

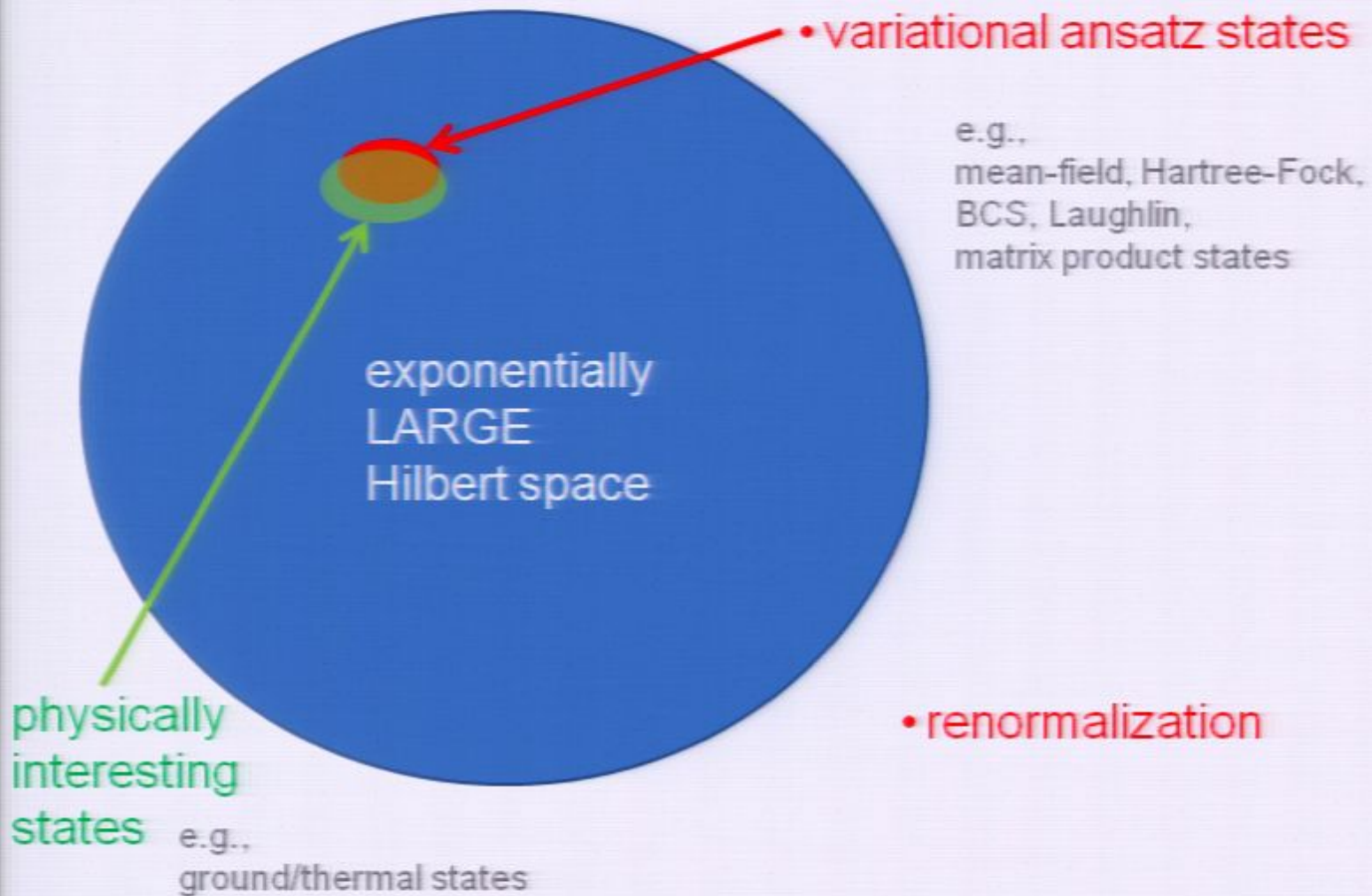
Title: Anyonic entanglement renormalization

Date: Oct 13, 2010 04:00 PM

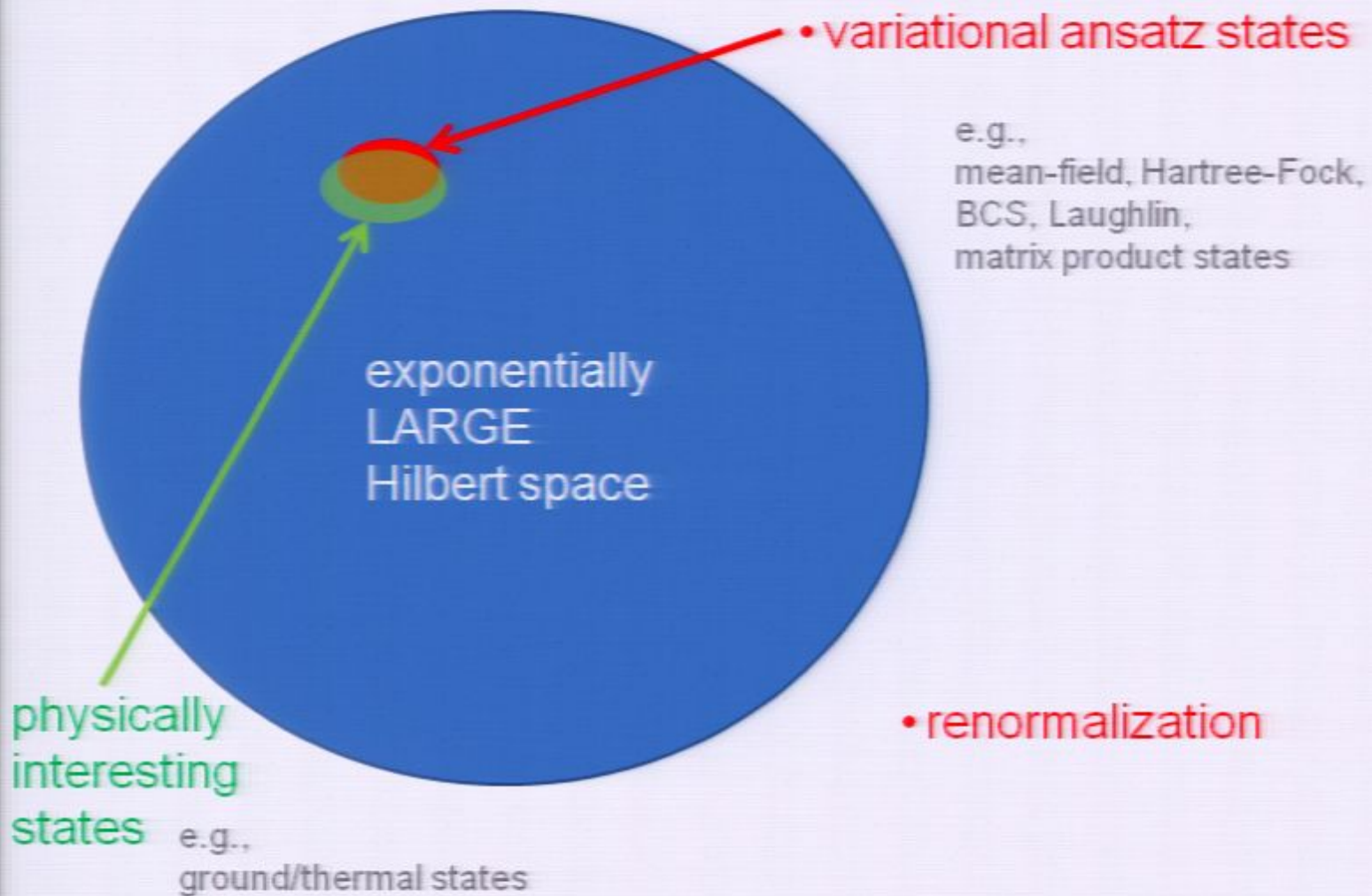
URL: <http://pirsa.org/10100045>

Abstract: We introduce a family of variational ansatz states for chains of anyons which optimally exploits the structure of the anyonic Hilbert space. This ansatz is the natural analog of the multi-scale entanglement renormalization ansatz for spin chains. In particular, it has the same interpretation as a coarse-graining procedure and is expected to accurately describe critical systems with algebraically decaying correlations. We numerically investigate the validity of this ansatz using the anyonic golden chain and its relatives as a testbed. This demonstrates the power of entanglement renormalization in a setting with non-abelian exchange statistics, extending previous work on qudits, bosons and fermions in two dimensions. This is joint work with Ersen Bilgin.

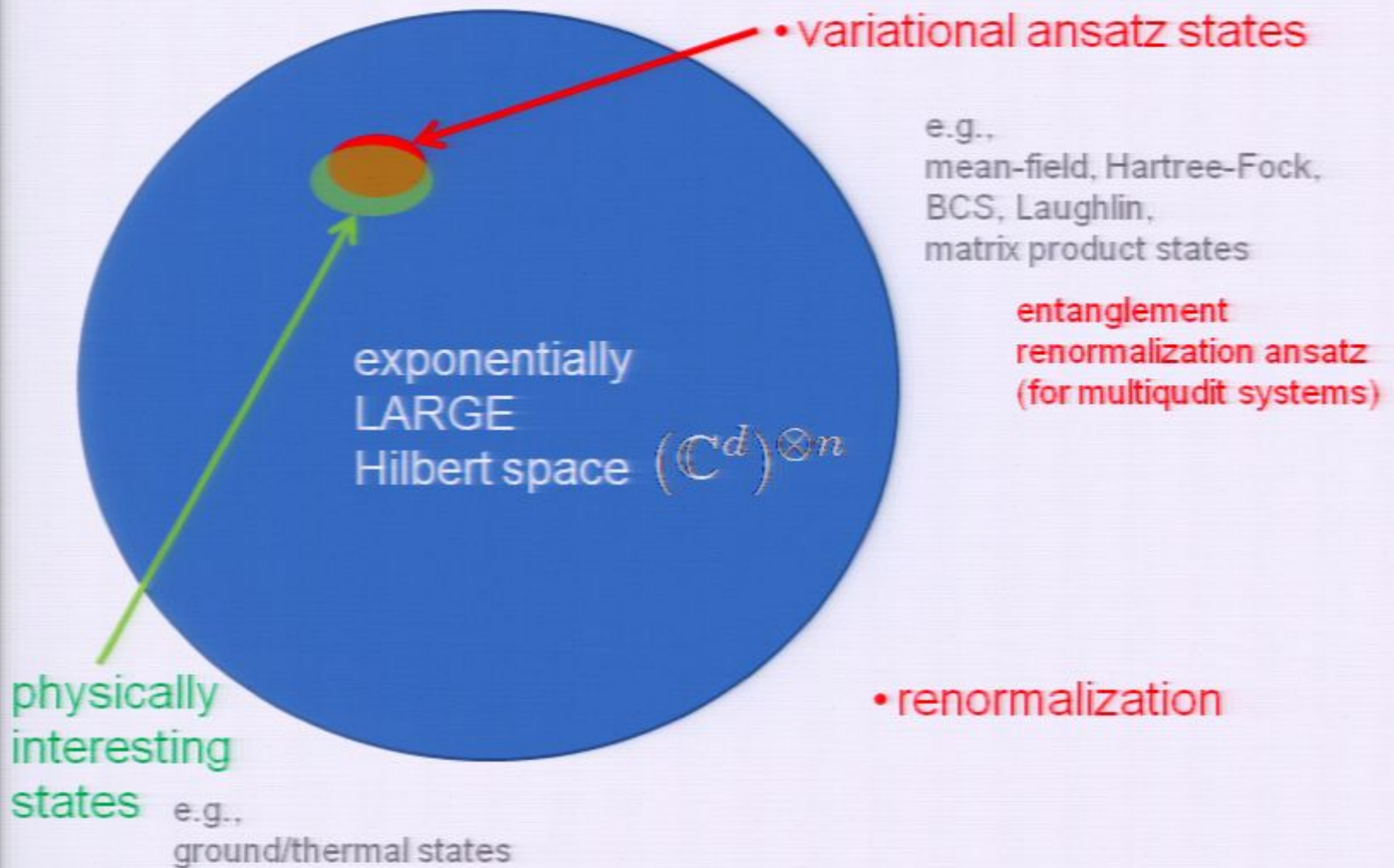
Many-body quantum physics: approaches



Many-body quantum physics: approaches



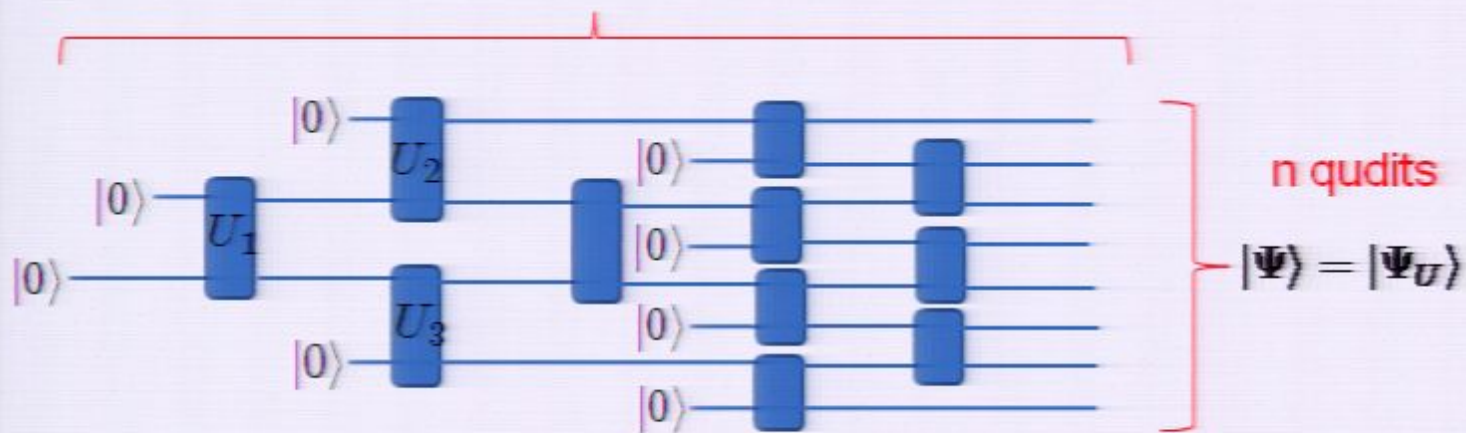
Many-body quantum physics: approaches



Entanglement renormalization for spins

G. Vidal, Phys. Rev. Lett. 99, 220405 (2007)

efficient $(O(\log n)\text{-depth})$ quantum circuit preparing n -qudit state

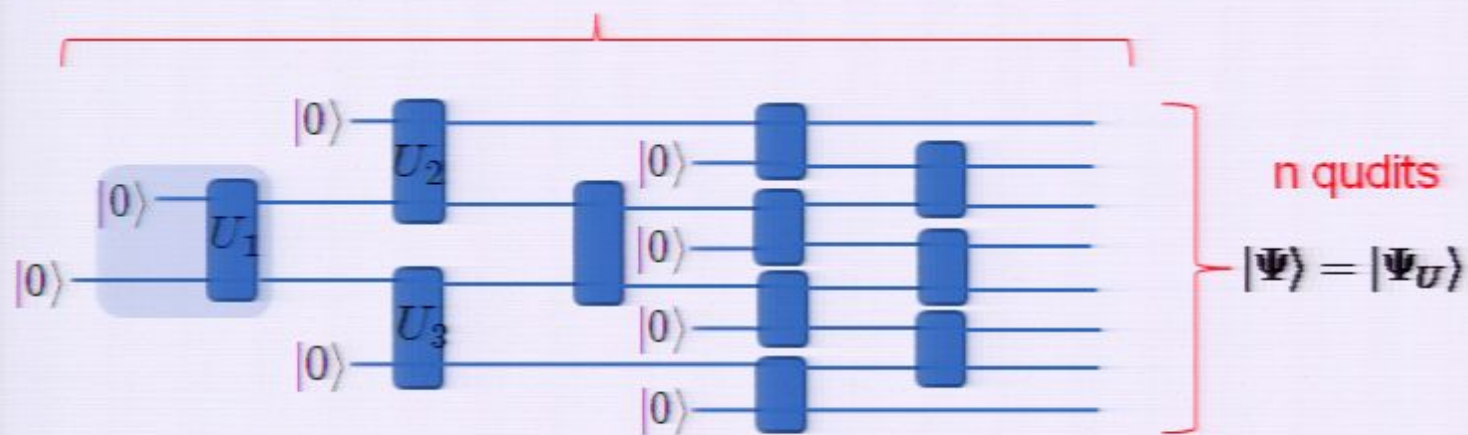


Varying over gates $U = \{U_i\}_i$ gives a variational family of states

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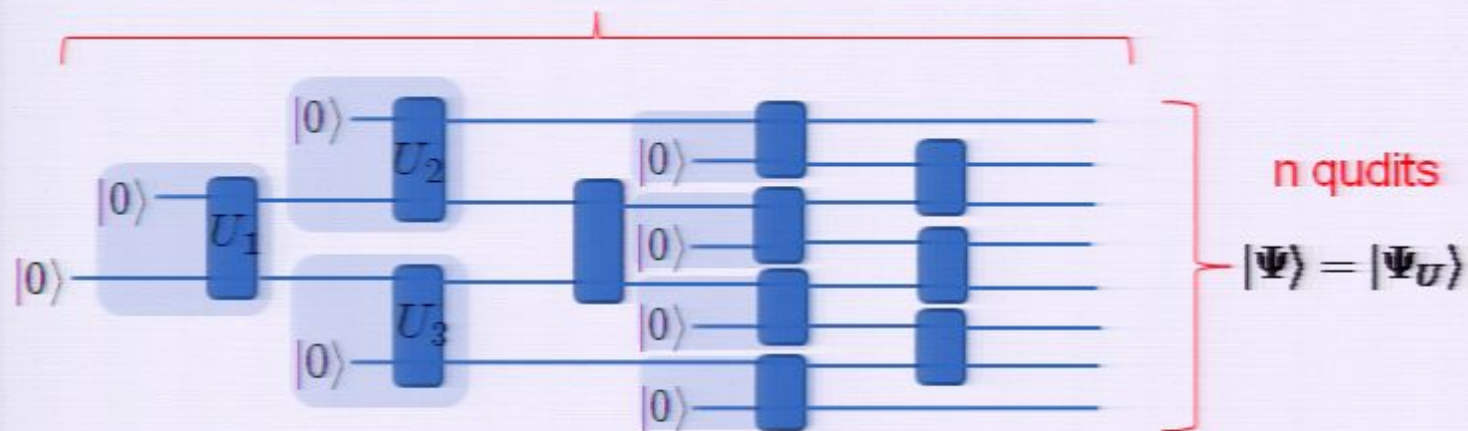


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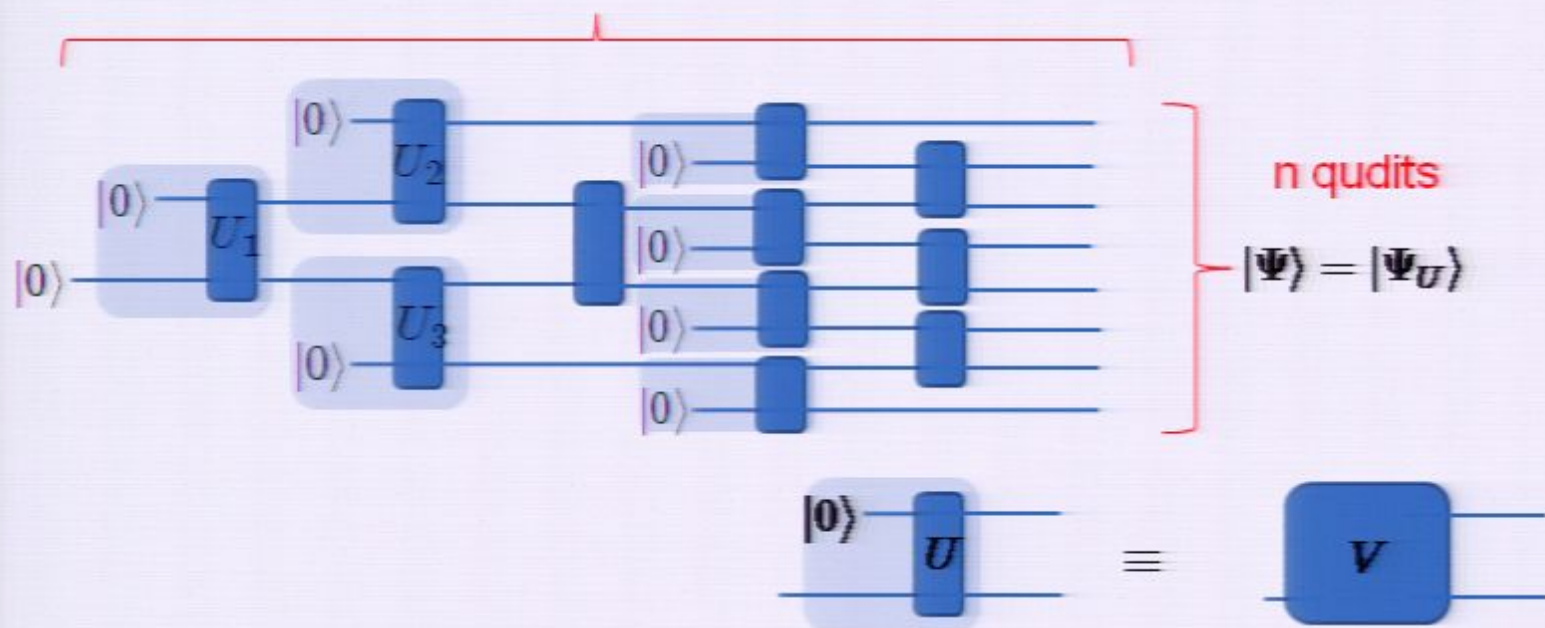


