

Title: Exotic topological order from quantum fractal code

Date: Jan 14, 2013 04:00 PM

URL: <http://pirsa.org/13010101>

Abstract: We present a family of three-dimensional local quantum codes with pairs of fractal logical operators. It has two polynomials over finite fields as input parameters which generate fractal shapes of anti-commuting logical operators, and possesses exotic topological order with quantum glassiness which is beyond descriptions of conventional topological field theory. A necessary and sufficient condition for being free from string-like logical operators is obtained under which the model works as marginally self-correcting quantum memory. We also provide theoretical tools to analyze physical and coding properties of translation symmetric stabilizer Hamiltonians.

Exotic topological order from quantum fractal code



Beni Yoshida (Caltech)

Exotic topological order from quantum fractal code



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Exotic topological order from quantum fractal code



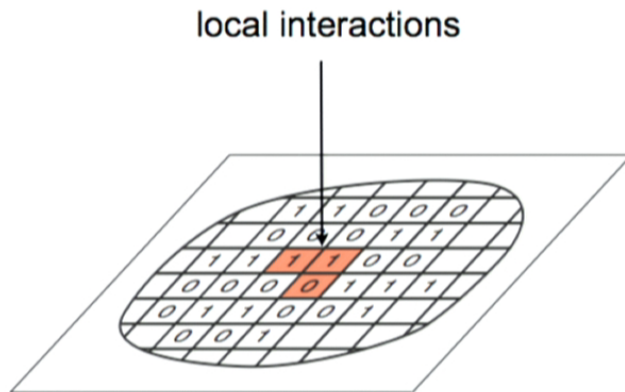
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Question:

- *What is the limit on the **amount** of information that can be **reliably** stored in a finite physical system?*

Local Code bound

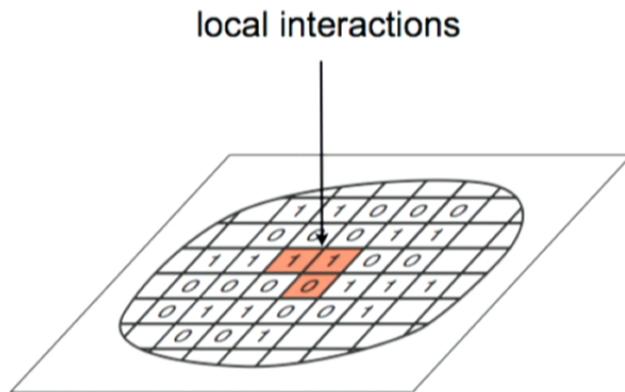
- Encode information into ground states of a **gapped local Hamiltonian** on a **D-dim lattice**



Local Code bound

- Encode information into ground states of a **gapped local Hamiltonian** on a **D-dim lattice**

Local Code Bound Bravyi, Terhal and Poulin (2009)



$$kd^{1/D} \leq O(n)$$

k : number of logical bits **Amount**

d : code distance **Reliability**

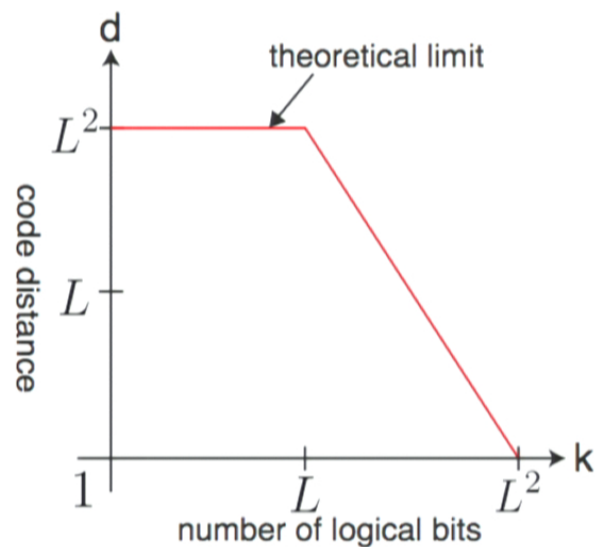
n : total number of spins

Saturation for discrete systems ?

- Previously found systems are far below the bound ...

Bound for $D=2$

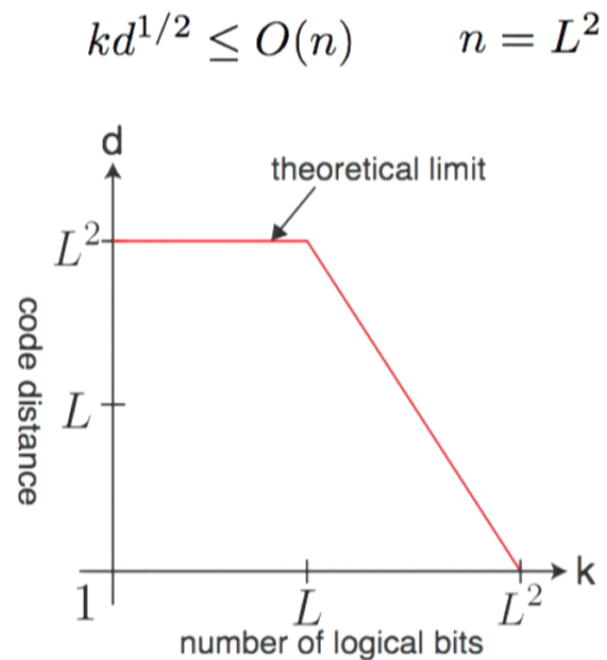
$$kd^{1/2} \leq O(n) \quad n = L^2$$



Saturation for discrete systems ?

