

Title: Non-contextual correlations in probabilistic models

Date: May 11, 2011 11:40 AM

URL: <http://pirsa.org/11050036>

Abstract: Non-contextuality is presented as an abstraction and at the same time generalisation of locality. Rather than in correlations, the underlying physical model leaves its signature in collections of expectation values, which are constrained by inequalities much like Bell's or Tsirelson's inequalities. These non-contextual inequalities reveal a deep connection to classic topics in graph theory, such as independence numbers, Lovasz numbers and other graph parameters. By considering the special case of bi-local experiments, we arrive at a semidefinite relaxation (and indeed a whole hierarchy of such relaxations) for the problem of determining the maximum quantum violation of a given Bell inequality.

Non-contextual correlations in probabilistic models

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[w/ A. Cabello, S. Severini — arXiv: 1010.2163]

1. Gleason theorem and Kochen-Specker
2. Compatibility structures
3. Example: pentagon
4. Models and inequalities
5. Commuting observables; Bell inequalities
6. ~~History~~ omitted
7. Conclusion

1. Gleason & Kochen-Specker theorems

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(2)

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(Classical) Hidden Variables, but of a weird kind:

Contextual — the same outcome P (e.g. a projector) is represented by multiple events $E \subset \mathcal{R}$ in the prob. space depending on the context in which it occurs, i.e. the different POVMs of which P is an element.

Gleason theorem. In $\dim \geq 3$, only functions $v(P) \in [0, 1]$ on projectors P s.t. $\left\{ \begin{array}{l} \text{whenever } \sum_{i=1}^n P_i = \mathbb{1}, \\ \text{then } \sum_{i=1}^n v(P_i) = 1, \end{array} \right.$ are $v(P) = \text{Tr } \rho P$ for a density operator $\rho \geq 0$. (3)

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Corollary (Bell): There doesn't exist a function $v(P) \in \{0,1\}$ on projectors s.t. $\sum_i P_i = \mathbb{1}$ implies $\sum_i v(P_i) = 1$.

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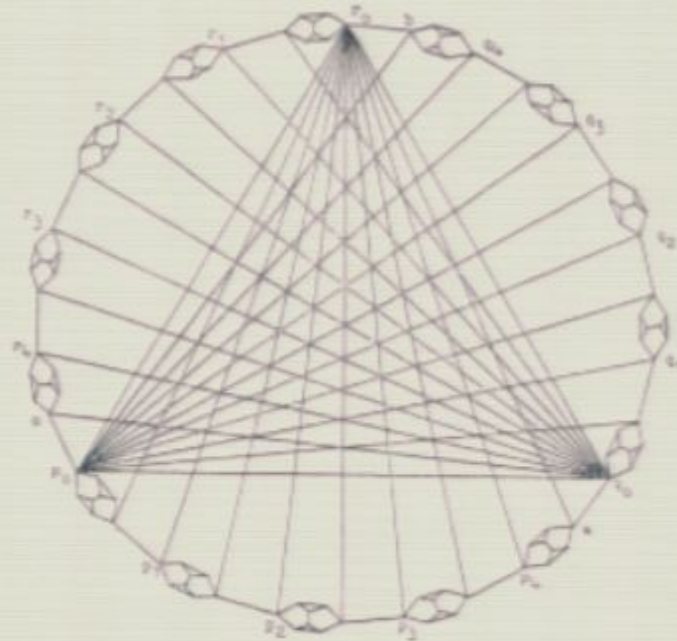
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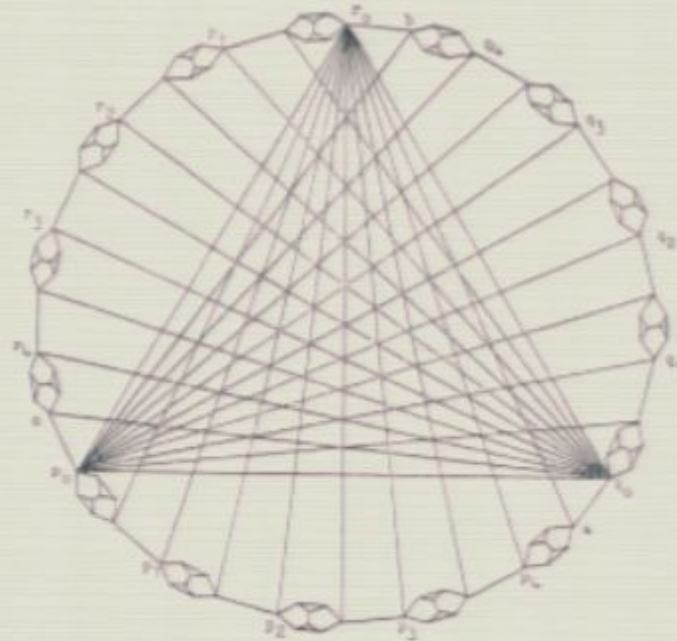
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Other arrangements:
Peres (33) } $d=3$
Conway-Kochen (31) }