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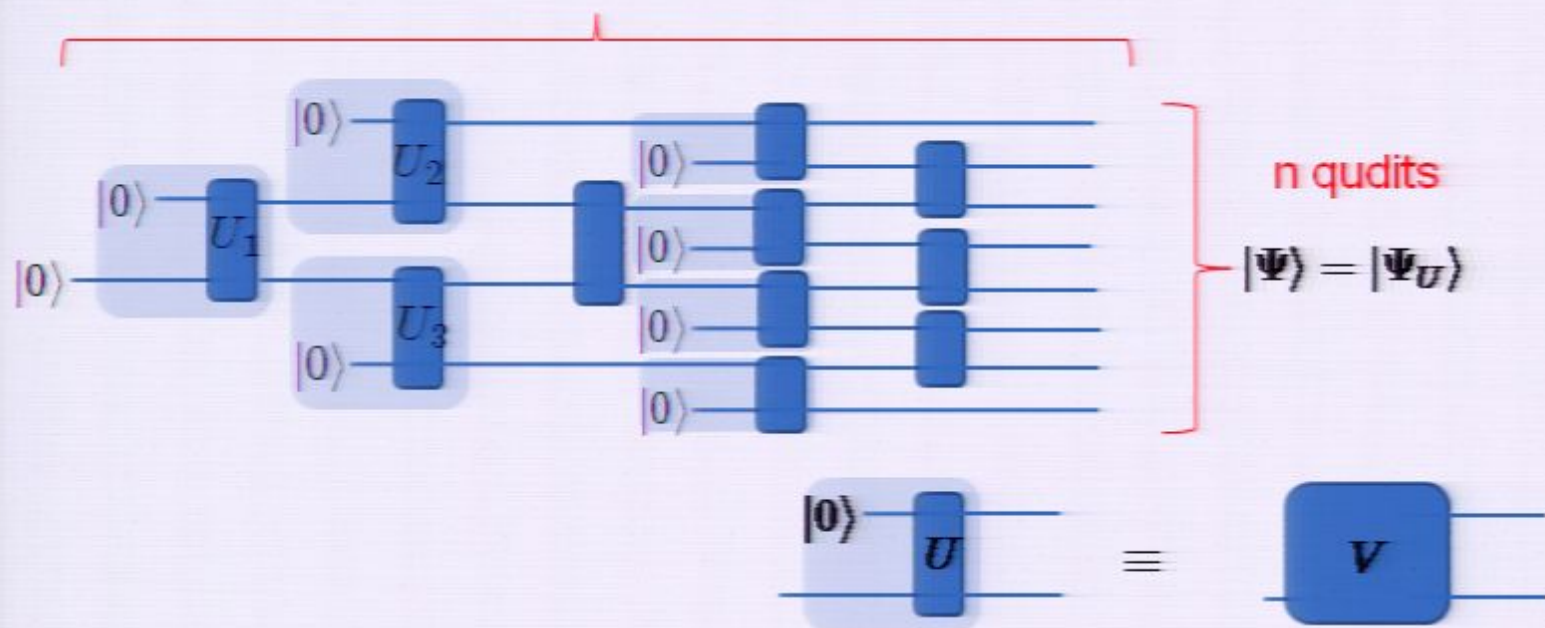


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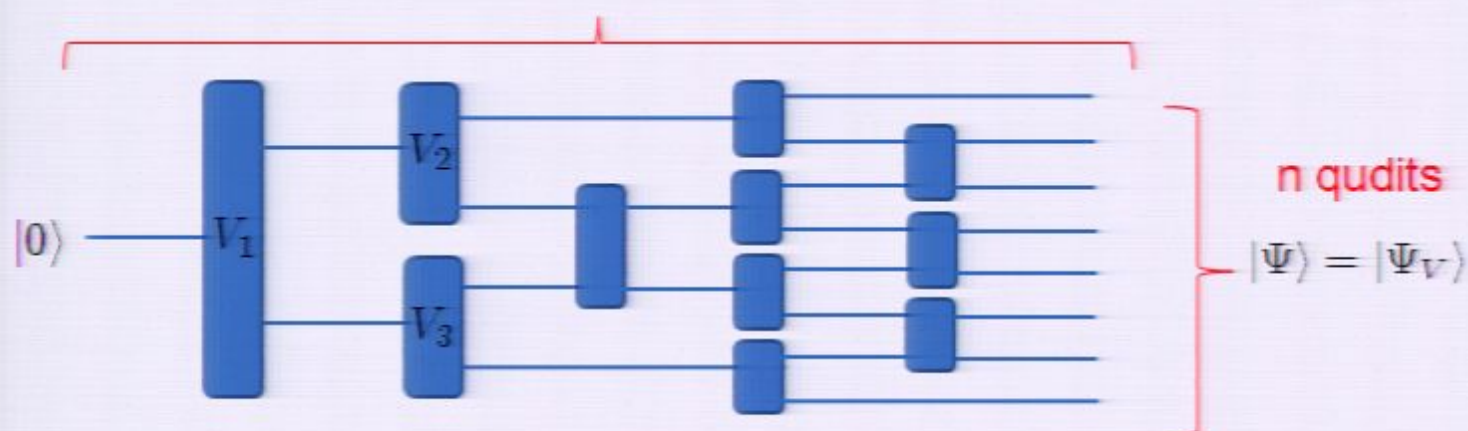


Varying over isometries $V = \{V_i\}_i$ gives a variational family of states

Entanglement renormalization for spins

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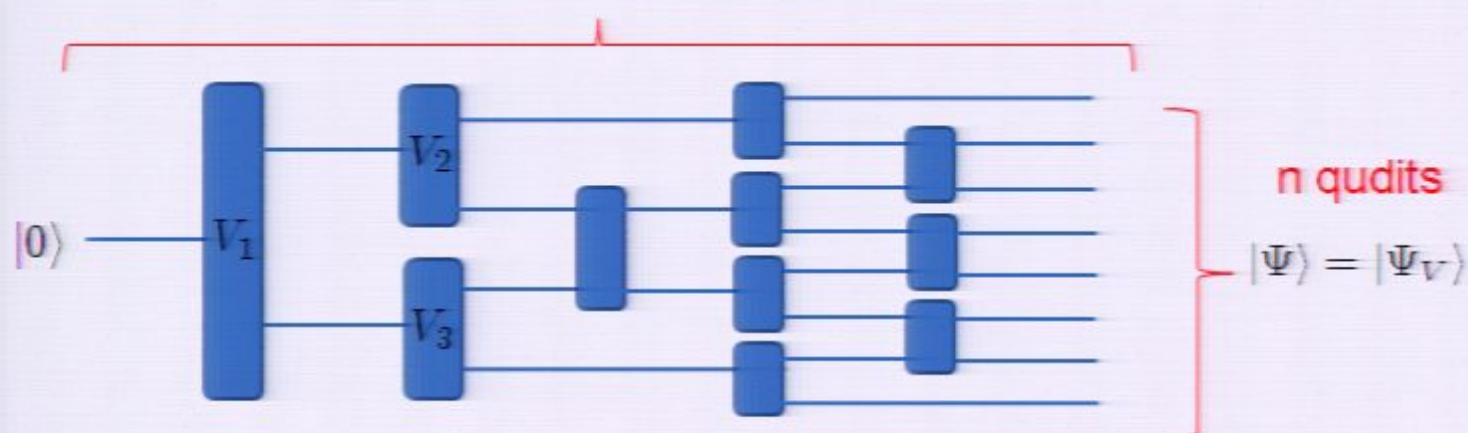


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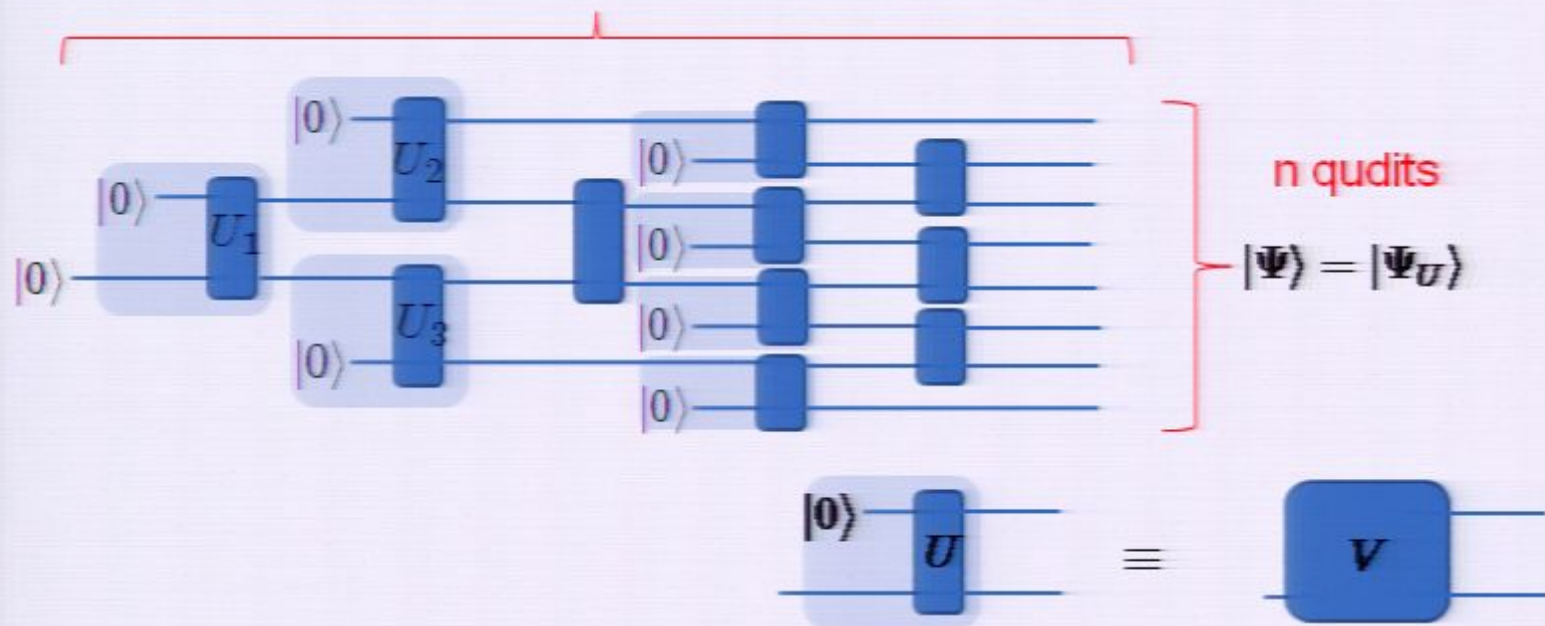


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Entanglement renormalization for spins

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Varying over **isometries** $V = \{V_i\}_i$ gives a variational family of states

Entanglement renormalization

Evidence that this is a **good ansatz**:



- **good approximations** to ground state energies & correlation functions of **critical systems**
(Evenbly, Vidal '09 & Vidal '07)

- describes states with $S(\rho_L) \approx \log L$ and algebraically decaying correlations in 1D

- **exact description** of interesting **topologically ordered states**

(Aguado/Vidal '07 & K. Reichardt, Vidal '08)

Useful as a numerical method because:

- **efficient evaluation** of expectation values
correlation function
critical exponents
- **generalizes** to $D > 1$ and to e.g., fermions

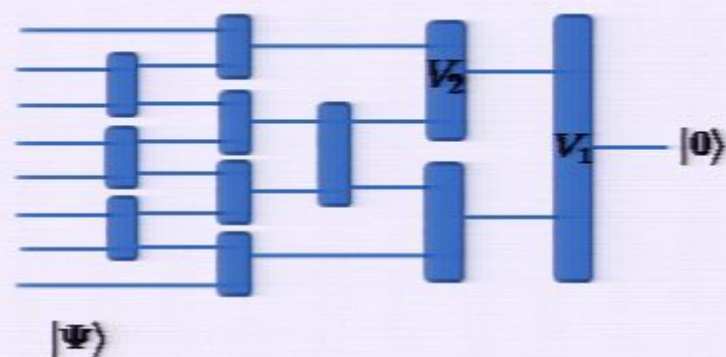
(Evenbly, Vidal '07, Pineda et al. '09)

- iterative **algorithms** for minimizing energy

Efficient evaluation of expectation values

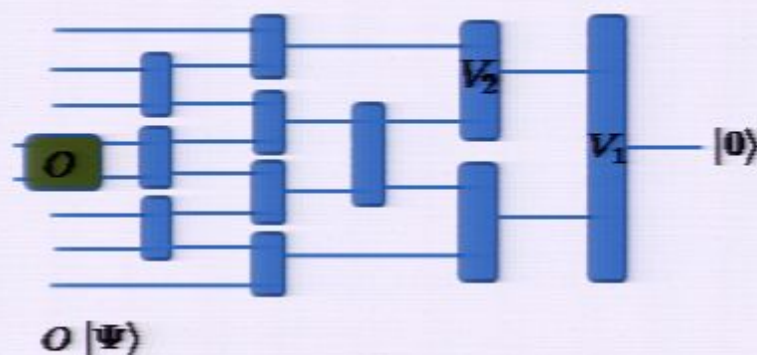
The expectation value of a *local observable* can be efficiently computed as only isometries within its “causal cone” contribute.

Efficient evaluation of expectation values



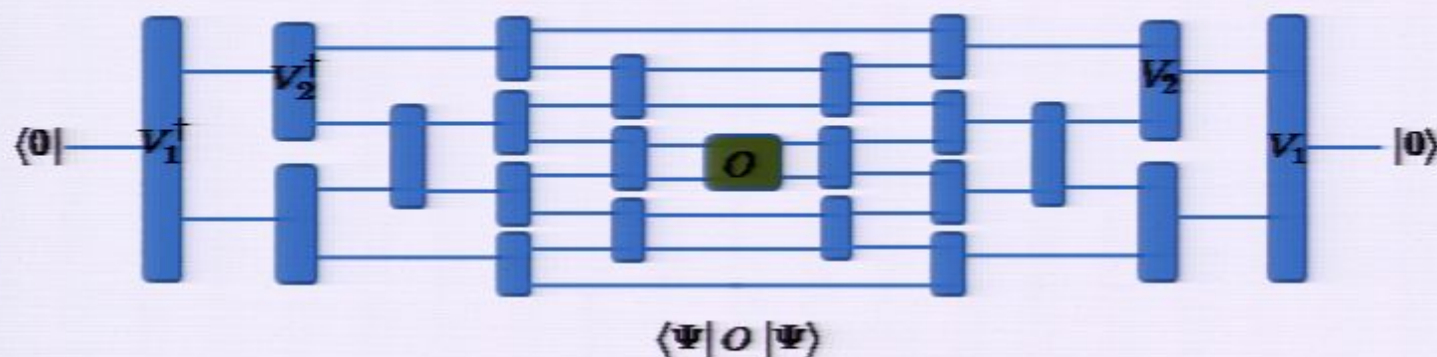
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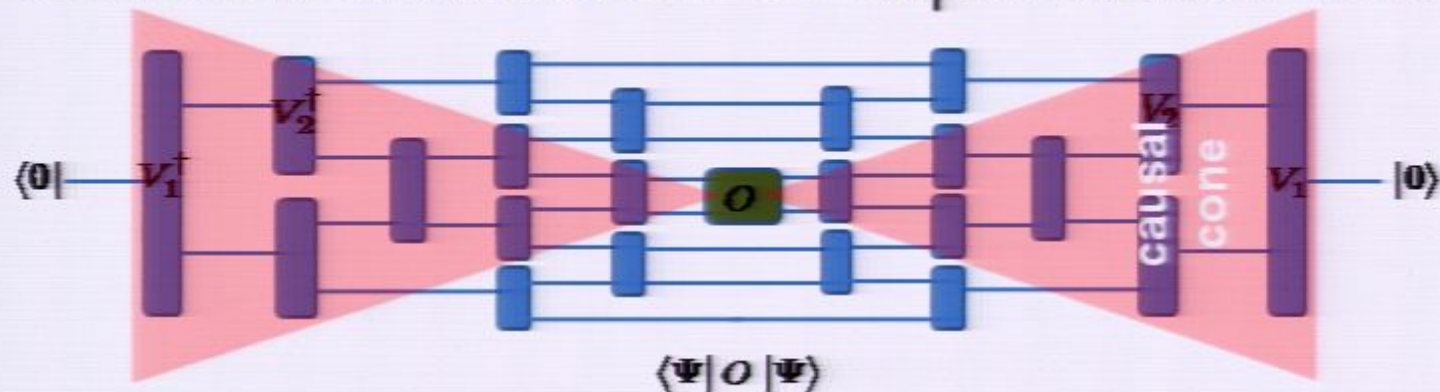
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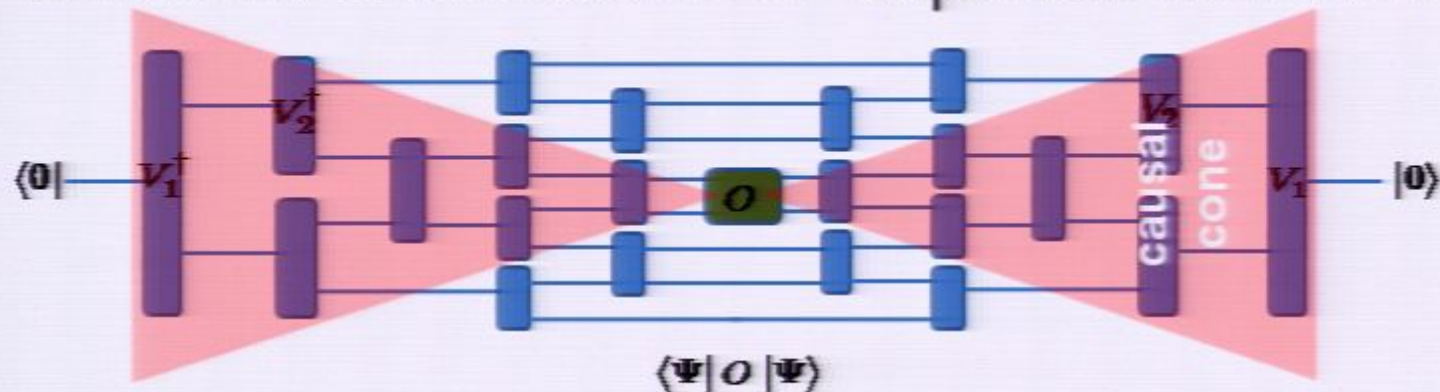
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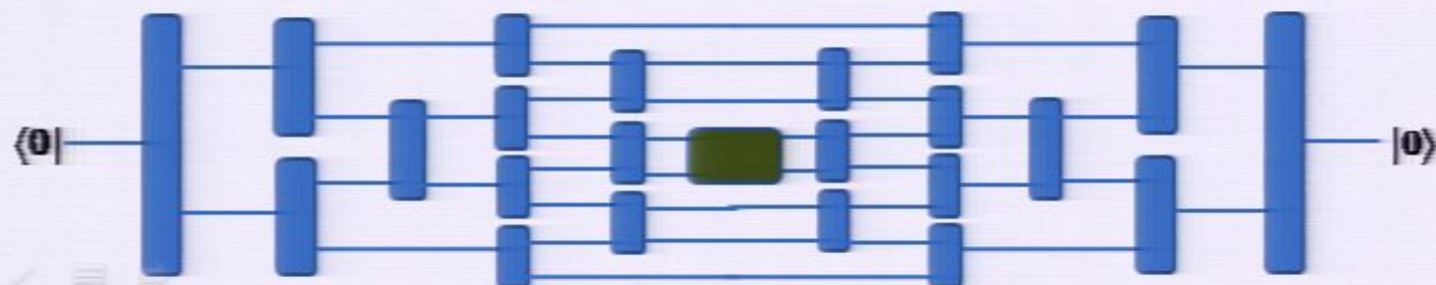
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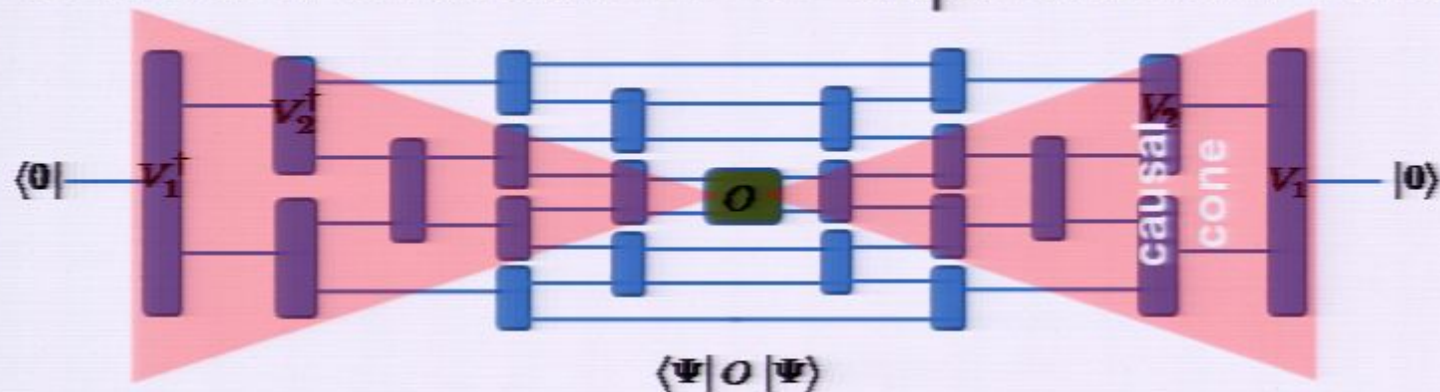
The expectation value of a *local observable* can be efficiently computed as only isometries within its "causal cone" contribute.

