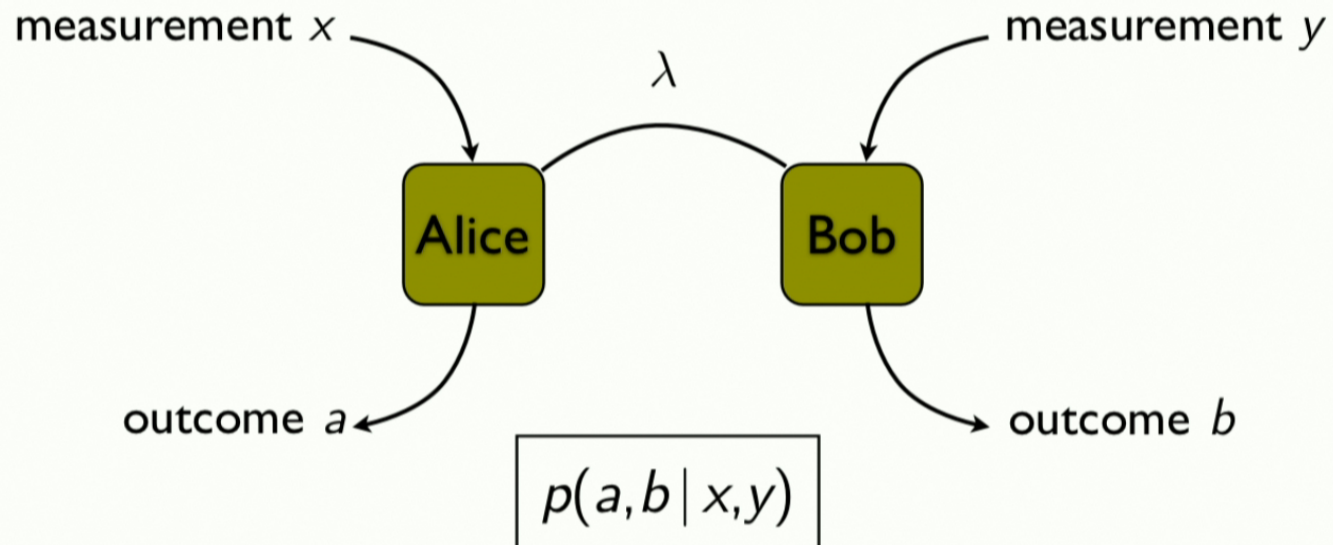
Quantum nonlocality



Can the theory be supplemented with local hidden variables (LHV)?

Can quantum-mechanical description of physical reality be considered complete?
Einstein, Podolsky and Rosen, 1935

Quantum nonlocality



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Einstein-Podolsky-Rosen



Albert Einstein



Boris Podolsky



Nathan Rosen

Entanglement



I would not call that one but rather the characteristic trait of quantum mechanics, the one that enforces its entire departure from classical lines of thought.