Peres (24) } $d=4$
Abello et al. (18) }

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This structure is then realised/interpreted in various probabilistic models.

6

Classical / Non-contextual model:

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Aka fractional packing of the hypergraph Γ .

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- \mathcal{O} is real vector space
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3. Example: the pentagon

$$V = \{1, 2, 3, 4, 5\}$$

$$\Gamma = \{12, 23, 34, 45, 15\}$$

A classical model: $\Omega = V = \{1, 2, 3, 4, 5\}$
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(A state would be any distribution on Ω ...)

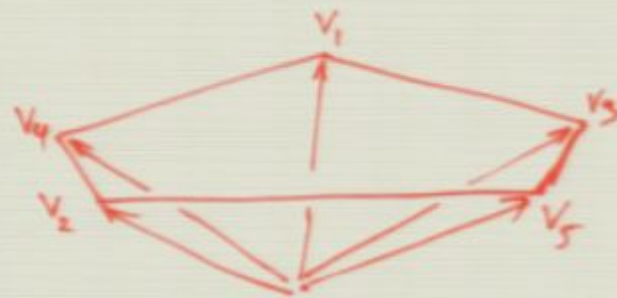
but note that always $\sum_i \langle P_i \rangle \leq 2$.

A quantum model :

[Lovász 1979; Wright 1978;
Klyachko et al. 2008]

(10)

$$\mathcal{K} = \mathbb{C}^3 \text{ and } P_i = |v_i\rangle\langle v_i|$$



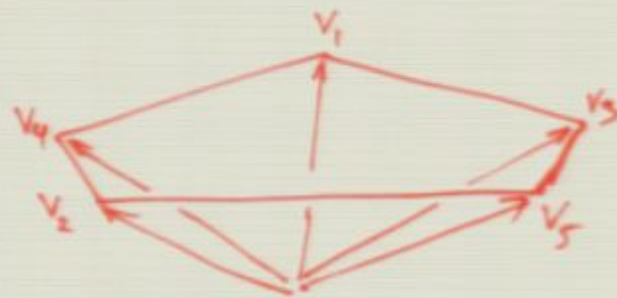
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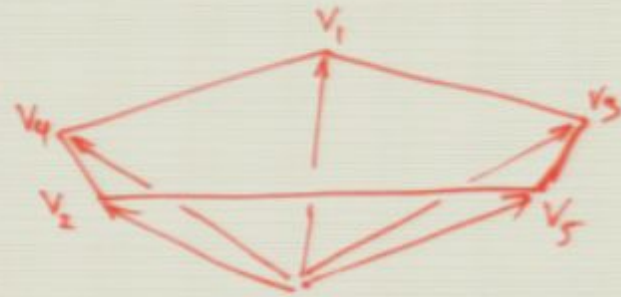
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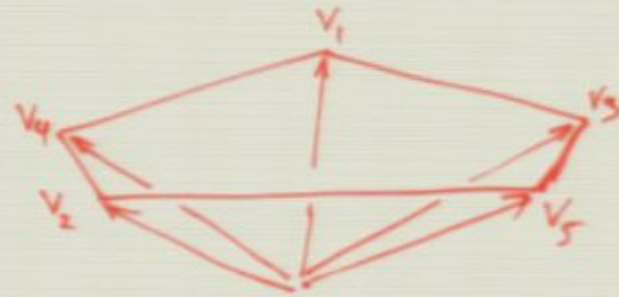
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Simply assign consistent weights $w_i \geq 0$ to $i = 1, 2, 3, 4, 5$:

E.g. $w_i = \frac{1}{2} \quad \forall i \in V$

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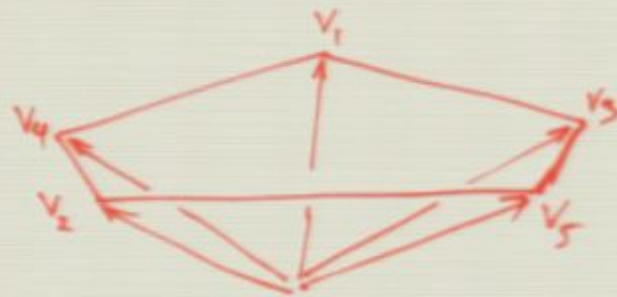
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$$\Sigma_C \subset \Sigma_{QM} \subset \Sigma_{QPT} \subset \Sigma_F \subset [0,1]^V$$

$$\Sigma_c < \Sigma_{QH} < \Sigma_{GPT} < \Sigma_F$$

(13)

$$E_C < E_{QM} < E_{GPT} = E_F$$

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$\beta_c(\Gamma)$, $\beta_{QM}(\Gamma)$, $\beta_{GPT}(\Gamma)$:
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Computations on Σ_c — in particular $\beta_c(\Gamma)$ — are NP-hard. (15)

In contrast, Σ_{GPT} is easy (polynomial time). E.g.

$\beta_{GPT}(\Gamma) = \alpha^*(\Gamma)$ "fractional packing number"

$$= \max \sum_{i \in V} w_i \quad \text{s.t. } w_i \geq 0 \\ \forall C \in \Gamma \sum_{i \in C} w_i \leq 1$$

Finally, quantum value (ρ):

$$\beta_{QH}(\Gamma) = \rho(G) \quad \text{"Lovász number"}$$

$$= \max \left\| \sum_{i \in V} |v_i\rangle \langle v_i| \right\| \quad \text{s.t. the } |v_i\rangle$$

are an orthonormal representation of G

(16)

Finally, quantum value(s):

$$\beta_{\text{QH}}(\Gamma) = \vartheta(G) \quad \text{"Lovász number"}$$

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are an orthonormal representation of G
 $\leftrightarrow [i \sim j \Rightarrow \langle v_i | v_j \rangle = 0]$

Finally, quantum value(s):

$$\beta_{\text{QH}}(\Gamma) = \nu(\mathcal{G}) \quad \text{"Lovász number"}$$

$$= \max \left\| \sum_{i \in V} |v_i\rangle\langle v_i| \right\| \quad \text{s.t. the } |v_i\rangle$$

are an orthonormal representation of \mathcal{G}
 $\left[i \neq j \Rightarrow \langle v_i | v_j \rangle = 0 \right]$

$\beta = \sum_i \langle P_i \rangle = \text{Tr} \rho \sum_i P_i$ is maximised on a pure state $\rho = |\psi\rangle\langle\psi|$. Then can restrict to 1-dim. projectors $P_i = |v_i\rangle\langle v_i|$ & compatibility enforces orthogonality.

$$\text{So, } \beta = \sum_i |\langle \psi | v_i \rangle|^2 = \langle \psi | \sum_i |v_i\rangle\langle v_i| | \psi \rangle$$

and maximum is largest eigenvalue of $\sum_i |v_i\rangle\langle v_i|$.

