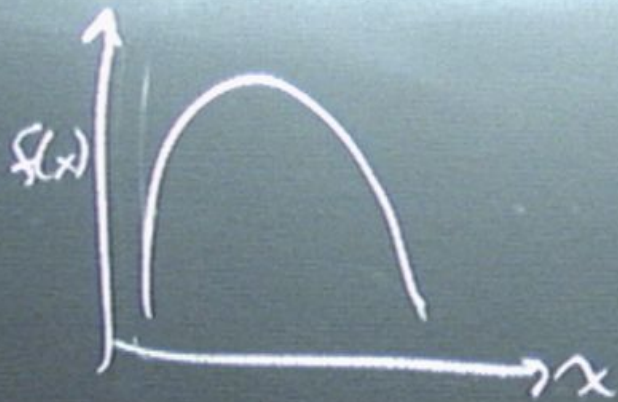
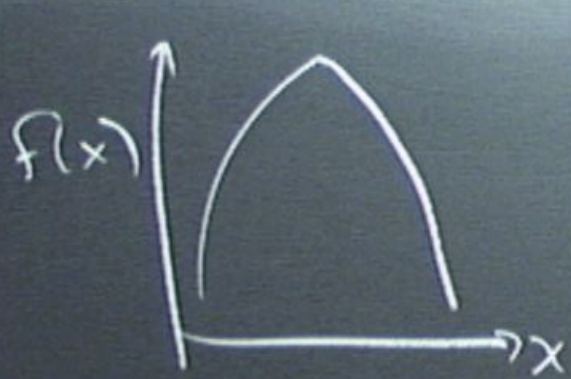


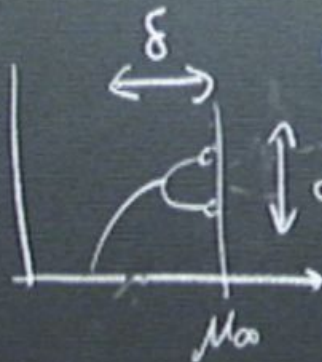
Title: Dynamical Systems - Review (PHYS 607) - Lecture 8

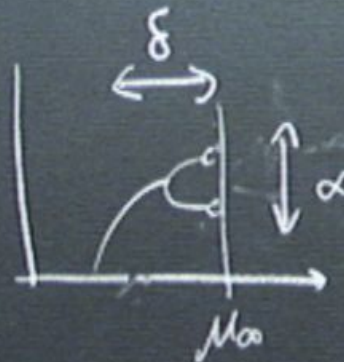
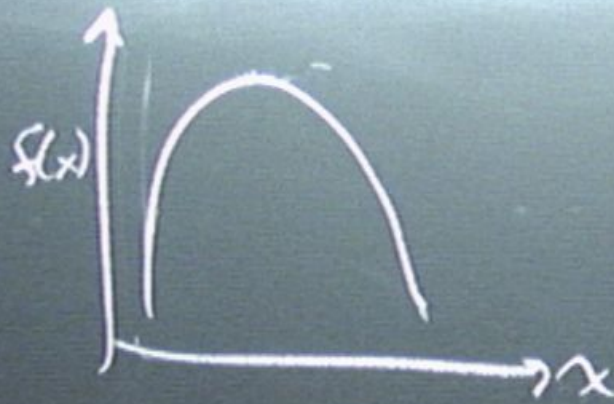
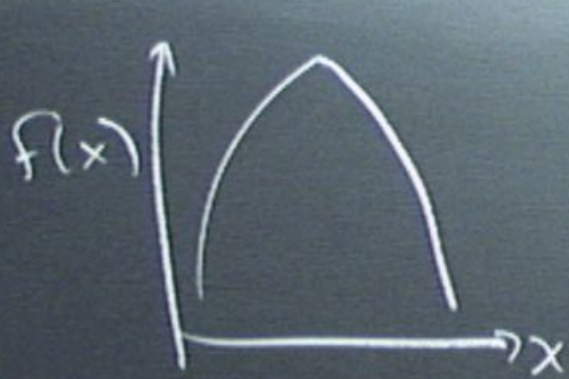
Date: Jan 13, 2010 09:00 AM

