

Title: Scientific Computation (PHYS 608) - Lecture 13

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Abstract:

# Iterative Methods



# Iterative Methods

$$\underline{A} x = \underline{b}$$

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$$\underline{\underline{A}} \underline{x} = \underline{b}$$

$$\underline{\underline{A}} = \underline{\underline{Q}} + \underline{\underline{A}} - \underline{\underline{Q}}$$

# Iterative Methods

$$\textcircled{1} \quad \underline{\underline{A}} x = \underline{\underline{b}}$$

$$\underline{\underline{A}} = \underline{\underline{Q}} + \underline{\underline{A}} - \underline{\underline{Q}}$$

Then  $\textcircled{1}$

$$Qx + (A - Q)x = b$$

or  $Qx =$

# Iterative Methods

$$\textcircled{1} \quad \underline{\underline{A}} x = \underline{\underline{b}}$$

$$\underline{\underline{A}} = \underline{\underline{Q}} + \underline{\underline{A}} - \underline{\underline{Q}}$$

Then  $\textcircled{1}$

$$Qx + (A - Q)x = b$$

or

$$Qx = (Q - A)x + b$$

# Iterative Methods

$$\textcircled{1} \quad \underline{\underline{A}} x = \underline{\underline{b}}$$

$$\underline{\underline{A}} = \underline{\underline{Q}} + \underline{\underline{A}} - \underline{\underline{Q}}$$

Then  $\textcircled{1}$

$$\underline{\underline{Q}}x + (\underline{\underline{A}} - \underline{\underline{Q}})x = \underline{\underline{b}}$$

or

$$\underline{\underline{Q}}x = (\underline{\underline{Q}} - \underline{\underline{A}})x + \underline{\underline{b}}$$

# Iterative Methods

$$\textcircled{1} \quad \underline{\underline{A}} x = \underline{\underline{b}}$$

$$\underline{\underline{A}} = \underline{\underline{Q}} + \underline{\underline{A}} - \underline{\underline{Q}}$$

Then  $\textcircled{1}$

$$\underline{\underline{Q}}x + (\underline{\underline{A}} - \underline{\underline{Q}})x = \underline{\underline{b}}$$

or

$$\underline{\underline{Q}}x = (\underline{\underline{Q}} - \underline{\underline{A}})x + \underline{\underline{b}} \quad \leftarrow \text{Iterating}$$



$$\underline{Q} \underline{x}^{(k+1)} = (\underline{Q} - \underline{A}) \underline{x}^{(k)} + \underline{b}$$

We hope

$$\underline{Q} \underline{x}^* = (\underline{Q} - \underline{A}) \underline{x}^* + \underline{b}$$

$$\underline{\underline{Q}} \underline{\underline{x}}^{(k+1)} = (\underline{\underline{Q}} - \underline{\underline{A}}) \underline{\underline{x}}^{(k)} + \underline{\underline{b}}$$

↓  
We hope

$$\underline{\underline{Q}} \underline{\underline{x}}^* = (\underline{\underline{Q}} - \underline{\underline{A}}) \underline{\underline{x}}^* + \underline{\underline{b}} \iff \underline{\underline{A}} \underline{\underline{x}}^* = \underline{\underline{b}}$$

$$\underline{Q} \underline{x}^{(k+1)} = (\underline{Q} - \underline{A}) \underline{x}^{(k)} + \underline{b}$$

↓  
We hope

$$\underline{Q} \underline{x}^* = (\underline{Q} - \underline{A}) \underline{x}^* + \underline{b} \iff \underline{A} \underline{x}^* = \underline{b}$$

How should choose  $\underline{Q}$

$$\underline{Q} \underline{x}^{(k+1)} = \underbrace{(\underline{Q} - \underline{A}) \underline{x}^{(k)} + \underline{b}}_{\text{We hope}}$$

$$\underline{Q} \underline{x}^{(k+1)} = \underline{g}$$

$$\underline{Q} \underline{x}^* = (\underline{Q} - \underline{A}) \underline{x}^* + \underline{b} \iff \underline{A} \underline{x}^* = \underline{b}$$

How should choose  $\underline{Q}$

$$\underline{Q} \underline{x}^{(k+1)} = \underline{(Q - A)} \underline{x}^{(k)} + \underline{b}$$

We hope

$$\underline{Q} \underline{x}^{(k+1)} = \underline{y}$$

value for  $\underline{x}^{(k+1)}$

$$\underline{Q} \underline{x}^* = \underline{(Q - A)} \underline{x}^* + \underline{b} \iff \underline{A} \underline{x}^* = \underline{b}$$

How should choose  $\underline{Q}$

- So that  $\underline{Q} \underline{x}^{(k+1)} = \underline{y}$  can be solved easily
- So that things converge

$$\underline{Q} \underline{x}^{(k+1)} = \underline{(Q - A)} \underline{x}^{(k)} + \underline{b}$$

We hope

$$\underline{Q} \underline{x}^* = \underline{(Q - A)} \underline{x}^* + \underline{b}$$

$$\Leftrightarrow \underline{A} \underline{x}^* = \underline{b}$$

$$\underline{Q} \underline{x}^{(k+1)} = \underline{y}$$

↑ Solve for  $\underline{x}^{(k+1)}$

How should choose  $\underline{Q}$

- So that  $\underline{Q} \underline{x}^{(k+1)} = \underline{y}$  can be solved easily