Lovász (1979): $\nu(G)$ is a semidefinite programme. (17)

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More generally, maximizing a linear function

$$\sum_i \lambda_i \omega_i \quad \text{over } \vec{\omega} \in \Sigma_{QH}(\Gamma) \quad [\text{w.l.o.g. } \lambda_i \geq 0!]$$

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$$\max_{\text{QH-model } i} \sum \lambda_i \langle \rho_i \rangle = \max \text{Tr } \Delta T \text{ s.t. } T \geq 0, \text{Tr } T = 1 \\ \text{and } i \neq j \Rightarrow T_{ij} = 0,$$

$$\text{where } \Delta_{ij} = \sqrt{\lambda_i \lambda_j}.$$

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Cor. Since these are semidefinite programmes, computations on Σ_{QH} are polynomial-time.

5. Commuting observables & Non-locality

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Also non-locality experiments give rise to
Compatibility structures : any observable of
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118

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Alice

settings $x \in X$
outcomes $a \in A$
projectors P_{ax}

Bob

settings $y \in Y$
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Also non-locality experiments give rise to compatibility structures : any observable of Alice commutes with any of Bob.

Alice

settings $x \in X$
outcomes $a \in A$
projectors P_{ax}

Bob

settings $y \in Y$
outcomes $b \in B$
projectors P_{by}



atomic events $ab|xy$ with associated projectors $P_{ab|xy} = P_{ax} P_{by} = P_{by} P_{ax}$

Compatibility structure :

$$V = \{ abxy \} = \mathcal{A} \times \mathcal{B} \times \mathcal{X} \times \mathcal{Y}$$

\mathcal{G} has an edge $abxy \sim a'b'x'y'$ iff $(x=x' \text{ and } a \neq a')$
or $(y=y' \text{ and } b \neq b')$

Hypergraph Γ of all cliques of \mathcal{G}

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Hypergraph Γ of all cliques of G

Note: G (and Γ) captures coexistence and mutual
exclusiveness of elementary events $abxy$. But
not that for given xy , one of $abxy$ ($a \in \mathcal{A}, b \in \mathcal{B}$)
has to occur with probability 1...

So, add the constraint that $\forall xy \sum_{ab} w_{ab|xy} = 1 :$

$$\sum_{\cup} \rho_{PT}(\Gamma)$$

$$\sum_{\cup} \rho_{QT}(\Gamma)$$

$$\sum_c(\Gamma)$$

(20)

So, add the constraint that $\forall xy \sum_{ab} w_{ab|xy} = 1$:

$$\sum_{\cup} \sigma_{\Gamma PT}(\Gamma) > \sum_{\cup} \sigma'_{\Gamma PT}(\Gamma)$$

$$\sum_{\cup} \sigma_{\Omega \Pi}(\Gamma) > \sum_{\cup} \sigma'_{\Omega \Pi}(\Gamma)$$

$$\sum_c \sigma_c(\Gamma) > \sum_c \sigma'_c(\Gamma)$$

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$$\sum_{\cup} \rho_{QT}(\Gamma) > \sum_{\cup} \rho_{QT}(\Gamma) > \sum_{\cup} \mathbb{1}_{QT}(\Gamma)$$

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$$\forall xy \sum_{ab} P_{ab|xy} = \mathbb{1}$$

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$$\Sigma_{\text{qPT}}(\Gamma) \supset \Sigma_{\text{qPT}}^1(\Gamma)$$

$$\Sigma_{\text{QM}}(\Gamma) \supset \Sigma_{\text{QM}}^1(\Gamma) \supset \Sigma_{\text{QM}}^{\perp}(\Gamma)$$

$$\Sigma_C(\Gamma) \supset \Sigma_C^1(\Gamma)$$

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- Proposition :
- $\Sigma_C^1(\Gamma)$ are exactly the local (hidden variable) corr.
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(21)

Furthermore, maximum values of expressions like

$$\sum_{abxy} \lambda_{abxy} P\{ab|xy\}$$

[w.l.o.g. with $\lambda_{abxy} \geq 0$] over local / no-signalling correlations, i.e. over $\Sigma_C / \Sigma_{\text{NPT}}$ are the same as maximised over $\Sigma_C / \Sigma_{\text{GPT}}$.