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Bound for $D=2$



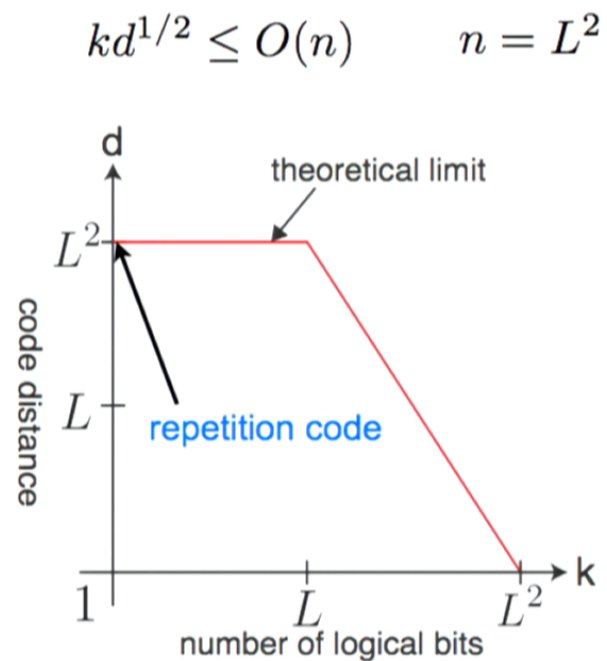
Ex: Repetition code

0	0	0	0	0	0	0
	0	0	0	0	0	0
	0	0	0	0	0	0
	0	0	0	0	0	0
	0	0	0	0	0	0
	0	0	0	0	0	0
$k = 1$						
1	1	1	1	1	1	1
	1	1	1	1	1	1
	1	1	1	1	1	1
	1	1	1	1	1	1
	1	1	1	1	1	1
	1	1	1	1	1	1
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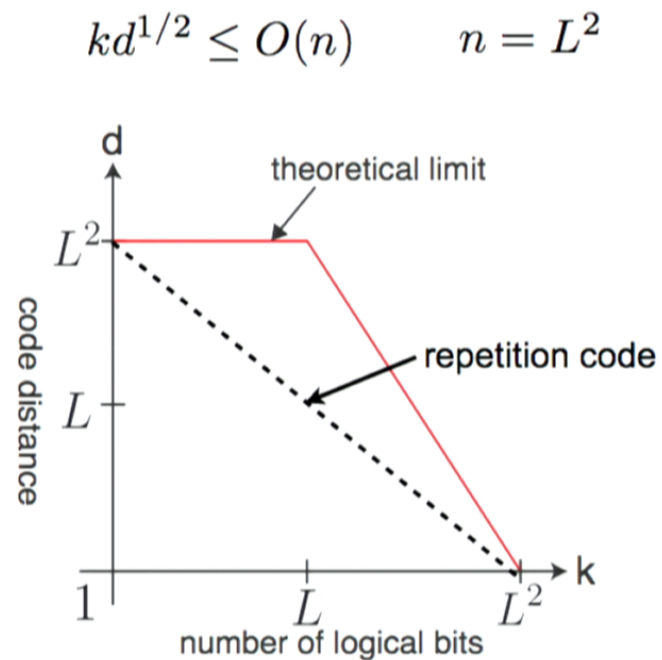
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0	0	0	0	0	0	0
	0	0	0	0	0	0
	0	0	0	0	0	0
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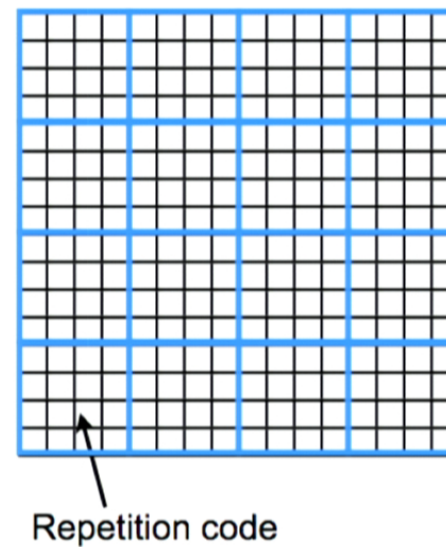
Saturation for discrete systems ?

- Previously found systems are far below the bound ...

Bound for $D=2$



Ex: Copies of repetition codes



Multiple ground states

- Physical realization ? (Newman, Moore 1999)

1	0	0	0	0	0	0	0	0
1	1	0	0	0	0	0	0	0
1	0	1	0	0	0	0	0	0
1	1	1	1	0	0	0	0	0
1	0	0	0	1	0	0	0	0
1	1	0	0	1	1	0	0	0
1	0	1	0	1	0	1	0	0
1	1	1	1	1	1	1	1	1
⋮								
⋱								