Erwin Schrödinger, 1935

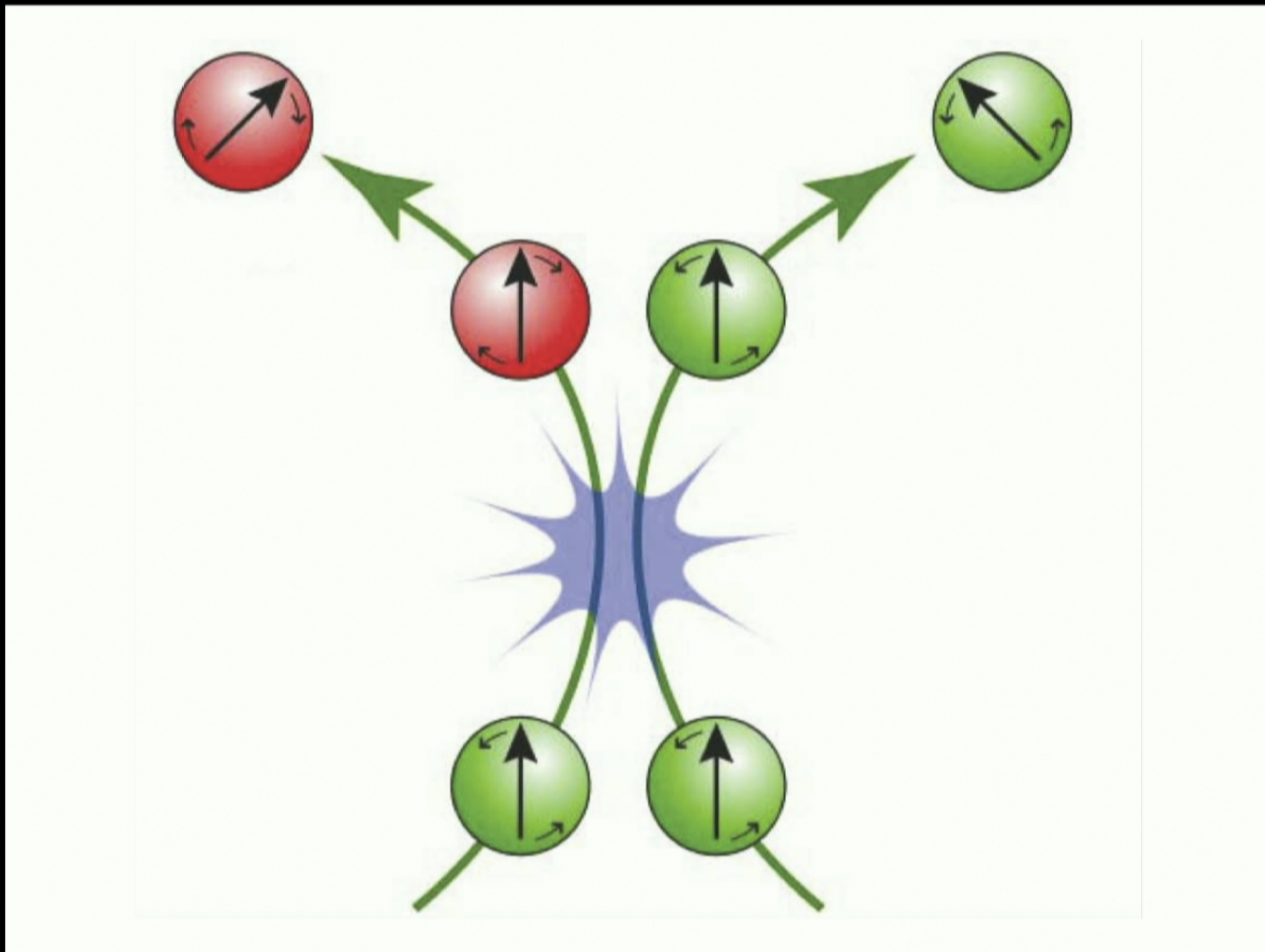
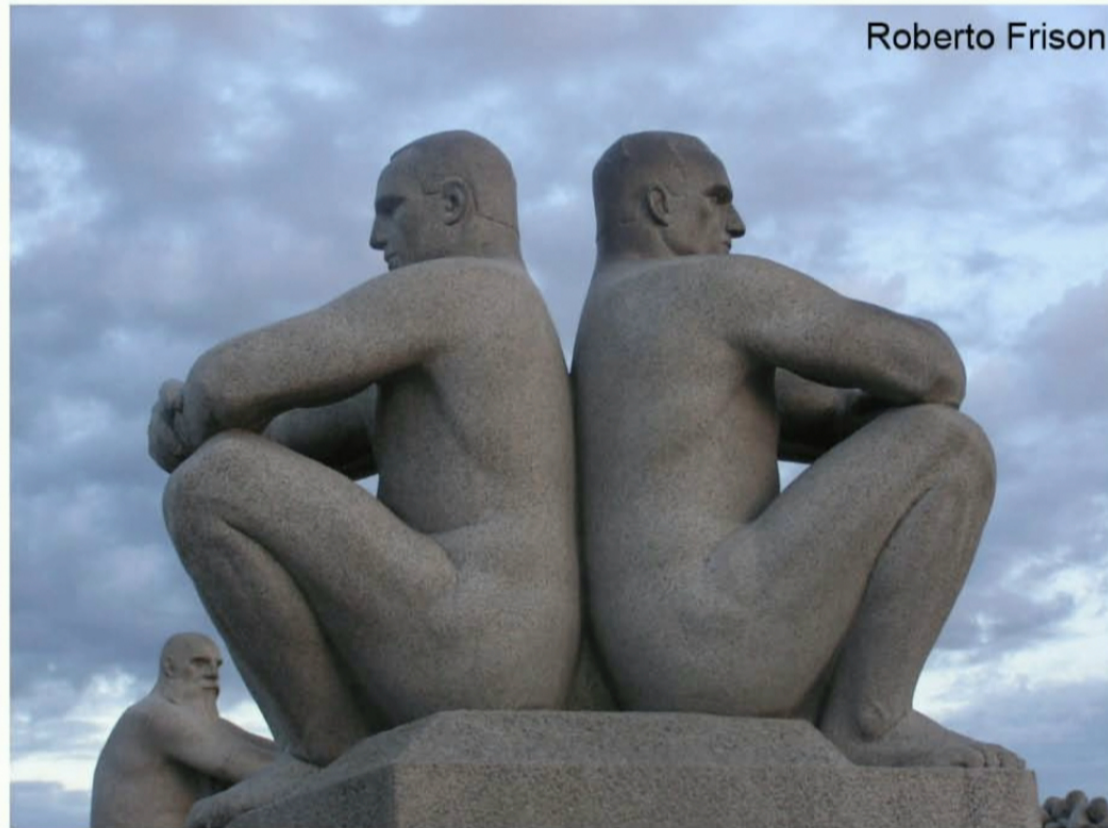


