

Title: Quantum Factoring with Trapped Ions - A test of scalability

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Abstract: Shor's algorithm can be a meaningful test for experimental quantum processing systems, when suitably realized. I present results from a recent implementation of quantum factoring using trapped ion qubits, demonstrating feed-forward control, use of quantum memory during computation, and cascaded three-qubit gates. Such capabilities are necessary ingredients for a future large-scale, fault-tolerant quantum computing system.



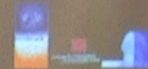
Quantum Factoring with Trapped Ions

– A TEST OF SCALABILITY

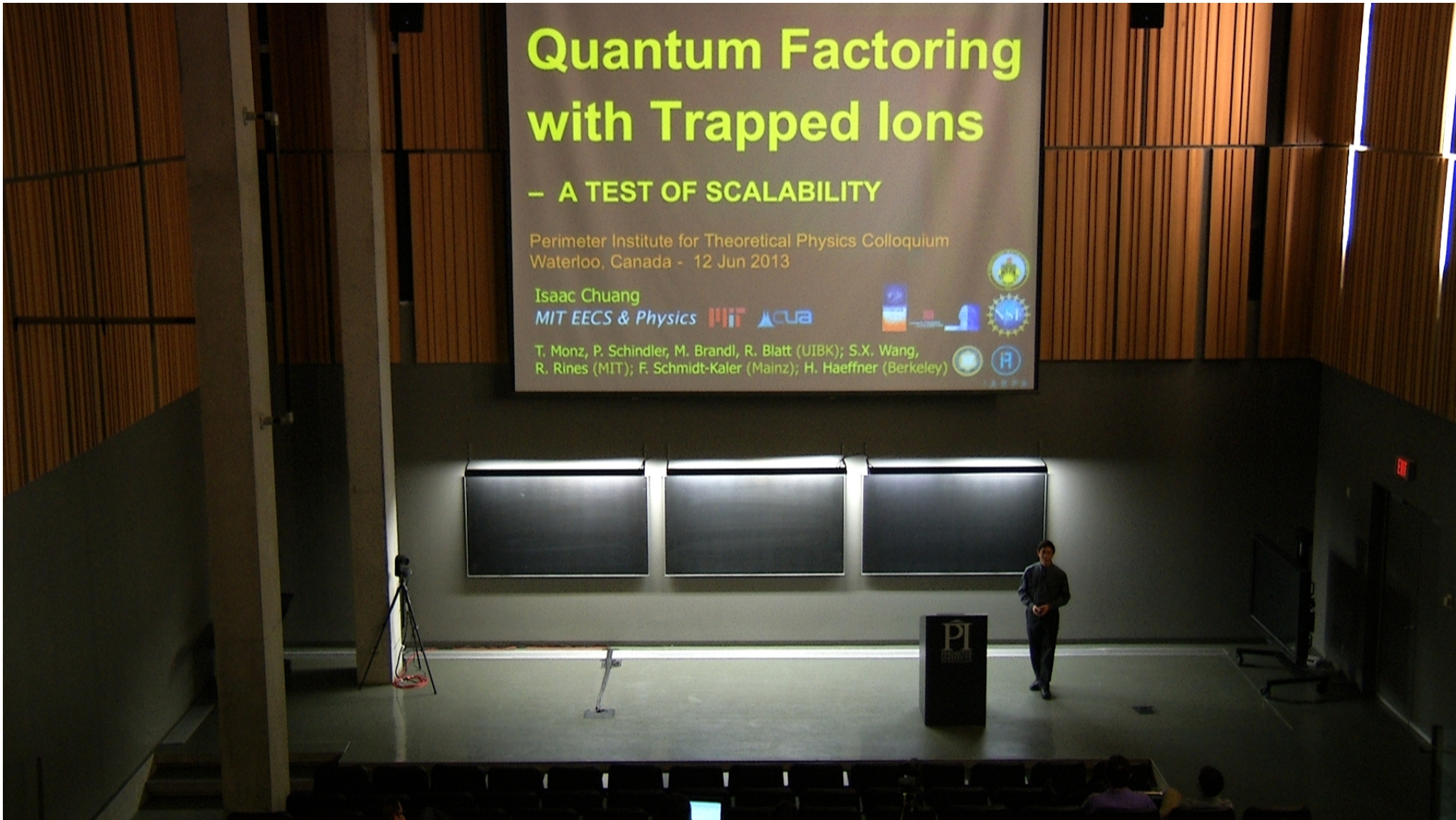
Perimeter Institute for Theoretical Physics Colloquium
Waterloo, Canada - 12 Jun 2013

Isaac Chuang

MIT EECS & Physics



T. Monz, P. Schindler, M. Brandl, R. Blatt (UIBK); S.X. Wang,
R. Rines (MIT); F. Schmidt-Kaler (Mainz); H. Haeflner (Berkeley)



Shor's Q. Factoring Algorithm

$$f(x) = a^x \bmod N$$

↑ ↑
coprime with N composite number

Results from number theory:

- f is periodic in x (period r)
- $\gcd(a^{r/2} \pm 1, N)$ is a factor of N



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Quantum factoring: find r

Complexity of factoring
numbers of length L :

Quantum: $\sim L^3$ P. Shor (1994)

Classically: $\sim e^{L/3}$

Widely used crypto systems (RSA) would become insecure.

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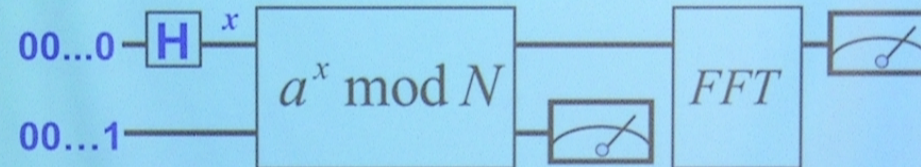
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The Quantum Factoring Algorithm: Quick Review

Theoretical Algorithm (Shor, 1994)



Three main steps:

1. Input superposition preparation
2. Modular exponentiation (multi-qubit gates required)
3. Quantum Fourier Transform
4. Classical pre- and post-processing

