

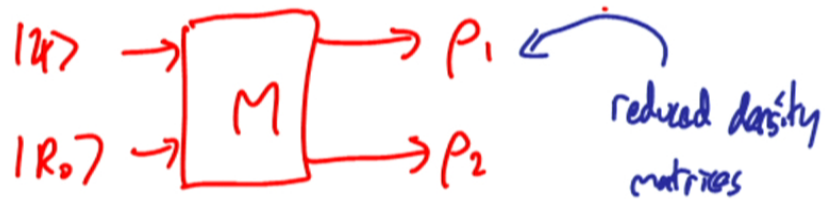
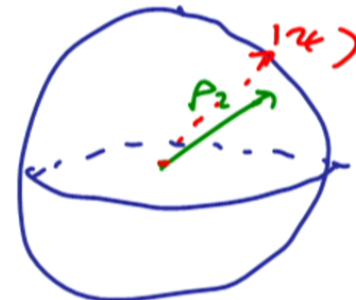
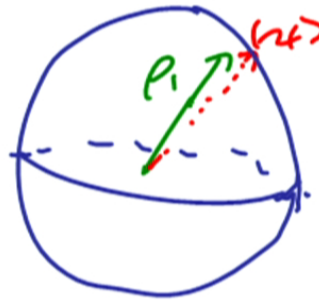
Title: Quantum Theory - Lecture 14

Date: Sep 28, 2011 03:15 PM

URL: <http://pirsa.org/11090056>

Abstract:

Extending the no-cloning theorem II What about approximate quantum cloning?



$$\text{Tr}_2(\rho) = \rho_1, \quad \text{Tr}_1(\rho) = \rho_2$$

ρ_1, ρ_2 close to $|\psi\rangle\langle\psi|$.

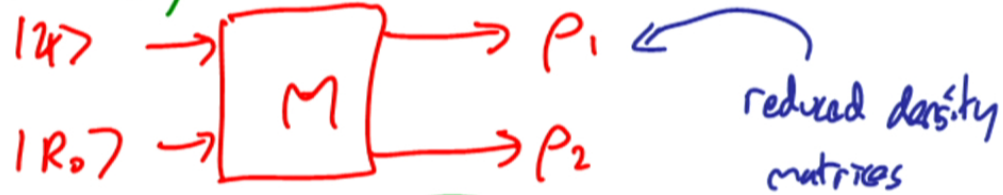
We need a measure of closeness: a natural one is the fidelity $\langle\psi|\rho_i|\psi\rangle$ which gives the probability of outcome "yes" if we measure $|\psi\rangle\langle\psi|$ on ρ_i .

$$\begin{array}{ccc}
 |\psi\rangle & |R_0\rangle & |M\rangle \longrightarrow \underline{|\psi'\rangle} \\
 \text{qubit} & \text{qubit} & \mathbb{C}^d \\
 & & \text{entangled state} \\
 & & \mathbb{C}^2 \otimes \mathbb{C}^2 \otimes \mathbb{C}^d \\
 \rho_1 = \text{Tr}_{H_2 \otimes H_M} (|\psi\rangle\langle\psi|) & & H_1 \otimes H_2 \otimes H_M \\
 \rho_2 = \text{Tr}_{H_1 \otimes H_M} (|\psi'\rangle\langle\psi'|) & &
 \end{array}$$

If you measure $P_{\psi} = |\psi\rangle\langle\psi|$ on ρ_1 , what's the prob you get outcome 1?

$$\text{tr}(P_{\psi}\rho_1) = \langle\psi|\rho_1|\psi\rangle$$

Effectively:



$$| \psi \rangle | R_0 \rangle | M \rangle \rightarrow | \psi' \rangle$$

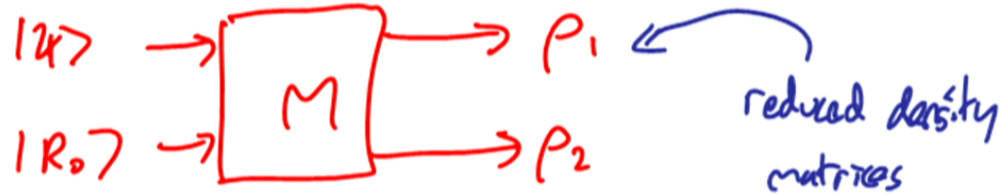
Full picture:

State of all
3, maybe
entangled.

