

Title: Anderson localization and adiabatic quantum optimization

Date: Jul 04, 2010 11:15 AM

URL: <http://pirsa.org/10070006>

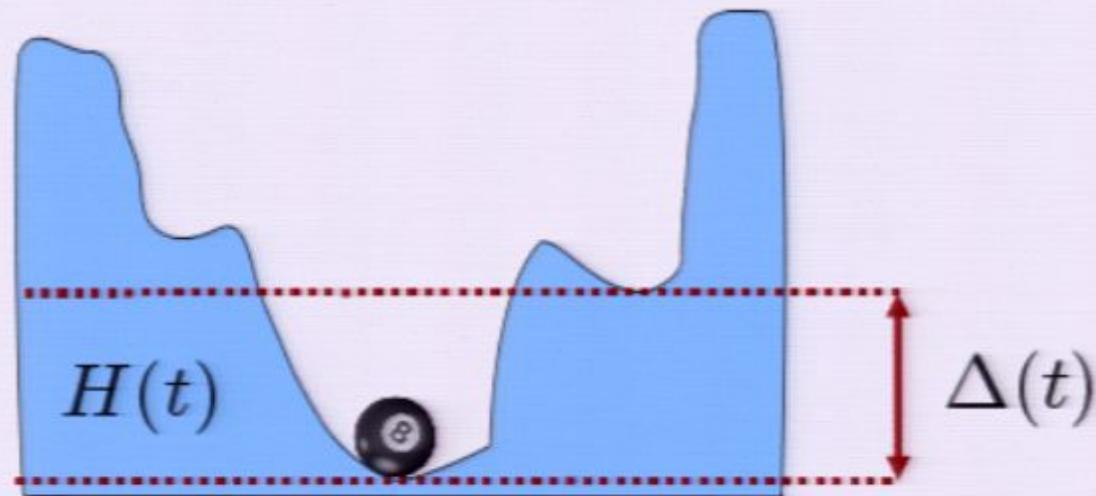
Abstract: Understanding NP-complete problems is a central topic in computer science. This is why adiabatic quantum optimization has attracted so much attention, as it provided a new approach to tackle NP-complete problems using a quantum computer. The efficiency of this approach is limited by small spectral gaps between the ground and excited states of the quantum computer's Hamiltonian.

We show that the statistics of the gaps can be analyzed in a novel way, borrowed from the study of quantum disordered systems in statistical mechanics. It turns out that due to a phenomenon similar to Anderson localization, exponentially small gaps appear close to the end of the adiabatic algorithm for large random instances of NP-complete problems. We show that this effect makes adiabatic quantum optimization fail, as the system gets trapped in one of the numerous local minima. We will also discuss recent developments including the effect of the exponential number of solutions and Hamiltonian path change.

Joint work with Boris Altshuler and Hari Krovi

Based on arXiv:0908.2782 and arXiv:0912.0746

Adiabatic evolution



Slowly varying $H(t)$ → Stays close to ground state

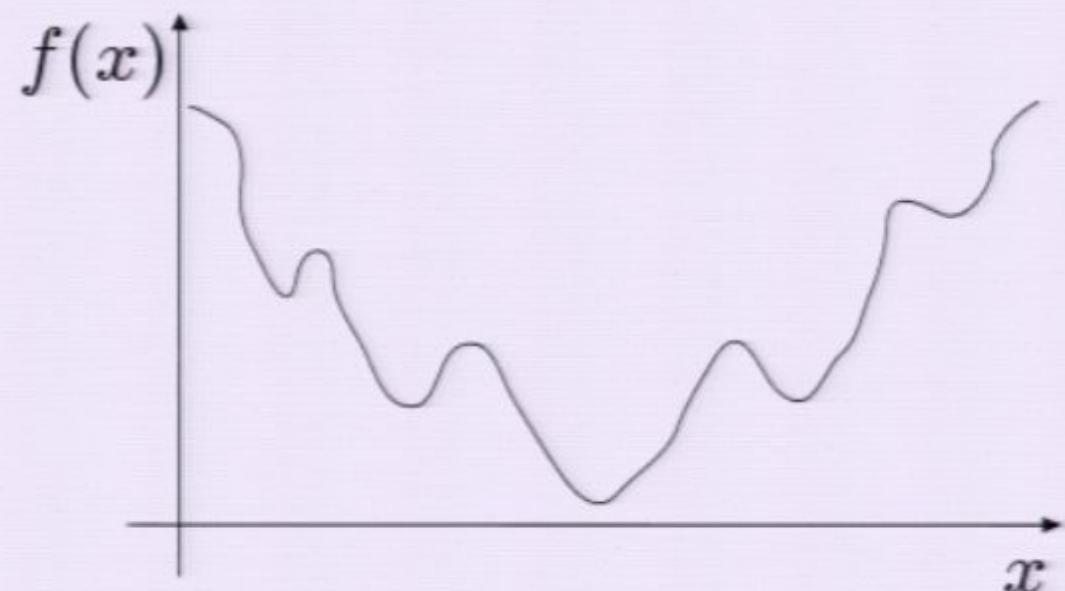
Probability of excitation depends on

- Total time T (slower is better)
- Gap $\Delta(t)$ (larger gap is better)

Adiabatic quantum optimization

[Farhi *et al.* '00]

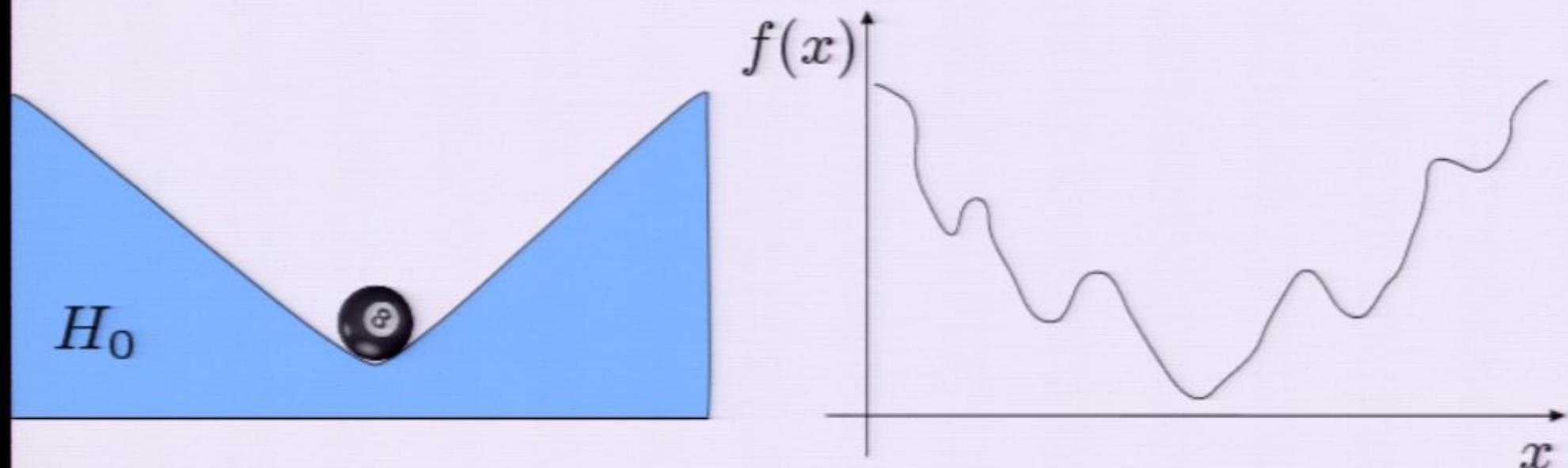
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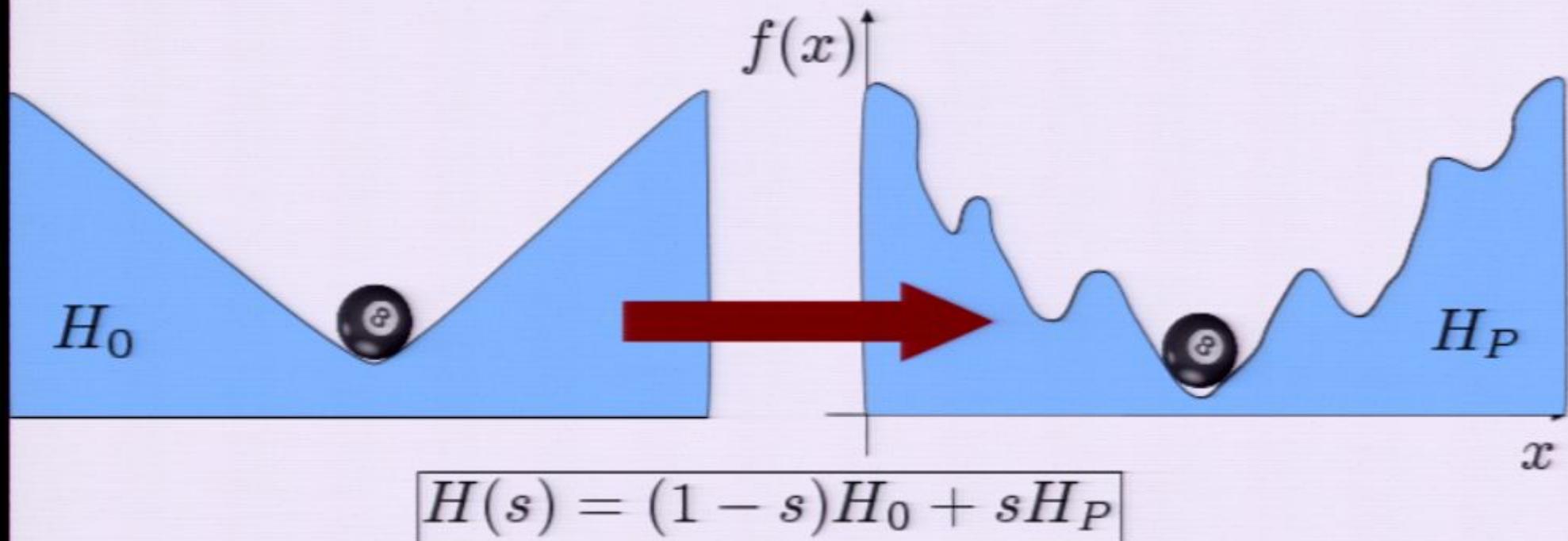
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Adiabatic quantum optimization

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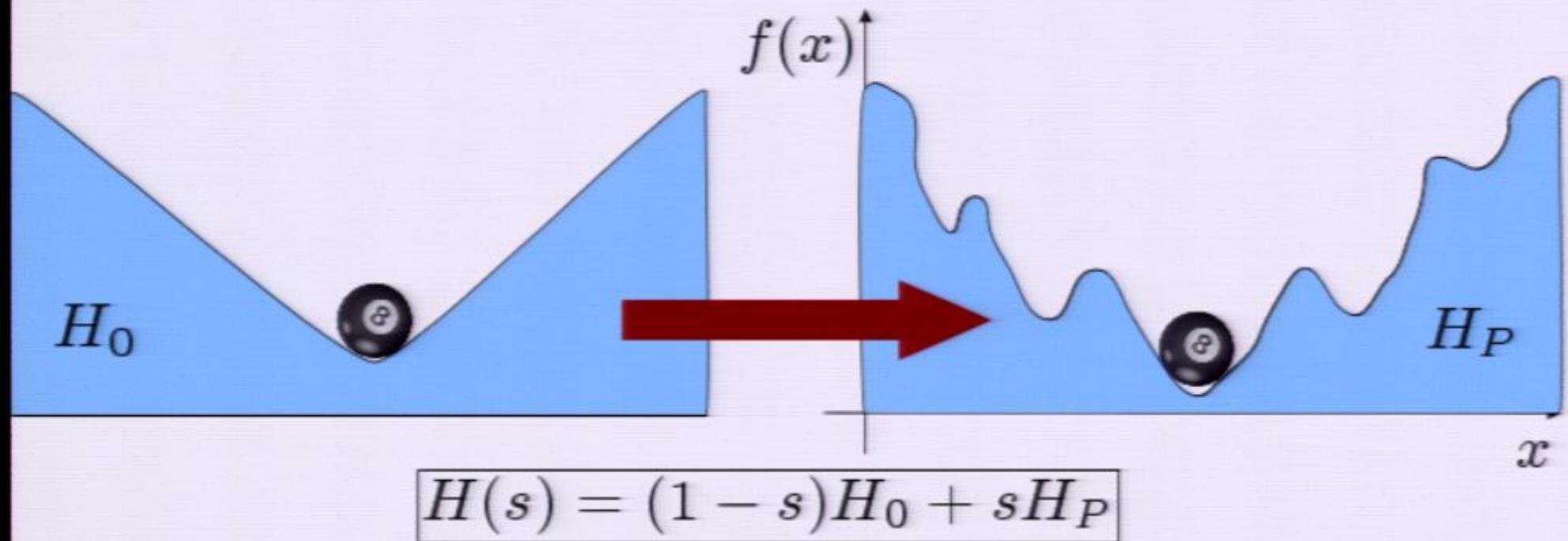


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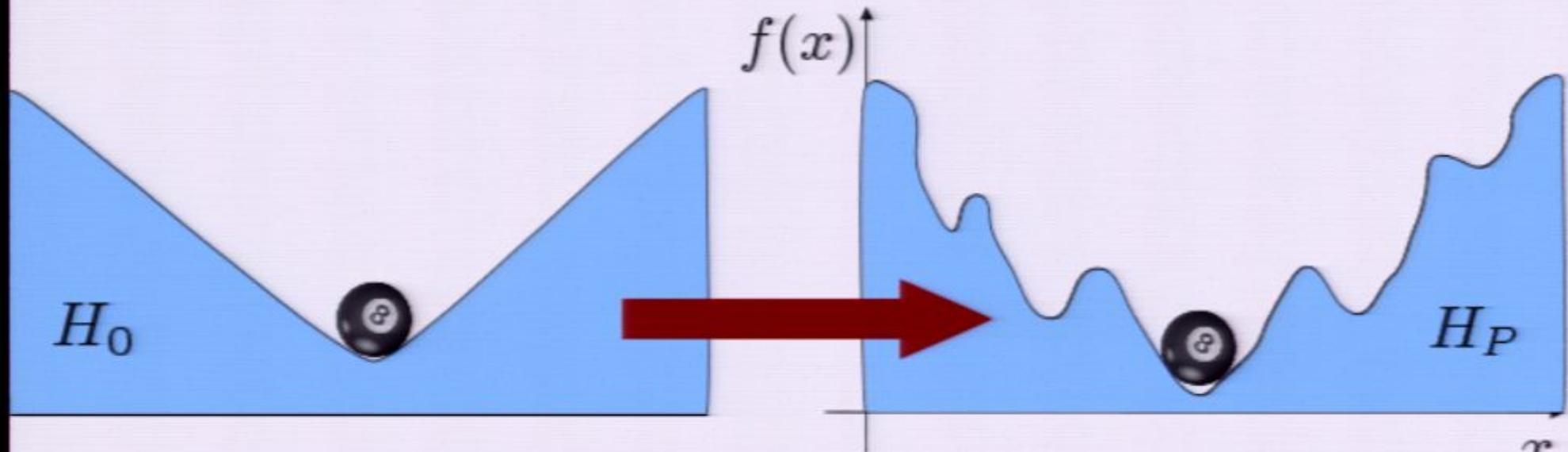
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- It is quantum! Unstructured search in time $O(\sqrt{N})$ (cf Grover)
[vanDam-Mosca-Vazirani'01,Roland-Cerf'02] 

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$$H(s) = (1 - s)H_0 + sH_P$$

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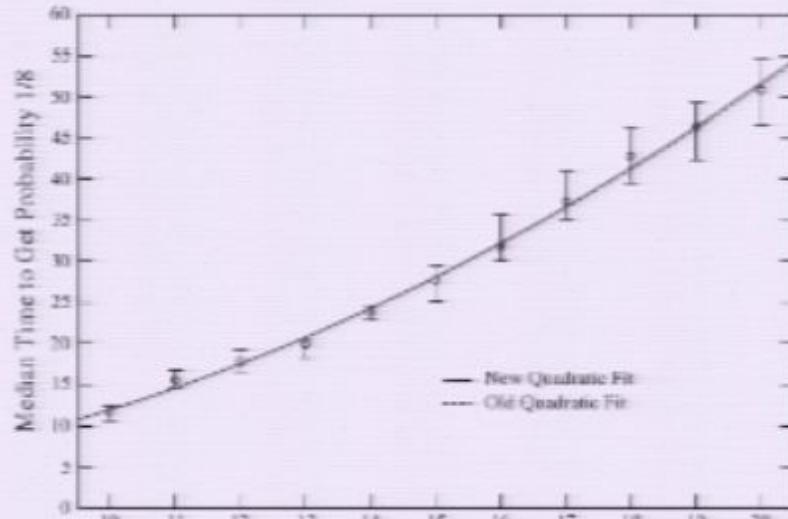
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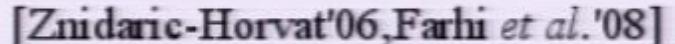
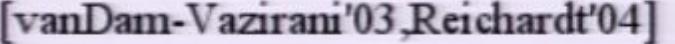
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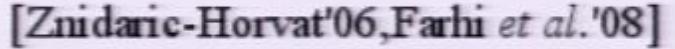
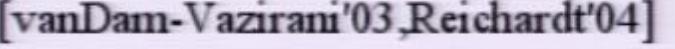
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But maybe typical gaps are only polynomial?

Exact-Cover 3 (EC3)

- NP complete problem (similar to 3-SAT)

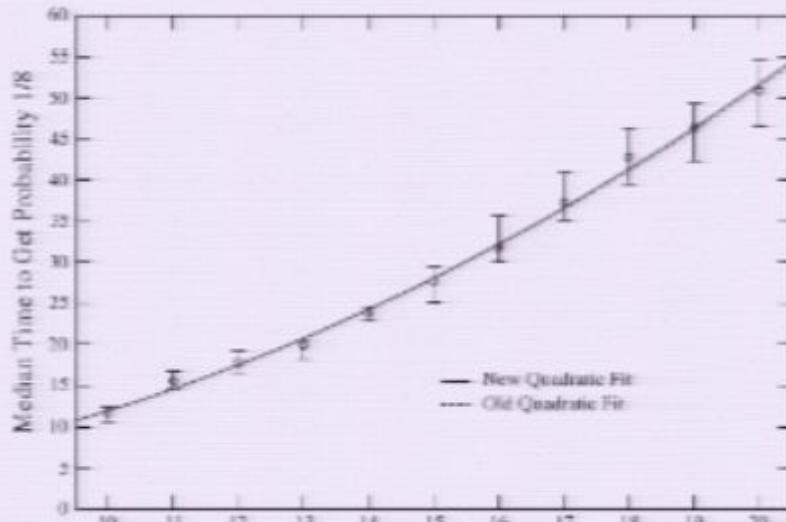
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- M clauses of 3 bits:
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 $= M - \sum_i B_i x_i + 2 \sum_{i,j} J_{ij} x_i x_j$