Illustration of Entanglement



Vigeland Park, Oslo, Norway

Do you believe in God?

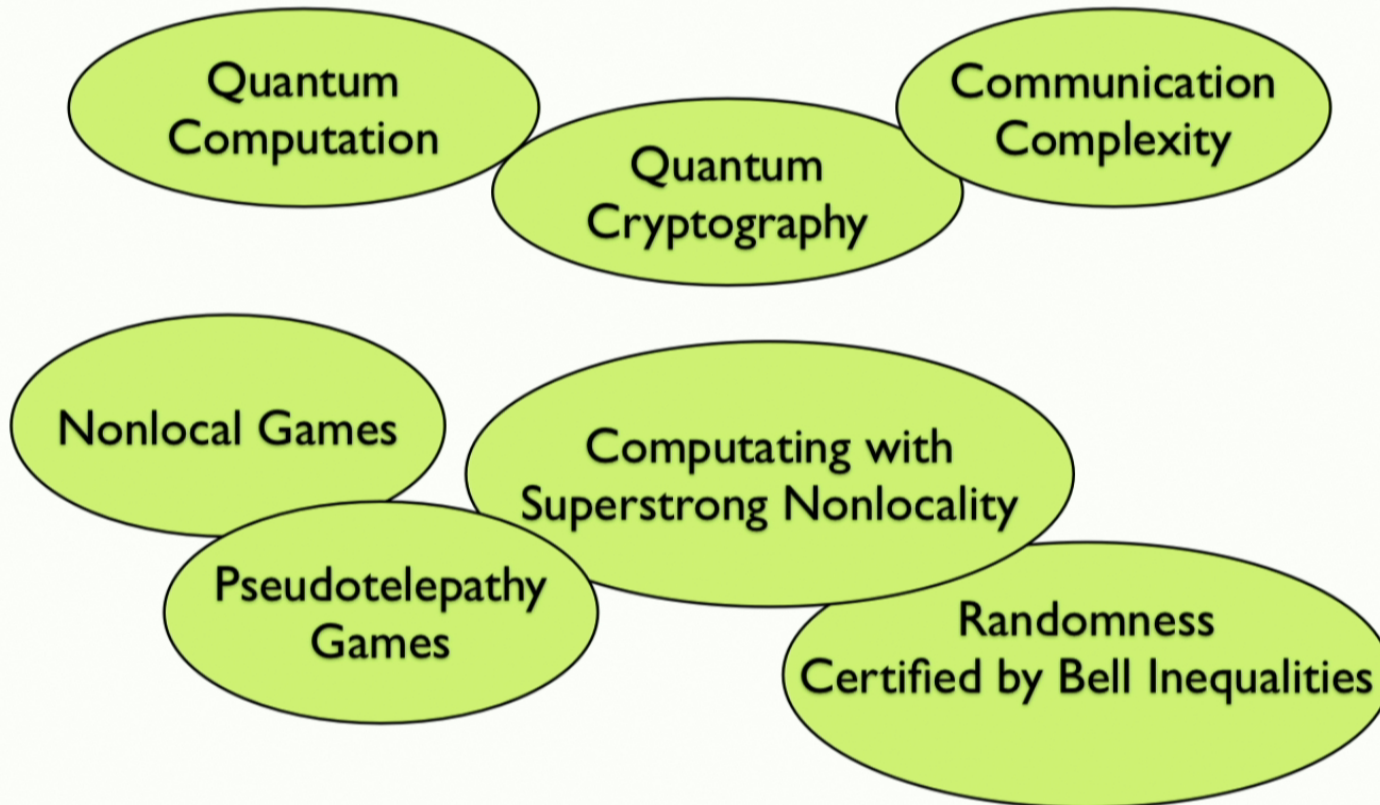
Yes!

