Title: Adiabatic Gate Teleportation and Topological Quantum Computing

Date: Sep 14, 2009 03:00 PM

URL: http://pirsa.org/09090004

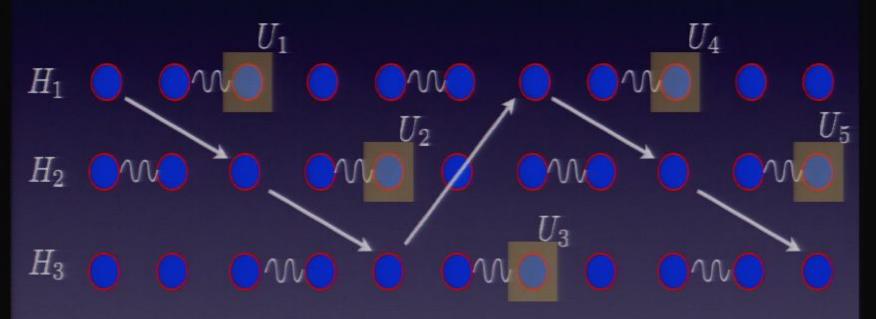
Abstract: TBA

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# Adiabatic Gate Teleportation & Its Applications

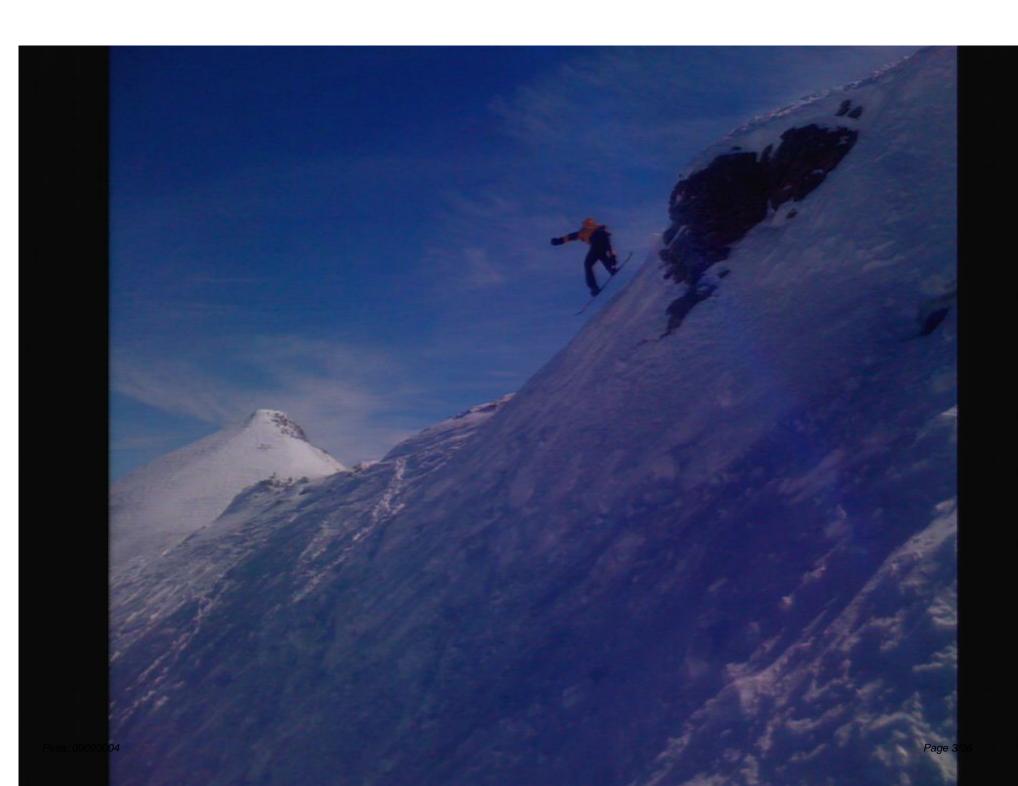
Dave Bacon

Department of Computer Science & Engineering University of Washington, Seattle, WA USA