This follows from the defining property of isometries (resp. unitaries)

$$V^\dagger V = \text{id}$$



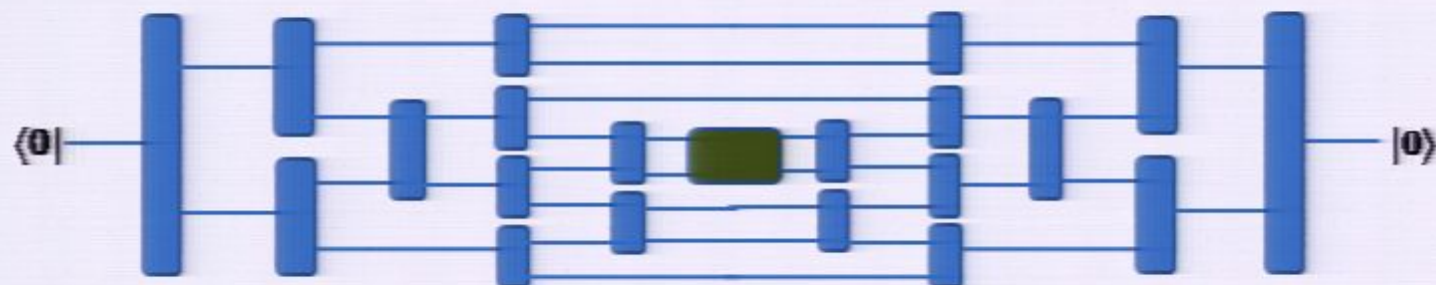
Efficient evaluation of expectation values



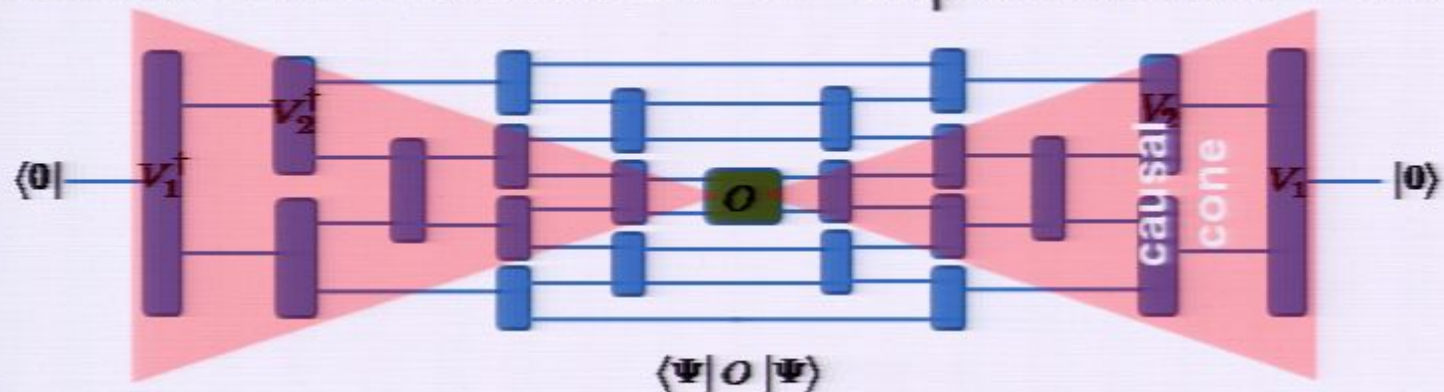
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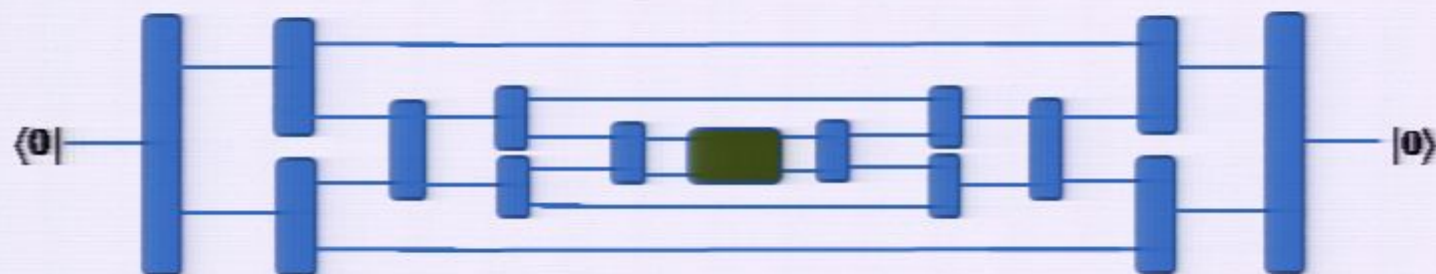
Efficient evaluation of expectation values



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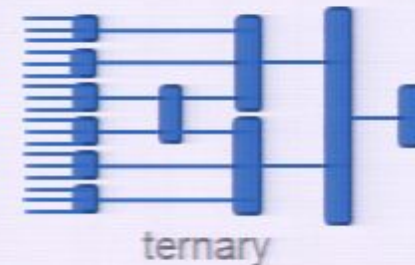
$$V^\dagger V = \text{id}$$



Entanglement renormalization: Summary

1. Choose a suitable tree-like graph

necessary condition: sufficiently "contracting" (causal cone of finite width)



2. Associate variational parameters



qudit state
 $|0\rangle^{\otimes 3} \in (\mathbb{C}^d)^{\otimes 3}$



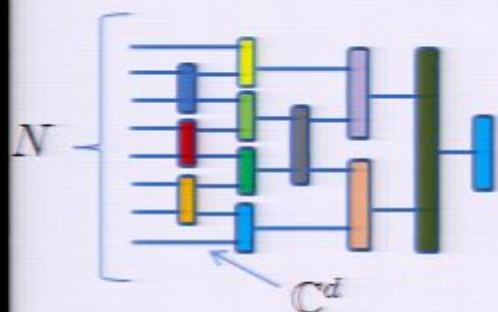
isometry
 $V: (\mathbb{C}^d)^{\otimes 2} \rightarrow (\mathbb{C}^d)^{\otimes 3}$

3. Compute expectation values from using substitutions



and **contracting** resulting
 simplified **tensor network**

Parameter counting



general

$$O(\text{poly}(d)N \log N)$$



translation-invariant

$$O(\text{poly}(d) \log N)$$



scale-invariant

$$O(\text{poly}(d))$$

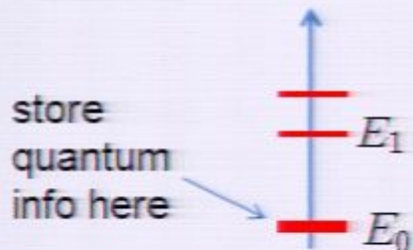
Introduction to anyons

- Background
- The anyonic Hilbert space
- Conservation of topological charge
- Anyonic Hamiltonians
- Computational rules

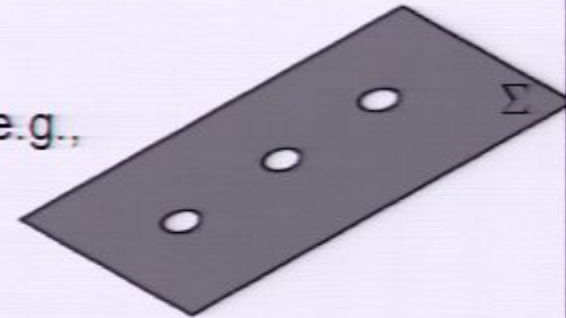
Topological order

Characteristic properties:

- ground space **degeneracy** depending on Σ
- **finite gap**, “stable” under local perturbations



2D-surface, e.g., plane with punctures



1) Embed “quantum degrees of freedom” into Σ , e.g.,

- qudits placed on edges of a triangulation of Σ
- electrons confined to Σ

2) Assume certain “topological” Hamiltonian, e.g.,

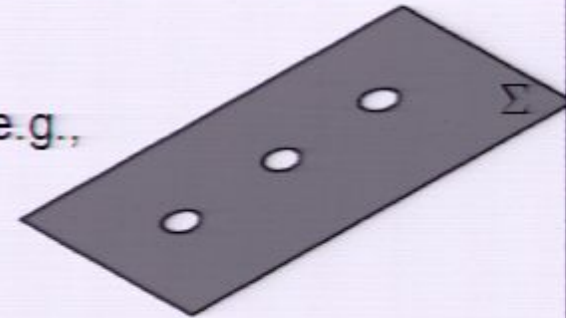
- Levin-Wen-stringnet or Kitaev-toric code Hamiltonian
- Chern-Simons-action

Topological order

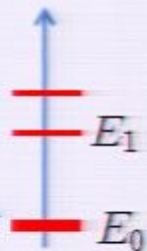
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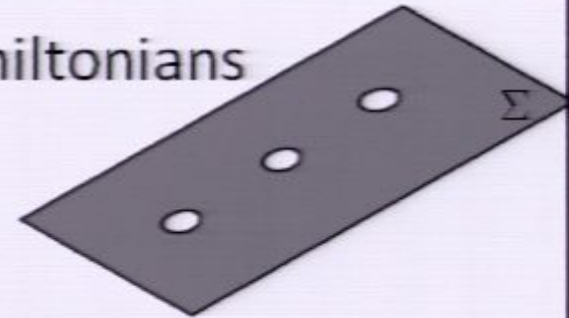
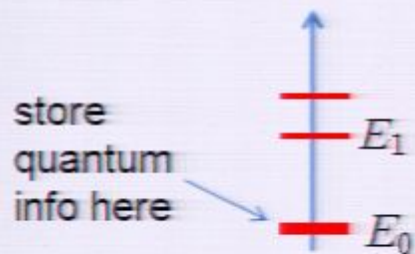
store
quantum
info here



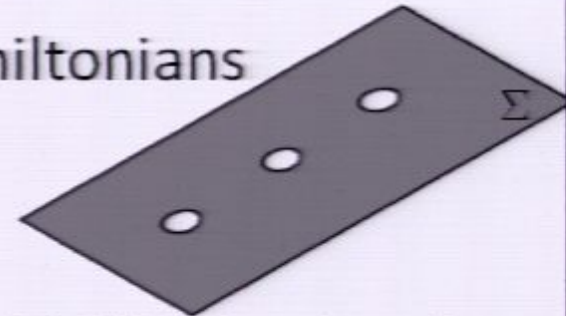
Local perturbations and effective Hamiltonians

Characteristic
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- **finite gap**, “stable” under local perturbations



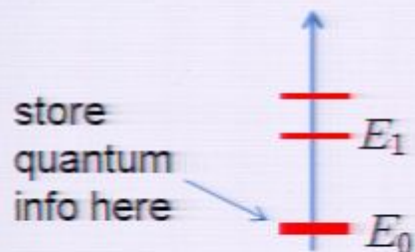
Local perturbations and effective Hamiltonians



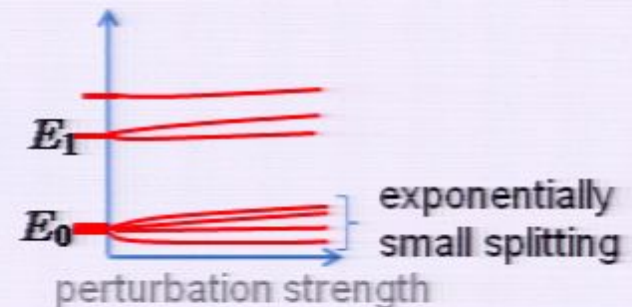
Characteristic properties:

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(Bravyi, Hastings & Michalakis'10 have rigorously shown gap stability for a certain class of Hamiltonians. Perturbation theory suggests the same is true more generally.)



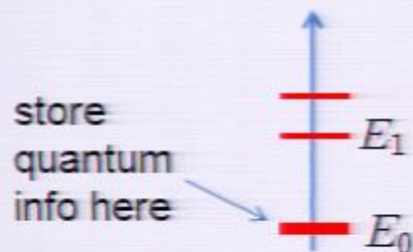
turn on local perturbations



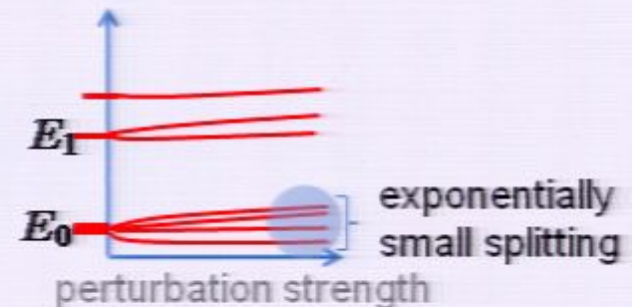
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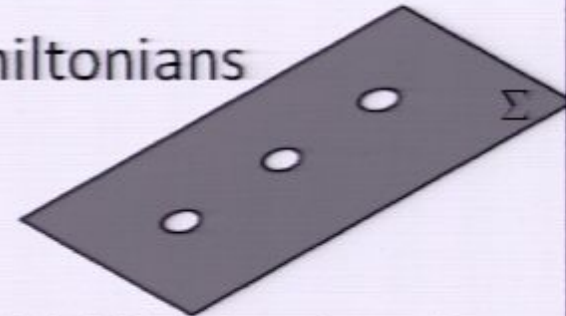
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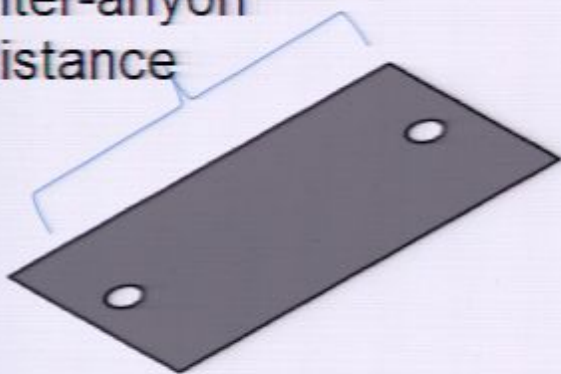
How do low-energy degrees of freedom behave?

What is a good *effective model* (Hamiltonian)?

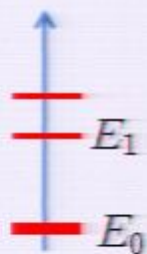
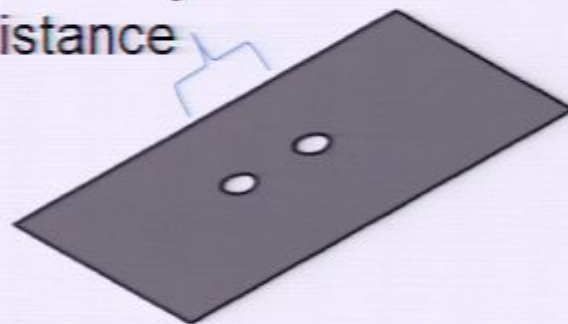


Topological charge tunneling

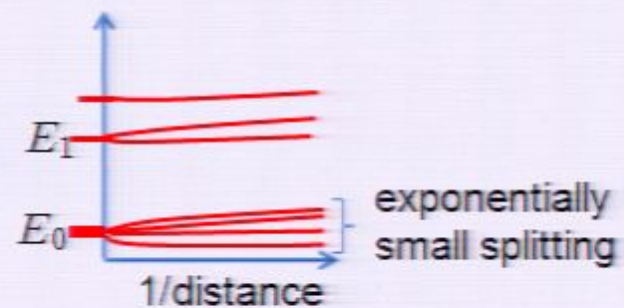
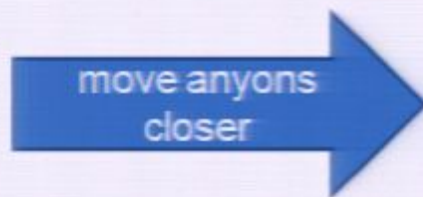
large
inter-anyon
distance



small
inter-anyon
distance

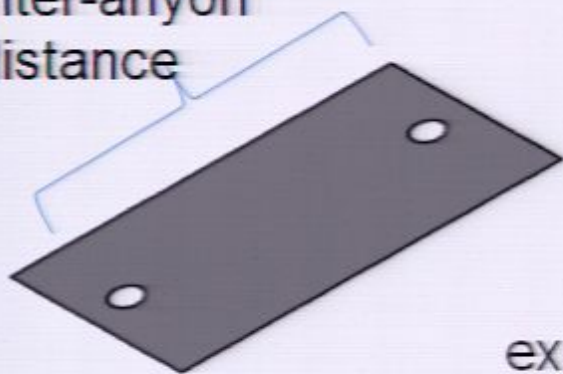


move anyons
closer

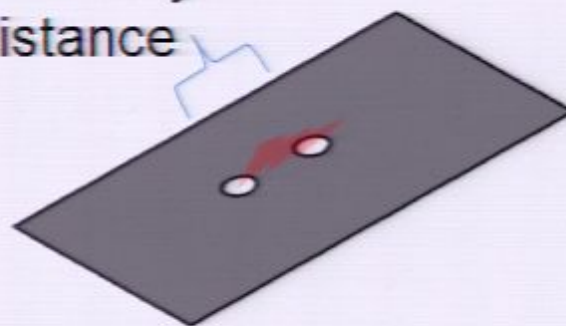


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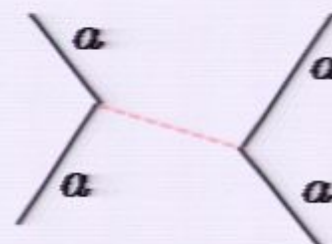
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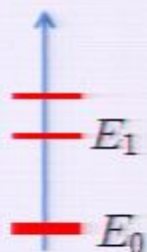
small
inter-anyon
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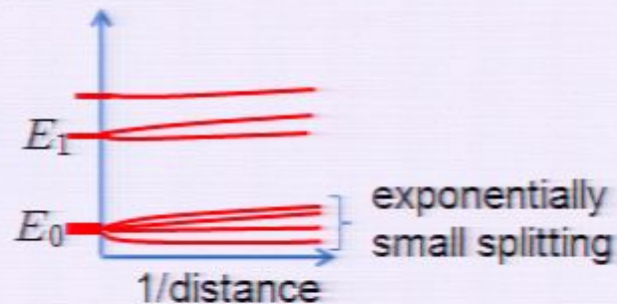
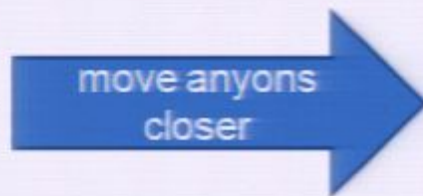
exchange of
topological charge
is responsible for
energy splitting



Bonderson,
Phys.Rev.Lett.103:110403,2009



move anyons
closer



Ground space \mathcal{H}_Σ

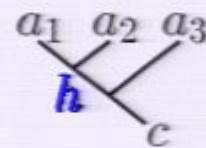
depends on topology of surface Σ

Underlying algebraic data: (UTMF/TQFT)

- set of labels “anyons” $\{a, b, c, \dots\}$
- fusion rules: set of triples of labels
-

Σ

Basis vectors of \mathcal{H}_Σ



(a_1, a_2, h)

(h, a_3, c)

fusion-
consistent

Ground space \mathcal{H}_Σ

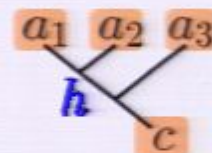
depends on topology of surface Σ
and a labeling of boundary components

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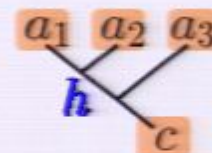
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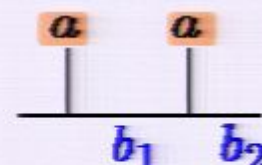
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-

Σ

Basis vectors of \mathcal{H}_Σ



(a_1, a_2, h)
 (h, a_3, c)
fusion-
consistent



(b_2, a, b_1)
 (b_1, a, b_2)
fusion-
consistent

Ground space \mathcal{H}_Σ

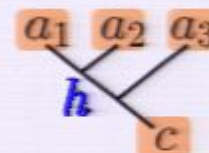
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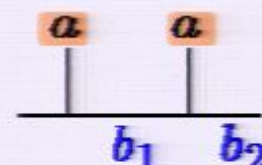
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-

Σ

Basis vectors of \mathcal{H}_Σ



(a_1, a_2, h)
 (h, a_3, c)
fusion-
consistent



(b_2, a, b_1)
 (b_1, a, b_2)
fusion-
consistent



(b_j, a, b_{j+1})
fusion-
consistent

Example state space of N Fibonacci anyons

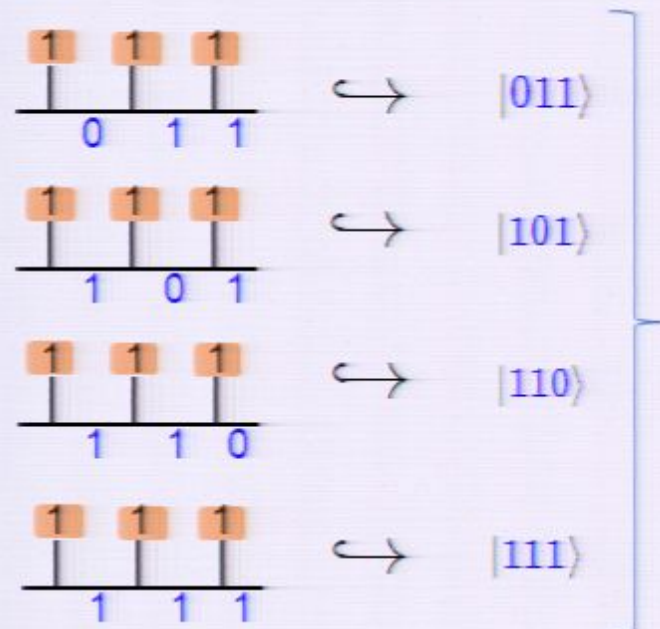
torus Σ_N with N holes



Fibonacci TQFT

- label set $\{0, 1\}$
- fusion rules
 $(0, 0, 0), (1, 1, 0), (1, 1, 1) + \text{permutations}$

basis for N=3 Fibonacci anyons:



Every N-bit string **without consecutive 0's** represents a basis state.