Could we arrange that $|\langle \psi | \rho_1 | \psi \rangle| \geq 1 - \epsilon$, $|\langle \psi | \rho_2 | \psi \rangle| \geq 1 - \epsilon$

for all possible inputs $| \psi \rangle$?

And can we arrange this for any $\epsilon > 0$?



Could we arrange that $|\langle \psi | P_1 | \psi \rangle| \geq 1 - \epsilon$, $|\langle \psi | P_2 | \psi \rangle| \geq 1 - \epsilon$

for all possible inputs $|\psi\rangle$?

And can we arrange this for any $\epsilon > 0$?

The no-signalling argument - remembers Herbert's F.L.A.S.H. scheme - implies not. We can distinguish the mixtures $\frac{1}{2}(|\uparrow\rangle|\uparrow\rangle\langle\uparrow|\langle\uparrow| + |\downarrow\rangle|\downarrow\rangle\langle\downarrow|\langle\downarrow|)$
 $\frac{1}{2}(|\rightarrow\rangle|\rightarrow\rangle\langle\rightarrow|\langle\rightarrow| + |\leftarrow\rangle|\leftarrow\rangle\langle\leftarrow|\langle\leftarrow|)$

So by continuity we must be able to distinguish mixtures of sufficiently close approximations.

If you measure $P_1 = |\psi\rangle\langle\psi|$ on ρ_1 , what's the prob you get outcome 1?

$$F = \text{Tr}(P_1 \rho_1) = \langle\psi|\rho_1|\psi\rangle$$

$$\sigma_1 = \frac{1}{2} (|\uparrow\rangle\langle\uparrow| + |\downarrow\rangle\langle\downarrow|)$$

$$\sigma_2 = \frac{1}{2} (|\rightarrow\rangle\langle\rightarrow| + |\leftarrow\rangle\langle\leftarrow|)$$

If you measure $P_\psi = |\psi\rangle\langle\psi|$ on ρ_1 , what's the prob you get outcome 1?

$$F = \text{Tr}(P_\psi \rho_1) = \langle\psi|\rho_1|\psi\rangle$$

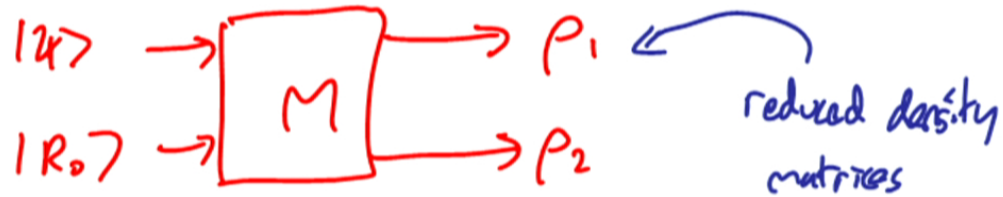
approx closer
to values

$$G_1^\epsilon \approx G_1 = \frac{1}{2} (|\uparrow\rangle\langle\uparrow| + |\downarrow\rangle\langle\downarrow|)$$

$$G_2^\epsilon \approx G_2 = \frac{1}{2} (|\rightarrow\rangle\langle\rightarrow| + |\leftarrow\rangle\langle\leftarrow|)$$

$$\text{so } G_1^\epsilon \neq G_2^\epsilon \quad G_1 \neq G_2$$

for small ϵ .



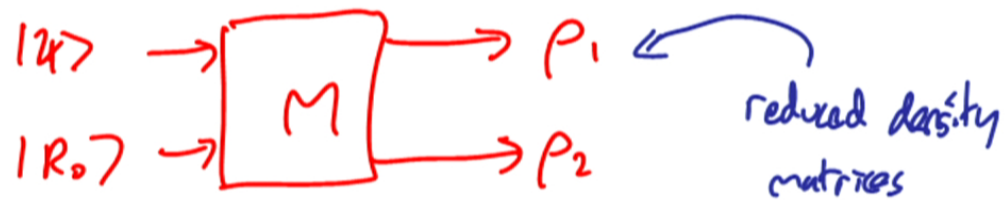
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So by continuity we must be able to distinguish mixtures of sufficiently close approximations.



OK, we can't achieve $\langle \psi | \rho_1 | \psi \rangle = \langle \psi | \rho_2 | \psi \rangle = 1$.

Nor can we achieve $\langle \psi | \rho_1 | \psi \rangle = \langle \psi | \rho_2 | \psi \rangle = (1 - \epsilon)$ for arbitrary $\epsilon > 0$.

