Title: Epistemology and the laws of nature

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Abstract:

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# EPISTEMOLOGY AND THE LAWS OF NATURE

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## ROADMAP

- 1) Shortcomings of current impossibility results concerning laws of Nature
  - 2) Knowledge operators and their shortcoming
- 3) Formalize mathematical structure shared by observation and prediction: inference devices

4) Elementary properties of inference devices

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## **COMPUTATION AND PHYSICS**

- ) Impossibility results of Moore, Pour-El and Richards, etc. rely on uncountable number of states of universe.
  - What if universe is countable, or even finite?
  - What if there exist oracles, so Halting theorem (the basis of those results) is irrelevant?

- 2) Impossibility results of Lloyd rely on current model of laws of physics (e.g., no superluminal travel).
  - Pirsa: 10050057 What if laws are actually different?

## COMPUTATION AND PHYSICS

- 3) To apply Godel's incompleteness theorem presumes physical laws are "written in predicate logic"
  - What if universe is written in a different language?
  - What if there are no "laws" at all, just a huge list of events, which just happen to appear to have patterns?
  - What if Godel-style intuitionism is correct?

What if our models are wrong?

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## COMPUTATION AND PHYSICS

- 3) To apply Godel's incompleteness theorem presumes physical laws are "written in predicate logic"
  - What if universe is written in a different language?
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  - What if Godel-style intuitionism is correct?

What if our model is wrong?

Is there some model more fundamental, (almost) impossible not to accept?

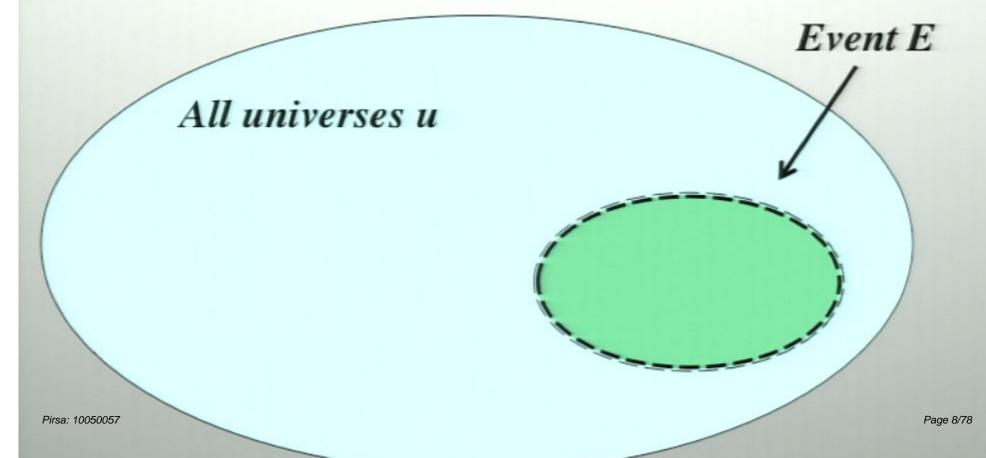
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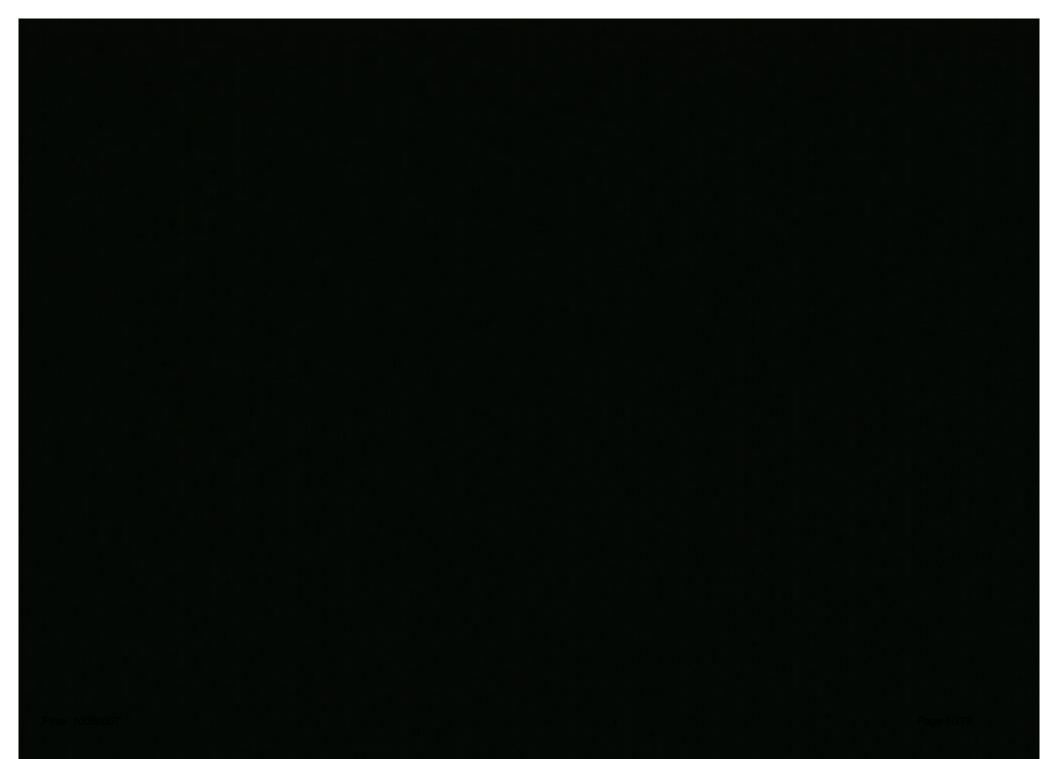
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- ) Laws are patterns, relating "events I know"
- ?) What does it mean to "know" an event?