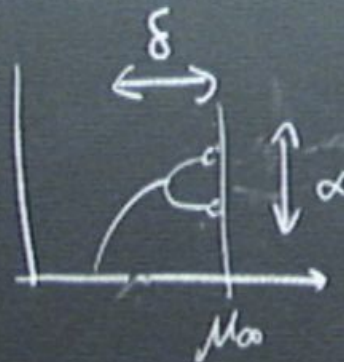
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Abstract:



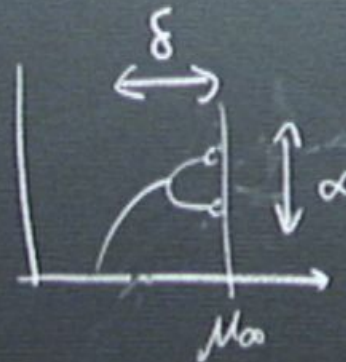






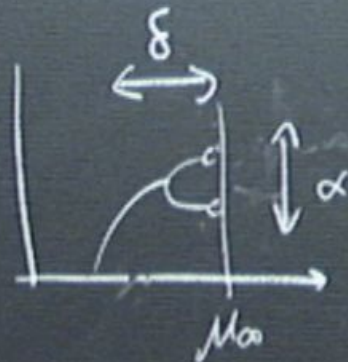
$$g(x) = A_0 + A_n x^n +$$



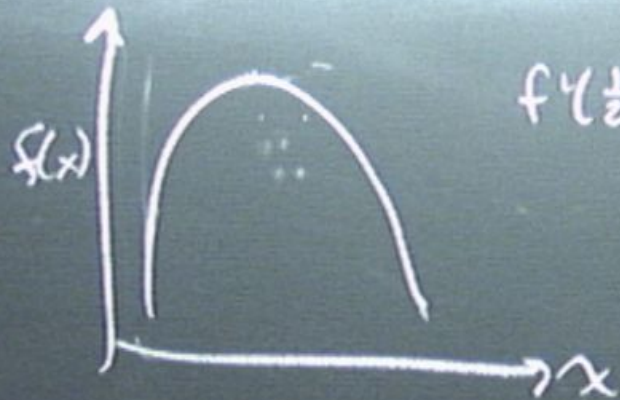


$$g(x) = A_0 + A_1 x^1 + \dots$$



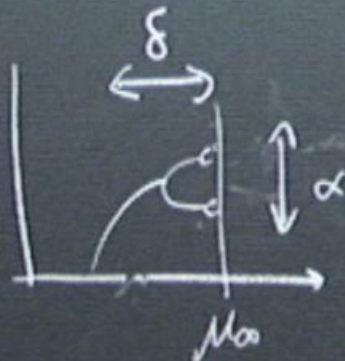


$$g(x) = A_0 + A_1 x^1 + \dots$$



$$f'(\frac{1}{2}) = 0$$

^



$$g(x) = A_0 + A_n x^n + \dots$$

Shannon entropy
quantify information

Set of n possible events

probs of occurrence

p_1, \dots, p_n [known]

Shannon entropy
quantify information

Set of n possible events
probs of occurrence

p_1, \dots, p_n [known]

define measure
 $S(p_1, \dots, p_n)$ of how
much we learn when ~~we find~~
out that one event has
been chosen

requirements

requirements on S :

1) S continuous

2) S monotonically increases with n if n choices are equally likely

3) Breaking down a choice into steps doesn't change S .