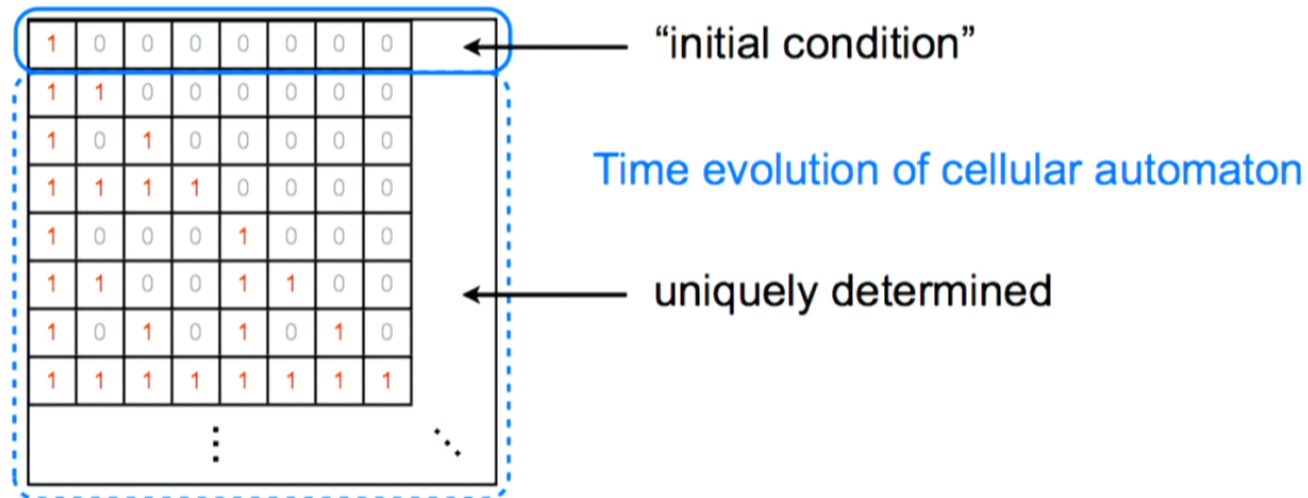
0	1	0	0	0	0	0	0	0
0	1	1	0	0	0	0	0	0
0	1	0	1	0	0	0	0	0
0	1	1	1	1	0	0	0	0
0	1	0	0	0	1	0	0	0
0	1	1	0	0	1	1	0	0
0	1	0	1	0	1	0	1	0
0	1	1	1	1	1	1	1	1
⋮								
⋱								

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Strong breaking of translation symmetries

Multiple ground states

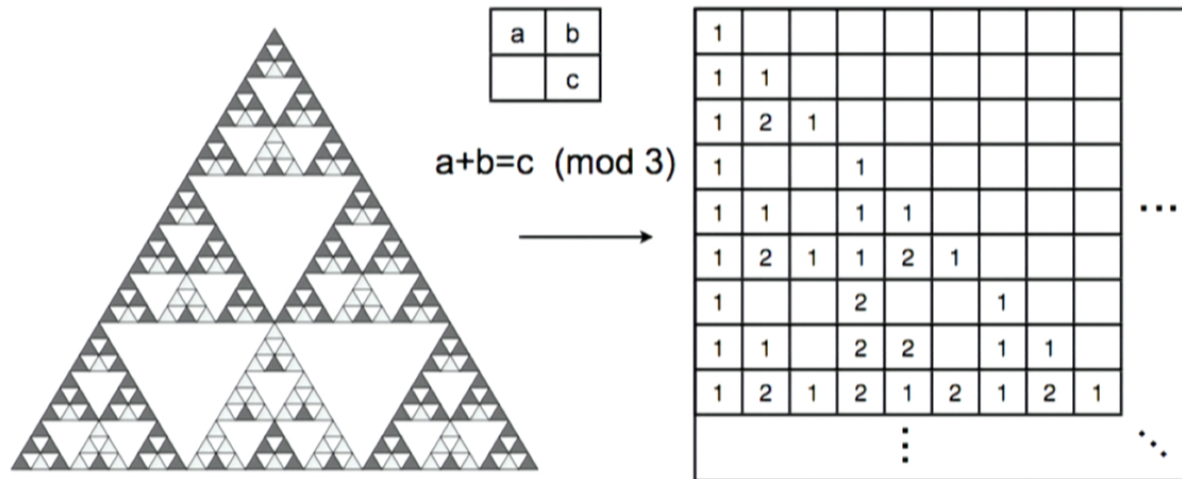
- Physical realization ? (Newman, Moore 1999)



Strong breaking of translation symmetries

The Sierpinski triangle (generalized)

- Physical Realization ?



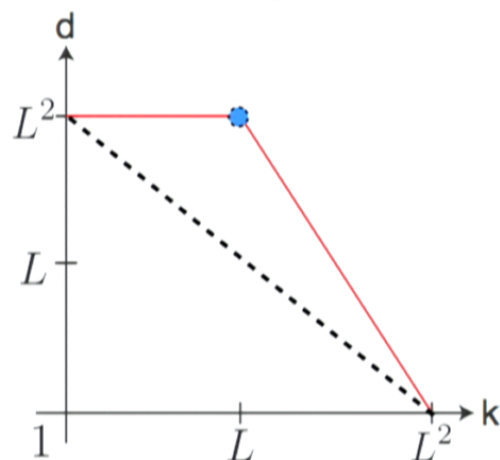
Physically realizable with 3-dimensional spins !

Asymptotic saturation (D=2)

- Sierpinski triangle with **p**-dim spins

Fractal dimension

$$\frac{\log\left(\frac{p(p+1)}{2}\right)}{\log p} \rightarrow 2 \quad \text{for } p \rightarrow \infty.$$

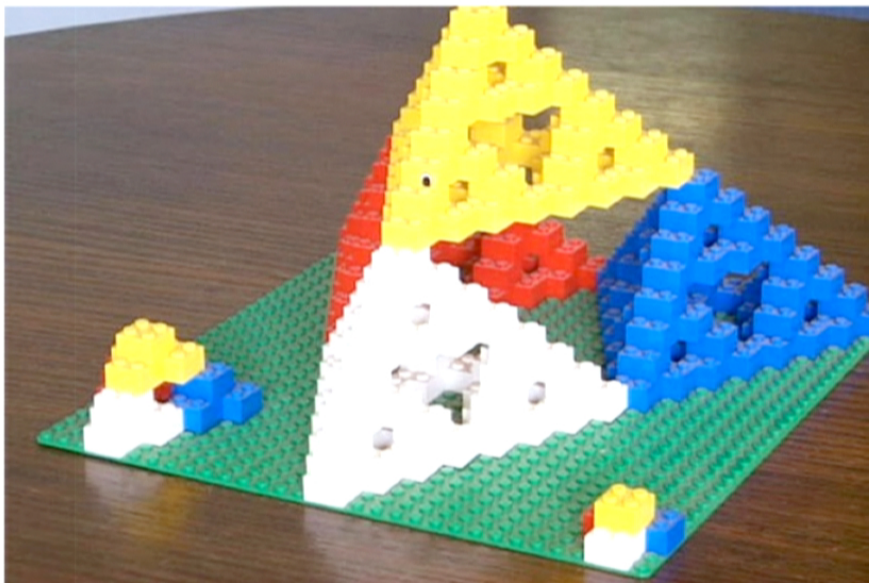


$$k \sim O(L), \quad d \sim O(L^{2-\epsilon})$$

Asymptotically saturate the bound !