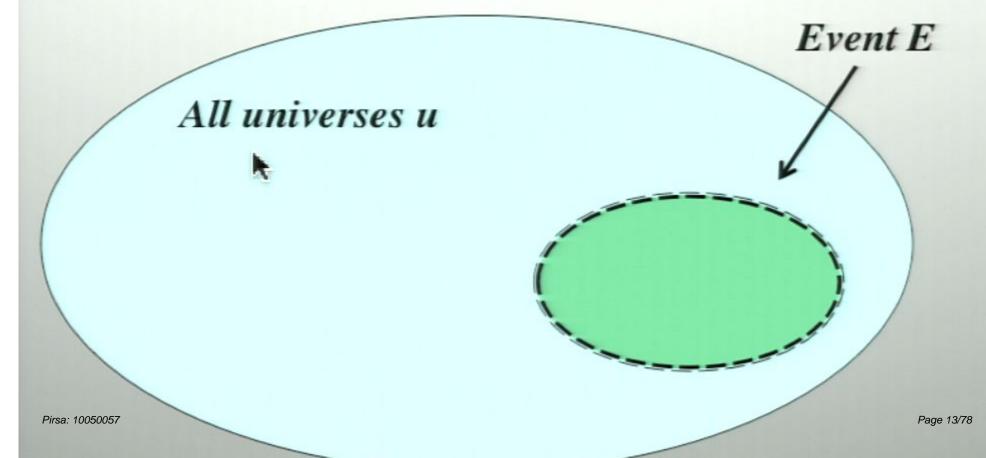
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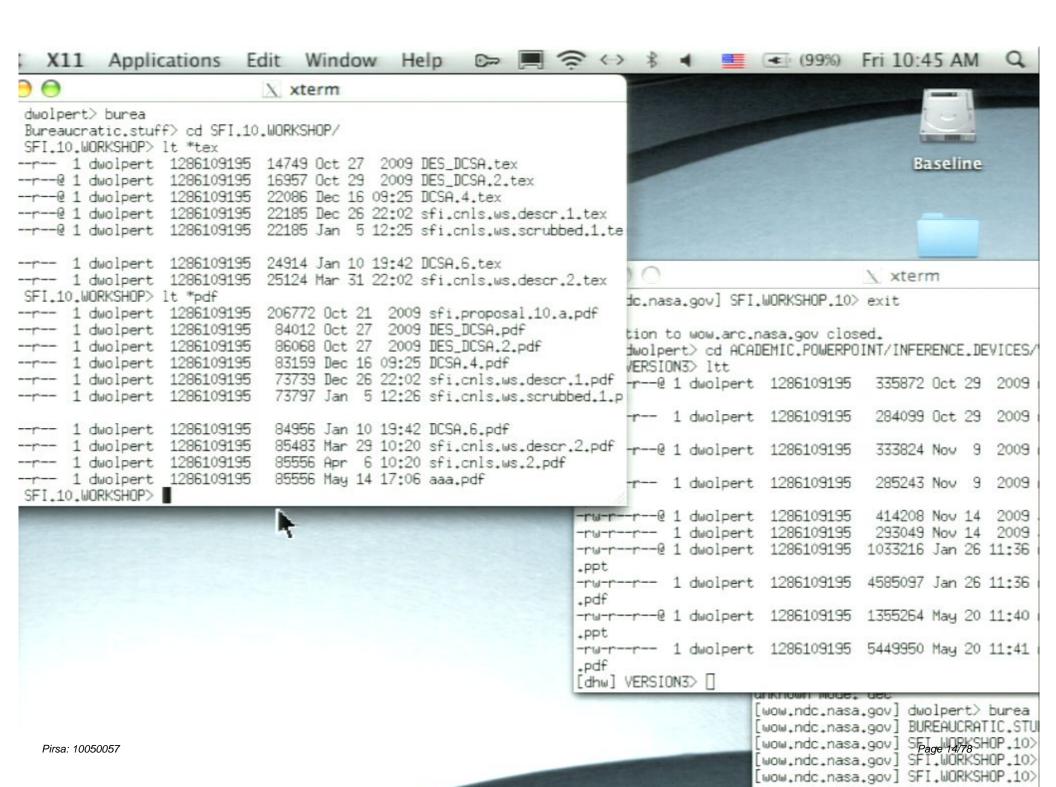


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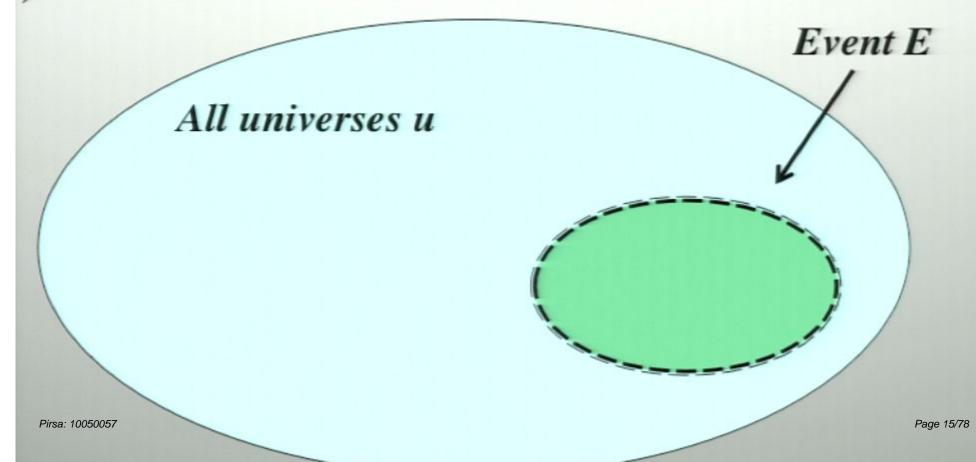
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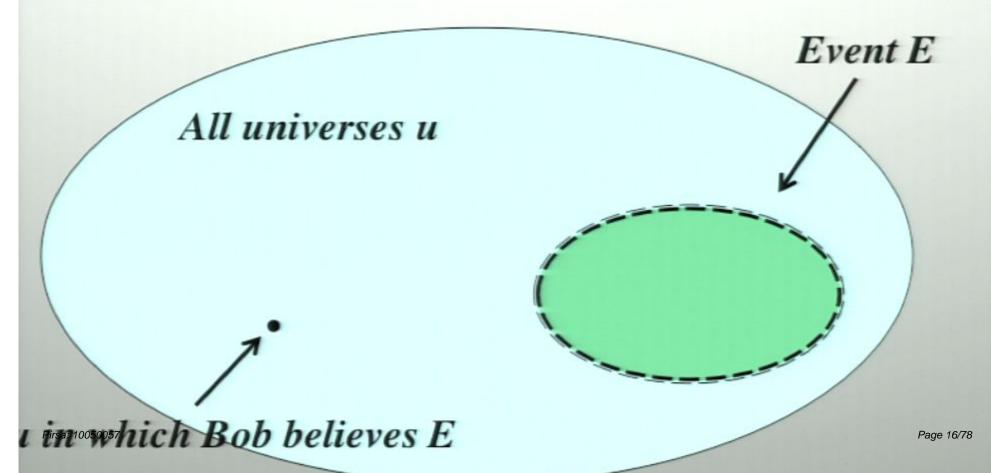




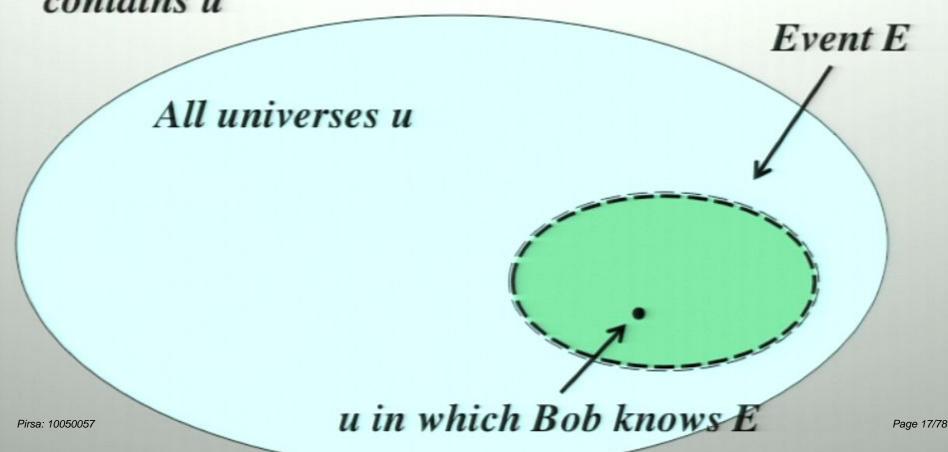
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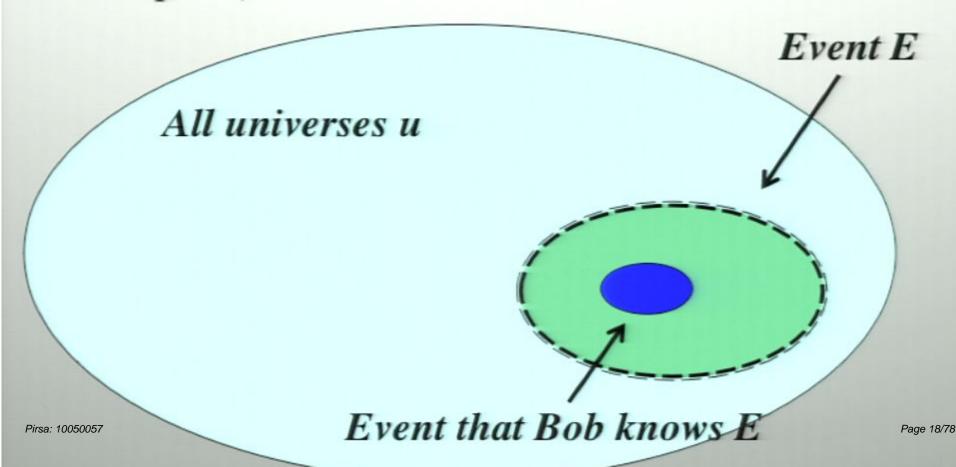
- ) At indicated u, Bob has belief that he's in E.
- ?) Belief is a function from u to subsets of {u}