- So that things converge

Q is diagonal

Q is diagonal

$$Q = \text{diag}(A)$$



Q is diagonal

$$Q = \text{diag}(A)$$

Jacobi

Q is diagonal

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Jacobi

Q is diagonal

$$Q = \text{diag}(A)$$

Jacobi iteration

Q is diagonal

$$Q = \text{diag}(A)$$

Jacobi iteration

Gauss-Seidel Iteration

Q is diagonal

$$Q = \text{diag}(A)$$

Jacobi iteration

Gauss-Seidel Iteration

Q = lower/upper triangular part of A

Q is diagonal

$$Q = \text{diag}(A)$$

Jacobi iteration

Gauss-Seidel Iteration

Q = lower/upper triangular part of A  
including the diagonal

# Analysis

$$\underline{x}^{(k+1)} = Q^{-1} \left( (Q - A)x^{(k)} + \underline{r} \right)$$
$$= Cx^{(k)} + \underline{g}$$

## Analysis

$$\underline{x}^{(k+1)} = Q^{-1} \left( (Q - A)x^{(k)} + \underline{r} \right)$$

$$= \underline{C} x^{(k)} + \underline{g}$$

$$\underline{C} = I - Q^{-1}A$$



## Analysis

$$\underline{x}^{(k+1)} = Q^{-1} \left( (Q - A)x^{(k)} + b \right)$$

$$= C x^{(k)} + g$$

$$g = Q^{-1}b$$

$$C = I - Q^{-1}A$$

## Analysis

$$\underline{x}^{(k+1)} = Q^{-1} \left( (Q - A)x^{(k)} + \underline{b} \right)$$

$$y = Q^{-1} \underline{b}$$

$$= \underline{x}^{(k)} + y$$

$$= I - Q^{-1}A$$

$$e^{(k)} = x^{(k)} - x$$

$$= (I - Q^{-1}A)x^{(k+1)} + Q^{-1}\underline{b} - x$$

# Analysis

$$\underline{x}^{(k+1)} = Q^{-1} \left( (Q - A)x^{(k)} + b \right)$$

$$y = Q^{-1}b$$

$$= \left( x^{(k)} + y \right)$$

$$\left( = I - Q^{-1}A \right)$$

$$e^{(k)} = x^{(k)} - x$$

$$= (I - Q^{-1}A)x^{(k+1)} + Q^{-1}b - x$$

$$= (I - Q^{-1}A)x^{(k+1)} + (Q^{-1}A - I)x$$

$$= (I - Q^{-1}A)$$

# Analysis

$$\underline{x}^{(k+1)} = Q^{-1} \left( (Q - A)x^{(k)} + b \right)$$

$$y = Q^{-1}b$$

$$= \left( x^{(k)} + y \right)$$

$$\left( = I - Q^{-1}A \right)$$

$$e^{(k)} = x^{(k)} - x$$

$$= (I - Q^{-1}A)x^{(k-1)} + Q^{-1}b - x$$

$$= (I - Q^{-1}A)x^{(k-1)} + (Q^{-1}A - I)x$$

$$= (I - Q^{-1}A)(x^{(k-1)} - x) =$$

## Analysis

$$\underline{x}^{(k+1)} = Q^{-1} \left( (Q - A)x^{(k)} + b \right)$$

$$y = Q^{-1}b$$

$$= \left( x^{(k)} + y \right)$$

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$$e^{(k)} = x^{(k)} - x$$

$$= (I - Q^{-1}A)x^{(k+1)} + Q^{-1}b - x$$

$$= (I - Q^{-1}A)x^{(k)} + (Q^{-1}A - I)x$$

$$= (I - Q^{-1}A)(x^{(k)} - x) = (I - Q^{-1}A)e^{(k)}$$

## Analysis

$$\underline{x}^{(k+1)} = Q^{-1} \left( (Q - A)x^{(k)} + b \right) \quad y = Q^{-1}b$$

$$= \left( x^{(k)} + y \right) \quad (= I - Q^{-1}A)$$

$$e^{(k)} = x^{(k)} - x$$

$$= (I - Q^{-1}A)x^{(k-1)} + Q^{-1}b - x$$

$$= (I - Q^{-1}A)x^{(k-1)} + (Q^{-1}A - I)x$$

$$= (I - Q^{-1}A)(x^{(k-1)} - x) = (I - Q^{-1}A)e^{(k-1)}$$

$$\underline{Q} \underline{x}^{(k+1)} = \underline{(Q - A)} \underline{x}^{(k)} + \underline{b}$$

We hope

$$\underline{Q} \underline{x}^{(k+1)} = \underline{g}$$

↑ Solve for  $\underline{x}^{(k+1)}$

$$\underline{Q} \underline{x}^* = \underline{(Q - A)} \underline{x}^* + \underline{b} \iff \underline{A} \underline{x}^* = \underline{b}$$

How should choose  $\underline{Q}$

- So that  $\underline{Q} \underline{x}^{(k+1)} = \underline{g}$  can be solved easily

- So that things converge  
 $\underline{Q}^{-1} \underline{A}$  should be "close" to  $\underline{I}$

# Convergence Theorem

For

$$Qx^{(k+1)} = (Q-A)x^{(k)} + G$$

l



# Convergence Theorem

For

$$Qx^{(k+1)} = (Q-A)x^{(k)} + b$$

to independent of  $x^{(0)}$

# Convergence Theorem

For

$$Qx^{(k+1)} = (Q-A)x^{(k)} + b$$

to independent of  $x^{(0)}$

it is a necessary and sufficient condition that all eigenvalues of  $I - Q^{-1}A$