Can we achieve anything that can sensibly be called approximate cloning?

$$\langle \psi | \rho_1 | \psi \rangle = \langle \psi | \rho_2 | \psi \rangle = \eta \quad \text{for } \eta > 0? \\ \eta > \frac{1}{2}? \\ \dots$$

for small ϵ ,

Measure S_z ? $\begin{matrix} |\uparrow\rangle \\ |\downarrow\rangle \end{matrix} \rightarrow \begin{matrix} |\uparrow\rangle|\uparrow\rangle \\ |\downarrow\rangle|\downarrow\rangle \end{matrix}$

--- Gives fidelity 1 for input $|\uparrow\rangle$ or $|\downarrow\rangle$
 $\frac{1}{2}$ for input $|\rightarrow\rangle$ or $|\leftarrow\rangle$

State-dependent fidelity



Simple (trivial?) cloning strategies ① random measurement.

Initial state $|a\rangle = \frac{1}{2}(1 + a \cdot \underline{e}) \equiv |a\rangle$

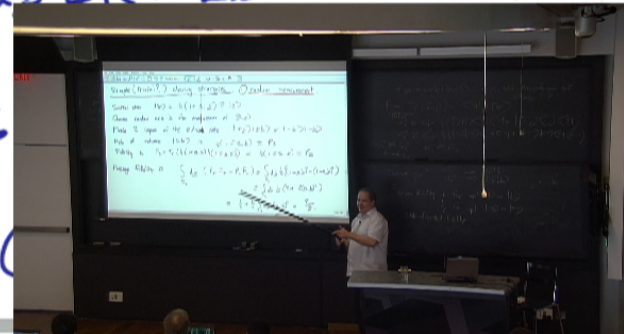
Choose random axis \underline{b} for measurement of $(\underline{e} \cdot \underline{b})$

Make 2 copies of the outcome state $|\pm b\rangle|\pm b\rangle$ or $|-b\rangle|-b\rangle$

Prob of outcome $|\pm b\rangle$ is $\frac{1}{2}(1 \pm a \cdot b) \equiv P_{\pm}$

Fidelity is $F_{\pm} = \text{Tr}(\frac{1}{2}(1 + a \cdot \underline{e}) \frac{1}{2}(1 \pm b \cdot \underline{e})) = \frac{1}{2}(1 \pm a \cdot b) = P_{\pm}$

Average fidelity is $\int_{S_2} d\underline{b} (P_+ F_+ + P_- F_-) = \int_{S_2} d\underline{b} \frac{1}{2}((1 + a \cdot b)^2 + (1 - a \cdot b)^2)$
 $= \int_{S_2} d\underline{b}$
 $= \frac{1}{2} + \frac{1}{2} \int_{S_2} d\underline{b} ($



Simple (trivial) approximate cloning strategies: ② add a random qubit

$$\text{Initial state } |\psi\rangle = \frac{1}{2}(|+a\rangle \otimes |0\rangle)$$

$$\text{New state } |\psi'\rangle = \frac{1}{2}(|+b\rangle \otimes |0\rangle) \quad (\text{random } b)$$

"Flip a coin" to mix them: prepare $\frac{1}{\sqrt{2}}(|0\rangle |\psi\rangle |\psi'\rangle + |1\rangle |\psi'\rangle |\psi\rangle)$

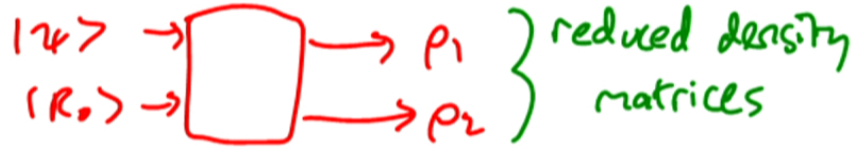
$$P_1 = P_2 = \frac{1}{2} (|\psi\rangle\langle\psi| + |\psi'\rangle\langle\psi'|)$$

$$\text{Averaged over } b, \text{ fidelity } \langle\psi|P_1|\psi\rangle = \frac{1}{2} + \frac{1}{4} = \frac{3}{4}$$

Could we do better? What is the max possible?

Firming up the no-signalling bound on approximate cloning (Gisin, 1998)

Let's focus on cloning qubits:



We can assume symmetric cloning: $\langle \psi | \rho_1 | \psi \rangle = \langle \psi | \rho_2 | \psi \rangle$
 (if not, randomize outputs).

