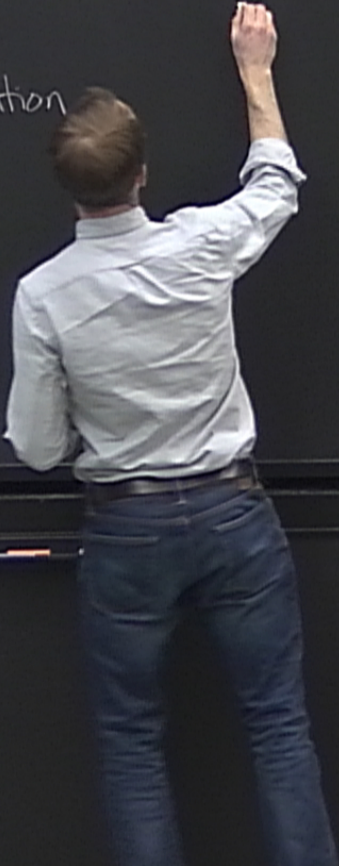


Title: 13/14 PSI - Quantum Information Review - Lecture 2

Date: Feb 19, 2014 11:30 AM

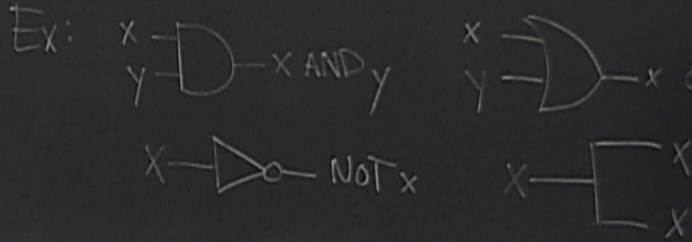
URL: <http://pirsa.org/14020073>

Abstract:

- 
- A man in a white shirt and blue jeans is standing with his back to the camera, writing on a blackboard. He is holding a piece of chalk in his right hand, which is raised towards the top of the board. The blackboard is dark and has three lines of text written on it in white chalk. The text is a list of topics: '- Circuits', '- Reversible computation', and '- Universality'. The man is positioned in the lower-left quadrant of the frame, facing the blackboard which occupies most of the background.
- Circuits
 - Reversible computation
 - Universality

- Circuits
- Reversible computation
- Universality

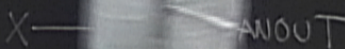
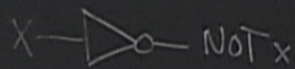
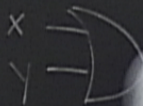
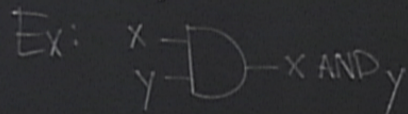
Classical circuit model:
 consider a set of logic gates
 acting on bits



the

Classical circuit model:

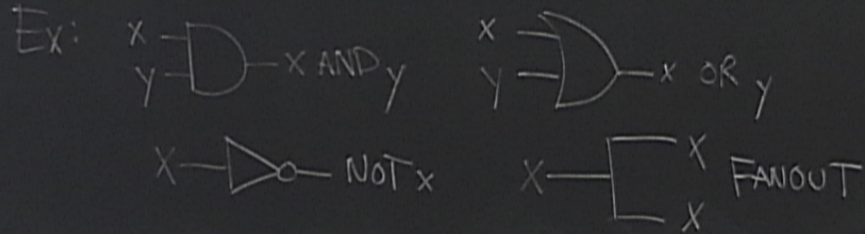
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this set is universal: it can compute any function $f: \{0,1\}^n \rightarrow \{0,1\}^m$

Classical circuit model:

consider a set of logic gates acting on bits

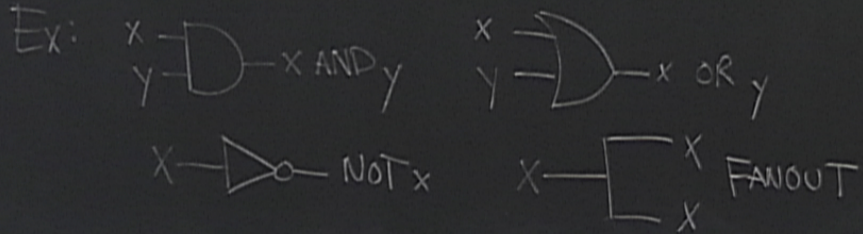


this set is universal: it can compute any function $f: \{0,1\}^n \rightarrow \{0,1\}^m$

computation

Classical circuit model:

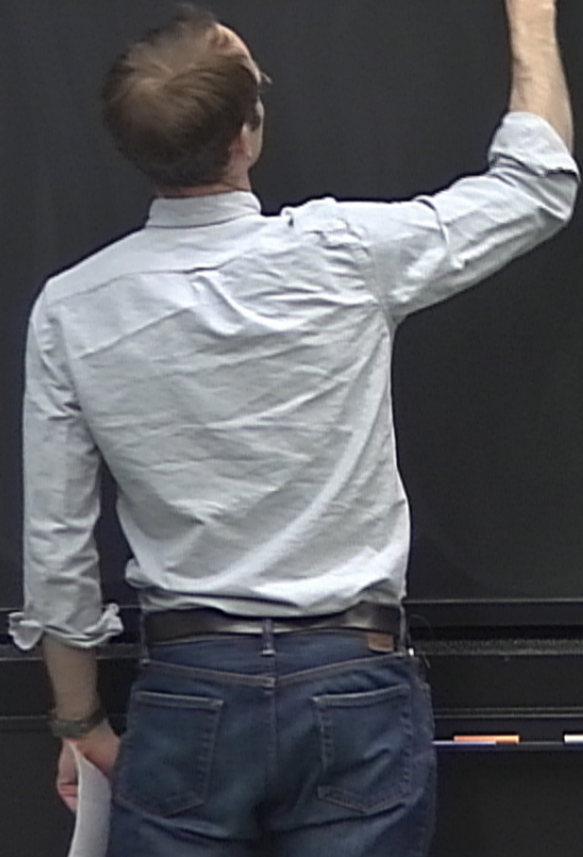
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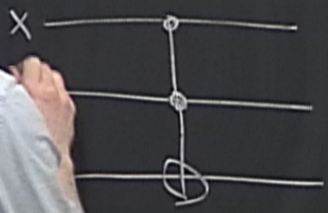
Ex: NAND $\begin{matrix} x \\ y \end{matrix} \rightarrow \text{NOT}(x \text{ AND } y)$ and FANOUT

To make a reversible w



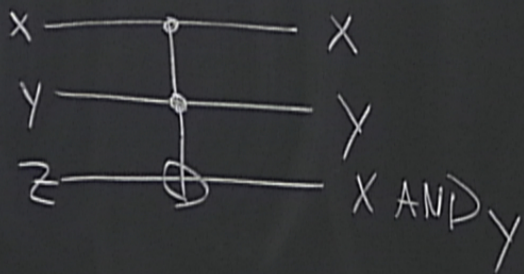
To make a reversible model,
keep the input around.

Toffoli gate



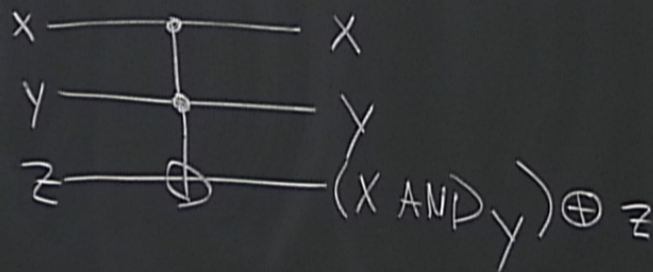
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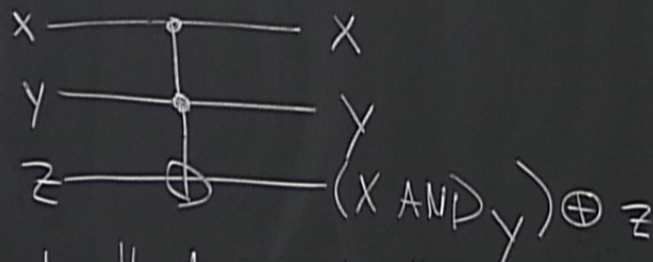
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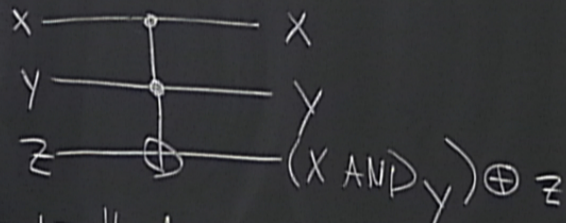