Factoring Algorithm: Universality

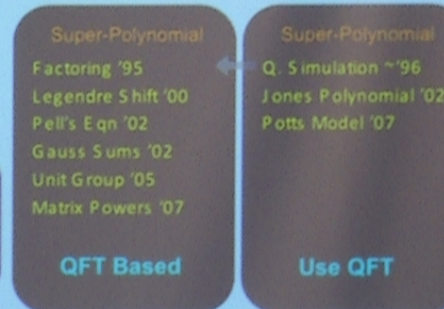
- The QFT structure of Shor's algorithm is universal to virtually all exponentially-fast quantum algorithms
- Beyond factoring: linear equations

• **Problem: given a linear system of equations**

$$A\vec{x} = \vec{b} \quad \text{estimate } \vec{x}^\dagger M \vec{x}$$

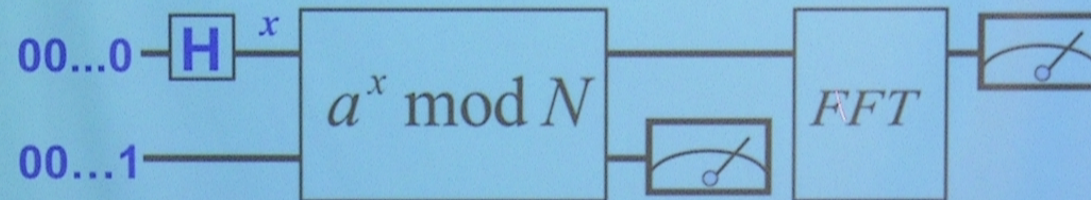
← dim n ; sparse, well-conditioned
← Some matrix

Uses quantum phase estimation (QPE)



Simplest Meaningful Factoring?

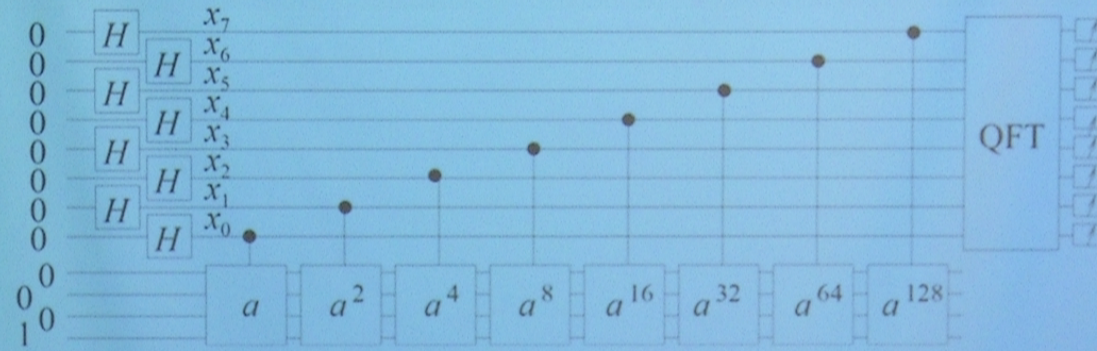
- Quantum Factoring of $N=15$



Generic factoring circuit

Simplest Meaningful Factoring?

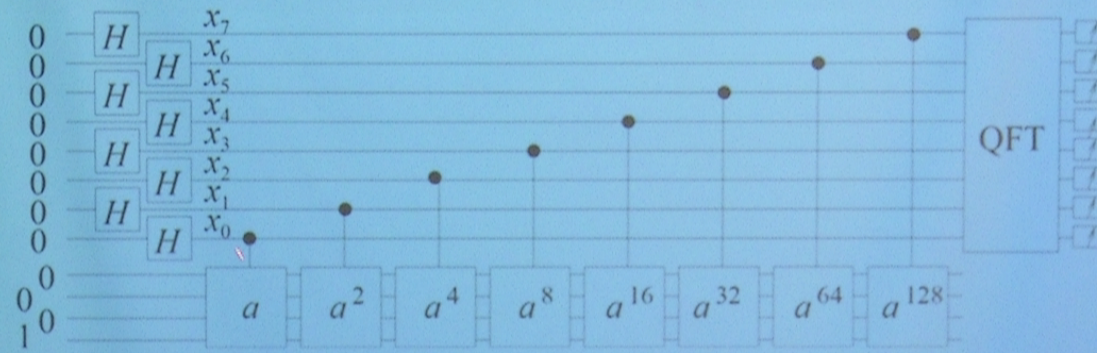
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$a=7$ (hard) or $a=11$ (easy)

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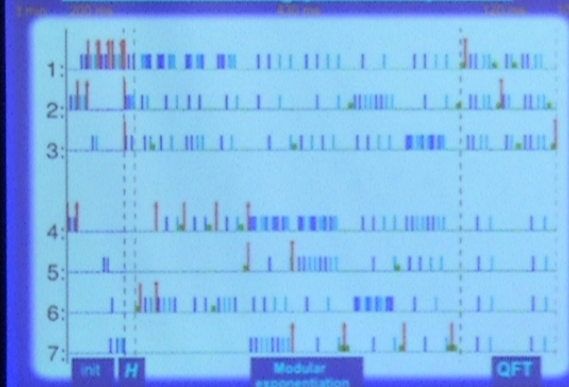
Quantum Factoring

(Vandersypen, et al, Nature, Dec. 2001)

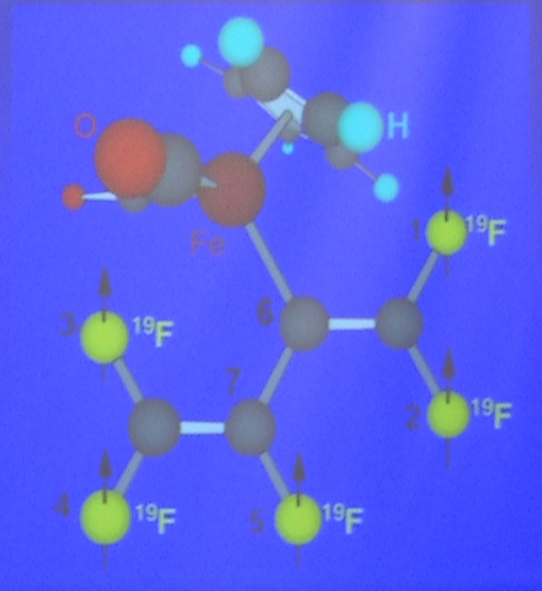
- Expt. demonstration of Shor's factoring algorithm