We can also assume universal cloning: $\langle \psi | \rho_i | \psi \rangle = F$ independent of $|\psi\rangle$.

(For if we have a non-universal machine, we can rotate the input and output states by a randomly chosen U and U^{-1} , making a universal cloner.)

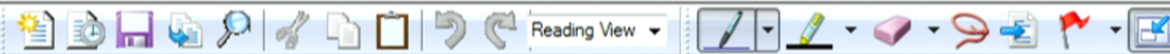
It can be shown this also makes the outputs rotationally invariant functions of the inputs.

$$|\psi\rangle = \frac{1}{2}(1 + \underline{a} \cdot \underline{\sigma}) \Rightarrow \rho_i = \frac{1}{2}(1 + \eta \underline{a} \cdot \underline{\sigma})$$

where the fidelity $F = \frac{1}{2}(1 + \eta)$

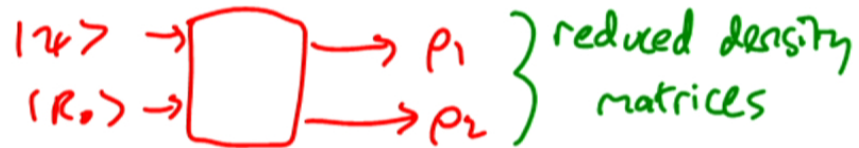


Bloch sphere "shrinking"



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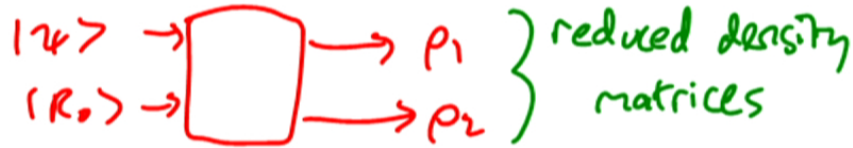
$$\psi = \frac{1}{2}(1 + \underline{a} \cdot \underline{\sigma}) \Rightarrow \rho_i = \frac{1}{2}(1 + \eta \underline{a} \cdot \underline{\sigma})$$

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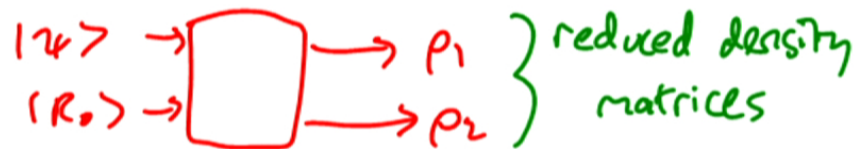
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Bloch sphere "shrinking"

Summary ① We analyse machines for approximately cloning qubits:



② Symmetry arguments allow us to restrict attention to machines of a quite restricted type: $\rho_1 = \rho_2$ and ρ_i have the same Bloch vector as ψ , up to a scaling factor. The machine outputs are "noisy" copies of ψ : mixtures of ψ and the uniformly mixed state $\frac{1}{2}I$.

$$\psi = \frac{1}{2}(1 + \underline{a} \cdot \underline{\sigma}) \Rightarrow \rho_i = \frac{1}{2}(1 + \eta \underline{a} \cdot \underline{\sigma})$$

where the fidelity $F = \frac{1}{2}(1 + \eta)$



$$P_{out}(\underline{m}) = \frac{1}{4}(I \otimes I + \eta(\underline{m} \cdot \underline{\epsilon} \otimes I + I \otimes \underline{m} \cdot \underline{\epsilon})) + \sum_{j,k=1}^3 t_{jk} \epsilon_j \otimes \epsilon_k$$

We can constrain the t_{jk} and η using ① rotational invariance,

To see this explicitly, consider e.g. $P_{out}(\uparrow)$. Rotational invariance about z axis means $t_{xz} = t_{yz} = t_{zx} = t_{zy} = 0$, $t_{xx} = t_{yy}$, $t_{xy} = -t_{yx}$ (think of t as a 3-d tensor invariant under rotations about \underline{z})

$$\text{So } P_{out}(\uparrow) = \frac{1}{4}(I \otimes I + \eta(G_z \otimes I + I \otimes G_z)) + t_{xx}(G_x \otimes G_x + G_y \otimes G_y) + t_{zz}(G_z \otimes G_z) + t_{xy}(G_x \otimes G_y - G_y \otimes G_x).$$

$$\text{and } P_{out}(\uparrow) + P_{out}(\downarrow) = \frac{1}{2}(I \otimes I + t_{xx}(G_x \otimes G_x + G_y \otimes G_y) + t_{zz}(G_z \otimes G_z))$$