Example state space of N Fibonacci anyons

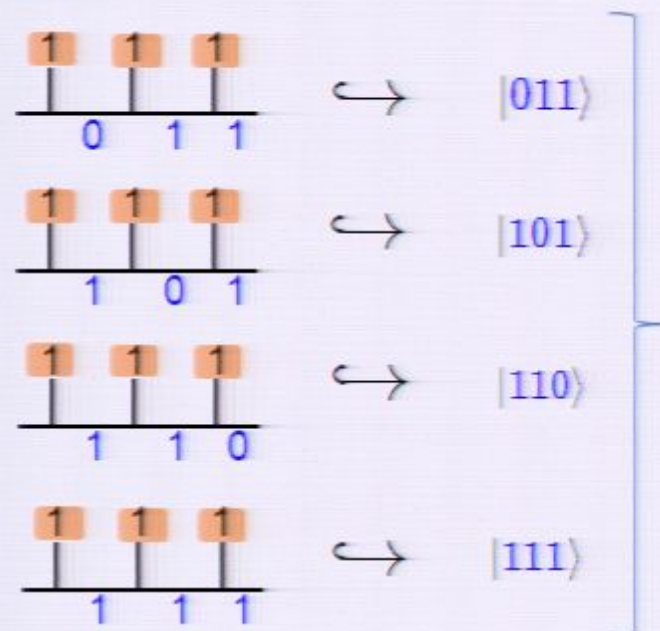
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Fibonacci TQFT

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basis for N=3 Fibonacci anyons:



$$\dim \mathcal{H}_{\Sigma_N} \in O\left(\left(\frac{\sqrt{5}+1}{2}\right)^N\right) \ll 2^N$$

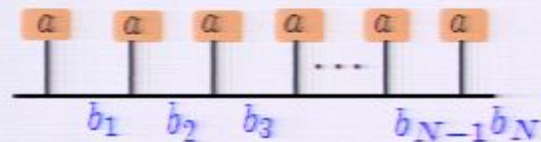
≈ 1.62

Treating Fibonacci anyons in terms of qubits is computationally disadvantageous!

Every N-bit string *without consecutive 0's* represents a basis state.

Example: state space of N anyons of type a

torus Σ_N with N holes each with label a



$|b_1, b_2, \dots, b_N\rangle$

(b_j, a, b_{j+1}) fusion-consistent

\mathcal{H}_{Σ_N}



$(\mathbb{C}^d)^{\otimes N}$

anyonic states

qudit states $d = \# \text{particle types}$

$$\dim \mathcal{H}_{\Sigma_N} \in O(d_a^N) \ll d^N$$

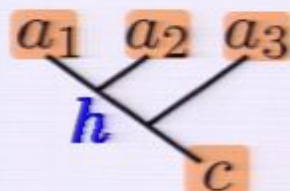
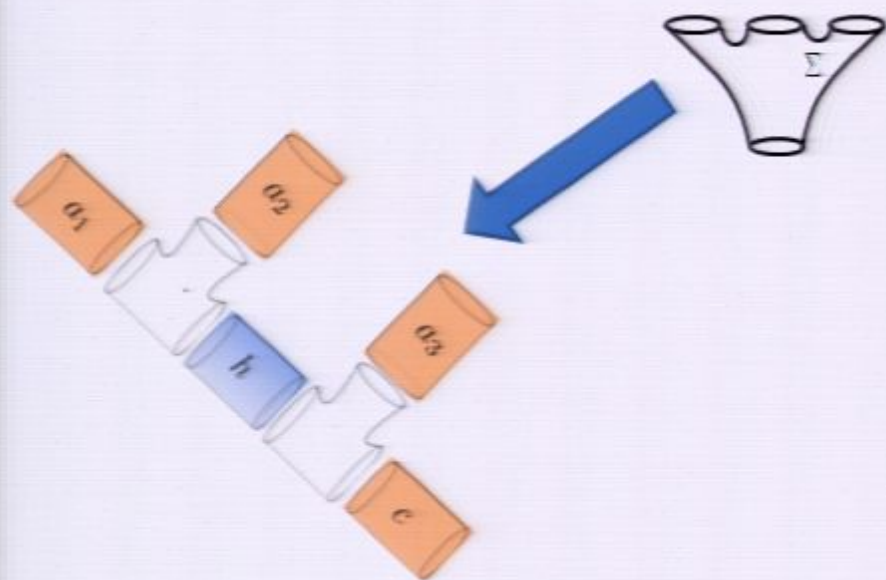
“quantum dimension” of particle a
(not necessarily integer!)

Treating anyons in terms of qudits is
computationally disadvantageous!

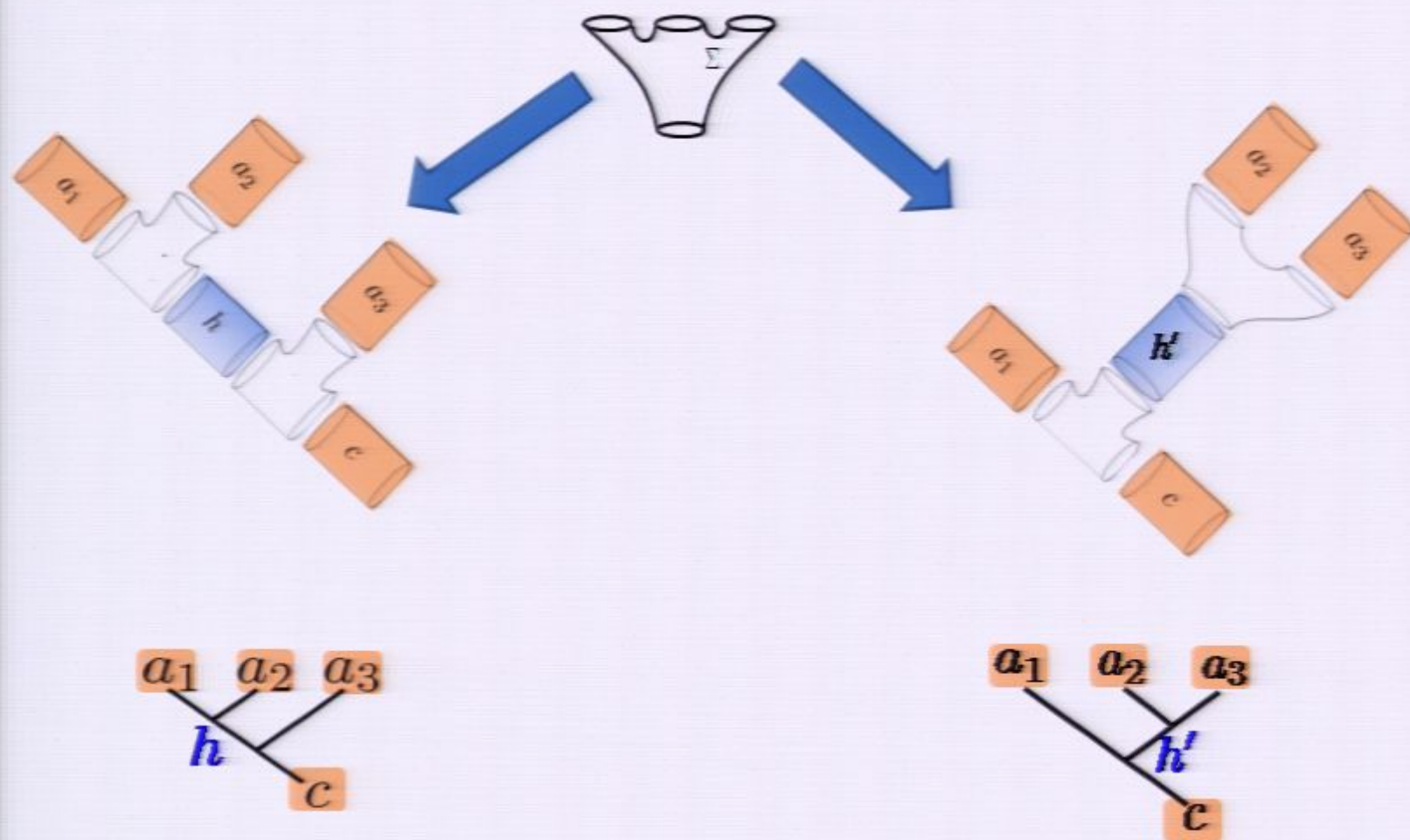
“Locally related” pants decompositions



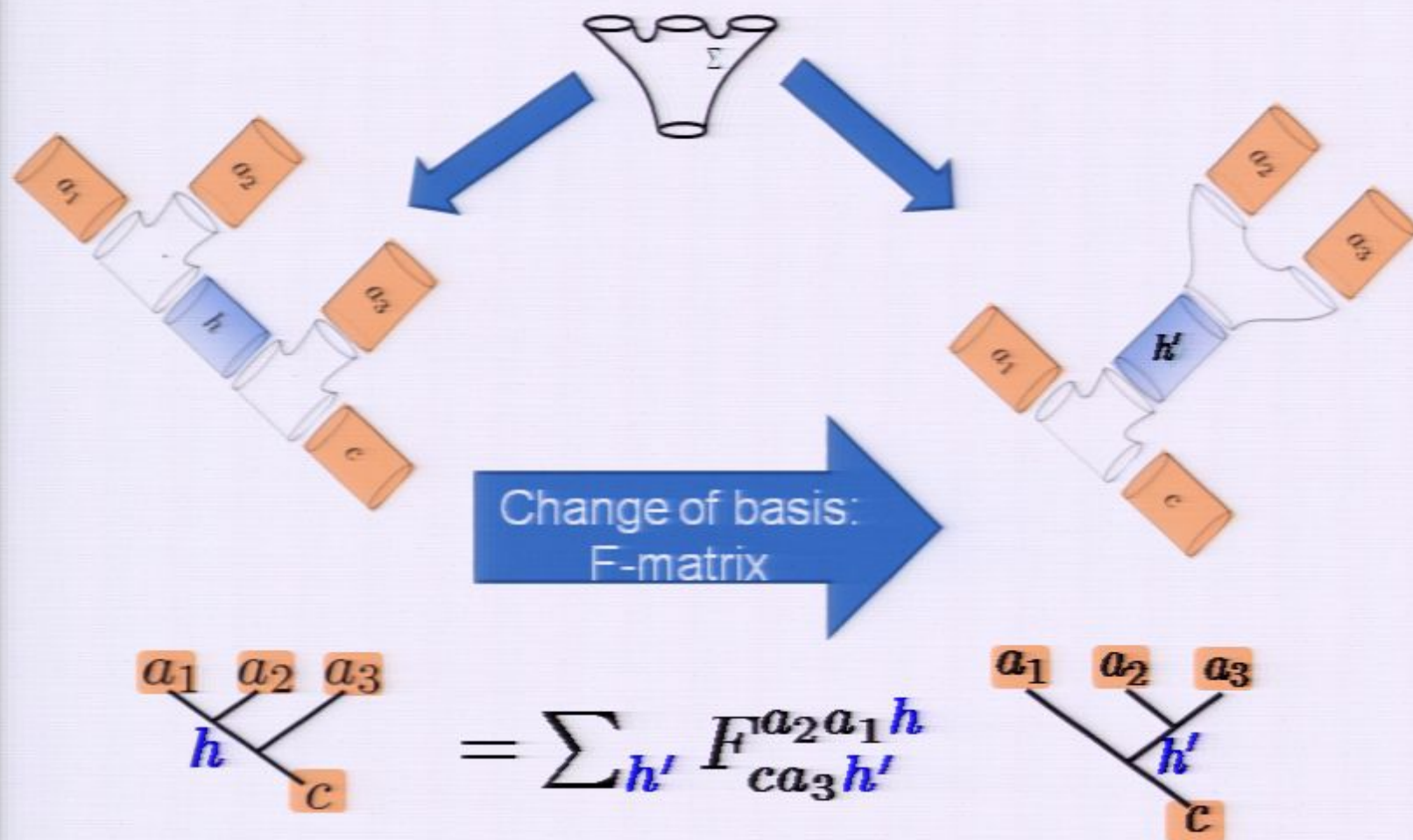
“Locally related” pants decompositions



“Locally related” pants decompositions



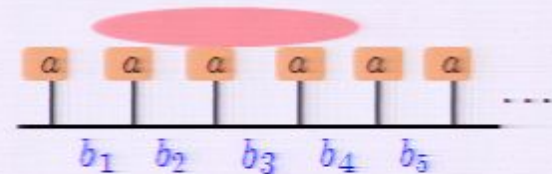
“Locally related” pants decompositions



Topological charge conservation by local operators



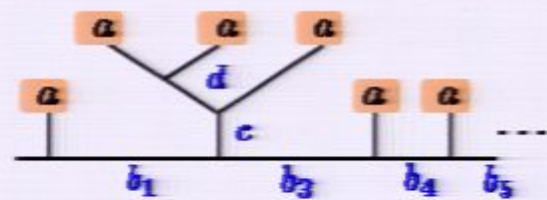
A local operator O acting on a region \mathcal{A} (e.g., with 3-anyons) should only affect topological variables “inside” \mathcal{A} .



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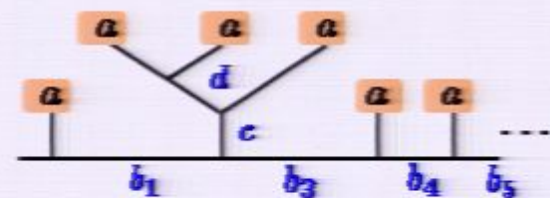


Consider a pants decomposition where the boundary $\partial\mathcal{A}$ is the boundary of a pant segment.

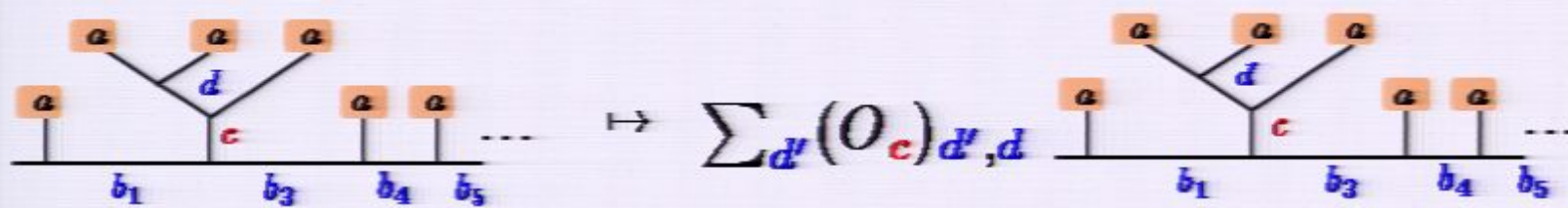
Topological charge conservation by local operators



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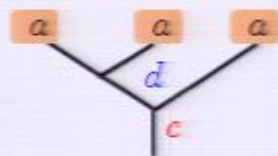


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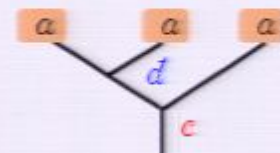


The operator preserves only affects/depends on the subtree, and preserves its total charge c (\rightarrow block-diagonal).

Ribbon graph representation of local operators



$$\mapsto \sum_{d'} (O_c)_{d', d}$$



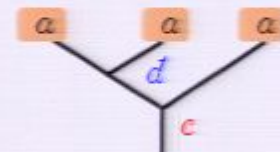
The operator preserves only affects/depends on the subtree, and preserves its total charge c (\rightarrow block-diagonal).

Ribbon graph representation of local operators

$$O = \bigoplus_{\mathbf{c}} \left(\sum_{\mathbf{d}, \mathbf{d}'} (O_{\mathbf{c}})_{\mathbf{d}', \mathbf{d}} \left| \begin{array}{c} \text{a} \quad \text{a} \quad \text{a} \\ \diagdown \quad \diagup \\ \text{d} \\ | \\ \text{c} \end{array} \right\rangle \left\langle \begin{array}{c} \text{a} \quad \text{a} \quad \text{a} \\ \diagdown \quad \diagup \\ \text{d} \\ | \\ \text{c} \end{array} \right| \right)$$



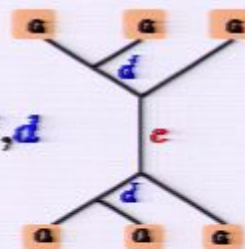
$$\mapsto \sum_{\mathbf{d}'} (O_{\mathbf{c}})_{\mathbf{d}', \mathbf{d}}$$



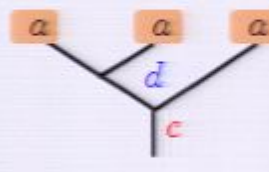
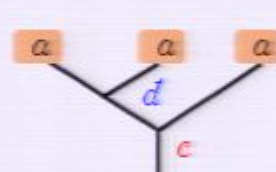
The operator preserves only affects/depends on the **subtree**, and preserves its total charge **c** (-> block-diagonal).

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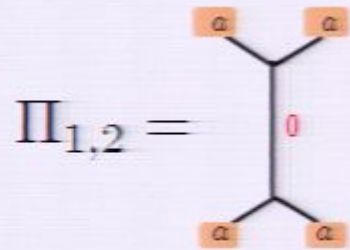
$$= \sum_{\mathbf{c}, \mathbf{d}, \mathbf{d}'} (O_{\mathbf{c}})_{\mathbf{d}', \mathbf{d}}$$


This graphical representation of operators incorporates charge conservation and obeys simple rules for taking adjoints, multiplication, ...

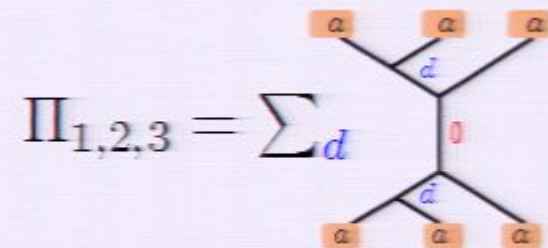

 $\mapsto \sum_{\mathbf{d}'} (O_{\mathbf{c}})_{\mathbf{d}', \mathbf{d}}$


The operator preserves only affects/depends on the **subtree**, and preserves its total charge **c** (-> block-diagonal).