$$15 = 3 \times 5$$

NMRQC factoring pulse sequence



- The Molecule $T_2 > 0.3$ sec
~ 200 gates



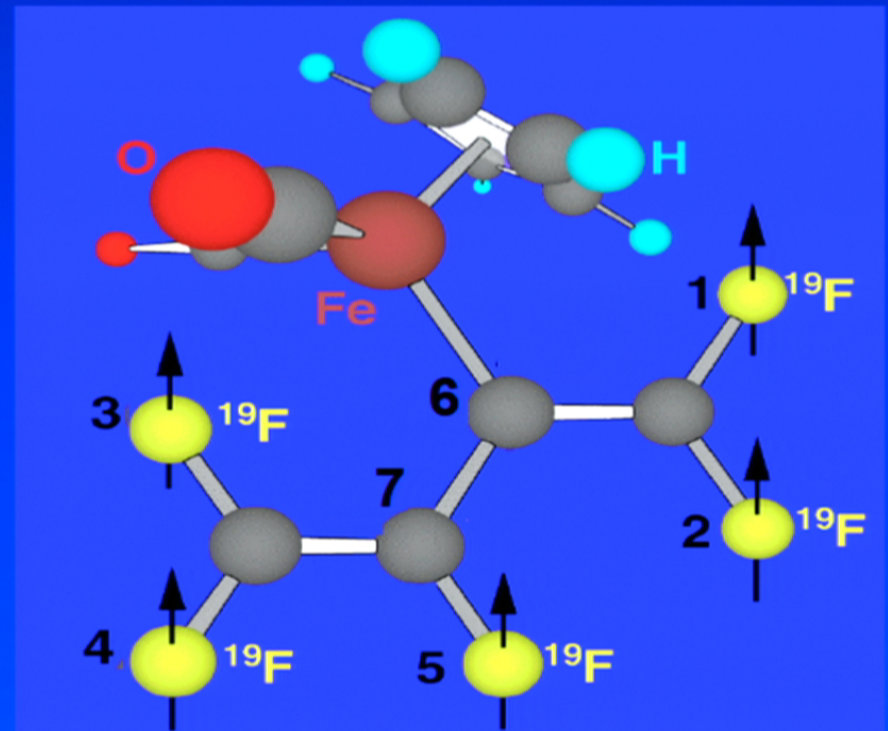
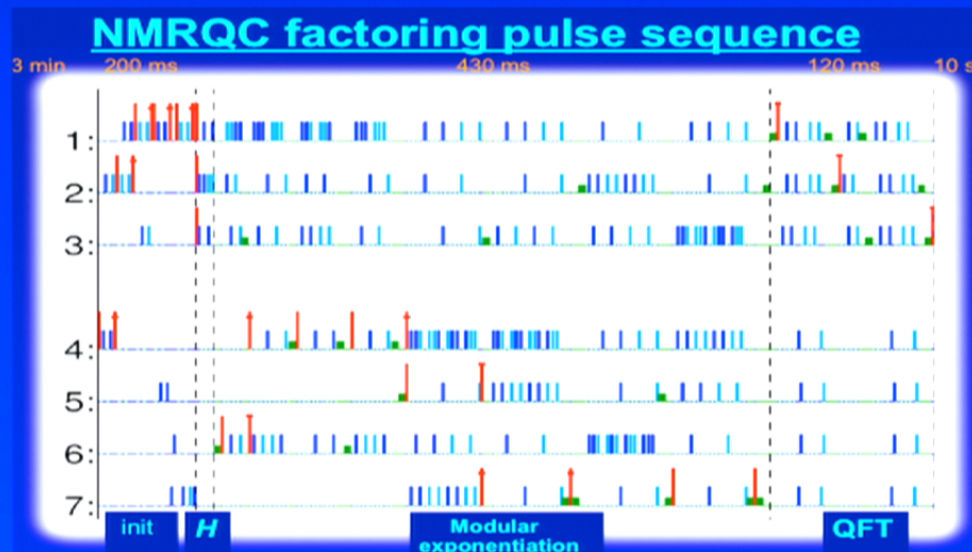
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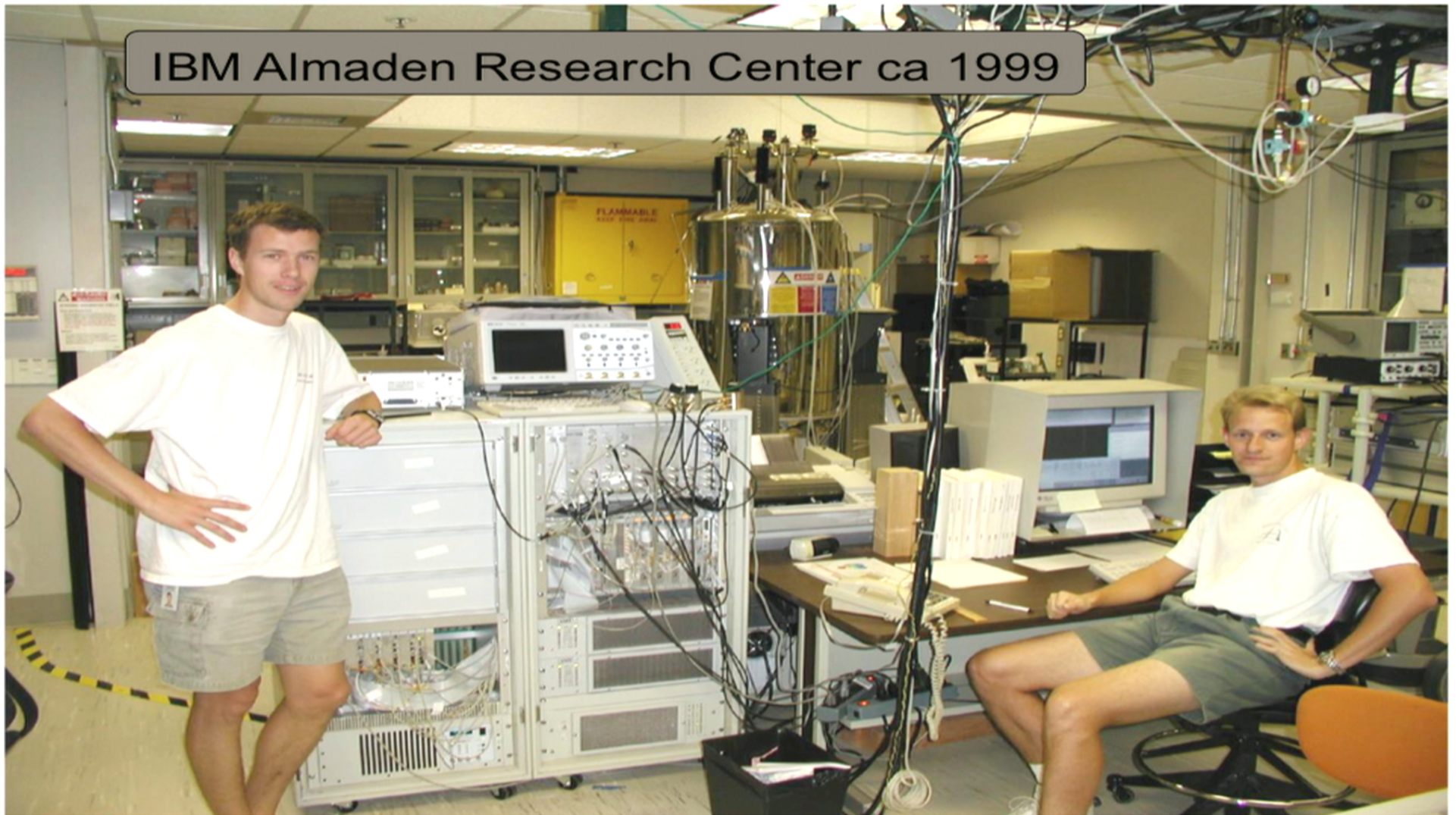
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IBM Almaden Research Center ca 1999