#clauses

#clauses with bit i

#clauses with bits i, j

Random instances

- Pick M clauses uniformly at random

Random instances

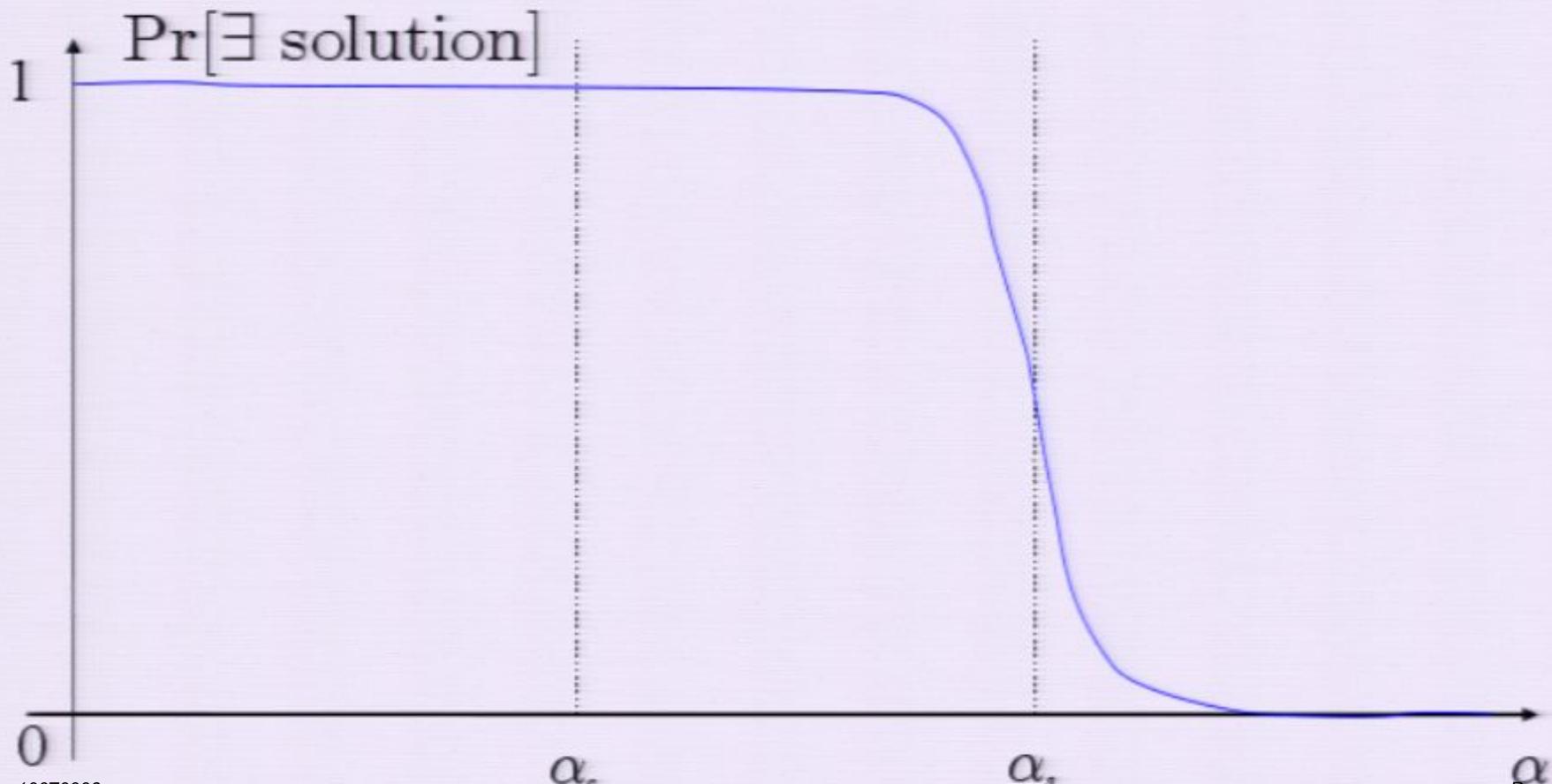
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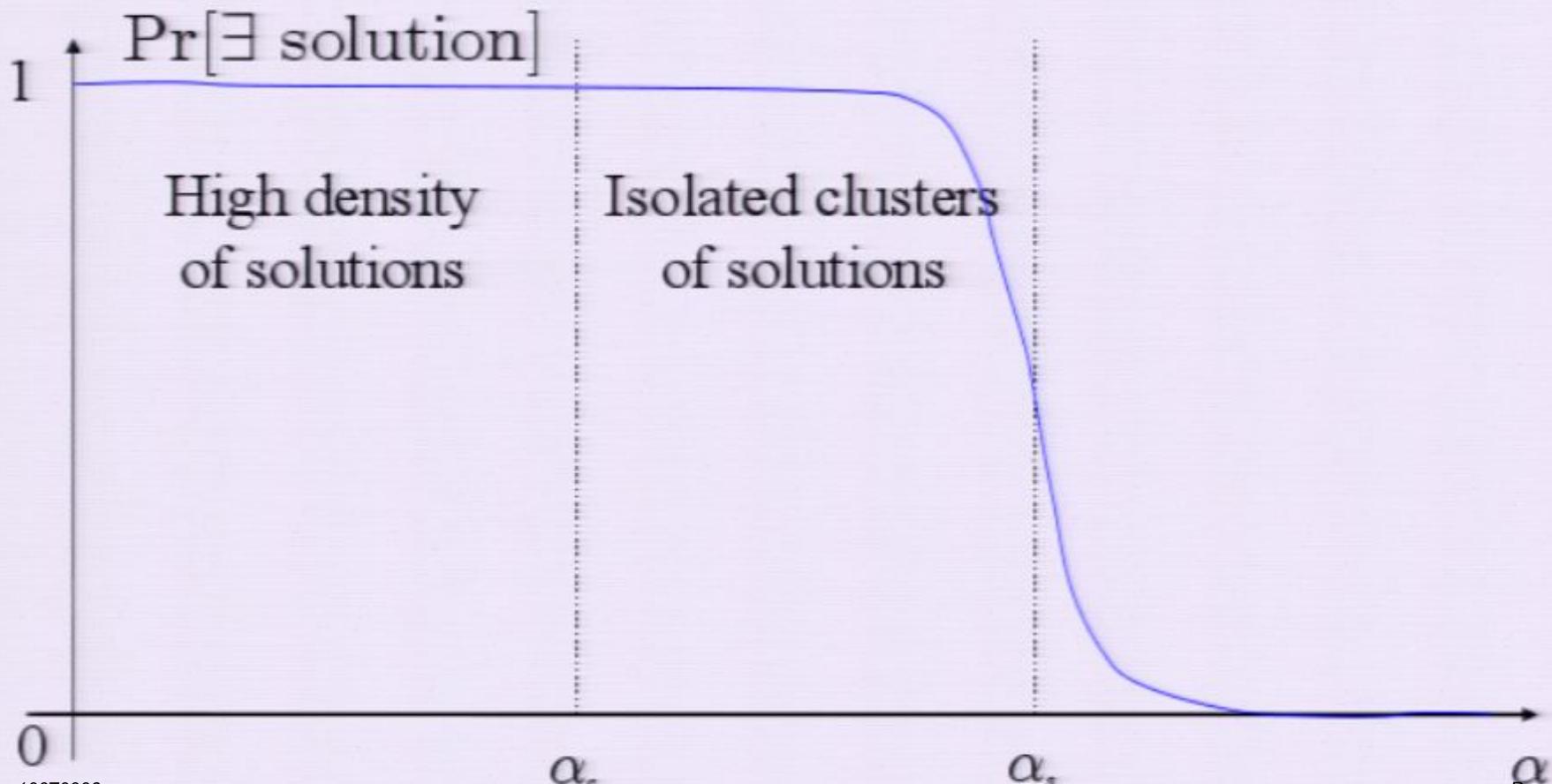
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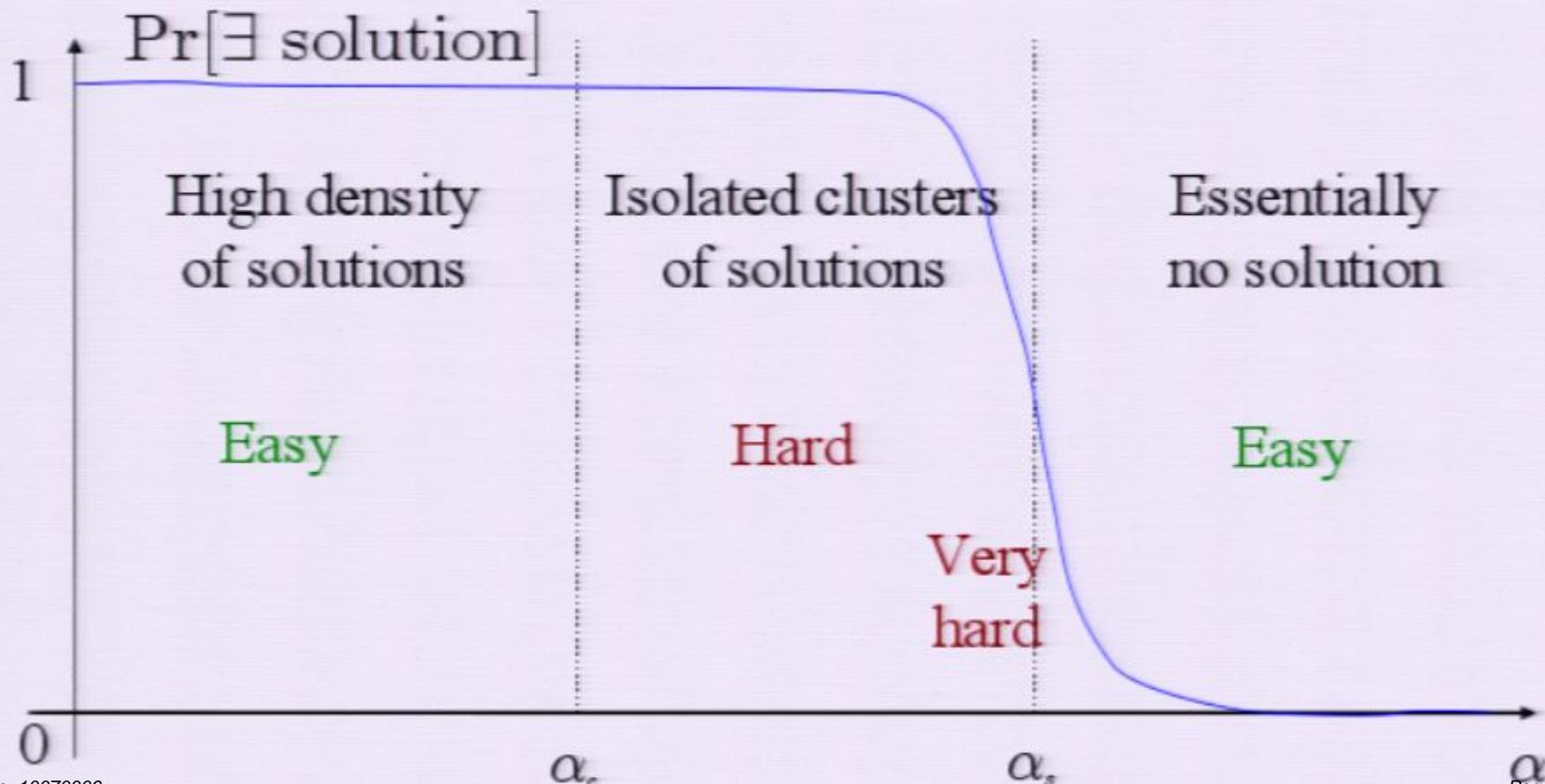
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The adiabatic algorithm

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Connection to Anderson localization

To study the spectrum of $H(s)$ close to $s=1$, we consider

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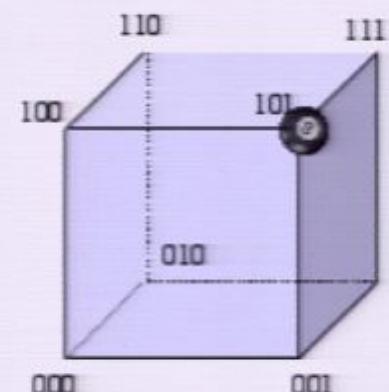
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⇒ Particle hopping on a hypercube

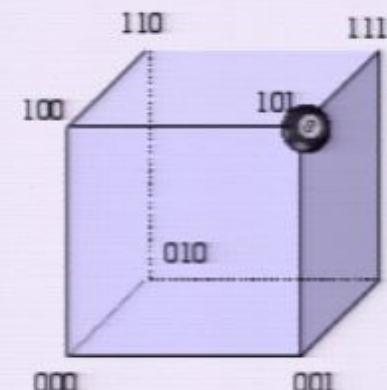


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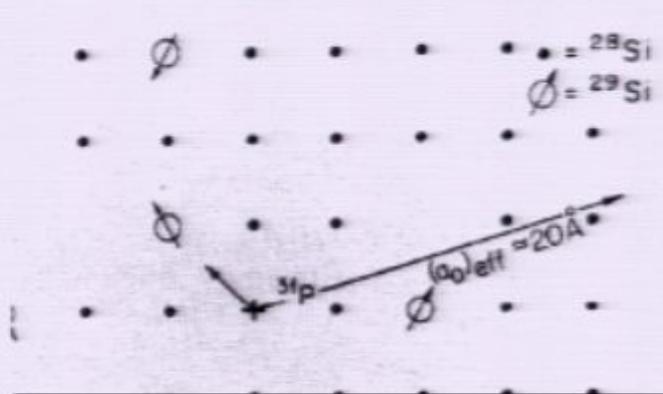
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- ⇒ Particle hopping on a hypercube
- ⇒ Similar to Anderson's tight binding model

Anderson localization

“Extended states become localized due to disorder”



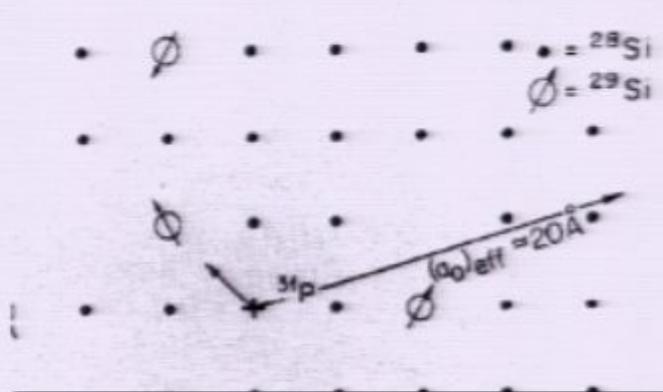
P. Anderson
Nobel Prize
Physics 1977

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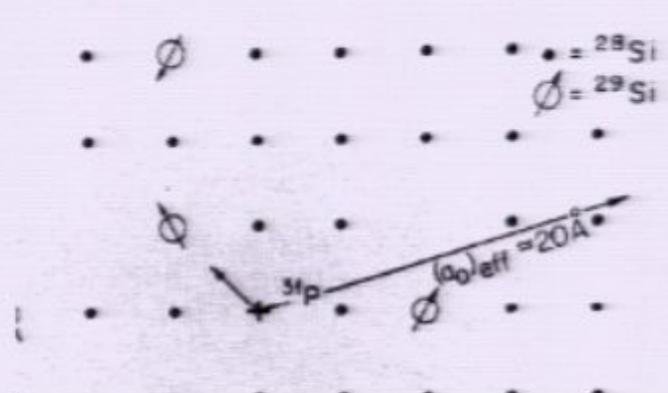
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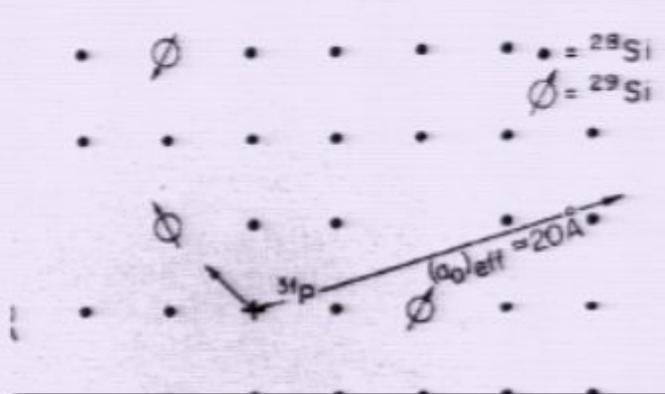
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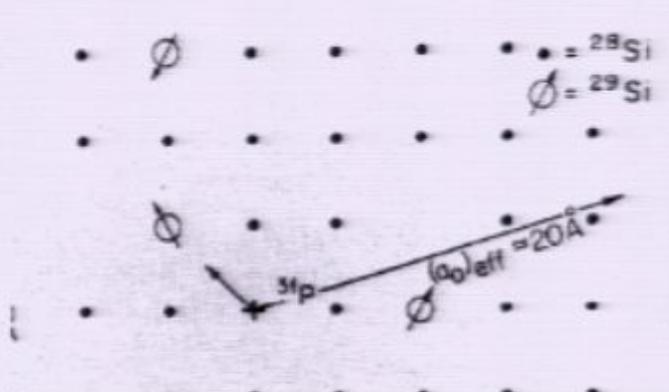
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$\lambda > \lambda_c \rightarrow$ Extended state \rightarrow Metal

$\lambda < \lambda_c \rightarrow$ Localized state \rightarrow Insulator



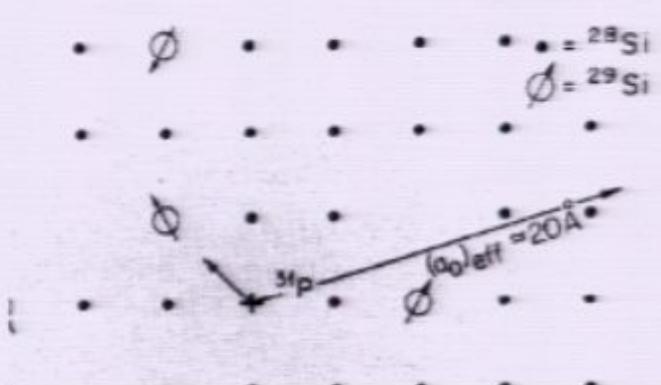
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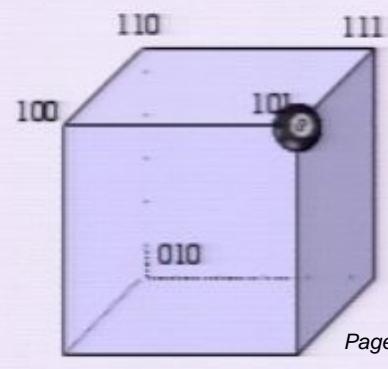
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In our case:

- Hypercube with coupling λ
- Energies from random Exact-Cover 3



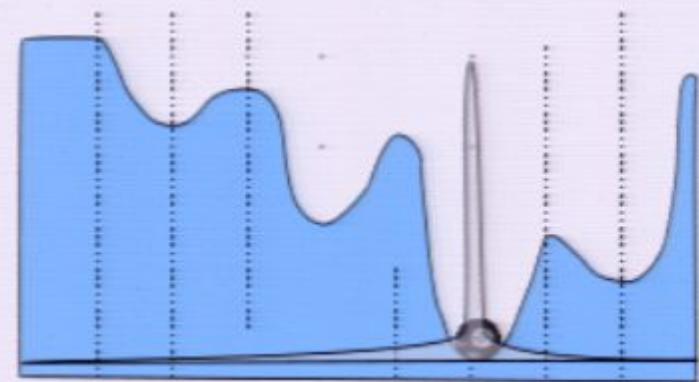
Localized and extended states

NEC Laboratories
America
Relentless passion for innovation

Localized and extended states

$$\lambda = 0$$

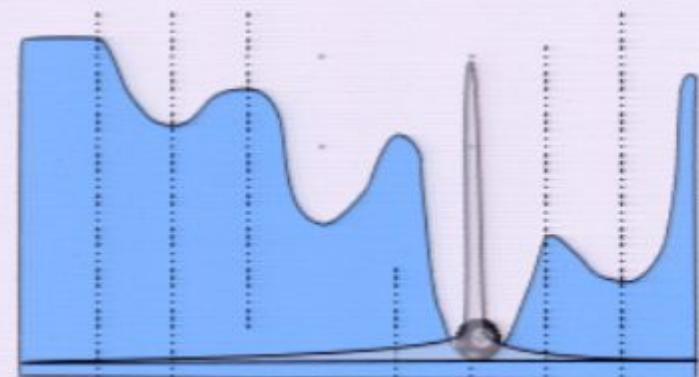
- State is localized



Localized and extended states

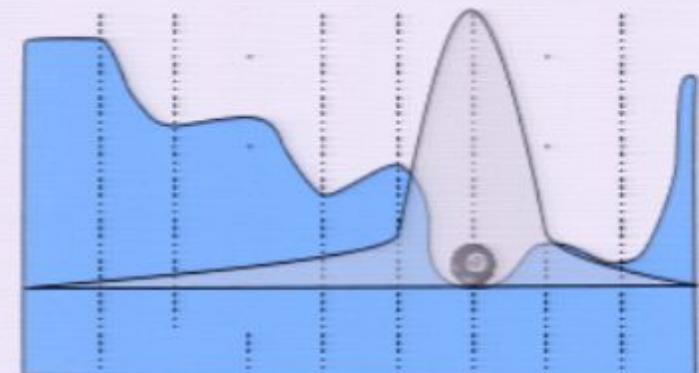
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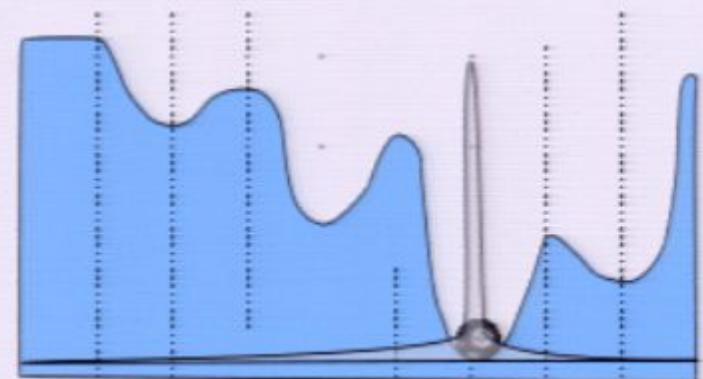
- Transverse field “spreads” the state



Localized and extended states

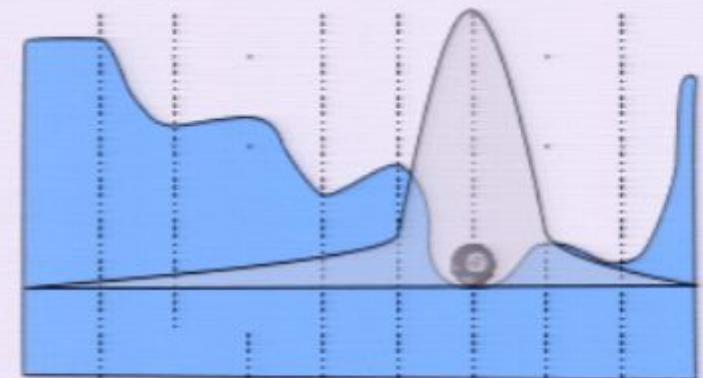
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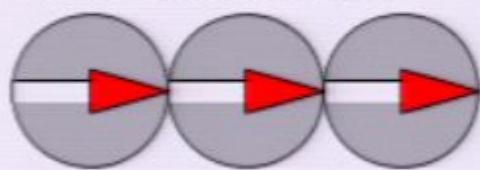
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$$\lambda \gg 1$$

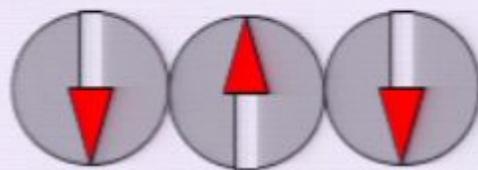
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Localized and extended states