Projection onto trivial charge and the golden chain

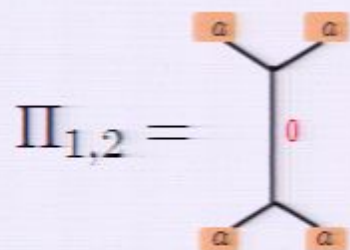


pair-interaction

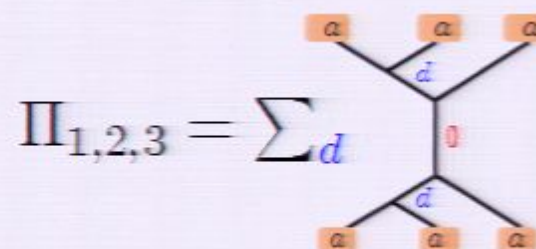


3-anyon interaction

Projection onto trivial charge and the golden chain



pair-interaction



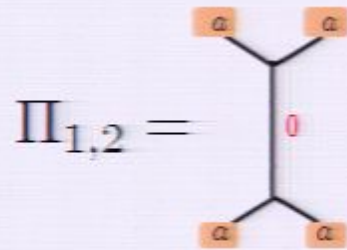
3-anyon interaction

translation-invariant effective “anyonic chain” model

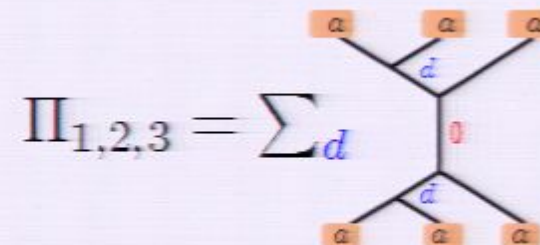


$$H = -J_1 \sum_i \Pi_{i,i+1} - J_2 \sum_i \Pi_{i,i+1,i+2}$$

Projection onto trivial charge and the golden chain



pair-interaction



3-anyon interaction

translation-invariant effective “anyonic chain” model



$$H = -J_1 \sum_i \Pi_{i,i+1} - J_2 \sum_i \Pi_{i,i+1,i+2}$$

For Fibonacci-anyons:

- exact solution known for $J_2 = 0$

Feiguin et al., Phys. Rev. Lett. 98, 160409 (2007)

- phase diagram (numerical)

Trebst et al., Phys. Rev. Lett. 101, 050401

Rest of this talk: (1) variational ansatz
(2) numerical test of validity

various generalizations:

- $\text{su}(2)_k$, 2D arrangements of anyons

Gils et al., Phys. Rev. Lett. 103, 070401

Ludwig et al., arxiv:1003.3453

- random couplings

Fidkowski et al., Phys. Rev. B 78, 22204

Fidkowski et al., Phys. Rev. B 79, 155120

Anyonic formalism: Computational rules

$$|\Psi\rangle = \sum_{b_1, \dots, b_N} \Psi_{b_1, \dots, b_N} \begin{array}{c} \text{a} \quad \text{a} \quad \text{a} \quad \text{a} \quad \text{a} \\ | \quad | \quad | \quad | \quad | \\ b_1 \quad b_2 \quad \dots \quad b_{N-1} \quad b_N \end{array}$$

$$O = \sum_{c, d, d'} (O_c)_{d', d}$$



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• **adjoints:** invert & conjugate

$$\langle \Psi | = \sum_{b_1, \dots, b_N} \Psi_{b_1, \dots, b_N}^* \begin{array}{c} b_1 \quad b_2 \quad \dots \quad b_{N-1} \quad b_N \\ | \quad | \quad | \quad \dots \quad | \\ \text{a} \quad \text{a} \quad \text{a} \quad \dots \quad \text{a} \end{array}$$

$$O^\dagger = \sum_{c, d, d'} (O_c)^*_{d', d} \begin{array}{c} \text{a} \quad \text{a} \quad \text{a} \\ \diagdown \quad \diagup \quad \diagdown \\ \text{d} \quad \text{d}' \quad \text{d} \\ | \\ \text{c} \\ \diagup \quad \diagdown \quad \diagup \\ \text{a} \quad \text{a} \quad \text{a} \end{array}$$

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• **multiplication:** stack

$$O|\Psi\rangle = \sum_{\substack{b_1, \dots, b_N \\ c, d, d'}} (O_c)_{d', d} \Psi_{b_1, \dots, b_N} \begin{array}{c} \text{a} \quad \text{a} \quad \text{a} \\ \diagdown \quad \diagup \quad | \\ \quad \quad \quad \text{d} \\ \diagup \quad \diagdown \quad | \\ \text{a} \quad \text{a} \quad \text{a} \\ | \quad | \quad | \quad \dots \quad | \\ b_1 \quad b_2 \quad \dots \quad b_{N-1} \quad b_N \end{array}$$

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$$O = \sum_{c, d, d'} (O_c)_{d', d} \begin{array}{c} \text{a} \quad \text{a} \quad \text{a} \\ \diagdown \quad \diagup \quad | \\ \text{d} \quad \text{d}' \\ \diagup \quad \diagdown \quad | \\ \text{c} \quad \text{c} \quad \text{c} \end{array}$$

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& simplify

$$= \sum_{b'_1, \dots, b'_N} \Psi_{b'_1, \dots, b'_N} \begin{array}{c} \text{a} \quad \text{a} \quad \text{a} \quad \text{a} \quad \text{a} \\ | \quad | \quad | \quad | \quad | \\ b'_1 \quad b'_2 \quad \dots \quad b'_{N-1} \quad b'_N \end{array}$$

some sequence of subgraph deformations/substitutions

Simplification rules for anyonic diagrams

$$\langle i = i \rangle$$

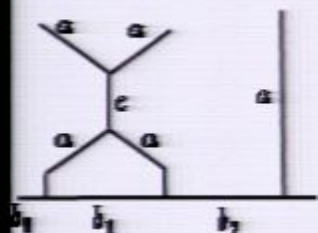
$$\begin{array}{c} i \\ | \\ | \\ j \end{array} = \begin{array}{c} i \quad j \\ \diagdown \quad \diagup \\ 0 \\ \diagup \quad \diagdown \\ i \quad j \end{array}$$

$$\bigcirc i = d_i$$

$$\begin{array}{c} i \\ \text{---} \bigcirc j \end{array} = 0 \quad i \neq j$$

$$\begin{array}{c} i \quad t \\ \diagdown \quad \diagup \\ m \\ \diagup \quad \diagdown \\ j \quad k \end{array} = \sum_n F_{ktn}^{ijm} \begin{array}{c} i \quad t \\ \diagdown \quad \diagup \\ n \\ \diagup \quad \diagdown \\ j \quad k \end{array}$$

Simplification rules for anyonic diagrams



$$\langle i = i \rangle$$

$$i \parallel j = \begin{array}{c} i \quad j \\ \diagdown \quad \diagup \\ \quad 0 \\ \diagup \quad \diagdown \\ i \quad j \end{array}$$

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Simplification rules for anyonic diagrams

$$\begin{array}{c} \text{a} \quad \text{a} \\ \diagdown \quad \diagup \\ \text{c} \\ \diagup \quad \diagdown \\ \text{a} \quad \text{a} \\ \diagup \quad \diagdown \\ b_1 \quad b_1 \end{array} \quad \begin{array}{c} | \\ \text{a} \\ | \end{array} = \sum_n F_{b_2 a n}^{a b_0 b_1} \begin{array}{c} \text{a} \quad \text{a} \\ \diagdown \quad \diagup \\ \text{c} \\ \diagup \quad \diagdown \\ \text{a} \quad \text{a} \\ \diagup \quad \diagdown \\ b_1 \quad b_2 \end{array} \quad \begin{array}{c} | \\ \text{a} \\ | \end{array}$$

$$\langle i = i \rangle$$

$$\begin{array}{c} i \\ | \end{array} \quad \begin{array}{c} | \\ j \end{array} = \begin{array}{c} i \quad j \\ \diagdown \quad \diagup \\ 0 \\ \diagup \quad \diagdown \\ i \quad j \end{array}$$

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Simplification rules for anyonic diagrams

$$\begin{array}{c} a \\ \diagup \quad \diagdown \\ c \\ \diagdown \quad \diagup \\ a \\ \hline b_1 \quad b_1 \quad b_2 \end{array} \quad \begin{array}{c} a \\ | \\ a \end{array} = \sum_n F_{b_2 a n}^{a b_0 b_1} \begin{array}{c} a \\ \diagup \quad \diagdown \\ e \\ \bigcirc \\ n \\ \diagdown \quad \diagup \\ b_1 \quad b_2 \end{array} \quad \begin{array}{c} a \\ | \\ a \end{array}$$

$$= \sum_{n,l} F_{b_2 a n}^{a b_0 b_1} F_{n a l}^{a c a} \begin{array}{c} a \\ \diagup \quad \diagdown \\ e \\ \begin{array}{c} l \\ \bigcirc \end{array} \\ n \\ \diagdown \quad \diagup \\ b_1 \quad b_2 \end{array} \quad \begin{array}{c} a \\ | \\ a \end{array}$$

$$\langle i = i \rangle$$

$$\begin{array}{c} i \\ | \end{array} \quad \begin{array}{c} j \\ | \end{array} = \begin{array}{c} i \quad j \\ \diagdown \quad \diagup \\ 0 \\ \diagup \quad \diagdown \\ i \quad j \end{array}$$

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Simplification rules for anyonic diagrams

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$$= \sum_{n,l} F_{b_2 a n}^{a b_0 b_1} F_{n a l}^{a c a} \begin{array}{c} a \\ \diagup \quad \diagdown \\ e \\ \diagdown \quad \diagup \\ l \quad \bigcirc \quad a \\ n \\ \diagdown \quad \diagup \\ b_1 \quad b_2 \end{array} \quad \begin{array}{c} a \\ | \end{array}$$

$$= F_{b_2 a e}^{a b_0 b_1} F_{c a 0}^{a c a} \begin{array}{c} a \\ \diagup \quad \diagdown \\ c \\ \diagdown \quad \diagup \\ \bigcirc \quad a \\ b_1 \quad b_2 \end{array} \quad \begin{array}{c} a \\ | \end{array}$$

$$= F_{b_2 a e}^{a b_0 b_1} F_{c a 0}^{a c a} d_a \begin{array}{c} a \\ \diagup \quad \diagdown \\ c \\ \diagdown \quad \diagup \\ \bigcirc \quad a \\ b_1 \quad b_2 \end{array} \quad \begin{array}{c} a \\ | \end{array} = \sum_m F_{b_2 a e}^{a b_0 b_1} F_{c a 0}^{a c a} d_a F_{b_0 b_2 m}^{a a c} \begin{array}{c} a \\ | \end{array} \quad \begin{array}{c} a \\ | \end{array} \quad \begin{array}{c} a \\ | \end{array}$$

$b_1 \quad m \quad b_2$

$$\langle i \rangle = i$$

$$i \quad | \quad j = \begin{array}{c} i \quad j \\ \diagdown \quad \diagup \\ 0 \\ \diagup \quad \diagdown \\ i \quad j \end{array}$$

$$\bigcirc i = d_i$$

$$i \quad \bigcirc \quad j = 0 \quad i \neq 0$$

$$\begin{array}{c} i \quad l \\ \diagdown \quad \diagup \\ m \quad k \\ \diagup \quad \diagdown \\ j \end{array} = \sum_n F_{k l n}^{i j m} \begin{array}{c} i \quad l \\ \diagdown \quad \diagup \\ n \\ \diagup \quad \diagdown \\ j \quad k \end{array}$$

Simplification rules for anyonic diagrams

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$$= \sum_{n,l} F_{b_2 a n}^{a b_0 b_1} F_{n a l}^{a c a} \begin{array}{c} a \\ \diagup \quad \diagdown \\ e \\ \diagup \quad \diagdown \\ l \quad \bigcirc \quad a \\ \diagdown \quad \diagup \\ b_1 \quad b_1 \end{array} \Big| \begin{array}{c} a \\ | \\ a \end{array}$$

$$= F_{b_2 a e}^{a b_0 b_1} F_{c a 0}^{a c a} \begin{array}{c} a \\ \diagup \quad \diagdown \\ c \\ \diagdown \quad \diagup \\ b_1 \quad b_1 \end{array} \Big| \begin{array}{c} a \\ \diagup \quad \diagdown \\ \bigcirc \\ a \end{array}$$

$$= F_{b_2 a e}^{a b_0 b_1} F_{c a 0}^{a c a} d_a \begin{array}{c} a \\ \diagup \quad \diagdown \\ c \\ \diagdown \quad \diagup \\ b_1 \quad b_1 \end{array} \Big| \begin{array}{c} a \\ | \\ a \end{array} = \sum_m F_{b_2 a e}^{a b_0 b_1} F_{c a 0}^{a c a} d_a F_{b_0 b_2 m}^{a a c} \begin{array}{c} | \\ | \\ | \\ b_1 \quad m \quad b_2 \end{array}$$

$$\left(\begin{array}{c} i \\ | \end{array} \right) = \left(\begin{array}{c} i \\ | \end{array} \right)$$

$$\begin{array}{c} i \\ | \end{array} \begin{array}{c} j \\ | \end{array} = \begin{array}{c} i \quad 0 \quad j \\ \diagdown \quad \diagup \\ i \quad j \end{array}$$

$$\bigcirc i = d_i$$

$$\begin{array}{c} i \\ | \end{array} \bigcirc j = 0 \quad i \neq 0$$

$$\begin{array}{c} i \quad l \\ \diagdown \quad \diagup \\ j \quad k \end{array} = \sum_n F_{k l n}^{i j m} \begin{array}{c} i \quad l \\ \diagdown \quad \diagup \\ n \\ \diagdown \quad \diagup \\ j \quad k \end{array}$$



Anyonic formalism *versus* tensor networks

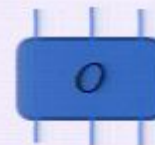
state

$$|\Psi\rangle = \sum_{\vec{b}} \Psi_{\vec{b}} \quad \begin{array}{c} \text{a} \quad \text{a} \quad \text{a} \quad \text{a} \quad \text{a} \\ | \quad | \quad | \quad | \quad | \\ \text{b}_1 \quad \text{b}_2 \quad \dots \quad \text{b}_{N-1} \quad \text{b}_N \end{array}$$



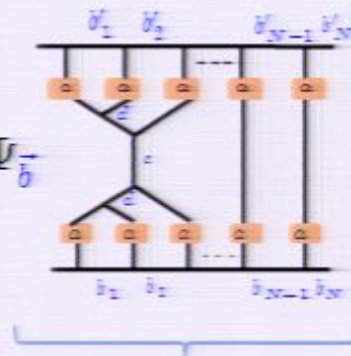
local operator

$$O = \sum_{c,d,d'} (O_c)_{d',d} \quad \begin{array}{c} \text{a} \quad \text{a} \quad \text{a} \\ \diagdown \quad \diagup \quad \diagdown \\ \quad \quad c \quad \quad \diagup \\ \diagup \quad \diagdown \quad \diagdown \\ \text{a} \quad \text{a} \quad \text{a} \end{array}$$

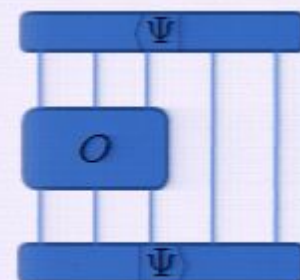


expectation value

$$\langle \Psi | O | \Psi \rangle = \sum_{\substack{\vec{b}, \vec{b}' \\ c, d, d'}} \Psi_{\vec{b}'}^* \cdot (O_c)_{d', d} \cdot \Psi_{\vec{b}}$$



evaluated by local rules



evaluated by tensor contraction

Anyonic entanglement renormalization:
a variational ansatz which


- exploits structure of Hilbert space
- allows straightforward generalization to 2D
- has operational meaning

Anyonic entanglement renormalization: Definition

1. Choose a suitable tree-like graph



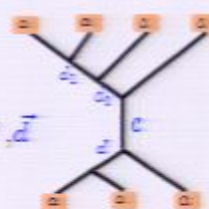
2. Associate (variational) parameters



$$|\varphi\rangle = \sum_{\vec{b}} \varphi_{\vec{b}} \begin{array}{c} \text{a} \quad \text{a} \quad \text{a} \\ | \quad | \quad | \\ \vec{b}_1 \quad \vec{b}_2 \quad \vec{b}_3 \end{array}$$

(3-)anyon state



$$V = \sum_{c, \vec{d}, \vec{d}'} (V_c)_{\vec{d}, \vec{d}'}$$


topological charge-preserving isometry

3. Substitute boxes by graphical representation

Result: formal linear combination,
represents variational state $|\Psi\rangle$

Anyonic entanglement renormalization: Definition

1. Choose a suitable tree-like graph



2. Associate (variational) parameters



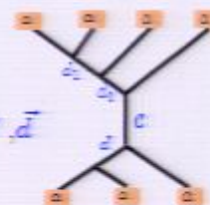
$$|\varphi\rangle = \sum_{\vec{b}} \varphi_{\vec{b}}$$



(3-)anyon state



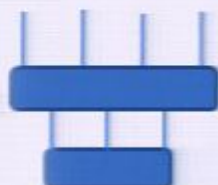
$$V = \sum_{c, \vec{d}, \vec{d'}} (V_c)_{\vec{d}, \vec{d'}}$$



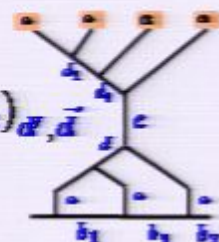
topological charge-preserving isometry

3. Substitute boxes by graphical representation

example:



$$|\Psi\rangle = \sum_{\vec{b}, c, \vec{d}, \vec{d'}} \varphi_{\vec{b}} \cdot (V_c)_{\vec{d}, \vec{d'}}$$



Result: formal linear combination,
represents variational state $|\Psi\rangle$

Anyonic entanglement renormalization: Definition

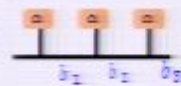
1. Choose a suitable tree-like graph



2. Associate (variational) parameters



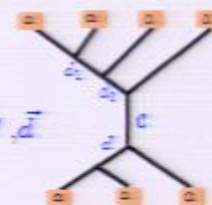
$$|\varphi\rangle = \sum_{\vec{b}} \varphi_{\vec{b}}$$



(3-)anyon state



$$V = \sum_{c, \vec{d}, \vec{d'}} (V_c)_{\vec{d}, \vec{d'}}$$

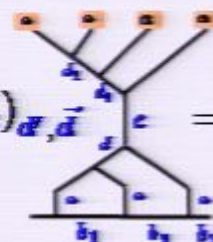


topological charge-preserving isometry

3. Substitute boxes by graphical representation



$$|\Psi\rangle = \sum_{\vec{b}, c, \vec{d}, \vec{d'}} \varphi_{\vec{b}} \cdot (V_c)_{\vec{d}, \vec{d'}}$$



$$= \sum_{\vec{b}} \Psi_{\vec{b}}$$

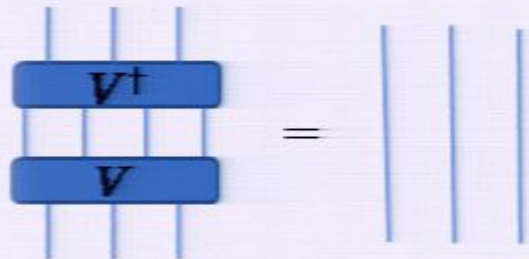


Computing this coefficient is generally hard, but we don't need to!

Result: formal linear combination,
represents variational state $|\Psi\rangle$

Topological charge-preserving isometries

The familiar rule

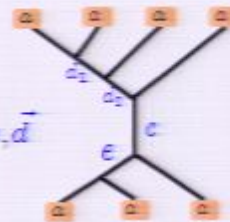


applies in the
anyonic context

if V is a

charge-preserving
isometry:

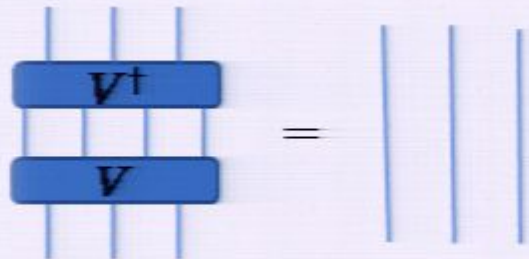
$$V = \bigoplus_c \sum_{\vec{d}, \vec{e}} (V_c)_{\vec{e}, \vec{d}}$$



defined by a family $\{V_c\}_c$
of isometries, i.e., $V_c^\dagger V_c = \text{id}_c$

Topological charge-preserving isometries

The familiar rule



applies in the anyonic context

if V is a

charge-preserving isometry:

$$V = \bigoplus_c \sum_{\vec{d}, e} (V_c)_{e, \vec{d}}$$

defined by a family $\{V_c\}_c$ of isometries, i.e., $V_c^\dagger V_c = \text{id}_c$

This is because such operators satisfy the diagrammatic identity

$$V^\dagger V = \bigoplus_{c, c'} \sum_{\vec{d}, \vec{d}', e, e'} (V_{c'})_{e', \vec{d}'} (V_c)_{e, \vec{d}}$$

Anyonic entanglement renormalization

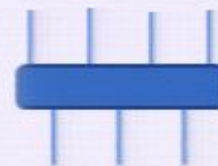
1. Choose a suitable tree-like graph



2. Associate (variational) parameters

$$|\varphi\rangle = \sum_{\vec{b}} \varphi_{\vec{b}} \text{ (diagram with three orange boxes labeled } b_1, b_2, b_3 \text{)}$$

(3-)anyon state



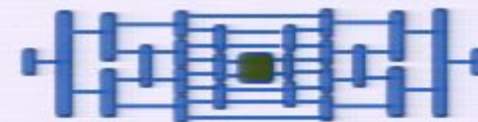
$$V = \sum_{c, \vec{d}, \vec{d'}} (V_c)_{\vec{d}, \vec{d'}} \text{ (diagram with a central node } c \text{ and branches } d_1, d_2, d_3 \text{)}$$

topological charge-preserving isometry

3. Compute expectation values from using substitutions

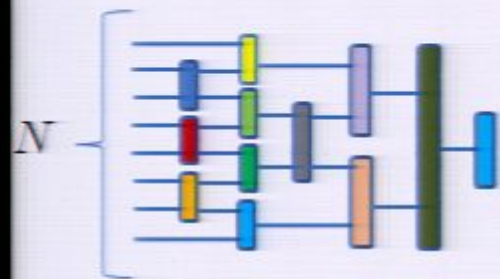
$$\text{(diagram with two blue blocks)} = \text{(simplified diagram)}$$

and evaluating resulting simplified anyon diagram



With these modifications, techniques developed for spin chains (e.g., for numerically varying over isometries, extracting critical exponents/CFT data) can be applied to anyons.