Shor Algorithm Realizations: Prior Art & Perspective

Prior Art

- NMR: $N=15$ / 7 qubits / 200 pulses / Both **Easy** and **Hard** cases
Vandersypen et al. (2001)
- Linear optics: $N=15$ / 3 qubits / 13 gates / Only **Easy** case /
post-selected
Lanyon et al. (2007)
- Superconductors: $N=15$ / 3 qubits / 2 CNOTs / Only **Easy** case
Lucero et al. (2012)

Shor Algorithm Realizations: Prior Art & Perspective

Prior Art

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Vandersypen et al. (2001)
- Linear optics: $N=15$ / 3 qubits / 13 gates / Only **Easy** case /
post-selected
Lanyon et al. (2007)
 $N=21$ / 4 qubits / 26 gates / Only **Easy** case
Lopez et al. (2011)
- Superconductors: $N=15$ / 3 qubits / 2 CNOTs / Only **Easy** case
Lucero et al. (2012)

The Challenge:

10^6
gates

Devices → Systems

- **Fault-tolerant QC**

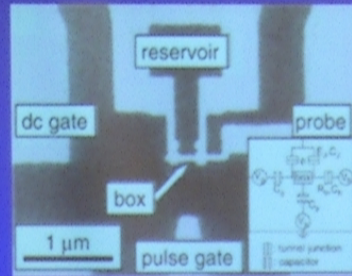
- Encoded qubits
- Memory / CPU
- “power supply” states
- classical control

Factoring;
Search;
Quantum
simulations

10^4
qubits

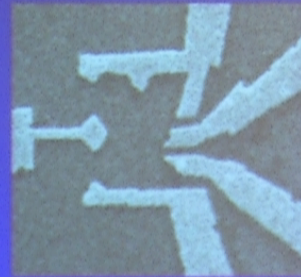
QIS&T Devices: ~2002

- Superconductor



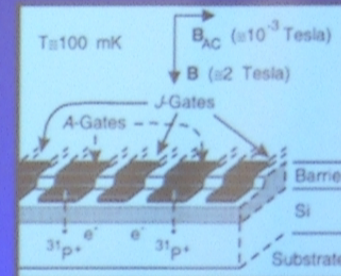
(Nakamura)

- Quantum Dots



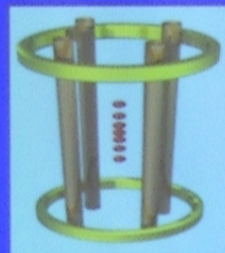
(Marcus / Tarucha)

- ^{31}P in Silicon



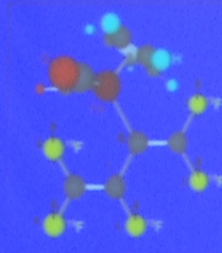
(Kane)

- Atoms



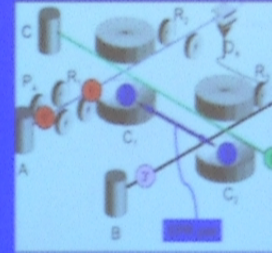
(Blatt / Wineland)

- NMR



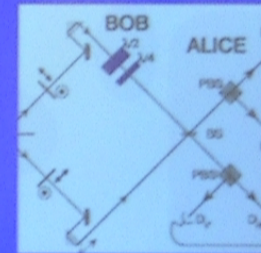
(Vandersypen et al)

- Cavity QED



(Brune / Haroche)

- Optics



(Zeilinger)

QIS&T Devices: ~2002

• Superconductor

2 qubits
1 two-qubit gate

(Nakamura)

• Quantum Dots

1 qubit
0 two-qubit gates

(Marcus / Tarucha)

• ^{31}P in Silicon

1 qubit
0 two-qubit gates

(Kane)

• Atoms

2 qubits
1 two-qubit gate

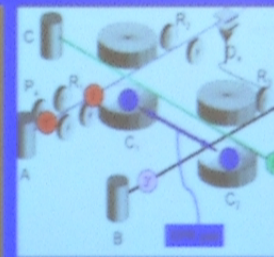
(Blatt / Wineland)

• NMR

7 qubits
20 two-qubit gates

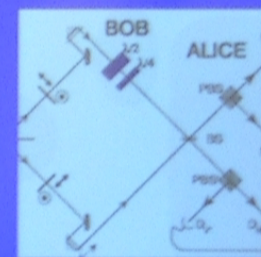
(Vandersypen et al)

• Cavity QED



(Brune / Haroche)

• Optics



(Zeilinger)

QIS&T Devices: ~2012

- Ions

14 qubits
10 two-qubit,
4 three-qubit
gates

(Blatt / Wineland ...)