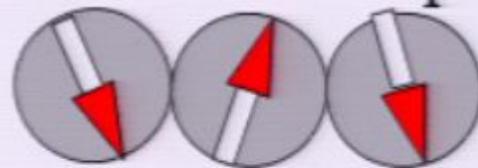
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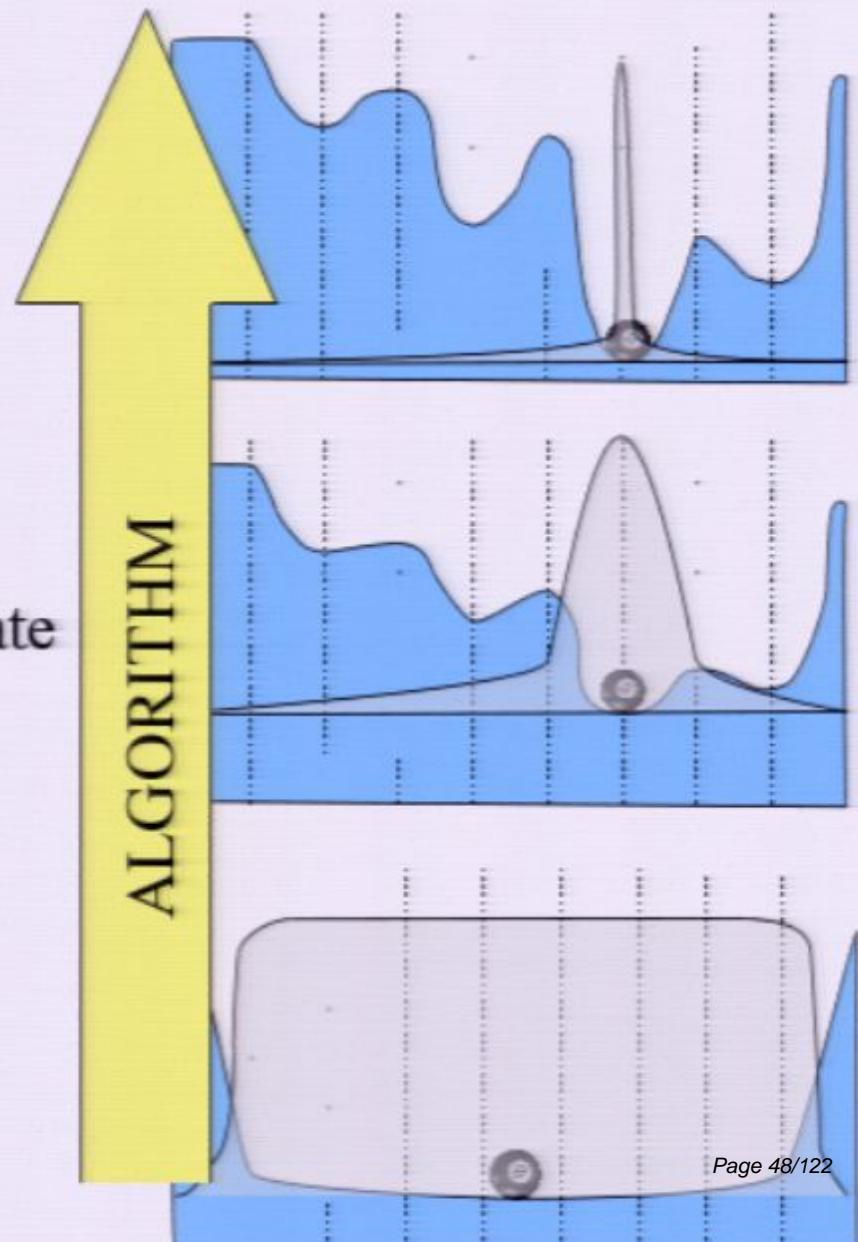
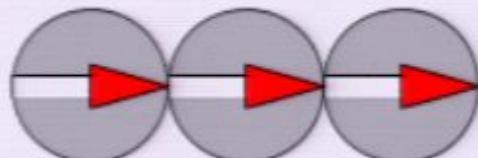
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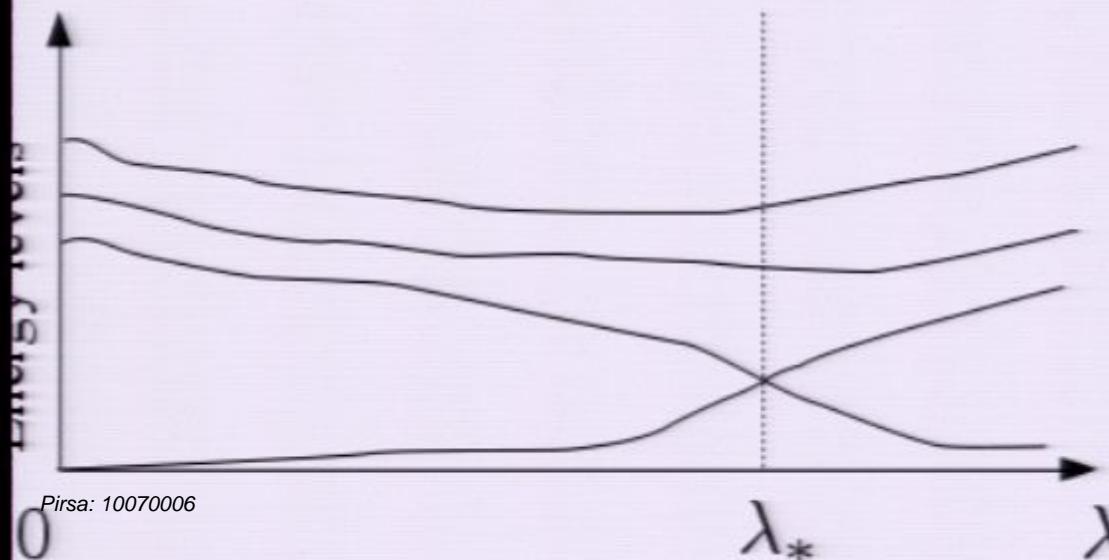
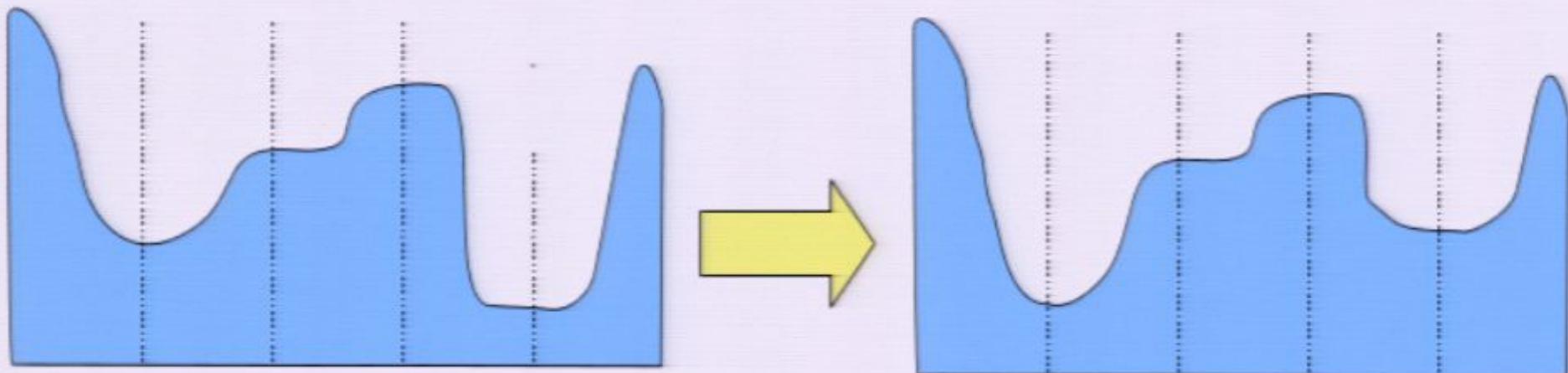


Tunneling: extended state

What if a local minimum later becomes the global minimum?

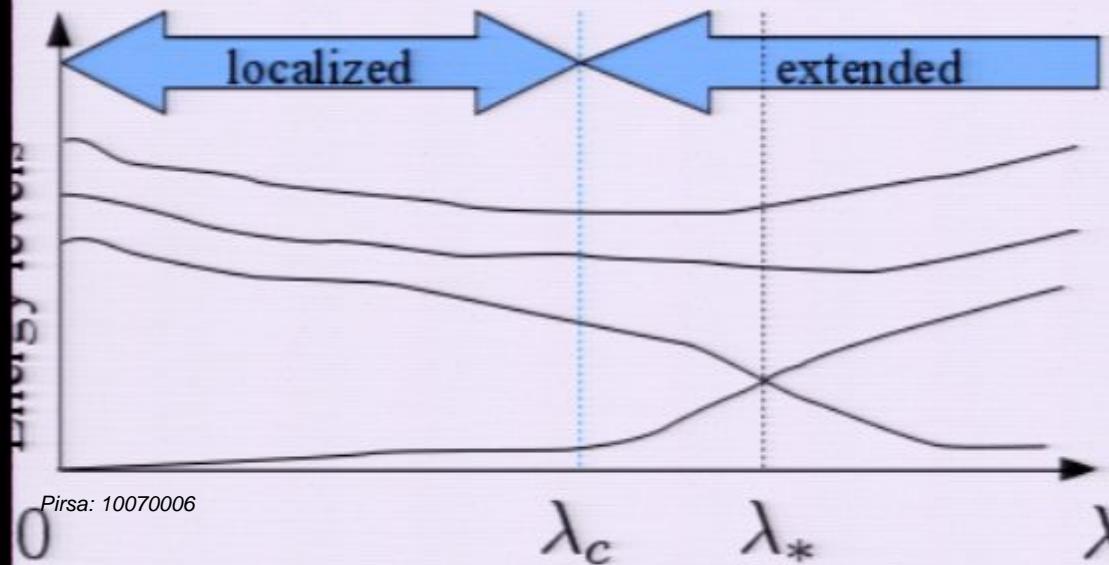
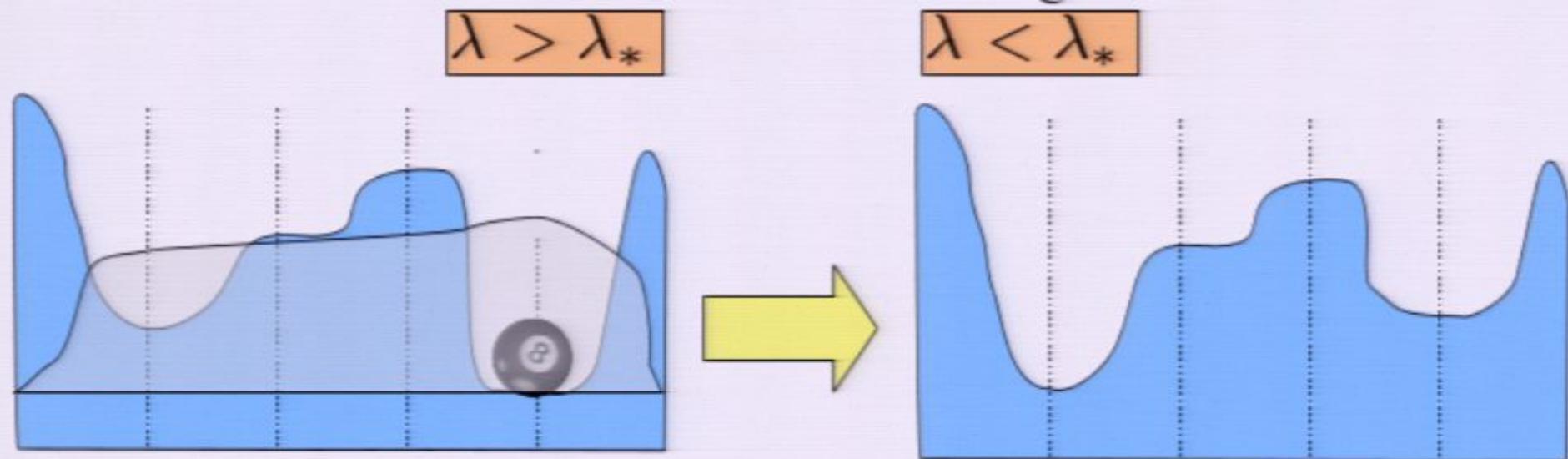
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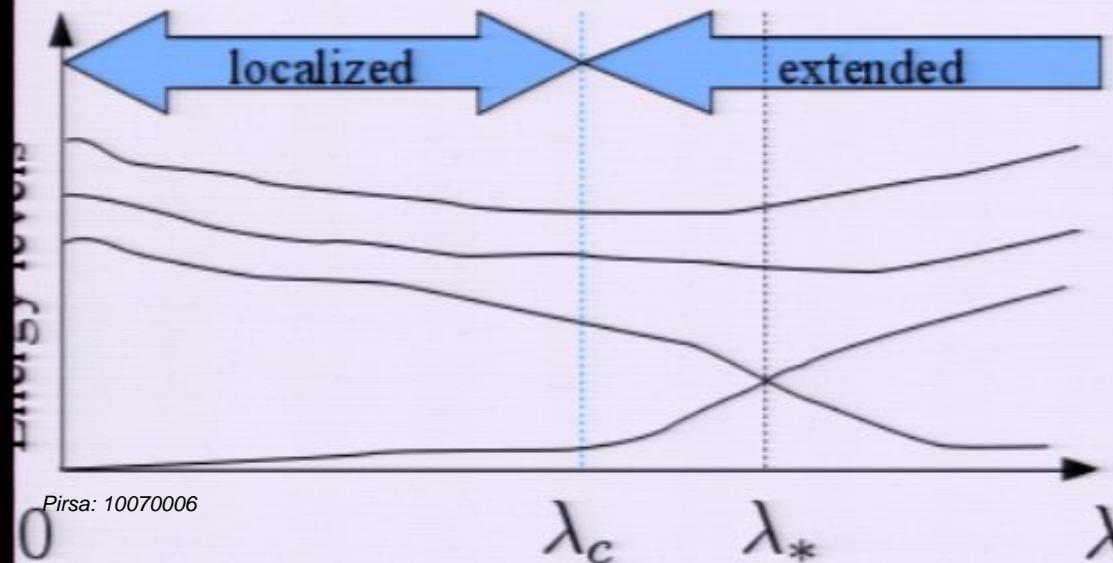
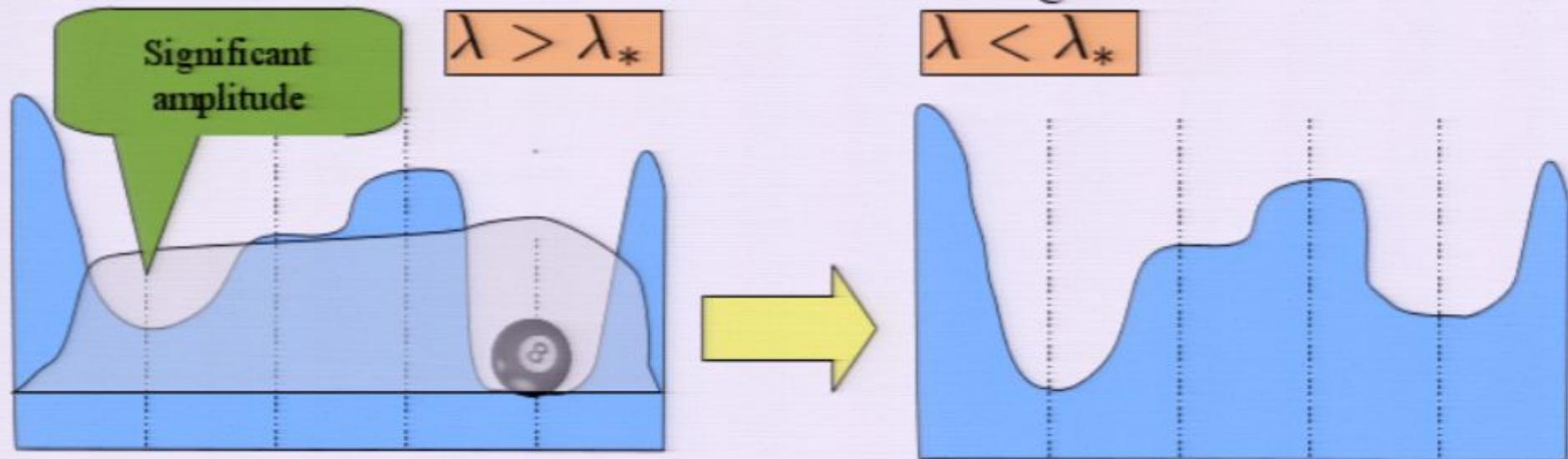
Tunneling: extended state

What if a local minimum later becomes the global minimum?



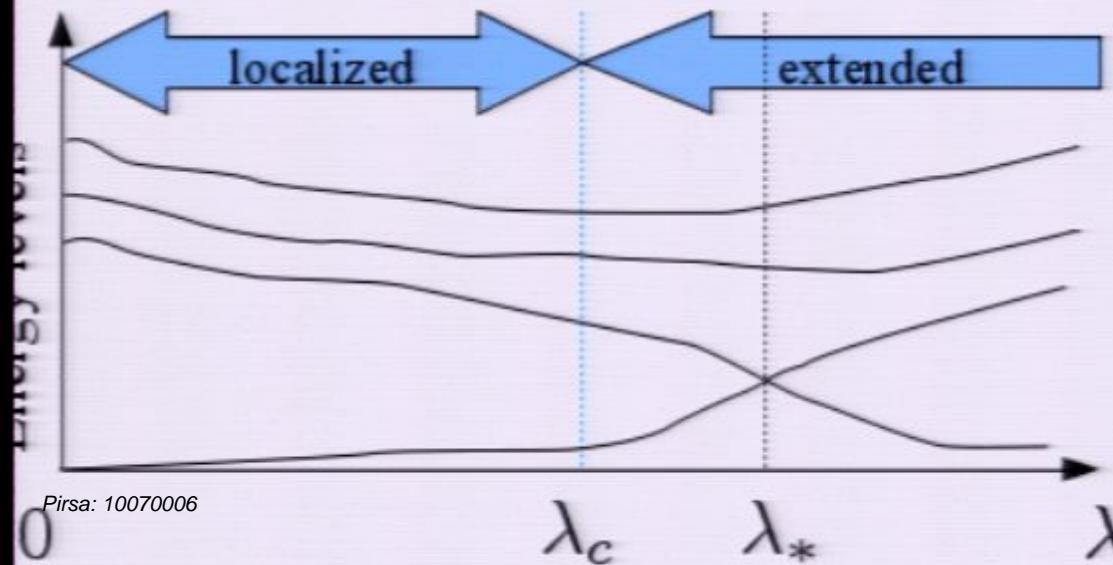
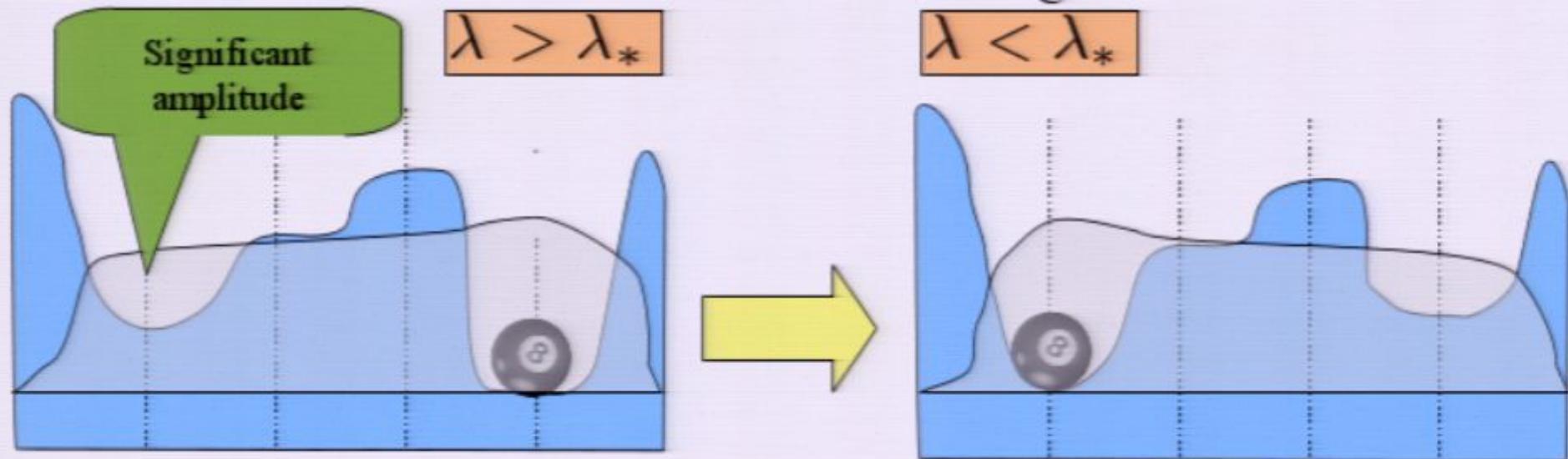
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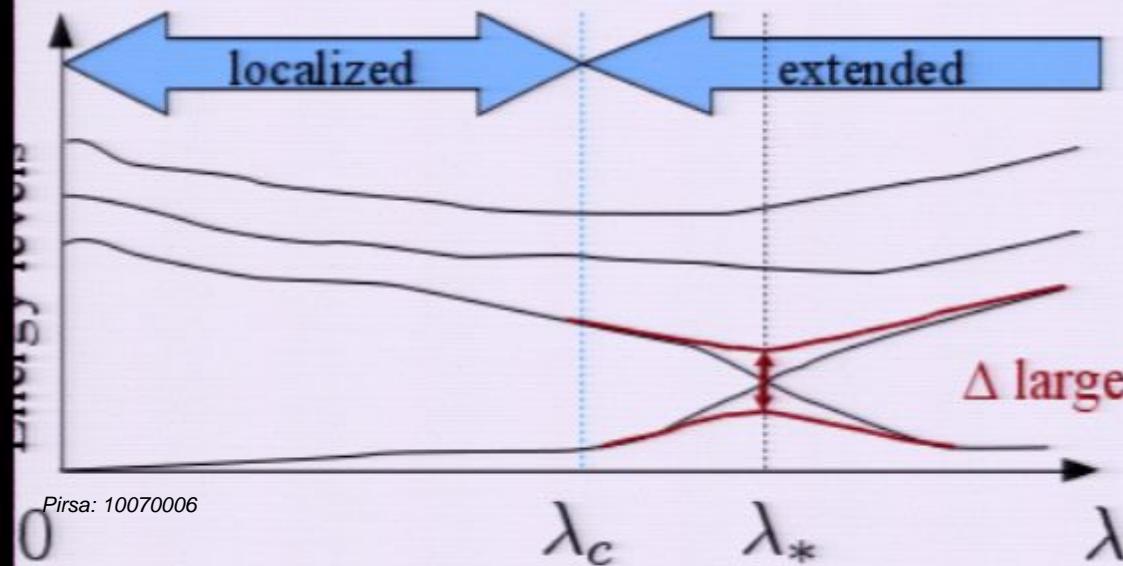
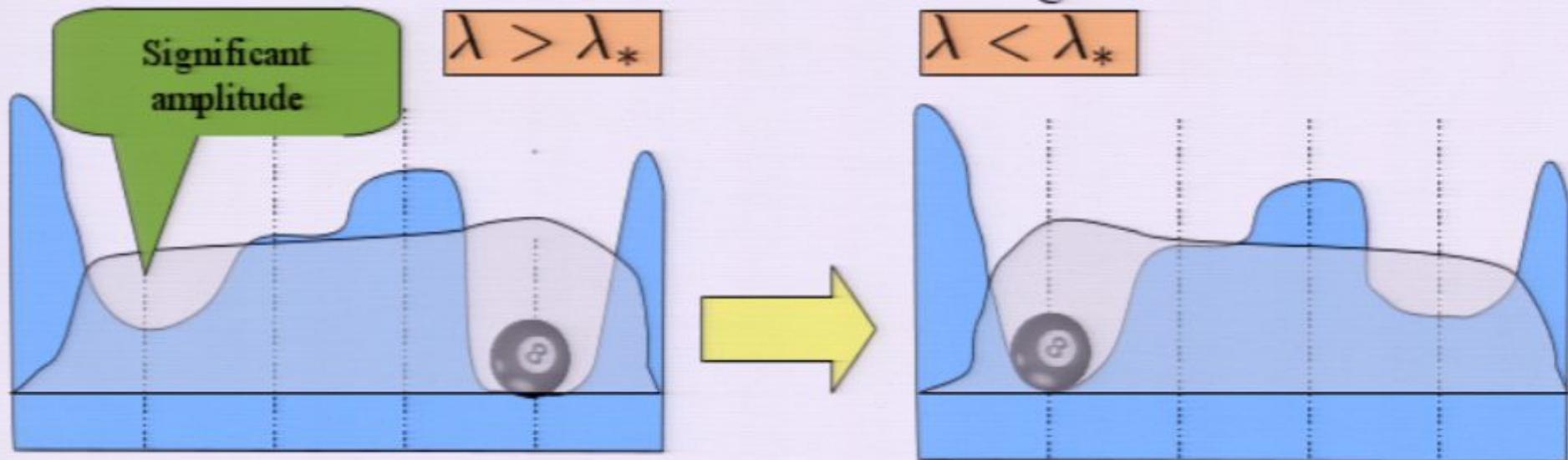
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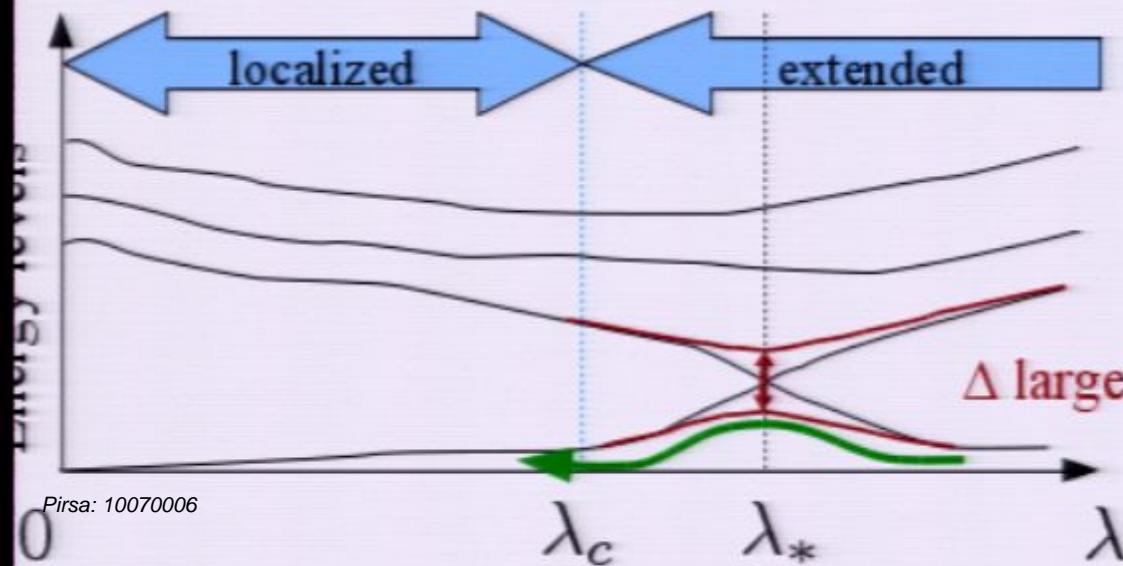
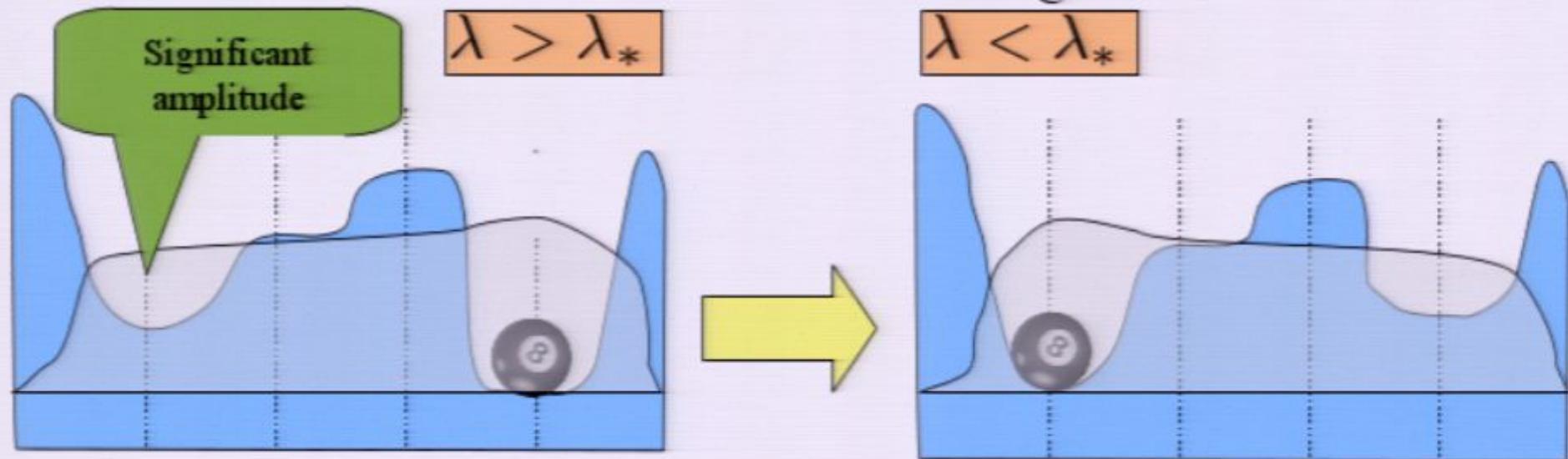
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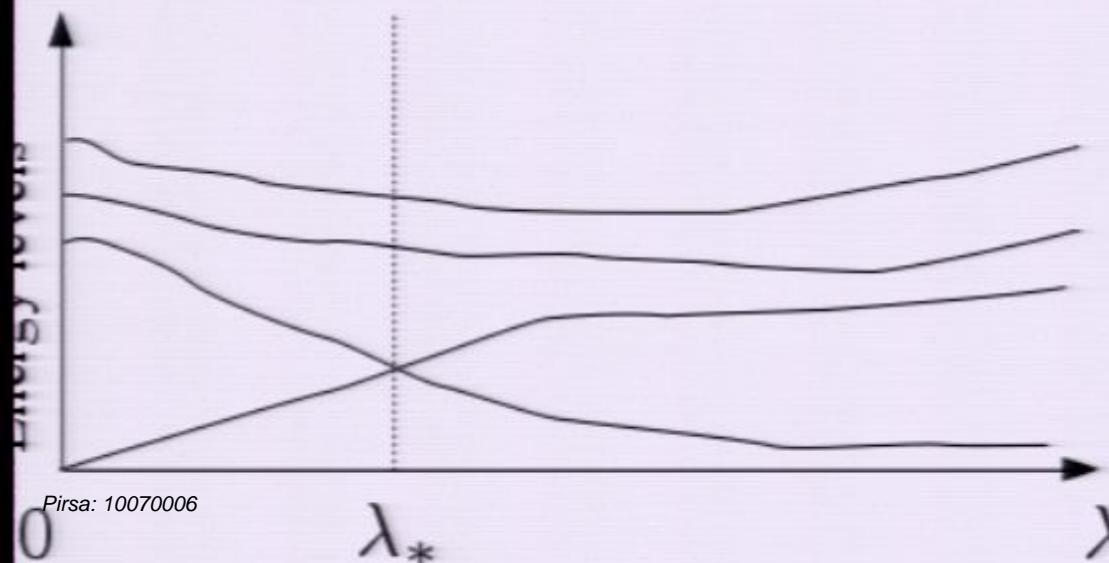
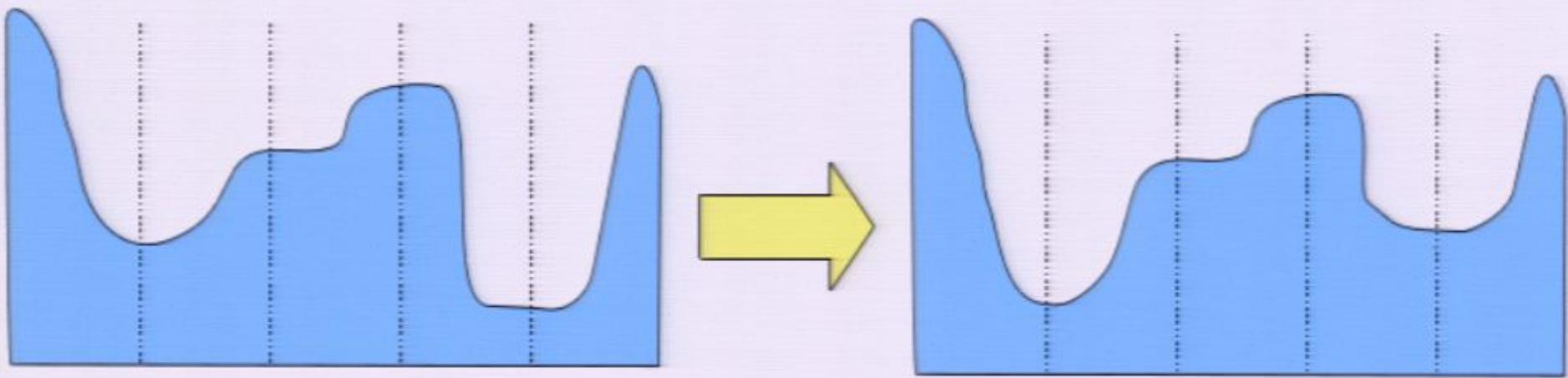