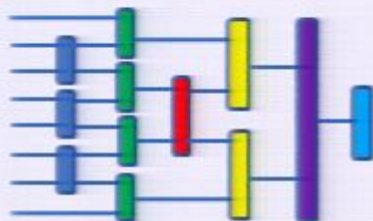
Parameter counting: N anyons of type a



general

$$O(\text{poly}(d_a) N \log N)$$

quantum dimension



translation-invariant

$$O(\text{poly}(d_a) \log N)$$



scale-invariant

$$O(\text{poly}(d_a))$$

Numerical test

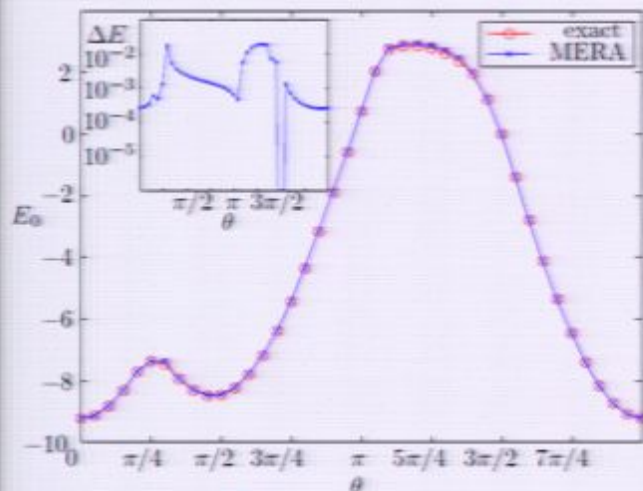
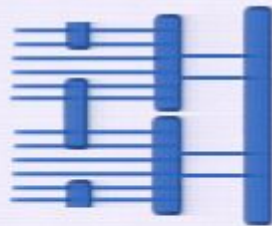
Application to the golden chain



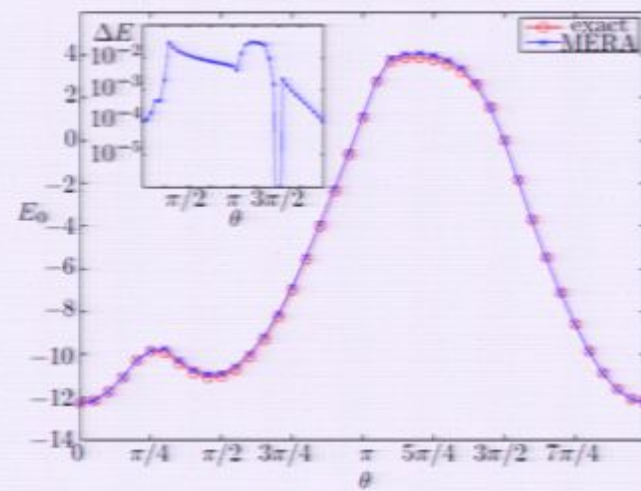
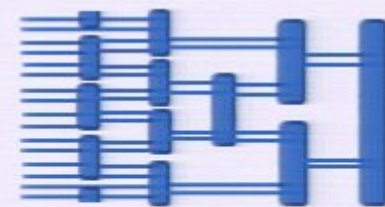
$$H = -J_1 \sum_i \Pi_{i,i+1} - J_2 \sum_i \Pi_{i,i+1,i+2}$$

$$J_1 = \cos \theta, J_2 = \sin \theta$$

Ground state energy computed by anyonic MERA:



12 anyons



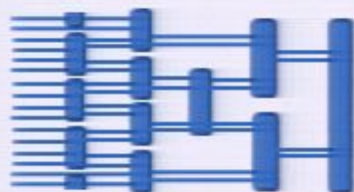
16 anyons

Application to the golden chain

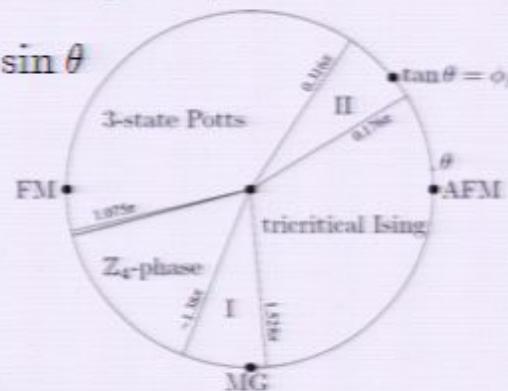


$$H = -J_1 \sum_i \Pi_{i,i+1} - J_2 \sum_i \Pi_{i,i+1,i+2}$$

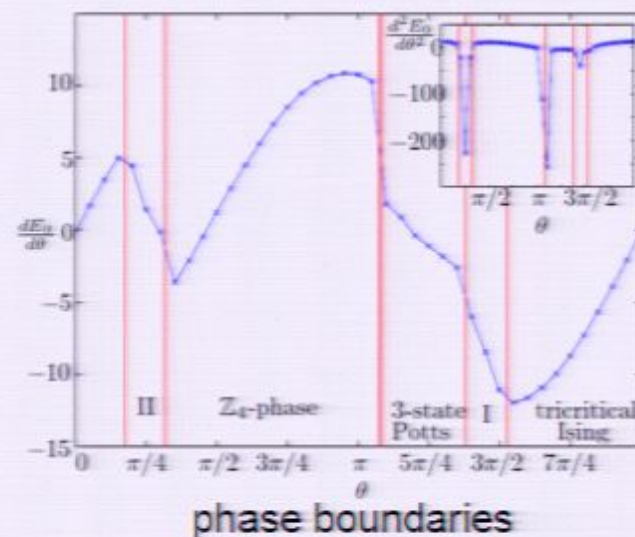
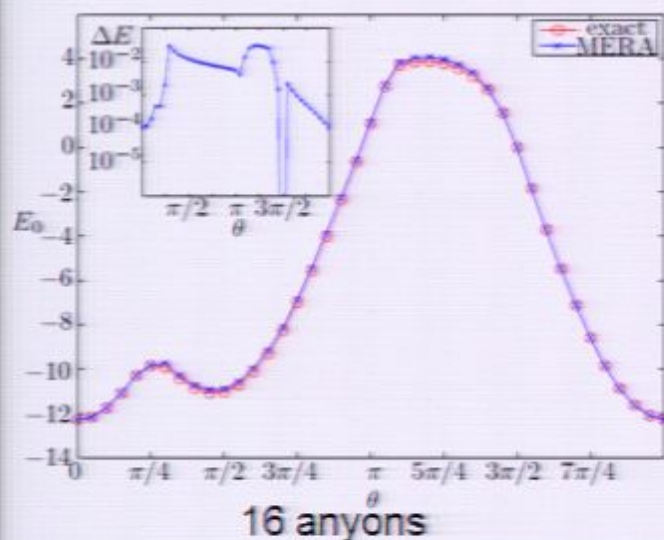
$$J_1 = \cos \theta, J_2 = \sin \theta$$



ground state energy
computed by anyonic MERA



phase diagram obtained
by Trebst et al.,
Phys. Rev. Lett. 101,
050401

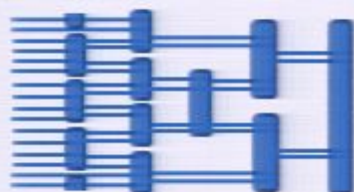


Application to the golden chain

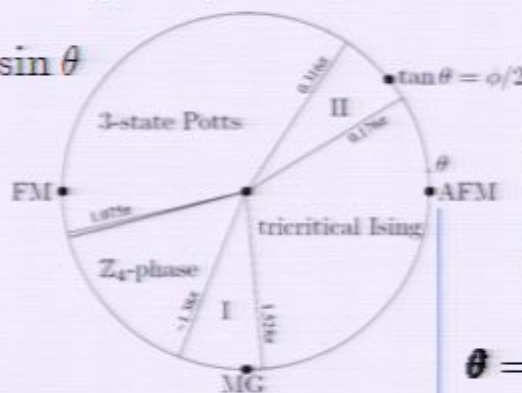


$$H = -J_1 \sum_i \Pi_{i,i+1} - J_2 \sum_i \Pi_{i,i+1,i+2}$$

$$J_1 = \cos \theta, J_2 = \sin \theta$$

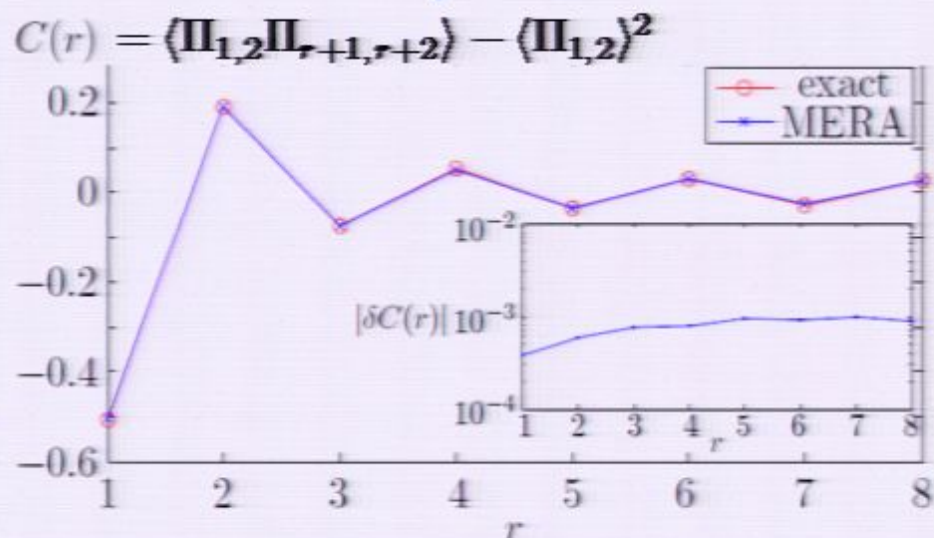
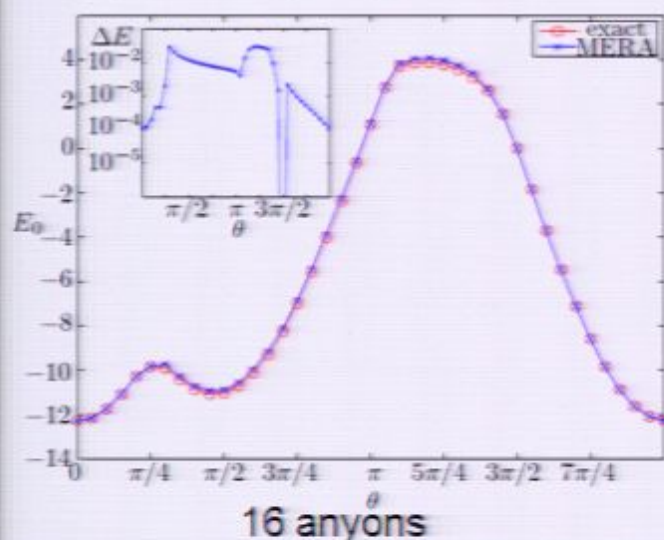


ground state energy
computed by anyonic MERA

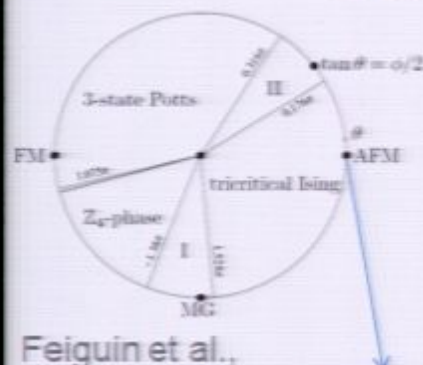


phase diagram obtained
by Trebst et al.,
Phys. Rev. Lett. 101,
050401

$\theta = 0$



Large systems and CFT



Feiguin et al.,
Phys. Rev. Lett. 98,
160409 (2007)

CFT
identified: $c = 7/10$

spectrum of periodic 1D critical system of length N

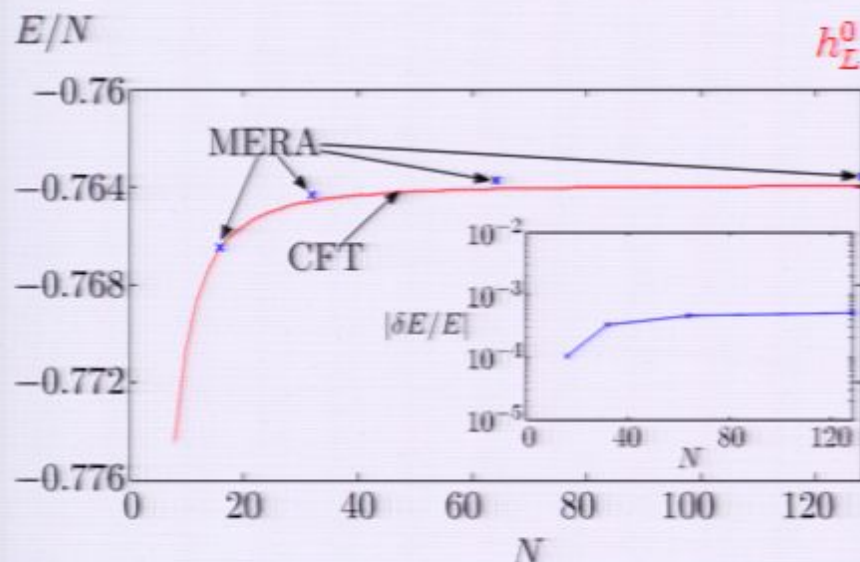
$$E_{h_L, h_R} = \epsilon N + \frac{2\pi v}{N} \left(h_L + h_R - \frac{c}{12} \right)$$

$$h_L = h_L^0 + m_L$$

$$h_R = h_R^0 + m_R$$

$$m_L, m_R \in \mathbb{N} \cup \{0\}$$

$$h_L^0 = h_R^0 \in \{0, 3/80, 1/10, 7/16, 3/5, 3/2\}$$



use these to estimate non-universal constants ϵ, v

agrees with CFT in thermodynamic limit

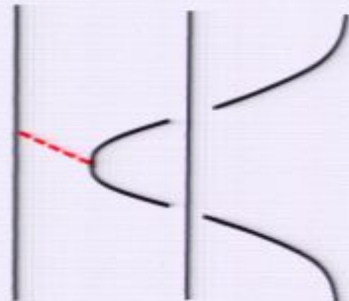
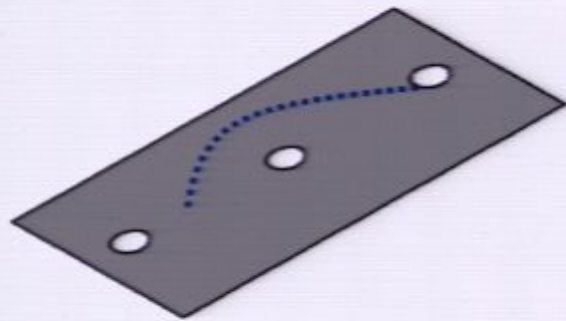
(see also recent related
work by Pfeifer et al.,
arXiv:1006.3532)

2D arrangements of anyons

- Hamiltonians
- evaluation of expectation values by braiding

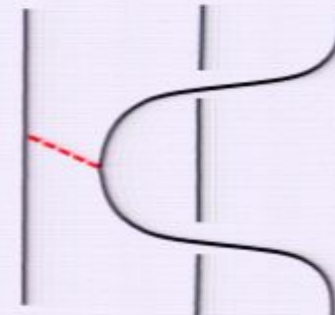
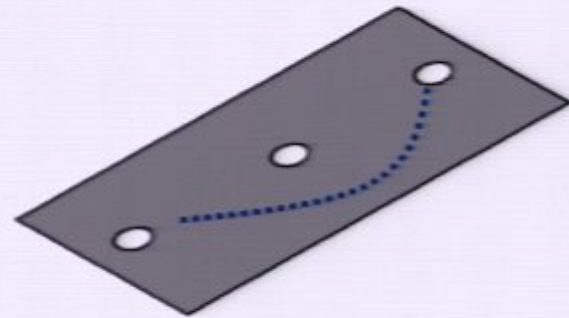
Non-abelian exchange statistics

Next-to-nearest neighbor interactions:



$$B^{-1}HB$$

$$B = \begin{array}{c} \diagdown \quad \diagup \\ \diagup \quad \diagdown \end{array}$$

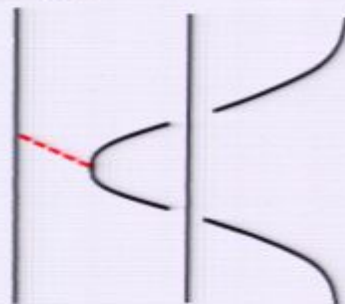
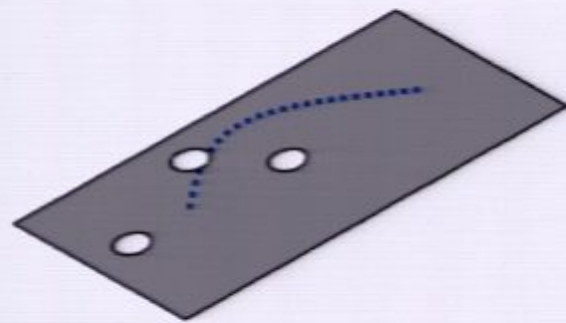


$$\tilde{B}^{-1}H\tilde{B}$$

$$\tilde{B} = \begin{array}{c} \diagup \quad \diagdown \\ \diagdown \quad \diagup \end{array}$$