- Circuit QED + SuperC

4 qubits
3 two-qubit,
1 three-qubit
gates

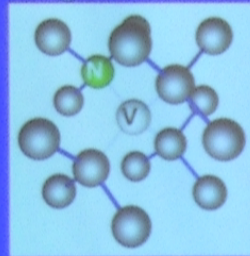
(Schoof / Harland / Staffor ...)

- Quantum Dots

2 qubits
1 two-qubit,
0 three-qubit
gates

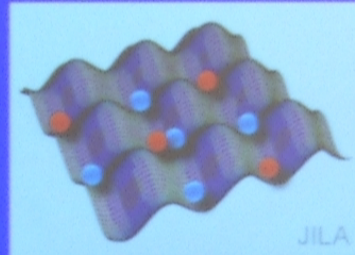
(Petráš / Kouřil ...)

- Nitrogen Vacancies



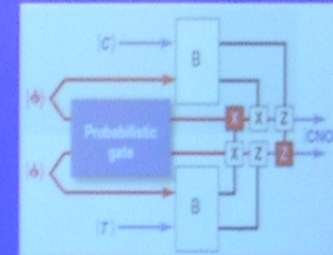
(Lukin / Wrachupt ...)

- Optical Lattices



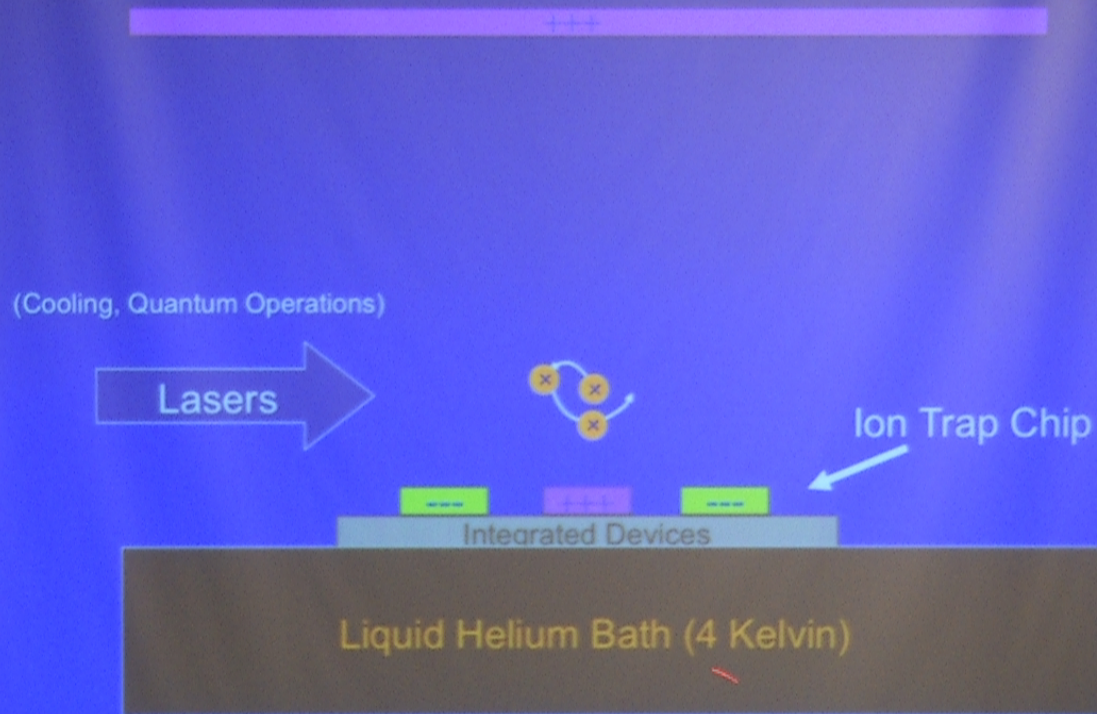
(Porto / Weiss / Saffman ...)

- Linear Optics

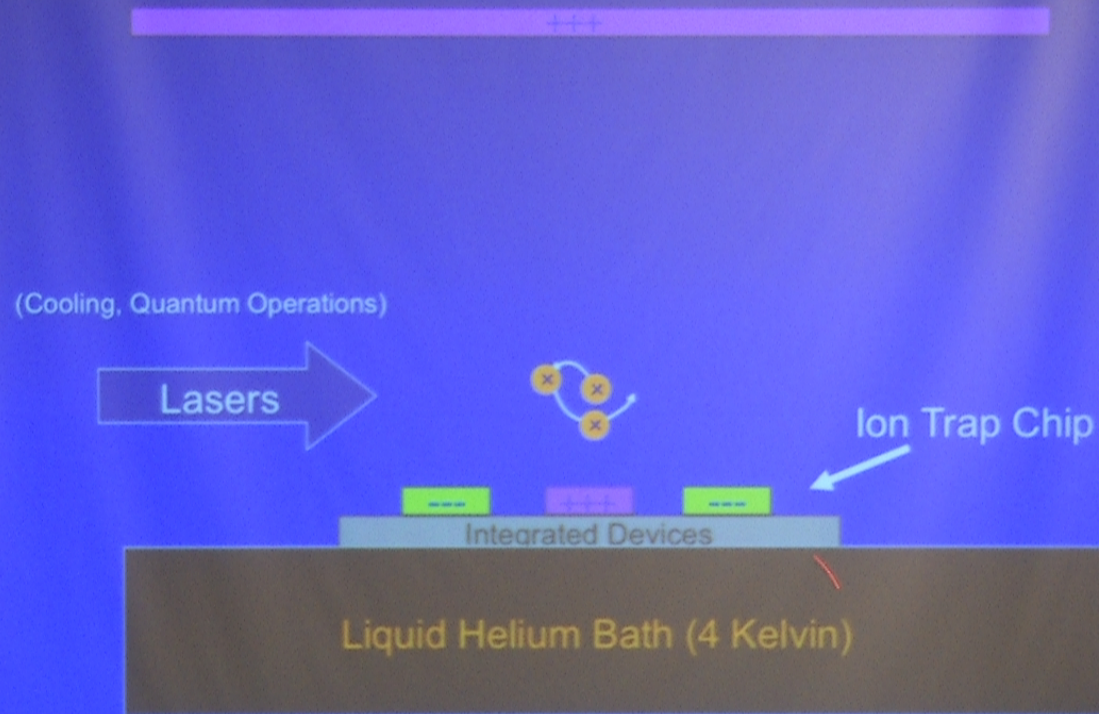


(White / Zeilinger / Pan ...)