Asymptotic saturation ($D > 2$)

- Higher-dimensions ?



Fractal dimension

$$\frac{\log \left(\frac{p(p+1) \cdots (p+D-1)}{D!} \right)}{\log(p)} \xrightarrow[p \rightarrow \infty]{} D$$

Fractal codes saturate the bound for $D > 2$ too !

$$k \sim O(L^{D-1}), \quad d \sim O(L^{D-\epsilon}),$$

Area-like Volume-like

Quantum Fractal Code ?

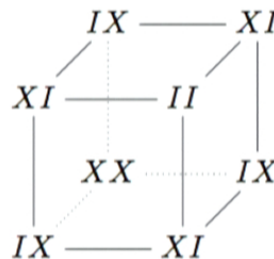
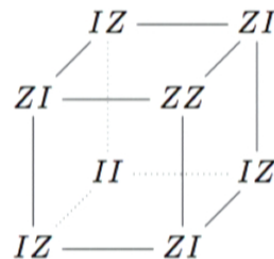
Classical bound

$$kd^{1/D} \leq O(n) \longleftarrow \text{Classical fractal code}$$

Quantum bound

$$kd^{\frac{1}{D-1}} \leq O(n) \longleftarrow \text{Quantum fractal code ?}$$

- The Cubic code has pairs of fractal logical operators. (Haah 2011)



general construction ?

Plan of Part 2

- Translation symmetric stabilizer Hamiltonians
 - via **polynomials over finite fields** (a new theoretical framework)
- Quantum fractal code
 - **fractal** logical operators
- **Properties** of quantum fractal code
 - ground states (exotic topo order, entanglement entropy)
 - quasi-particle excitations
 - code distances (lower bounds)

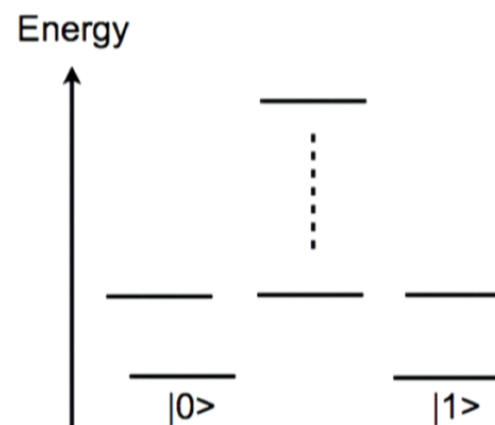
Stabilizer Hamiltonians (Gottesman 1996)

- Encode qubits into degenerate ground states with **mass gap**

$$H = - \sum_j S_j, \quad [S_j, S_{j'}] = 0$$

$$S_j |\psi\rangle = |\psi\rangle$$

Pauli Operators



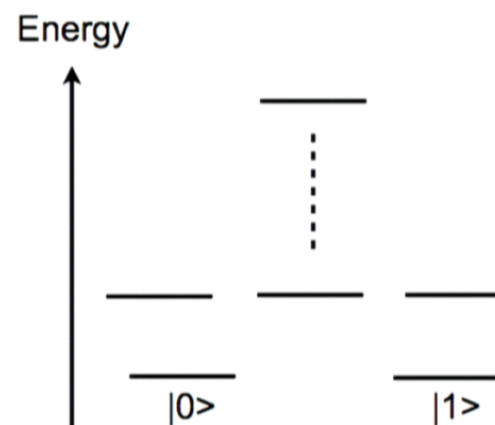
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Pauli Operators



Pauli operators by polynomials over finite fields

Pauli operators by polynomials (1 dim)

$$f = \sum_{j=-\infty}^{\infty} c_j x^j, \quad c_j = 0, 1. \quad \text{-----} \quad \begin{matrix} -2 & -1 & 0 & 1 & 2 \\ \circ & \circ & \circ & \circ & \circ \end{matrix} \quad \text{-----}$$

$$Z(f) = \prod_{j=-\infty}^{\infty} Z_j^{c_j}, \quad X(f) = \prod_{j=-\infty}^{\infty} X_j^{c_j} \quad (\text{eg}) \quad f = 1 + x + x^2$$

$$Z(f) = Z_0 Z_1 Z_2$$

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Translations of Pauli operators