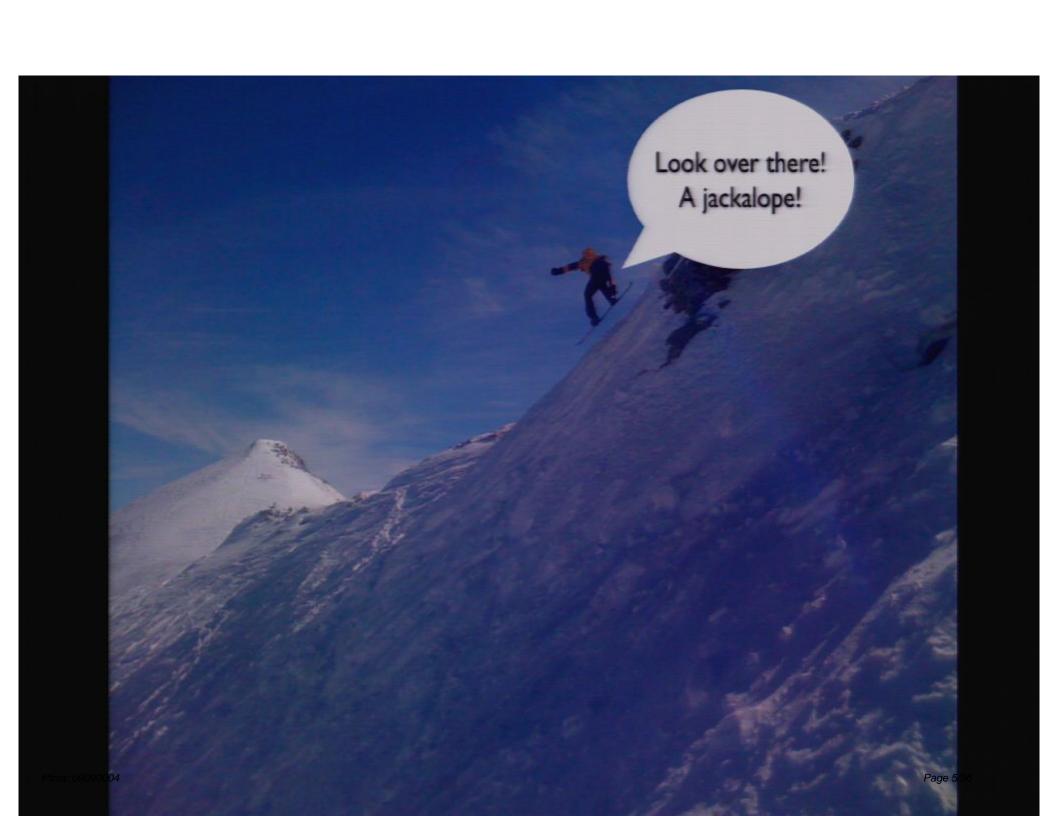


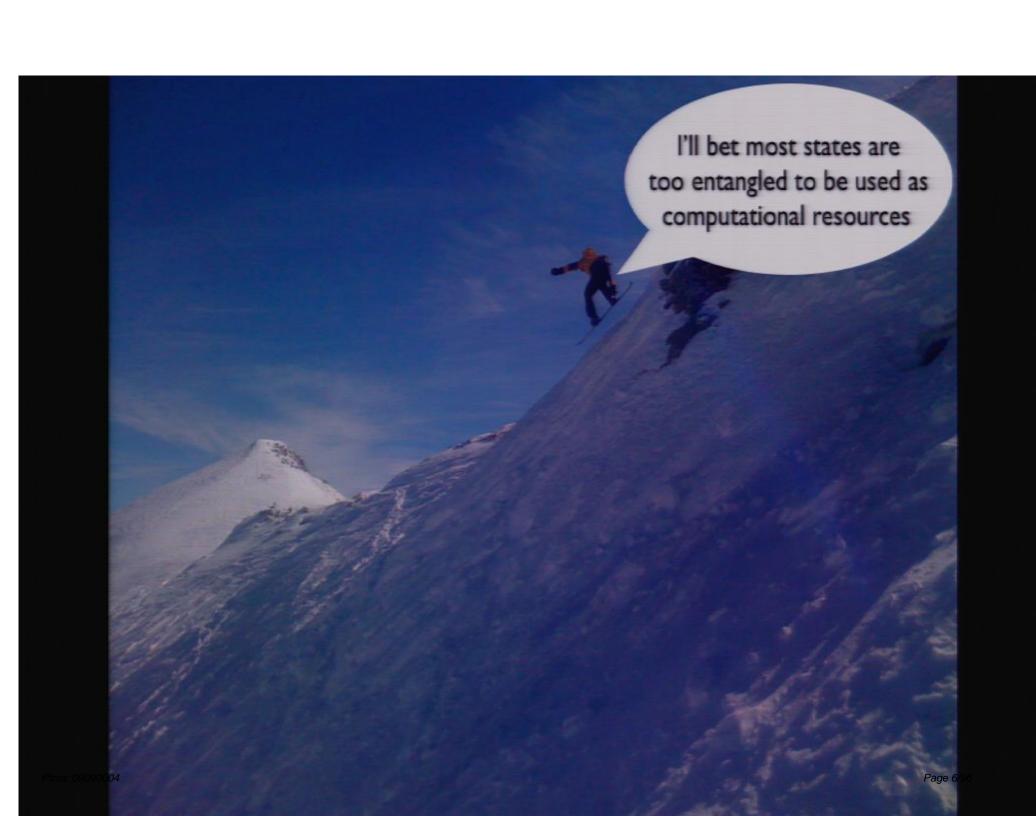
[joint work with Steve Flammia (Perimeter), Alice Neels (Washington), Andrew Landahl (Sandia Labs)]

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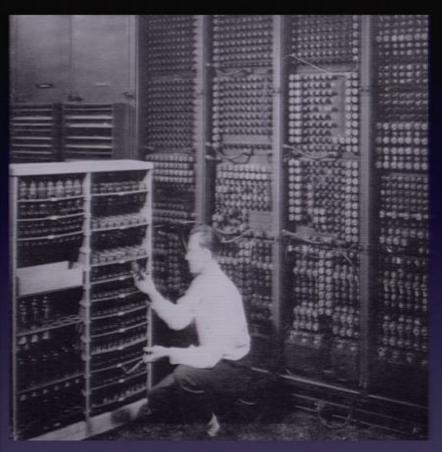






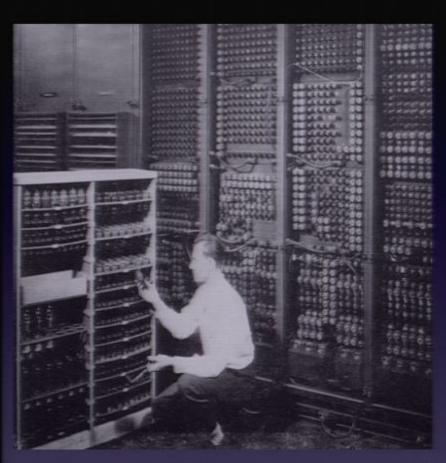


#### ENIAC, 1946



- 17,468 vacuum tubes
- 7,200 crystal diodes
- 1,500 relays
- 70,000 resistors
- 10,000 capacitors
- ~5,000,000 hand-soldered joints

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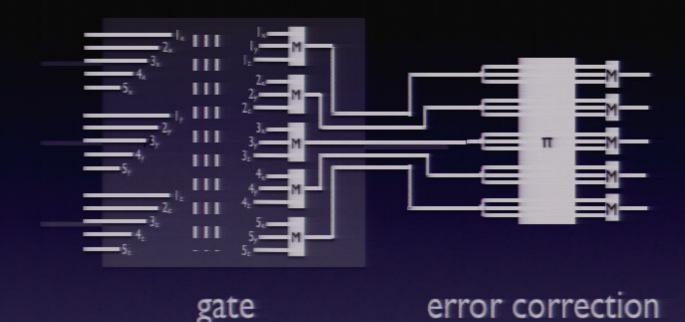
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Could we build a computer out of unreliable components?

J. von Neumann, "Probabilistic Logics and the Synthesis of Reliable Organism from Unreliable Components" 1956

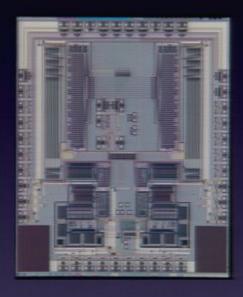
#### Von Neumann's Solution



Theorem (von Neumann): A circuit of g gates can be made to fail with probability  $\epsilon$  using O(g log (g/ $\epsilon$ )) gates which fail with probability p, assuming p<pth>pth (vN estimated pth ~ I%)

# Engineering? We Don't Need Your Stickin' Engineering





There are distinct physical reasons why robust classical computation is possible

#### Kitaev's Idea



"Magnetism arise from spins of individual atoms. Each spin is quite sensitive to thermal fluctuations. But the spins interact with each other and tend to be oriented in the same direction. If some spin flips to the opposite direction, the interaction forces it to flip back to the direction of other spins. This process is quite similar to the standard error correction procedure for the repetition code. We may say that errors are being corrected at the physical level. Can we propose something similar in the quantum case? Yes, but it is not so simple."

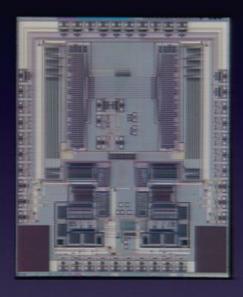
Alexei Kitaev

Kitaev, "Fault-tolerant quantum computation by anyons" (1997, published 2003)

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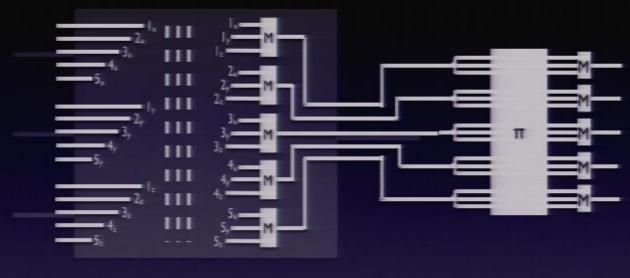
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gate

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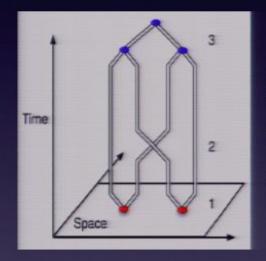
# Topological Quantum Computing

non-abelian Anyons used to build a universal quantum computer

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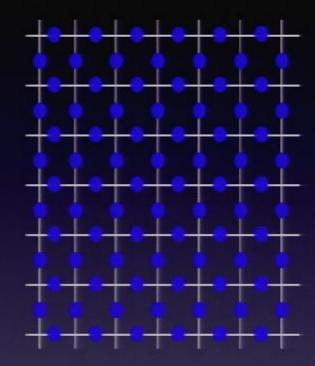
$$H = \begin{bmatrix} \frac{1}{\sqrt{2}} & \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} & -\frac{1}{\sqrt{2}} \end{bmatrix}$$



quantum computing using representations of braid group

Kitaev, "Fault-tolerant quantum computation by anyons" (1997, published 2003)
Freedman, "Quantum computation and the localization of modular functors" (2001) 2009 22