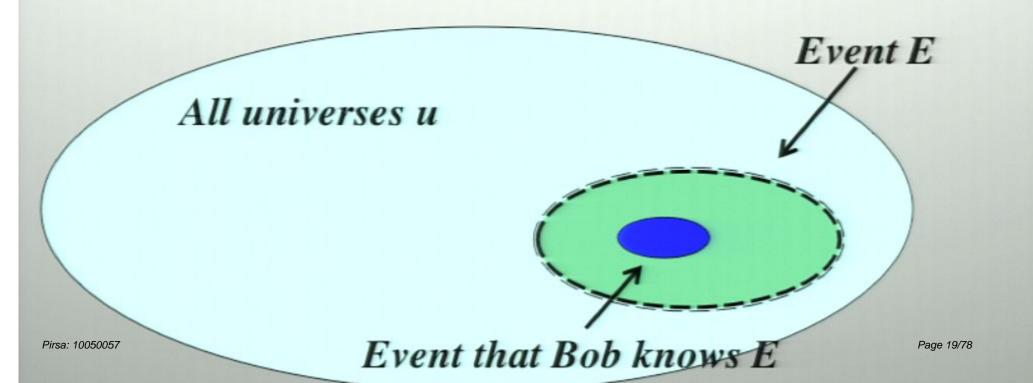
- ) At indicated u, Bob knows he's in E.
- 2) Knowledge is a belief function where the image of u contains u



(At other u in E, Bob instead knows some E' that overlaps E.)



- Need physical manifestation of such knowledge, if Bob's having it involves the laws of nature.
- 2) So need Bob to be able to physically answer questions about what he knows.



- 1) What does it means for Bob to be able to physically answer questions about what he knows?
- 2) To formalize this, analyze physical phenomena where Bob knows an event.
- 3) These are phenomena where information outside Bob gets inside Bob.
- () Examples:
  - Observation
  - Prediction
  - Memory

## ROADMAP

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4) Elementary properties of inference devices

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## EXAMPLE OF PHYSICAL KNOWLDEGE: OBSERVATION

Present a stylized example of observation.

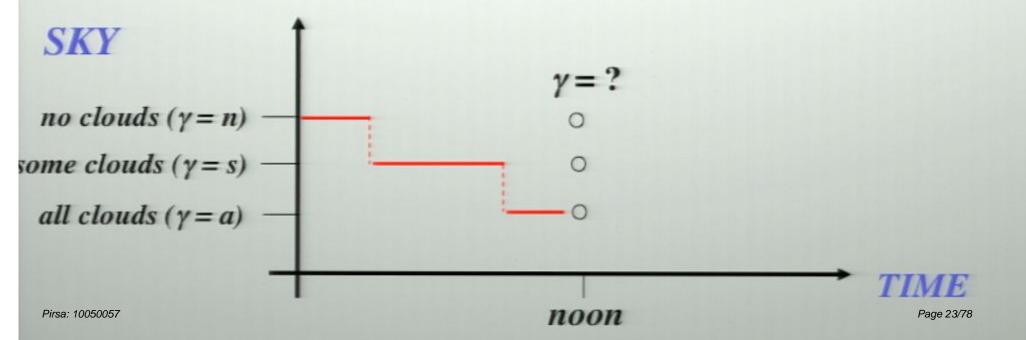
Emphasize features of that example found in all "observations"

Why those features are always found in observations:

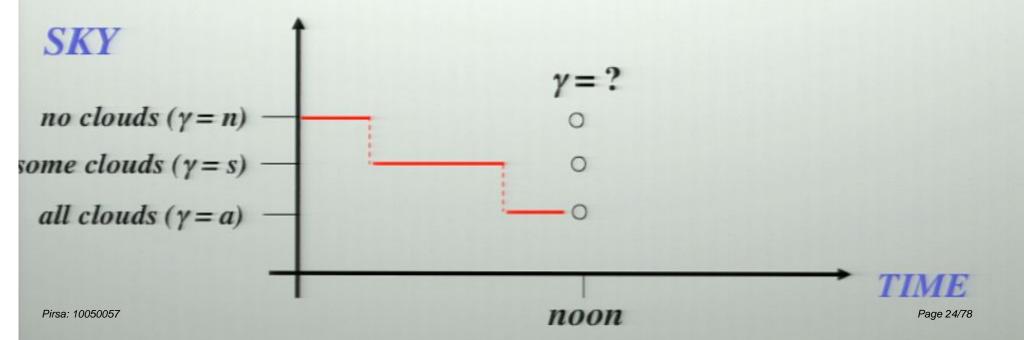
Without those features, the observation conveys no semantic information

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Want to observe γ, state of sky at noon tomorrow

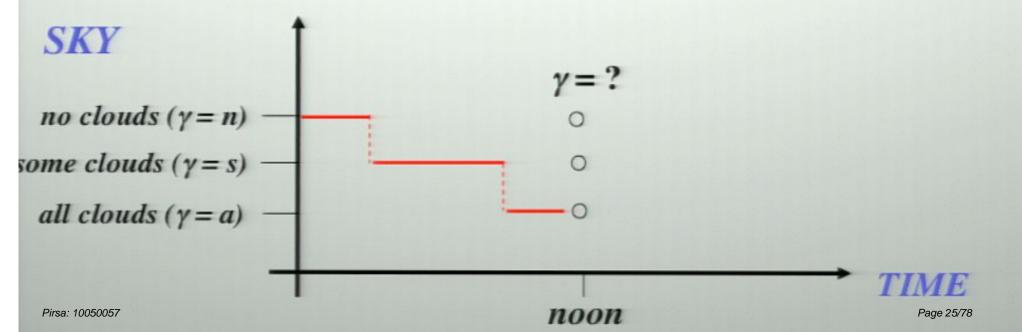


Bob claims to be able to make that observation



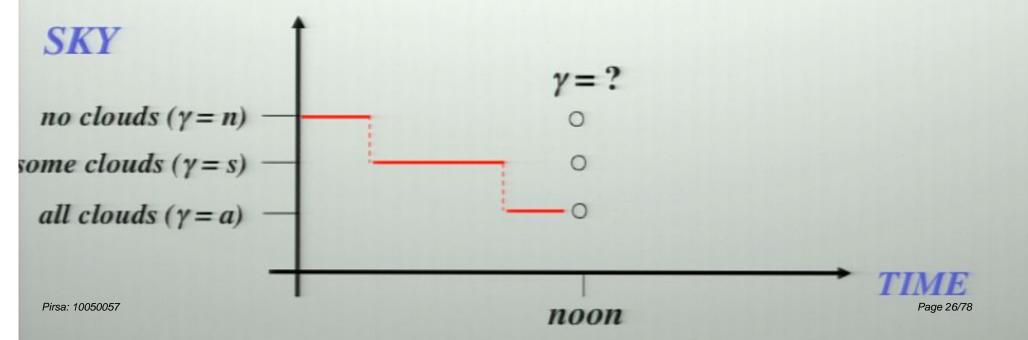
If Bob's claim is true, he can correctly answer three questions that could be posed to him:

- i) Does  $\gamma = 'n'$ ? (Yes / no)
- ii) Does  $\gamma = s'$ ? (Yes / no)
- iii) Does  $\gamma = 'a'$ ? (Yes / no)



## What does this mean physically?

 Restrict attention to universes where Bob and the sky exist; Bob considers one of the three binary questions; then observes γ; then gives his honest answer to that question.