"controlled-controlled-NOT"

To make a reversible model,
keep the input around.

this gate alone is sufficient
for universality

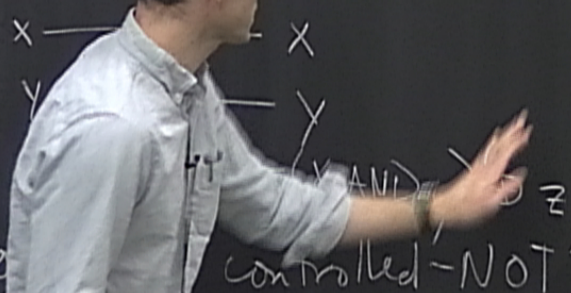
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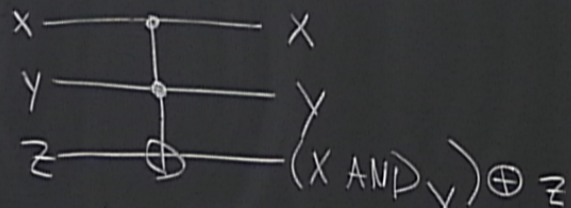
(may need workspace, aka ancillas)

NOT: set $x=y=1$

FANOUT: set $y=1, z=0$

To make a reversible model,
keep the input around.

Ex: Toffoli gate



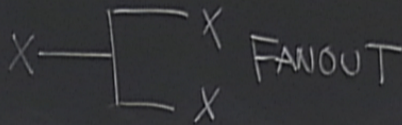
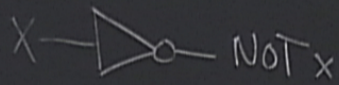
"controlled-controlled-NOT"

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NOT: set $x=y=1$

FANOUT: set $y=1, z=0$



e.g., can't infer (x, y) from $x \text{ AND } y$

model,
around.

this gate alone is sufficient
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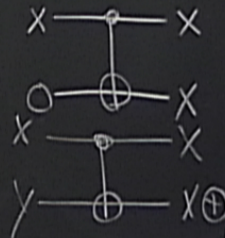
(may need workspace, aka ancillas)

NOT: set $x=y=1$

FANOUT: set $y=1, z=0$

$(x \text{ AND } y) \oplus z$
"controlled-NOT"

controlled-NOT gate:

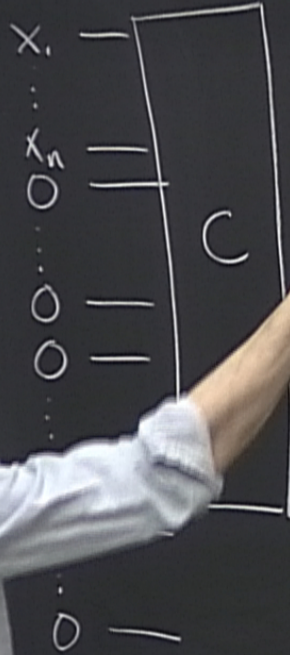


Any function can be computed reversibly:

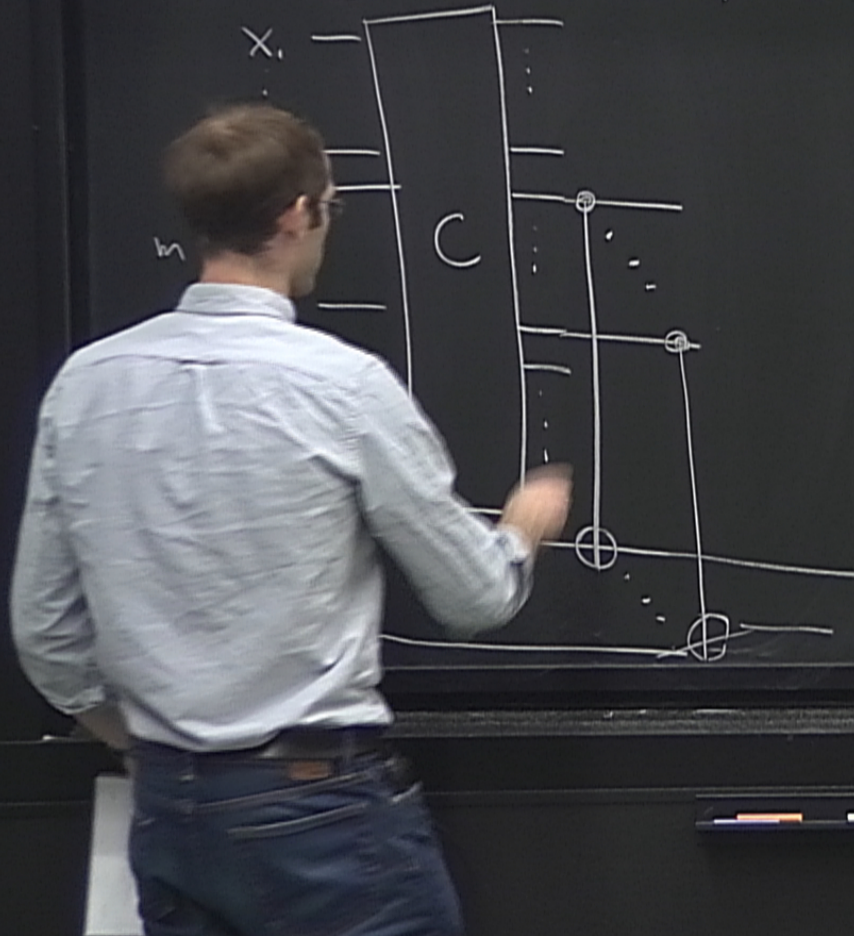
x_1 —
⋮
 x_n —



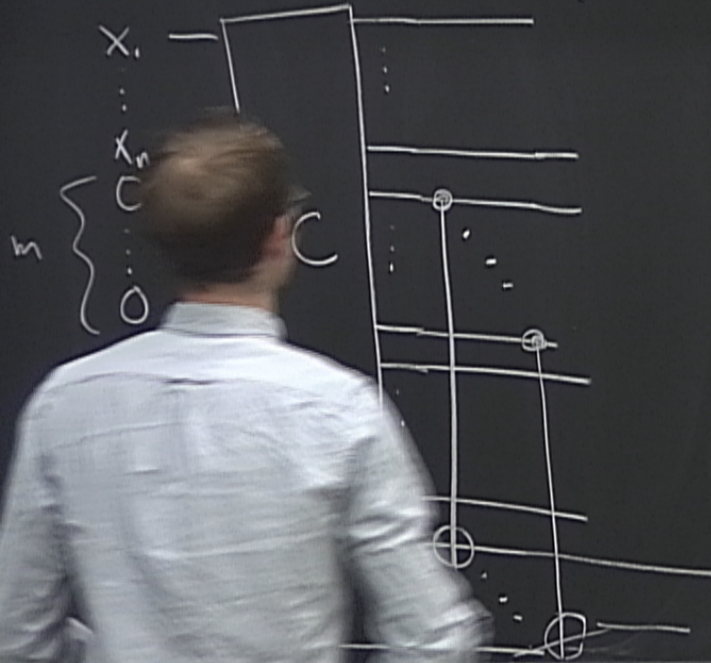
to reset the workspace.