# Kitaev's Toric System



periodic boundary conditions (torus)

plaquette operator

$$S_p = \prod_{e \in \delta P} Z_e$$

vertex operator

$$S_v = \prod_{e|v \in e} X_e$$
$$[S_p, S_v] = 0 \quad S_p^2 = S_v^2 = I$$

Kitaev Hamiltonian:

$$H = -\frac{\Delta}{4} \left[ \sum_{p} S_{p} + \sum_{v} S_{v} \right]$$

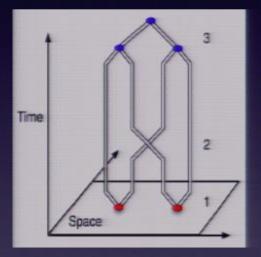
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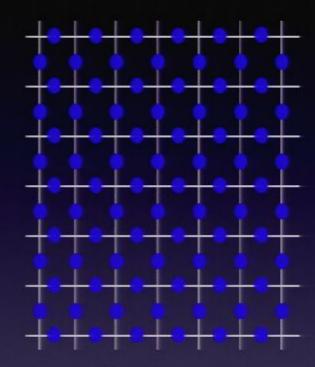
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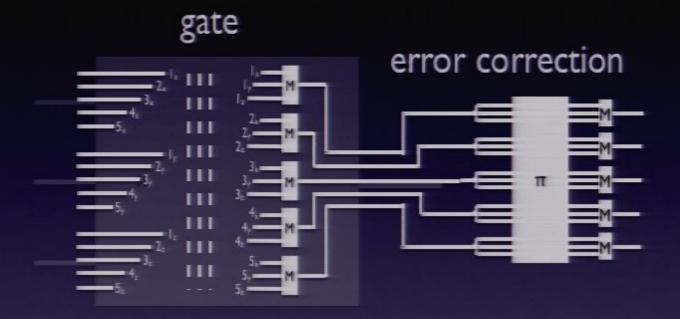
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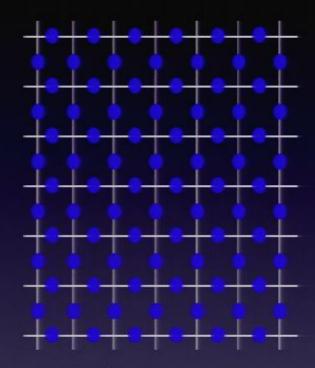
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Theorem (von Neumann): A circuit of g gates can be made to fail with probability  $\varepsilon$  using O(g log (g/ $\varepsilon$ )) gates which fail with probability p, assuming p<p<sub>th</sub> (vN estimated p<sub>th</sub> ~ I%)

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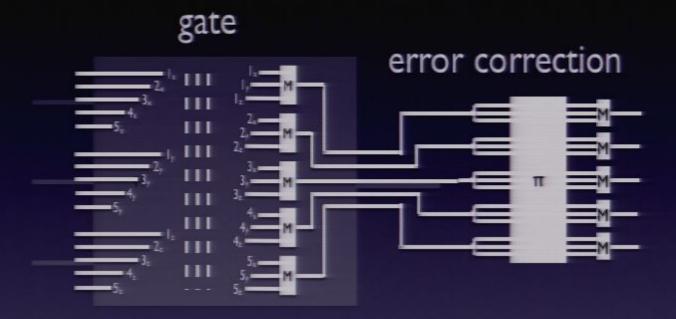
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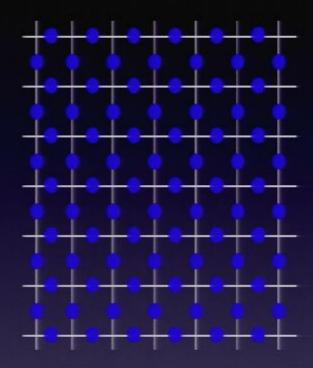
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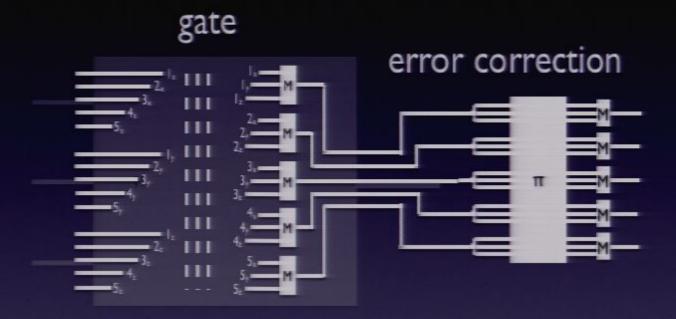
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### <u>Controversy</u>

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If we use this definition of fault-tolerant, then I believe the answer is NO!

Theorem (von Neumann): A circuit of g gates can be made to fail with probability  $\varepsilon$  using O(g log (g/ $\varepsilon$ )) gates which fail with probability p, assuming p<p<sub>th</sub> (vN estimated p<sub>th</sub> ~ I%)

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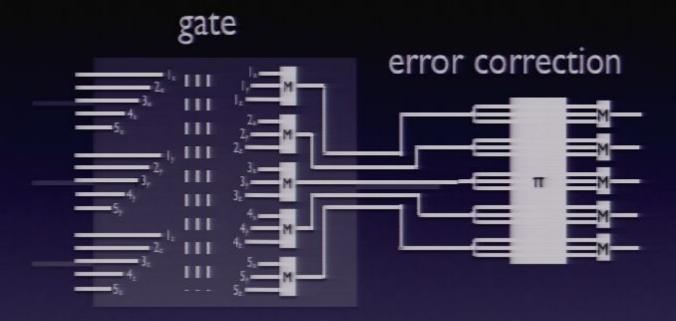
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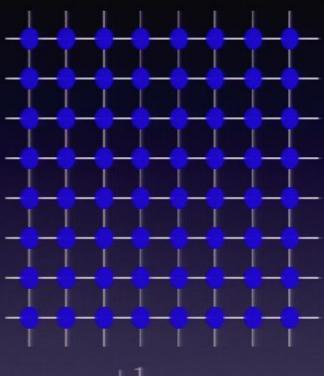
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# Cartoon Guide to Classical Physical Fault-Tolerance



) +. -:

spin

 $s_i \in \{+1, -1\}$ 

Energy

$$E = -J \sum_{\langle i,j \rangle} s_i s_j$$

neighbor

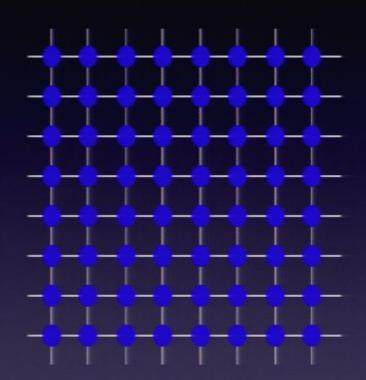
total magnetization

$$M = \sum_{i} s_{i}$$

"0" as 
$$M > N_0$$

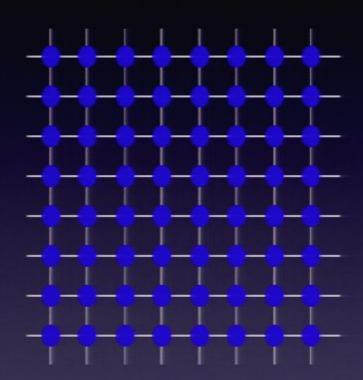
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Simple noise model