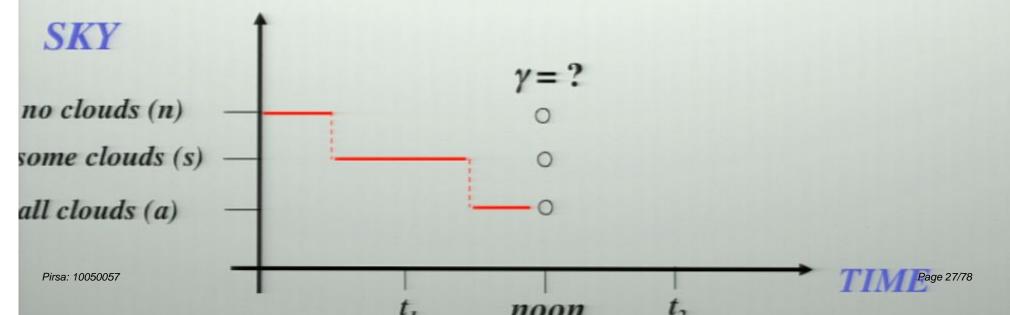
3ob can observe γ if for each of the three questions, q,

The universe has property  $x_q$ :

"At some  $t_1 < noon Bob considers q$ "

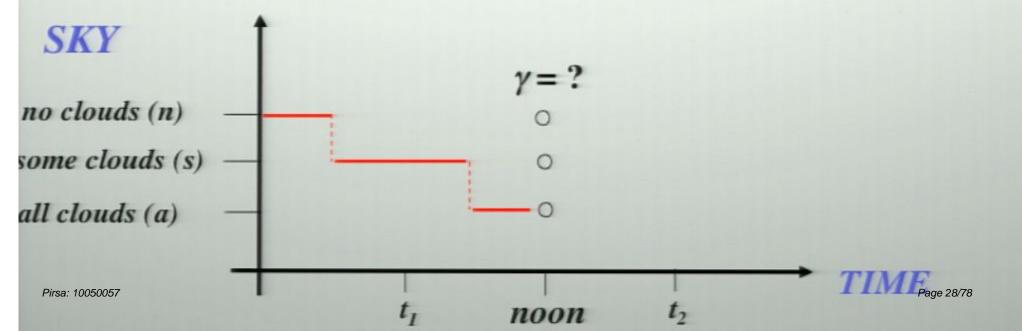


y, the binary answer Bob gives at some  $t_2 > noon$ , equals correct answer to q



 $U = \{all\ universe\ histories\ consistent\ with\ physics,\ in\ which\ Bob\ and\ the\ sky\ exist;\ at\ t_1\ Bob\ considers\ a\ q;\ then\ observes\ \gamma;\ then\ gives\ honest\ answer\ to\ that\ q\}$ 

State of sky at noon is fixed by  $u \in U$ , the actual universe-history. So  $\gamma = \Gamma(u)$  for some function  $\Gamma$ 



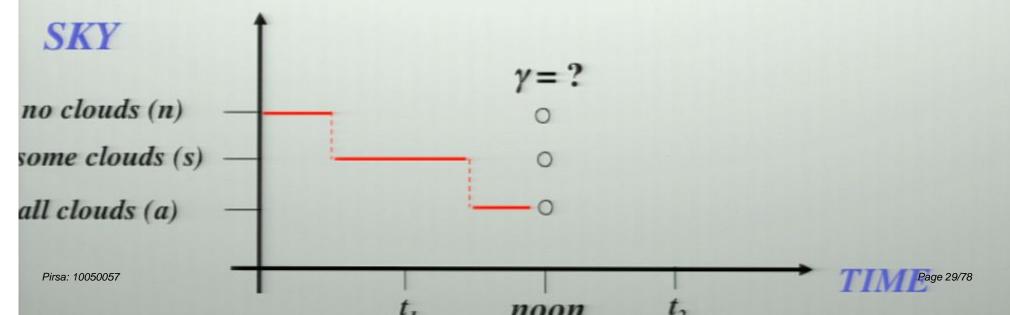
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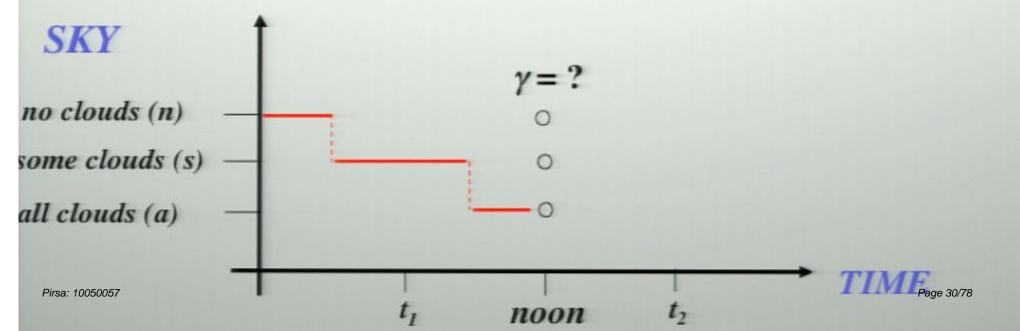


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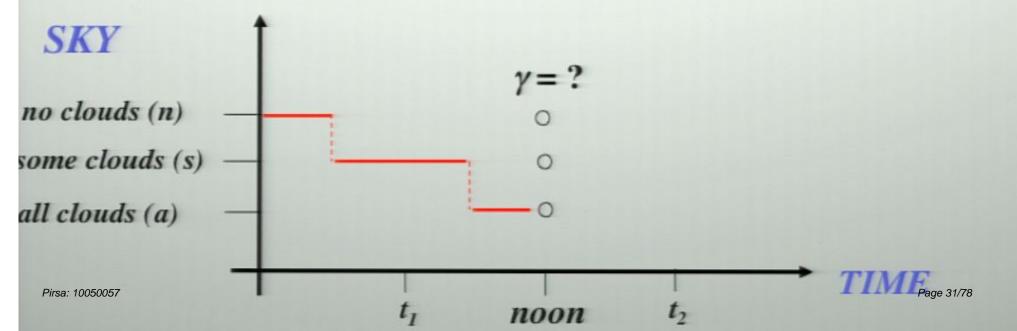
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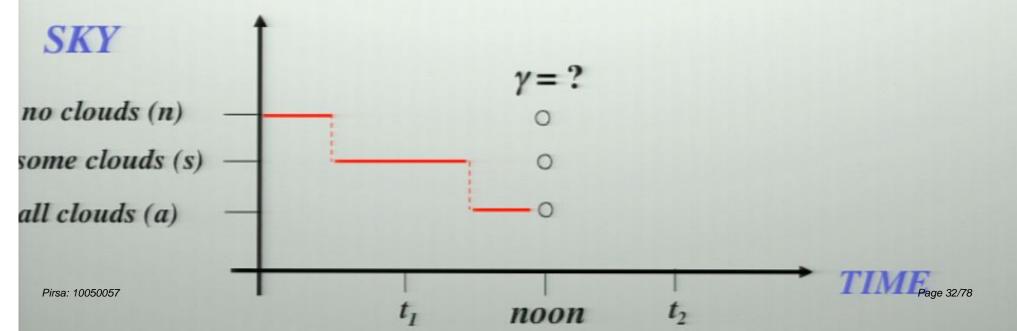
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The question Bob considers at  $t_1$  is given by the actual universe-history  $u \in U$ , so x = X(u) for some function X