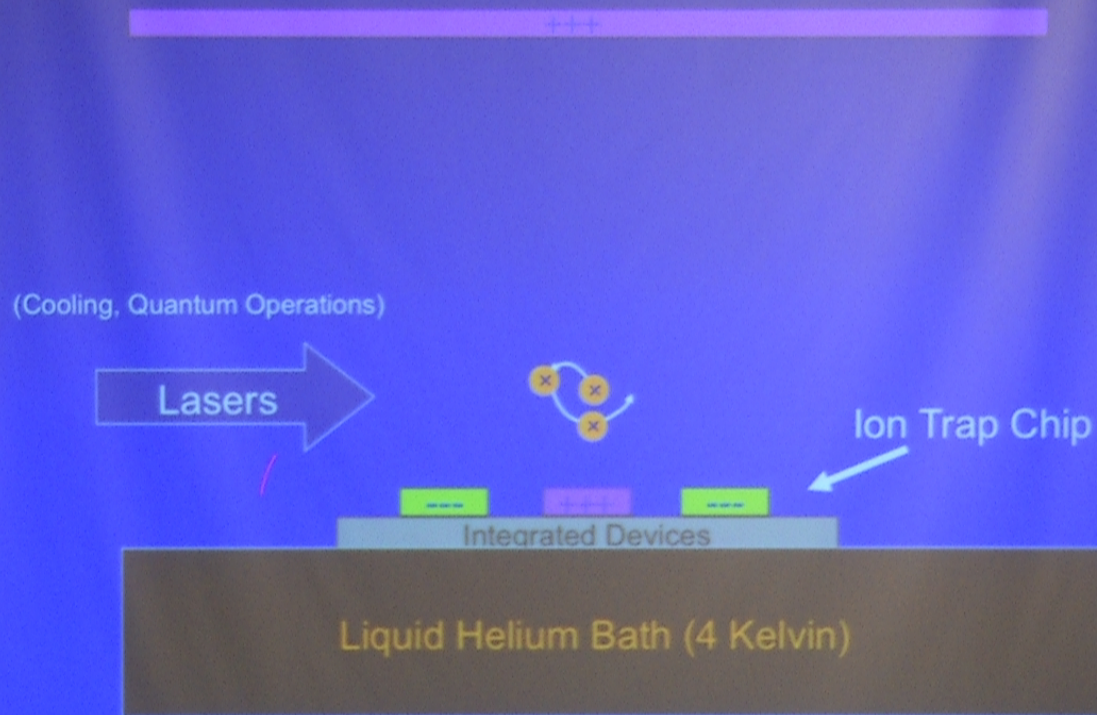
How to trap an atom on a chip



How to trap an atom on a chip

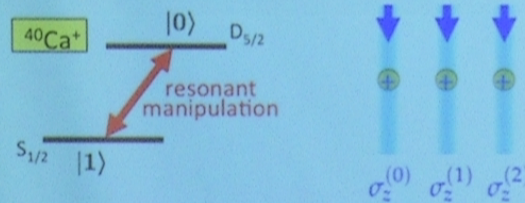


How to trap an atom on a chip



Trapped Ion QC Quantum Operations Toolbox

1. Addressed single-qubit rotations

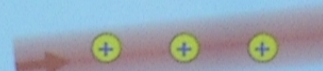


Generate rotations about z axis

$$U_{\sigma_z^{(i)}}(\theta) = e^{-i\frac{\theta}{2}\sigma_z^{(i)}}$$

2. Collective gates

Monochromatic:



$$S_x = \sigma_x^{(0)} + \sigma_x^{(1)} + \sigma_x^{(2)}$$

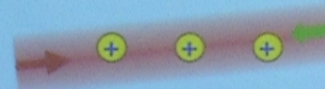
$$S_y = \sigma_y^{(0)} + \sigma_y^{(1)} + \sigma_y^{(2)}$$

Bichromatic (entangles):

$$\omega_r = \omega_0 - (\nu + \epsilon)$$

$$\omega_b = \omega_0 + (\nu + \epsilon)$$

$$\omega_b + \omega_r = 2\omega_0$$



$$S_x^2 = \sigma_x^{(0)}\sigma_x^{(1)} + \sigma_x^{(1)}\sigma_x^{(2)} + \sigma_x^{(0)}\sigma_x^{(2)}$$