# Convergence Theorem

For

$$Qx^{(k+1)} = (Q-A)x^{(k)} + b$$

to independent of  $x^{(0)}$

it is a necessary and sufficient condition that all eigenvalues of  $I - Q^{-1}A$  lie in the open unit circle

$$\rho(I - Q^{-1}A) < 1$$

## Analysis

$$\underline{x}^{(k+1)} = Q^{-1} \left( (Q - A)x^{(k)} + b \right)$$

$$y = Q^{-1}b$$

$$= \left( x^{(k)} + y \right)$$

$$\left( = I - Q^{-1}A \right)$$

$$e^{(k)} = x^{(k)} - x$$

$$= (I - Q^{-1}A)x^{(k-1)} + Q^{-1}b - x$$

$$= (I - Q^{-1}A)x^{(k-1)} + (Q^{-1}A - I)x$$

$$= (I - Q^{-1}A)(x^{(k-1)} - x) = (I - Q^{-1}A)e^{(k-1)}$$

# Convergence Theorem

For

$$Qx^{(k+1)} = (Q-A)x^{(k)} + b$$

to independent of  $x^{(0)}$

it is a necessary and sufficient condition that all eigenvalues of  $I - Q^{-1}A$  lie in the open

unit circle

$$\rho(I - Q^{-1}A) < 1$$

Spectral radius theorem

# Diagonally dominance

$$|a_{ii}| > \sum_{\substack{j=1 \\ j \neq i}} |a_{ij}|$$

## Diagonally dominance

$$|a_{ii}| > \sum_{\substack{j=1 \\ j \neq i}} |a_{ij}|$$

If  $A$  is diagonally dominant  
then Jacobi and G-S converge  
for any  $x^{(0)}$

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6 web pages



3

6

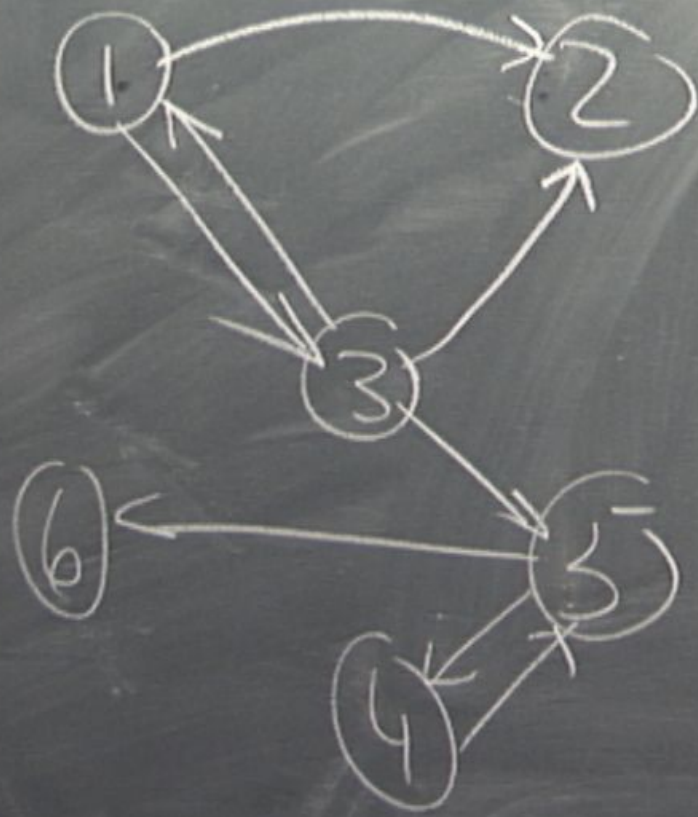
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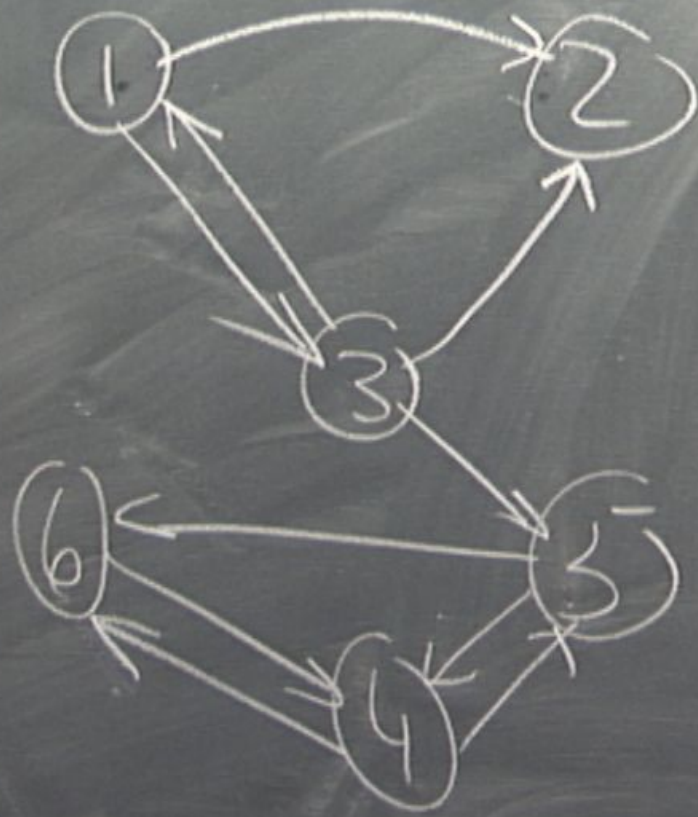
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→ Page rank