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I.e. local (and no-signalling) values of a Bell-type correlator can be interpreted 1-1 as non-contextual (and Gleason) values of a canonically associated compatibility structure.

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I.e. local (and no-signalling) values of a Bell-type correlator can be interpreted 1-1 as non-contextual (and Gleason) values of a canonically associated compatibility structure.

Maximum quantum violations can be upper bounded by the (efficiently computable!) maximum over Σ'_{QFT}

}??

Example 1: CHSH ($\mathcal{A} = \mathcal{B} = \mathcal{X} = \mathcal{Y} = \{0, 1\}$)

Expression to maximise is

$$\sum_{\substack{a, b, x, y \\ \text{s.t.} \\ a \oplus b = xy}}$$

$w_{ab|xy}$, I.e.

$$\lambda_{ab|xy} = \begin{cases} 1 & \text{if } a \oplus b = xy, \\ 0 & \text{otherwise.} \end{cases}$$

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$$\sum_{\substack{abxy \\ \text{s.t.} \\ a \oplus b = xy}} w_{ab|xy} \cdot I.e. \quad \lambda_{abxy} = \begin{cases} 1 & \text{if } a \oplus b = xy \\ 0 & \text{otherwise} \end{cases}$$

Turns out that local, quantum and no-signalling values of this function are just

$$\omega(\mathcal{F}) = 3,$$

$$\omega(\mathcal{F}) = 2 + \sqrt{2},$$

$$\omega^*(\mathcal{F}) = 4, \text{ resp.}$$

of the following 8-vertex graph $\mathcal{F} = C_8(1, 4)$:



Example 2 : I_{3322} [Collins/Gisin, J. Phys A 2004] (23)
 ($\mathcal{A} = \mathcal{B} = \{0, 1\}$; $\mathcal{X} = \mathcal{Y} = \{0, 1, 2\}$)

Maximise $\sum_{a,b,x,y} \lambda_{ab|xy} w_{ab|xy}$

where:
 $\lambda =$

$y \backslash x$	00	01	10	11	20	21
00	1	0	1	0	1	0
01	0	0	1	1	1	1
10	1	1	1	0	0	1
11	0	1	1	1	1	1
20	1	0	0	1	0	0
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Local value (max over Σ_c^1)
 is 6, the independence number
 of a $2a$ -vertex graph.

Maximum over Σ_{QH}^1 is ≈ 6.25747

[Currently best upper bound on quantum value is 6.25087556
 using the Navascués-Acín-Pironio hierarchy of SDP relaxations of
 Σ_{QH}^1 (Pal/Varkesi, arXiv:1006.3032), lower bound 6.25087538]

7. Conclusion / Questions

- Noncontextual inequalities give a new look on Gleason - and Kochen-Specker theorems.
- Unlike Bell inequalities, it is very easy to find examples of very large violations:

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- Noncontextual inequalities give a new look on Gleason - and Kochen-Specker theorems.
- Unlike Bell inequalities, it is very easy to find examples of very large violations:
 \exists graphs on n vertices with $\alpha(L^{(n)}) = O(\log n)$,
but $\beta(L^{(n)}) \approx \sqrt{n}$;
and $\alpha(C^{(n)}) = 3$, but $\beta(C^{(n)}) \approx 4\sqrt{n}$.

- 30
- Open: Can any violation of a non-contextual quantum inequality be converted into a (comparably large?) violation of a Bell inequality?

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- How different are Σ'_{QM} and Σ_{QM}^{\perp} in the non-local setting?

How good is the maximum over Σ'_{QM} of a Bell inequality as upper bound on the maximum quantum violation?

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VGA-1

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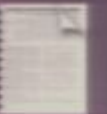
VGA-1

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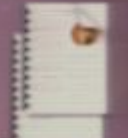
VGA-1



Macintosh HD



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RosarioK...0403.pdf



Monras -
fractalstate space.pdf



Increasing
entangle...111.pdf



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