

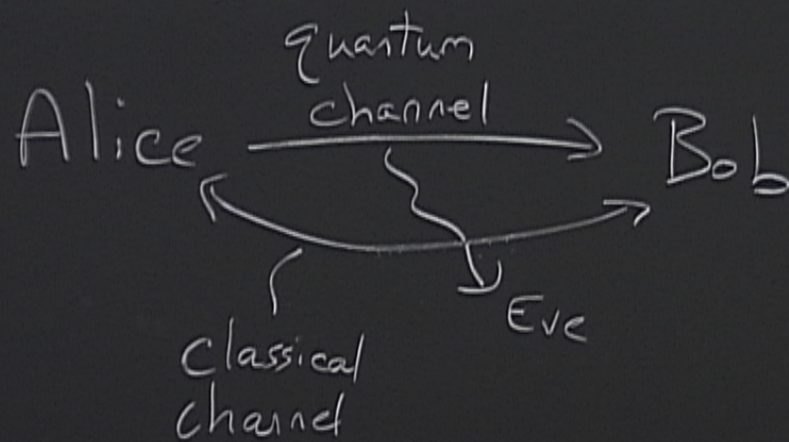
Title: Quantum Information (Review) - Lecture 14

Date: Mar 02, 2012 10:15 AM

URL: <http://www.pirsa.org/12030001>

Abstract:





BB84

- ① Alice generates N random bit pairs (a_i, r_i)
- ② Alice prepares qubits determined by each pair & sends the qubits to Bob

a_i	r_i	state
0	0	$ 0\rangle$
0	1	$ 1\rangle$

a_i	r_i	state
1	0	$ +\rangle$
1	1	$ -\rangle$

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- ③ Bob chooses N random bases b_i & measures qubit i in basis b_i , getting result s_i
- ④ Alice & Bob announce a_i & b_i over classical channel (but not r_i & s_i).
- If $a_i = b_i$, keep bit i
- If $a_i \neq b_i$, throw away bit i
- Have about $\sim N/2$ bits left.

Have about 10^{12} bits left.

Suppose Eve measures qubits
as they are sent.

E.g.: She might choose a basis
& measure in that basis, then send
state on to Bob

But Eve might measure in wrong basis

If Alice sends $|0\rangle$, but Eve chooses
 $|\pm\rangle$ basis Eve gets $|+\rangle$

But Eve might measure in wrong basis

If Alice sends $|0\rangle$, but Eve chooses $|\pm\rangle$ basis. Eve gets $|+\rangle$, sends that to Bob. If Bob measures in $|0\rangle/|1\rangle$ basis, Bob's outcome is random & is wrong 50% of the time

Eve's attempt to learn about the key introduces errors.

③ Bob chooses N random bases b_i & measures qubit i

⑤ Alice & Bob choose a random

is wrong 50% of the time

③ Bob chooses N random bases b_i & measures qubit i in basis b_i , getting result s_i

④ Alice & Bob announce a_i & b_i over classical channel (but not r_i & s_i).

If $a_i = b_i$, keep bit i

If $a_i \neq b_i$, throw away bit i

Have about $\sim N/2$ bits left.

⑤ Alice & Bob choose a random subset of bits & announce r_i & s_i for those bits.

They calculate error rate & abort if it is too high ($\geq 20\%$).

⑥ Error correction: Alice & Bob compute parities of their bits according to a classical ECC & announce parities to do error correction. Must discard 1 bit per parity announced.

⑦ Privacy amplification: Take random subsets of remaining bit strings & compute parities of those subsets. The final key is the value of the parities.

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$\exists$ security proofs that
say that either Eve is
caught with probability very
close to 1 or she has
essentially no information about the key.

Authenticated classical channel $\leftarrow$ Can be created using
from Alice to Bob. $\sim \log N$ bits of shared secret key.