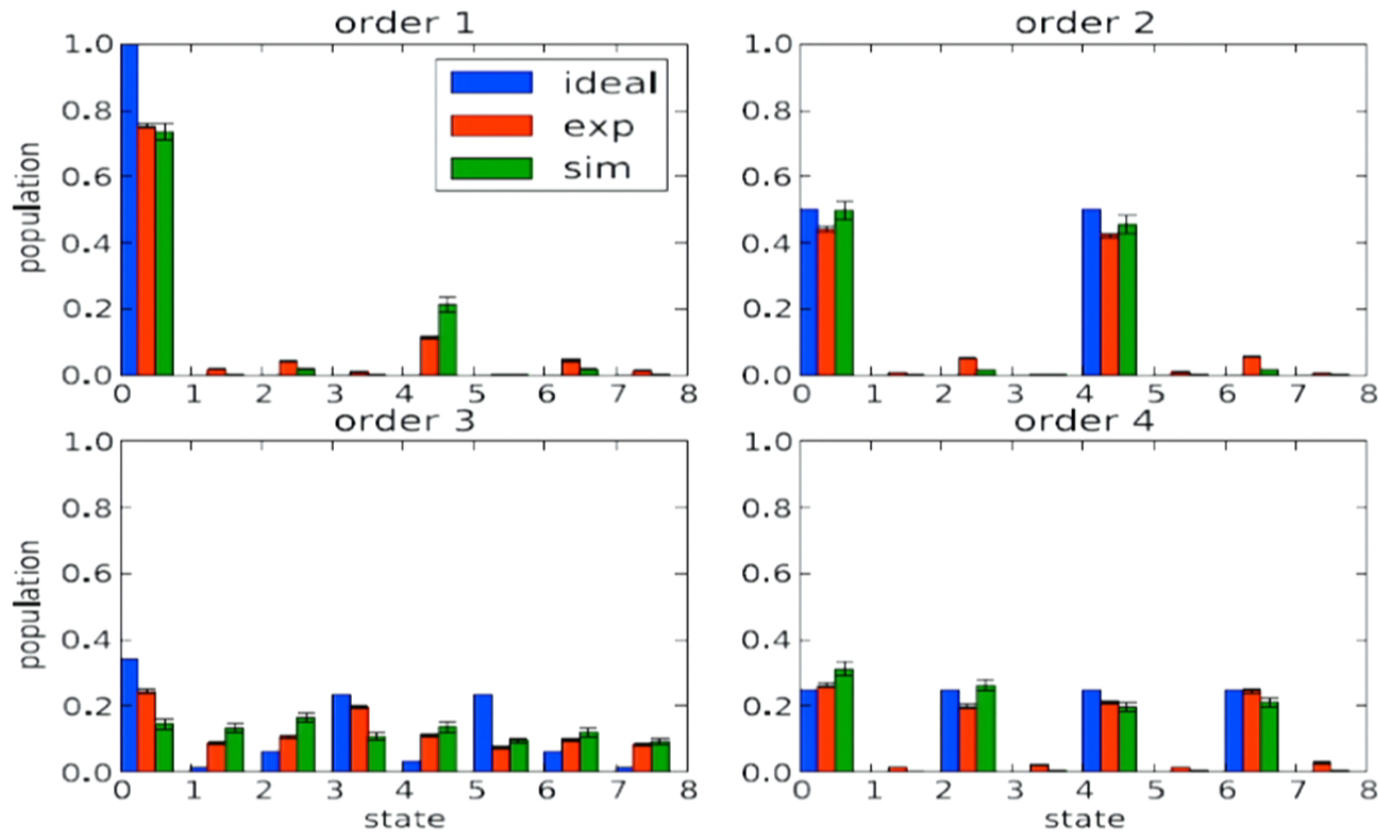
Mølmer-Sørensen gate

Generate rotations about x/y axis

$$U_{S_{x,y}}(\theta) = e^{-i\frac{\theta}{2}S_{x,y}}$$

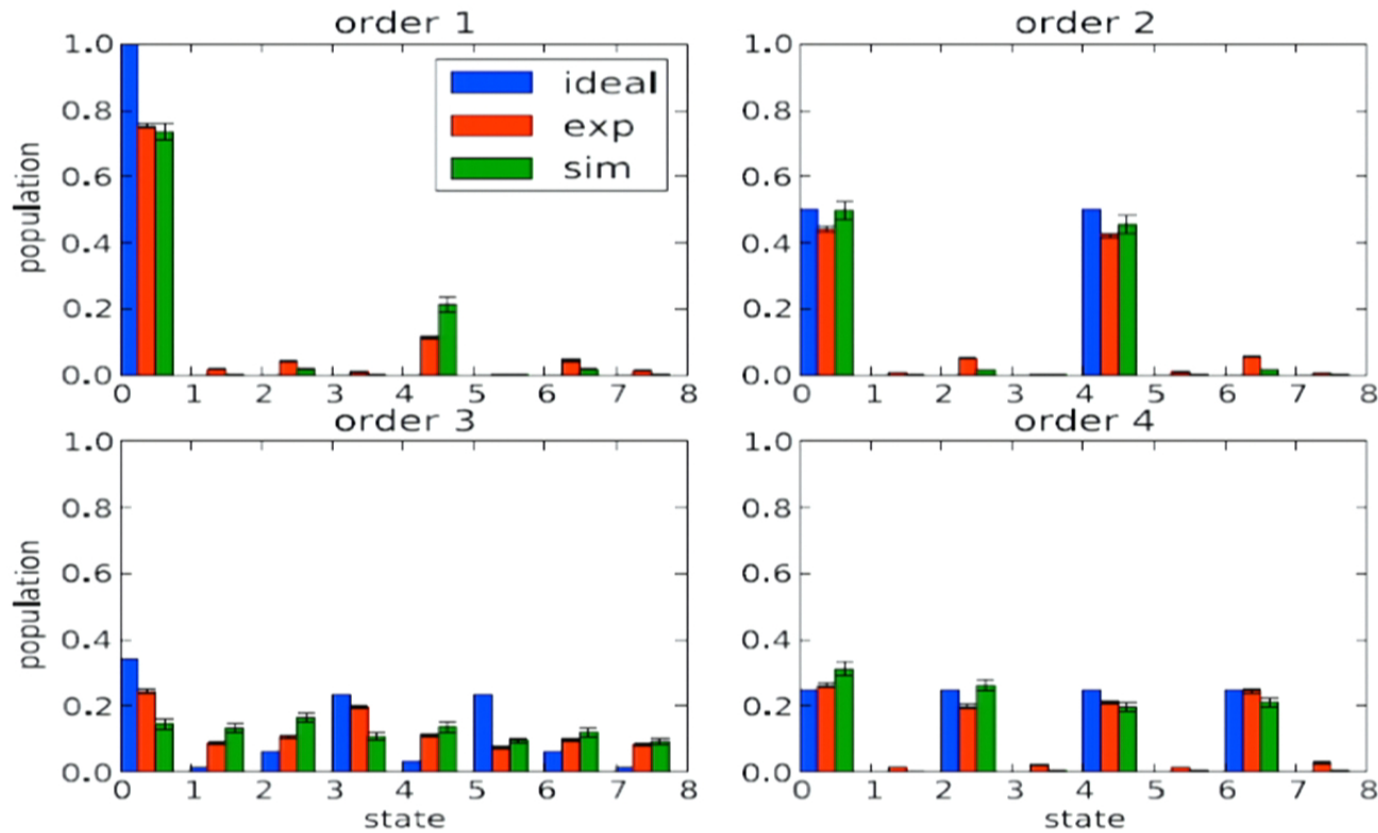
$$MS(\theta) = e^{-i\frac{\theta}{2}S_{x,y}^2}$$

Challenge 4: Predicting Scaling → TIQC-SPICE



No fit parameters!! Entirely based on system calibration

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Montgomery Multiplication

- Classical algorithm used for repeated modular multiplication

Goal

$$c = x b \pmod{N}$$

Montgomery
Reduction Method

$$\begin{aligned}\bar{c} &= \bar{x} \bar{b} R^{-1} \pmod{N} \\ &= \text{MR}(\bar{x} \bar{b})\end{aligned}$$