Thm

$$S = -K \sum_{i=1}^n P_i \log P_i$$

with $K > 0$

i) all p_i 's = $1/n$
 $S(1/n, 1/n, \dots, 1/n) = A(n)$

Shannon entropy
quantify information

Set of n possible
probabilities of occurrence

p_1, \dots, p_n

define measure

$S(p_1, \dots, p_n)$ of

much we learn when

out that one event

has been chosen

1) all p_i 's = $1/n$
 $S(1/n, 1/n, \dots, 1/n) = A(n)$

2) \Rightarrow must have
 $A(n) = K \log n$
 $K > 0$

Shannon entropy
quantify information

Set of n possible
probabilities of occurrence

$$p_1, \dots, p_n$$

define measure
 $S(p_1, \dots, p_n)$ of
much we learn when
out that one event
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1) all p_i 's = $1/n$
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1 2 3 4 5 6 7 8



$A(8)$

Shannon entropy
 quantify information

Set of n possible
 probs of occurrence

p_1, \dots, p_n

define measure

$S(p_1, \dots, p_n)$ of

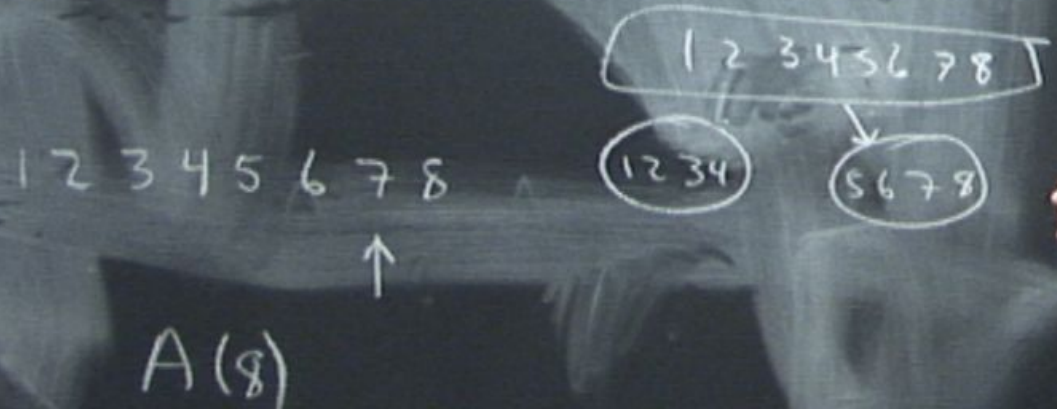
how much we learn when

that one event

is chosen

1) all p_i 's = $1/n$
 $S(1/n, 1/n, \dots, 1/n) = A(n)$

2) \Rightarrow must have
 $A(n) = K \log n$
 $K > 0$



$A(8)$

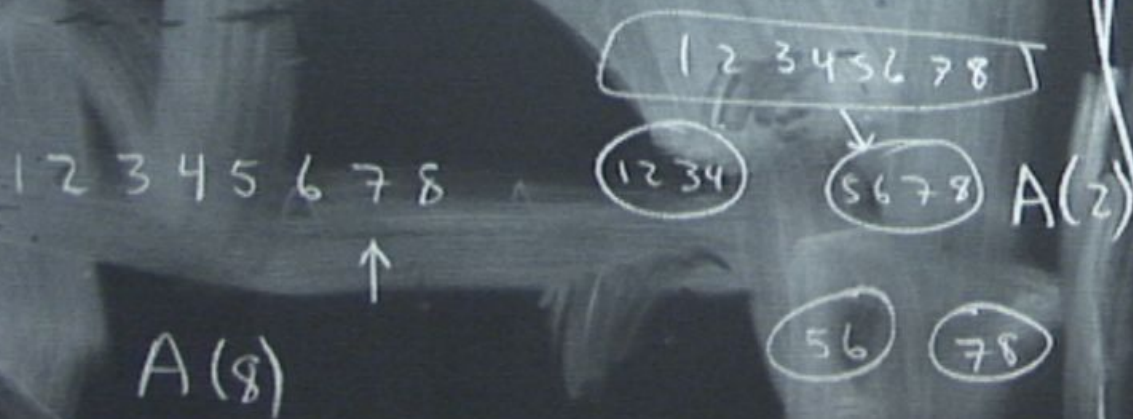
Shannon entropy
 quantify information

Set of n possible
 probs of occurrence

p_1, p_2, \dots, p_n
 a measure
 (p_i) of
 uncertainty when
 one event
 occurs

1) all p_i 's = $1/n$
 $S(1/n, 1/n, \dots, 1/n) = A(n)$

2) \Rightarrow must have
 $A(n) = K \log n$
 $K > 0$



Shannon entropy
 quantify information

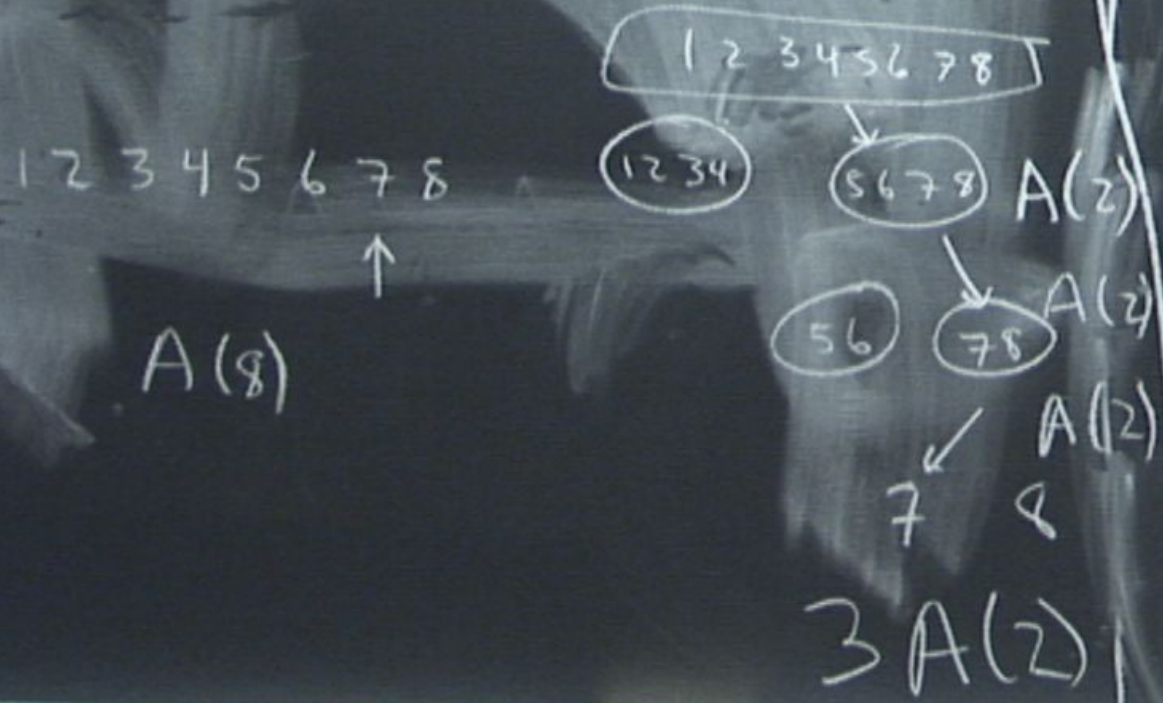
Set of n possible
 probs of occurrence

p_1, \dots, p_n

define measure
 $S(p_1, \dots, p_n)$ of
 how much we learn
 out that one
 has been chosen