$$f = 1 + x + x^2 \rightarrow \boxed{x}f = x + x^2 + x^3$$

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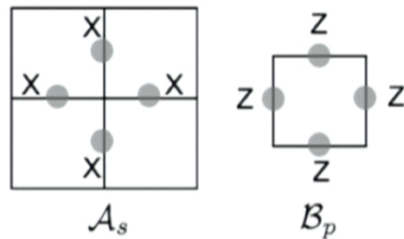
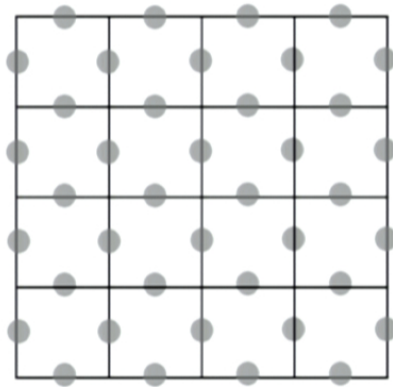
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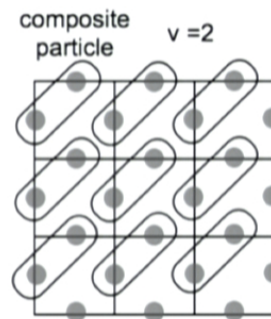
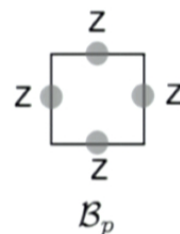
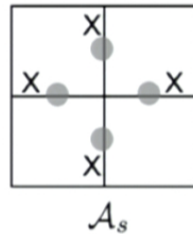
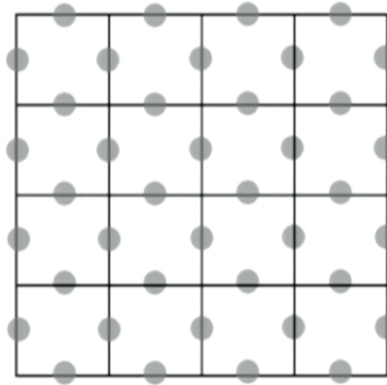
The Toric code via polynomials

The Toric code (qubits on edges)



The Toric code via polynomials

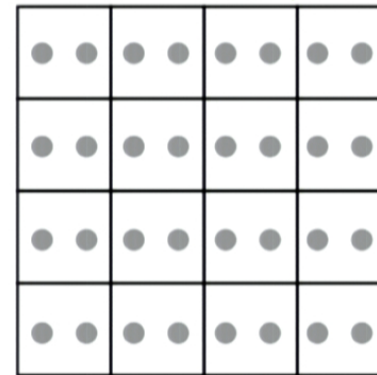
The Toric code (qubits on edges)



local unitary

The Toric code (square lattice)

(Yoshida 2010)



The Toric code via polynomials

	-1	0	1
-1			
0		$Z_A Z_B$	Z_A
1		Z_B	

The Toric code via polynomials

	-1	0	1
-1			
0		$Z_A Z_B$	Z_A
1		Z_B	

$$Z \begin{pmatrix} 1 + x \\ 1 + y \end{pmatrix}$$

	-1	0	1
-1		X_A	
0	X_B	$X_A X_B$	
1			

$$X \begin{pmatrix} 1 + y^{-1} \\ 1 + x^{-1} \end{pmatrix}$$

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Theory for model with duality

$$\begin{array}{ccc}
 & \text{duality} & \\
 Z \left(\begin{array}{c} \boxed{1+x} \\ \boxed{1+y} \end{array} \right) & \begin{array}{c} \swarrow \quad \searrow \\ \nwarrow \quad \nearrow \end{array} & X \left(\begin{array}{c} \boxed{1+y^{-1}} \\ \boxed{1+x^{-1}} \end{array} \right) \\
 & \boxed{x \rightarrow x^{-1}, y \rightarrow y^{-1}} &
 \end{array}$$

Theory for model with duality

- This duality is required for commutation of stabilizer generators

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- Model with duality

$$Z \begin{pmatrix} \alpha \\ \beta \end{pmatrix} \qquad X \begin{pmatrix} \bar{\beta} \\ \bar{\alpha} \end{pmatrix}$$

“theory” via algebra on polynomials over finite fields

See a forthcoming paper for discussion on codes with duality.

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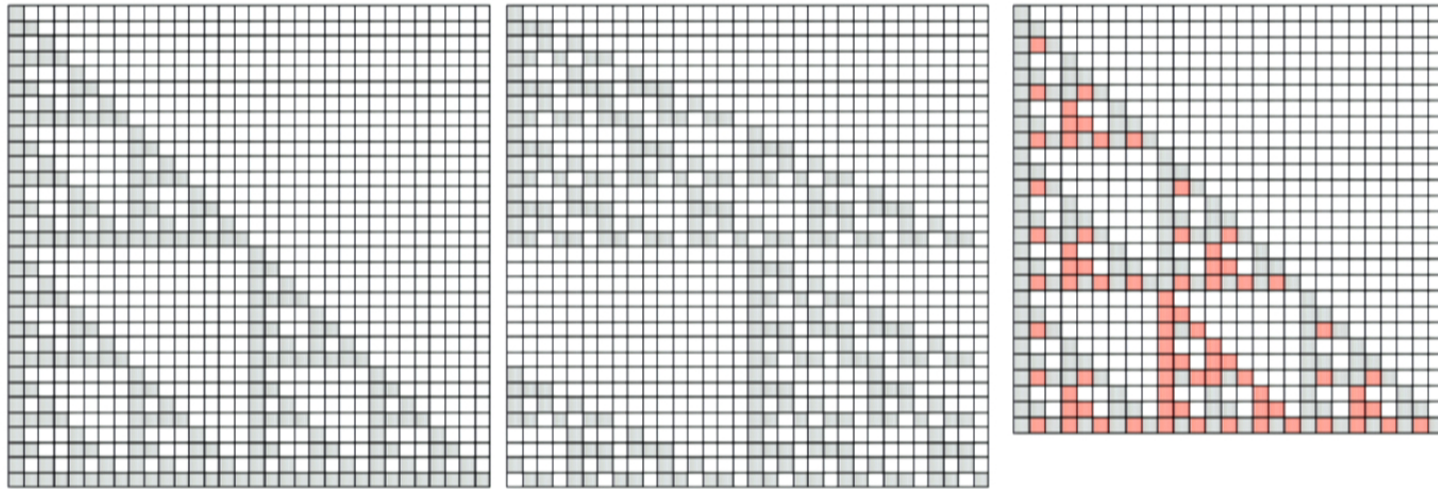
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Fractal geometries via polynomials