Simple noise model

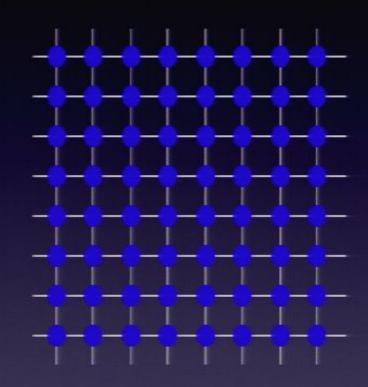
I. Pick spin at random

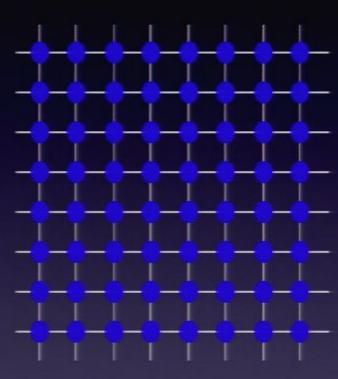


Simple noise model

1. Pick spin at random

2. If flipping spin would decrease energy a. flip spin





"bare" noise time step  $N_r$  runs of 1-3

Simple noise model

1. Pick spin at random

2. If flipping spin would decrease energy a. flip spin

If flipping spin would increase energy:

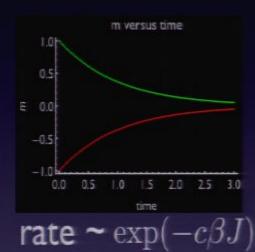
a. flip spin with probability

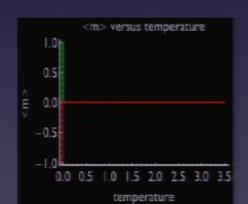
$$\exp(-\beta \Delta E)$$
  $\beta = T^{-P_{\text{alge 40}}}$ 

$$\beta = T^{-1}$$

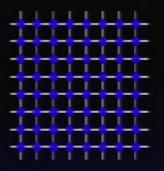
# Memory







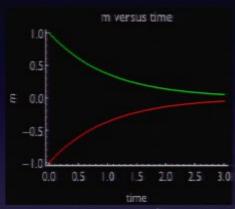
2D



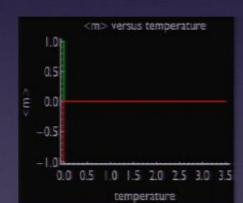
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# Memory

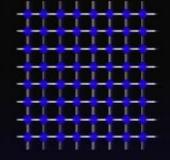






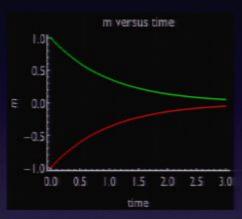


2D

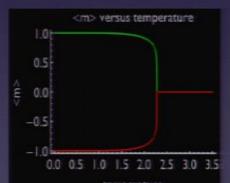




$$T < T_C$$



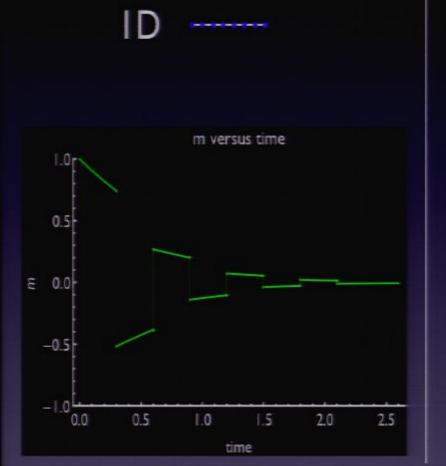
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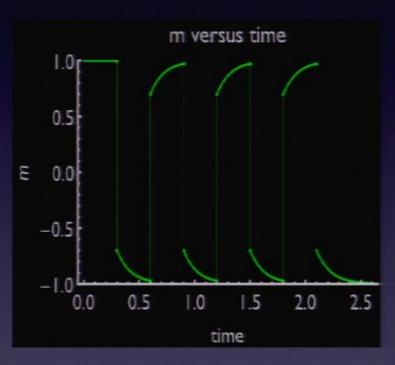
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### Fault-Tolerance

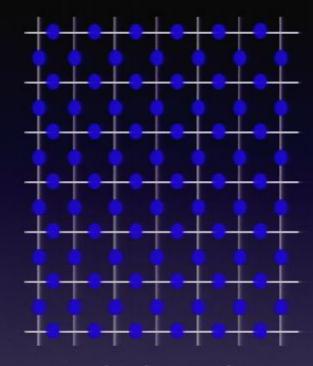






## Toy Model / Strawman





periodic boundary conditions (torus)

plaquette operator

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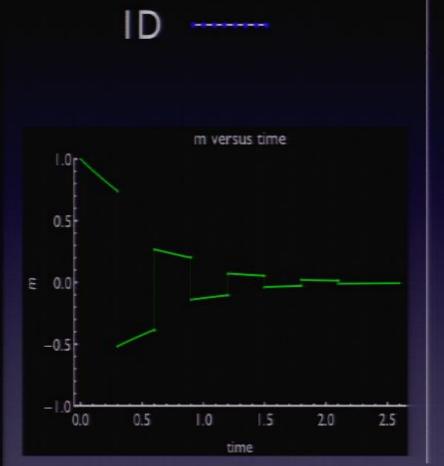
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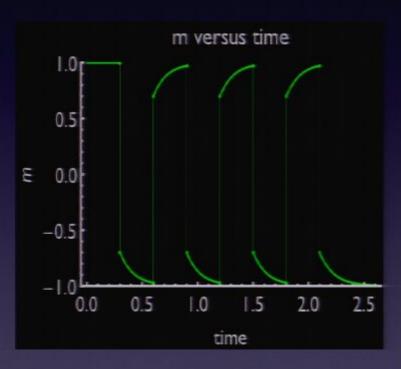
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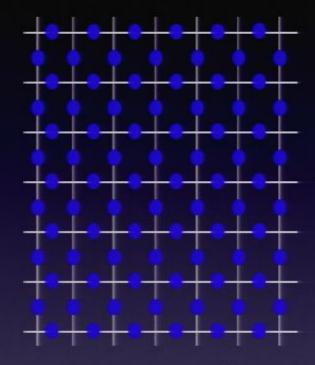






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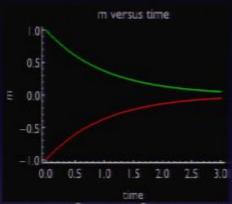
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### A Problem?

#### Toric code is like ID Ising:



"On Thermalization in Kitaev's 2D model" R. Alicki, M. Fannes, and M. Horodecki, arXiv:0810.4584

Decay rate for information encoded into ground state is independent of system size, but proportional to  $\exp(-c\beta J)$ 

Contrast with 4D model

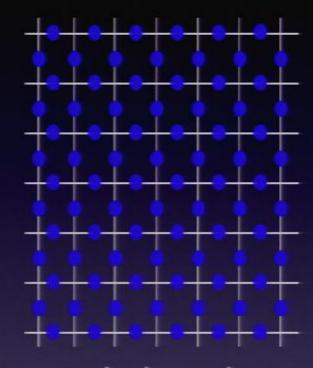
"On thermal stability of topological qubit in Kitaev's 4D model" R. Alicki, (M.+P. +R.) Horodecki, arXiv:0811.0033

$$\exp(-cN)$$

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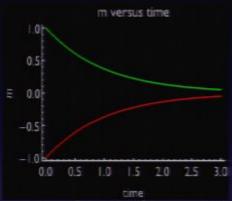
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### Not a Problem?