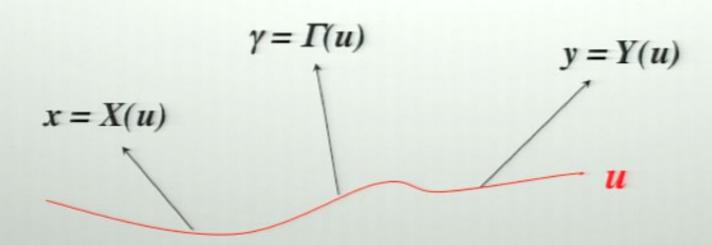


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Bob's answer at  $t_2$  is given by the actual universe-history  $u \in U$ , so y = Y(u) for some function Y

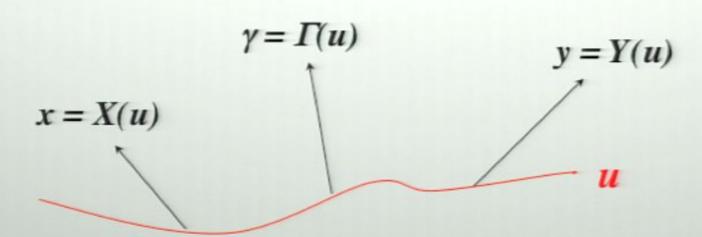


- $\gamma \in \{ (n', (s', (a')) = sky \text{ at noon} = I(u) \}$
- $x = (what \ q \ Bob \ considers \ at \ t_I) = X(u)$
- $y = (Bob's \ answer \ at \ t_2) = Y(u)$



## o:

- For each of three binary questions  $q_{\gamma}$ ,
- ∃x such that
- $X(u) = x \Rightarrow Y(u) = q_{\gamma}(\Gamma(u))$



- Nothing about observation process;
   all about what it means to successfully observe.
- · The 'what' of observation, not the 'how'.

For each of three binary questions  $q_{\gamma}$ ,  $\exists x \text{ such that } X(u) = x \implies Y(u) = q_{\gamma}(\Gamma(u))$ 

$$\gamma = \Gamma(u) \\
x = X(u)$$

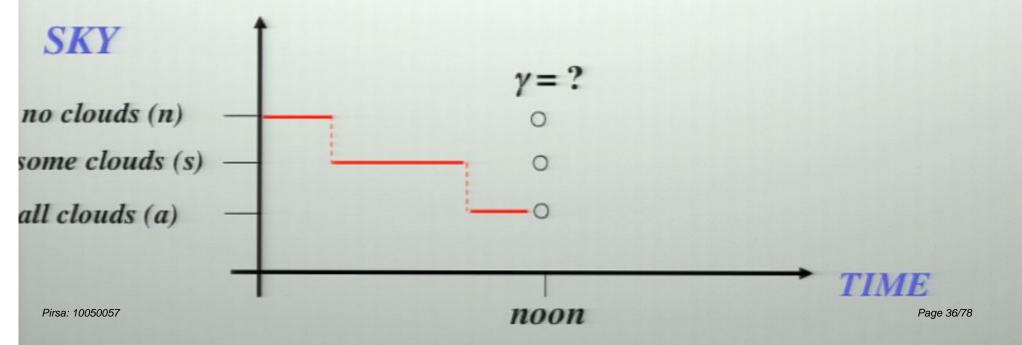
$$y = Y(u)$$

$$u$$

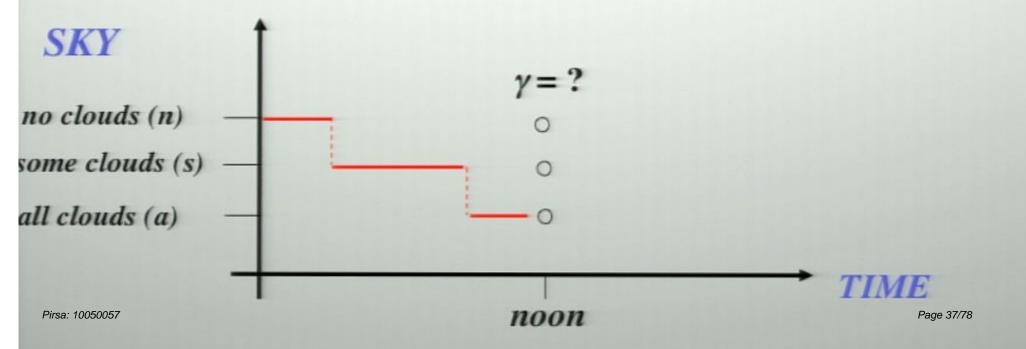
TIME

### **PREDICTION**

Want to predict γ, state of sky at noon tomorrow

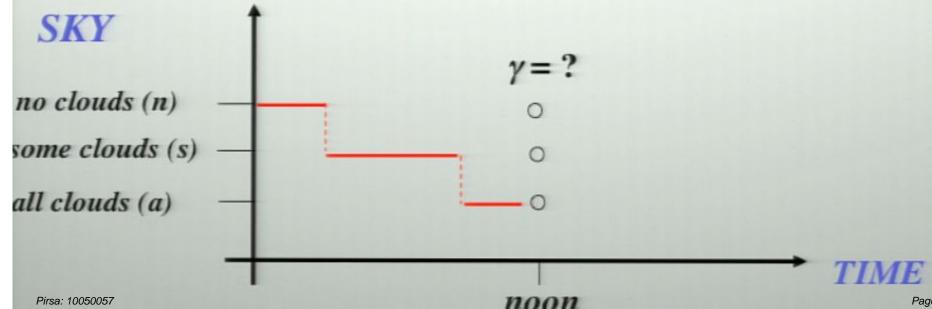


 Bob claims to have a laptop that he can program to make that prediction



If Bob's claim is true, he will be able to correctly answer three questions:

- i) Does  $\gamma = 'n'$ ? (Yes / no)
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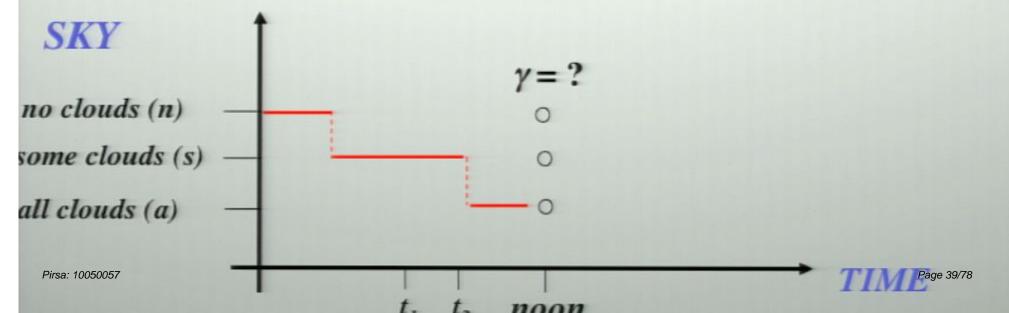


Bob can predict  $\gamma$  if for each of the three questions, q, The universe has property xq:

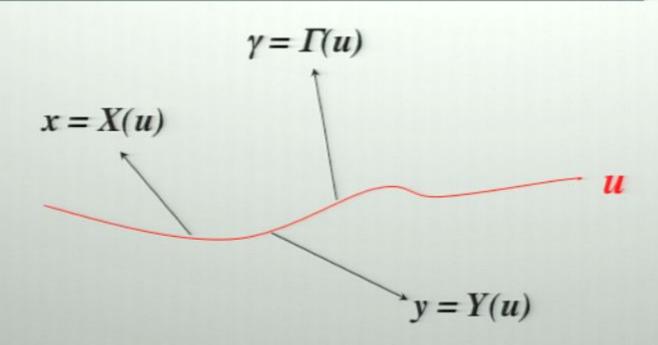
"At some  $t_1 < noon Bob programs the laptop to predict q"$ 



y, the binary answer Bob reads off at some  $t_2 < noon$ , equals correct answer to q