1) all p_i 's = $1/n$
 $S(1/n, 1/n, \dots, 1/n) = A(n)$

2) \Rightarrow must have
 $A(n) = K \log n$
 $K > 0$



Shannon entropy
 quantify information
 Set of n possible
 probs of occurrence
 p_1, \dots, p_n
 define measure
 $S(p_1, \dots, p_n)$ of
 how much we learn when
 one of them has
 been chosen

$$A(8) = 3 A(2)$$

$$8 = 2^3$$

$$A(8) = 3 A(2)$$

$$8 = 2^3$$

$$A(2^3) = 3 A(2)$$

$$A(n^y) = y A(n)$$

$$A(m) = K \log m$$

$$\rightarrow A(S^m) = mA(S) \quad A(8) = 3A(2)$$

$$8 = 2^3$$

$$A(2^3) = 3A(2)$$

$$A(n^y) = yA(n)$$

$$\rightarrow A(S^m) = mA(S)$$

$$A(8) = 3A(2)$$

$$8 = 2^3$$

$$A(2^3) = 3A(2)$$

$$A(n^y) = yA(n)$$

$$\rightarrow A(S^m) = mA(S)$$

$$A(8) = 3A(2)$$

$$8 = 2^3$$

$$\underline{A(2^3) = 3A(2)}$$

reverse continuity

1) Choose integers a, b, s, t

$$a^s < b^t$$

require continuity

1) Choose integers a, b, s, t

$$a^s < b^t < a^{s+1}$$

events

now]

we find
t has

require continuity

1) Choose integers a, b, s, t

$$a^s < b^t < a^{s+1}$$

$$s \ln a < t \ln b < (s+1) \ln a$$

$$\frac{s}{t} \ln a < \ln b < \frac{(s+1)}{t} \ln a$$

$$s \frac{\ln a}{\ln b} < \frac{t}{s} < \frac{s+1}{s} \frac{\ln a}{\ln b}$$

reverse continuity

1) Choose integers a, b, s, t

$$a^s < b^t < a^{s+1}$$

$$s \ln a < t \ln b < (s+1) \ln a$$

$$\frac{s}{t} \ln a < \ln b < \frac{(s+1)}{t} \ln a$$

$$s \frac{\ln a}{\ln b} < \frac{t}{s} < \frac{s+1}{s} \frac{\ln a}{\ln b}$$

$$A(a^s) < A(b^t) < A(a^{s+1})$$

reverse continuity

1) Choose integers a, b, s, t

$$a^s < b^t < a^{s+1}$$

$$s \ln a < t \ln b < (s+1) \ln a$$

$$\frac{s}{t} \ln a < \ln b < \frac{(s+1)}{t} \ln a$$

$$s \frac{\ln a}{\ln b} < t < \frac{s+1}{s} \frac{\ln a}{\ln b}$$

$$A(a^s) < A(b^t) < A(a^{s+1})$$

$$sA(a) < tA(b) < (s+1)A(a)$$

require continuity

Choose integers a, b, s, t

$$a^s < b^t < a^{s+1}$$

$$s \ln a < t \ln b < (s+1) \ln a$$

$$\frac{s}{t} \ln a < \ln b < \frac{(s+1)}{t} \ln a$$

$$s \frac{\ln a}{\ln b} < t < \frac{s+1}{s} \frac{\ln a}{\ln b}$$

$$A(a^s) < A(b^t) < A(a^{s+1})$$

$$sA(a) < tA(b) < (s+1)A(a)$$

$$S(A(a)) < S(A(b)) < S(A(a|a))$$

A

Shannon entropy
quantify information

Set of n possible
probabilities of occurrence

p_1, \dots, p_n

define measure

$S(p_1, \dots, p_n)$

much used

ent

been chosen

$$s A(a) < t A(b) < (s+1) A(a)$$

$$\frac{A(a)}{A(b)} < \frac{t}{s} < \frac{s+1}{s} \frac{A(a)}{A(b)}$$

Shannon entropy
quantify information

Set of n possible
probabilities of occurrence

p_1, \dots, p_n

define measure

$S(p_1, \dots, p_n)$ of

much uncertainty

ent

been chosen

$$s A(a) < t A(b) < (s+1) A(a)$$

$$\frac{A(a)}{A(b)} < \frac{t}{s} < \frac{s+1}{s} \frac{A(a)}{A(b)}$$

$$\textcircled{1} \quad \frac{A(a)}{A(b)} < \frac{\ln a}{\ln b} \left(1 + \frac{1}{s}\right)$$

$$\textcircled{2} \quad \frac{A(a)}{A(b)} > \frac{\ln a}{\ln b} \left(1 + \frac{1}{s}\right)$$