No!

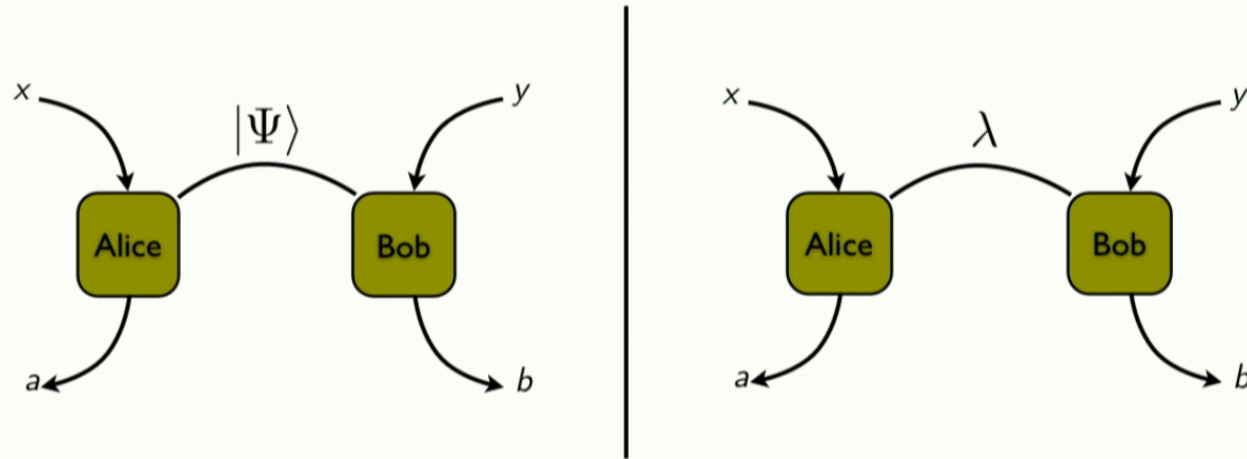


Applications of quantum nonlocality

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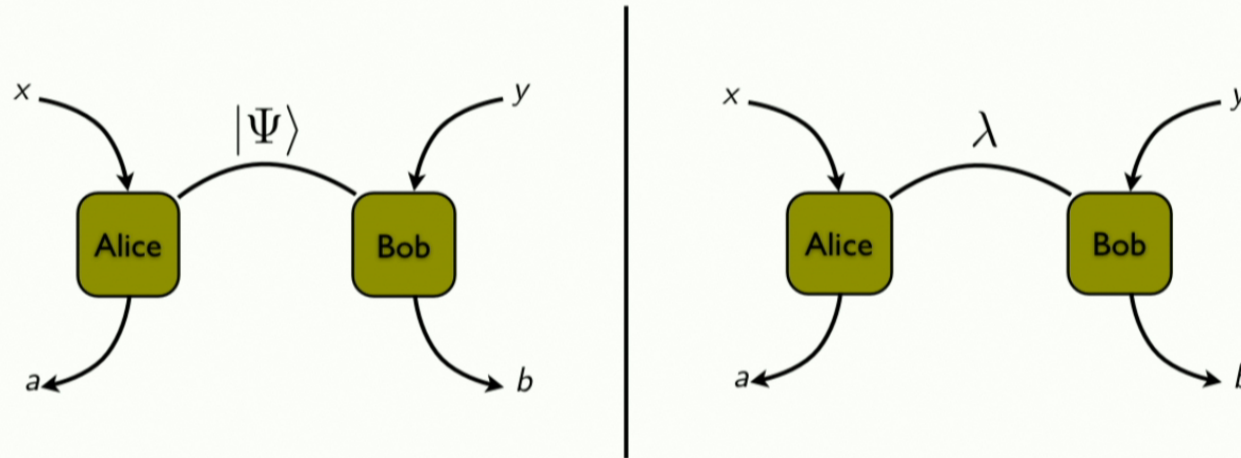