Non-abelian exchange statistics

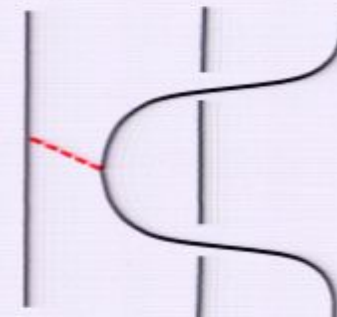
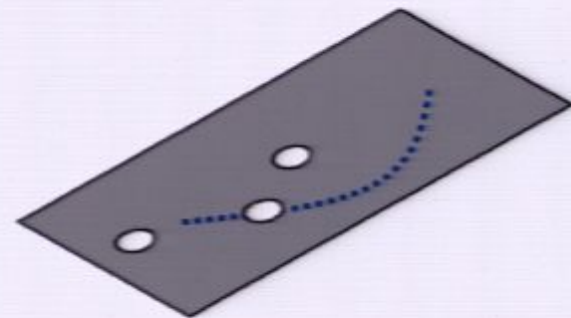
Next-to-nearest neighbor interactions:



$$B^{-1}HB$$

additional
rule:

$$B = \begin{array}{c} \diagup \quad \diagdown \\ \diagdown \quad \diagup \end{array}$$



$$\tilde{B}^{-1}H\tilde{B}$$

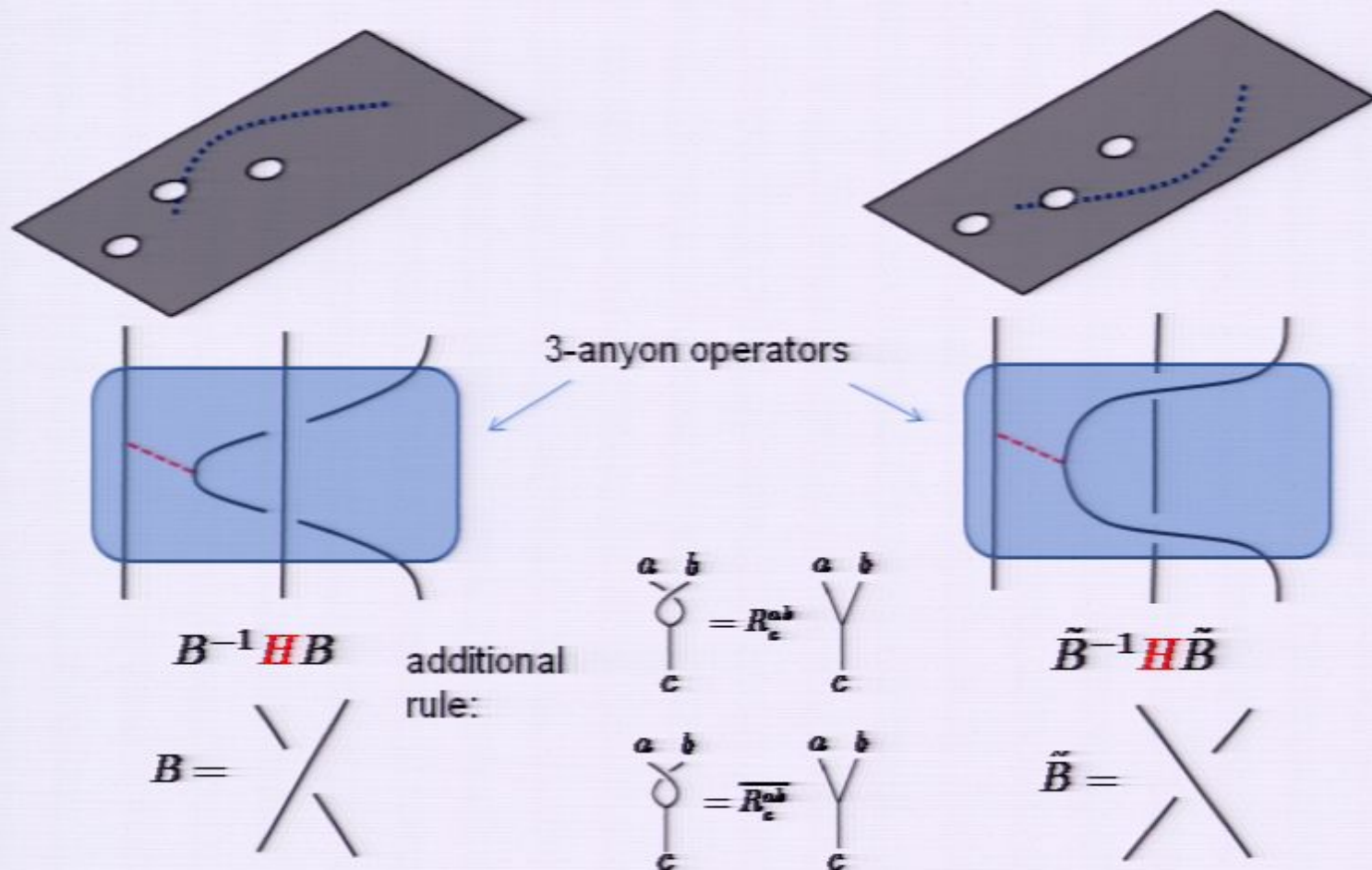
$$\tilde{B} = \begin{array}{c} \diagdown \quad \diagup \\ \diagup \quad \diagdown \end{array}$$

$$\begin{array}{c} a \quad b \\ \diagdown \quad \diagup \\ c \end{array} = R_c^{ab} \begin{array}{c} a \quad b \\ \diagup \quad \diagdown \\ c \end{array}$$

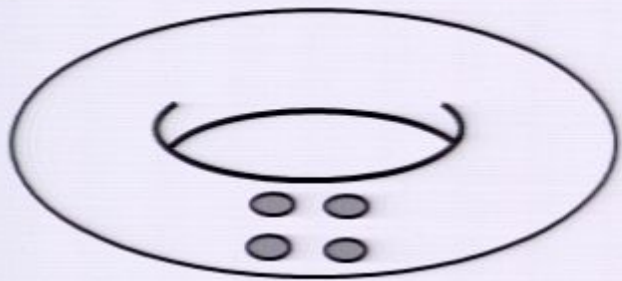
$$\begin{array}{c} a \quad b \\ \diagup \quad \diagdown \\ c \end{array} = \overline{R}_c^{ab} \begin{array}{c} a \quad b \\ \diagdown \quad \diagup \\ c \end{array}$$

Non-abelian exchange statistics

Next-to-nearest neighbor interactions:

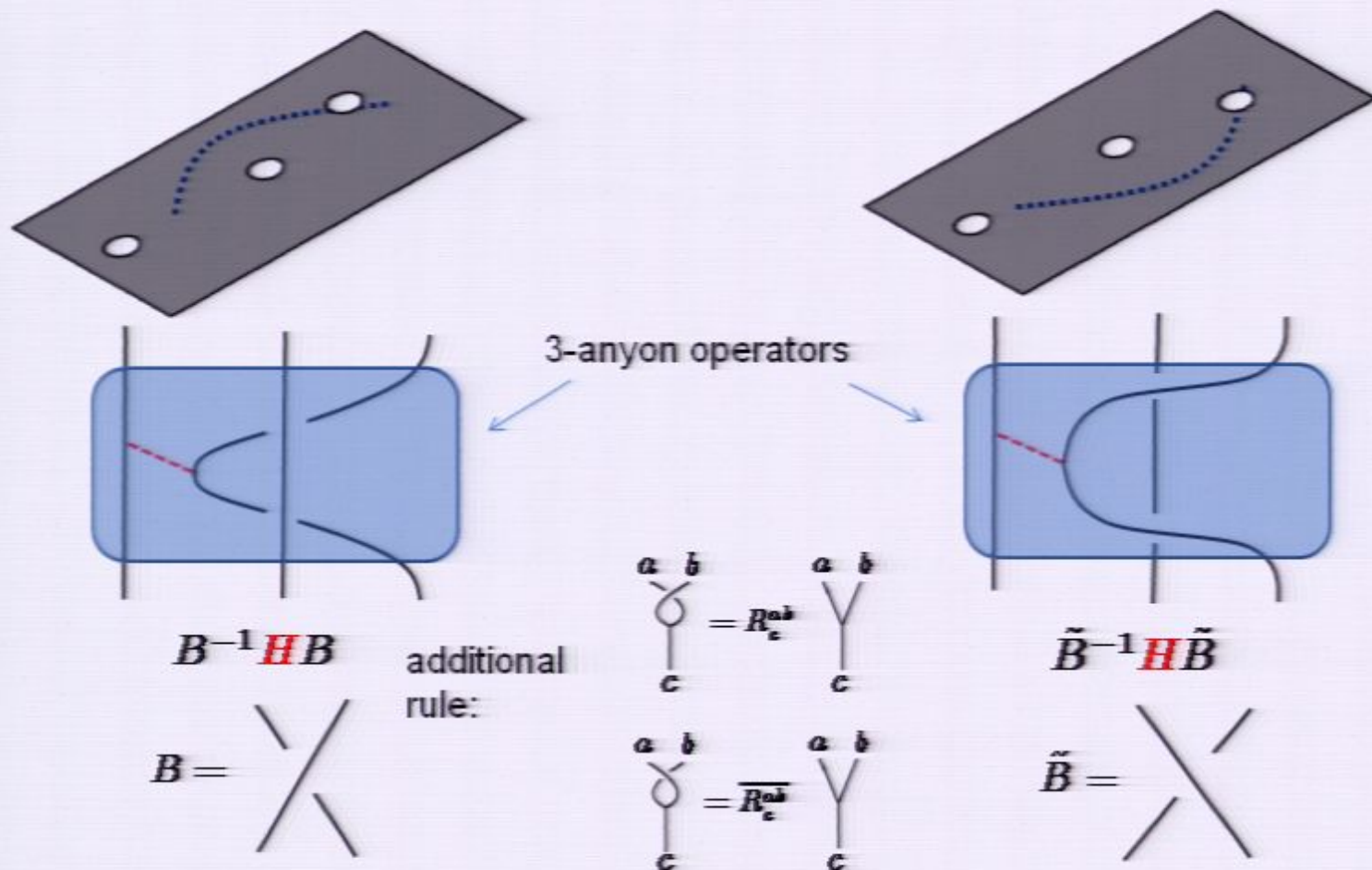


Example: anyonic array

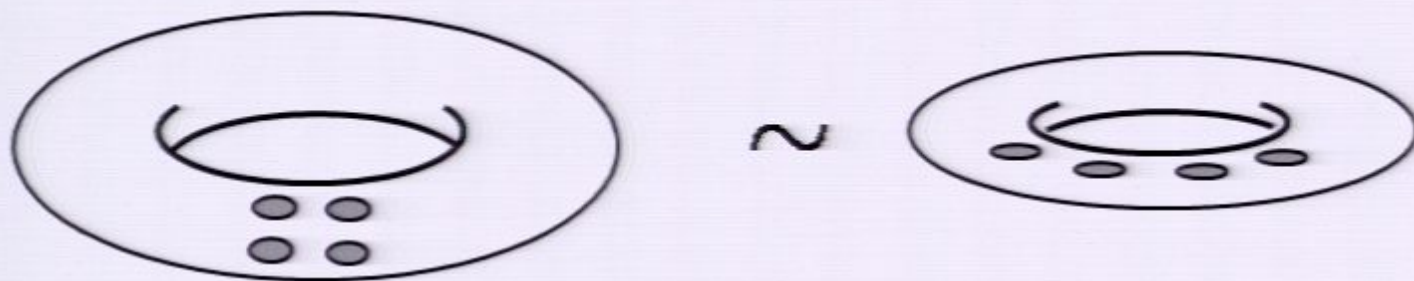


Non-abelian exchange statistics

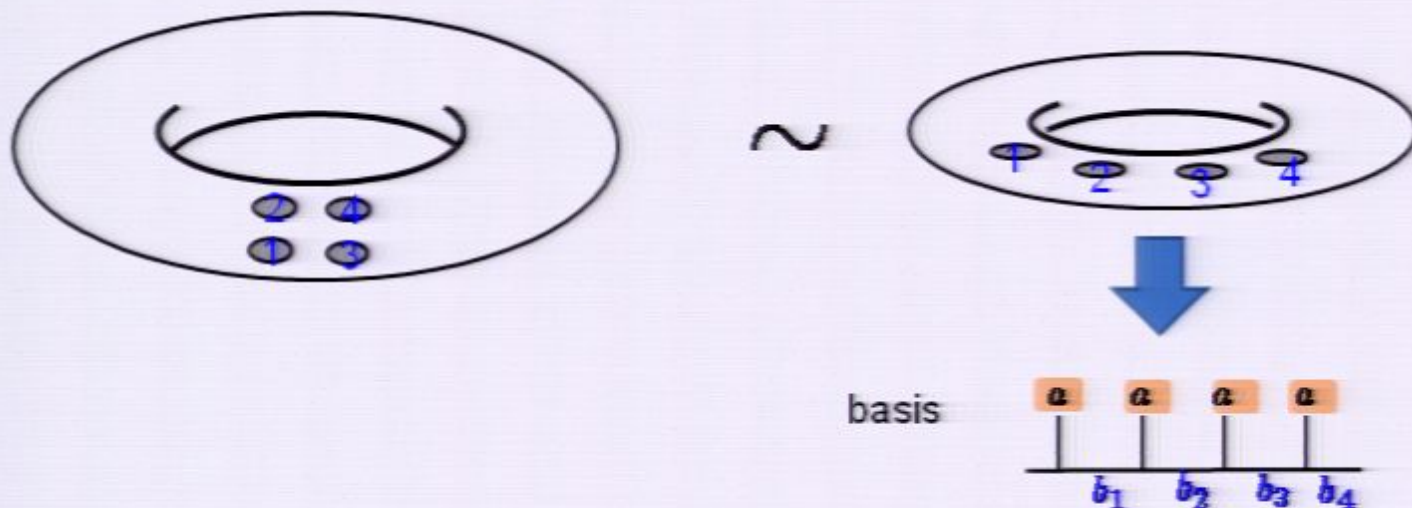
Next-to-nearest neighbor interactions:



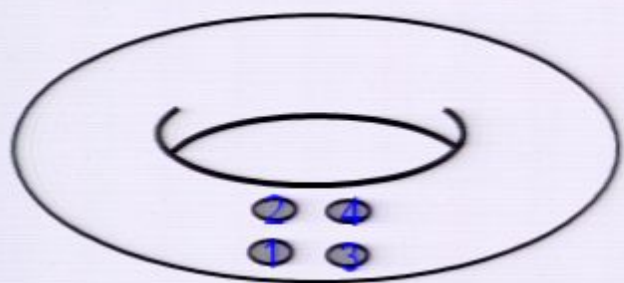
Example: anyonic array



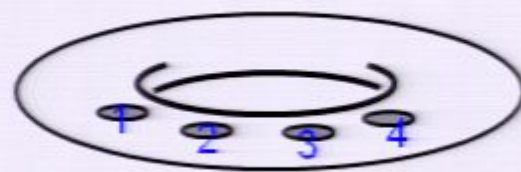
Example: anyonic array



Example: anyonic array



\sim

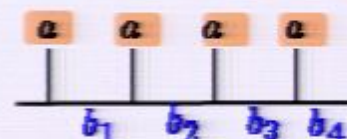


nearest-
neighbor
interaction

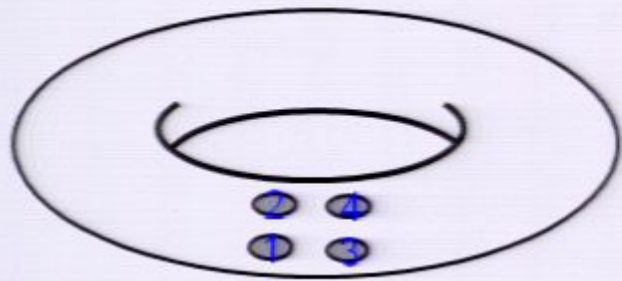
$\Pi =$



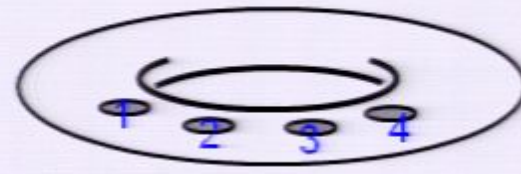
basis



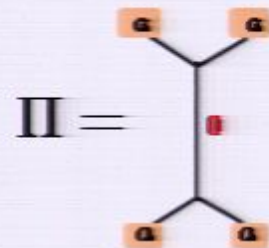
Example: anyonic array



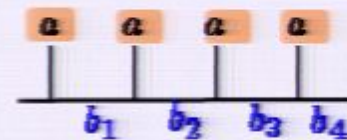
\sim



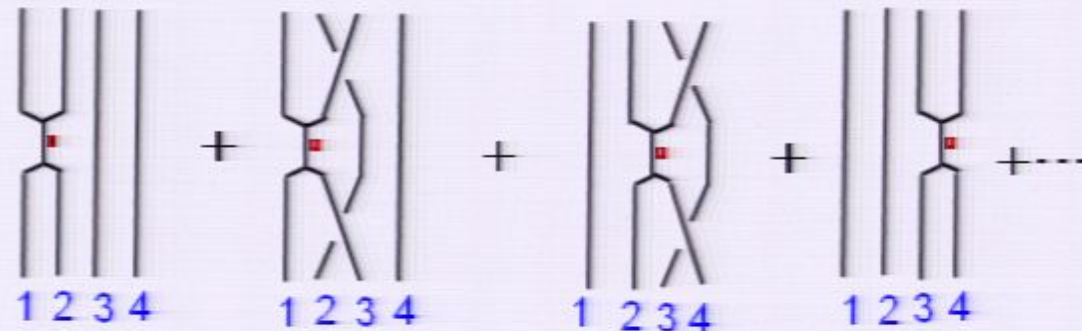
nearest-neighbor interaction



basis



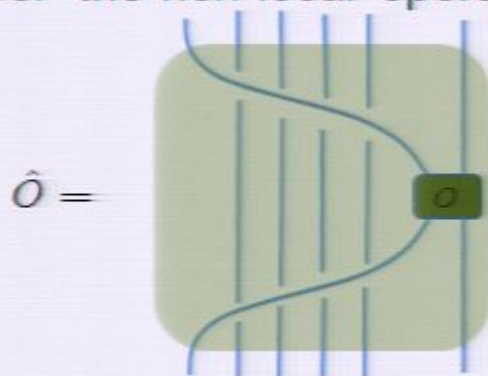
Hamiltonian:



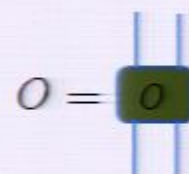
Braiding-type non-local operators

$$\langle \Psi_{\{V\}} | \hat{O} | \Psi_{\{V\}} \rangle = \langle \Psi_{\{W\}} | O | \Psi_{\{W\}} \rangle$$

for the non-local operator



for the “local” operator



isometries: $\{V\} \mapsto \{W\}$

computable with complexity of order $O(\text{poly}(\#crossings))$
(on causal cone)

This provides a technique for treating 2D anyon-arrangements (similar to fermionic tensor networks).

Rules for resolving crossings...

$$\begin{array}{c} a \quad b \\ \diagdown \quad \diagup \\ \diagup \quad \diagdown \\ \quad c \end{array} = R_c^{ab} \begin{array}{c} a \quad b \\ \diagdown \quad \diagup \\ \quad c \end{array}$$

$$\begin{array}{c} a \quad b \\ \diagup \quad \diagdown \\ \diagdown \quad \diagup \\ \quad c \end{array} = \overline{R_c^{ab}} \begin{array}{c} a \quad b \\ \diagdown \quad \diagup \\ \quad c \end{array}$$

compatibility with F-matrix implies the identities

$$\begin{array}{c} \diagdown \quad \diagup \\ \diagup \quad \diagdown \\ \quad | \end{array} = \begin{array}{c} \diagdown \quad \diagup \\ \quad | \end{array}$$

$$\begin{array}{c} \diagdown \quad \diagup \\ \quad | \end{array} = \begin{array}{c} \diagdown \quad \diagup \\ \diagup \quad \diagdown \\ \quad | \end{array}$$

$$\begin{array}{c} | \end{array} \begin{array}{c} \diagdown \quad \diagup \\ \diagup \quad \diagdown \end{array} = \begin{array}{c} 0 \\ | \end{array} \begin{array}{c} \diagdown \quad \diagup \\ \diagup \quad \diagdown \end{array} = \begin{array}{c} 0 \\ \diagdown \quad \diagup \end{array} \begin{array}{c} \diagdown \quad \diagup \\ \diagup \quad \diagdown \end{array} = \begin{array}{c} \diagdown \quad \diagup \\ \diagup \quad \diagdown \end{array}$$

Rules for resolving crossings...

$$\begin{array}{c} a \quad b \\ \diagdown \quad \diagup \\ \diagup \quad \diagdown \\ c \end{array} = R_c^{ab} \begin{array}{c} a \quad b \\ \diagdown \quad \diagup \\ c \end{array}$$

$$\begin{array}{c} a \quad b \\ \diagup \quad \diagdown \\ \diagdown \quad \diagup \\ c \end{array} = \overline{R}_c^{ab} \begin{array}{c} a \quad b \\ \diagdown \quad \diagup \\ c \end{array}$$

compatibility with F-matrix implies the identities

$$\begin{array}{c} \diagdown \quad \diagup \\ \diagup \quad \diagdown \\ | \end{array} = \begin{array}{c} \diagdown \quad \diagup \\ | \quad \diagup \\ | \end{array}$$

$$\begin{array}{c} \diagdown \quad \diagup \\ | \quad \diagdown \\ | \end{array} = \begin{array}{c} \diagdown \quad \diagup \\ \diagdown \quad \diagup \\ | \end{array}$$

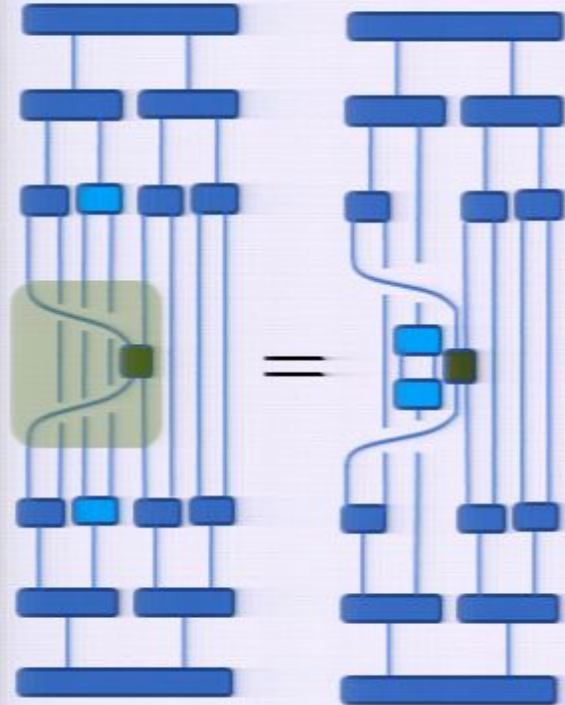
$$\begin{array}{c} | \quad \diagdown \quad \diagup \\ \diagup \quad \diagdown \end{array} = \begin{array}{c} \diagdown \quad \diagup \\ \diagdown \quad \diagup \end{array}$$