Shannon entropy
quantify information

Set of n possible
probabilities of occurrence

p_1, \dots, p_n

define measure
 $S(p_1, \dots, p_n)$ of
much uncertainty
has been chosen

$$s A(a) < t A(b) < (s+1) A(a)$$

$$\frac{A(a)}{A(b)} < \frac{t}{s} < \frac{s+1}{s} \frac{A(a)}{A(b)}$$

$$\textcircled{1} \quad \frac{A(a)}{A(b)} < \frac{\ln a}{\ln b} \left(1 + \frac{1}{s}\right)$$

$$\textcircled{2} \quad \frac{A(a)}{A(b)} > \frac{\ln a}{\ln b} \left(1 + \frac{1}{s}\right)$$

true for any s

$$\Rightarrow \text{true as } s \rightarrow \infty \Rightarrow \frac{A(a)}{A(b)} = \frac{\ln a}{\ln b}$$

Shannon entropy
quantify information

Set of n possible
probabilities of occurrence

p_1, \dots, p_n

define measure
 $S(p_1, \dots, p_n)$ of
uncertainty

entirely
been chosen

$$\Rightarrow A(a) = K \ln a$$

$$\left[K = \frac{A(a)}{\ln b} \text{ for some } b \right]$$

require continuity

Choose integer

$$a^s < b$$

$$s \ln a < t$$

<

$$\Rightarrow A(a) = K \ln a$$

$$\left[K = \frac{A(b)}{\ln b} \text{ for some } b \right]$$

alternative proof.

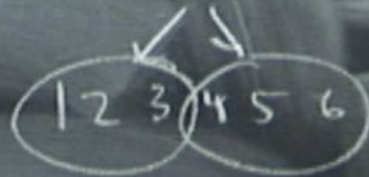
$$A(mn) = A(m) + A(n)$$

$$\Rightarrow A(a) = K \ln a$$

$$\left[K = \frac{A(b)}{\ln b} \text{ for some } b \right]$$

alternative proof.

$$A(mn) = A(m) + A(n)$$



$$\Rightarrow A(a) = K \ln a$$

$$\left[K = \frac{A(b)}{\ln b} \text{ for some } b \right]$$

alternative proof.

$$A(mn) = A(m) + A(n)$$

$$\text{let } x = \ln m$$

$$y = \ln n$$

$$\tilde{A}(x) = A(e^x)$$

$$\Rightarrow A(a) = K \ln a$$

$$\left[K = \frac{A(b)}{\ln b} \text{ for some } b \right]$$

alternative proof.