Quantifying nonlocality



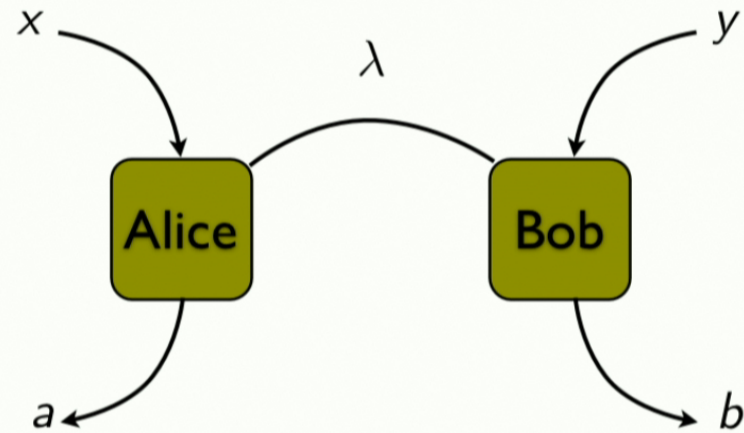
How far are these two models?

Quantifying nonlocality

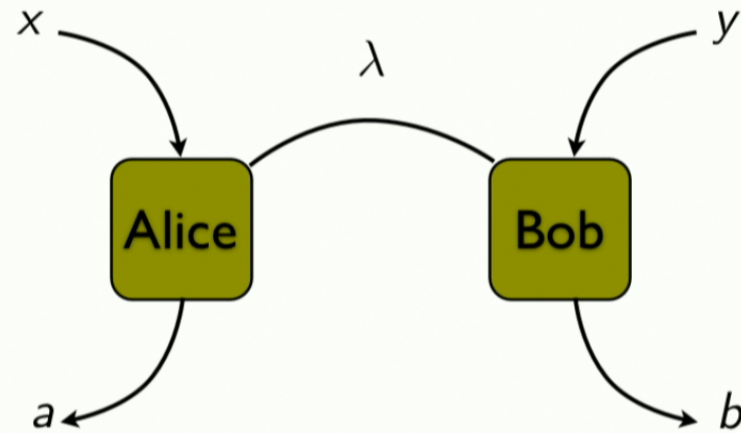


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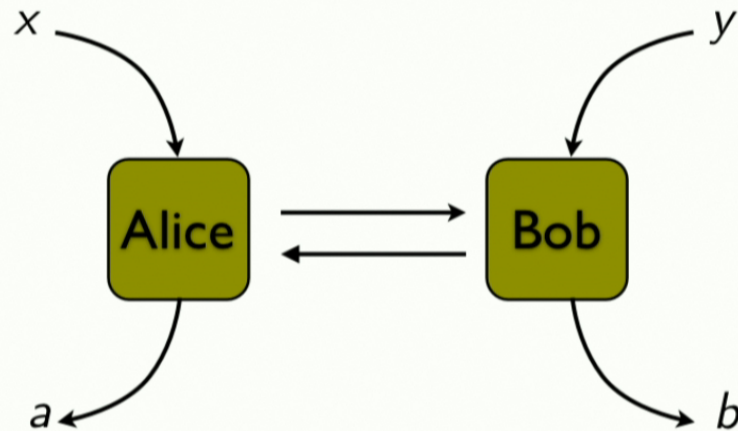
Quantifying nonlocality



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How much communication is required to reproduce the predictions of quantum mechanics?

What do we want to simulate?

We can't hope to simulate classically **everything** made possible by entanglement.

Quantum Teleportation



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But can we simulate **the effect of measurements?**

Quantifying nonlocality

Simulating binary observables on bipartite states

[Maudlin'92] 1.17 expected bits

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Quantifying nonlocality

Simulating binary observables on **bipartite** states

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[Cerf-Gisin-Massar'00] 1.19 expected bits

[Toner-Bacon'03] 1 bit worst case (\uparrow two maximally entangled qubits \downarrow)

[Regev-Toner'07] 2 bits worst case (arbitrary bipartite state)

How about multipartite states?