Factor a 1000 bit number:

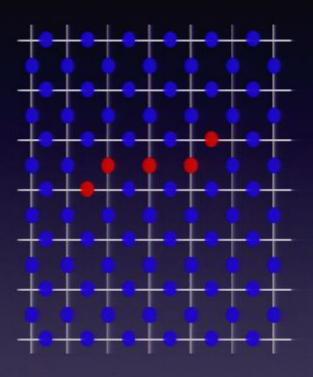
```
18819881292060796383869723946165043 98071635633794173827007633564229888 59715234665485319060606504743045317 38801130339671619969232120573403187 9550656996221305168759307650257059
```

Naive algorithm: ~109 gates Gap to temp ratio: ~20

Seems likely that exponential suppression as function of temp well worth working for

### A Problem?

What about during computing?



Creation and manipulation of
 excited states require
 operators like:

$$X^{\otimes 5}$$

How to do this?

$$\$^{\otimes 5}$$
 not  $X^{\otimes 5}$ 

How do you "create" and "move" anyons without creating errors = creating extra anyons

### TQC: Gates

"We now introduce into the system's Hamiltonian a scalar potential composed of many local "traps," each sufficient to capture exactly one quasiparticle. These traps may be created by impurities, by small gates, or by the potential created by tips of scanning microscopes. The quasiparticle's charge screens the potential introduced by the trap and the quasiparticle-tip combination cannot be observed by local measurements from far away. We denote the positions of these traps by R<sub>1</sub>,...,R<sub>k</sub>, and assume that these positions are well spaced from each other compared to the microscopic length scales. A state with quasiparticles at these positions can be viewed as an excited state of the Hamiltonian without the trap potential or, alternatively, as the ground state in the presence of the trap potential..."

C. Nayak, S. H. Simon, A. Stern, M. Freedman, and S. Das Sarma "Non-Abelian anyons and topological quantum computation" Rev. Mod. Phys. 80, 1083 (2008).

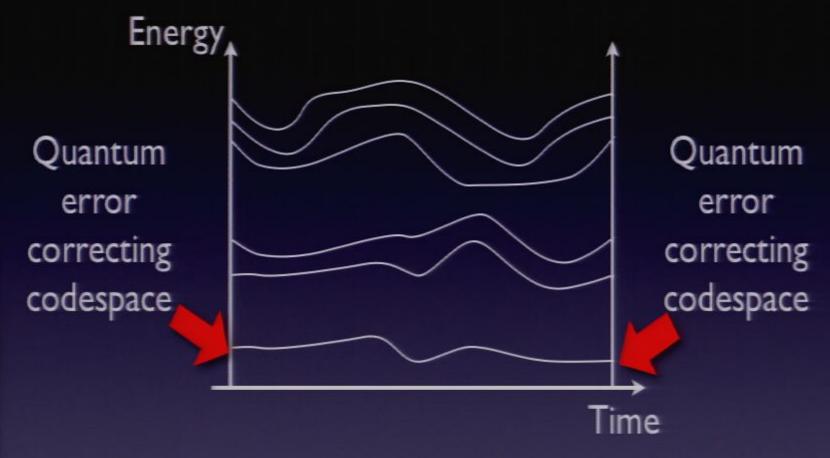
### TQC Problems?

How does one guarantee that one can trap one and only one anyon?

Every operation the previous paragraph needs to be demonstrated that it can be done with EXCEPTIONALLY small error if this is to be called physical fault-tolerance.

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# Adiabatic Dragging



Require: degenerate through entire evolution

"Open loop" holonomic evolution







Initial Code

M

$$|\Phi_{+}\rangle = \frac{1}{\sqrt{2}}(|00\rangle + |11\rangle)$$

Final Code





$$|\psi
angle\otimes|\Phi_{+}
angle$$

$$|\Phi_+
angle\otimes|\psi'
angle$$







#### Initial Code

$$|\Phi_{+}\rangle = \frac{1}{\sqrt{2}}(|00\rangle + |11\rangle)$$

#### Final Code





$$|\psi
angle\otimes|\Phi_{+}
angle$$

$$|\Phi_{+}\rangle\otimes|\psi'\rangle$$

#### Initial Stabilizer

$$S_1 = Z_2 Z_3$$

$$S_2 = X_2 X_3$$

$$S_i|\psi\rangle\otimes|\Phi_+\rangle=|\psi\rangle\otimes|\Phi_+\rangle$$

#### Final Stabilizer

$$S_1' = Z_1 Z_2$$

$$S_2' = X_1 X_2$$

$$S_i'|\Phi_+\rangle\otimes|\psi\rangle=|\Phi_+\rangle\otimes|\psi\rangle$$

Initial Code



$$|\psi
angle\otimes|\Phi_{+}
angle$$

Final Code





$$|\Phi_{+}\rangle\otimes|\psi'
angle$$

Initial Hamiltonian  $\Delta > 0$  Final Hamiltonian

$$\Delta > 0$$

$$H_i = -\Delta(X_2X_3 + Z_2Z_3)$$

$$H_i = -\Delta(S_1 + S_2)$$

$$H_i = -\Delta(X_2X_3 + Z_2Z_3)$$
  $H_f = -\Delta(X_1X_2 + Z_1Z_2)$ 

$$H_i = -\Delta(S_1' + S_2')$$

Evolution: 
$$H(t) = \left(1 - \frac{t}{T}\right)H_i + \frac{t}{T}H_f$$

$$ar{X}_1 = X_1 X_2 \qquad ar{Z}_1 = Z_2 Z_3 \ ar{X}_2 = X_2 X_3 \qquad ar{Z}_2 = Z_1 Z_2 \ ar{X}_3 = X_1 X_2 X_3 \qquad ar{Z}_3 = Z_1 Z_2 Z_3 \ ar{X}_3 = X_1 X_2 X_3 \qquad ar{Z}_3 = Z_1 Z_2 Z_3 \ ar{X}_3 =$$

Initial Hamiltonian

$$H_i = -\Delta(\bar{Z}_1 + \bar{X}_2)$$

$$H(s) = (1 - s)H_i + sH_f$$

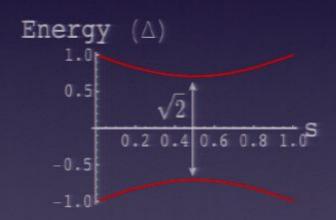
$$\bar{Z}_1 = Z_2 Z_3$$

$$\bar{Z}_2 = Z_1 Z_2$$

$$\bar{Z}_3 = Z_1 Z_2 Z_3$$

Final Hamiltonian

$$H_f = -\Delta(\bar{X}_1 + \bar{Z}_2)$$



$$\bar{X}_1 = X_1 X_2$$

$$\bar{X}_2 = X_2 X_3 \qquad \bar{Z}_2 = Z_1 Z_2$$

$$\bar{X}_3 = X_1 X_2 X_3 \qquad \bar{Z}_3 = Z_1 Z_2 Z_3$$

$$\bar{Z}_1 = Z_2 Z_3$$

$$\bar{Z}_2 = Z_1 Z_2$$

$$\bar{Z}_3 = Z_1 Z_2 Z_3$$

#### Initial Hamiltonian

$$H_i = -\Delta(\bar{Z}_1 + \bar{X}_2)$$

$$X_1 = \bar{X}_2 \bar{X}_3$$

$$Z_1 = \bar{Z}_1 \bar{Z}_3$$

#### Final Hamiltonian

$$H_f = -\Delta(\bar{X}_1 + \bar{Z}_2)$$

$$X_3 = \bar{X}_1 \bar{X}_3$$

$$Z_3 = \bar{Z}_2 \bar{Z}_3$$

$$|\psi
angle\otimes|\Phi_{+}
angle$$

$$|\Phi_{+}\rangle\otimes|\psi\rangle$$

$$H_i = -U_3 \Delta (X_2 X_3 + Z_2 Z_3) U_3^{\dagger}$$
  $H_f = -\Delta (X_1 X_2 + Z_1 Z_2)$ 