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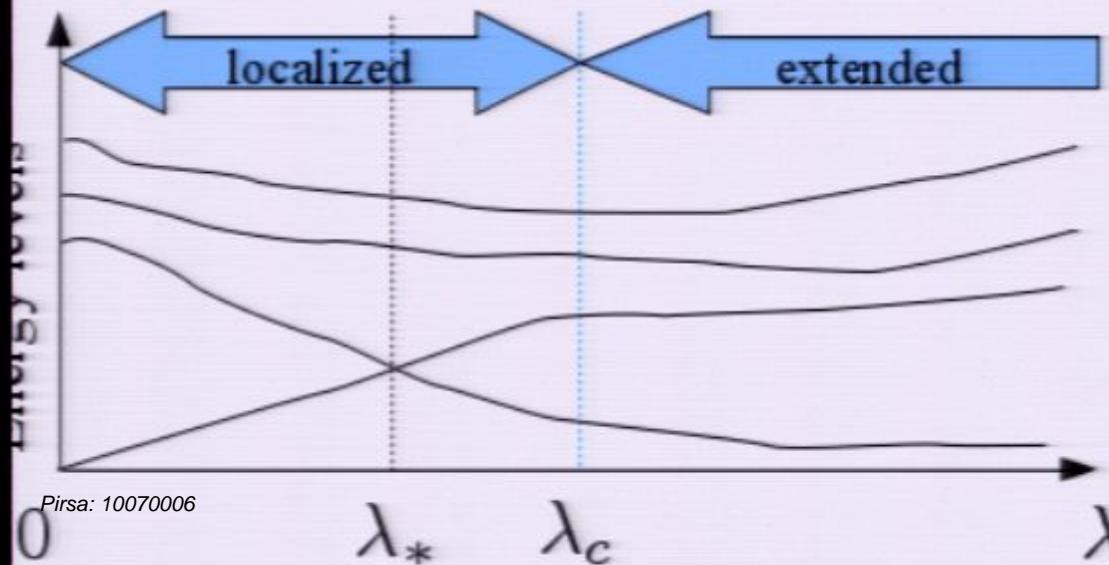
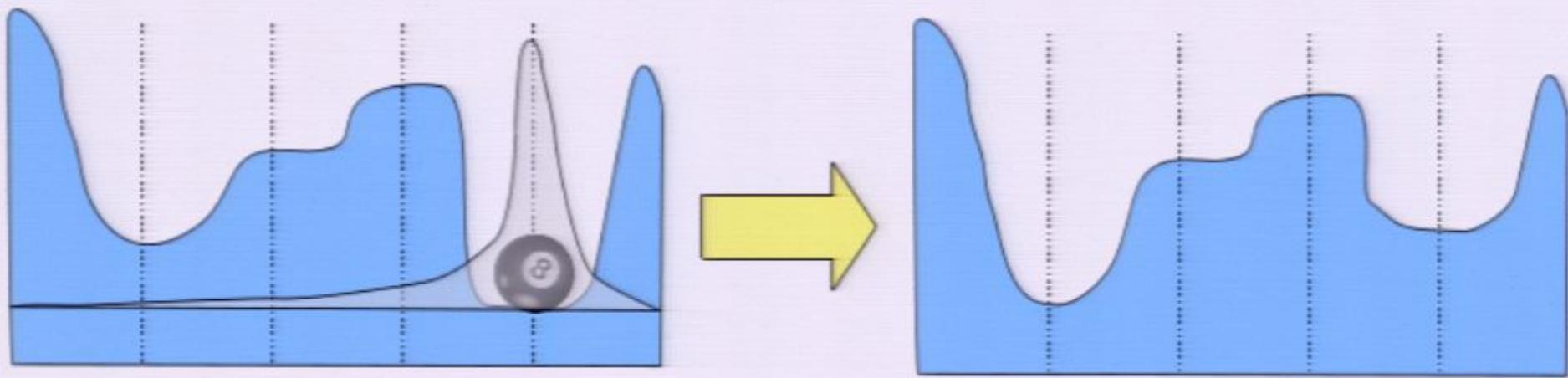


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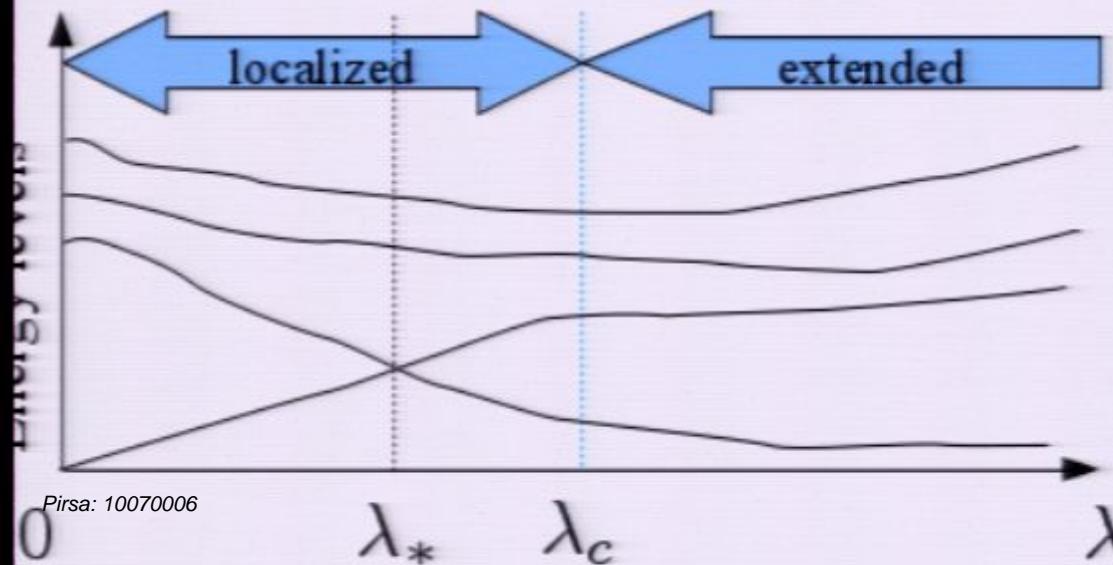
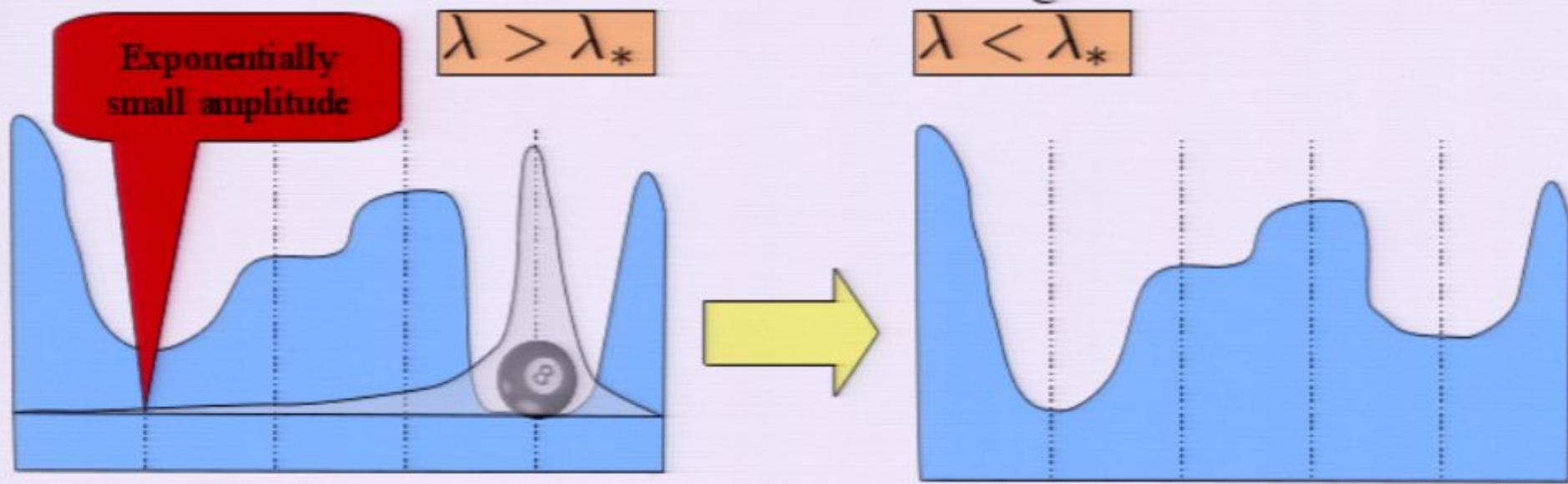
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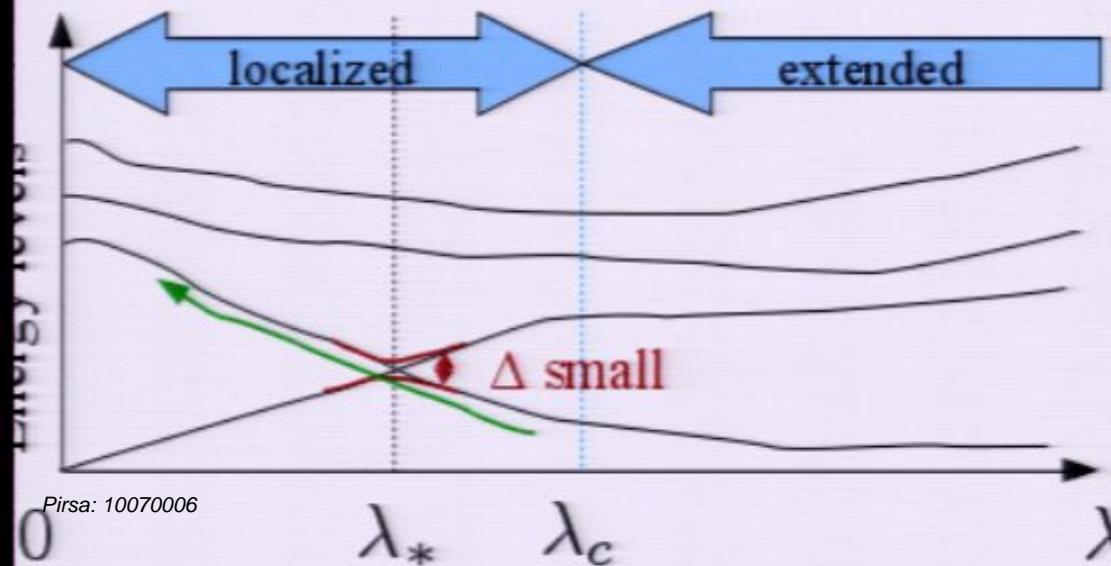
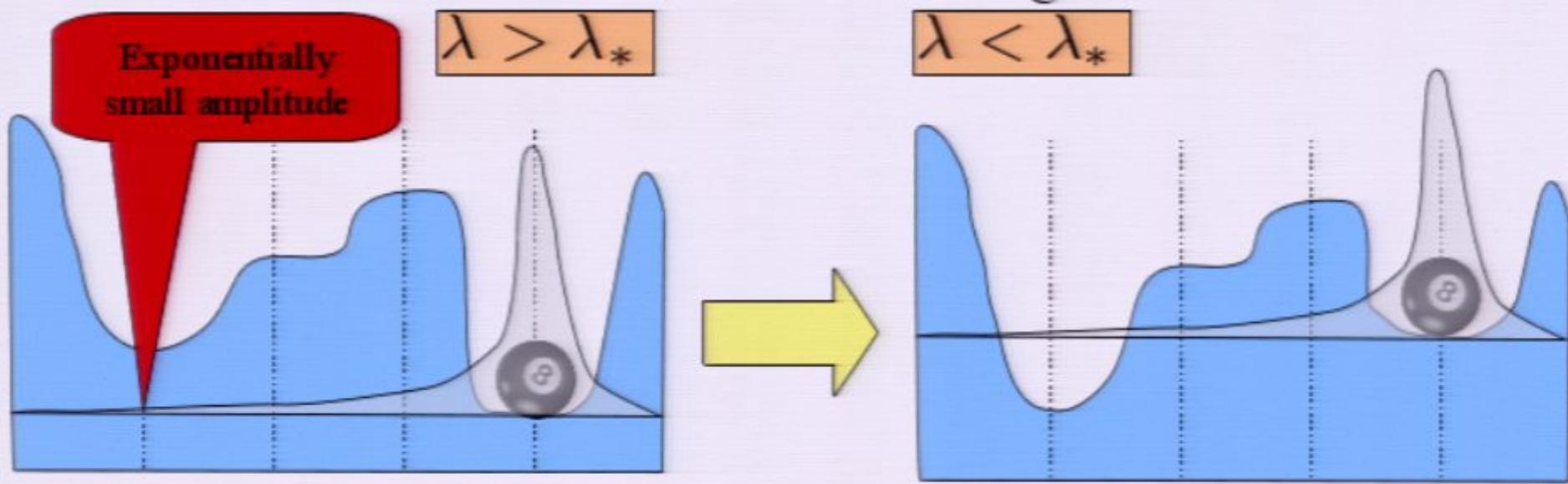
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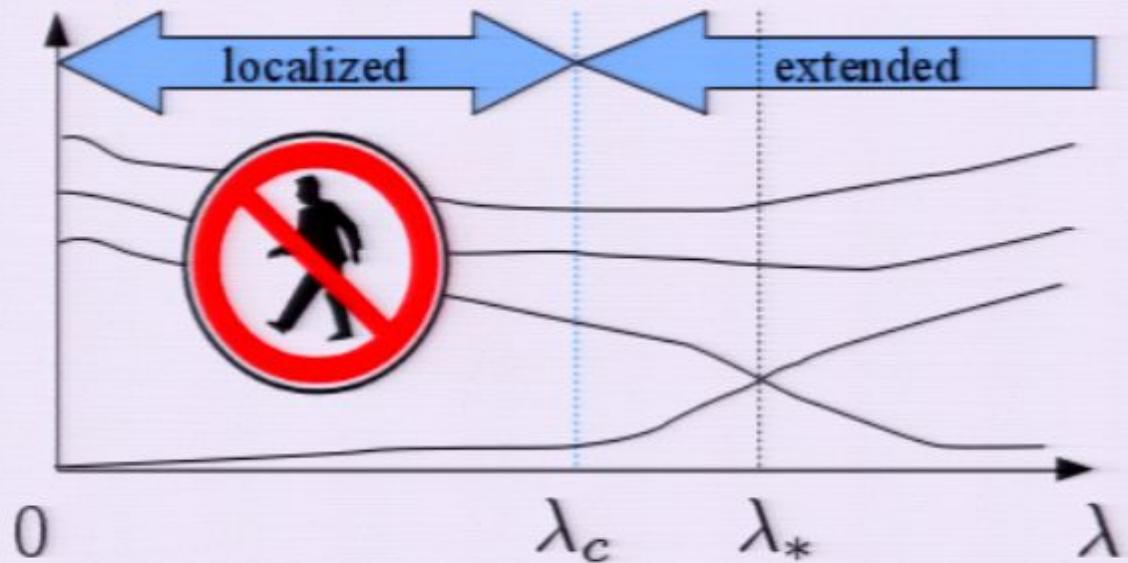
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Small anti-crossing gap
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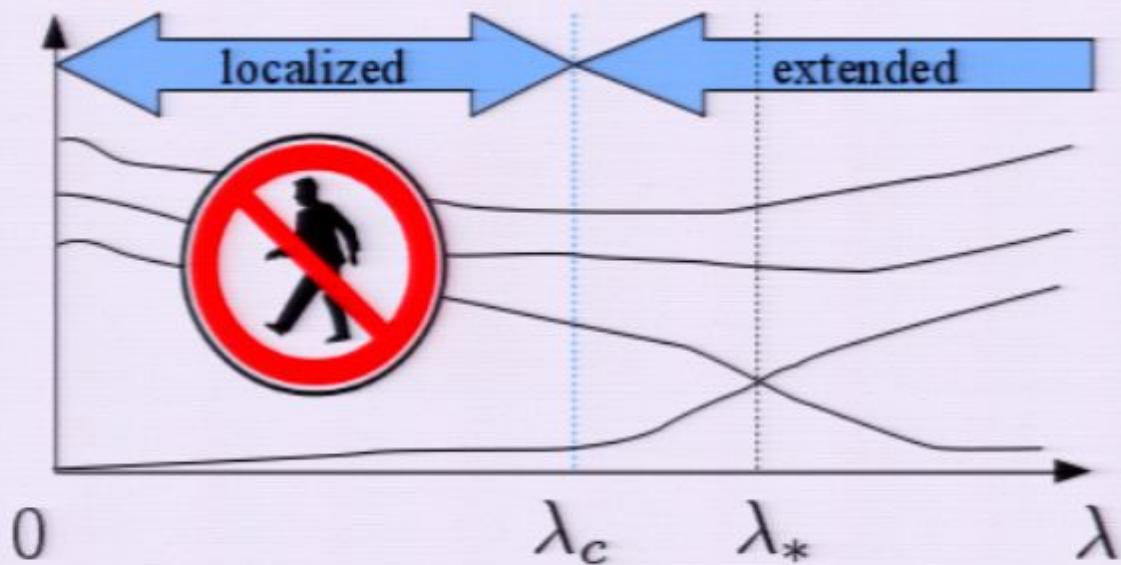