What do we want to simulate?

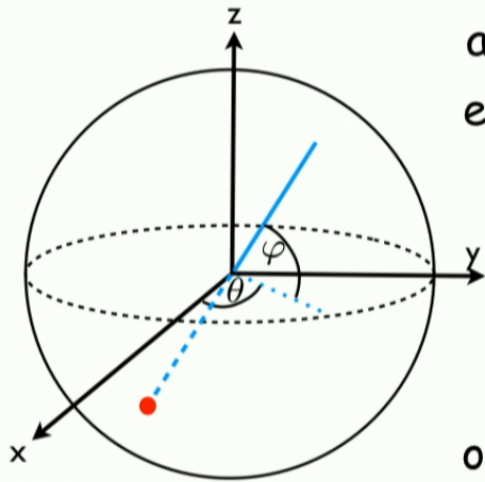
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Von Neumann measurement

Von Neumann measurement



azimuthal angle $0 \leq \theta < 2\pi$

elevation angle $-\frac{\pi}{2} \leq \varphi \leq \frac{\pi}{2}$

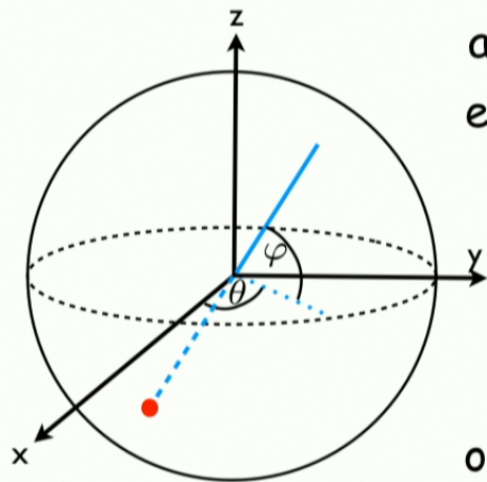
$$x = \cos \theta \cos \varphi$$

$$y = \sin \theta \cos \varphi$$

$$z = \sin \varphi$$

outcome of measurement is +1 or -1

Von Neumann measurement



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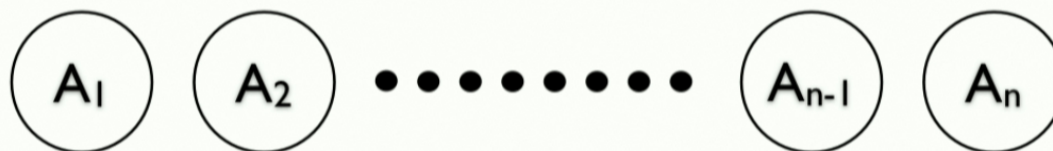
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GHZ States

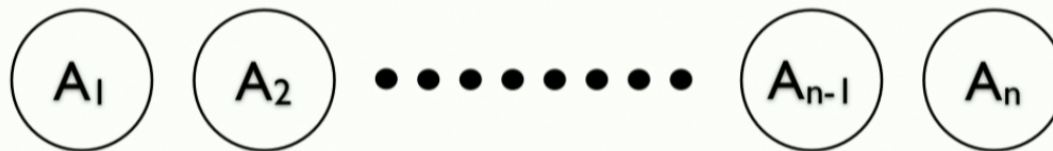
GHZ States

$$|\Psi_n\rangle = \frac{1}{\sqrt{2}}|00\dots 0\rangle + \frac{1}{\sqrt{2}}|11\dots 1\rangle$$