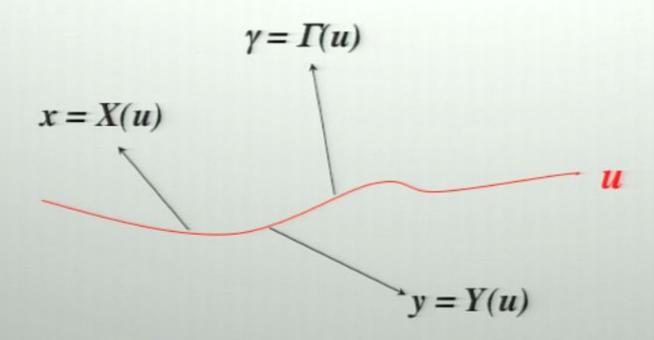


- $\gamma \in \{ (n', (s', (a')) = sky \text{ at noon} = \Pi(u) \}$
- $x = (laptop \ program \ at \ t_1) = X(u)$
- $y = (Bob's \ answer \ at \ t_2) = Y(u)$



## o:

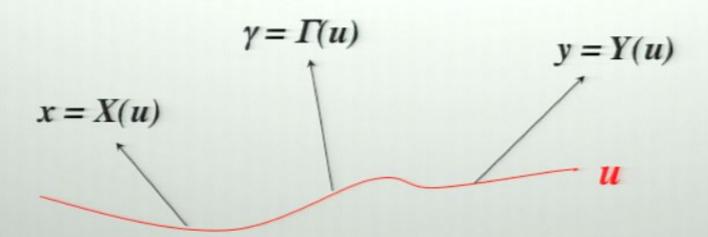
- For each of three binary questions  $q_{\gamma}$ ,
- If a such that
- $X(u) = x \Rightarrow Y(u) = q_{\gamma}(\Gamma(u))$



## **OBSERVATION**

## o:

- For each of three binary questions  $q_{\gamma}$ ,
- ∃x such that
- $X(u) = x \Rightarrow Y(u) = q_{\gamma}(\Gamma(u))$



#### KNOWLEDGE

- More generally, if at some time,
   "Bob knows the state of the sky at noon", γ, then he can answer three questions:
  - i) Does  $\gamma = 'n'$ ? (Yes / no)
  - ii) Does  $\gamma = s'$ ? (Yes / no)
  - iii) Does  $\gamma = 'a'$ ? (Yes / no)
- Note no chronological ordering. Just X (what question Bob considers),

Y (his answer),

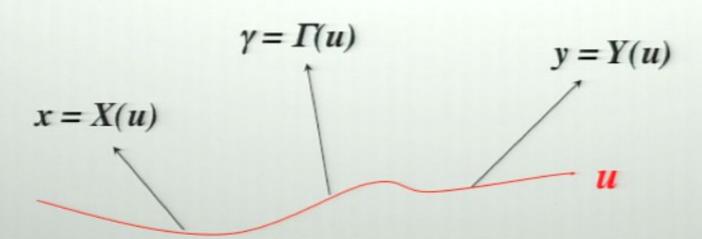
 $\Gamma$  (the sky's actual state at noon),

- all functions of  $u \in U$ 

### **OBSERVATION**

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Y (his answer),

 $\Gamma$  (the sky's actual state at noon),

- all functions of  $u \in U$ 

- An inference device is any two functions (X,Y) over U, where range of Y is binary.
- An inference device (X, Y) (weakly) infers a function Γ over U iff
   ∀ γ in Γ's range,
   ∃x such that X(u) = x ⇒ Y(u) = q,(Γ(u))
- A necessary condition to say that (X,Y) "observes", "predicts", or "knows"  $\Gamma$  is that (X,Y) weakly infers  $\Gamma$ .
- No claims of sufficiency; observation, prediction, knowledge, etc. involve much more than just weak inference.
- But even requiring weak inference restricts observation, Page 46/78

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- Note the use of counterfactual questions, asking whether Γ(u)
  has a value that it does not have. Analogous to the use of
  intervention to define causality in Bayes nets.
- Contrast inference with Aumann-style "knowledge operators"

- Advantages of using binary questions:
  - i) Formalism doesn't change if range of  $\Gamma$  changes
  - ii) Device never need give value  $\Gamma(u)$ , only confirm/reject suggested  $\Gamma(u)$ 's. (Cf. computational complexity)
  - iii) Formalizes semantic information (contrast Shannon)

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## INFERENCE DEVICES TERMINOLOGY

1) Setup function X over U

2) Conclusion binary-valued function Y over U

3) " $(X,Y) > \Gamma$ " means (X,Y) weakly infers  $\Gamma$ 

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# INFERENCE DEVICES AND THE LAWS OF NATURE

- 1) A reality is a space U, a set of devices defined over U, and a set of functions the devices might infer.
- 2) So a reality is a quadruple,  $(U, \{X_j, Y_j\}, \{\Gamma_i\})$ .
- 3) As far as any device in a reality is concerned, U is irrelevant. It's only the inference graph relating the sets  $\{X_j, Y_j\}$  and  $\{\Gamma_i\}$  that matter:

The laws of Nature are patterns in the inference graph of a reality

# ROADMAP

- 1) Shortcomings of current impossibility results concerning laws of Nature
  - 2) Knowledge operators and their shortcoming
- 3) Formalize mathematical structure shared by observation and prediction: inference devices

4) Elementary properties of inference devices

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#### ELEMENTARY PROPERTIES OF INFERENCE

1) Inference need not be transitive:

$$(X_1, Y_1) > Y_2$$
 and  $(X_2, Y_2) > Y_3$  does not mean  $(X_1, Y_1) > Y_3$ 

- 2) For any  $\Gamma$ ,  $\exists$  a device that infers  $\Gamma$ .
- 3) For any device,  $\exists$  a  $\Gamma$  it does not infer. (Impossibility result)
  - Intuition: X ~ initial configuration of a Turing machine.
     Y (a bit) ~ whether Turing machine halts or not.
     So apply Halting theorem-style reasoning

## IMPLICATIONS OF IMPOSSIBILITY RESULT

- ) For any simulator, there is always a prediction that cannot be guaranteed correct.
  - Laplace was wrong.
  - Results of Pour-El et al., Fredkin et al., Moore, etc. are far narrower than this general impossibility
- 2) For any observation apparatus, there is always an observation that cannot be guaranteed to be correct.
  - Non-quantum mechanical "uncertainty principle"

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### BREADTH OF IMPOSSIBILITY RESULT

1) Holds even for a countable U (even for a finite one).

- Holds even if current formulation of physics is wrong.
- 3) Holds even if C has Super-Turing capability

- 4) Holds even if laws of Nature are not written in predicate logic, or intuitionism is correct,
  - Pirsa: 10050057 even if there are no laws, just a huge list of eyents.