Our result



As the size of the problem N increases

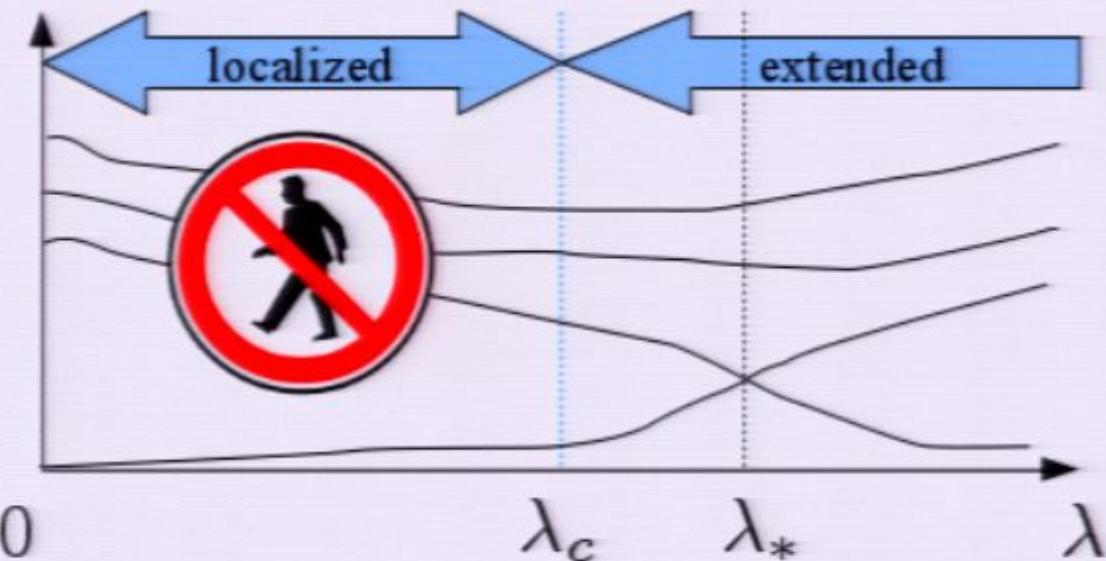
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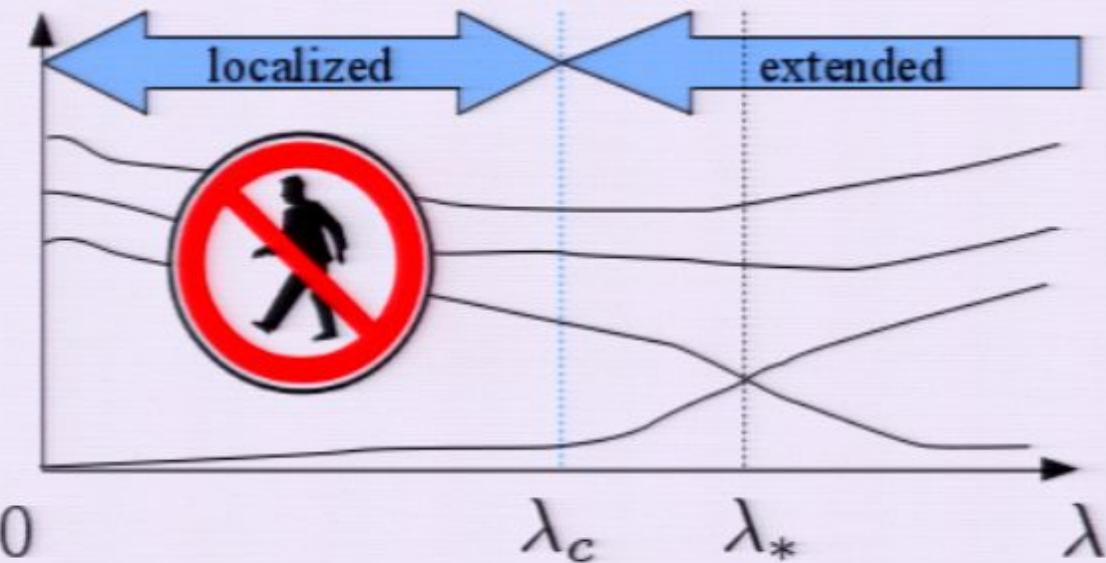
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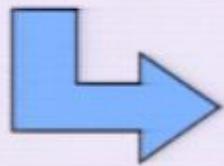
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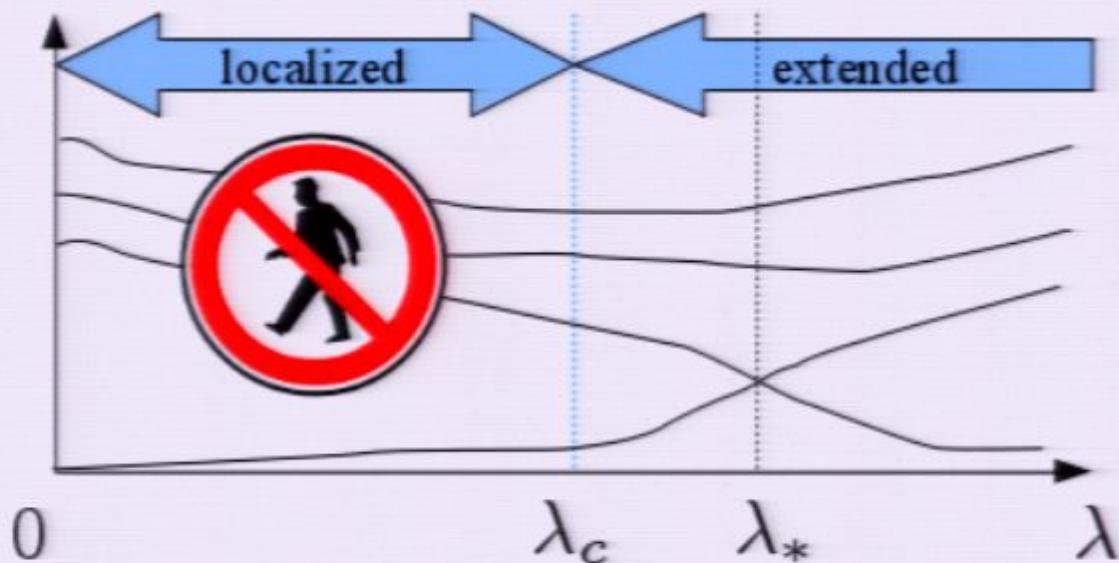
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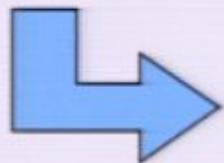
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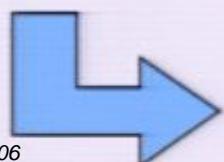


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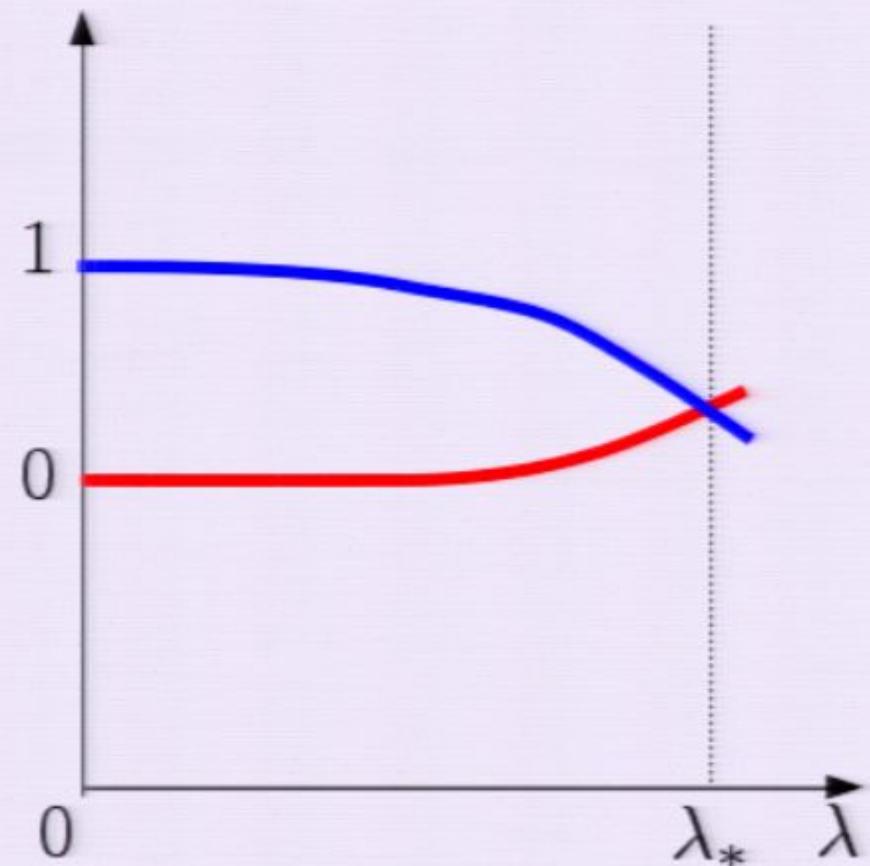
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The algorithm fails (stuck in a local minimum)

Level anti-crossings

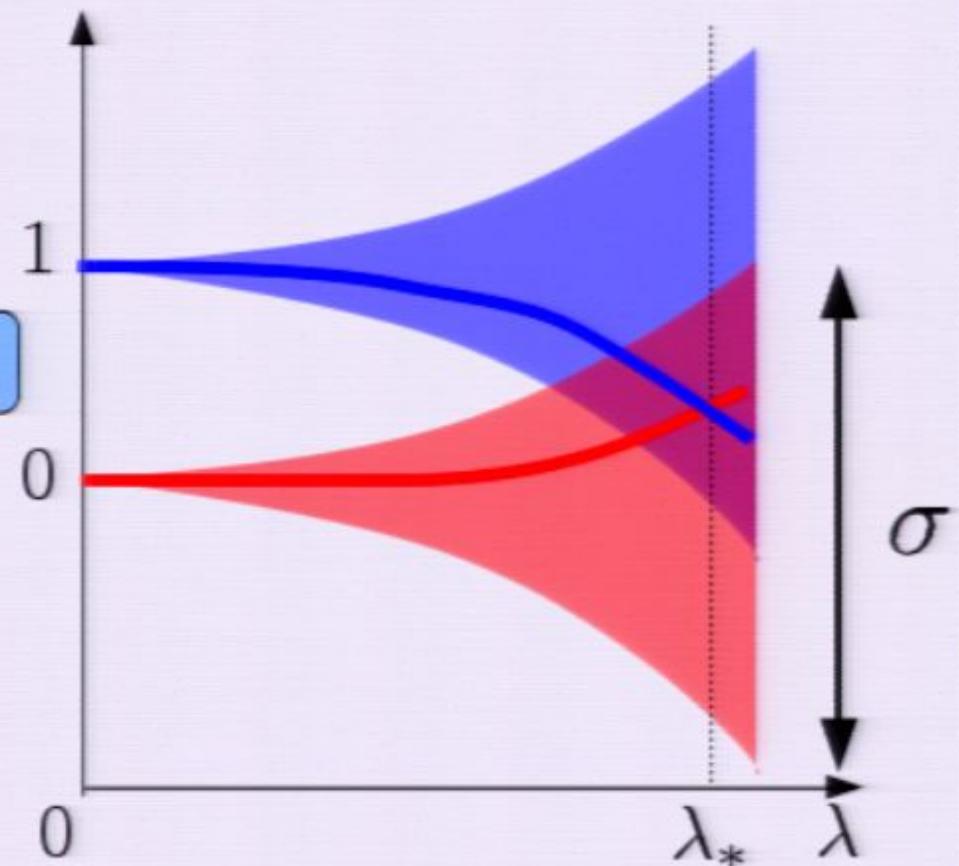
Study slopes by perturbation theory



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Random instances \Rightarrow random slopes!



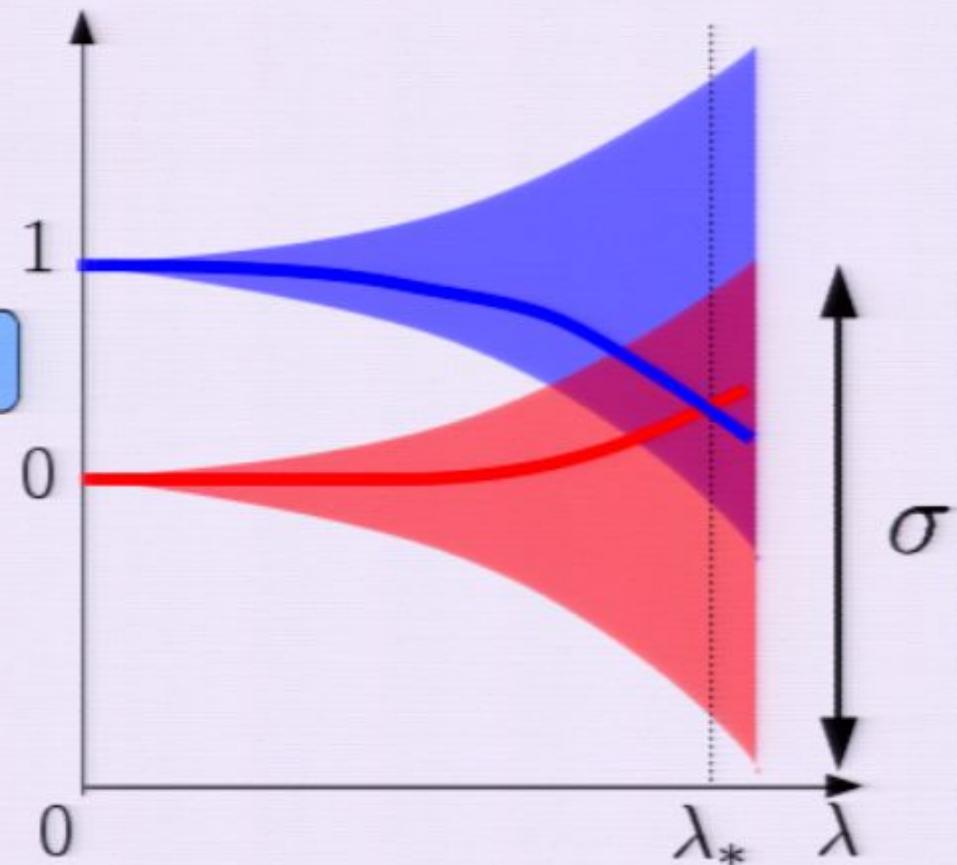
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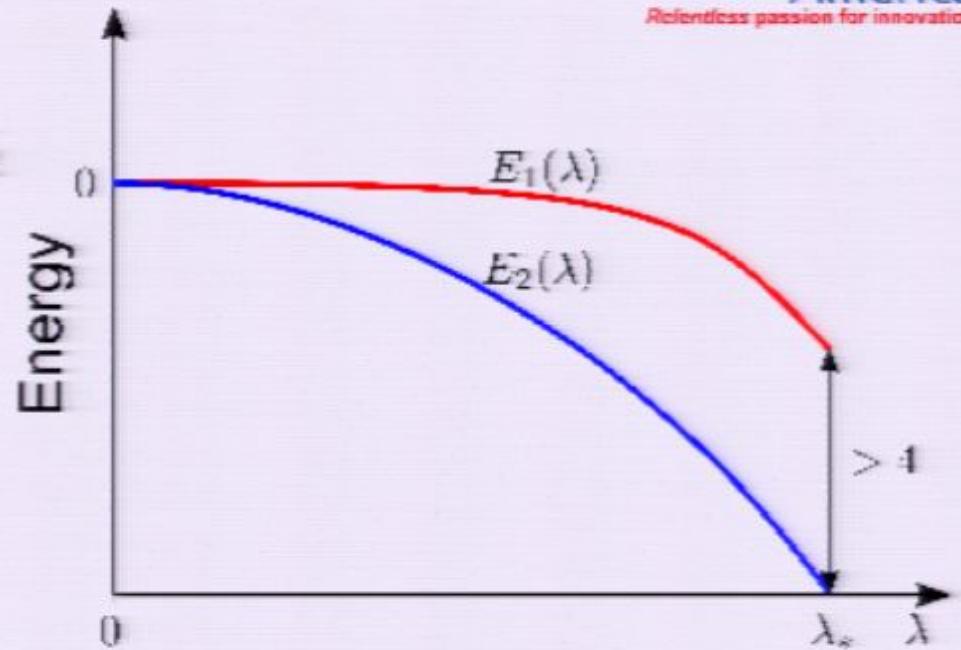
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$$E_1(0) = E_2(0) = 0$$