Fractal geometries via polynomials

- $f = 1 + x$ over $\mathbb{F}_2$

Sierpinski triangle

powers of “f”

$$f^0 = 1$$

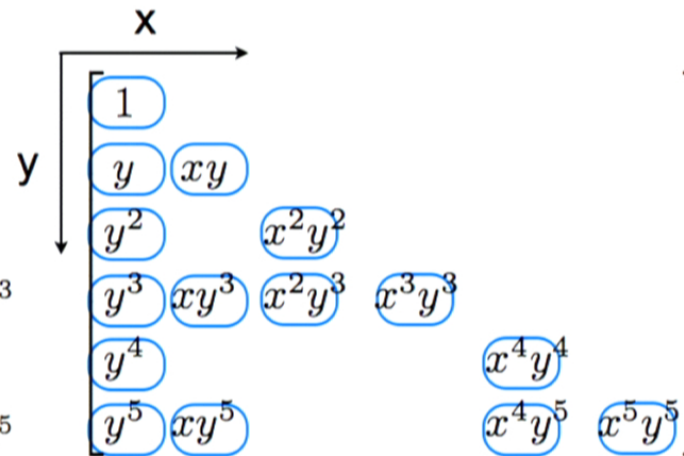
$$f^1 = 1 + x$$

$$f^2 = 1 + x^2$$

$$f^3 = 1 + x + x^2 + x^3$$

$$f^4 = 1 + x^4$$

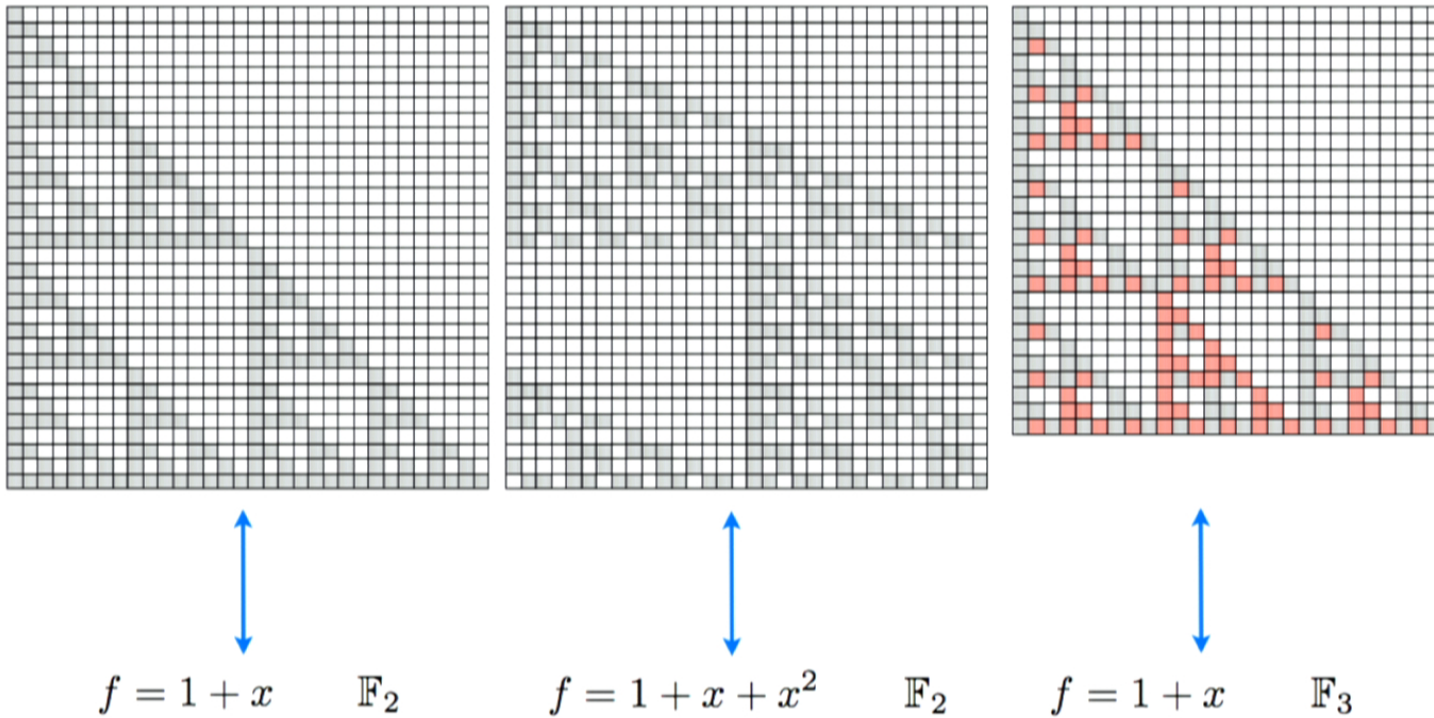
$$f^5 = 1 + x + x^4 + x^5$$



$$f(x, y) = 1 + fy + f^2y^2 + f^3y^3 + \dots$$

Quantum Fractal Code (3 dim)

Fractal geometries via polynomials



Quantum Fractal Code (3 dim)

- Two parameters $a(x)$ $b(x)$ over $\mathbb{F}_p$
- Interaction terms

$$Z \begin{pmatrix} 1 + a(x)y \\ 1 + b(x)z \end{pmatrix}, \quad X \begin{pmatrix} 1 + \bar{b}(x)\bar{z} \\ 1 + \bar{a}(x)\bar{y} \end{pmatrix}$$