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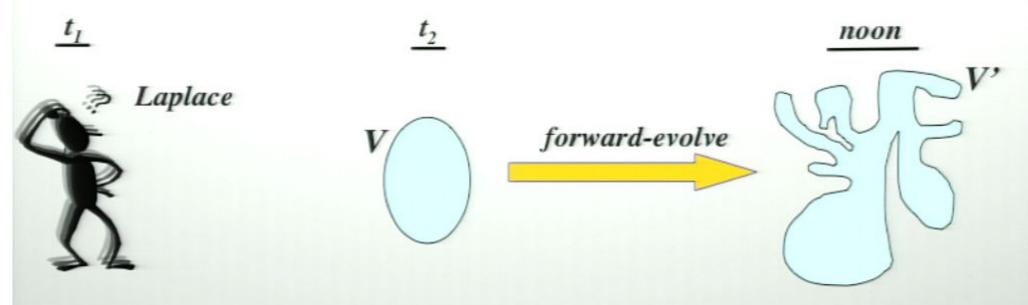
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## **EXAMPLE: PREDICTION FAILURE**



- 1. V = {all time-t<sub>2</sub> universes where Laplace is answering "yes" to his t<sub>1</sub> question}
- 2. V' = V evolved forward to noon
- 3. At t<sub>1</sub>, ask Laplace, "will universe be outside V' at noon?"

Trivially, Laplace's answer is wrong

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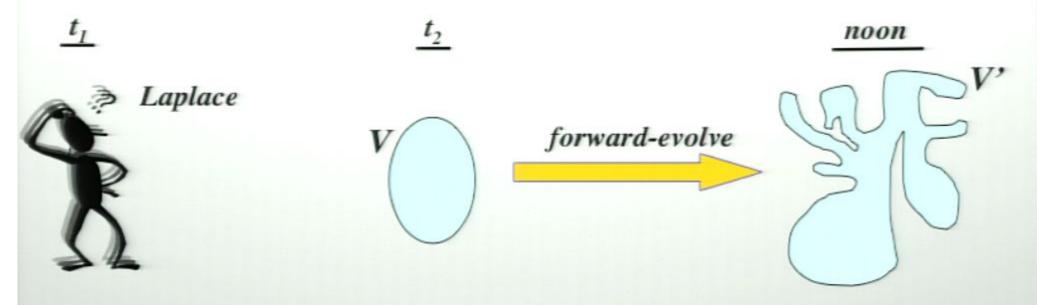
#### INFERENCE RELATIONS BETWEEN DEVICES

Often not interested in inference of arbitrary functions, but rather inference relation among a set of devices.

- I) Two devices (X, Y), (X', Y') are <u>pairwise distinguishable</u> iff every pair (x, x') occurs for some u
- II) A set of devices {(X<sub>i</sub>, Y<sub>i</sub>)} is <u>mutually distinguishable</u> iff every tuple (x<sub>1</sub>, x<sub>2</sub>, ...) occurs for some u

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#### INFERENCE RELATIONS BETWEEN DEVICES -2

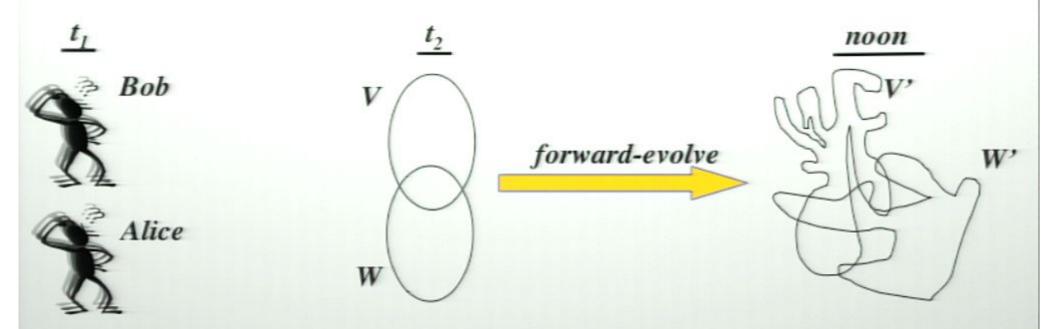
- 3) If all pairs of devices from {C<sub>i</sub>} are pairwise distinguishable, ∃ at most one k : C<sub>k</sub> > C<sub>j</sub> ∀j ≠ k. "Monotheism" theorem.
  - N.b., control is a special type of inference.

4) If all pairs of devices from  $\{C_i\}$  are pairwise distinguishable, can have  $C_1 > C_2 > ... C_1$ .

If the set of devices  $\{C_i\}$  is mutually distinguishable, cannot have  $C_1 > C_2 > ... C_1$ .

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#### MONOTHEISM EXAMPLE



 $V = \{time-t_2 \text{ universes where Bob is answering 'yes' to his } t_1 \text{ question}\}$ 

 $W = \{time-t_2 \text{ universes where Alice is answering 'yes' to her } t_1 \text{ question} \}$ 

V' = V evolved forward to noon

W'= W evolved forward to noon

At t, ask Bob, "will universe be in W' at noon?"

At t1, ask Alice, "will universe be outside of V' at noon?"

# INFERENCE KNOWLEDGE AND BOOLEAN ALGEBRA

nowledge defined in terms of weak inference obeys many of the properties of Boolean algebra:

- (X, Y) may know A, or may know ~A, but not both.
- If (X, Y) knows event A, and knows event  $A \Rightarrow B$ , then B is true.
  - However (X, Y) need not know B; no problem of knowing all truths via deduction.
- If (X, Y) knows  $A \Rightarrow B$  and (X, Y) knows  $B \Rightarrow C$ , then (X, Y) knows  $A \Rightarrow C$ .
- Plist 10050057, Y) knows A, then (X, Y) knows event "(X, Y) knows age 36/78.

#### STOCHASTIC INFERENCE

- What changes if there is probability measure P over U?
- Given a function Γ and device C = (X, Y), C infers Γ with <u>covariance accuracy</u>

$$\varepsilon(C,\Gamma) = \frac{\sum_{q_L} \max_{x} \left[ E_P(Yq_L(\Gamma)|x) \right]}{|\Gamma(U)|}$$

2) Can't instead use mutual information; that only captures syntactic content of distributions, not semantic content.

## EXAMPLE OF STOCHASTIC INFERENCE RESULT

For any probability distribution P over U,

$$\varepsilon((X,Y),\Gamma) \geq (2-n) \; \frac{\max_{x} \left[ E_{P}(Y|x) \right]}{n}$$

where  $n = |\Gamma(\mathbf{U})|$ 

) For any probability distribution P over U, there exists two devices  $(X_1, Y_1)$ ,  $(X_2, Y_2)$  where  $X_1$  and  $X_2$  are distinguishable, but both  $\varepsilon((X_1, Y_1), Y_2)$  and  $\varepsilon((X_1, Y_1), Y_2)$  are arbitrarily close to one;

Second Laplace impossibility theorem is "barely true"

# HOWEVER RELATED RESULTS ARE QUITE STRONG

- 1) Let  $C_1$  and  $C_2$  be two devices, where:
  - i) Both  $X_1(U)$  and  $X_2(U)$  are the binaries;
  - ii)  $C_1 > C_2$  with accuracy  $\varepsilon_1$ , and  $C_2 > C_1$  with accuracy  $\varepsilon_2$ .
  - iii)  $P(X_1 = -1) = \alpha$ , and  $P(X_2 = -1) = \beta$
- 2) Define H as the four-dimensional unit open hypercube, and
  - i)  $\forall z \in H, k(z) = z_1 + z_4 z_2 z_3;$
  - ii)  $\forall z \in H, m(z) = z_2 z_4;$
  - iii)  $\forall z \in H$ ,  $n(z) = z_3 z_4$ .
- 3)  $\varepsilon_1 \varepsilon_2 \le \max_{z \in H} |\alpha \beta[k(z)]^2 + ak(z)m(z) + bk(z)n(z) + m(z)n(z)|$
- 4) E.g., for  $\alpha = \beta = 1/2$ ,  $\varepsilon_1 \varepsilon_2 \le 1/4$ .