Suppose

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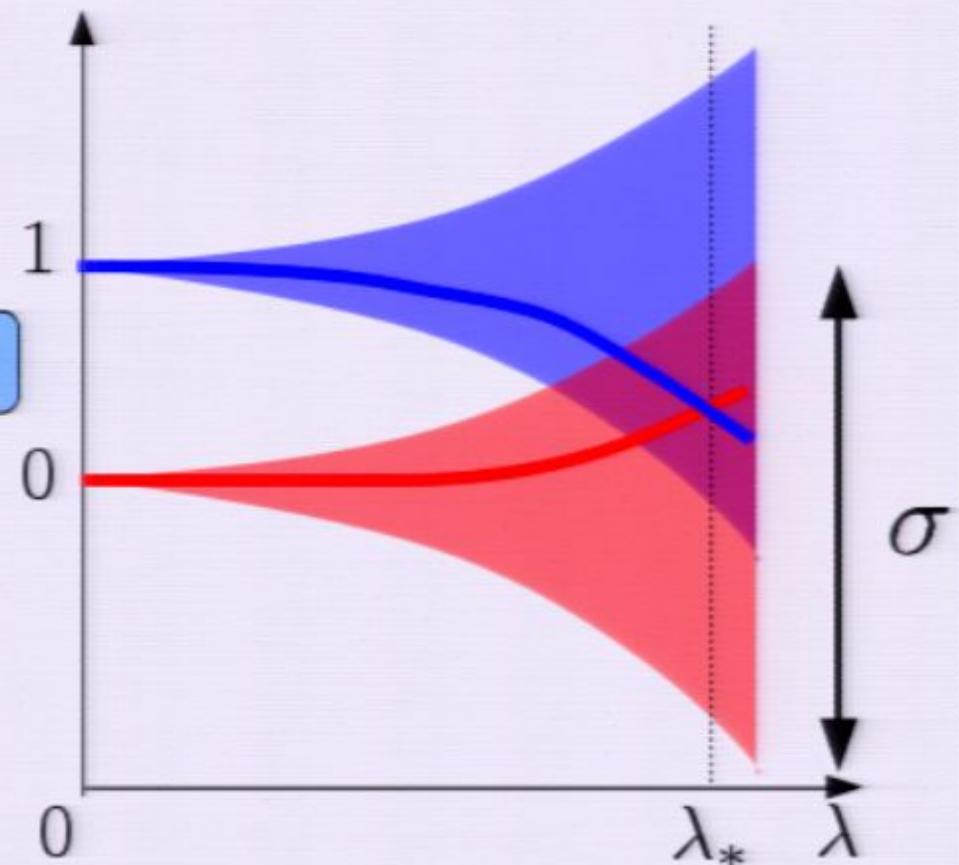
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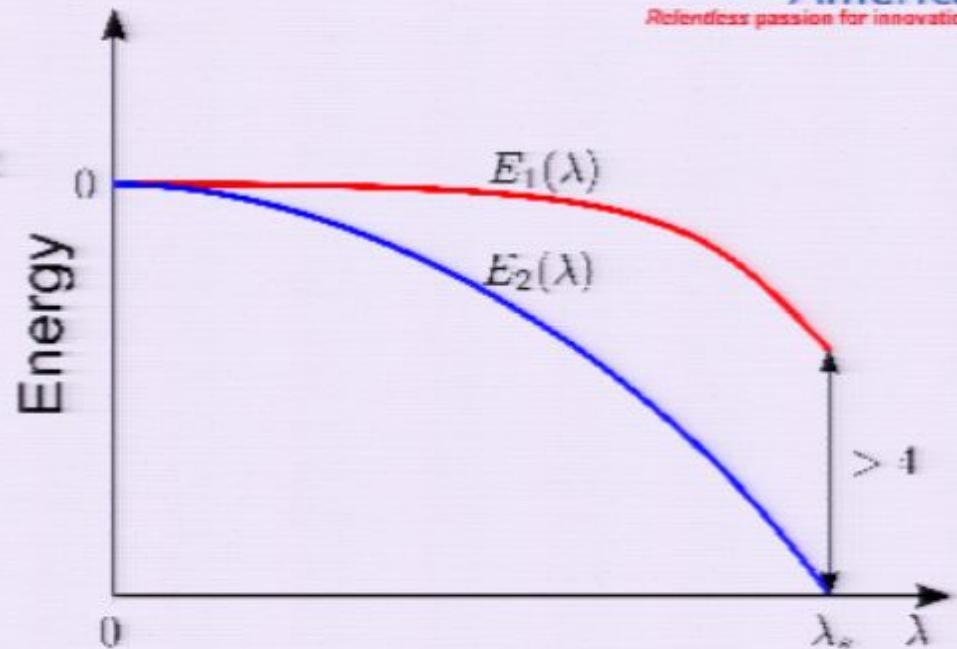
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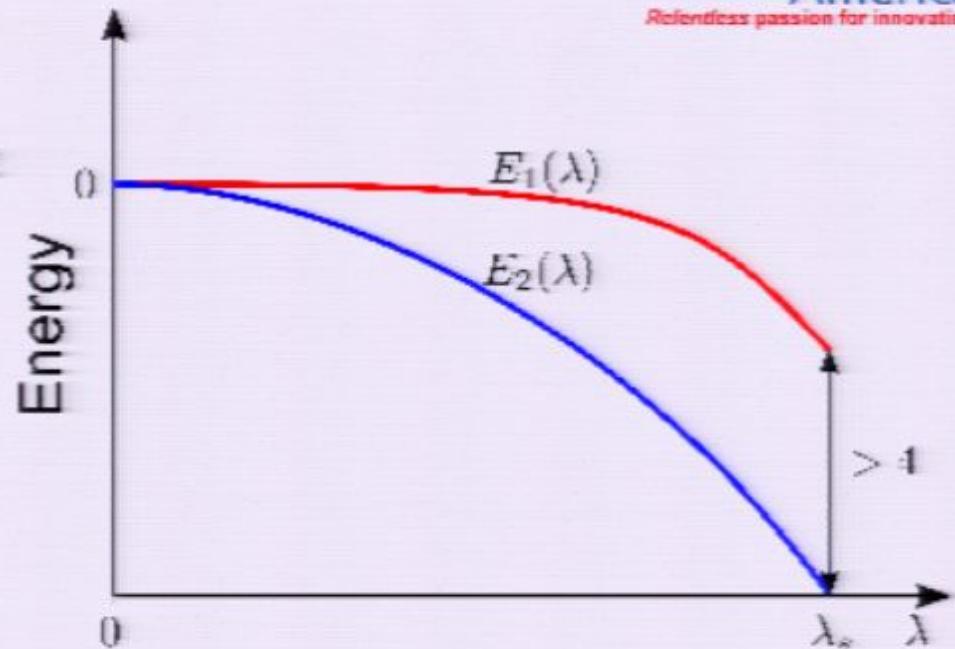
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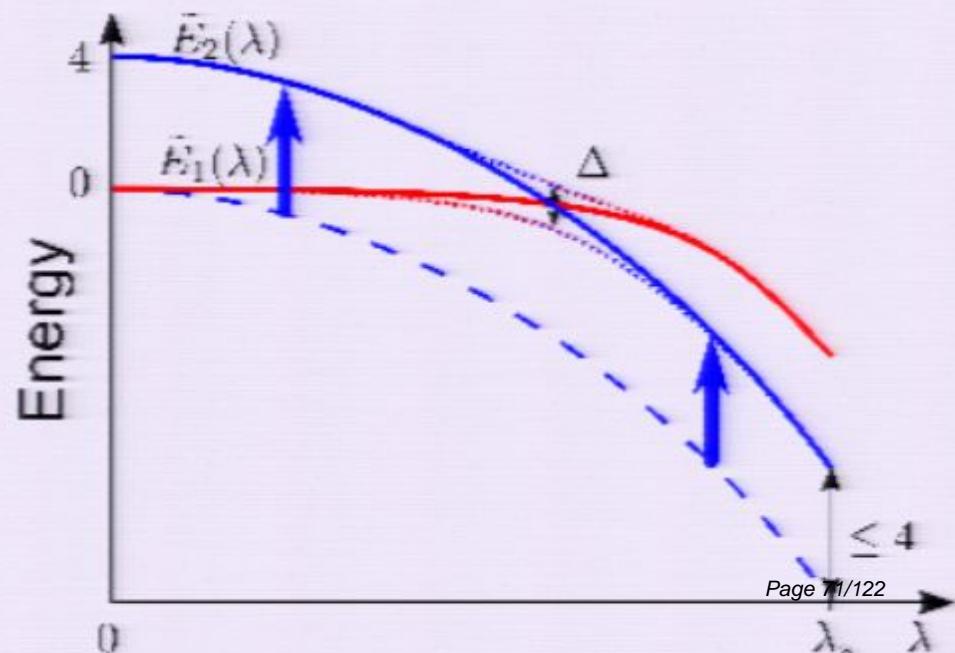
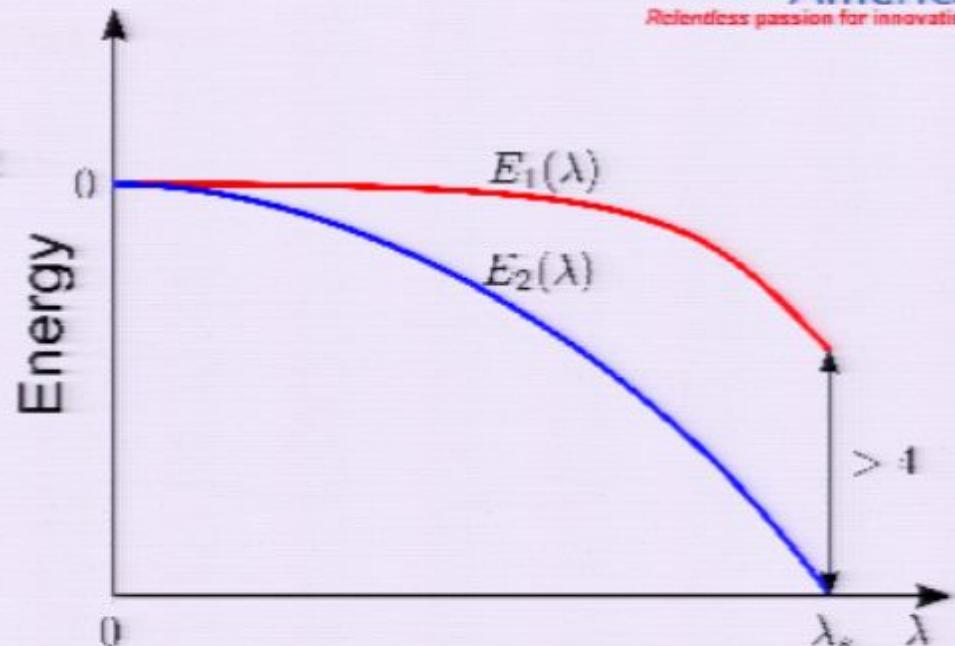
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anti-crossing



Perturbation theory

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$$E_{\vec{x}}(\lambda) = E_{\vec{x}}(0) + \sum_{m=1}^{\infty} \lambda^{2m} F_{\vec{x}}^{(m)}$$

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$$F_{\vec{x}}^{(m)} = O(N) \quad \forall m$$

Proof based on statistical properties of random instances

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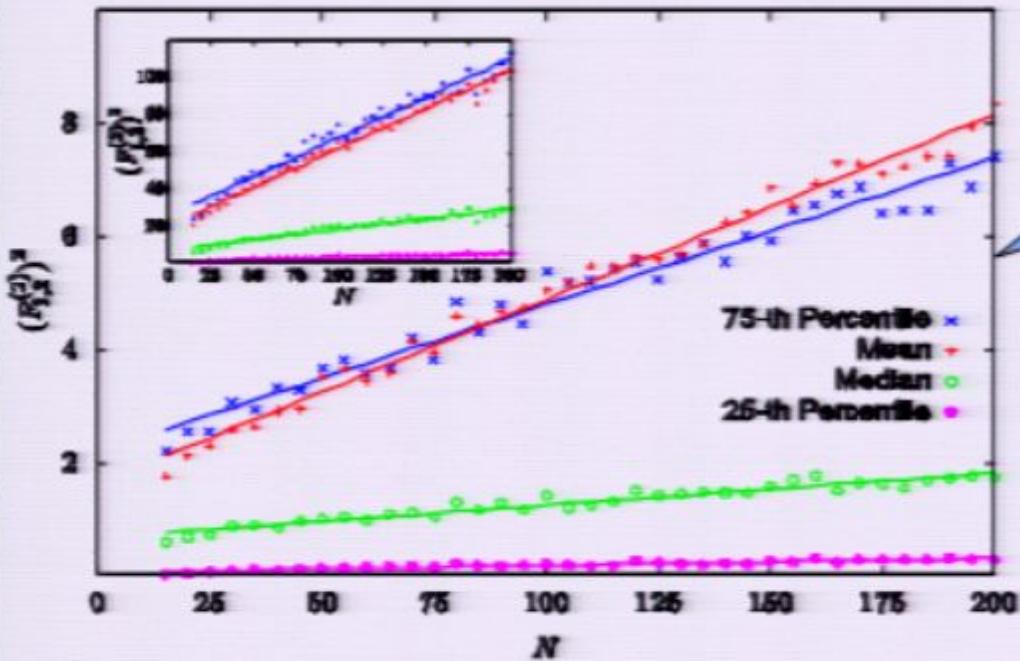
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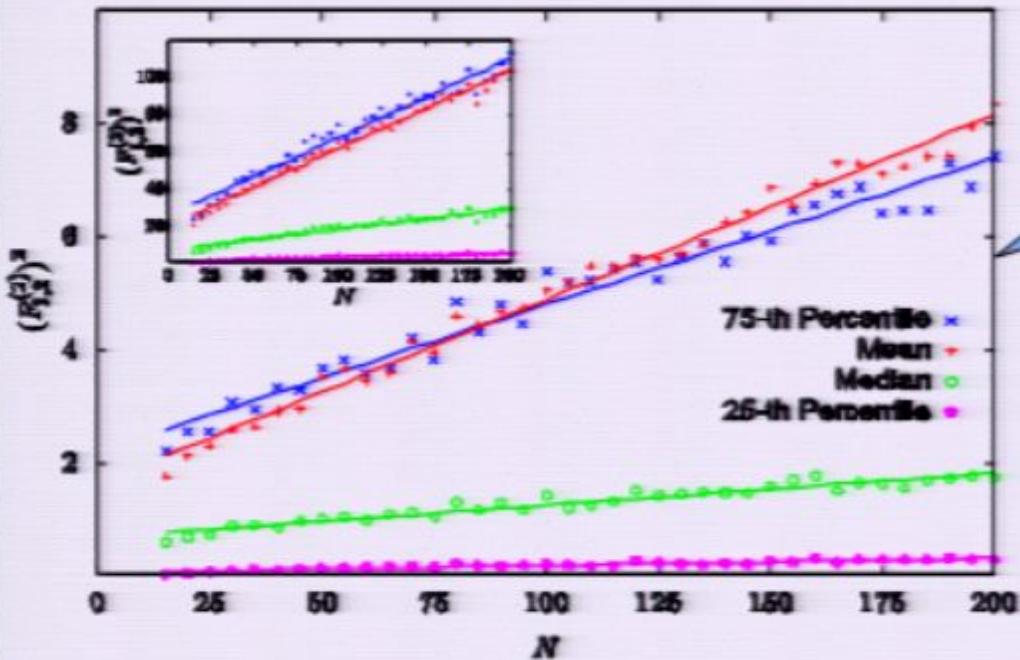


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We show that up to leading order in perturbation theory:

$$\Delta < (2\lambda_*)^n$$

Proof by reduction to the “Agree” problem:
2-bit clauses $(x_{i_C} = x_{j_C})$

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We have: $\boxed{\Delta = O(\exp(-N \log N))}$

Can we trust perturbation theory?

Anderson localization theory

⇒ Perturbation theory valid as long as states are localized

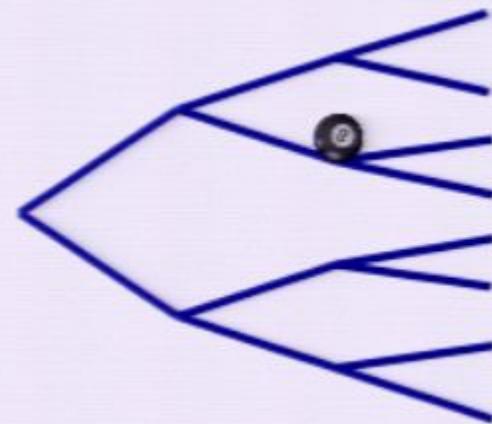
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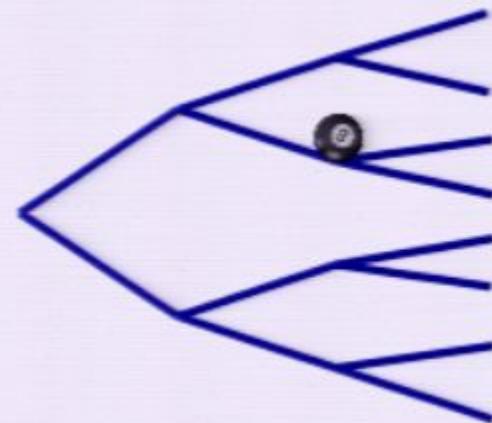
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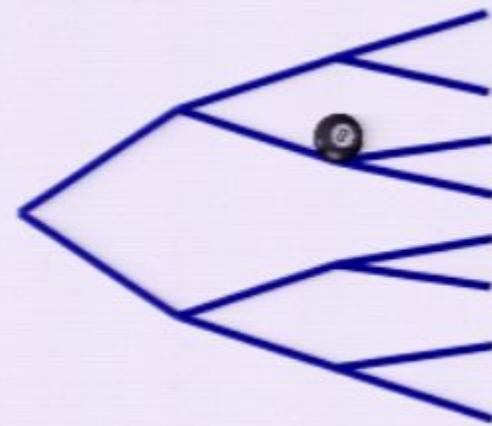
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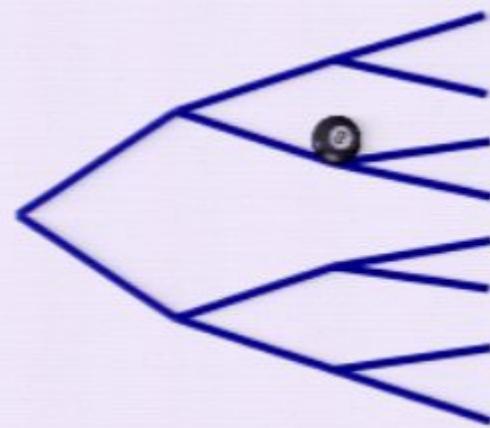
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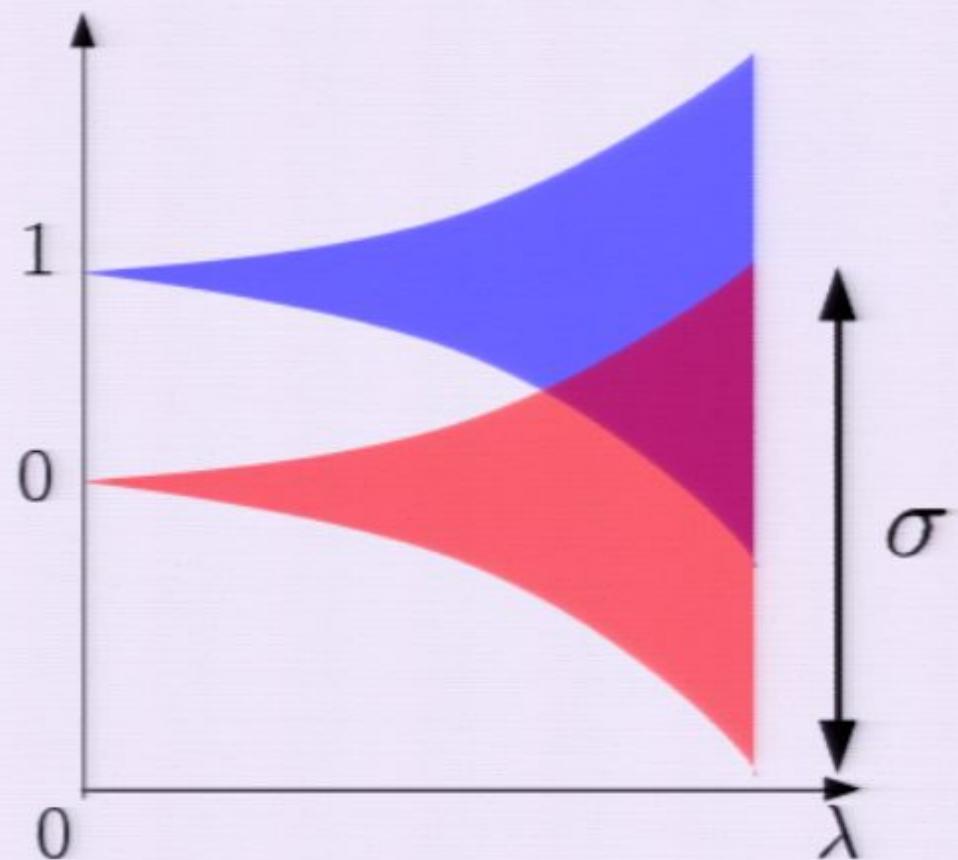
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Our estimation:

$$\sigma \sim \sqrt{N} \lambda^4$$



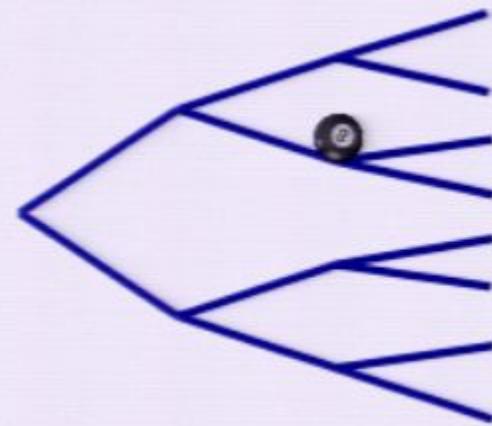
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Degeneracy of the ground state

- Our estimation: $\sigma \sim \sqrt{N} \lambda^4$

- Ground state is degenerate

$$S_1 \sim e^{\eta N} \Rightarrow \sigma_1 \sim \lambda^4$$

[Knysh-Smelyanskiy'10]

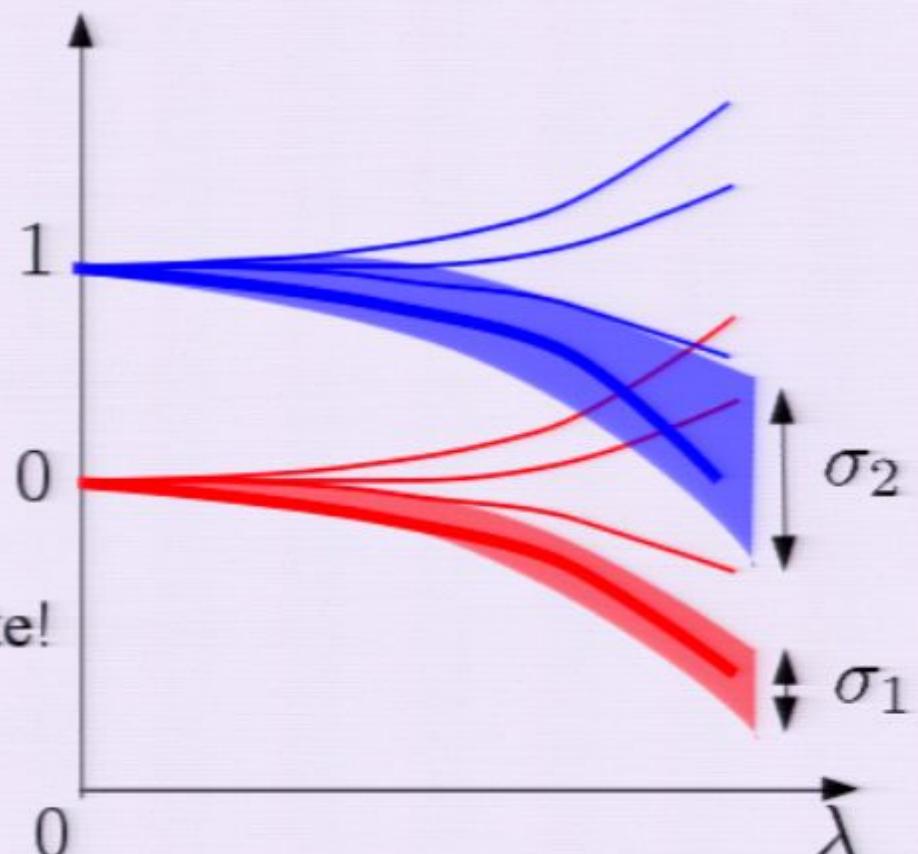
BUT

- First excited state is more degenerate!

$$S_2 \sim N e^{\eta N} \Rightarrow \sigma_2 \sim \frac{\log N}{\sqrt{\eta}} \lambda^4$$

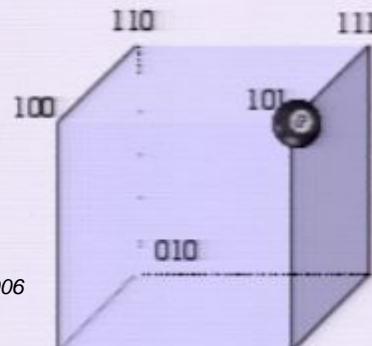
- Also: $\eta \rightarrow 0$ as $\alpha \rightarrow \alpha_s$

\Rightarrow Effect of degeneracy only appears for large N



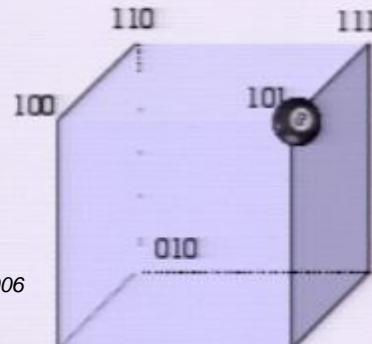
Conclusion

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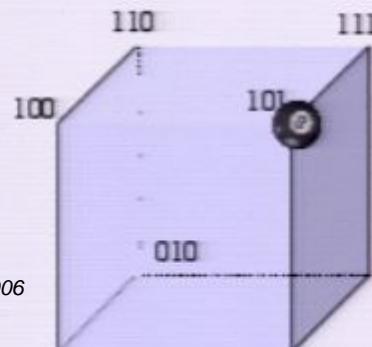
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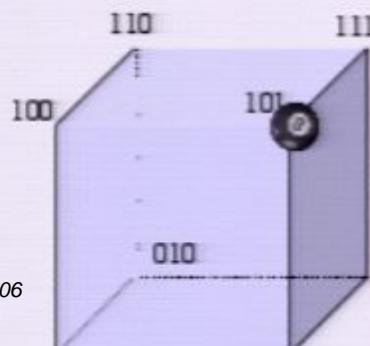
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- Important assumption: Localization on the hypercube



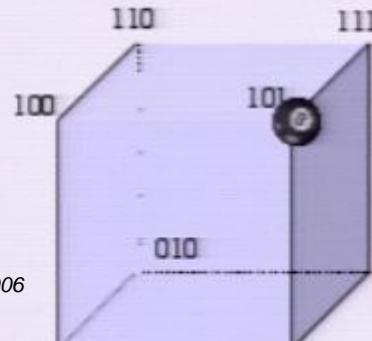
⇒ should be studied more closely

Thank you!