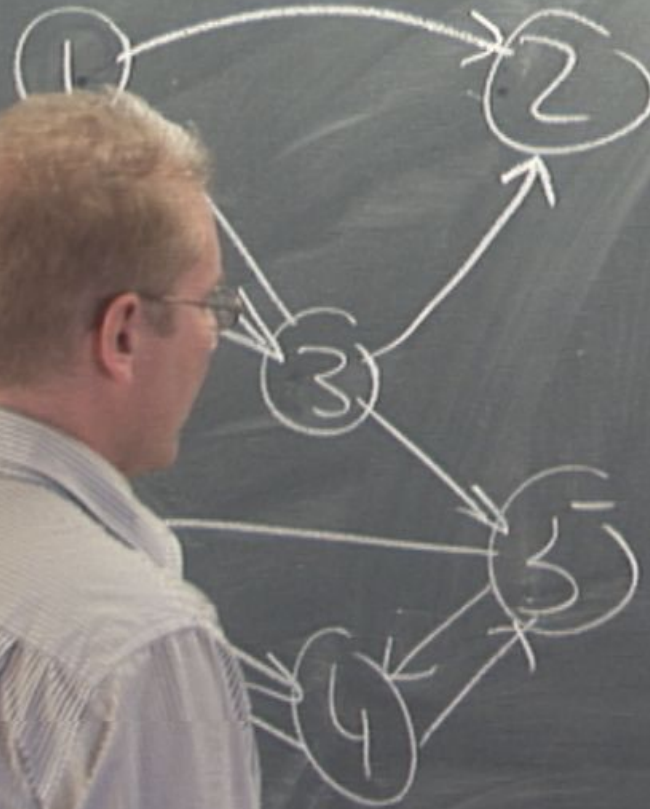
→ Page rank

$$r(P_i) = \sum_{\substack{P_j \\ \text{pointing to } P_i}} \bar{\phantom{r}}$$

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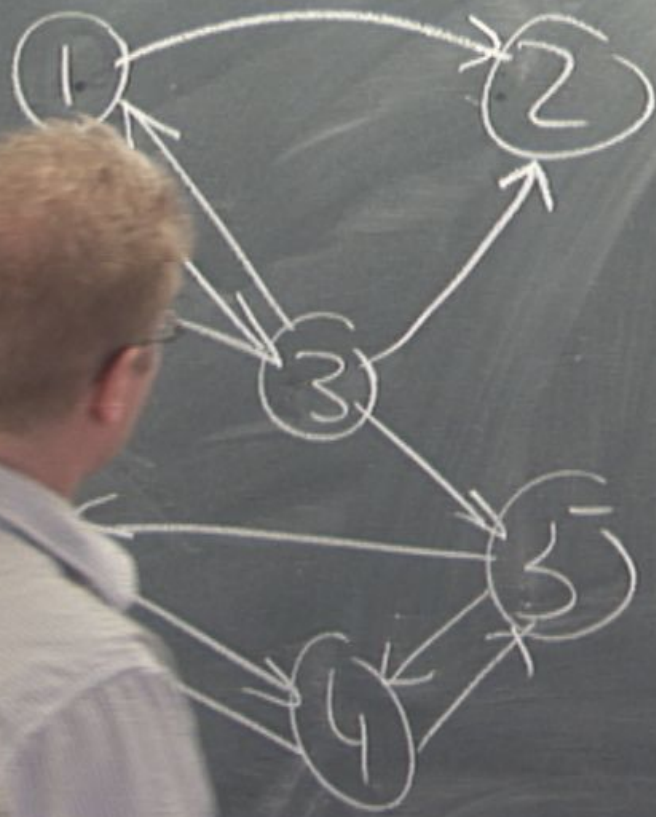


$|P_i| = \#$   
of outgoing  
links from

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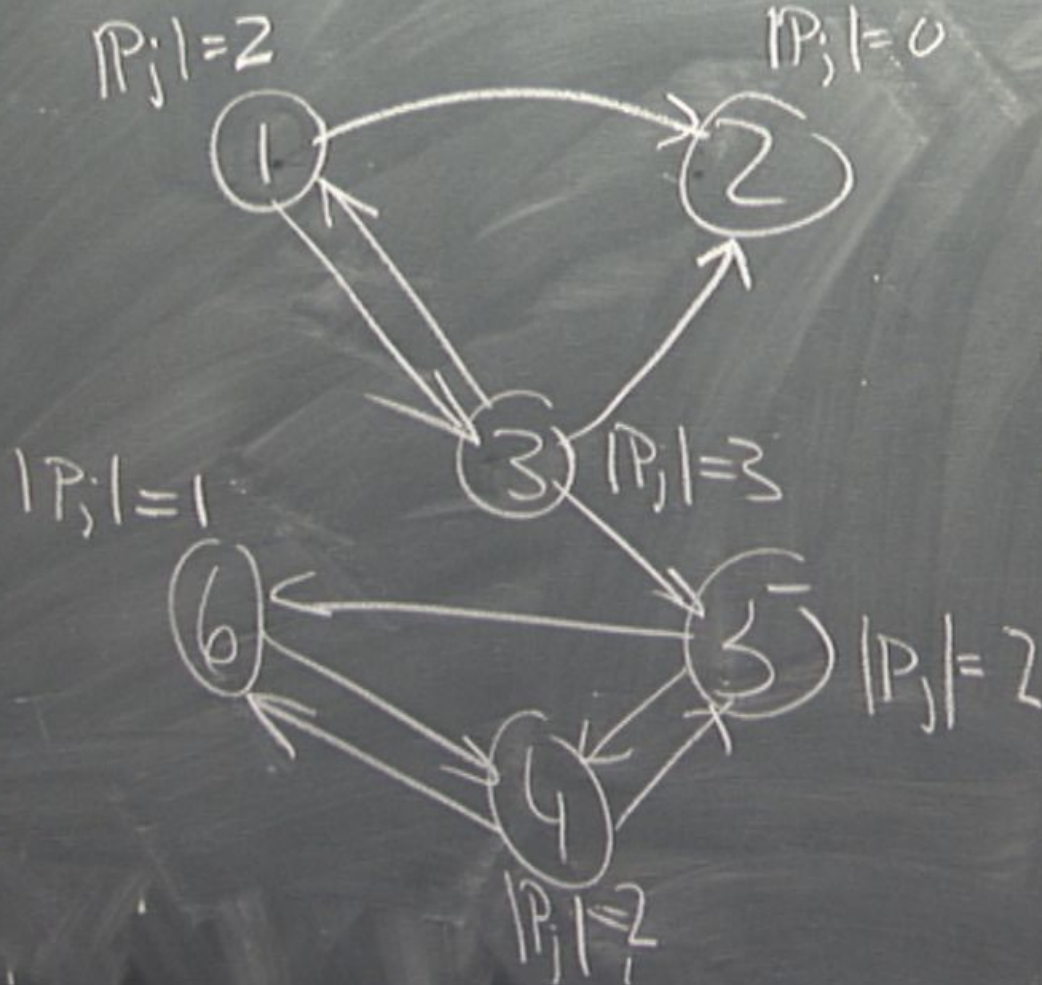
$|P_i| = \#$   
of outgoing  
links from  
 $P_i$