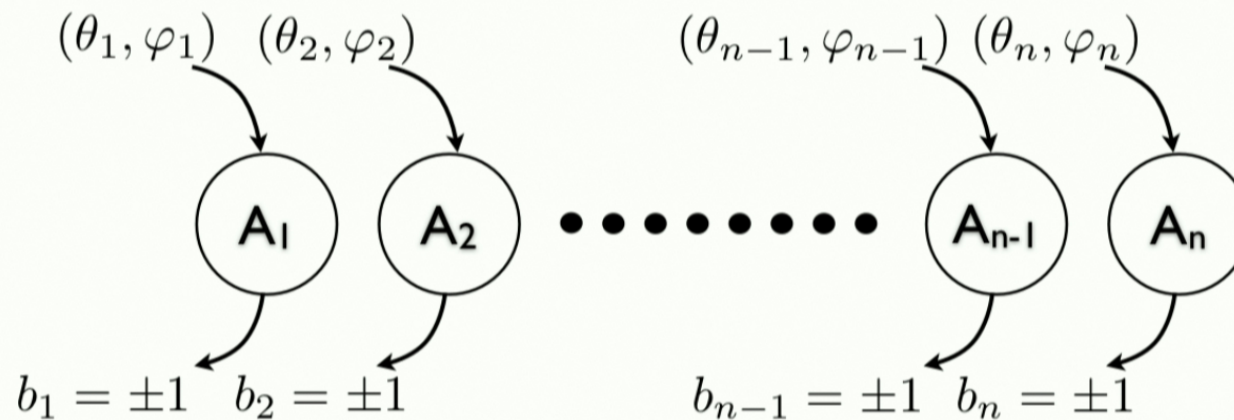
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GHZ States

Classical communication



$$p((b_1, b_2, \dots, b_n) \mid (\theta_1, \varphi_1), (\theta_2, \varphi_2), \dots, (\theta_n, \varphi_n))$$

Simulation of GHZ distribution

Simulating equatorial measurements

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[Bancal-Branciard-Gisin'10]:

10 expected bits for $n=3$, 20 expected bits for $n=4$.

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$\Omega(n \log n)$ lower bound in the worst case.

Simulating arbitrary von Neumann measurements?

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Simulating arbitrary von Neumann measurements?

$O(n^2)$ expected bits for arbitrary n ; $O(n \log n)$ parallel time.

$O(n \log n)$ expected bits for equatorial measurements;

$O(\log^2 n)$ parallel time.

GHZ probability distribution

$$p(b) = \cos^2\left(\frac{\theta}{2}\right) p_1(b) + \sin^2\left(\frac{\theta}{2}\right) p_2(b)$$

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It is a **convex decomposition**

Key observation

It is possible to sample **exactly**
a probability distribution whose
parameters are given as
arbitrarily precise **approximations**

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This would **not** be possible
if we wanted to perform
actual measurements!

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Example: Bernoulli (heads or tails)
with probability p of heads=0

Bernoulli(p)

Choose $U \in_r [0, 1)$

Bernoulli(p)

Choose $U \in_r [0, 1)$

If $U < p$ then return 0 else return 1

$$p = 0.01010101010101 = \frac{1}{3}$$

$$U = 0.01101101010100$$

Bernoulli(p)

Choose $U \in_r [0, 1)$

If $U < p$ then return 0 else return 1

$$\begin{aligned} p &= 0.01010101010101 &= \frac{1}{3} \\ U &= 0.01101101010100 \end{aligned}$$

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$$k + \lceil \lg n \rceil + 4 \text{ bits of each } \frac{\theta_j}{2} \implies k \text{ bits of } \cos^2 \frac{1}{2} \left(\sum \theta_j \right)$$

GHZ probability distribution

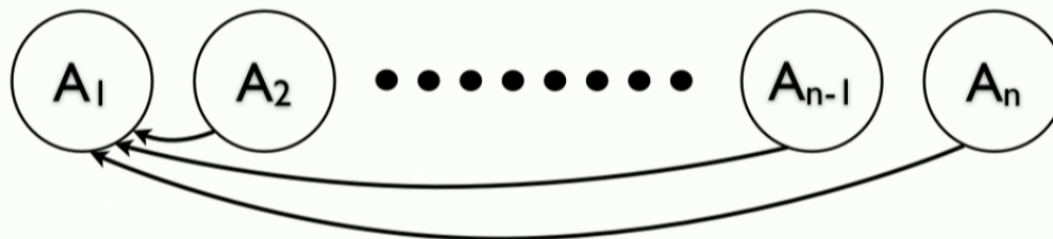
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$O(n \log n)$ expected bits of communication suffice
in order to decide whether to sample p_1 or p_2