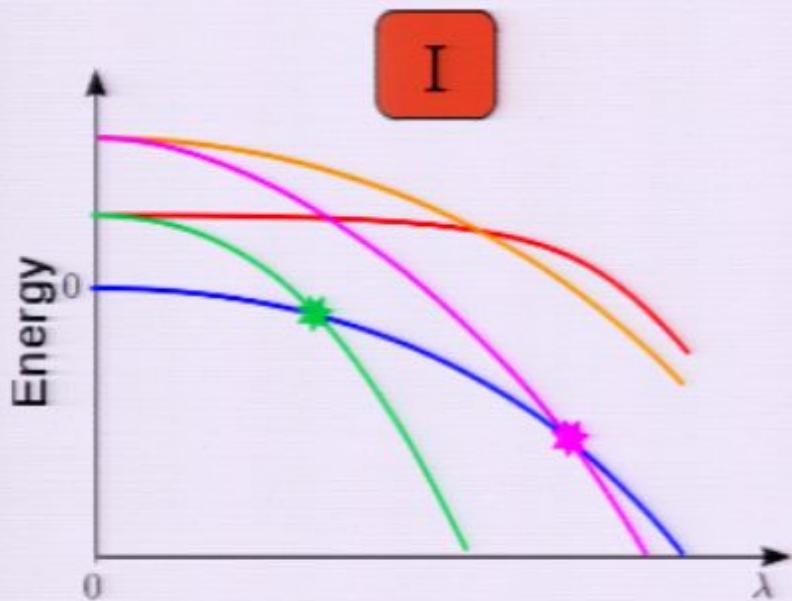


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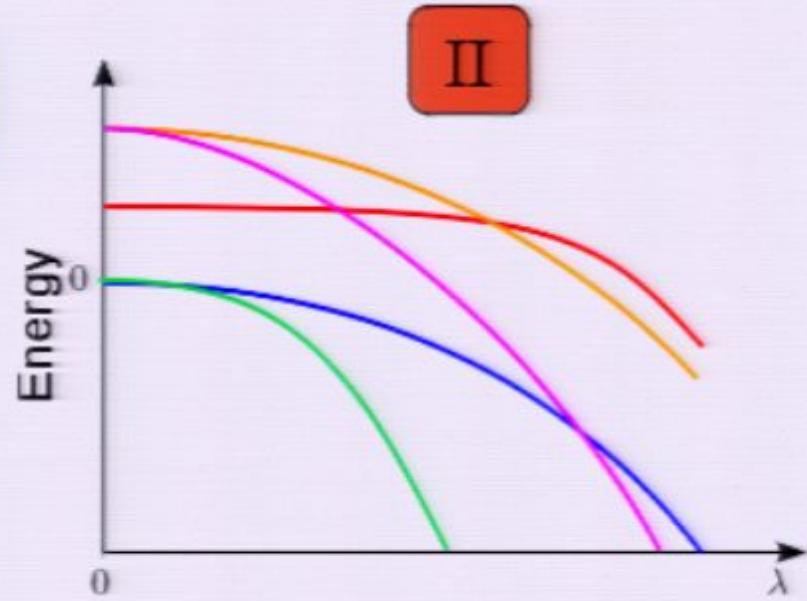
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Effect of path change



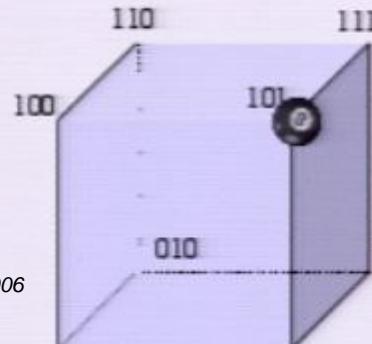
OR



- Idea: Pick random $H(s)$ ("path change") to obtain case II
[Farhi *et al.* '09]
- Avoid 1 crossing: $Pr[\text{"II"}] = \text{constant}$
- Avoid poly # of crossings: $Pr[\text{"II"}] = 1/\text{poly}$
- Estimated # of crossings: $\exp(N/\log^7 N)$

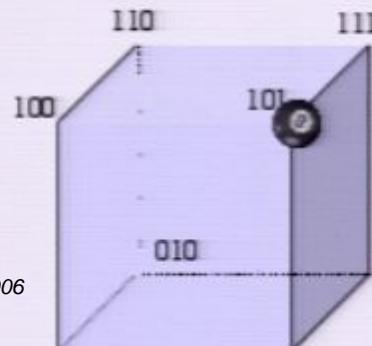
Conclusion

- Anderson localization causes exponentially small gaps in adiabatic quantum optimization

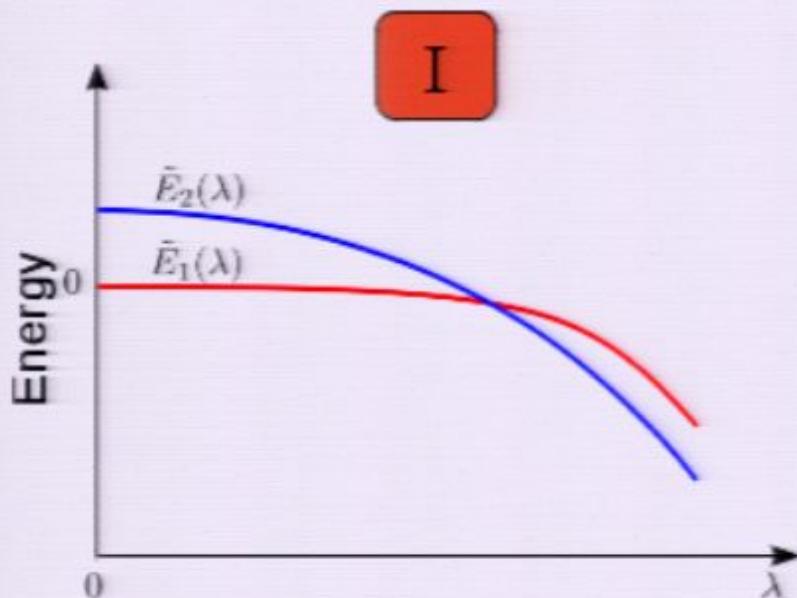


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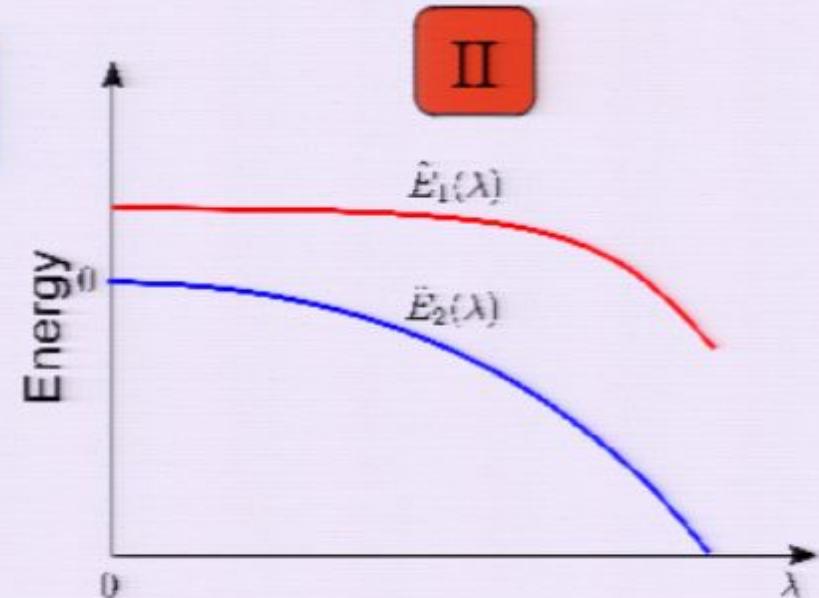
- Anderson localization causes exponentially small gaps in adiabatic quantum optimization



Effect of path change



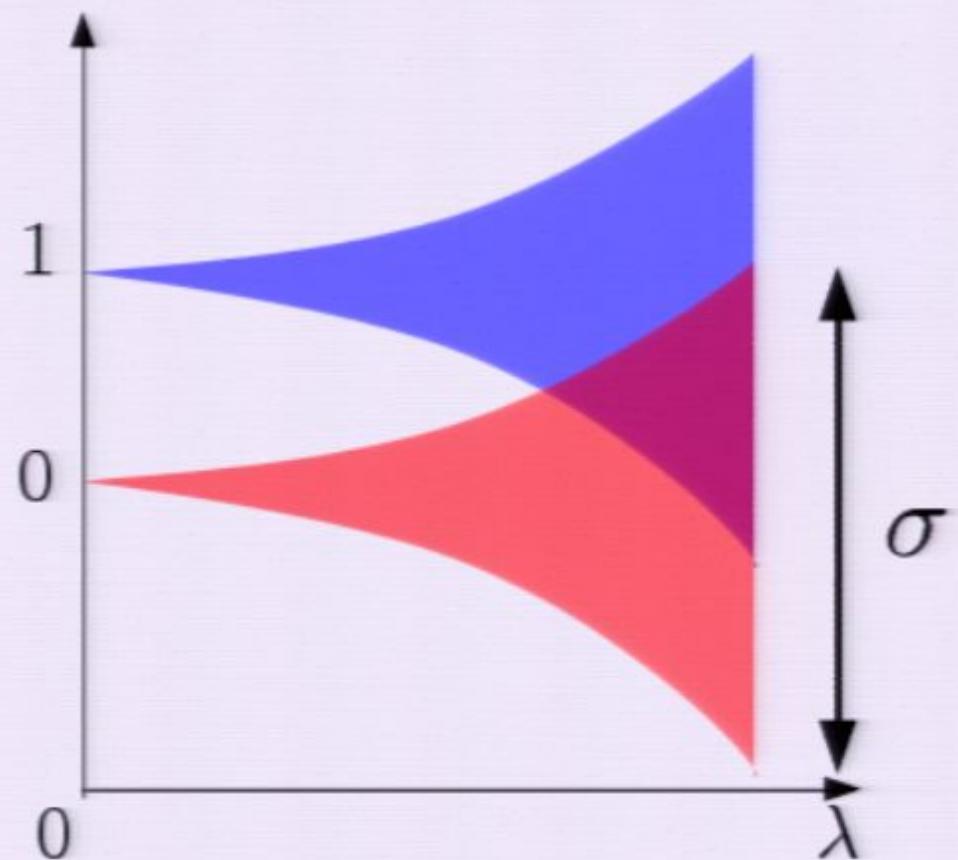
OR



- Idea: Pick random $H(s)$ (“path change”) to obtain case II
[Farhi *et al.* '09]
- Avoid 1 crossing: $Pr[\text{“II”}] = \text{constant}$

Degeneracy of the ground state

Our estimation: $\sigma \sim \sqrt{N}\lambda^4$



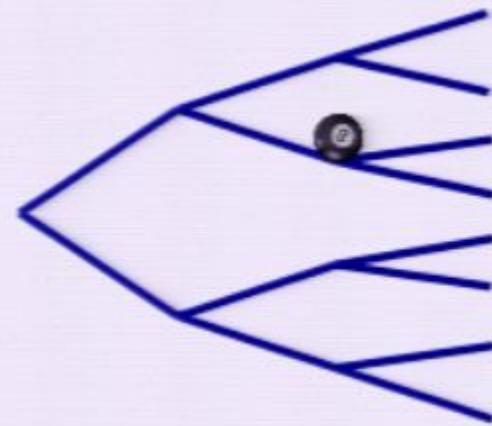
Can we trust perturbation theory?

Anderson localization theory

⇒ Perturbation theory valid as long as states are localized

Cayley tree with branching number K :

$$\lambda_c = \Theta\left(\frac{E}{K \log K}\right)$$



Here: Energy E and degree K are $\Theta(N)$, which would imply

$$\lambda_c = \Theta((\log N)^{-1}) \gg \Theta(N^{-1/8}) = \lambda_*$$

However, $F_{\vec{x}}^{(m)} = O(N) \quad \forall m$ suggests $\lambda_c = \Theta(1)$

How small is the gap?

We show that up to leading order in perturbation theory:

$$\boxed{\Delta < (2\lambda_*)^n}$$

Since: 1) level crossings appear at $\lambda_* = O(N^{-1/8})$

Can we trust perturbation theory?

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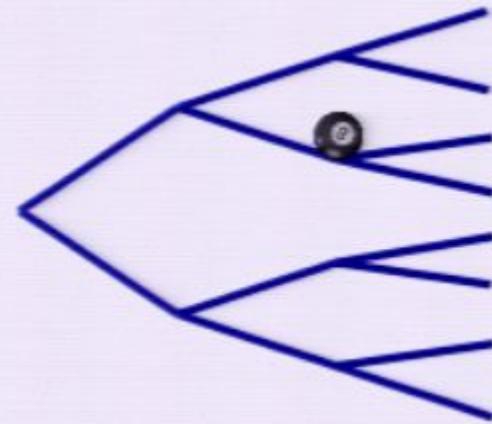
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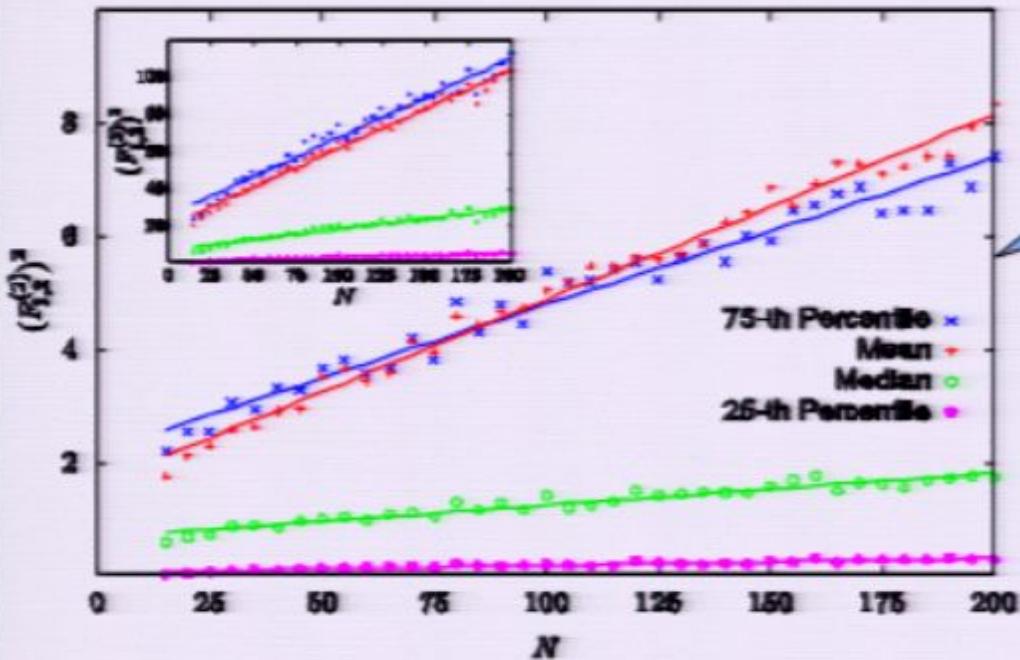
Cayley tree with branching number K :

$$\lambda_c = \Theta\left(\frac{E}{K \log K}\right)$$



How small is the gap?

- We generated EC3 random instances with >2 solutions
- then computed $E_1(\lambda) - E_2(\lambda)$ by order 4 perturbation theory



Each data point computed
from 2500 instances

$$(E_1(\lambda) - E_2(\lambda))^2 \approx CN\lambda^8$$

How small is the gap?

We show that up to leading order in perturbation theory:

$$\Delta < (2\lambda_*)^n$$

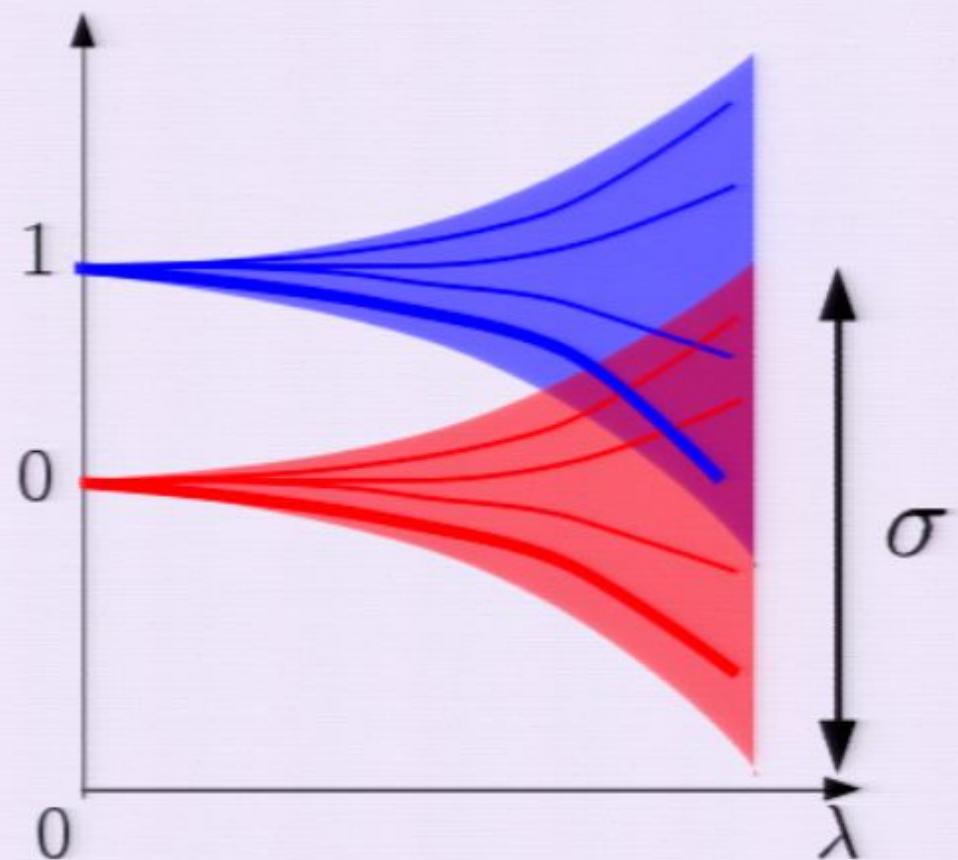
Since: 1) level crossings appear at $\lambda_* = O(N^{-1/8})$

2) typical distance between solutions is $n = \Theta(N)$

Degeneracy of the ground state

- Our estimation: $\sigma \sim \sqrt{N} \lambda^4$

- Ground state is degenerate



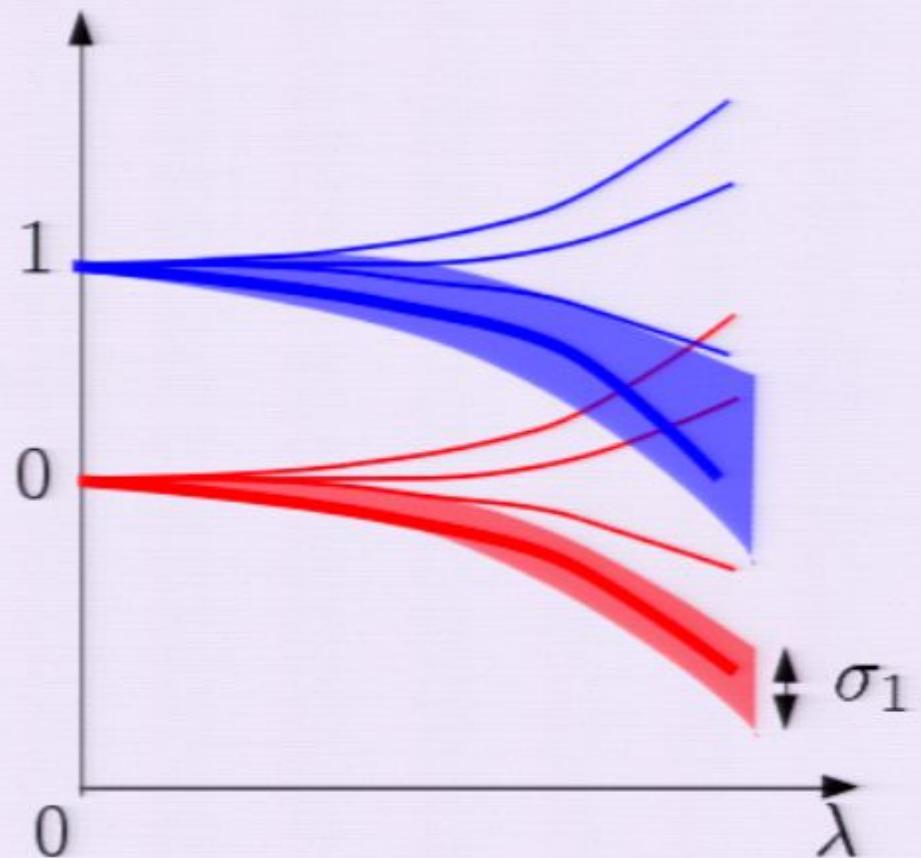
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$$S_1 \sim e^{\eta N} \Rightarrow \sigma_1 \sim \lambda^4$$

[Knysh-Smelyanskiy'10]



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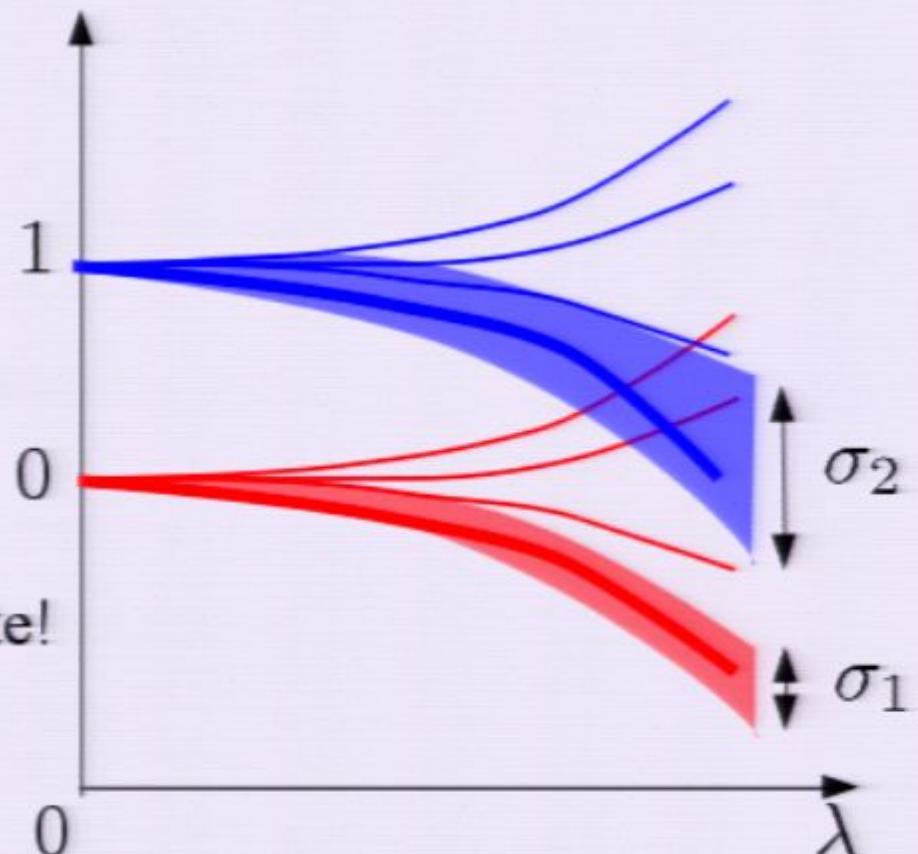
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BUT

- First excited state is more degenerate!