1998

Larry Page / S. Brin

6 web pages



$IP_j = \#$   
of outgoing  
links from  
 $P_j$

→ Page rank

$$r(P_i) = \sum_{\substack{P_j \\ \text{pointing to } P_i}} \frac{r(P_j)}{|P_j|}$$

$$\begin{pmatrix} 0 & 1/2 & 1/2 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \end{pmatrix}$$

→ Page rank

$$r(P_i) = \sum_{\substack{P_j \\ \text{Painting to } P_i}} \frac{r(P_j)}{|P_j|}$$

$$\begin{pmatrix} 0 & 1/2 & 1/2 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 1/3 & 1/3 & 0 & 0 & 1/3 & 0 \\ 0 & 0 & 0 & 0 & 1/2 & 1/2 \\ 0 & 0 & 0 & 1/2 & 0 & 1/2 \\ 0 & 0 & 0 & 1 & 0 & 0 \end{pmatrix}$$

→ Page rank

Solve by iteration

$$r(P_i) = \sum_{\substack{P_j \\ \text{painting to } P_i}} \frac{r(P_j)}{|P_j|}$$

$r(P)$

=

$$\begin{pmatrix} 0 & 1/2 & 1/2 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 1/3 & 1/3 & 0 & 0 & 1/3 & 0 \\ 0 & 0 & 0 & 0 & 1/2 & 1/2 \\ 0 & 0 & 0 & 1/2 & 0 & 1/2 \\ 0 & 0 & 0 & 1 & 0 & 0 \end{pmatrix}$$

$r(P_i)$

Suppose I had



Suppose I had



$$\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$$

Suppose I had



$$\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \begin{pmatrix} 0 \\ 1 \end{pmatrix} \begin{pmatrix} 1 \\ 0 \end{pmatrix}$$

Suppose I had limit cycle



$$\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \begin{pmatrix} 0 \\ 1 \end{pmatrix} \begin{pmatrix} 1 \\ 0 \end{pmatrix}$$



Suppose I had limit cycle



$$\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \begin{pmatrix} 0 \\ 1 \end{pmatrix} \begin{pmatrix} 1 \\ 0 \end{pmatrix}$$

$$\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} 1 \\ 1 \end{pmatrix} = \begin{pmatrix} 1 \\ 1 \end{pmatrix}$$

Suppose I had

limit cycle



$$\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \begin{pmatrix} 0 \\ 1 \end{pmatrix} \begin{pmatrix} 1 \\ 0 \end{pmatrix}$$

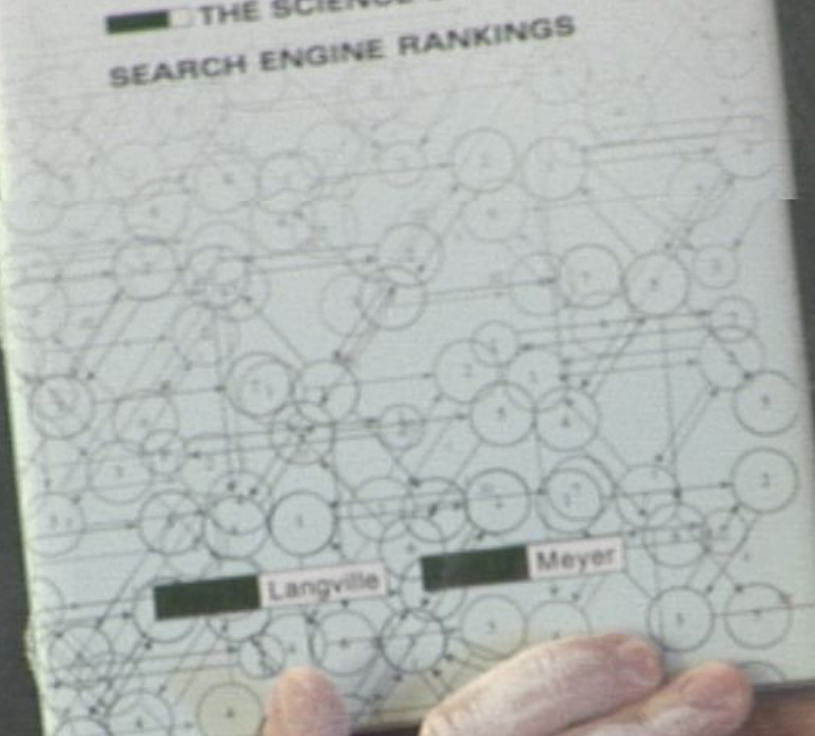
$$\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} 1 \\ 1 \end{pmatrix} = \begin{pmatrix} 1 \\ 1 \end{pmatrix}$$

↑ fixed point

Google's

PageRank and Beyond

THE SCIENCE OF  
SEARCH ENGINE RANKINGS



Langville

Meyer

$r(P_i)$

# Monte Carlo Simulations

Monte Carlo Simulations

Random Numbers

# Monte Carlo Simulations

## Random Numbers

— What is random?