$$A(mn) = A(m) + A(n)$$

$$\text{let } x = \ln m$$

$$y = \ln n$$

$$\tilde{A}(x) = A(e^x)$$

$$\tilde{A}(x+y) = \tilde{A}(x) + \tilde{A}(y)$$

a

$\frac{A(b)}{mb}$ for some b

revised

$$\frac{d}{dx} \tilde{A}(x+y) = \frac{d\tilde{A}(x)}{dx}$$

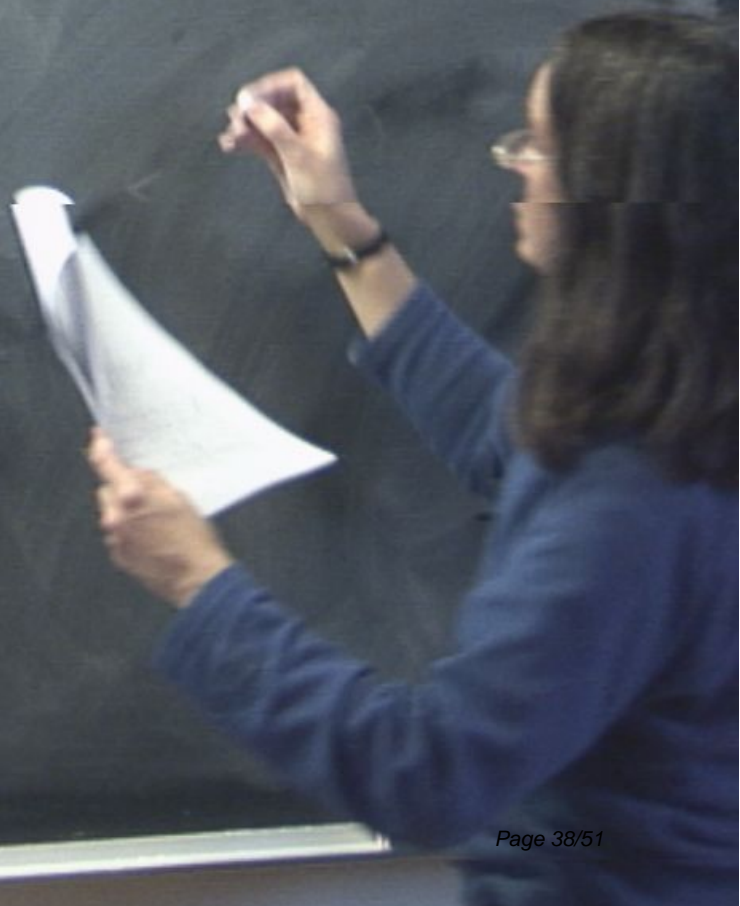
for all y

$$\Rightarrow \frac{d\tilde{A}}{dx}$$

$A(n) + A(n)$

$A(e^x)$

$\tilde{A}(x) + \tilde{A}(y)$



a.

$\frac{A(b)}{mb}$ for some b

homework

$$\frac{d}{dx} \tilde{A}(x+y) = \frac{d\tilde{A}(x)}{dx}$$

for all y

$\Rightarrow \frac{d\tilde{A}}{dx}$ is independent of x

$$A(n) + A(n)$$

$$\Rightarrow \tilde{A}(x) = Kx + b$$

$$\Rightarrow A(m) = K \ln m + b$$

$$A(e^x)$$

$$\tilde{A}(x) + \tilde{A}(y)$$

$$\Rightarrow A(a) = K \ln a$$

$$\left[K = \frac{A(b)}{\ln b} \text{ for some } b \right]$$

Alternative proof.