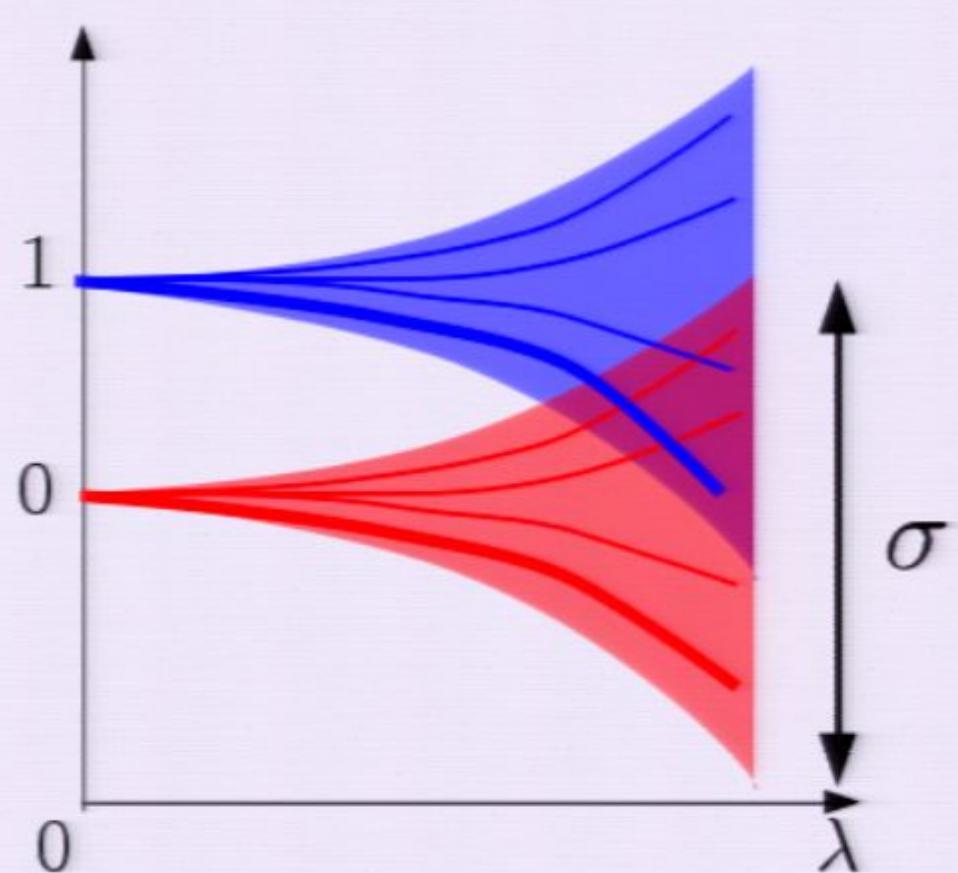
$$S_2 \sim N e^{\eta N} \Rightarrow \sigma_2 \sim \frac{\log N}{\sqrt{\eta}} \lambda^4$$



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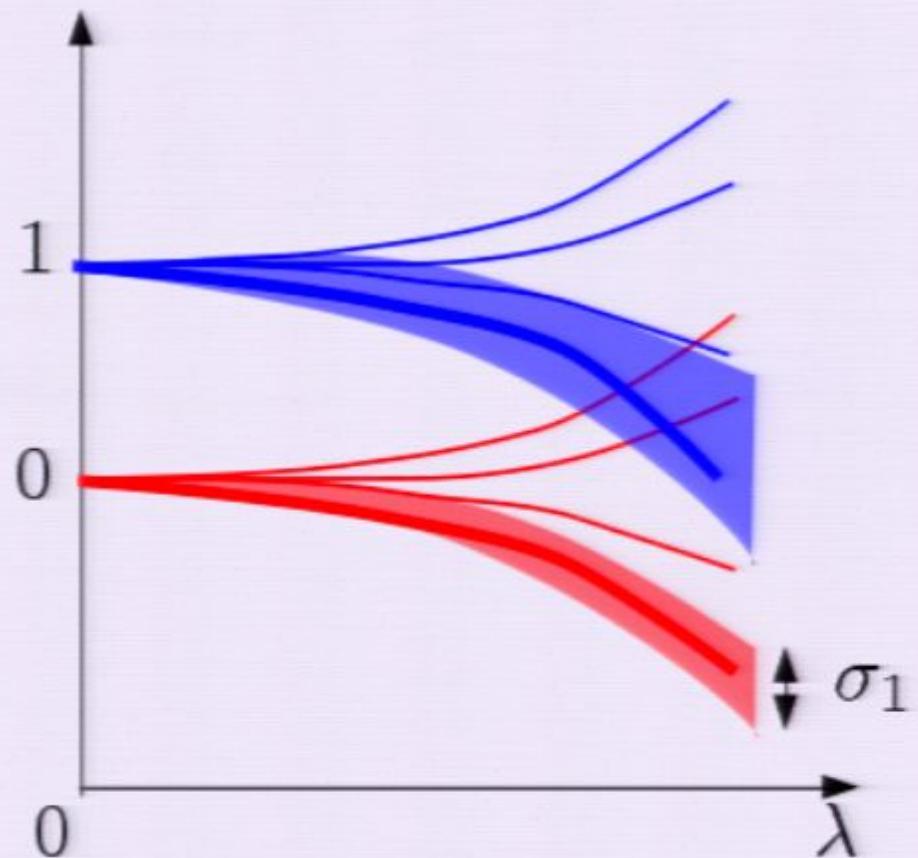
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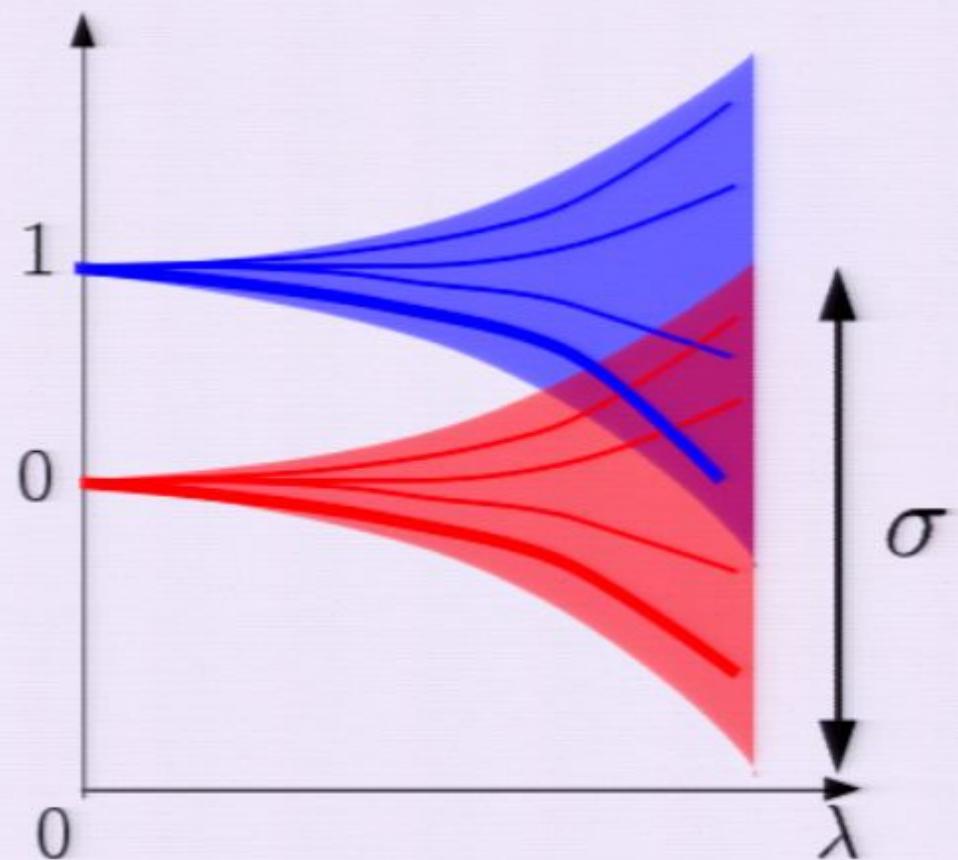
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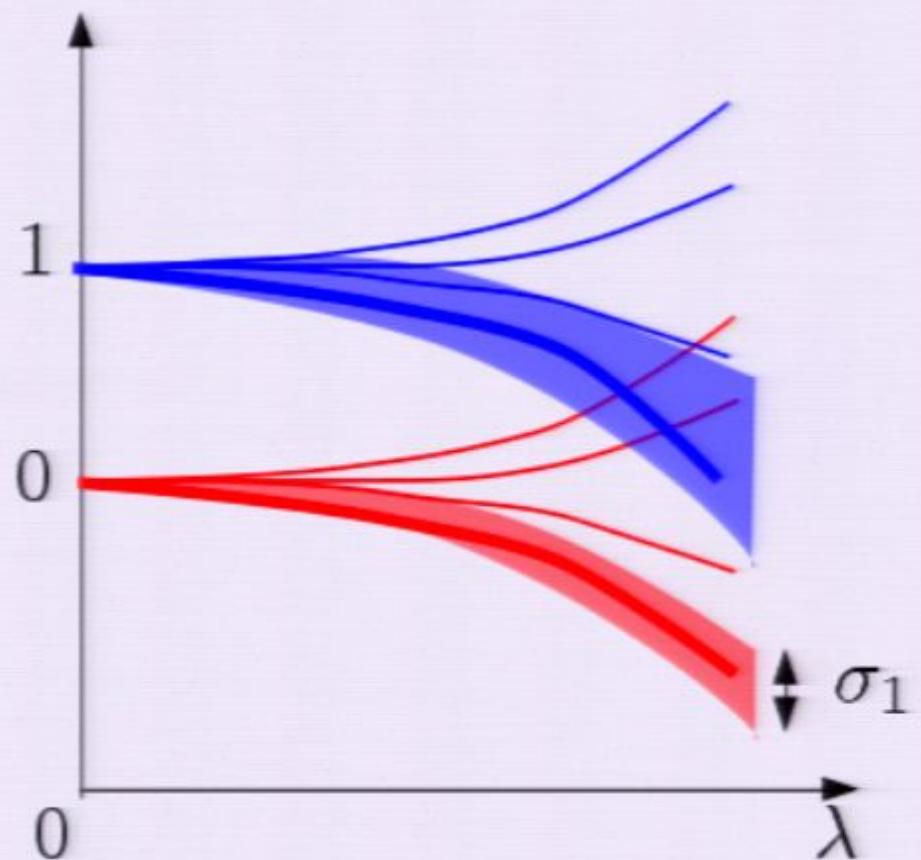
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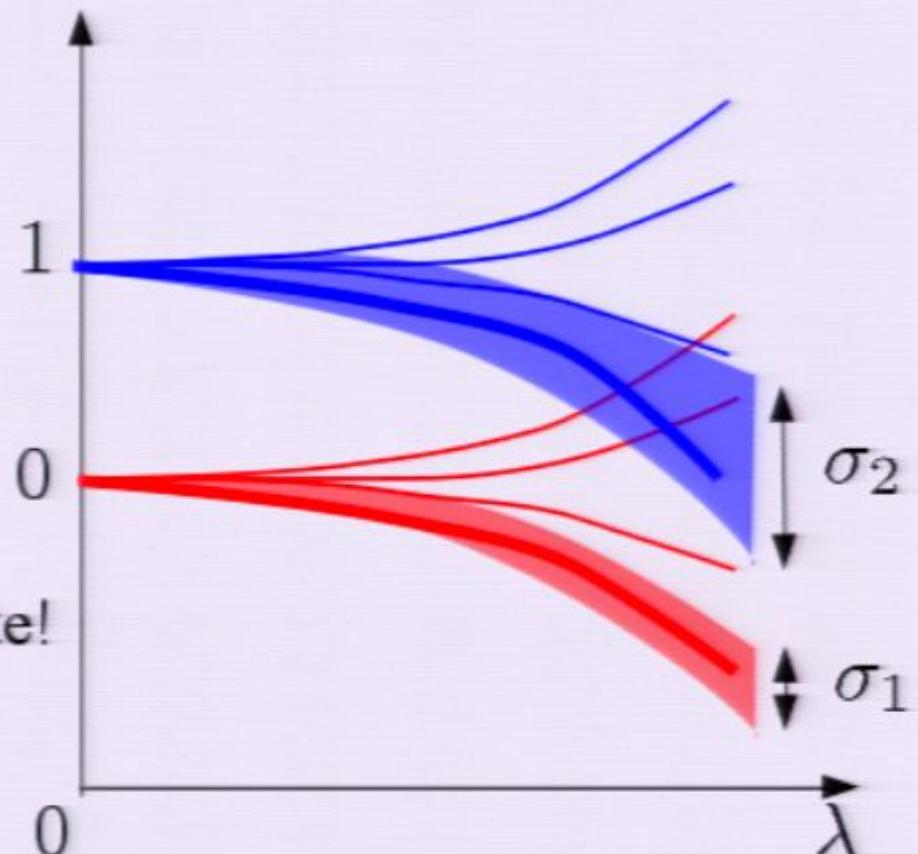
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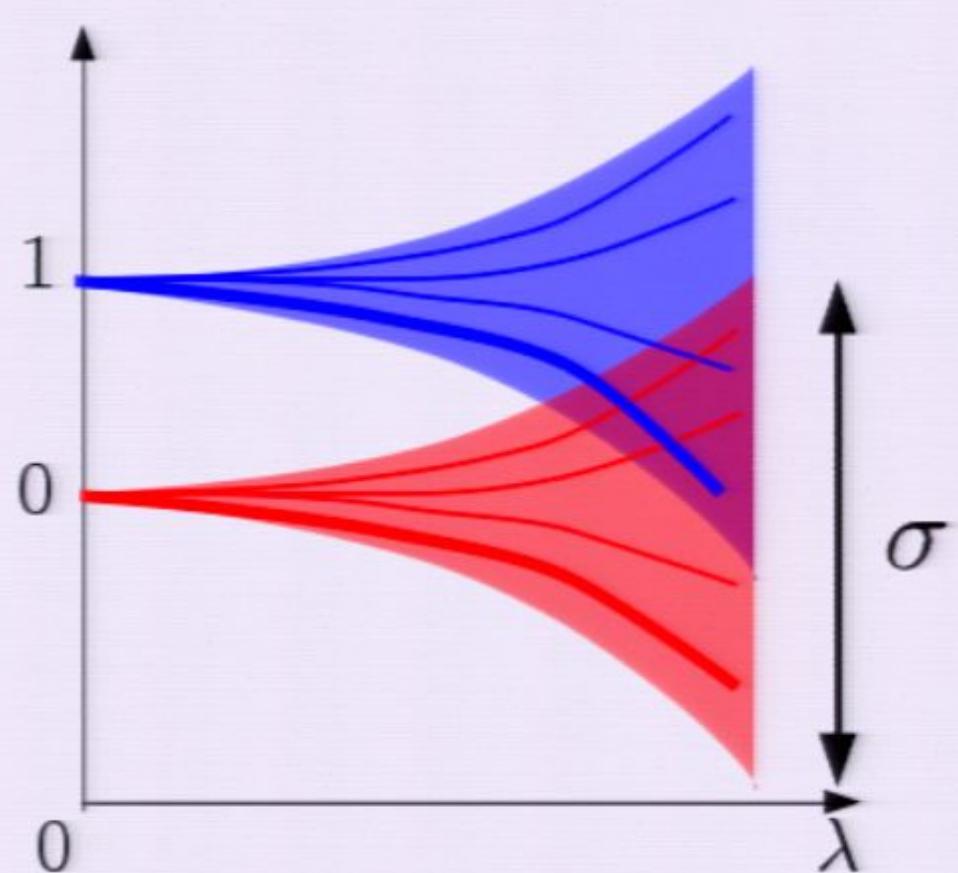
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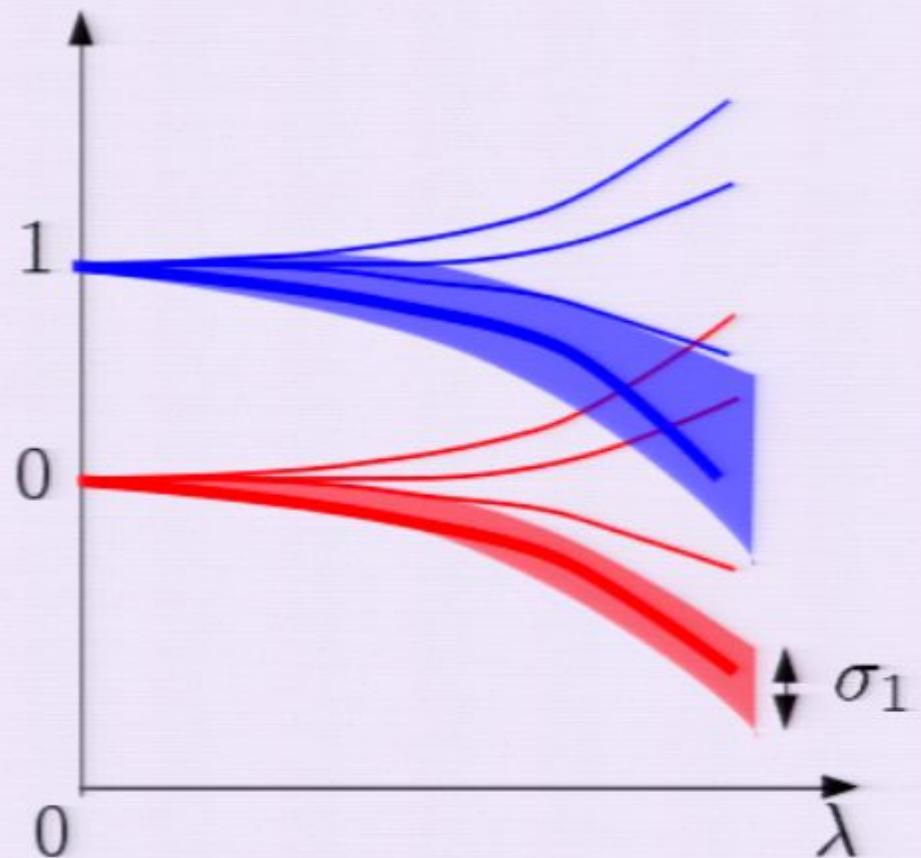
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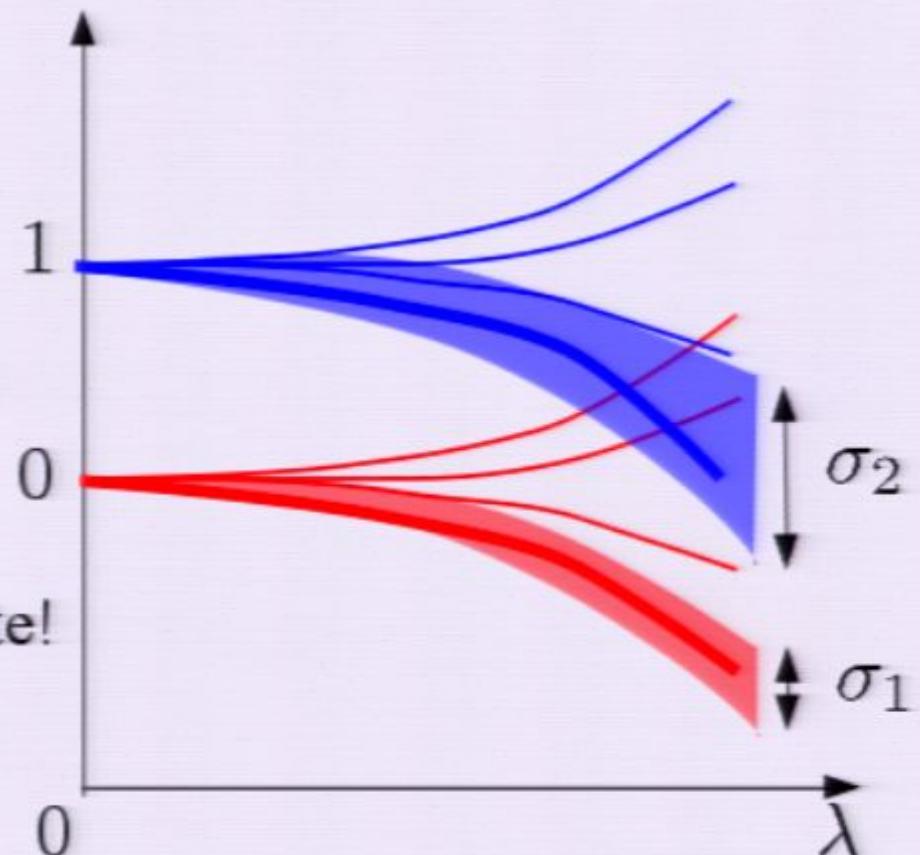
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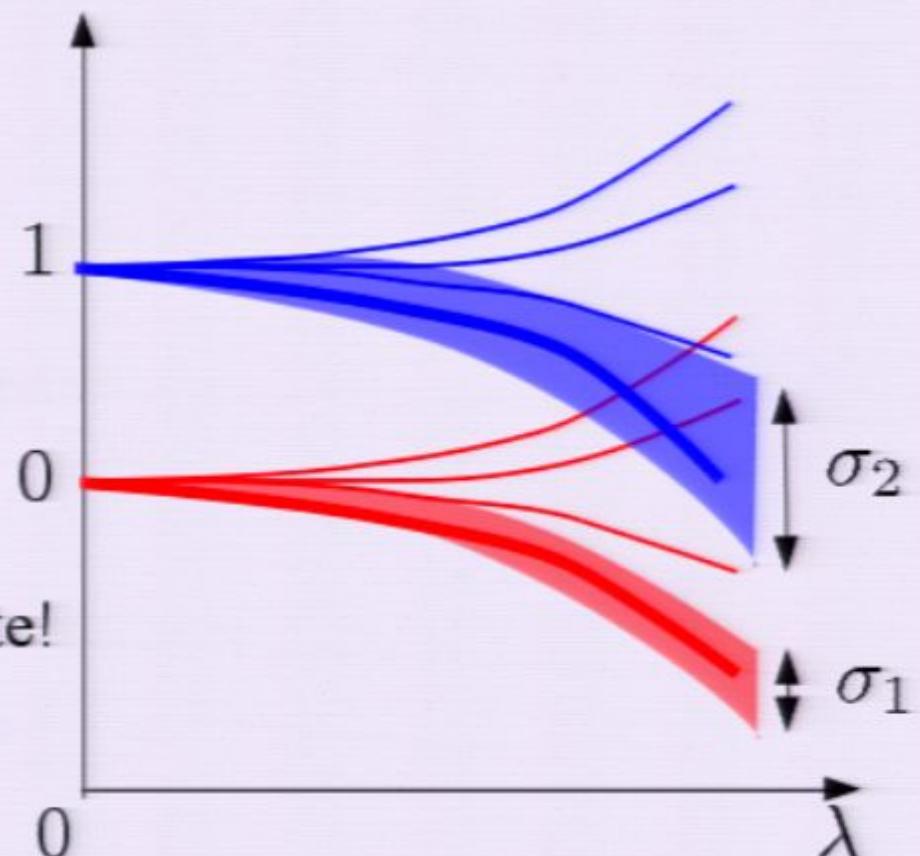
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- Also: $\eta \rightarrow 0$ as $\alpha \rightarrow \alpha_s$

\Rightarrow Effect of degeneracy only appears for large N