# Monte Carlo Simulations

## Random Numbers

$U(0, 1)$

— What is random?

# Monte Carlo Simulations

## Random Numbers

— What is random?

$U(0, 1)$

↳ Usually  
Integers between  
 $0 \leq I \leq I_{\max}$



# Monte Carlo Simulations

## Random Numbers

$U(0, 1)$

— What is random?

↳ Usually  
Integers between

— Pseudo-Random Numbers  $\% I_{max}$

# Monte Carlo Simulations

## Random Numbers

$U(0, 1)$

— What is random?

↳ Usually  
Integers between

— Pseudo-Random Numbers  
deterministic  $\% I_{\max}$   
 $I^{(k+1)}$   
from  $I^{(k)}$

# Monte Carlo Simulations

## Random Numbers

$U(0, 1)$

— What is random?

↳ Usually  
Integers between

— Pseudo-Random Numbers %  $I_{max}$   
deterministic

$I^{(k+1)}$

from  $I^{(k)}$

RNG we build in correlation

→

$$I^{(12)} = \underset{31}{1010111011} \dots 10$$

→

$$I^{(k)} = \underset{31}{1010111011} \dots 10$$

$$I^{(k+1)} = 0011101 \dots 0$$

→

$$I^{(k)} = 1010111011 \dots$$

31

$$I^{(k+1)} = 0011101 \dots$$

1  
0  
← bits changed  
0

→

$$I^{(k)} = \frac{1010111011 \dots}{31}$$

$$I^{(k+1)} = 0011101 \dots$$

1  
0

← bits  
changed  
50%

→

$$I^{(k)} = \frac{1010111011 \dots}{31}$$

$$I^{(k+1)} = 0011101 \dots$$

1  
0  
← bits changed 50%

Statistical tests



$$I^{(k)} = 1010111011 \dots$$

31 29 26

$$I^{(k+1)} = 0011101 \dots$$

← bits changed 50%

Statistical tests

$$I^{(k)} = 1010111011 \dots$$

31    29    26

$$I^{(k+1)} = 0011101 \dots$$

← bits changed 50%

### Statistical tests

$$I^{(k)} \quad |$$

31

$$101000$$

$$I^{(k+1)}$$

$$I^{(k)} = 1010111011 \dots$$

31    29    26

$$I^{(k+1)} = 0011101 \dots$$

← bits changed 50%

### Statistical tests

$$I^{(k)} \quad 1 \quad 31$$

$$I^{(k+1)} \quad 0 \quad 31$$

$$1010000$$

$$10000$$

newer change

# Monte Carlo Simulations

PRB 69  
3382 (1992)

Monte Carlo Simulations

PRB 69  
3382 (1999)

Linear Congruential Generator

# Monte Carlo Simulations

PRB 69  
3382 (1998)

## Linear Congruential Generator

$$I^{(k+1)} = aI^{(k)} + c$$

# Monte Carlo Simulations

PRB 69  
3382 (1990)

## Linear Congruential Generator

$$I^{(k+1)} = aI^{(k)} + C \pmod{m}$$

# Monte Carlo Simulations

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## Linear Congruential Generator

$$I^{(k+1)} = aI^{(k)} + c \pmod{m}$$

ANSI C standard



# Monte Carlo Simulations

PRB 69  
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## Linear Congruential Generator

$$I^{(k+1)} = aI^{(k)} + c \pmod{m}$$

ANSI C standard

$$a = 1103515245 \quad c =$$

# Monte Carlo Simulations

PRB 69  
3382 (1990)

## Linear Congruential Generator

$$I^{(k+1)} = aI^{(k)} + c \pmod{m}$$

ANSI C Standard

$$a = 1103515245$$

$$c = 12345$$

$$m = 2^{32}$$

# Monte Carlo Simulations

PRB 69  
3382 (1990)

## Linear Congruential Generator

$$I^{(k+1)} = aI^{(k)} + C \pmod{m}$$

ANSI C Standard

$$a = 1103515245$$

$$C = 12345$$

$$m = 2^{32}$$

Enb's RNG

$$a=3$$

$$c=2$$

$$m=32$$

Einh's RNG

$$a=3$$

$$c=2$$

$$m=32$$

$$I^{k+1} = 3 I^k + 2 \pmod{32}$$

Euler's RNG

$$a=3$$

$$c=2$$

$$m=32$$

$$I^{k+1} = 3 I^k + 2 \pmod{32}$$

$$1 \rightarrow 5 \rightarrow$$

Enb's RNG

$$a=3$$

$$c=2$$

$$m=32$$

$$I^{k+1} = 3 I^k + 2 \pmod{32}$$

1 → 5 → 17 ↘ 21 ↘ 1

Euler's RNG

$$a=3$$

$$c=2$$

$$m=32$$

$$I^{k+1} = 3 I^k + 2 \pmod{32}$$

1  $\rightarrow$  5  $\rightarrow$  17  $\rightarrow$  21  $\rightarrow$  1

cycle small !