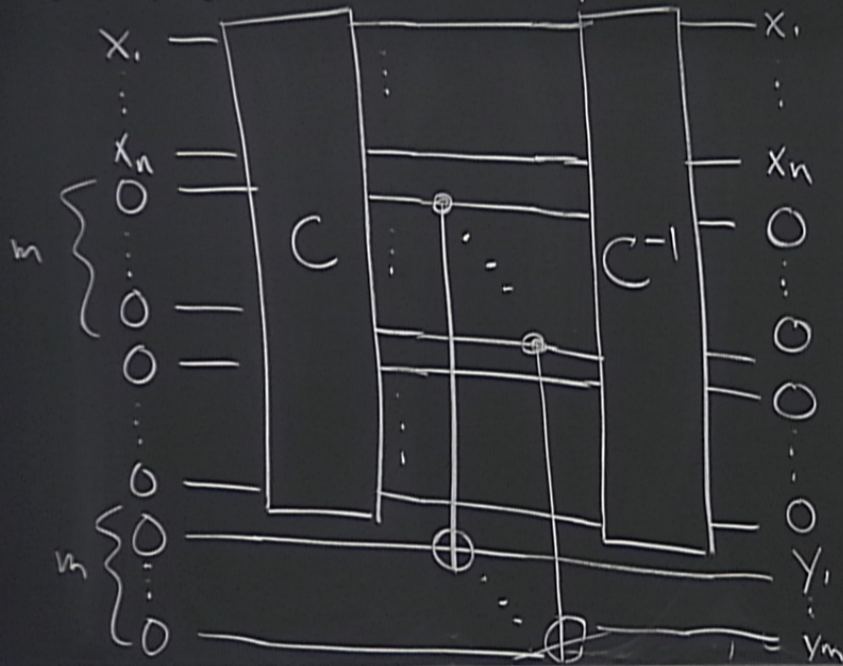
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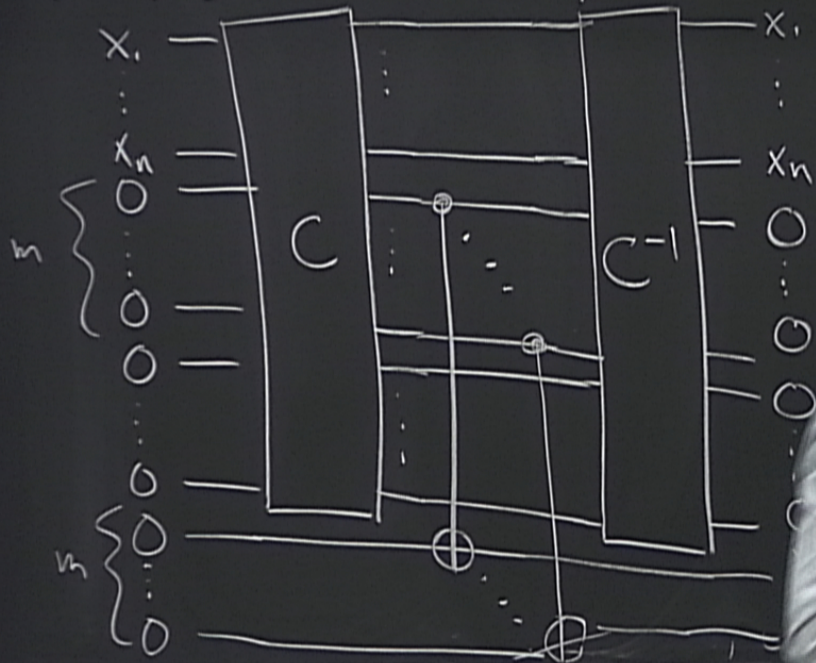
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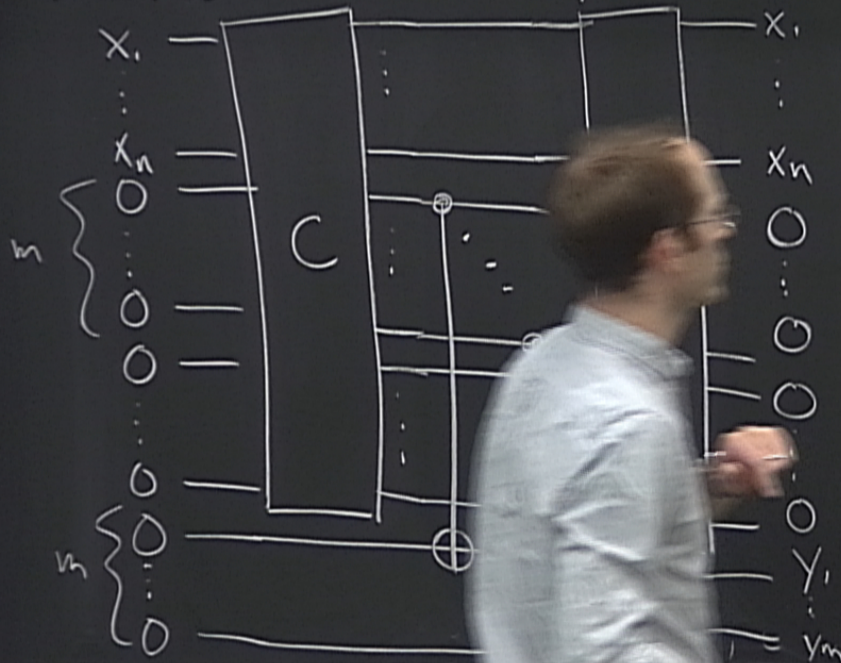
In general, we can compute

$$(x, z_0) \mapsto (x, z_0 \oplus f(x))$$

we can perform f efficiently



to reset the workspace.