$x \rightarrow x^{-1}, y \rightarrow y^{-1}, z \rightarrow z^{-1}$

- Fractal logical operators

$$Z \begin{pmatrix} 0 \\ x^i \mathbf{a}(x, y) \end{pmatrix} \quad Z \begin{pmatrix} x^i \mathbf{b}(x, z) \\ 0 \end{pmatrix}$$

$$X \begin{pmatrix} x^i \bar{\mathbf{a}}(x, y) \\ 0 \end{pmatrix} \quad X \begin{pmatrix} 0 \\ x^i \bar{\mathbf{b}}(x, z) \end{pmatrix}$$

Quantum Fractal Code (3 dim)

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- Fractal logical operators

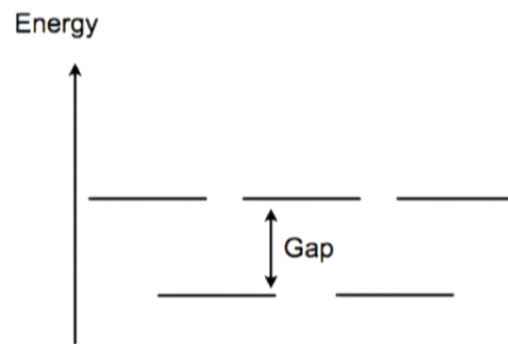
$$\begin{array}{cc} Z \begin{pmatrix} 0 \\ x^i \mathbf{a}(x, y) \end{pmatrix} \xleftarrow{a(x)} & Z \begin{pmatrix} x^i \mathbf{b}(x, z) \\ 0 \end{pmatrix} \xleftarrow{b(x)} \\ X \begin{pmatrix} x^i \bar{\mathbf{a}}(x, y) \\ 0 \end{pmatrix} \xleftarrow{\bar{a}(x)} & X \begin{pmatrix} 0 \\ x^i \bar{\mathbf{b}}(x, z) \end{pmatrix} \xleftarrow{\bar{b}(x)} \end{array}$$

Plan of Part 2

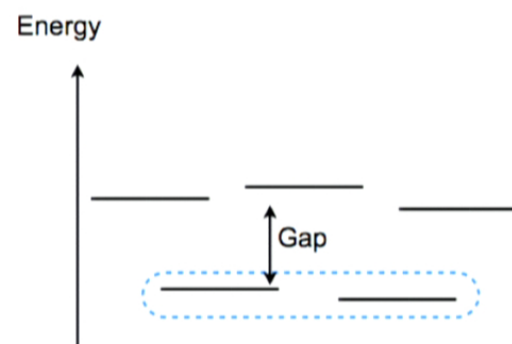
- Translation symmetric stabilizer Hamiltonians
 - via **polynomials over finite fields** (a new theoretical framework)
- Quantum fractal code
 - **fractal** logical operators
- **Properties** of quantum fractal code
 - **ground states** (exotic topo order, entanglement entropy)
 - quasi-particle excitations
 - code distances (lower bounds)

Topological Order (Stability against perturbation)

Def: stability against perturbations

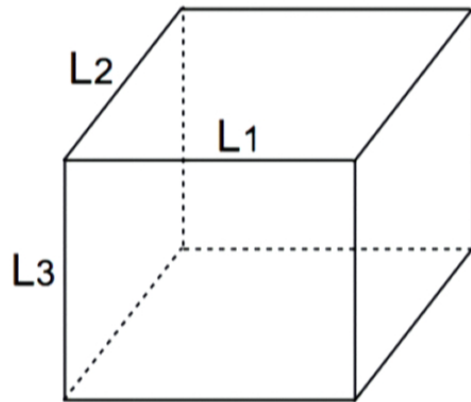


Original System : H



Perturbed System : $H+V$

Number of ground states

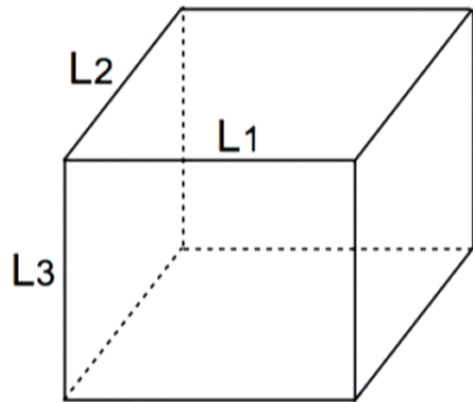


Number of solutions to

$$a(x)^{L_2} \gamma(x) = b(x)^{L_3} \gamma(x) = \gamma(x), \quad x^{L_1} = 1$$

$$= \frac{1}{2} \text{ Number of ground states}$$

Number of ground states



Number of solutions to

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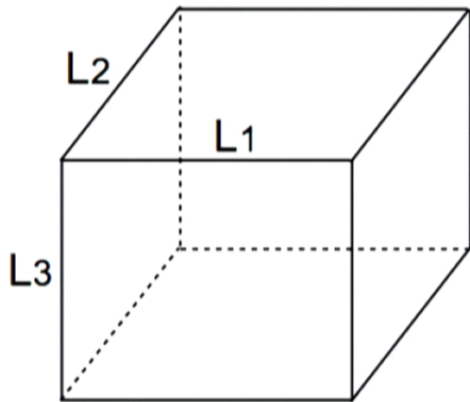
Discrete scale symmetries