Montgomery
Rep*

$$\bar{x} = x R$$

Can be done fast!
(half # ops as normal mod mult)

$$R = 2^n$$

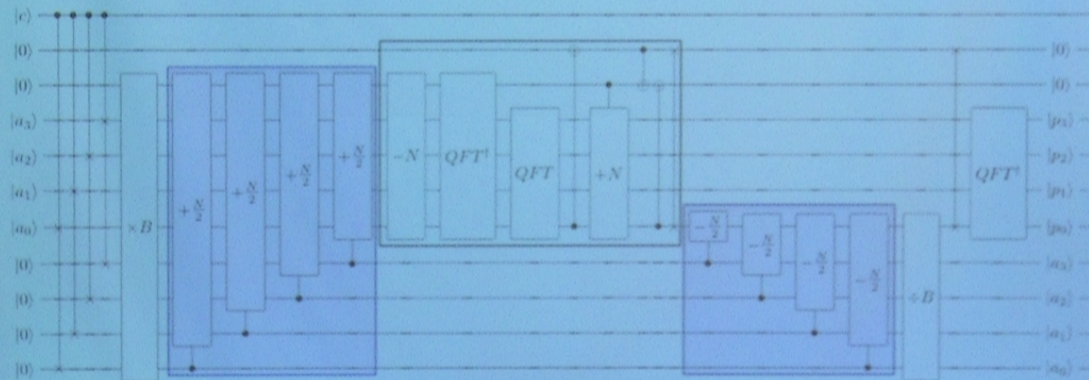
Quantum Fourier Montgomery Multiplication (QFMM)

- The quantum Fourier Montgomery multiplication operator reversibly performs the operation:

$$U_{MP}(b)|x\rangle|0\rangle = |bxR^{-1} \bmod N\rangle|x\rangle$$

- Erase its input:

$$U_{MP}^\dagger(b^{-1})|bxR^{-1} \bmod N\rangle|x\rangle = |bxR^{-1} \bmod N\rangle|0\rangle$$



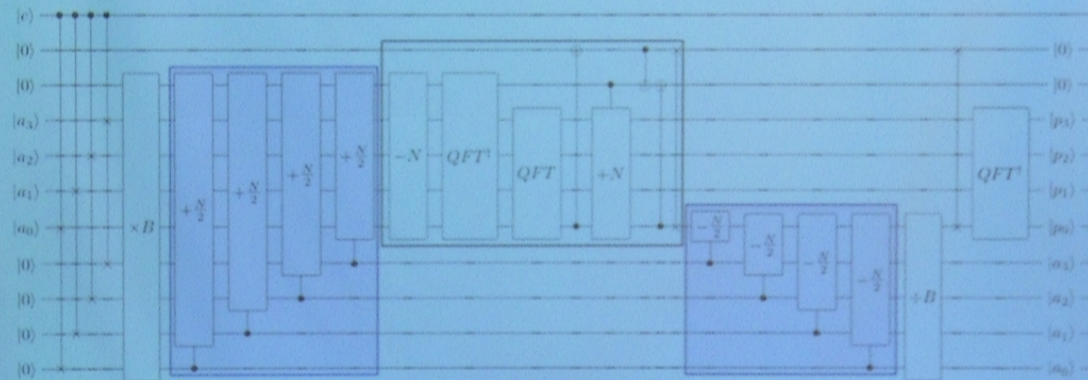
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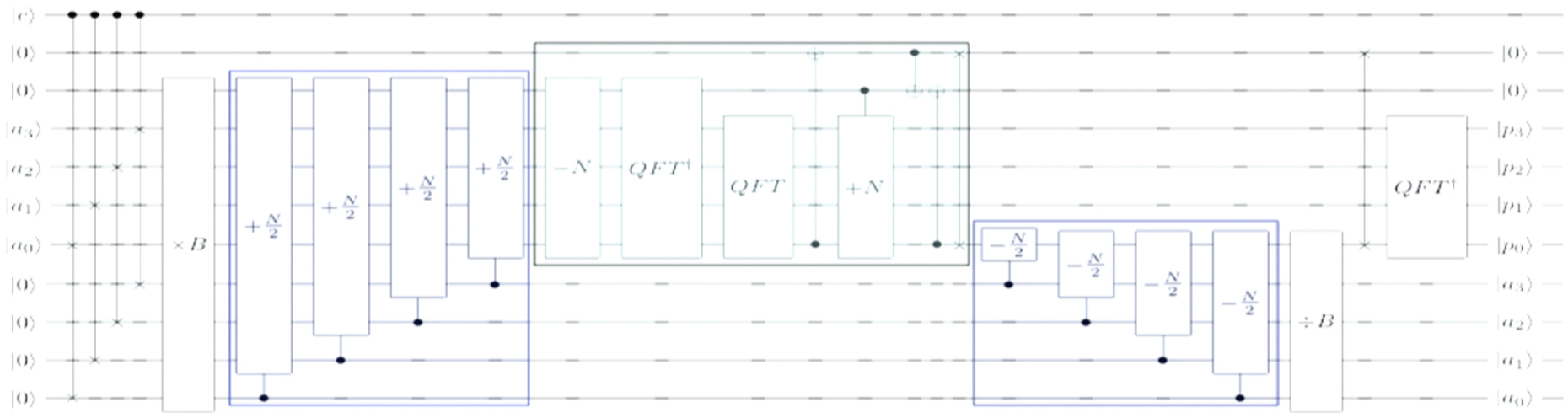
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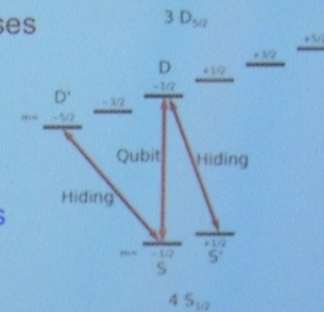


Quantum Factoring with Trapped Ions: a Test of Scalability – Conclusions

Experiment:

Factoring 15 hard and easy cases
> 90% Fidelity results (SSO)

- (1) Fast feed-forward control
- (2) Quantum memory with quantum computation
- (3) Two cascaded, deterministic, 3-qubit q. Fredkin gates
- (4) TIQC-SPICE predictive model of system scalability



Scalability limited by technical issues!