Photon loss over long distances
Quantum repeaters can extend distance

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from Alice

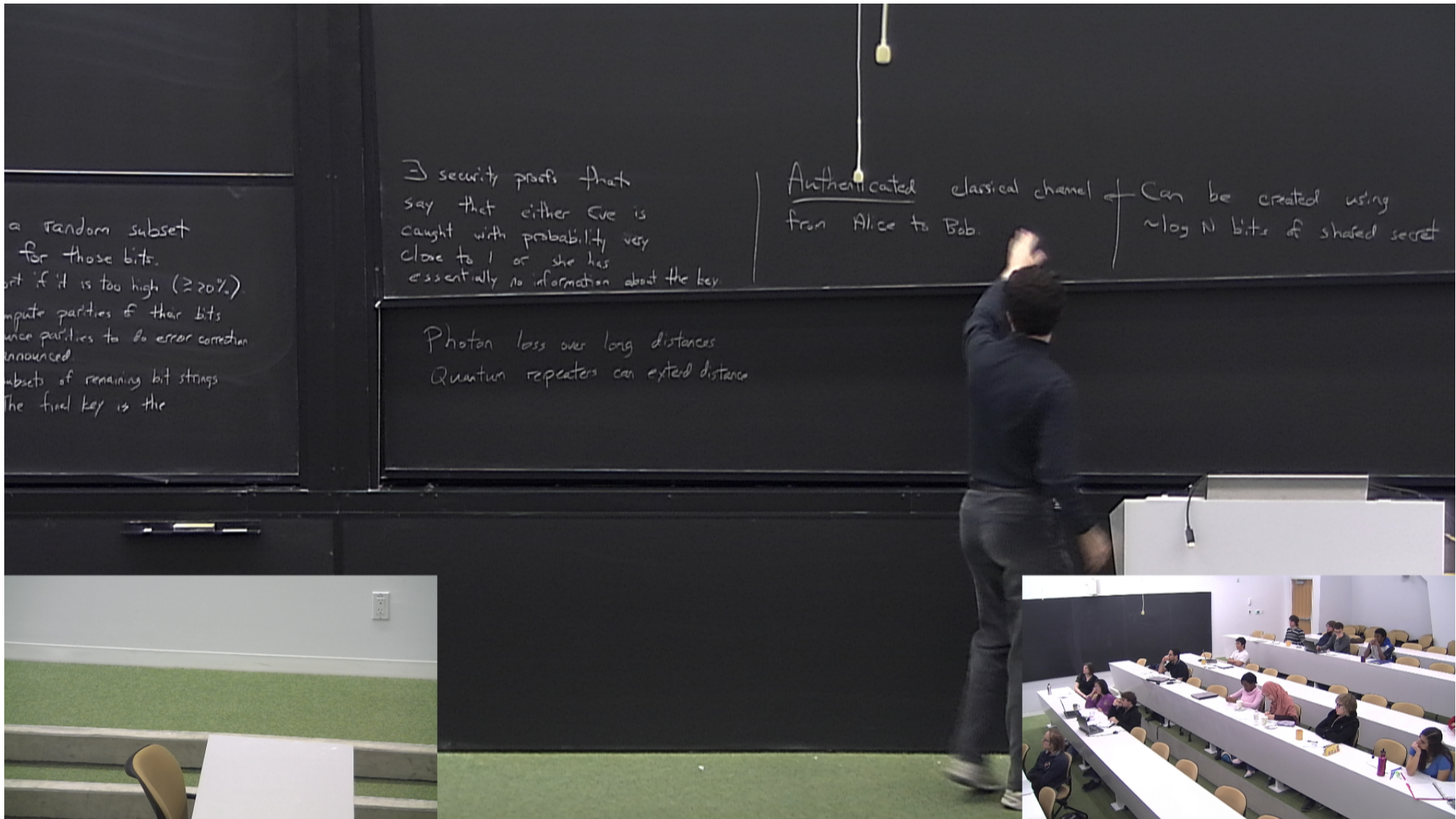
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