## BREADTH OF IMPOSSIBILITY RESULT

- 1) Holds even for a countable U (even for a finite one).
- 2) Holds even if current formulation of physics is wrong.
- 3) Holds even if C has Super-Turing capability

But I was triver!

- 4) Holds even if laws of Nature are not written in predicate logic, or intuitionism is correct,
  - Pirsa: 10050057 even if there are no laws, just a huge list of eyesonts.

# INFERENCE DEVICES AND THE LAWS OF NATURE

- 1) A reality is a space U, a set of devices defined over U, and a set of functions the devices might infer.
- 2) So a reality is a quadruple,  $(U, \{X_j, Y_j\}, \{\Gamma_i\})$ .
- 3) As far as any device in a reality is concerned, U is irrelevant. It's only the inference graph relating the sets  $\{X_j, Y_j\}$  and  $\{\Gamma_i\}$  that matter:

The laws of Nature are patterns in the inference graph of a reality

#### ELEMENTARY PROPERTIES OF INFERENCE

1) Inference need not be transitive:

$$(X_1, Y_1) > Y_2$$
 and  $(X_2, Y_2) > Y_3$  does not mean  $(X_1, Y_1) > Y_3$ 

- 2) For any  $\Gamma$ ,  $\exists$  a device that infers  $\Gamma$ .
- 3) For any device,  $\exists$  a  $\Gamma$  it does not infer. (Impossibility result)
  - Intuition: X ~ initial configuration of a Turing machine.
     Y (a bit) ~ whether Turing machine halts or not.
     So apply Halting theorem-style reasoning

#### INFERENCE RELATIONS BETWEEN DEVICES -2

- 3) If all pairs of devices from {C<sub>i</sub>} are pairwise distinguishable, ∃ at most one k : C<sub>k</sub> > C<sub>j</sub> ∀j ≠ k. "Monotheism" theorem.
  - N.b., control is a special type of inference.

4) If all pairs of devices from  $\{C_i\}$  are pairwise distinguishable, can have  $C_1 > C_2 > ... C_1$ .

If the set of devices  $\{C_i\}$  is mutually distinguishable, cannot have  $C_1 > C_2 > ... C_1$ .

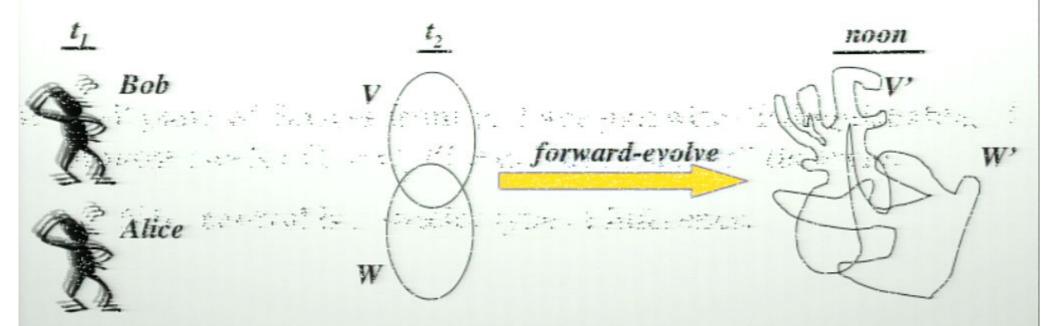
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# INFERENCE KNOWLEDGE AND BOOLEAN ALGEBRA

nowledge defined in terms of weak inference obeys many of the properties of Boolean algebra:

- (X, Y) may know A, or may know ~A, but not both.
- If (X, Y) knows event A, and knows event  $A \Rightarrow B$ , then B is true.
  - However (X, Y) need not know B; no problem of knowing all truths via deduction.
- If (X, Y) knows  $A \Rightarrow B$  and (X, Y) knows  $B \Rightarrow C$ , then (X, Y) knows  $A \Rightarrow C$ .
- Plus 10050057, Y) knows A, then (X, Y) knows event "(X, Y) knows age 14/78.

## MONOTHEISM EXAMPLE



V = {time-t, universes where Bob is answering 'yes' to his t, question}

W = {time-t2 universes where Alice is answering 'yes' to her t1 question}

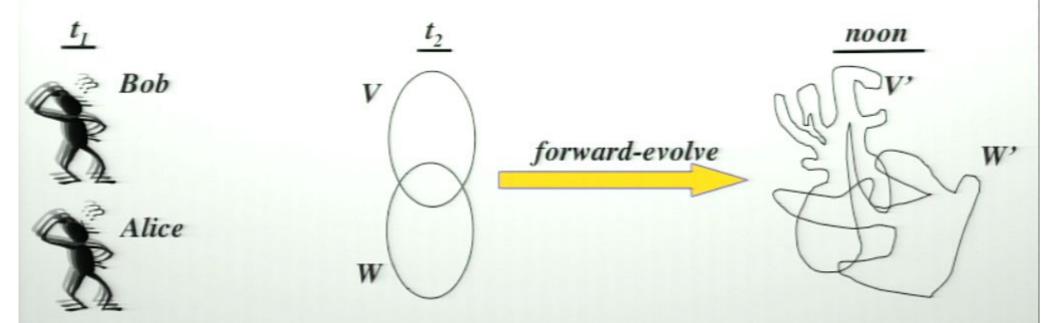
V' = V evolved forward to noon

W'= W evolved forward to noon

At t, ask Bob, "will universe be in W' at noon?"

At t, ask Alice, "will universe be outside of V' at noon?"

#### MONOTHEISM EXAMPLE



 $V = \{time-t_2 \text{ universes where Bob is answering 'yes' to his } t_1 \text{ question} \}$ 

 $W = \{time-t_2 \text{ universes where Alice is answering 'yes' to her } t_1 \text{ question} \}$ 

V' = V evolved forward to noon

W'= W evolved forward to noon

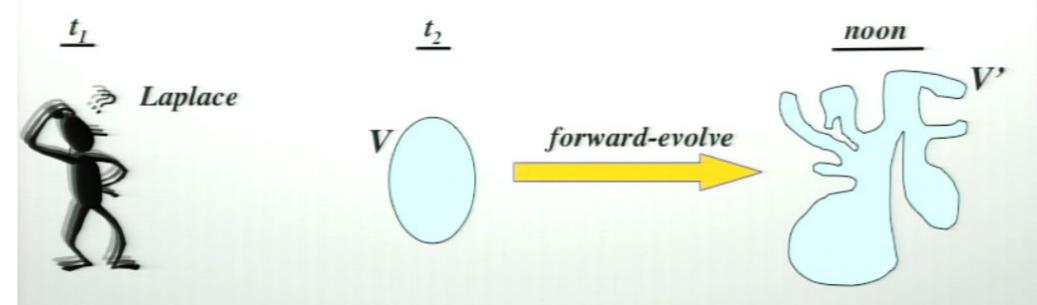
At t, ask Bob, "will universe be in W' at noon?"

At t1, ask Alice, "will universe be outside of V' at noon?"

# HOWEVER RELATED RESULTS ARE QUITE STRONG

- 1) Let  $C_1$  and  $C_2$  be two devices, where:
  - i) Both  $X_1(U)$  and  $X_2(U)$  are the binaries;
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- 4) E.g., for  $\alpha = \beta = 1/2$ ,  $\varepsilon_1 \varepsilon_2 \le 1/4$ .

## **EXAMPLE: PREDICTION FAILURE**



- 1. V = {all time-t<sub>2</sub> universes where Laplace is answering "yes" to his t<sub>1</sub> question}
- 2. V' = V evolved forward to noon
- 3. At t<sub>1</sub>, ask Laplace, "will universe be outside V' at noon?"

Trivially, Laplace's answer is wrong