Next we consider ② no-signalling: $P_{out}(\uparrow) + P_{out}(\downarrow) = P_{out}(\rightarrow) + P_{out}(\leftarrow)$

② non-signalling $p_{out}(\uparrow) + p_{out}(\downarrow) = p_{out}(\rightarrow) + p_{out}(\leftarrow)$ (*)

$$p_{out}(\uparrow) + p_{out}(\downarrow) = \frac{1}{2}(I \otimes I + t_{xx}(G_x \otimes G_x + G_y \otimes G_y) + t_{zz}(G_z \otimes G_z))$$

and similarly $p_{out}(\rightarrow) + p_{out}(\leftarrow) = \frac{1}{2}(I \otimes I + t_{yy}(G_y \otimes G_y + G_z \otimes G_z) + t_{zz}(G_x \otimes G_x))$

So (*) implies $t_{xx} = t_{zz} = t$

We now have

$$\begin{aligned} p_{out}(\uparrow) &= \frac{1}{4}(I \otimes I + m(G_z \otimes I + I \otimes G_z) + \cancel{t_{xx}}(G_x \otimes G_x + G_y \otimes G_y) \\ &\quad + \cancel{t_{zz}}(G_z \otimes G_z) + t_{xy}(G_x \otimes G_y - G_y \otimes G_x)) \\ &= \frac{1}{4}(I \otimes I + m(G_z \otimes I + I \otimes G_z) + t(G_x \otimes G_x + G_y \otimes G_y + G_z \otimes G_z) \\ &\quad + t_{xy}(G_x \otimes G_y - G_y \otimes G_x)) \end{aligned}$$

Finally we consider ③ non-negativity - all eigenvalues ≥ 0

$$P_{out}(\uparrow) = \frac{1}{4}(I \otimes I + \eta(G_2 \otimes I + I \otimes G_2)) + t(G_x \otimes G_x + G_y \otimes G_y + G_z \otimes G_z) + t_{xy}(G_x \otimes G_y - G_y \otimes G_x)$$

Finally we consider ③ non-negativity - all eigenvalues ≥ 0

$$P_{out}(\uparrow) \text{ has eigenvalues } \frac{1}{4}(1 \pm 2\eta + t) \Rightarrow \eta \leq \frac{t+1}{2}$$

$$\frac{1}{4}(1-t \pm 2(t^2 + t_{xy}^2)^{1/2}) \Rightarrow 2(t^2 + t_{xy}^2)^{1/2} \leq 1-t$$

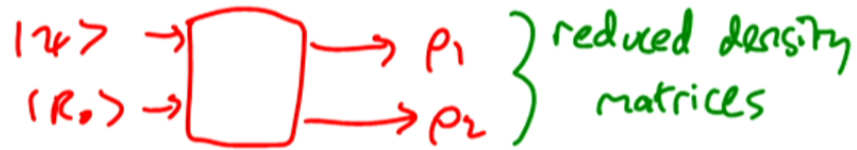
To maximize η we need to maximize t , which gives $t_{xy} = 0$

$$t = \frac{1}{3}$$

$$\eta = \frac{2}{3}$$

$$F = \left(\frac{1+\eta}{2}\right) = \frac{5}{6}$$

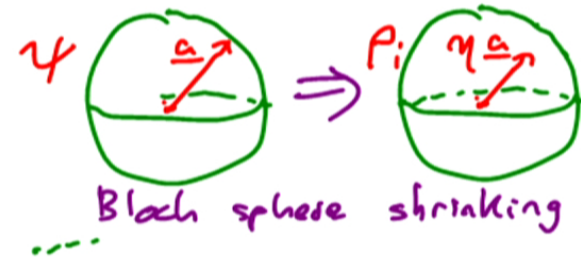
QED



$$\mathcal{U} = \frac{1}{2}(1 + \underline{a} \cdot \underline{\sigma}) \Rightarrow \rho_i = \frac{1}{2}(1 + \eta \underline{a} \cdot \underline{\sigma})$$

where the fidelity

$$F = \frac{1}{2}(1 + \eta) \leq \frac{5}{6}$$



So, no-signalling gives us an upper bound: a cloning scheme cannot reliably produce 2 copies both with fidelity $> \frac{5}{6}$.