Note: spectrum unchanged by unitary

$$|\psi\rangle\otimes|\Phi_{+}\rangle$$
  $|\Phi_{+}\rangle\otimes U_{3}|\psi\rangle$ 

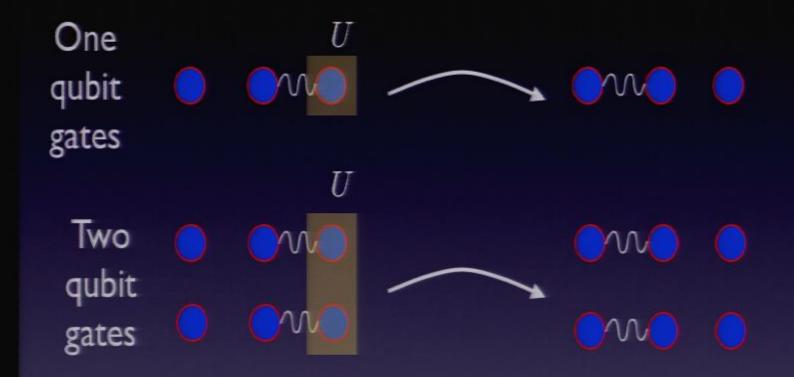
Example (Hadamard gate):

$$H_i = -\Delta(X_2Z_3 + Z_2X_3)$$
  $H_f = -\Delta(X_1X_2 + Z_1Z_2)$ 

$$|\psi\rangle \otimes |\Phi_{+}\rangle$$

$$H = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}$$

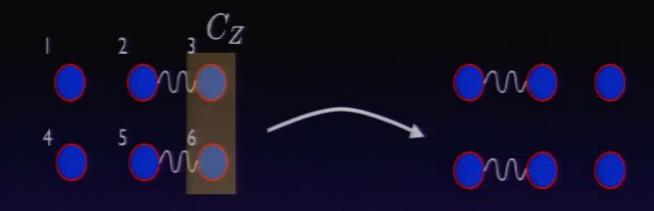
# Let's Build a Computer



Problem: initial Hamiltonian requires 3 qubit interactions

Pirea: 00000001

### Example



$$H_i = -\Delta(X_2 X_3 Z_6 + Z_2 Z_3 + Z_5 Z_6 + Z_3 X_5 X_6)$$

To get 2 qubit interactions: perturbation theory gadgets

# Gadgets



Replace 3,6 by encoded qubits:  $|0_L\rangle = |00\rangle, |1_L\rangle = |11\rangle$ 

$$H_{i} = -\lambda(X_{2}X_{3a} + Z_{2}Z_{3a}) - \lambda(X_{5}X_{6a} + Z_{5}Z_{6a})$$
$$-\omega Z_{3a}Z_{3b} - \lambda(X_{3b}Z_{6b} + Z_{3b}X_{6b}) \qquad \lambda \ll \omega$$

$$H_f = -\lambda (X_1 X_2 + Z_1 Z_2) - \lambda (X_4 X_5 + Z_4 Z_5) - \omega Z_{3a} Z_{3b}$$

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[S. D. Bartlett and T. Rudolph, Phys. Rev. A 74, 040302 (2006)]

# Gadgets

$$H_i = -\lambda(X_2X_{3a} + Z_2Z_{3a}) - \lambda(X_5X_{6a} + Z_5Z_{6a}) - \omega Z_{3a}Z_{3b} - \lambda(X_3bZ_{6b} + Z_3bX_{6b})$$
 $H_f = -\lambda(X_1X_2 + Z_1Z_2) - \lambda(X_4X_5 + Z_4Z_5) - \omega Z_{3a}Z_{3b}$ 

Minimum gap: 
$$O\left(\frac{\lambda^2}{\omega}\right)$$

Initial ground state fidelity:  $1 - O\left(\frac{\lambda^2}{\omega^2}\right)$ 

### Variations on a Theme

Exchange interactions:

$$H = -\Delta(X_1 X_2 + Y_1 Y_2 + Z_1 Z_2)$$

Adiabatically preparing rotated Hamiltonians:

$$H_i = -\Delta(X_1X_2 + Z_1Z_2)$$
  $H_f = -\Delta U_2(X_1X_2 + Z_1Z_2)U_2^{\dagger}$ 

Not all U's allow a gap

$$A = \frac{1}{2} \begin{bmatrix} 1 + i\sqrt{2} & 1 \\ 1 & -1 + 1\sqrt{2} \end{bmatrix} \qquad B = \begin{bmatrix} 1 & 0 \\ 0 & e^{i\pi} \end{bmatrix}$$

gap: 
$$\sqrt{2}\Delta$$
 gap:  $\sqrt{2}-\sqrt{2}\Delta$ 

# Gadgets

$$H_i = -\lambda(X_2X_{3a} + Z_2Z_{3a}) - \lambda(X_5X_{6a} + Z_5Z_{6a}) - \omega Z_{3a}Z_{3b} - \lambda(X_1X_2 + Z_1Z_2) - \lambda(X_4X_5 + Z_4Z_5) - \omega Z_{3a}Z_{3b}$$

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gap: 
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 gap:  $\sqrt{2}-\sqrt{2}\Delta$ 

# "Impossibility" Theorem

Is it possible to adiabatically SWAP with only one adiabatic interpolation?

$$H_i = \delta_1(|01\rangle\langle 01| + |11\rangle\langle 11|) + \delta_2|00\rangle\langle 00| + \delta_3|10\rangle\langle 10|$$

$$H_f = \gamma_1(|10\rangle\langle 10| + |11\rangle\langle 11|) + \gamma_2|00\rangle\langle 00| + \gamma_3|01\rangle\langle 01|$$

$$H(t) = \left(1 - \frac{t}{T}\right)H_i + \frac{t}{T}H_f$$

Diagonal in same basis so cannot transfer amplitude to perform SWAP

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# Gadgets

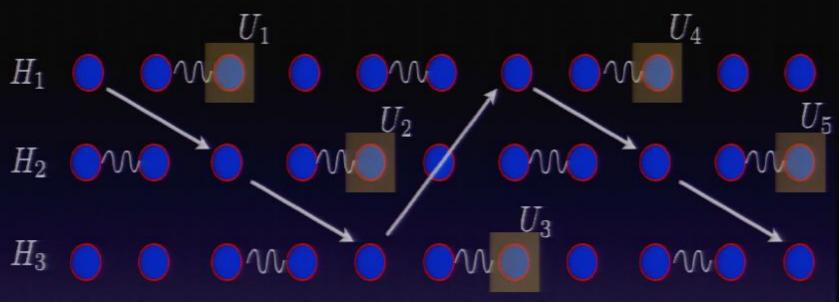
$$H_i = -\lambda(X_2X_{3a} + Z_2Z_{3a}) - \lambda(X_5X_{6a} + Z_5Z_{6a}) - \omega Z_{3a}Z_{3b} - \lambda(X_3bZ_{6b} + Z_3bX_{6b})$$
 $H_f = -\lambda(X_1X_2 + Z_1Z_2) - \lambda(X_4X_5 + Z_4Z_5) - \omega Z_{3a}Z_{3b}$ 