Example



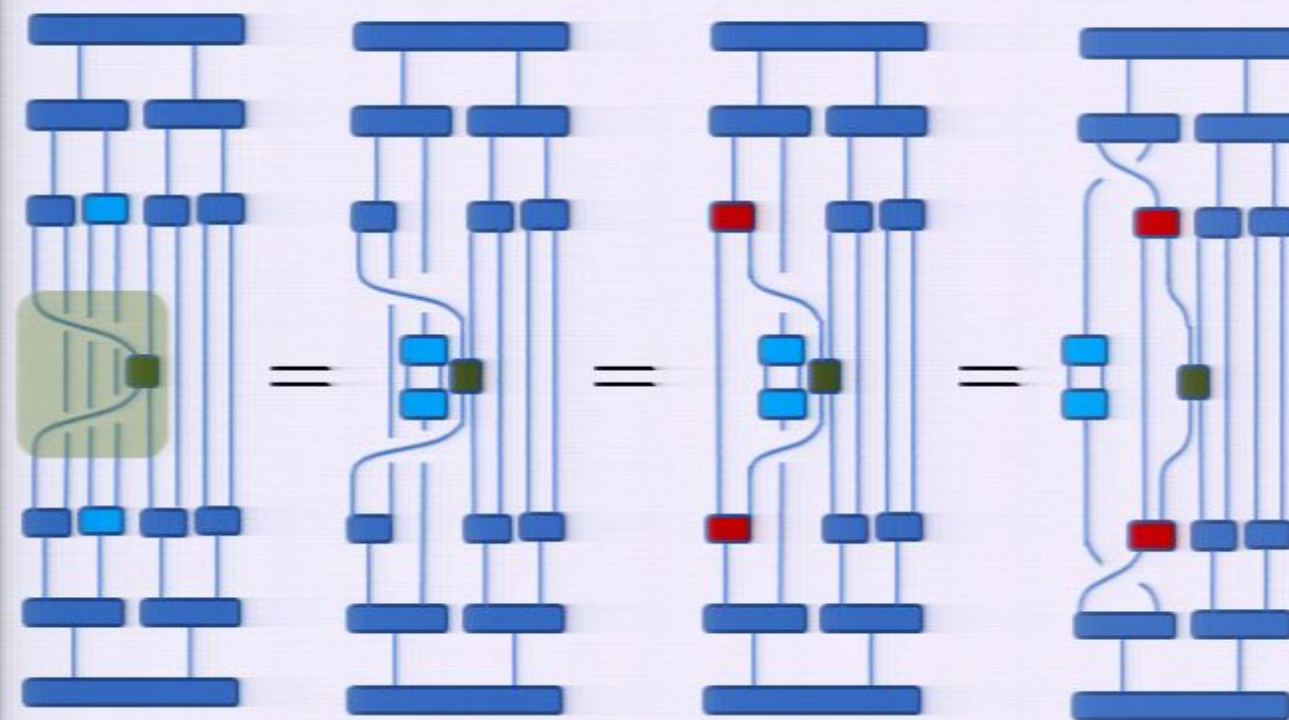
$$\langle \Psi_{\{v\}} | \hat{O} | \Psi_{\{v\}} \rangle$$

Example



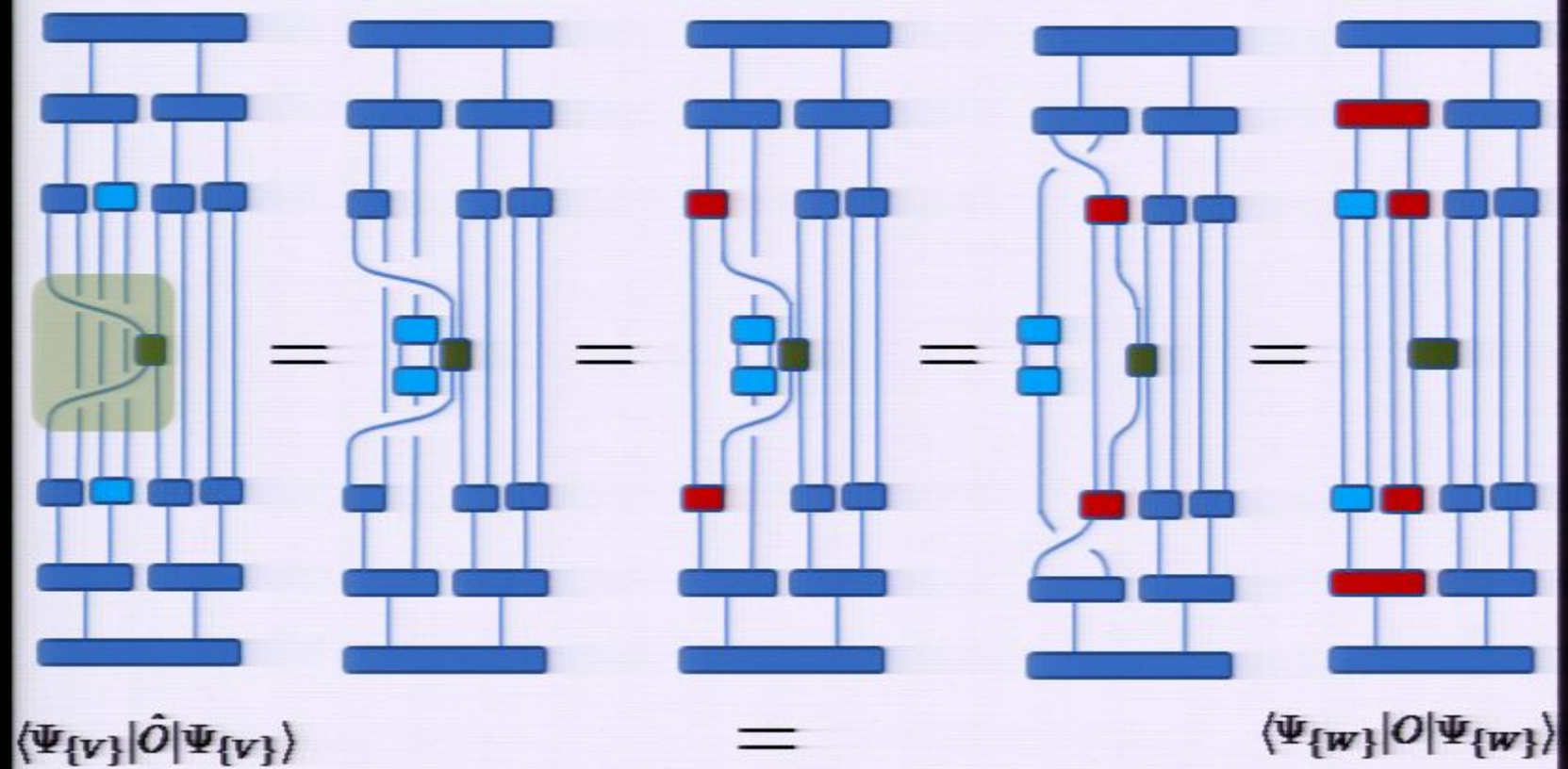
$$\langle \Psi_{\{v\}} | \hat{O} | \Psi_{\{v\}} \rangle$$

Example



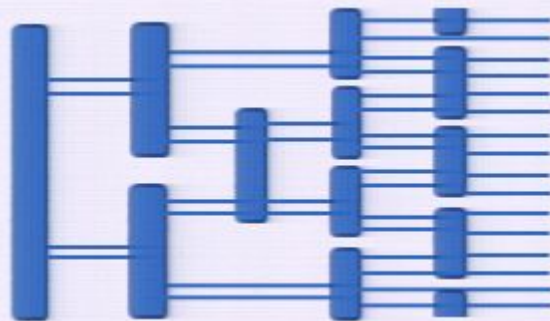
$$\langle \Psi_{\{v\}} | \hat{O} | \Psi_{\{v\}} \rangle$$

Example



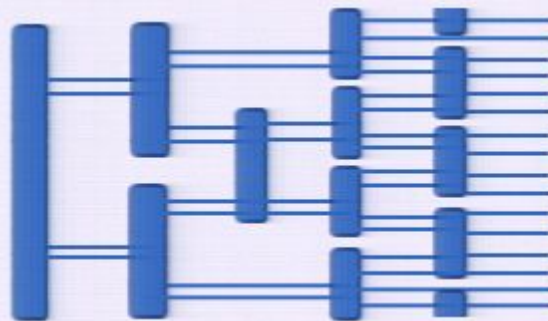
A new operational interpretation:

anyonic state preparation



A new operational interpretation:

anyonic state preparation



distillation of composite anyons

Universal computation with Fibonacci anyons

Typically used encoding:

$$|0\rangle \mapsto \begin{array}{c} 1 \quad 1 \quad 1 \quad 1 \\ \diagdown \quad \diagup \quad \diagdown \quad \diagup \end{array}$$

$$|1\rangle \mapsto \begin{array}{c} 1 \quad 1 \quad 1 \quad 1 \\ \diagdown \quad \diagup \quad \diagdown \quad \diagup \\ \quad \quad \quad \diagdown \quad \diagup \\ \quad \quad \quad \quad \quad 1 \end{array}$$

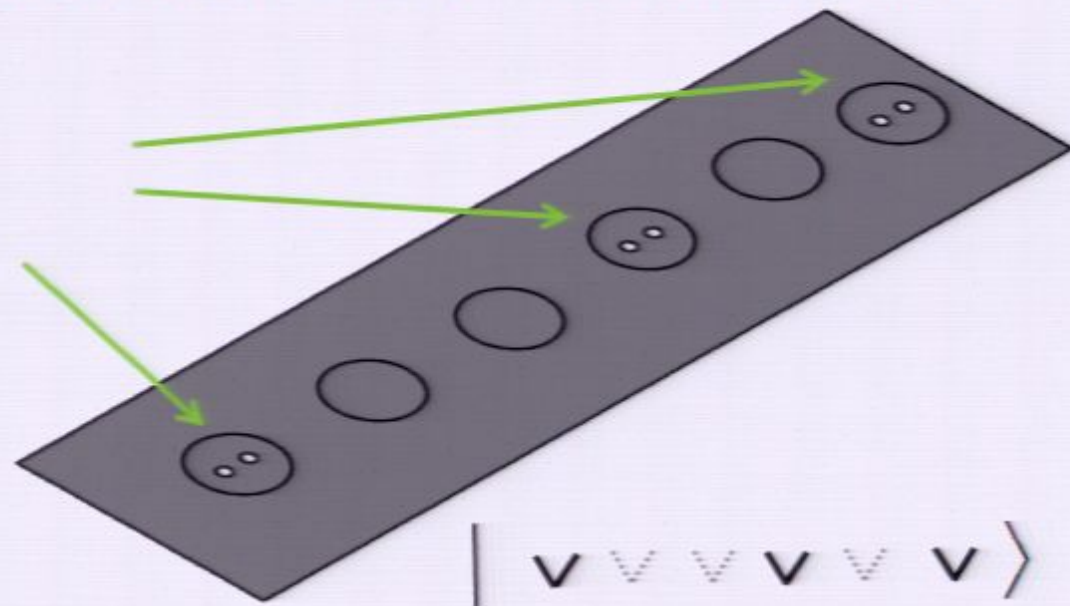
Need to be able to:

- prepare $|0\rangle^{\otimes n} = \overbrace{\begin{array}{c} 1 \quad 1 \quad 1 \quad 1 \\ \diagdown \quad \diagup \quad \diagdown \quad \diagup \end{array} \dots \begin{array}{c} 1 \quad 1 \\ \diagdown \quad \diagup \end{array}}^{2n \text{ pairs}}$
- execute braids/twists
- measure anyonic charge (repeat)

Dealing with noisy initial states

idealized scenario: Suppose **pair creation is probabilistic**

pair creation only
successful in these
(unknown) locations



How to deal with this situation?

Proposal 1: make **measurements** to determine presence/absence of pairs

Proposal 2: concentrate entropy into a certain sector by **reversible gates**

Sweeping dirt into the bin in a closed house

$\mathcal{H}_{\text{pairs}}$: subspace spanned by states of *arbitrary* (non-zero) *number of particle pairs* distributed among specified ("potential defect") locations

Example state: $\left| \begin{array}{ccccccc} \vee & \vee & \vee & \vee & \vee & \vee & \vee \end{array} \right\rangle$

Assumption ρ (mixed state) on $\mathcal{H}_{\text{pairs}} \subset \mathcal{H}$

Sweeping dirt into the bin in a closed house

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after
distillation: $U \rho U^\dagger = \left| \begin{array}{c} \vee \end{array} \right\rangle \left\langle \begin{array}{c} \vee \end{array} \right|$

pure state of one

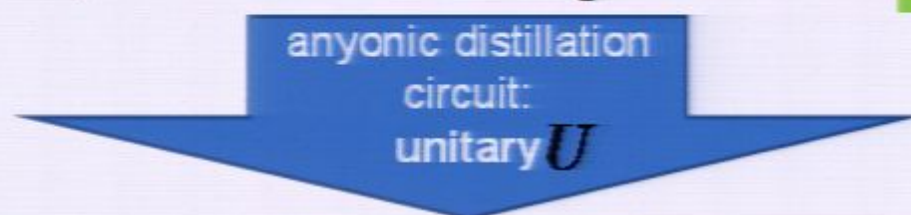
pair

Sweeping dirt into the bin in a closed house

$\mathcal{H}_{\text{pairs}}$: subspace spanned by states of *arbitrary* (non-zero) *number of particle pairs* distributed among specified ("potential defect") locations

Example state: $\left| \begin{array}{ccccccc} \vee & \vee & \vee & \vee & \vee & \vee & \vee \end{array} \right\rangle$

Assumption ρ (mixed state) on $\mathcal{H}_{\text{pairs}} \subset \mathcal{H} \cong \mathcal{H}_{\text{encoded}} \otimes \mathcal{H}_{\text{bin}}$



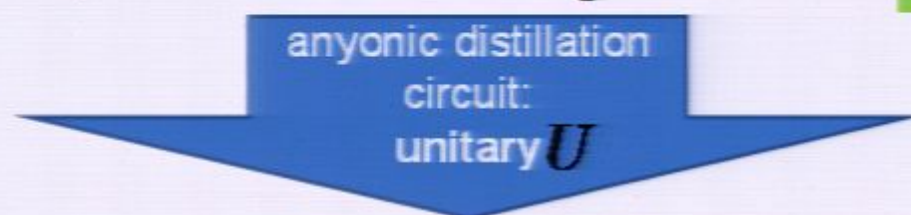
after distillation: $U \rho U^\dagger = \underbrace{|\vee\rangle\langle\vee|}_{\text{pure state of one encoded pair}} \otimes \underbrace{\sigma(\rho)_{\text{bin}}}_{\text{mixed state on "bin"}}$

Sweeping dirt into the bin in a closed house

$\mathcal{H}_{\text{pairs}}$: subspace spanned by states of *arbitrary* (non-zero) *number of particle pairs* distributed among specified ("potential defect") locations

Example state: $\left| \begin{array}{ccccccc} \vee & \vee & \vee & \vee & \vee & \vee & \vee \end{array} \right\rangle$

Assumption ρ (mixed state) on $\mathcal{H}_{\text{pairs}} \subset \mathcal{H} \cong \mathcal{H}_{\text{encoded}} \otimes \mathcal{H}_{\text{bin}}$



after distillation: $U \rho U^\dagger = \left| \vee \right\rangle \left\langle \vee \right| \otimes \sigma(\rho)_{\text{bin}}$

Main property
(subsystem
code):

$$U \mathcal{H}_{\text{pairs}} \subset \left| \vee \right\rangle \otimes \mathcal{H}_{\text{bin}}$$

Concatenated codes

physical (qu)bits



logical bits $0_L = 000$

(repetition) code $1_L = 111$



logical bits of
concatenated
(repetition) code

$0_{L'} = 0_L 0_L 0_L$

$1_{L'} = 1_L 1_L 1_L$

Composite anyons

$$\mathcal{H} \cong \mathcal{H}_{\text{encoded}} \otimes \mathcal{H}_{\text{bin}}$$

“physical”
anyon pairs
(defects)

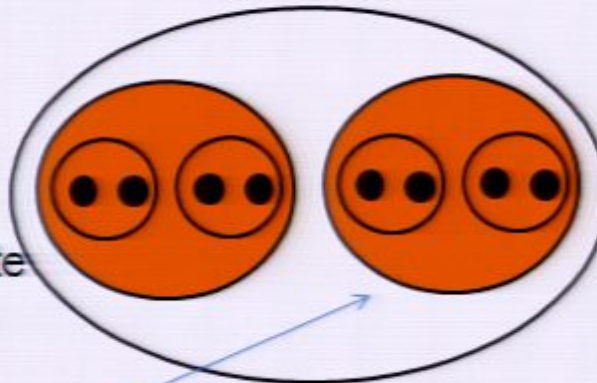


\vee

pair of
“encoded”
composite anyons



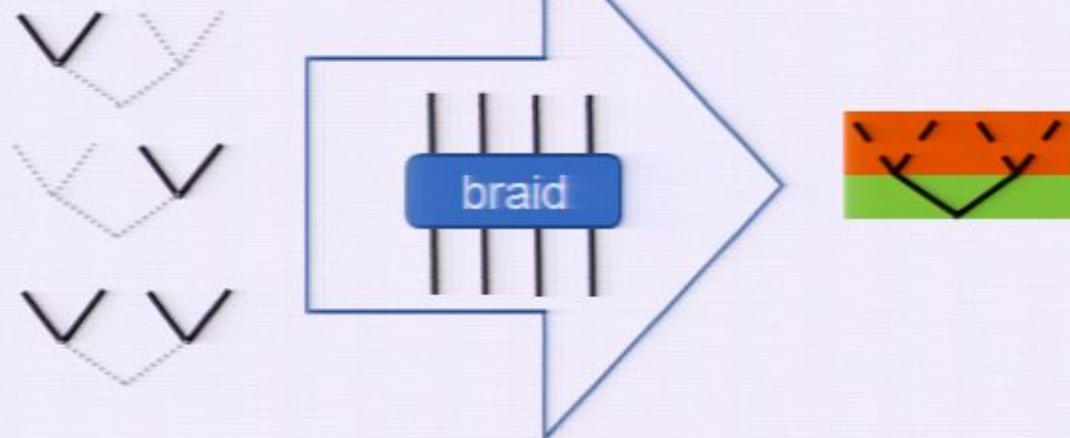
pairs of
(doubly)
composite
anyons



composite anyons are equivalent to bare anyons wrt. computation
– complexity increase only quadratically in circumference

Constructing an anyonic distillation circuit

Main building block:

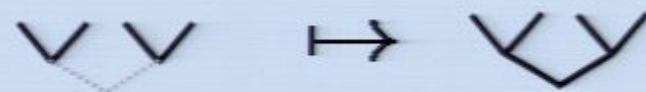


Constructing an anyonic distillation circuit

Main building block:



Construction: use purebraid

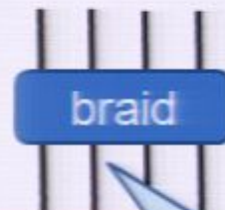


+ injection weave

S. Simon, N. Bonesteel, M. Freedman, N. Petrovic, L. Hormozi:
"Topological Quantum Computing with Only One Mobile Quasiparticle"

Constructing an anyonic distillation circuit

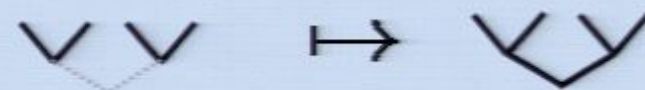
Main building block:



Final circuit:
recursive construction



Construction: use purebraid



+ injection weave

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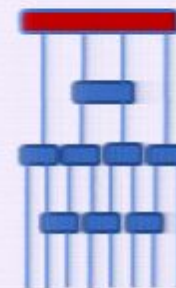
Relation to anyonic entanglement renormalization

Consider an anyonic entanglement RG with a fixed toplevel state:

e.g.,  = $\overset{1}{\vee} \overset{1}{\vee} \overset{1}{\vee} \overset{1}{\vee}$

This is the “target” state of composite anyons

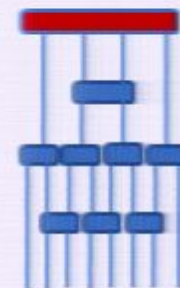
anyonic entanglement RG = composite anyon distillation scheme
 -> characterization of “distillable” states



Relation to anyonic entanglement renormalization

Consider an anyonic entanglement RG with a fixed toplevel state:


e.g.,  = 




This is the “target” state of composite anyons

anyonic entanglement RG = composite anyon distillation scheme
 -> characterization of “distillable” states

Assume that every isometry/unitary is composed of *physically realizable* processes

fusion  = $\sum_{a,b,c} a \begin{array}{c} c \\ \swarrow \downarrow \end{array} b$

braiding  = $\sum_{a,b} a \begin{array}{c} \swarrow \searrow \\ \swarrow \searrow \end{array} b$ (+inverse)

Anyonic entanglement renormalization

Variational family for anyons with operational interpretation:

- state preparation circuit
- renormalization group scheme, connects to composite anyon distillation

Evidence for its descriptive power:

- exact description of multi-anyon chains at certain coupling strengths
- accurate numerical estimation of ground state energies & correlation functions for finite-size systems
- good agreement with CFT predictions in thermodynamic limit

*.....compare
to successes
of non-anyonic
version!*

Provides a new numerical toolbox:

- extends to 2D-arrangements of anyons
- inherits various optimization/evaluation algorithms

Computational
savings
achieved by
exploiting
Hilbert space
structure