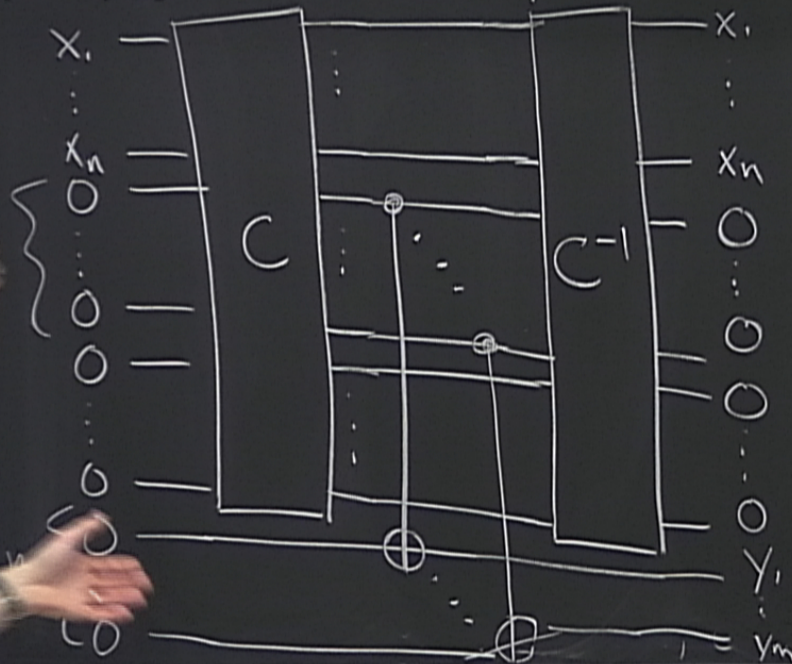


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to reset the workspace.

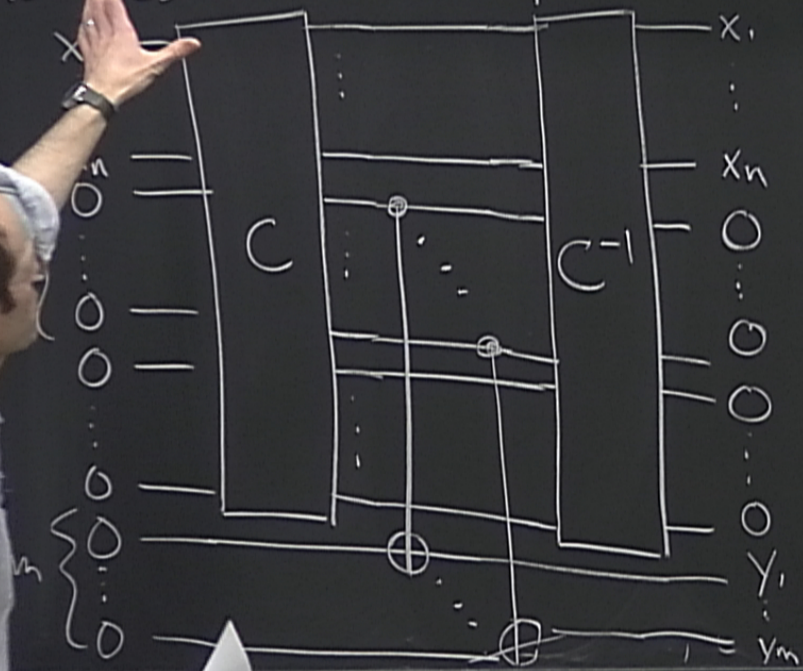


In general, we can compute

$$(x, z_0) \mapsto (x, z_0 \oplus f(x))$$

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In general, we can compute

$$(x, z, 0) \mapsto (x, z \oplus f(x), 0)$$

if we can perform f efficiently

we can do the same thing with g circuits:

$$|x, z, 0\rangle \mapsto |x, z \oplus f(x), 0\rangle$$

compute

$z \oplus f(x)$

f efficiently

with g

$|x, z \oplus f(x), 0\rangle$

Quantum universality

Q. universality means the ability to do any unitary
in $U(2^n)$

Can show

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Quantum universality

Q. universality means the ability to do any unitary
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Can show that CNOT +

compute

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Can show that CNOT + all single-qubit gates is a universal set