Can we achieve this bound? Yes!! Bužek-Hillery (1996) gave an explicit construction of a fidelity $\frac{5}{6}$ universal cloner.

The Bužek-Hillery approximate cloning machine a machine defined by just 1 qubit!

$$\text{Let } |0\rangle |R_0\rangle |M_0\rangle \rightarrow \sqrt{\frac{2}{3}} |0\rangle |0\rangle |1\rangle - \sqrt{\frac{1}{3}} |\Psi^+\rangle |0\rangle$$

$$|1\rangle |R_0\rangle |M_0\rangle \rightarrow -\sqrt{\frac{2}{3}} |1\rangle |1\rangle |0\rangle + \sqrt{\frac{1}{3}} |\Psi^+\rangle |1\rangle$$

$$\text{where } |\Psi^+\rangle = \frac{1}{\sqrt{2}} (|0\rangle |1\rangle + |1\rangle |0\rangle).$$

$$\text{Then } (a|0\rangle + b|1\rangle) |R_0\rangle |M_0\rangle \rightarrow \sqrt{\frac{2}{3}} a |0\rangle |0\rangle |1\rangle + \sqrt{\frac{1}{3}} a |\Psi^+\rangle |0\rangle \\ - \sqrt{\frac{2}{3}} b |1\rangle |1\rangle |0\rangle + \sqrt{\frac{1}{3}} b |\Psi^+\rangle |1\rangle$$

$$\begin{aligned} |\Psi\rangle &= a|0\rangle + b|1\rangle \\ |\Psi^+\rangle &= a^*|1\rangle - b^*|0\rangle \end{aligned} = \dots = \sqrt{\frac{2}{3}} (a|0\rangle + b|1\rangle)(a|0\rangle + b|1\rangle)(a^*|1\rangle - b^*|0\rangle) \\ - \sqrt{\frac{1}{3}} ((a|0\rangle + b|1\rangle)(a^*|1\rangle - b^*|0\rangle) + (a^*|1\rangle - b^*|0\rangle)(a|0\rangle + b|1\rangle)) (a|0\rangle + b|1\rangle)$$

$$\text{So as promised } \langle \Psi | \rho_1 | \Psi \rangle = \langle \Psi | \rho_2 | \Psi \rangle = \frac{5}{6} = \sqrt{\frac{2}{3}} \langle \Psi | \Psi \rangle \langle \Psi^+ | \Psi^+ \rangle - \sqrt{\frac{1}{3}} (\langle \Psi | \Psi^+ \rangle + \langle \Psi^+ | \Psi \rangle) \langle \Psi | \Psi \rangle$$

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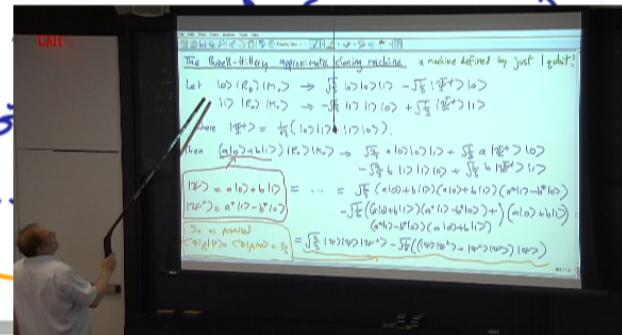
$$\text{where } |\Psi^+\rangle = \frac{1}{\sqrt{2}} (|0\rangle |1\rangle + |1\rangle |0\rangle).$$

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$$|\Psi\rangle = a|0\rangle + b|1\rangle = \dots = \sqrt{\frac{2}{3}} (a|0\rangle + b|1\rangle) (a|0\rangle + b|1\rangle) (a^*|1\rangle - b^*|0\rangle) \\ |\Psi^+\rangle = a^*|1\rangle - b^*|0\rangle$$

$$\text{So as promised} \\ \langle \Psi | \rho_1 | \Psi \rangle = \langle \Psi | \rho_2 | \Psi \rangle = \frac{5}{6}$$

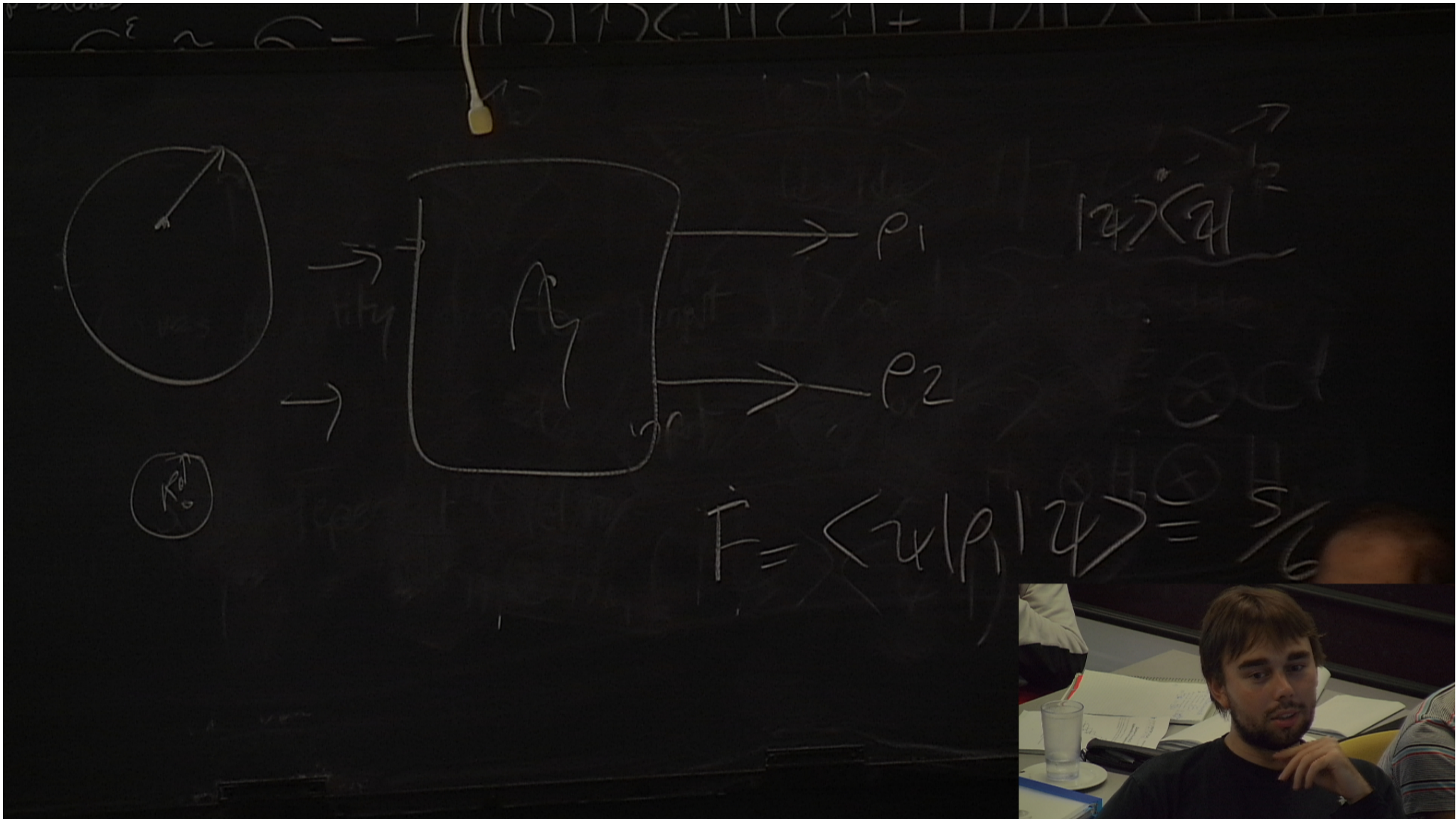
$$= \sqrt{\frac{2}{3}} |\Psi\rangle |\Psi\rangle |\Psi^+\rangle - \sqrt{\frac{1}{3}} (|\Psi\rangle |\Psi^+\rangle)$$



Note that while the Buzek-Hillery optimal universal cloner is simple and elegant, and while it's satisfying that it achieves precisely the no-signalling bound, it's not that much better than the trivial add-a-random-qubit cloner ($F = 5/6$ compared to $F = 3/4$).

This turns out to be generally true: in higher dimensions, or $M \rightarrow N$ copy cloning, you can do a bit better than trivial strategies, but not much.

(Bad news if you're thinking of cloning as a way of "backing up" your quantum states. But good news if you want to ensure that other people can't copy them!)

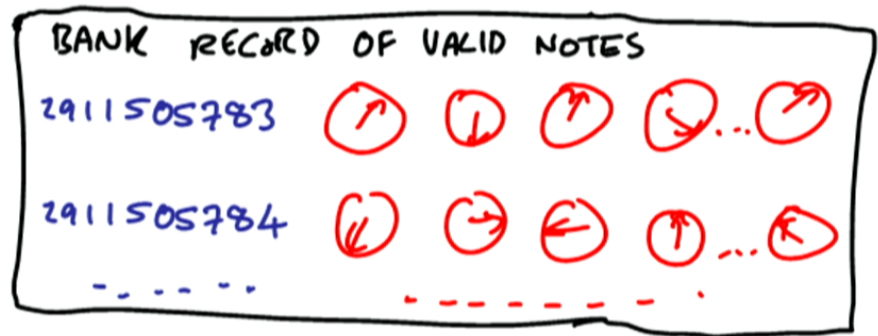


Quantum Money (Wiesner c.1970, unpublished till 1983)

Classical money is basically classical information: it is copiable (forgeable) with arbitrary precision. A sufficiently skilled forger can make any number of copies.



Wiesner's solution: incorporate randomly chosen quantum states, known to the bank but not to customers (or forgers)

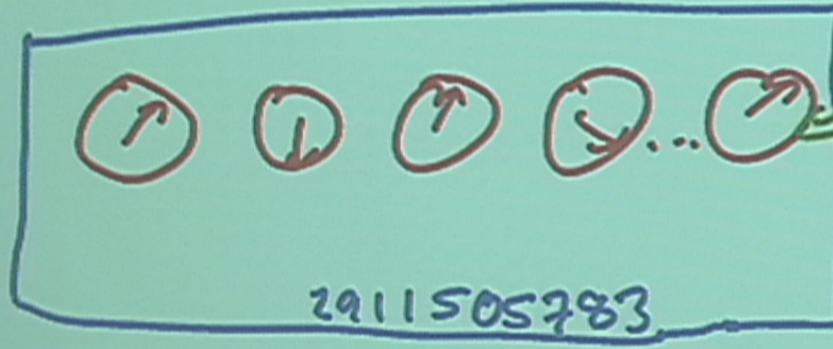


To create even two valid notes from one, the forger needs to try to clone N unknown qubits.

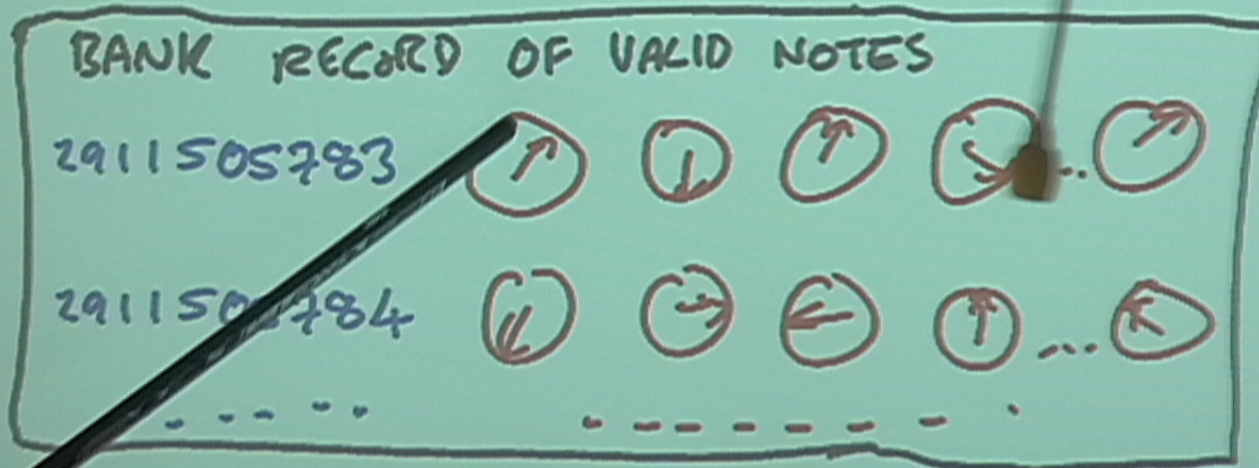
We've seen her success probability is $\leq (\frac{3}{4})^N$. I.E. SECURE FOR LARGE N

any number of copies.

N qubits



Wiesner's solution
quantum states,
to customers (0



To
from
clone
We've
≤