Minimum gap: 
$$O\left(\frac{\lambda^2}{\omega}\right)$$

Initial ground state fidelity:  $1 - O\left(\frac{\lambda^2}{\omega^2}\right)$ 

### Architectures





Sequence of Hamiltonians:  $H_1, H_2, H_3, H_1, H_2, H_3, ...$ 

Sequence of single qubit gates applied:  $U_1, U_2, U_3, U_4, U_5, \dots$ 

Generalizing to more than I qubit:

Universal QC by cycling between only 3 Hamiltonians Page 70

# "Impossibility" Theorem

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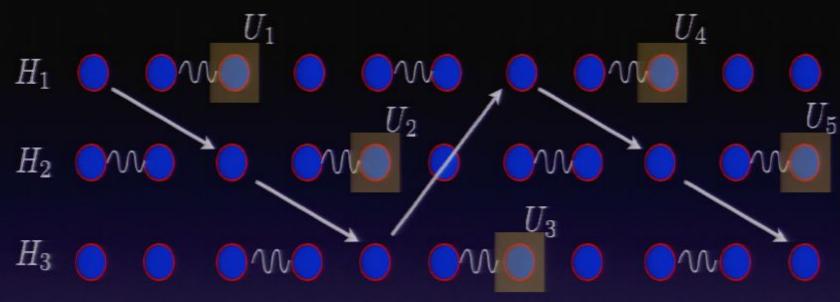
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Diagonal in same basis so cannot transfer amplitude to perform SWAP

Pirsa: 09090004

### Architectures





Sequence of Hamiltonians:  $H_1, H_2, H_3, H_1, H_2, H_3, ...$ 

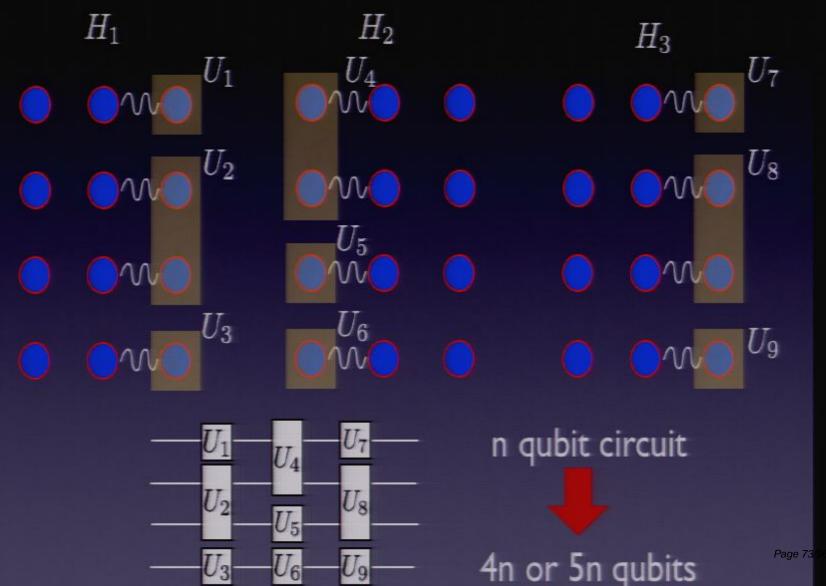
Sequence of single qubit gates applied:  $U_1, U_2, U_3, U_4, U_5, \dots$ 

Generalizing to more than I qubit:

Universal QC by cycling between only 3 Hamiltonians Page 72

### Architectures





# Comparison

Holonomic quantum computing, but open loop open loop holomic QC:
[D. Kult, J. Aberg, E. Sjoqvist, Phys. Rev. A 74, 022106 (2006)]

Adiabatic universal quantum computing, but with piecewise evolution

D. Aharonov et al., in 45th Annual IEEE FOCS (2004)]

Path robustness in holonomy is subsystem dependent compare:

[O. Oreshkov, T.A. Brun, D.A. Lidar, Phys. Rev. Lett. 102, 070502 (2009)]

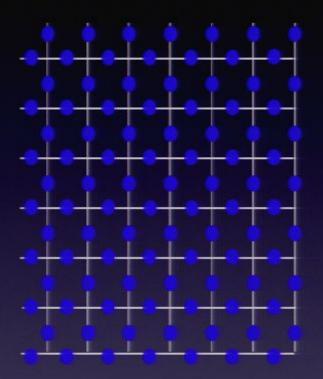
#### Piecewise keeps gap constant:

compare spin-1 chain transmission of:

[K. Eckert, O. Romero-Isart, and A. Sanpera, New J. Phys. 9, 155 (2007)]

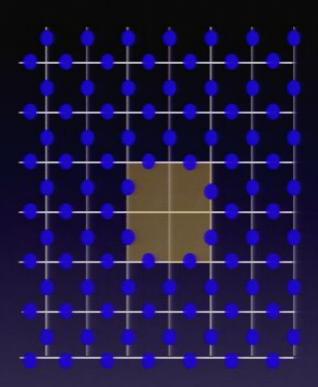
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# Back to TQC!