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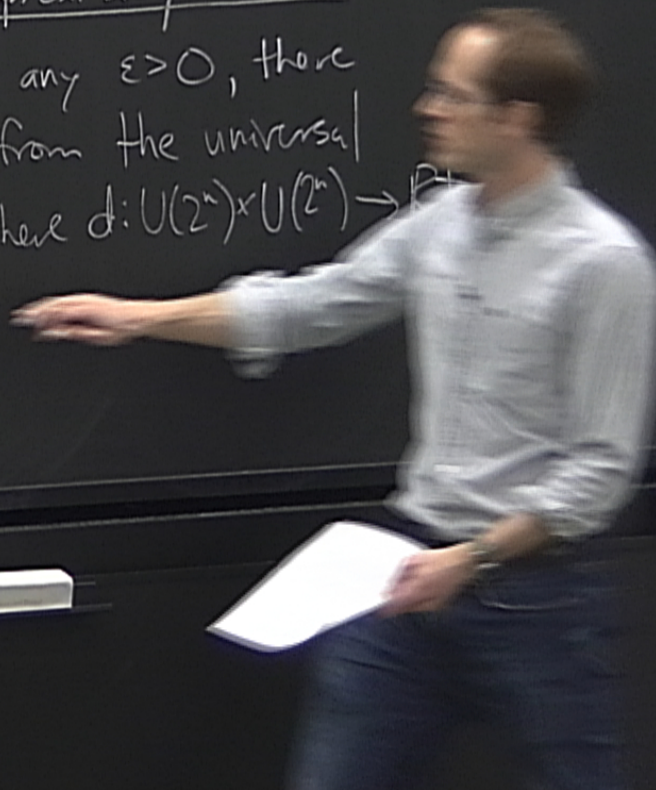
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But we'd like to have a finite gate set

To get this, we only demand approximate realization

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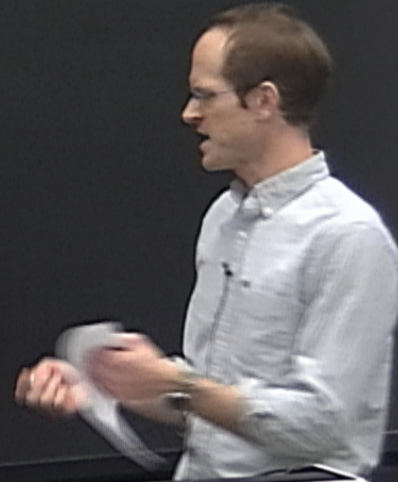
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We say a set of gates is approximately universal if, for any $U \in U(2^n)$ and any $\epsilon > 0$, there is a circuit V using gates from the universal set so that $d(U, V) \leq \epsilon$, where $d: U(2^n) \times U(2^n) \rightarrow \mathbb{R}^+$ is some distance measure.

Ex: Could use

$$d(U, V) = \|U - V\| \quad (\text{spectral norm})$$
$$= \max_{\|x\|=1} \|Ux - Vx\|$$

$\rightarrow \mathbb{R}^+$



ization

universal

, there

universal

$U(2^n) \rightarrow \mathbb{R}^+$

Ex: Could use

$$d(U, V) = \|U - V\| \quad (\text{spectral norm})$$
$$= \max_{|\psi\rangle} \|\langle \psi | U - V | \psi \rangle\|$$

Since CNOT + all 1-qubit gates is universal, it suffices to approximate any 1-qubit gate using gates from a finite set



A set of gates that suffices:

$$\{H, T\} \quad H = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} \text{ Hadamard}$$

universal

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So $\{H, T\}$ is a universal set
for

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So $\{CNOT, H, T\}$ is a universal set for any # of qubit.

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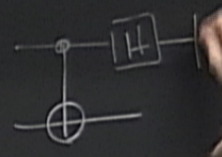
... by irrational multiples of 2π

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Symbols:



1001 by irrational multiples of π

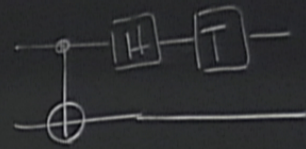
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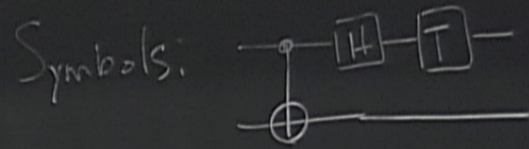
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Another common notation:

