

Title: Local states and channels in causal theories

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Abstract: We will analyze different aspects of locality in causal operational probabilistic theories. We will first discuss the notion of local state and local objective information in operational probabilistic theories, and define an operational notion of discord that coincides with quantum discord in the case of quantum theory. Using such notion, we will show that the only theory in which all separable states have null discord is the classical one. We will then analyze locality of transformations, reviewing some general properties of no-signaling channels in causal theories. We will show that it is natural to define transformations on no-signaling channels that cannot be extended to all bipartite channels, and discuss the consequences of this fact on information processing.

# Locality and causality in operational theories

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## In collaboration with

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- G. M. D'Ariano



- G. Chiribella



- S. Facchini



- M. Zaopo



# Outline

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- Causal theories
- Locality (different operational notions)

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- Non-locality without entanglement and local objective information

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- Discord = non-classicality
- Local discrimination

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- Discord = non-classicality
- Local discriminability
- Localizable and non-signalling channels in causal theories with I. d.



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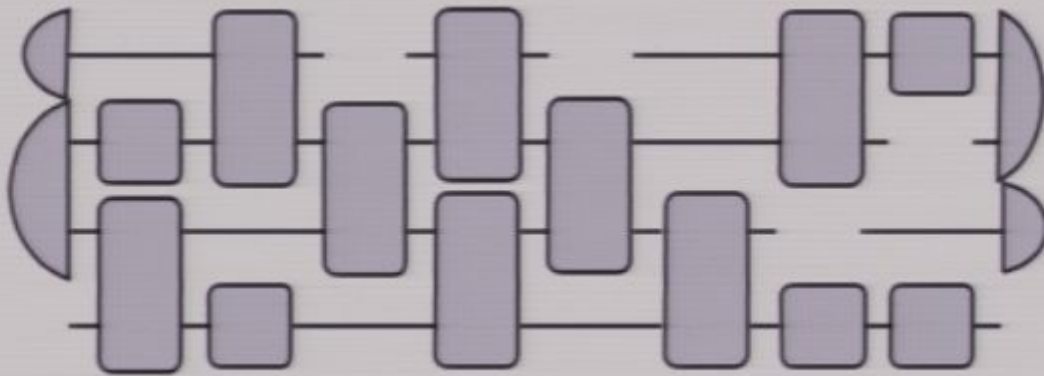
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- Causal theories
- Locality (different operational notions)
- Non-locality without entanglement and local objective information
- Discord = non-classicality
- Local discriminability
- Localizable and non-signalling channels in causal theories with I. d.
- Combs and higher order transformations in theories with purification

# The operational language

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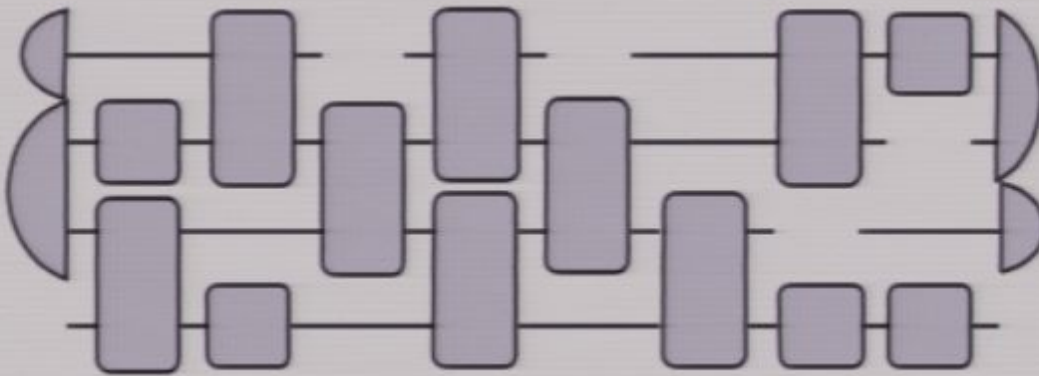
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# The operational language

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- Operational theory: tests with composition rules

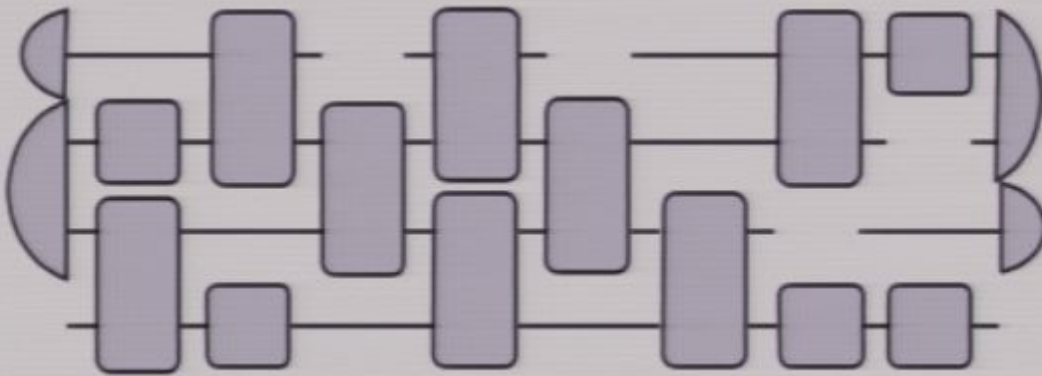


A : input label

B : output label

# The operational language

- Operational theory: tests with composition rules



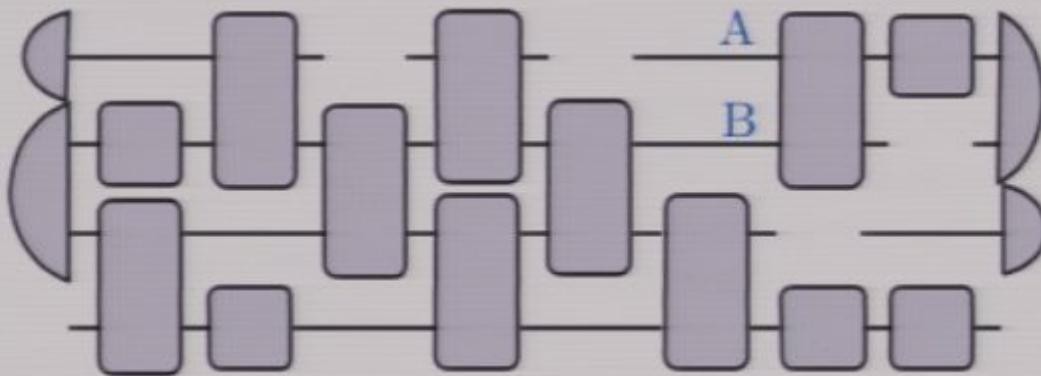
$$\frac{I}{\rho_i} \frac{A}{\rho_i} = \frac{A}{\rho_i}$$

$$\frac{B}{a_i} \frac{I}{a_i} = \frac{B}{a_i}$$

# The operational language

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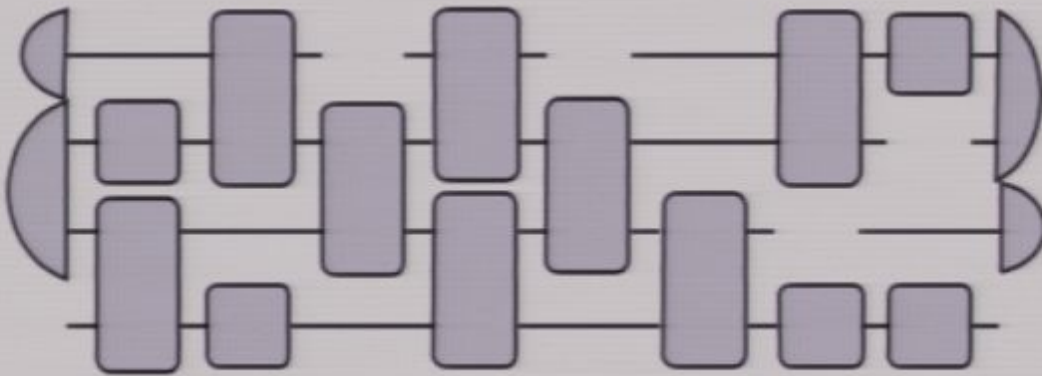
- Operational theory: tests with composition rules



- $C := AB = BA$
- $(AB)C = A(BC)$
- $AI = IA = A$

# The operational language

- Operational theory: tests with composition rules



- For any system  $A$  there exists a unique test  $\mathcal{I}_A$  such that

$$\begin{array}{c} A \\ \text{---} \end{array} \boxed{\mathcal{C}_i} \begin{array}{c} B \\ \text{---} \end{array} = \begin{array}{c} A \\ \text{---} \end{array} \boxed{\mathcal{I}_A} \begin{array}{c} A \\ \text{---} \end{array} \boxed{\mathcal{C}_i} \begin{array}{c} B \\ \text{---} \end{array} = \begin{array}{c} A \\ \text{---} \end{array} \boxed{\mathcal{C}_i} \begin{array}{c} B \\ \text{---} \end{array} \boxed{\mathcal{I}_B} \begin{array}{c} B \\ \text{---} \end{array}$$



# The probabilistic structure

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- Probabilistic theory

Every test of type  $I \rightarrow I$  is a probability distribution

$$\begin{array}{|c|} \hline i_1 \\ \hline \end{array} \text{---} \begin{array}{|c|} \hline i_2 \\ \hline \end{array} = p(i_1, i_2)$$

States are functionals on effects and viceversa



Real vector spaces  $\text{St}_{\mathbb{R}}(A), \text{Eff}_{\mathbb{R}}(A)$

$\mathcal{T}_{\mathbb{R}}(A, B)$  transformations are collections of linear maps

# Causality

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$$p_a(\rho_i) := \sum_j \left( \rho_i \text{---} a_j \right) = p(\rho_i)$$

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- Uniqueness of the deterministic effect

$$\sum_j \text{---} a_j = \sum_k \text{---} b_k = \text{---} e$$

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# Causality

- Causality

$$p_a(\rho_i) := \sum_j \left( \rho_i \text{---} a_j \right) = p(\rho_i)$$

- Uniqueness of the deterministic effect

$$\sum_j \text{---} a_j = \sum_k \text{---} b_k = \text{---} e$$

- All states are proportional to deterministic ones
- Unrestricted conditioning

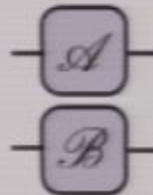
- Deterministic transformations

$$\text{---} \boxed{\mathcal{C}} \text{---} e = \text{---} e$$

# Locality properties of operational boxes

- Operationally locality of channels is classified by different notions:

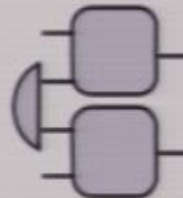
- Factorized



- LOSR\*

$$\sum_i p_i \begin{array}{c} \boxed{A_i} \\ \boxed{B_i} \end{array}$$

- Localizable\*\*



- Non-signalling\*\*





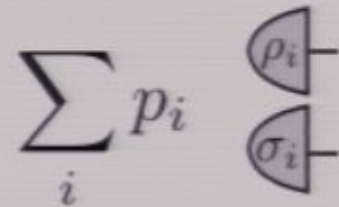
# Locality properties of states in causal theories

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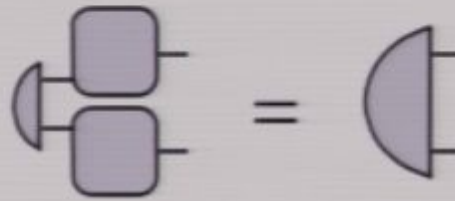
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- LOSR



- Localizable, non-signalling and general bipartite states coincide





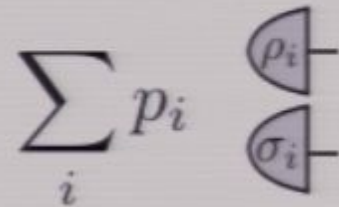
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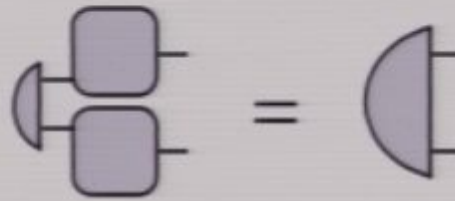


- LOSR



Separability

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- Localizable, non-signalling and general bipartite states coincide



# Non-locality without entanglement

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- In “Quantum Nonlocality without entanglement”<sup>\*</sup> the authors introduce a different notion of locality
- This definition is based on locality of the measurement of the eigenbasis
- The classical information encoded by a random source of distinguishable states can/cannot be accessed by LOCC
- How can we define a similar kind of non-locality in causal operational theories?
- A state is local if it encodes **locally readable objective information**

<sup>\*</sup> C. H. Bennett, D. P. DiVincenzo, C. A. Fuchs, T. Mor, E. Rains, P. W. Shor, J. A. Smolin, and H. K. Wootters, Phys. Rev. A 62, 102301 (2000)

# Objective information

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- Einstein, Podolski and Rosen: **sufficient criterion** for **elements of reality**

*If, without in any way disturbing a system, we can predict with certainty (i.e., with probability equal to unity) the value of a physical quantity, then there exists an element of physical reality corresponding to this physical quantity.*

- A state  $\rho$  encodes **objective information** about the test  $\{\mathcal{A}_i\}$  if

- The test is repeatable  $\mathcal{A}_i \mathcal{A}_j = \delta_{ij} \mathcal{A}_i$

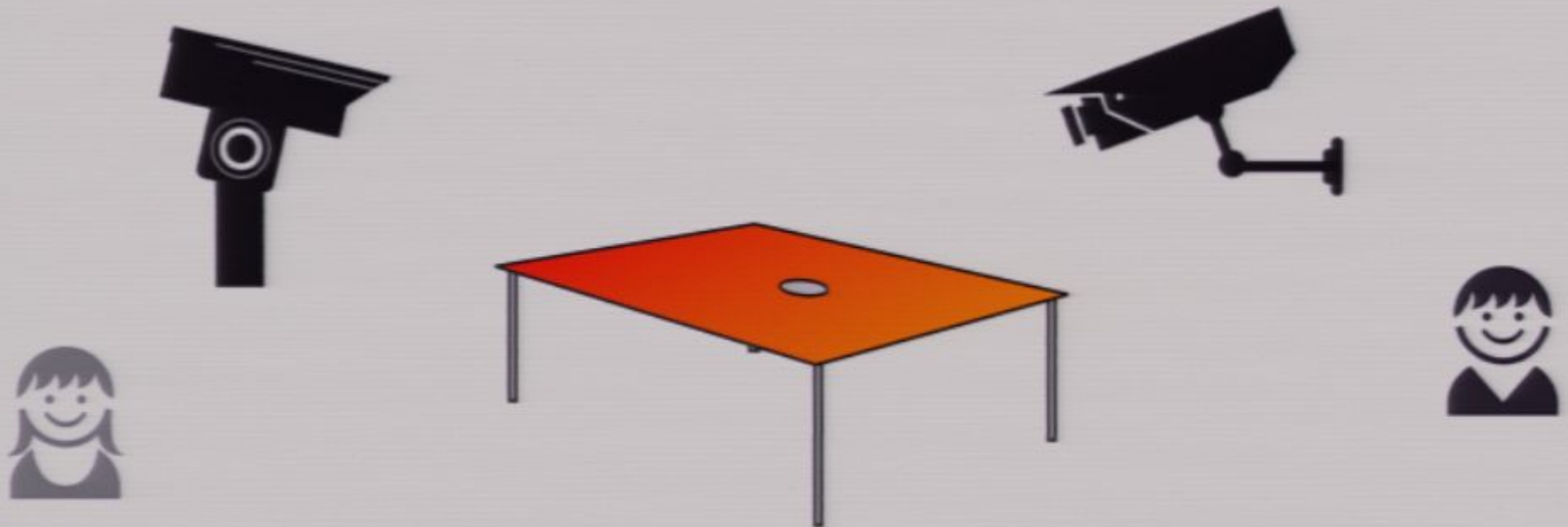
- The test does not disturb the state  $\mathcal{A} \rho = \rho$   $\mathcal{A} := \sum_i \mathcal{A}_i$

- The objective information is complete if  $\mathcal{A}_i \rho$  is pure for every  $i$

# Example

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- Consider a tossed coin before the {heads, tails} test has been performed



- Information about the upper side of the coin is objective

## Consequences of the definition

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$$\cup = \cup \boxed{\phantom{x}} \boxed{\phantom{x}} - \boxed{\phantom{x}}$$

- A state carries objective information if and only if

$$\rho = \sum_i p_i \rho_i \quad \mathcal{A}_i \rho_j = \delta_{ij} \rho_j$$

- A state carries complete objective information if and only if

$$\rho = \sum_i p_i \psi_i \quad \mathcal{A}_i \psi_j = \delta_{ij} \psi_j$$



# Local objective information

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- Local state in the sense of N. L. W. E.: a bipartite state  $\rho$  s. t.
  - The state  $\rho$  encodes complete objective information about  $\{\mathcal{A}_i\}$
  - The test  $\{\mathcal{A}_i\}$  can be measured by a LOCC procedure
- Conditions for locality/non locality without entanglement?
  - Work in progress
- The notion of objective information can be used to define **discord**

# The standard notion of discord

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- Definition

$$\delta(S:\mathcal{A})_{\{\Pi_j^{\mathcal{A}}\}} = I(S:\mathcal{A}) - J(S:\mathcal{A})_{\{\Pi_j^{\mathcal{A}}\}}$$

$$J(S:\mathcal{A})_{\{\Pi_j^{\mathcal{A}}\}} = H(S) - H(S \mid \{\Pi_j^{\mathcal{A}}\})$$

- Problems in general theories

- Entropy is not uniquely defined
- Entropy does not enjoy the same properties as in CPT and QT

# Objective information and discord

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- Null discord states: system + pointer after a measurement interaction
  - Complete objective information encoded in the pointer

- In causal theories

**Definition 11** *In a causal operational probabilistic theory, a bipartite state  $\rho_{AB}$  has **null discord** if and only if it satisfies the following conditions*

1.  $\rho_{AB}$  is separable,
2. there exists a test  $\{\mathcal{A}_k\}_{k \in X}$  on system A that provides complete objective information about the state  $\rho_A$ , and such that  $\{\mathcal{A}_k \otimes \mathcal{I}\}_{k \in X}$  provides objective information on  $\rho_{AB}$

- Operational notion of discord

$$\mathcal{D}(\rho_{AB}) := \min_{\sigma \in \Omega_{AB}} \|\rho_{AB} - \sigma\|_{op}$$



# Theorem

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- Hypothesis: a state is separable if and only if it has null discord
- Thesis: the theory is simplicial



- Consequence: discord is the weakest signature of non-classicality
  - Shared by any theory

# Locality and separability of channels

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- For channels separability is not a relevant criterion
- There are non-separable localizable channels

Example from quantum theory: PR box

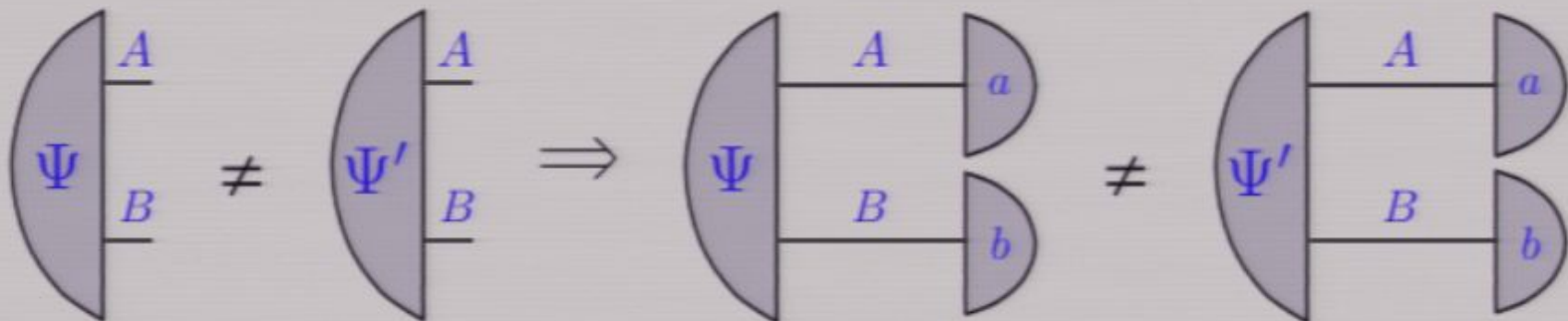
- There are separable channels that are signalling

$$\frac{1}{2} \sum_{ij} \begin{array}{|c|c|} \hline a_i & \rho_i \\ \hline b_j & \sigma_j \\ \hline \end{array} + \frac{1}{2} \sum_{ij} \begin{array}{|c|c|} \hline a_i & \rho_j \\ \hline b_j & \sigma_i \\ \hline \end{array}$$



# Local discriminability

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## A property of channels in causal theories with l. d.

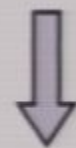
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- In a causal theory a physical transformation is a channel if and only if

$$e_B * \mathcal{C} = e_A$$

- Every  $\mathcal{T}$  satisfying the same condition can be decomposed as follows

$$\mathcal{T} \in \mathfrak{T}_{\mathbb{R}}(A, B), \quad e_B * \mathcal{T} = e_A$$



$$\mathcal{T} = a\mathcal{T}_+ - b\mathcal{T}_- \quad \mathcal{T}_{\pm} \in \mathfrak{T}(A, B)$$

$$e_B * \mathcal{T}_{\pm} = e_A$$

# Linear span of local boxes

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- Bipartite channels in causal theories with local discriminability

$$\mathcal{T} = \sum_i \mathcal{A}_i \otimes \mathcal{B}_i$$

- The transformations  $\{\mathcal{A}_i\}, \{\mathcal{B}_i\}$  can be taken to be linearly independent

- Non-signalling implies  $e * \mathcal{A}_i = \lambda_i e$        $e * \mathcal{B}_i = \mu_i e$

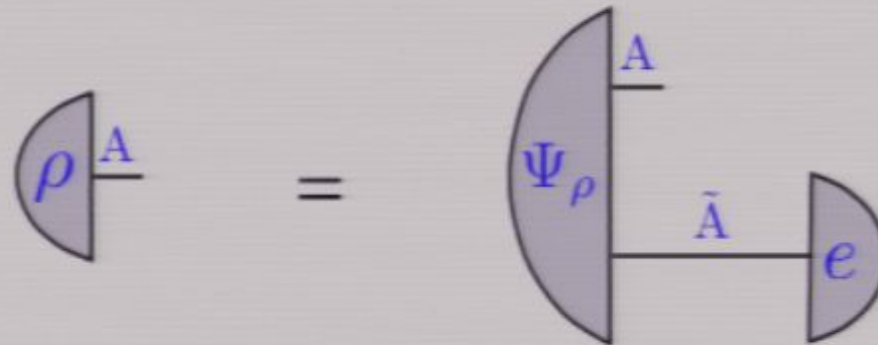
$$\mathcal{A}_i = a_i^+ \mathcal{A}_i^+ - a_i^- \mathcal{A}_i^- \quad \mathcal{B}_i = b_i^+ \mathcal{B}_i^+ - b_i^- \mathcal{B}_i^-$$

- The span of local channels contains non-signalling channels

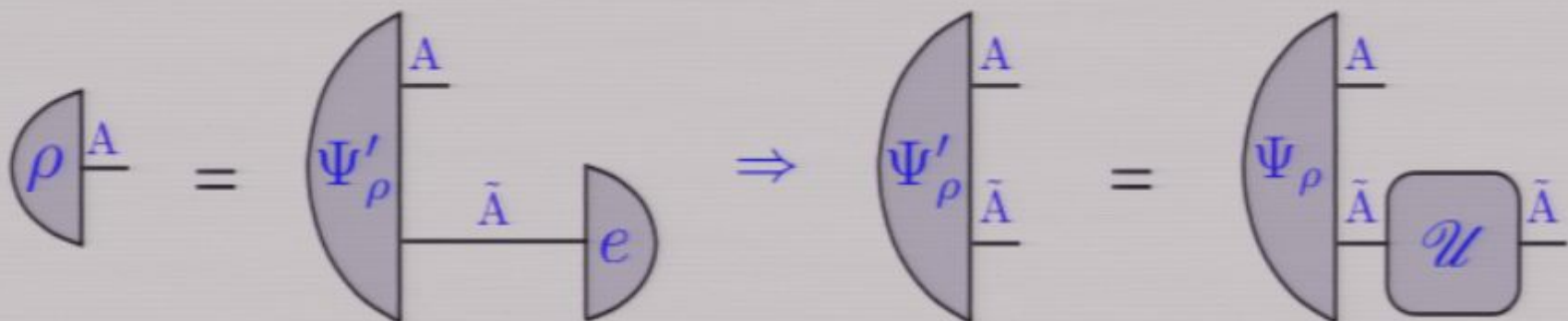
$$\mathcal{T} = \sum_j \mathcal{C}_j \otimes \mathcal{D}_j - \sum_k \mathcal{C}'_k \otimes \mathcal{D}'_k$$

# Purification

- For any state  $\rho$  there exists a **purifying** system  $\tilde{A}$  such that



- The purification is unique up to reversible transformations





# Choi correspondence

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# Choi correspondence

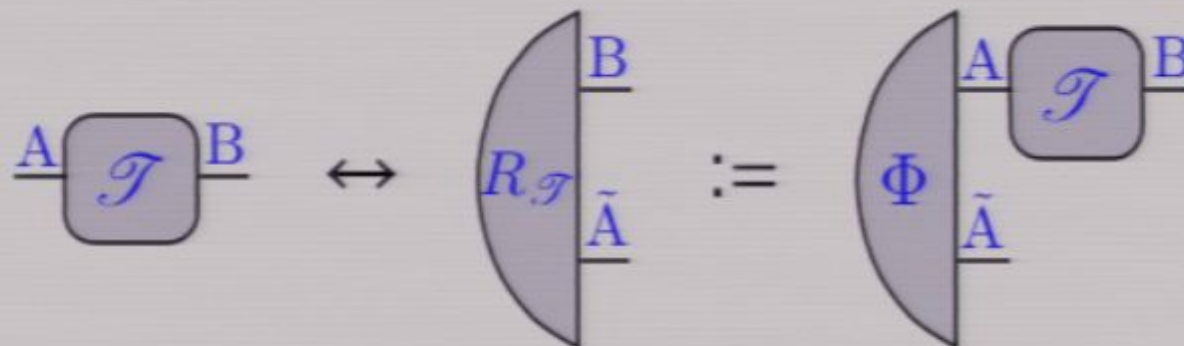
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- Correspondence between bipartite states  $B\tilde{A}$  and transformations  $A \rightarrow B$



# Choi correspondence

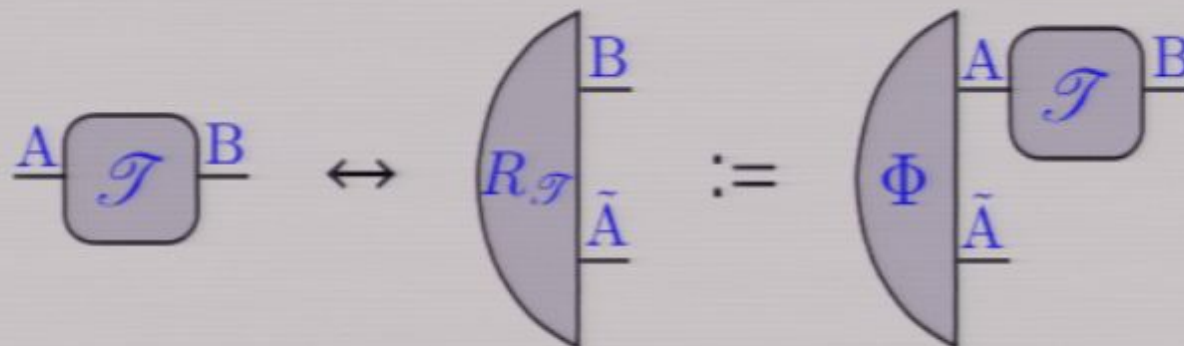
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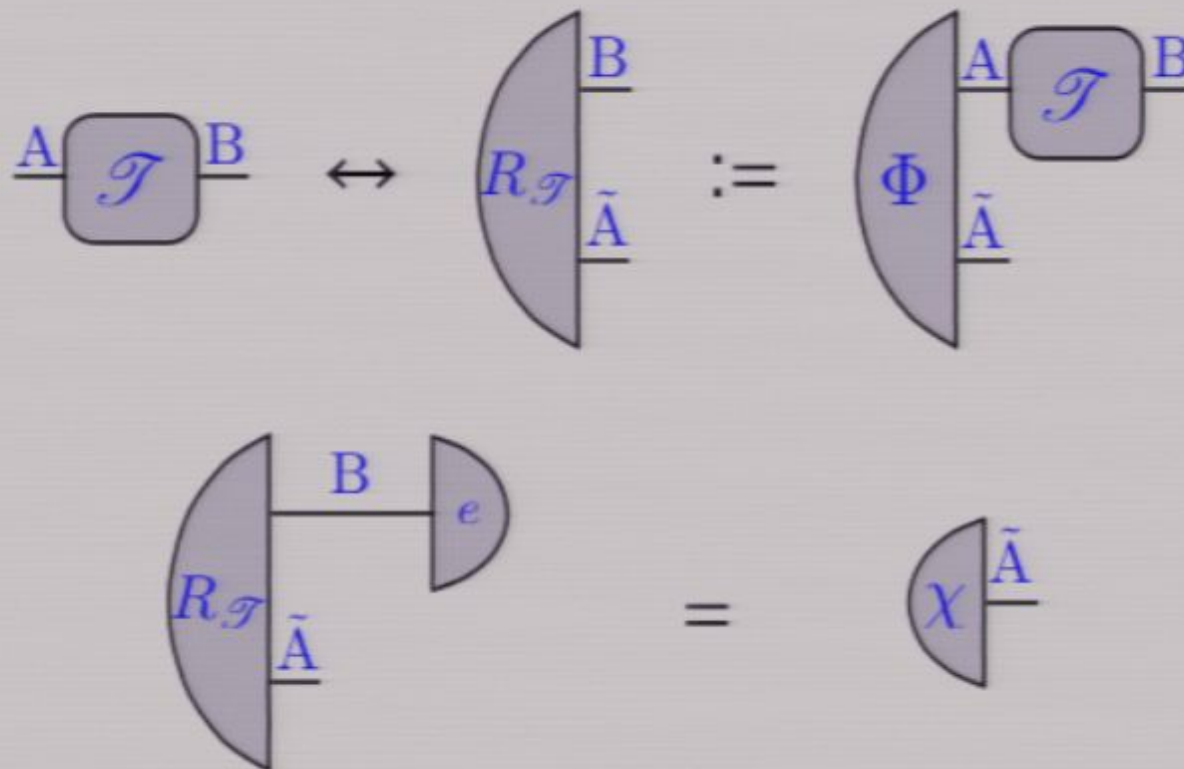
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- Correspondence between bipartite states  $B\bar{A}$  and transformations  $A \rightarrow B$ 
  - Deterministic transformations are in correspondence with **some** states



# Choi correspondence

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# Faithful states

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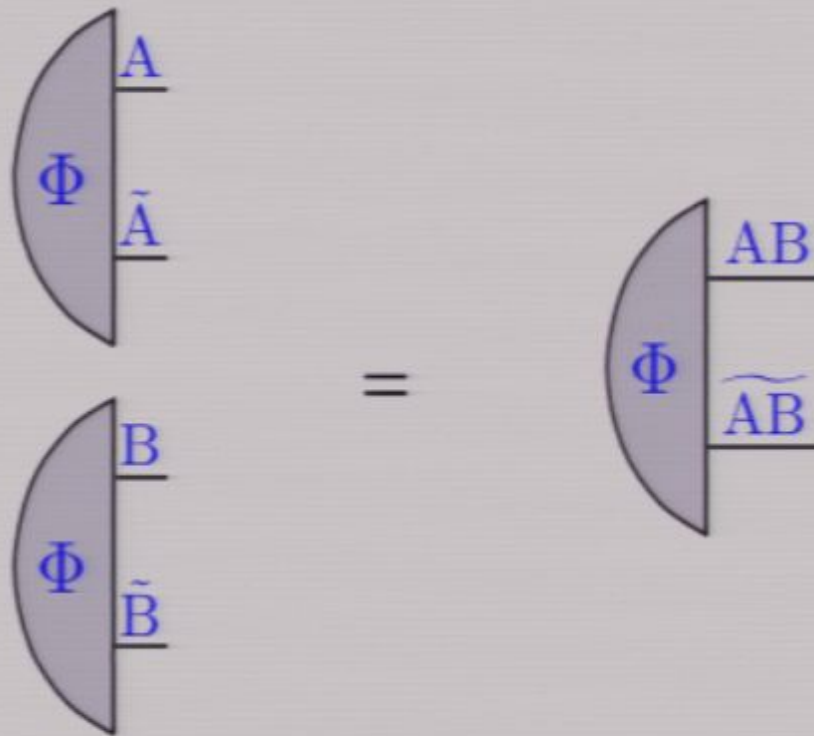
- The Choi correspondence holds through a **faithful** state

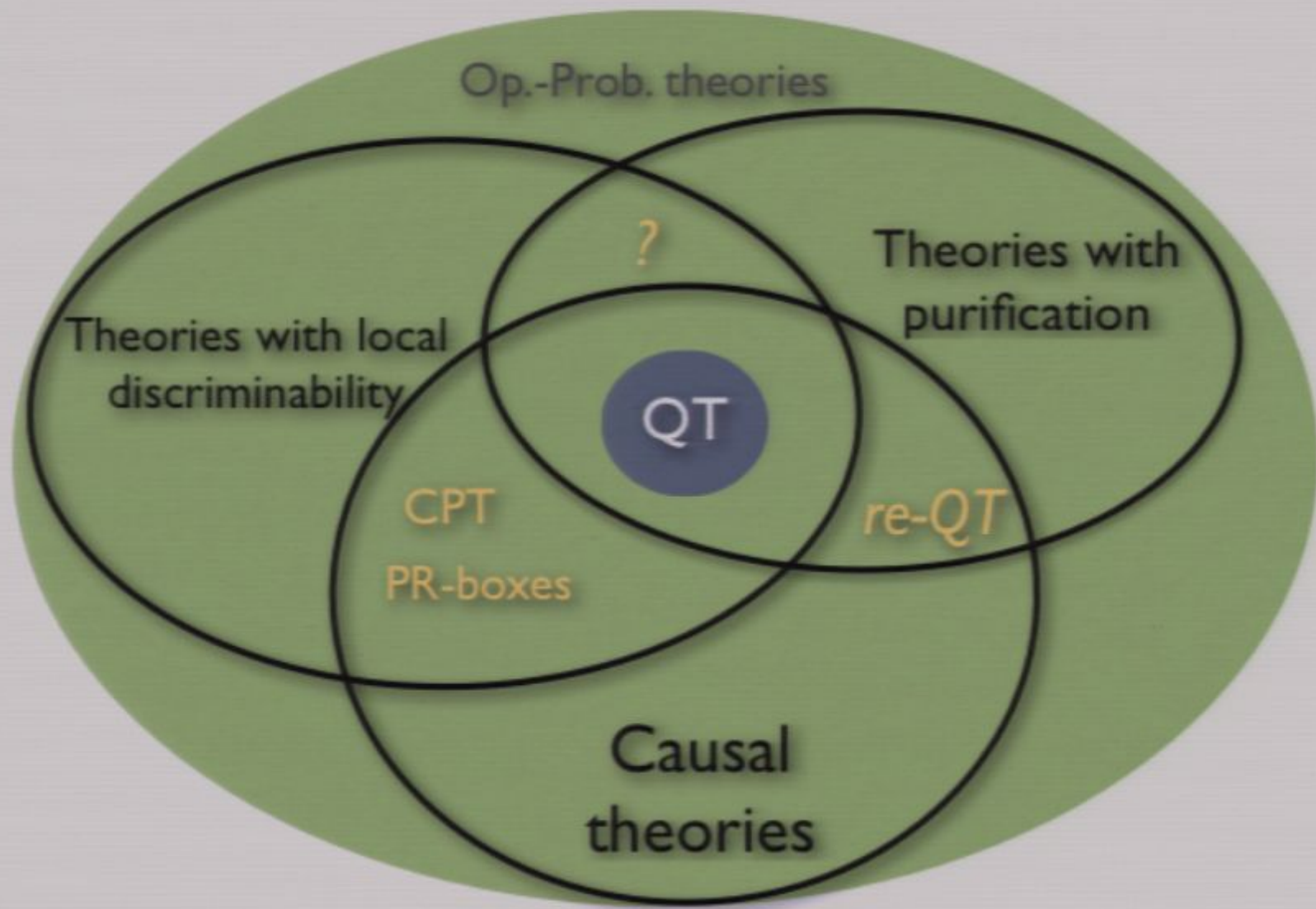


# Faithful states

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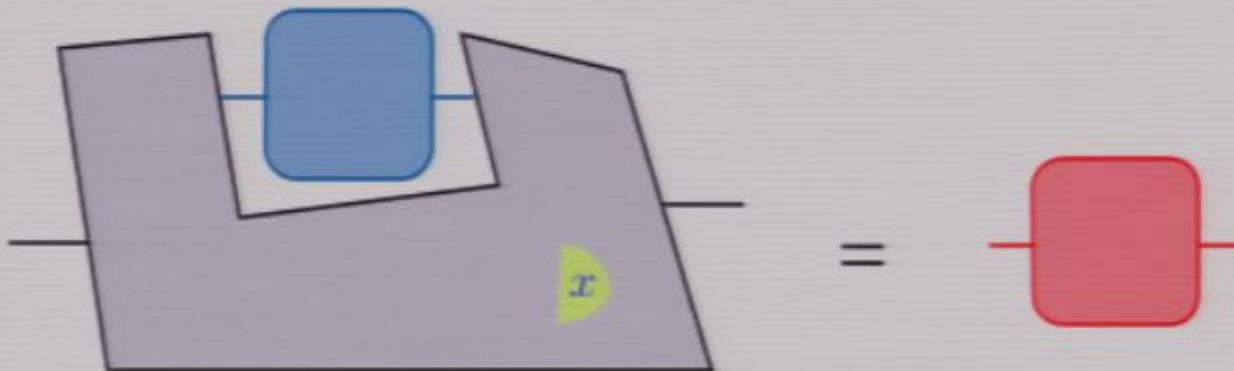
- The Choi correspondence holds through a **faithful** state
- The composition of two faithful states is faithful





# Maps on transformations: supermaps

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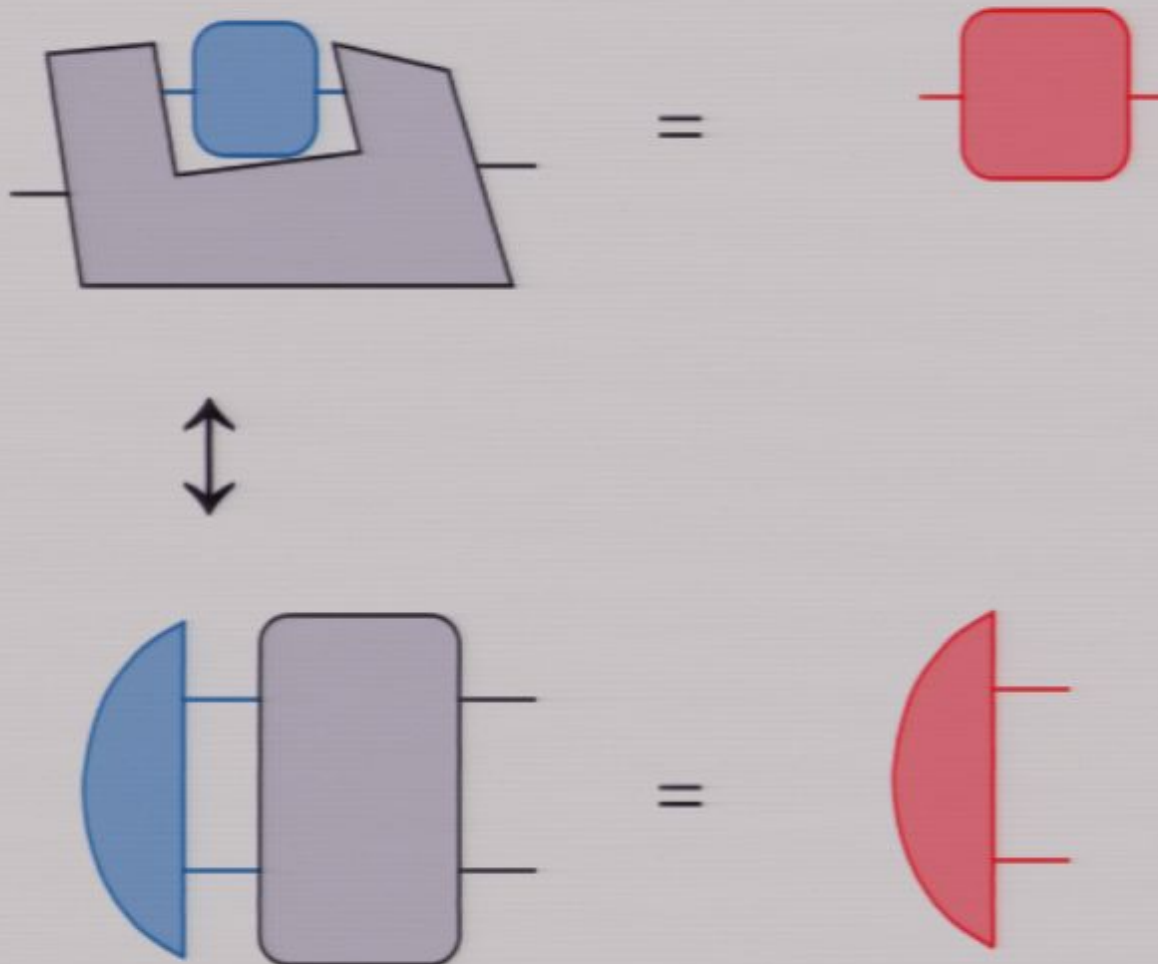
# Maps on transformations: supermaps

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## Admissibility conditions

# Maps on transformations: supermaps

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## Admissibility conditions

- Linear  $\rightarrow$  preservation of convex combinations (probabilities)

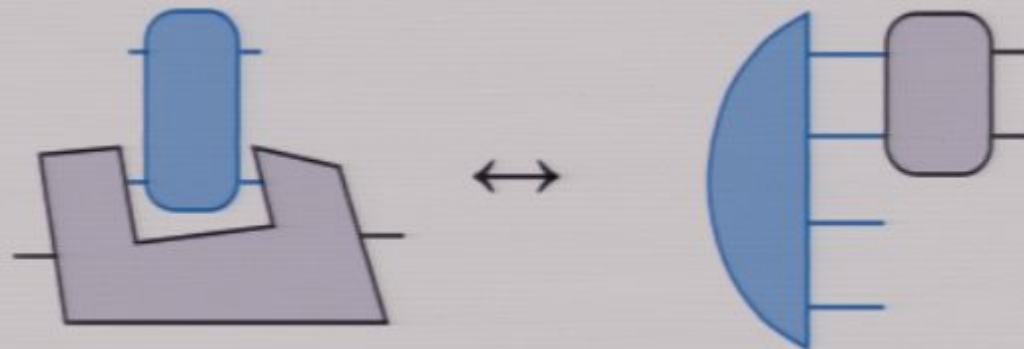
# Maps on transformations: supermaps

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## Admissibility conditions

- Linear  $\rightarrow$  preservation of convex combinations (probabilities)

- Completely Positive



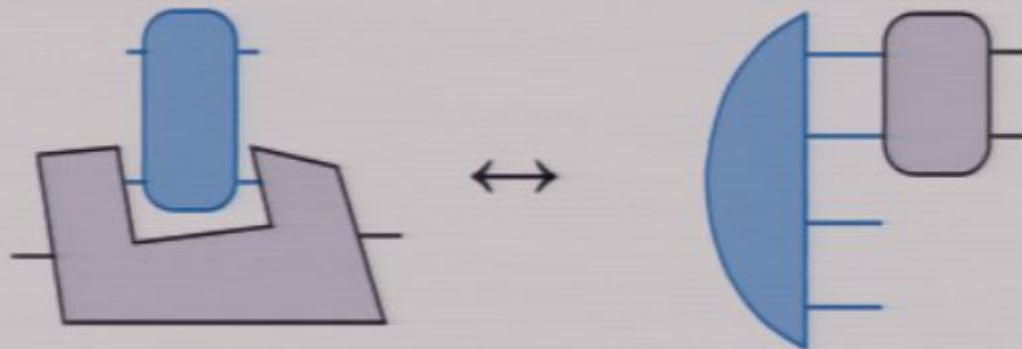
# Maps on transformations: supermaps

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## Admissibility conditions

- Linear  $\rightarrow$  preservation of convex combinations (probabilities)

- Completely Positive



- Deterministic  $\rightarrow$  preservation of normalisation

# Realisation theorem

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Admissibility conditions



$\mathbb{R}$

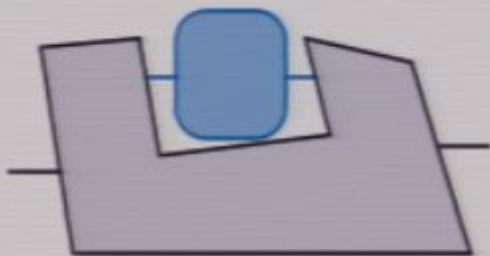




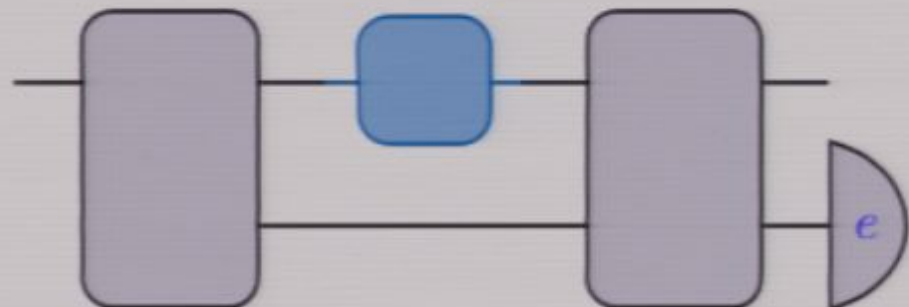
# Realisation theorem

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IR



# Realisation theorem

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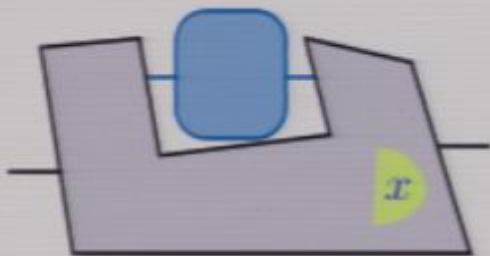
$\mathbb{R}$



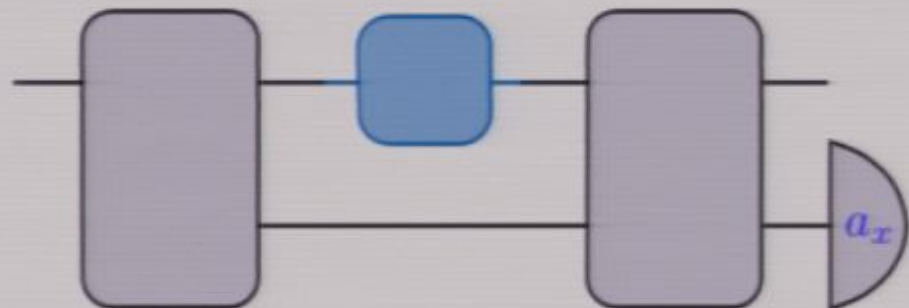
# Realisation theorem

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## Admissibility conditions



IR



# Testers

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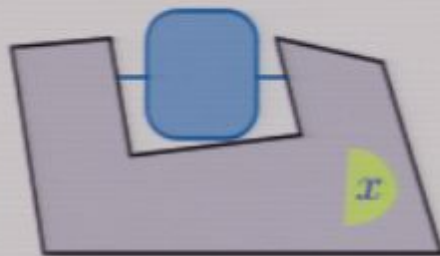
- A measurement on a transformation provides probabilities at the output
- A probability is a transformation of the trivial system  $I$
- Realisation theorem: collection of supermaps with the following scheme



# Testers

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- A measurement on a transformation provides probabilities at the output
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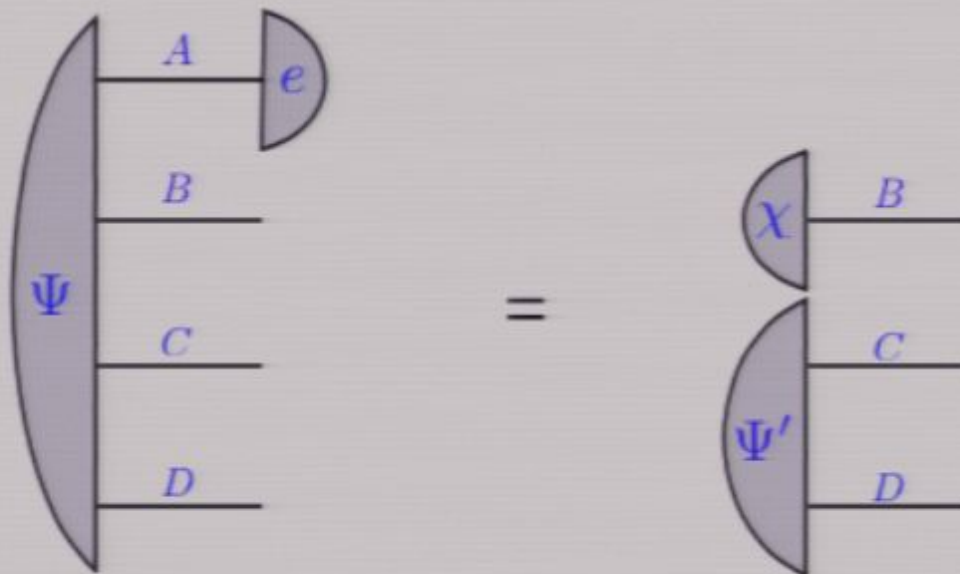


$\equiv$



# Supermaps and states

- Supermaps are in correspondence with states
  - Deterministic supermaps are in correspondence with **some** states



- The cones coincide

# Second-order theory

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# Second-order theory

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- System type

# Second-order theory

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- System types are  $A \rightarrow B$

$$\underline{A \rightarrow B} \rightarrow \underline{A} \quad \underline{B}$$

- States of the system  $A \rightarrow B$  are transformations from  $A$  to  $B$

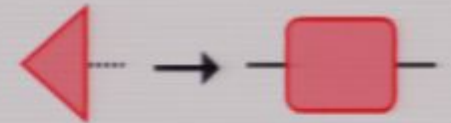
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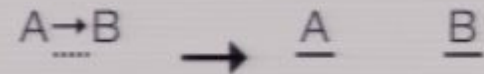
$$\underline{A \rightarrow B} \rightarrow \underline{A} \quad \underline{B}$$

- States of the system  $A \rightarrow B$  are transformations from A to B

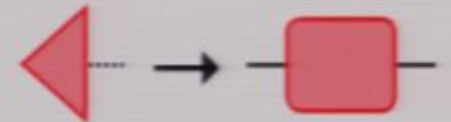


# Second-order theory

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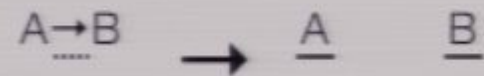


- Transformations  $(A \rightarrow B) \rightarrow (C \rightarrow D)$  are supermaps

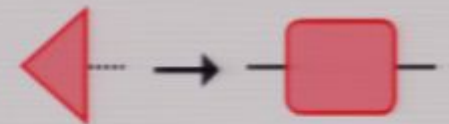


# Second-order theory

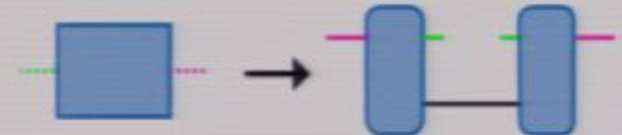
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- States of the system  $A \rightarrow B$  are transformations from  $A$  to  $B$



- Transformations  $(A \rightarrow B) \rightarrow (C \rightarrow D)$  are supermaps



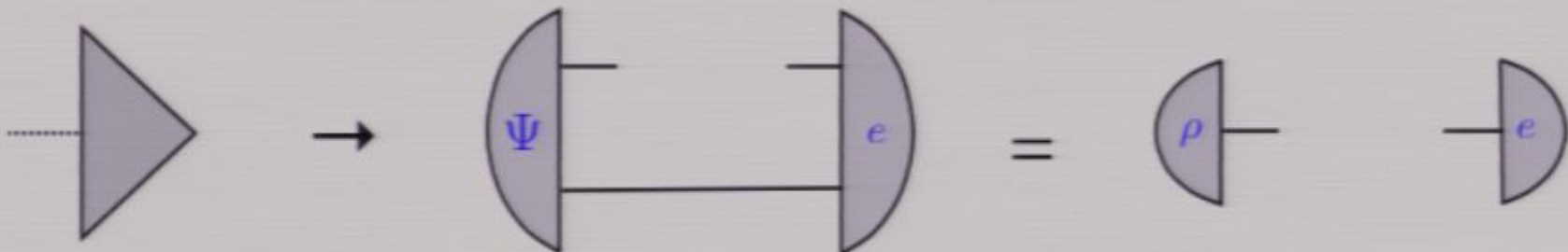
- Effects are testers



# Second-order theory

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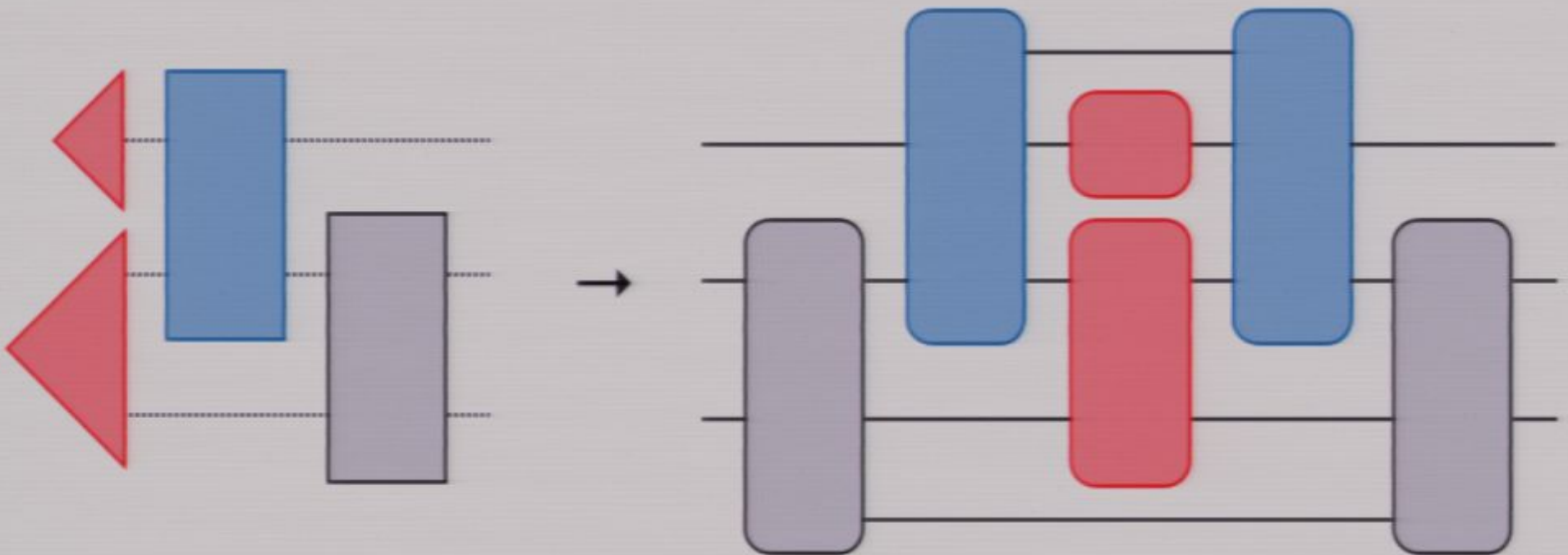
- Deterministic effect  $\leftrightarrow$  deterministic tester
- The second order theory is **non causal**



## Second order theory

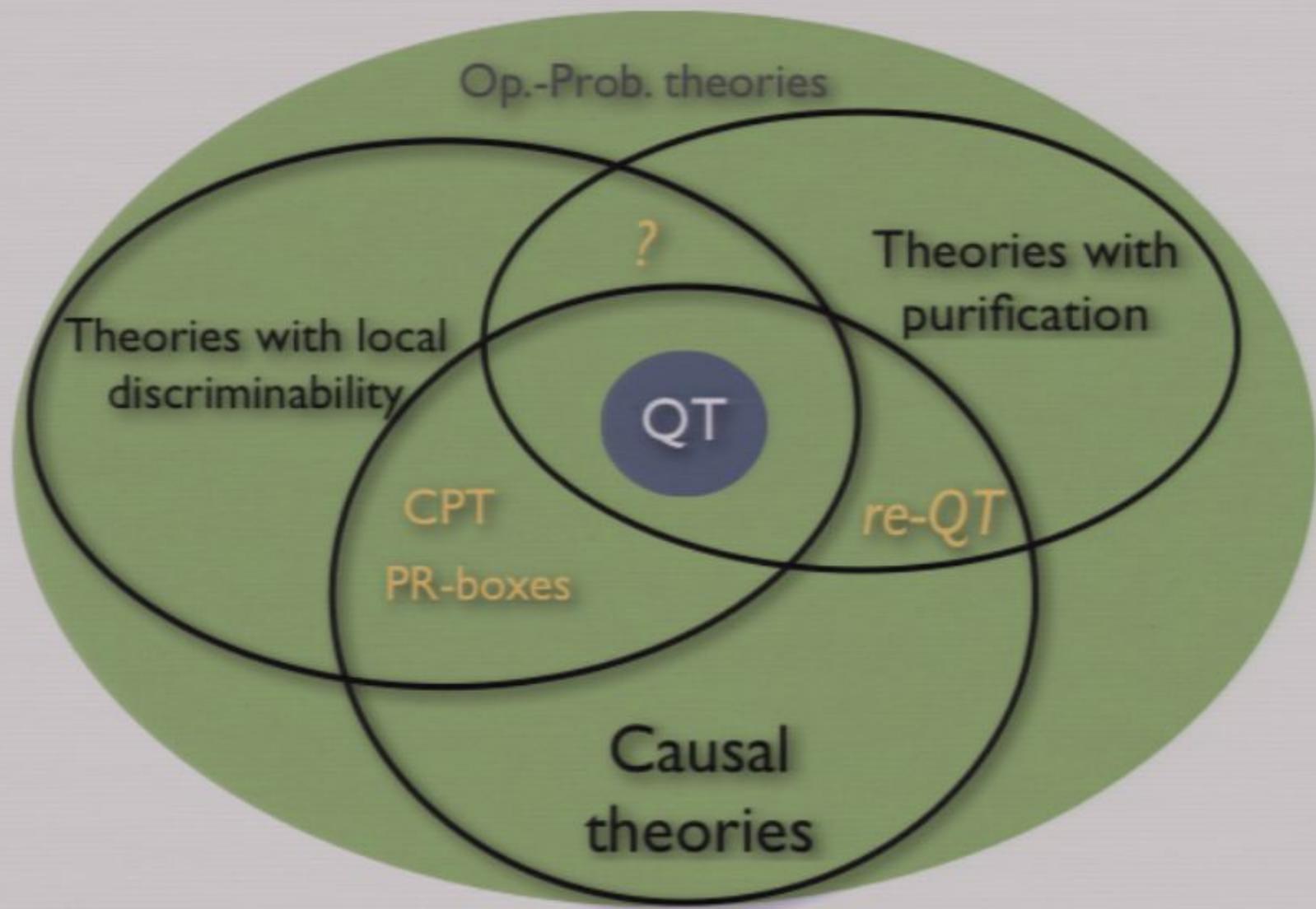
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- Circuits of a second order theory are perfectly simulated in a **causal** theory



Are all non-causal theories of this kind?





# The hierarchy of combs

- Consider the following recursively defined hierarchy of transformations
  - **1-Combs:** transformations in a causal theory with purification
  - **N-Combs:** transformations from N-1-combs to 1-Combs



- Example: 2-Combs are supermaps



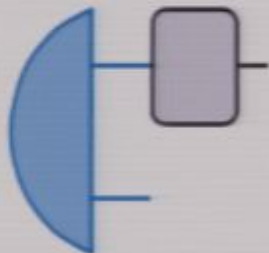
# The hierarchy of combs

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- 1-Combs are in correspondence with states
- If N-1-combs are in correspondence with states, N-combs are in correspondence with transformations, hence with states
- Admissibility conditions:

- Linear

- CP



- Deterministic  $\rightarrow$  Deterministic comb mapped to Deterministic transformations

# Realisation Theorem

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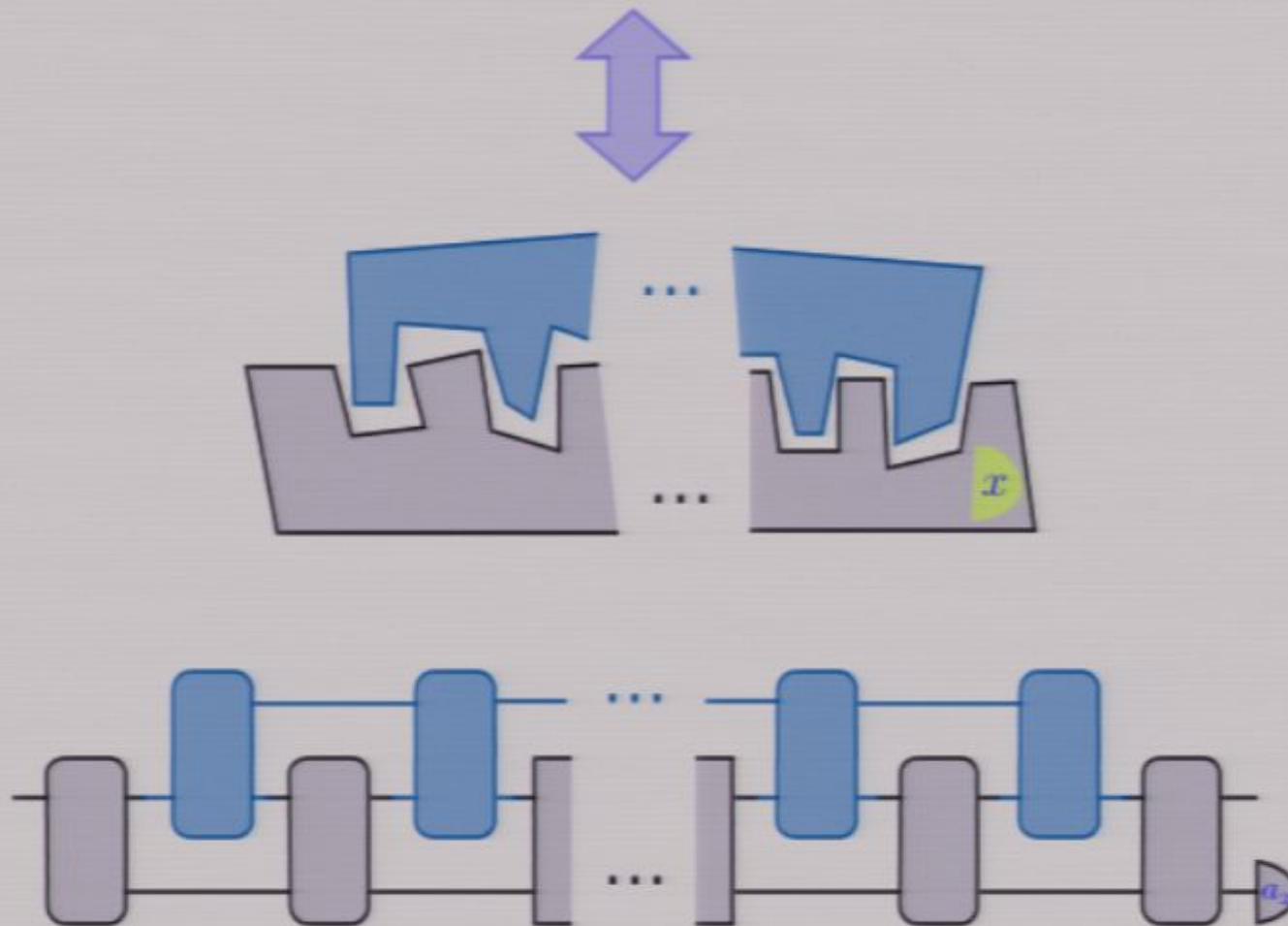
Admissibility conditions



# Realisation Theorem

---

Admissibility conditions



# Higher-order maps

---

- We want to define maps  $g$  from  $N$ -combs  $\{x\}$  to  $M$ -combs  $\{g_x\}$ 
  - $g_x$  is a map from  $(M-1)$ -combs  $\{y\}$  to transformations
- We use an “Uncurrying” procedure

$$h(x, y) := g_x(y)$$

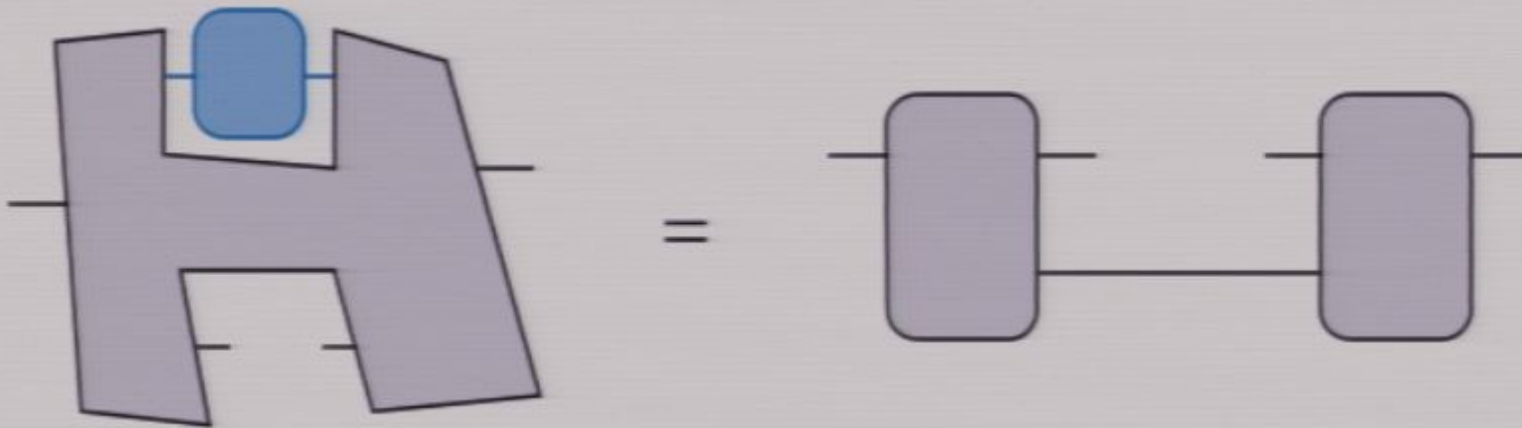
- A map from  $N$ -combs  $x$  to  $M$ -combs  $g_x$  is equivalent to a map from couples  $(x, y)$  to transformations (1-combs)



# Example

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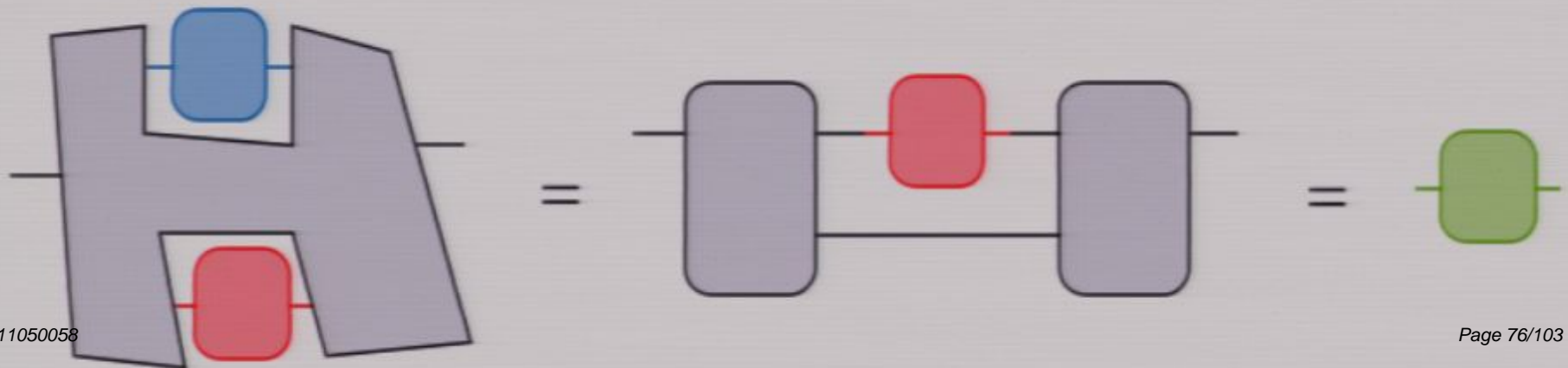
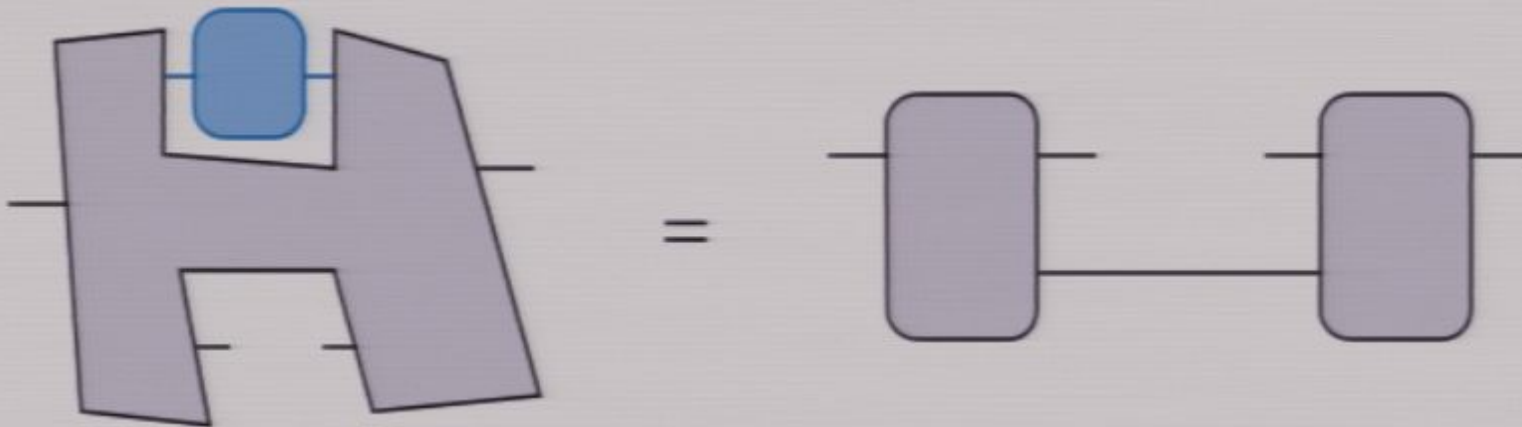
- Maps from 1-combs to 2-combs are equivalently defined as maps from couples of transformations to transformations





# Example

- Maps from 1-combs to 2-combs are equivalently defined as maps from couples of transformations to transformations



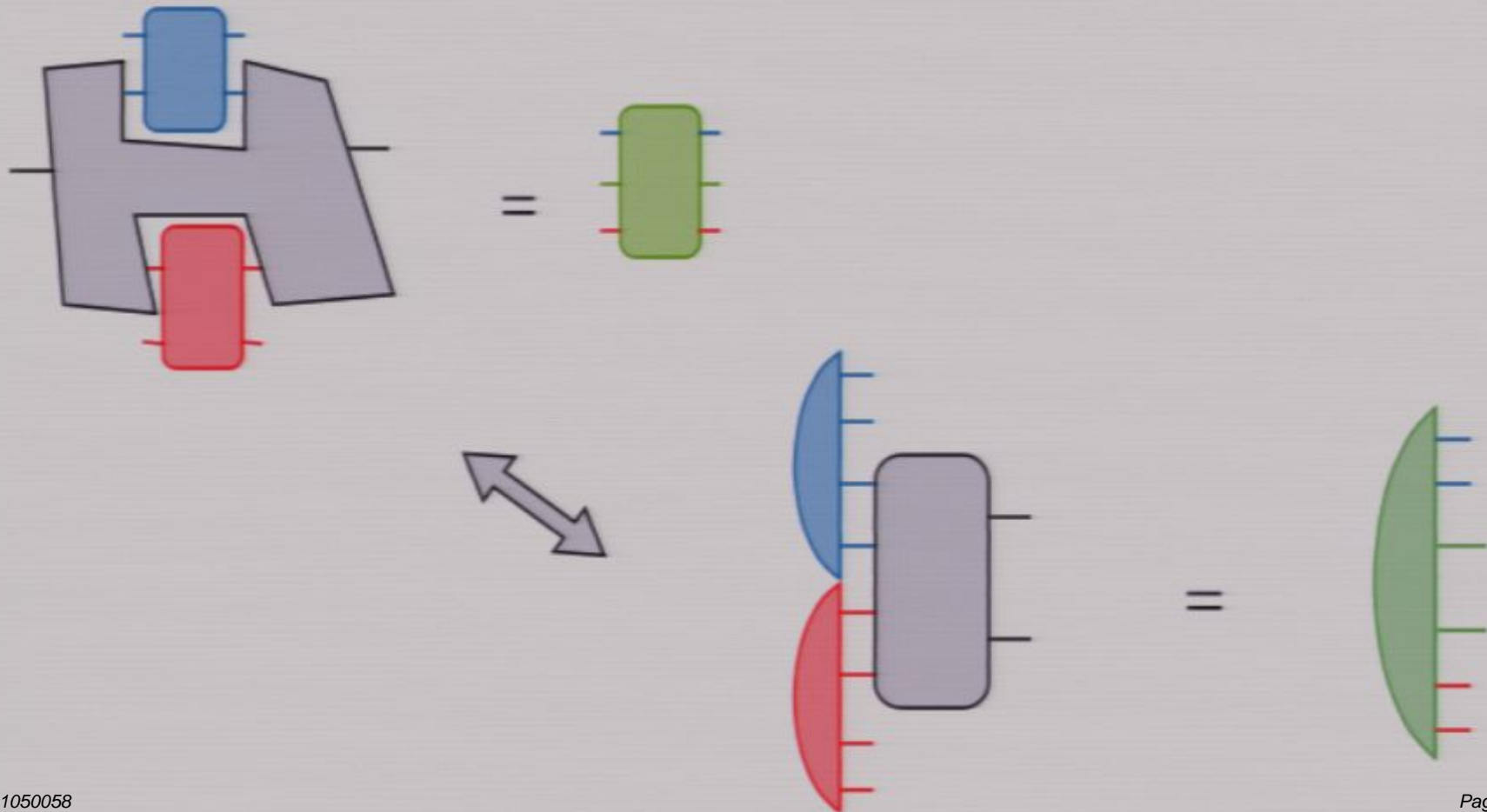
# The type of $1 \rightarrow 2$ maps

---

- Admissibility conditions on the uncurried map
  - Linearity
  - Complete Positivity
  - Normalization
- Imposed on **factorized transformations**

# Admissibility on non-signalling channels

- Complete positivity using Choi and parallel composition of faithful states



# Admissibility on non-signalling channels

---

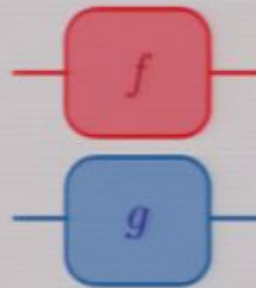
- Normalization is a linear constraint
- The non-signalling channels belong to the linear span of factorized channels
- Admissible maps are normalized on non-signalling channels

# The switch algorithm

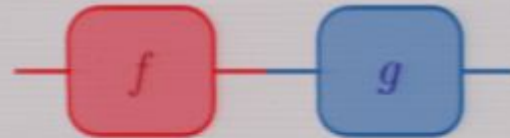
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- Input:

$x \in \{0, 1\}$ ,



- If  $x=1$ , then do



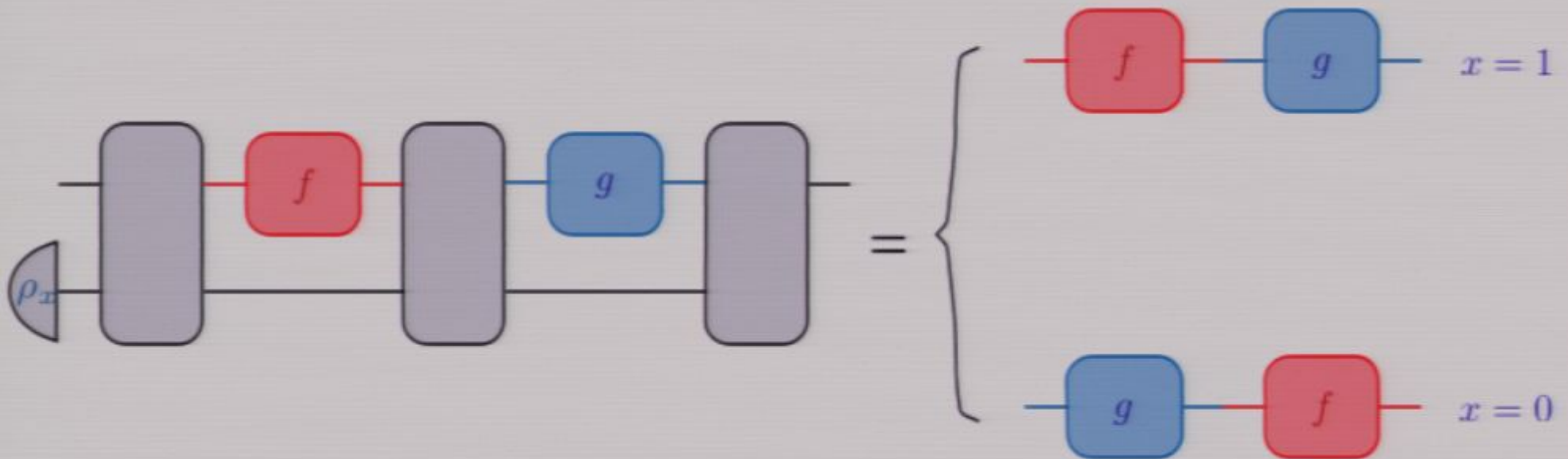
- If  $x=0$ , then do



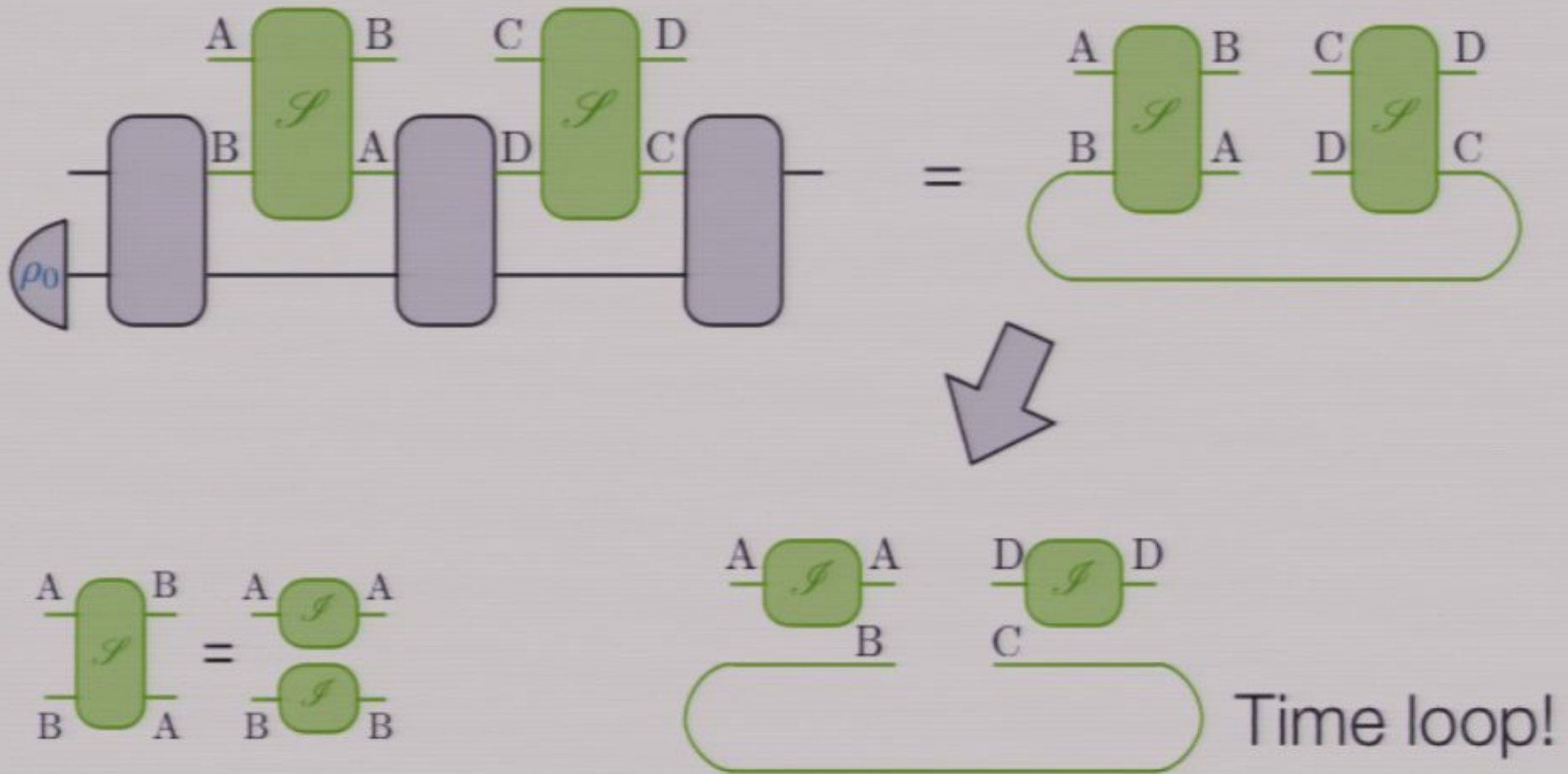
# No-switch theorem

---

Suppose a circuit exists that performs the SWITCH



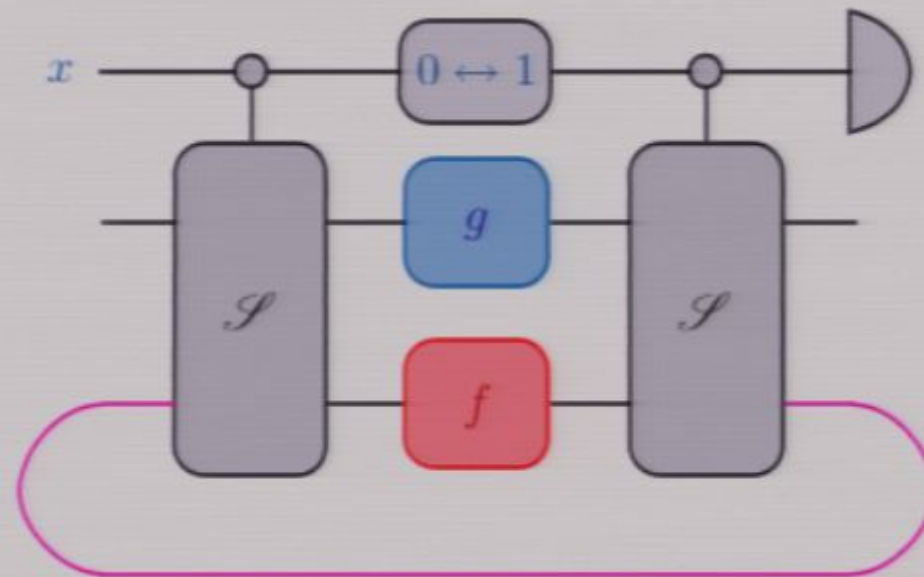
## No-switch theorem





# Equivalence of switch and time loops

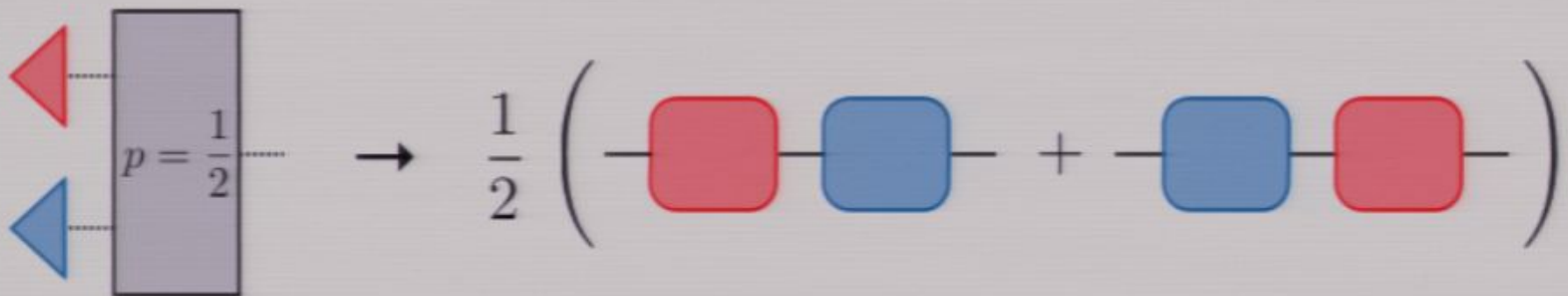
- If we had access to a **time loop** we could make a circuit for the SWITCH



# Higher order theory

---

- Higher order maps **are not** perfectly simulated in the underlying causal theory
- There exist non-circuital maps that are operationally well defined
- We lack an operational representation for convex combinations of circuits



- Analogously for superpositions

- Is there a categorical solution to the problem?

# Conjecture

---

# Conjecture

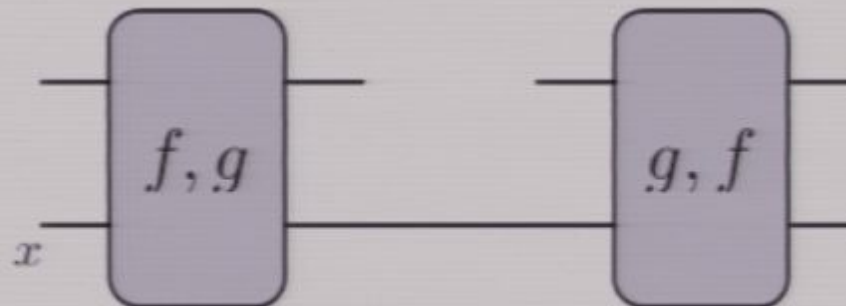
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- The operational resource: transformations  $f$  and  $g$  controlled by the input  $x$

# Conjecture

---

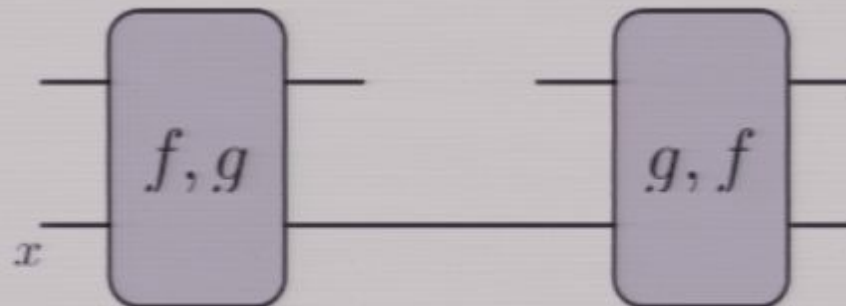
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- Operational representation comes through an oracle providing a circuit



# Conjecture

---

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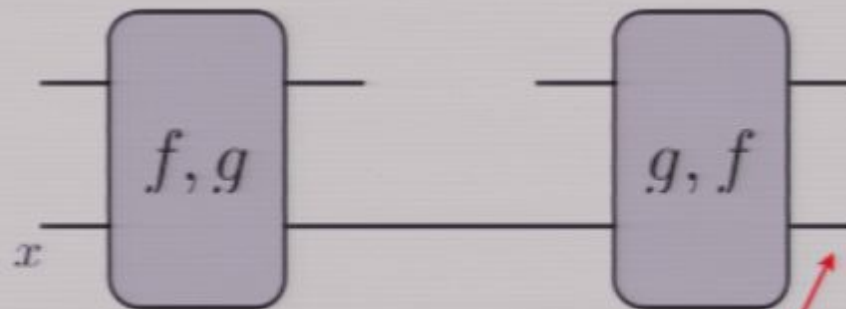




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preserves purity

# Conjecture

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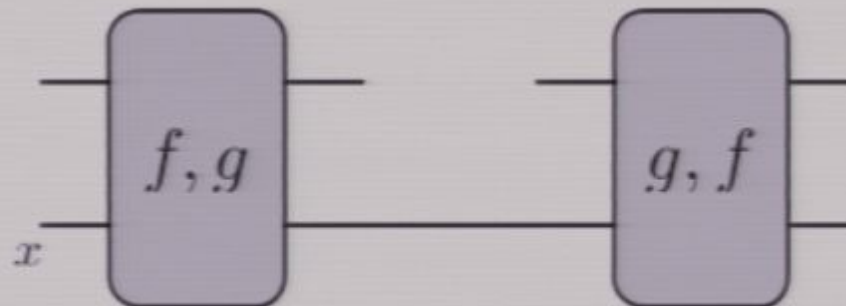


preserves purity

# Conjecture

---

- The operational resource: transformations  $f$  and  $g$  controlled by the input  $x$ 
  - Operational representation comes through an oracle providing a circuit
- The implementation of the switch becomes very simple



Can all non-causal maps be obtained by combs provided we allow for this special “oracle”?

# Higher-order maps

---

- Higher-order maps are in correspondence with multipartite states
- The purifications of such states are still admissible higher-order maps
  - Higher-order maps are not only combs
  - Higher-order maps are not only convex combinations of combs having different causal structures
  - Are all admissible maps “operational”?

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Classification: open problem

# Conclusion

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- Locality, local objective information, discord
- Factorized and non-signalling channels
- Supermaps and combs: non-causal theories
- The switch algorithm and the universality conjecture
- Non-causal theories without an immediate causal representation



# Conjecture

---

- The operational resource: transformations  $f$  and  $g$  controlled by the input  $x$



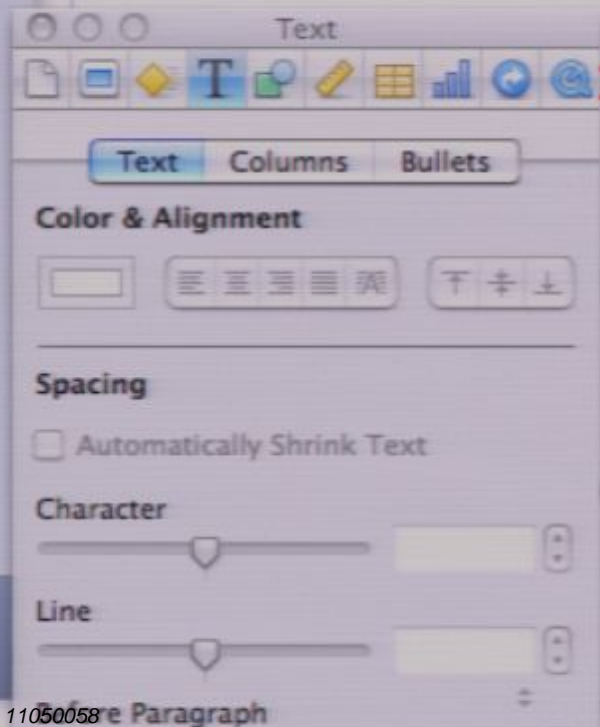


# Conjecture

Personal resource: transformations  $f$  and  $g$  control

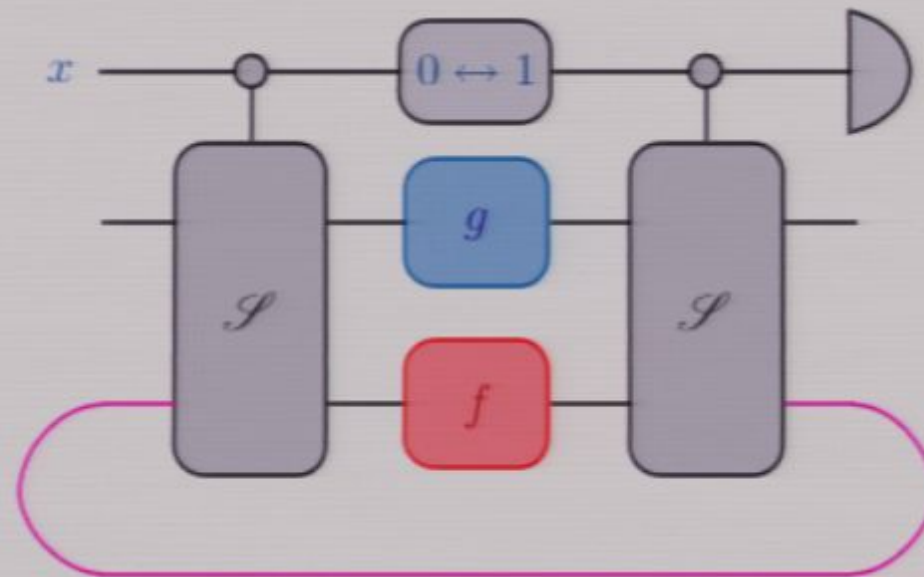
Personal representation comes through an oracle  $p$

Representation of the switch becomes very simple



# Equivalence of switch and time loops

- If we had access to a **time loop** we could make a circuit for the SWITCH

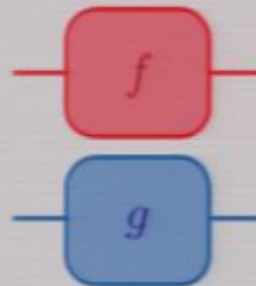


# The switch algorithm

---

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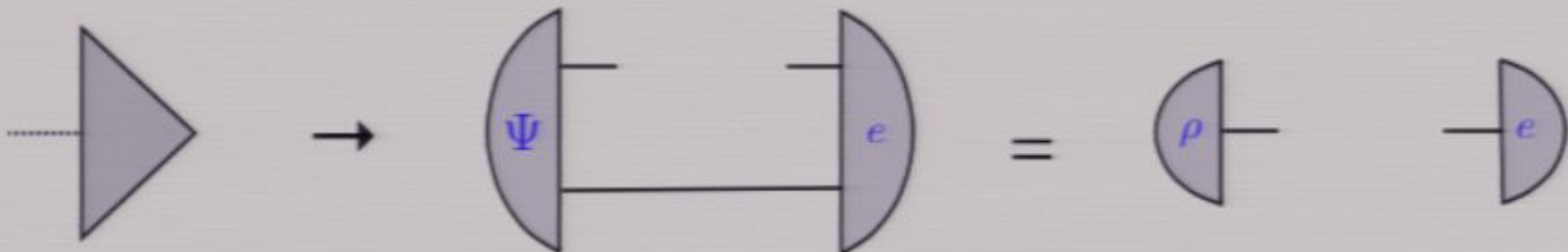
- If  $x=0$ , then do



# Second-order theory

---

- Deterministic effect  $\leftrightarrow$  deterministic tester
- The second order theory is **non causal**



# Second-order theory

---

- System types are

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# Second-order theory

- System types are  $A \rightarrow B$
- States of the system  $A \rightarrow B$  are transformations f

$A \rightarrow B$  →

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reality of a physical quantity is the possibility of predicting ... of the problem of making predictions concerning a system



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# Second-order theory

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