$$A(mn) = A(m) + A(n)$$

$$\text{let } x = \ln m$$

$$y = \ln n$$

$$\tilde{A}(x) = A(e^x)$$

$$\tilde{A}(x+y) = \tilde{A}(x) + \tilde{A}(y)$$

$$\frac{d}{dx} \tilde{A}(x+y) =$$

for a

$$\Rightarrow \frac{d\tilde{A}}{dx}$$

$$\Rightarrow \tilde{A}(x) =$$

$$\Rightarrow A(m) =$$

$$K \ln(mn) + b = K$$

$$\Rightarrow b$$

Now consider

p_i that are not all equal

consider p_i rational

$$\Rightarrow \underline{A(a) = \dots}$$

alternative procedure

$$A(mn) = \dots$$

$$\text{let } x = \ln$$

$$y = \ln$$

$$\tilde{A}(x)$$

$$\hat{A}(x+y)$$

Now consider

p_i that are not all equal

consider p_i rational

$$p_i = \frac{n_i}{\sum_1 n_i}$$

1 2 3 4 5 6 7 8 9 10

$$\Rightarrow \underline{A(a)} = \dots$$

alternative prob

$$A(mn) = \dots$$

$$\text{let } x = \ln$$

$$y = \ln$$

$$\tilde{A}(x)$$

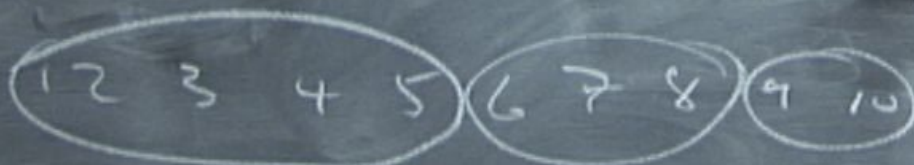
$$\hat{A}(x+y)$$

Now consider

p_i that are not all equal

consider p_i rational

$$p_i = \frac{n_i}{\sum_x n_x}$$



$$p_1 = 1/2$$

$$p_2 = 3/10 \quad p_3 = 1/5$$

$$\Rightarrow A(a) = \dots$$

alternative procedure

$$A(mn) = \dots$$

$$\text{let } x = \ln$$

$$y = \ln$$

$$\tilde{A}(x)$$

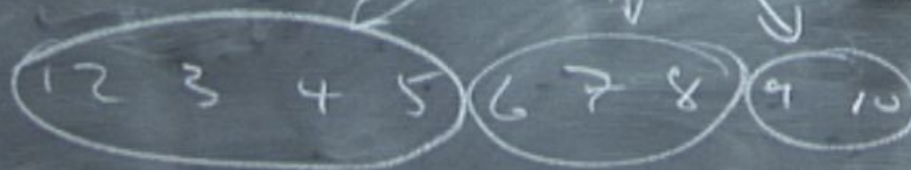
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$$p_2 = 3/10 \quad p_3 = 1/5$$

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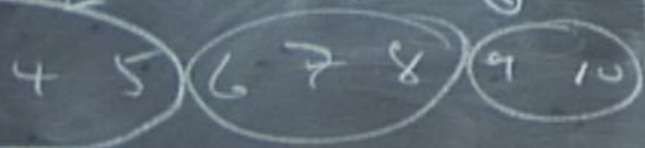
alternative prob

$$A(mn) = \dots$$

are not all equal

rational

$$\frac{n_i}{\sum_1 n_i}$$



$$P_1 = 1/2 \quad P_2 = 3/10 \quad P_3 = 1/5$$



want to find
 $S(p_1, \dots, p_n)$

want to find
 $S(p_1, \dots, p_n)$

$$K \ln \left(\sum_i n_i \right)$$

want to find
 $S(p_1, \dots, p_n)$

$$K \ln \left(\sum_i n_i \right) = S(p_1, \dots, p_n) + \sum$$

want to find
 $S(p_1, \dots, p_n)$

$$K \ln \left(\sum_i n_i \right) = S(p_1, \dots, p_n) + \sum_i p_i \ln n_i$$

at to find

(p_1, \dots, p_n)

$$H(\sum_i n_i) = S(p_1, \dots, p_n) + \sum_i p_i \ln n_i$$

$$S(p_1, \dots, p_n) = \ln(\sum_i n_i) (\sum_i p_i) - \sum_i p_i \ln n_i$$

nt. to find
(p_1, \dots, p_n)

$$m(\sum_i n_i) = S(p_1, \dots, p_n) + \sum_i p_i \ln n_i$$

$$S(p_1, \dots, p_n) = K \left[\ln(\sum_i n_i) (\sum_i p_i) - \sum_i p_i \ln n_i \right]$$
$$= -K \sum_i p_i \ln \left(\frac{n_i}{\sum_i n_i} \right) = -K \sum_i p_i \ln$$

nt. to find
(p_1, \dots, p_n)

$$m(\sum_i n_i) = S(p_1, \dots, p_n) + K \sum_i p_i \ln n_i$$

$$S(p_1, \dots, p_n) = K \left[\ln(\sum_i n_i) (\sum_i p_i) - \sum_i p_i \ln n_i \right]$$

$$= -K \sum_i p_i \ln \left(\frac{n_i}{\sum_i n_i} \right) = -K \sum_i p_i \ln$$