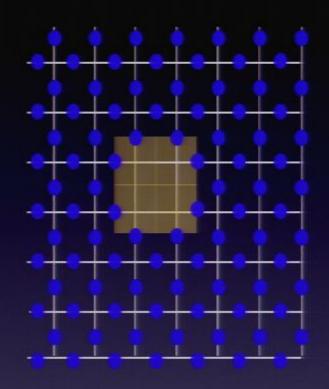


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### **Encoded Qubits**

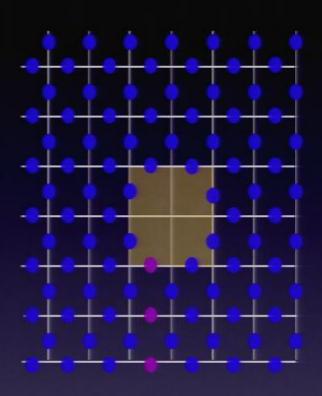


Smooth qubit

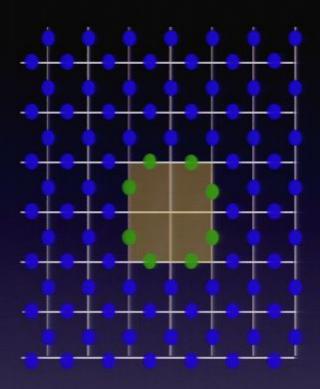


Rough qubit

### Back to TQC!

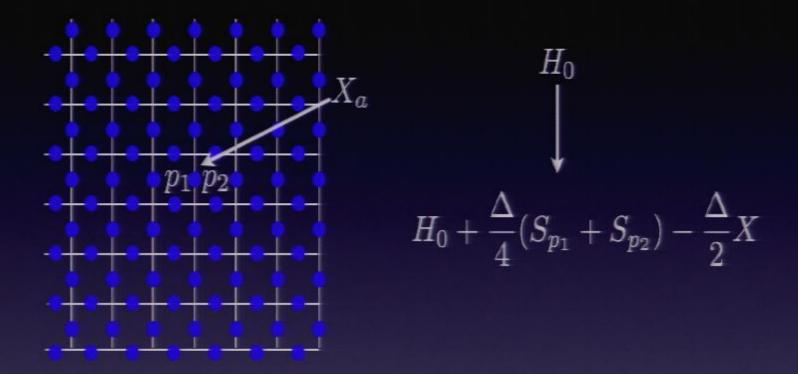


$$\bar{X} = X^{\otimes c_1}$$

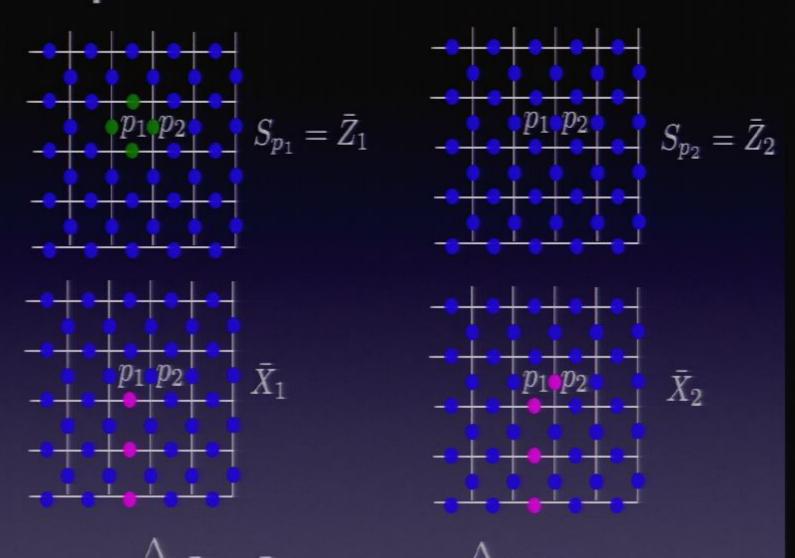


$$\bar{Z} = Z^{\otimes c_2}$$

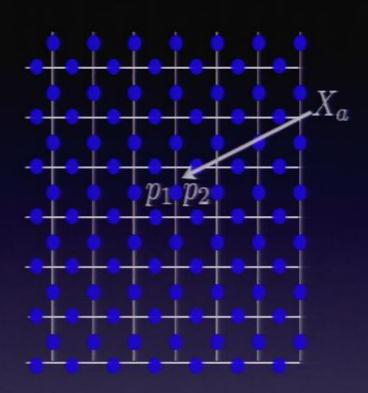
Smooth qubit



Turn off 2 plaquette operators adjacent to a qubit where X is turned on.



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$$H_0$$

$$\downarrow$$
 $H_0 + \frac{\Delta}{4}(S_{p_1} + S_{p_2}) + \frac{\Delta}{2}X_a$ 

$$-\frac{\Delta}{4}(\bar{Z}_1 + \bar{Z}_2) \qquad \bar{Z}_1\bar{Z}_2$$

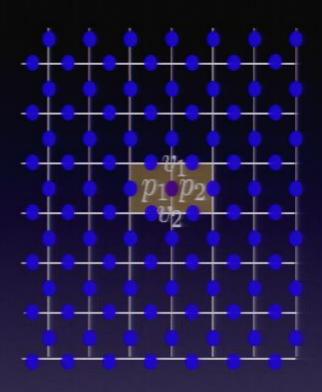
$$\downarrow \qquad \qquad \text{commutes}$$

$$-\frac{\Delta}{2}\bar{X}_1\bar{X}_2 \qquad \text{with H(t)}$$

+1 eigenspace of  $\bar{Z}_1\bar{Z}_2$ 

$$-\frac{\Delta}{2}X \longrightarrow -\frac{\Delta}{2}Z$$

end up in +1 e.v. of smooth qubit encoded Z



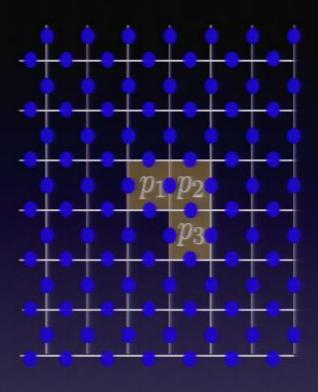
$$H_0 \downarrow$$

$$H_0 + \frac{\Delta}{4}(S_{p_1} + S_{p_2}) + \frac{\Delta}{2}X_a$$

$$H_0 + \frac{\Delta}{4}(S_{p_1} + S_{p_2}) - \frac{\Delta}{2}X_a + \frac{\Delta}{4}(S_{v_1} + S_{v_2}) - \frac{\Delta}{4}(S'_{v_1} + S'_{v_2}))$$

Smooth qubit in +1 eigenvalue of encoded Z

#### Deform Smooth Puncture



Turn on

$$-\frac{\Delta}{2}X_b$$

While turning off

$$-\frac{\Delta}{4}S_{p_3}$$

+ relevant vertex operators

Rough qubits: Ditto

Rough qubit prepare +1 e.v. of encoded X

#### **Toolkit**

- 1. Create small smooth qubit in +1 e.v. of Z
- 2. Create small rough qubit in +1 e.v. of X
- 3. Enlarge, move both punctures
- 4. Braid puncture: this give a  $C_X$  target = rough

These things all commute with each other, as of course these are Abelian anyons.

<sup>p</sup>irsa: 090900<mark>0</mark>4 Page 83<mark>.9</mark>

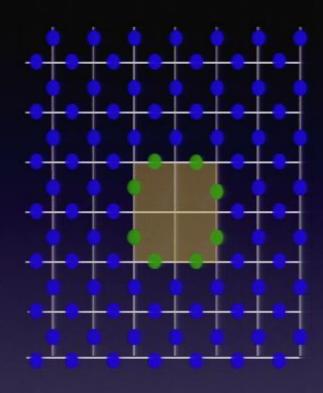
#### Universal Toolkit

- I. Create small smooth qubit in +I e.v. of Z
- 2. Create small rough qubit in +1 e.v. of X
- 3. Enlarge, move both punctures
- 4. Braid puncture: this give a  $C_X$  target = rough
- 5. Measure encoded X,Z of qubits
- 6. Create magic states

Requires distillation.

<sup>p</sup>irsa: 090900<mark>0</mark>4 Page 84

#### Measurement



$$\bar{Z} = Z^{\otimes c_2}$$

Simple, right?

But how to perform why Hamiltonian still on?

Note: destructive Pauli measurements are tolerant to errors because (1) complementary errors are large and (2) we can remove region being measured

#### Universal Toolkit

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5+6 not tolerant to errors in this model

<sup>p</sup>irsa: 090900<mark>0</mark>4 Page 86<mark>.9</mark>

# Summary

Adiabatically deforming between different
Hamiltonians which have stabilizer codewords as
energy eigenstates yields interesting protocols like
AGT

Using AGT as a motivator you can understand how topological quantum computing can or cannot robustly compute (and thus can provide rigorous error bounds.)

rsa: 09090004 Page 87/9

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<sup>p</sup>irsa: 090900<mark>0</mark>4 Page 88<mark>/9</mark>

# Summary

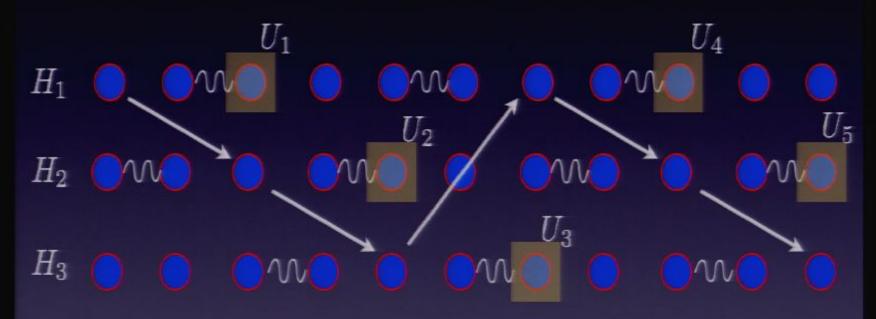
Adiabatically deforming between different Hamiltonians which have stabilizer codewords as energy eigenstates yields interesting protocols like AGT

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# Adiabatic Gate Teleportation

Dave Bacon University of Washington, Seattle, WA USA



[joint work with Steve Flammia (Perimeter), Alice Neels (Washington), Andrew Landahl (Sandia Labs)]