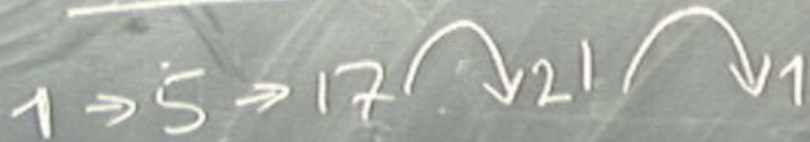
Enb's RNG

$a=3$

$c=2$

$m=32$

$I^{k+1} = 3I^k + 2 \pmod{32}$



cycle small

cycle length 4

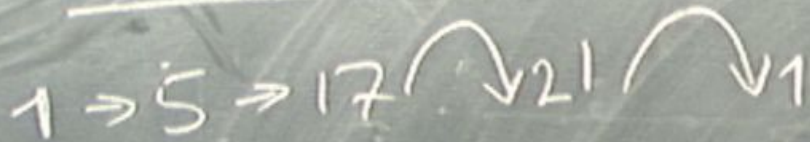
Enb's RNG

$a=3$

$C=2$

$m=32$

$I^{k+1} = 3I^k + 2 \pmod{32}$



cycle small

cycle length 4

$P \leq m$

Euler's RNG

$a=3$

$C=2$

$m=32$

$I^{k+1} = 3I^k + 2 \pmod{32}$

$1 \rightarrow 5 \rightarrow 17 \rightarrow 21 \rightarrow 1$

cycle small

cycle length 4

$2 \rightarrow 8 \rightarrow 26 \rightarrow 16 \rightarrow 18 \rightarrow 24 \rightarrow 10 \rightarrow 0 \rightarrow 2$

$P \leq m$

# Monte Carlo Simulations

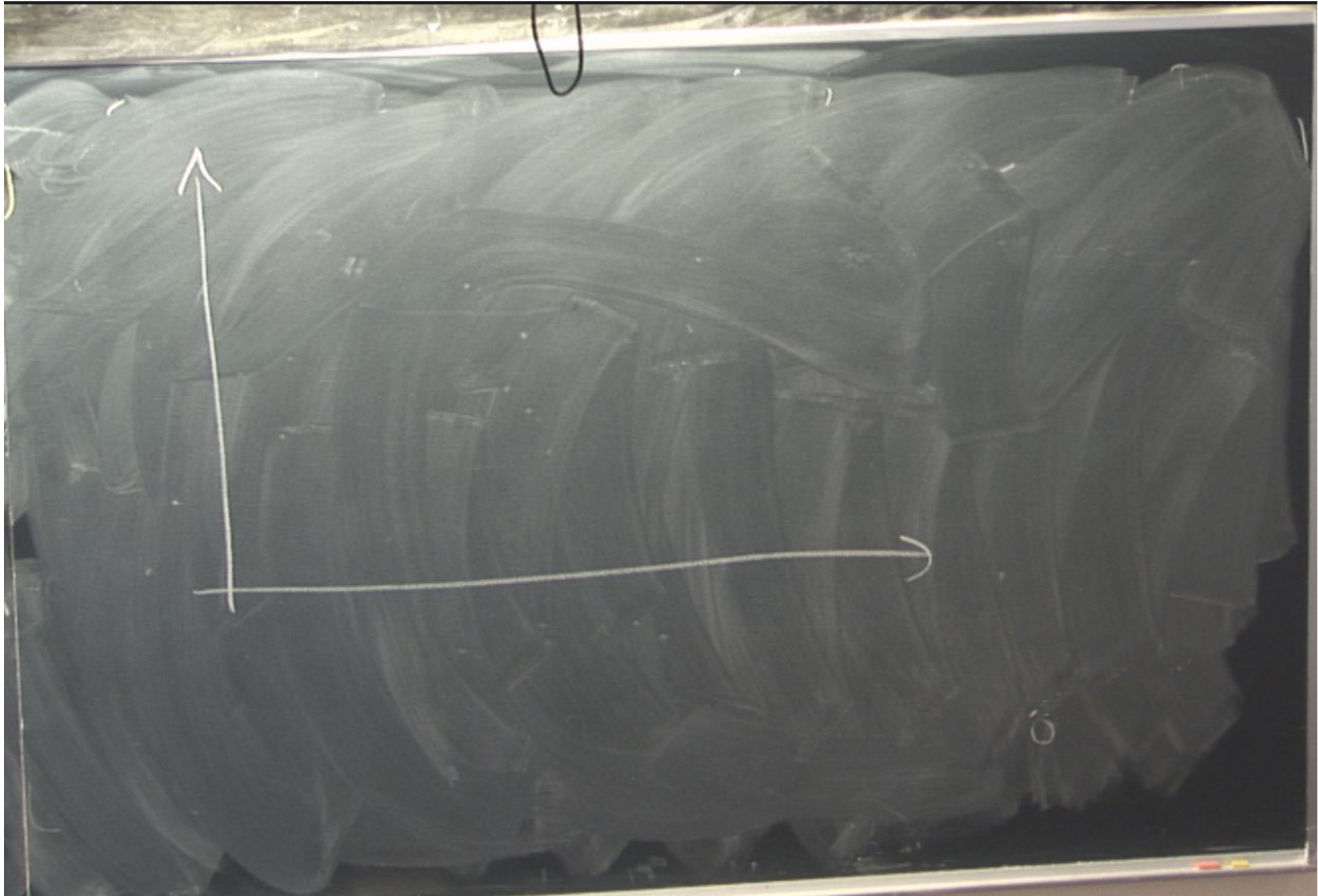
PRB 69  
3382 (1999)