GHZ probability distribution

$$p(b) = \cos^2\left(\frac{\theta}{2}\right) p_1(b) + \sin^2\left(\frac{\theta}{2}\right) p_2(b)$$

$$p_1(b) = \frac{1}{2} (a_1(b) + a_2(b))^2 \quad p_2(b) = \frac{1}{2} (a_1(b) - a_2(b))^2$$

$$a_1(b) = \prod_{j|b_j=+1} \alpha_j \prod_{j|b_j=-1} \beta_j \quad a_2(b) = \prod_{j|b_j=+1} \beta_j \prod_{j|b_j=-1} -\alpha_j$$

$$\alpha_j = \sin\left(\frac{1}{2}\left(\varphi_j + \frac{\pi}{2}\right)\right) \quad \beta_j = \cos\left(\frac{1}{2}\left(\varphi_j + \frac{\pi}{2}\right)\right)$$

$$\alpha_j^2 + \beta_j^2 = 1$$

GHZ probability distribution

$$p(b) = \cos^2\left(\frac{\theta}{2}\right) p_1(b) + \sin^2\left(\frac{\theta}{2}\right) p_2(b)$$

$$p_1(b) = \frac{1}{2} (a_1(b) + a_2(b))^2 \quad p_2(b) = \frac{1}{2} (a_1(b) - a_2(b))^2$$

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$$\alpha_j = \sin\left(\frac{1}{2}(\varphi_j + \frac{\pi}{2})\right) \quad \beta_j = \cos\left(\frac{1}{2}(\varphi_j + \frac{\pi}{2})\right)$$

$$p_1(b) \leq p_1(b) + p_2(b) = a_1^2(b) + a_2^2(b) = 2 \frac{a_1^2(b) + a_2^2(b)}{2}$$

Very easy to sample!

Von Neumann Rejection Algorithm

Say we want to sample distribution $f(x)$

We know how to sample $g(x)$

There exist c such that $(\forall x) f(x) \leq c g(x)$

Sample X according to distribution g

Sample U according to uniform distribution on $[0, 1)$

Von Neumann Rejection Algorithm

Say we want to sample distribution $f(b) = p_1(b)$

We know how to sample $g(b)$

There exist c such that $(\forall b) f(b) \leq c g(b)$

Sample B according to distribution g

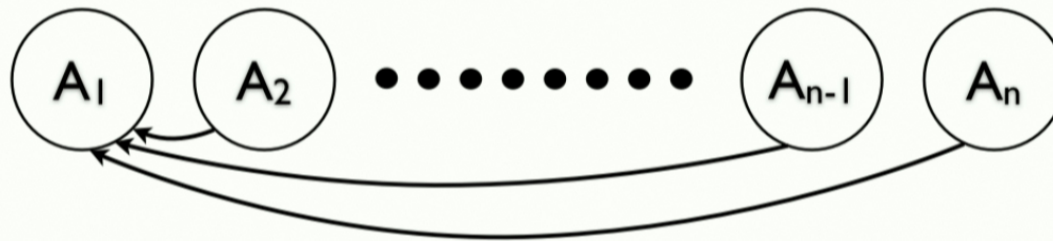
Sample U according to uniform distribution on $[0, 1)$

If $f(B) \leq c g(B) U$ go back to “Sample B ”

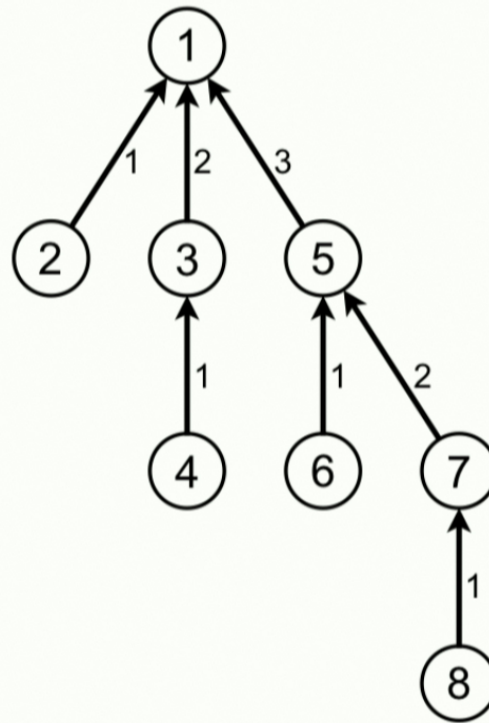
B is now sampled according to distribution f !

The expected number of times round the loop is c

Parallel version



Parallel version



Open problems

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- $O(n \log n)$ expected communication?
- Worst case communication?
- More general multipartite states?