$$k(pL_1, pL_2, pL_3) = pk(L_1, L_2, L_3)$$

→ topological order with **extensive degeneracy**

→ **beyond** topological field theory

Number of ground states



Number of solutions to

$$a(x)^{L_2} \gamma(x) = b(x)^{L_3} \gamma(x) = \gamma(x), \quad x^{L_1} = 1$$

$$= \frac{1}{2} \text{ Number of ground states}$$

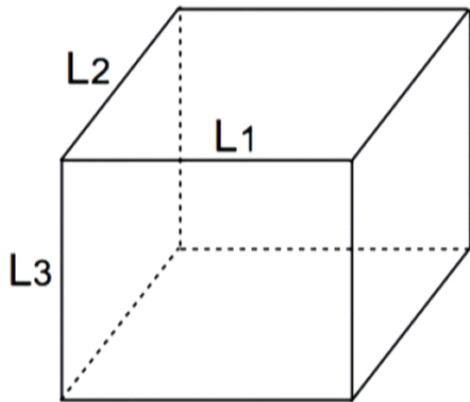
Discrete scale symmetries

$$k(pL_1, pL_2, pL_3) = pk(L_1, L_2, L_3)$$

→ topological order with **extensive degeneracy**

→ **beyond** topological field theory

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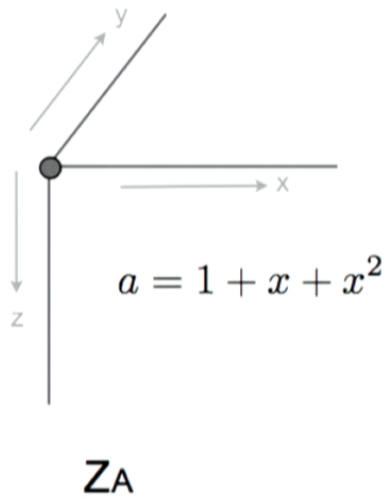
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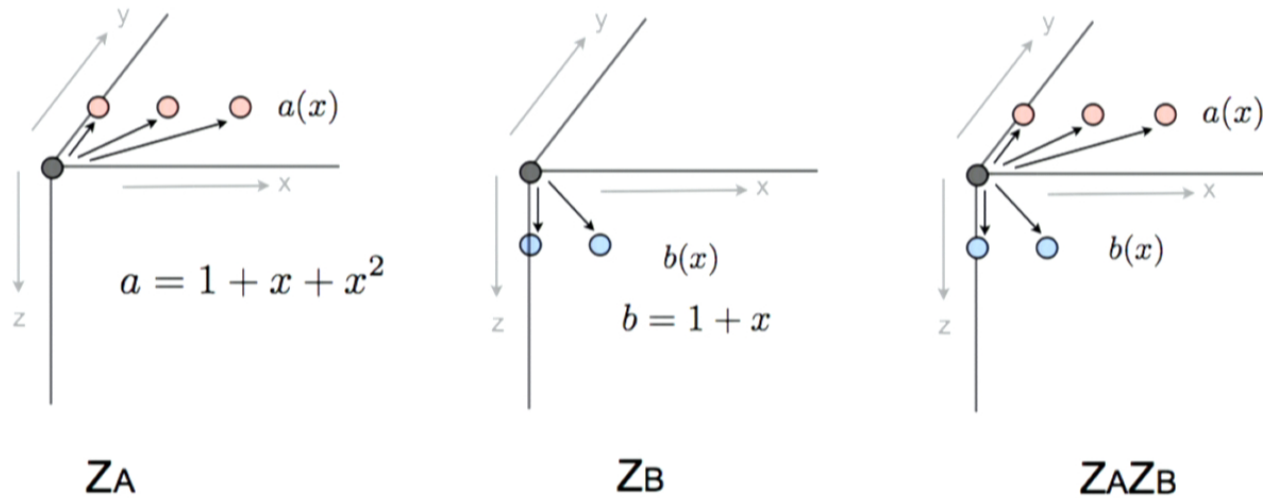
Quasi-particle excitations in quantum fractals



$$a = 1 + x + x^2$$

$$b = 1 + x$$

Quasi-particle excitations in quantum fractals



- Propagations of quasi-particles are highly constrained (not free).
- Propagations can be discussed concisely via polynomials.

No string logical operators

No string logical operators
(no point-like quasi-particle excitations)



Marginally self-correcting (Logarithmic energy barrier)

(Bravyi, Haah 2012)

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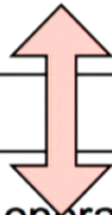
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No string logical operators

(c_a , c_b) does not exist :

$$a(x)^{c_a} = b(x)^{c_b}$$



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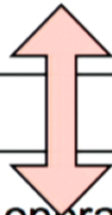
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Plan of Part 2

- Translation symmetric stabilizer Hamiltonians
 - via **polynomials over finite fields** (a new theoretical framework)
- Quantum fractal code
 - **fractal** logical operators
- **Properties** of quantum fractal code
 - ground states (exotic topo order, entanglement entropy)
 - quasi-particle excitations
 - **code distances** (lower bounds)

Polynomial decomposition problem

Theorem (lower bound)

Consider all the tensors $C_{ij}^{(\ell)}$ satisfying the following equation:

$$\gamma b^\ell = \sum_{ij} x^i a^j C_{ij}^{(\ell)}, \quad \text{for all } \ell.$$

and consider all the tensors $D_{i\ell}^{(j)}$ satisfying the following equation:

$$\gamma a^\ell = \sum_{ij} x^i b^j D_{ij}^{(\ell)}, \quad \text{for all } \ell.$$

$$\text{Weight}(C_{ij}^{(\ell)}) \geq W \quad \text{Weight}(D_{i\ell}^{(j)}) \geq W \quad \text{for } \gamma \neq 0 \quad \Rightarrow \quad d \geq W.$$

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- Reduced to some mathematical problem
- Still difficult to handle this problem

Summary

Quantum fractal code

$$Z \begin{pmatrix} 1 + a(x)y \\ 1 + b(x)z \end{pmatrix}, \quad X \begin{pmatrix} 1 + \bar{b}(x)\bar{z} \\ 1 + \bar{a}(x)\bar{y} \end{pmatrix}$$
$$x \rightarrow x^{-1}, y \rightarrow y^{-1}, z \rightarrow z^{-1}$$

- Is this universal ?
- Cubic code as quantum fractal code ?

Coding properties ?

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