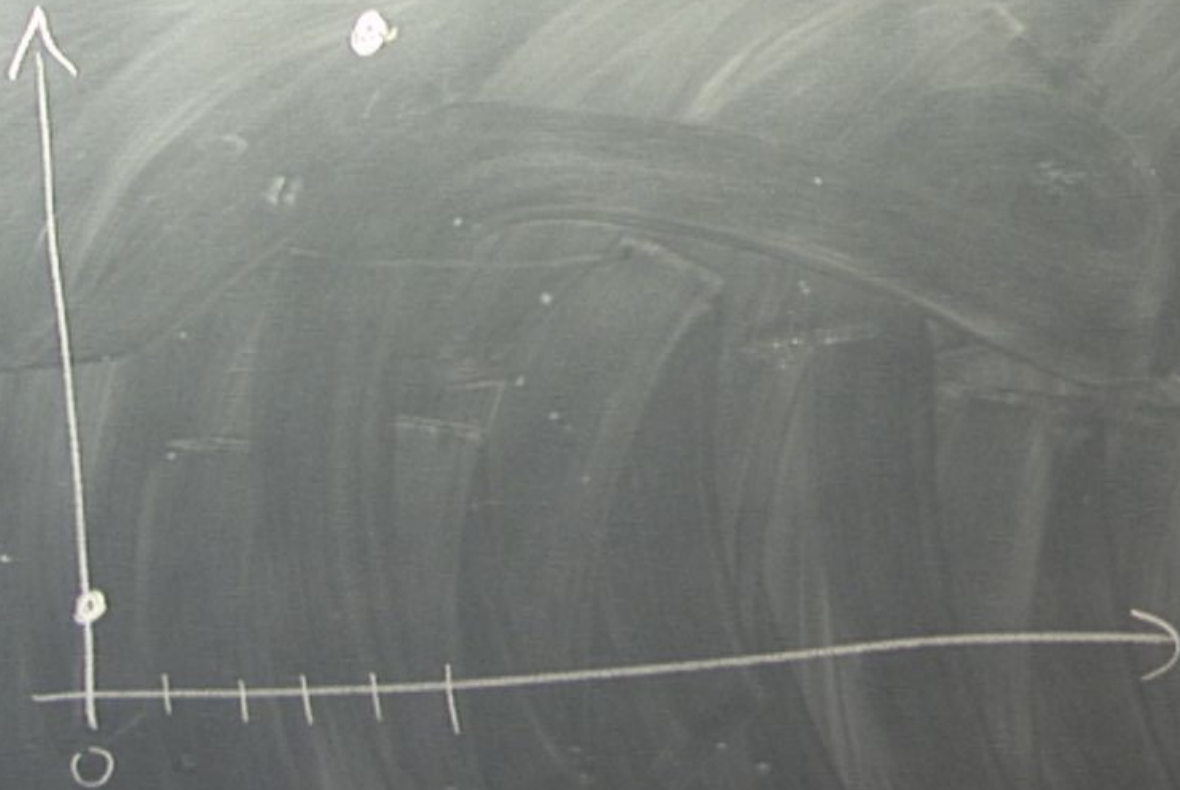
$I_1$	$I_2$
0	2
1	5
2	8
3	11
4	14
5	17
6	20
7	23
8	26
9	29
10	
11	

# Monte Carlo Simulations

PRB 69  
3382 (1999)

$I_1$	$I_2$
0	2
1	5
2	8
3	11
4	14
5	17
6	20
7	23
8	26
9	29
10	0
11	







$$y = 3x + 2$$





$$y = 3x + 2 \pmod{32}$$

3 lines in the 2D plane

ANSI C  $m = 2^{36}$

$(r_1, r_2, r_3)$  to put points in  $\mathbb{R}^3$

PRB 69

3382 (1990)

ANSI C  $m = 2^{36}$

$(r_1, r_2, r_3)$  to put points in  $\mathbb{R}^3$

1600 planes

PRB 69  
3382 (199)

ANSI C

$$m = 2^{36}$$

$(r_1, r_2, r_3)$  to put points in  $\mathbb{R}^3$

1600 planes

Fibonacci Generator

PRB 69

3382 (1998)

ANSI C

$$m = 2^{36}$$

$(r_1, r_2, r_3)$  to put points in  $\mathbb{R}^3$

1600 planes

PRB 69

3382 (1998)

Fibonacci Generator

$$I^{(k)} = I^{k-1} + I^{k-2} \quad \text{mod } c$$

ANSI C  $m = 2^{36}$

$(r_1, r_2, r_3)$  to put points in  $\mathbb{R}^3$

1600 planes

PRB 69  
3382 (1999)

## Fibonacci Generator

$$I^{(k)} = I^{(k-1)} + I^{(k-2)} \quad \text{mod } c$$

Lagged Fibonacci

$$I^{(k)} = I^{(k-r)} + I^{(k-s)}$$

ANSI C  $m = 2^{36}$

$(r_1, r_2, r_3)$  to put points in  $\mathbb{R}^3$   
1600 planes

PRB 69  
3382 (199)

## Fibonacci Generator

$$I^{(k)} = I^{(k-1)} + I^{(k-2)} \quad \text{mod } c$$

Lagged Fibonacci

$$I^{(k)} = I^{(k-r)} + I^{(k-s)} \quad \text{mod } c$$



$$y = 3x + 2 \pmod{32}$$

DIEHARD

3 lines in the 2D plane