

Title: Computational difficulty of simulation methods: Density Functional Theory, DMRG, and beyond

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Abstract: We analyze how quantum complexity poses bounds to the simulation of quantum systems. While methods as Density Functional Theory (DFT) and the Density Matrix Renormalization Group (DMRG) work very well in practice, essentially nothing on the formal requirements is known. In this talk, we consider these methods from a quantum complexity perspective: First, we discuss DFT which encapsulates the difficulty of solving the Schrodinger equation in a universal functional and show that this functional cannot be efficiently computed unless several complexity classes collapse. Second, we consider DMRG, a method to deal with quantum spin chains, and show that even under reasonable assumptions -- a polynomial gap and matrix product ground states -- finding the ground state is still a computationally hard problem. Beyond pinpointing the limitations of the methods, this helps us to understand under which assumptions we might be able to prove their convergence.

Complexity of sim.  $q$ -many-body systems

$q$ -many-body systems



Simulate

# Complexity of sim. $q$ -many-body systems

$q$ -many-body systems



simulate

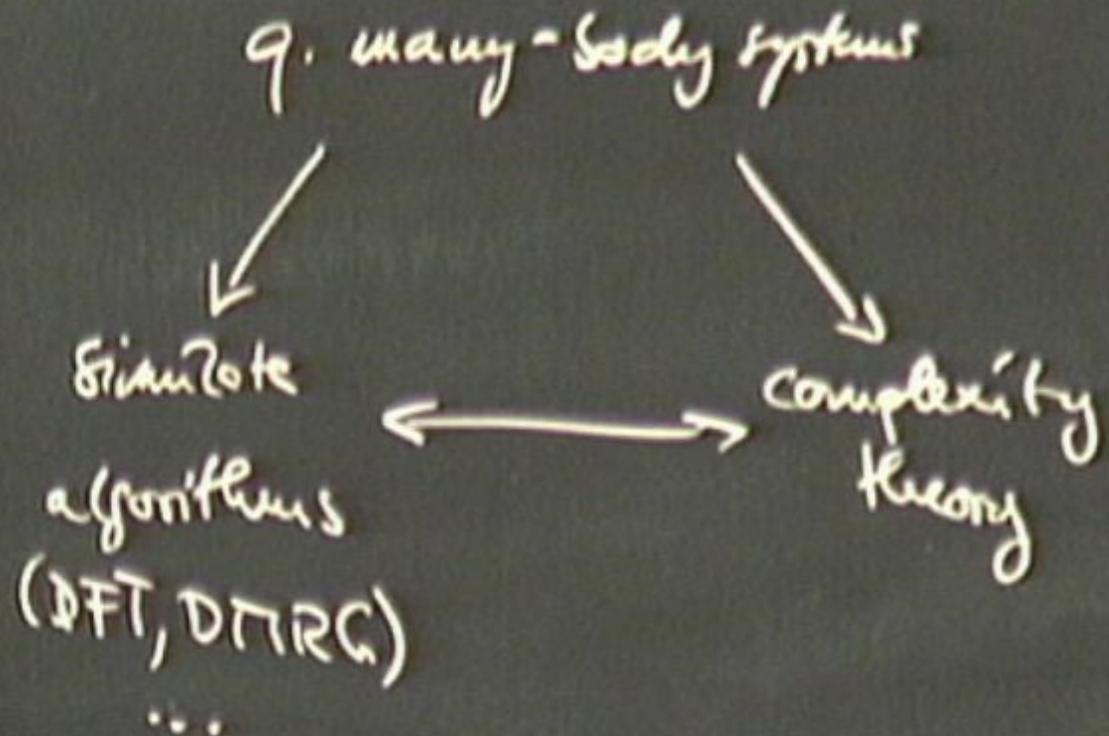
algorithms

(DFT, DMRG)

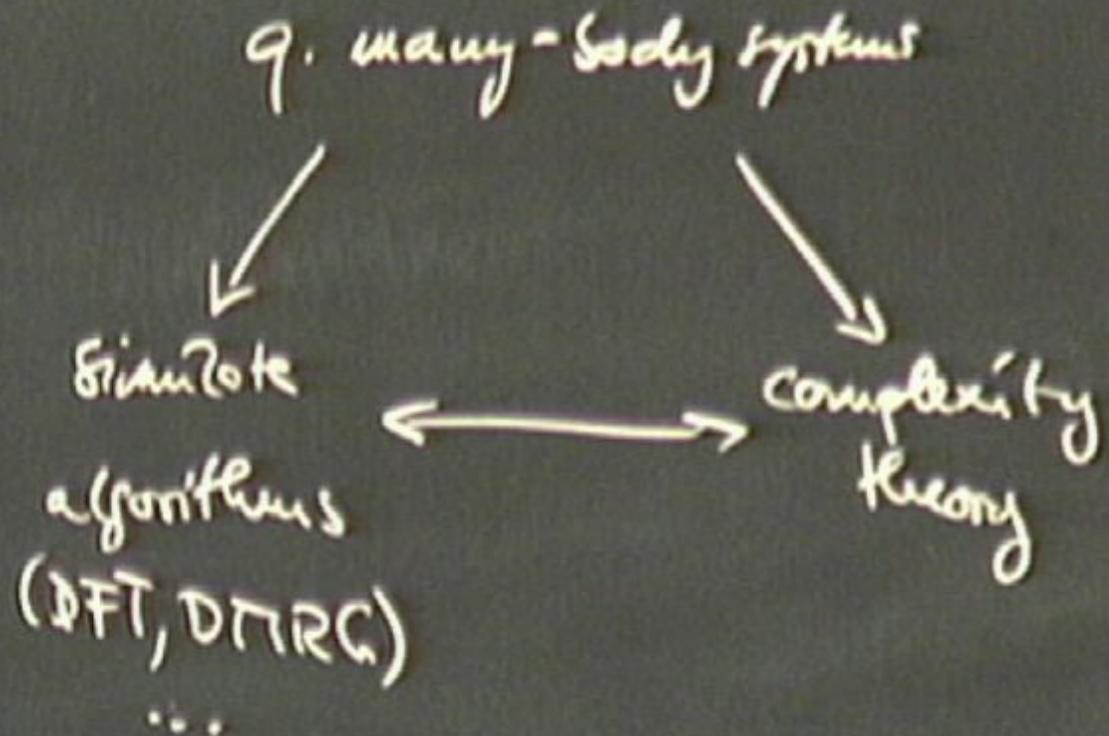
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# Complexity of sim. $q$ -many-body systems



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# Density Functional Theory



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$N$  electrons in an ext. potential

→ ground state energy

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$$H = -\frac{1}{2} \sum_i \Delta_i + \sum_{i,j} \frac{e^2}{|\mathbf{r}_i - \mathbf{r}_j|} + \sum_i V(\mathbf{r}_i)$$

# Density Functional Theory

$N$  electrons in an ext. potential

→ ground state energy

$$H = \underbrace{-\frac{1}{2} \sum_i \Delta_i}_{=: T} + \underbrace{\sum_{i,j} \frac{e^2}{|r_i - r_j|}}_{=: I} + \underbrace{\sum_i V(r_i)}_{=: V}$$

$$H = \min$$

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$$\min_{\Omega} [\text{tr}[(T+I+V)\rho]]$$

$\Omega$   
↑  
antisym.  $N$ -el.  
density matrix

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$$E_0 = \min_{\Omega} [\text{tr}[(T+I+V)\Omega]] = \min$$

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$$E_{gs} = \min_{\Omega} [\text{tr}[(T+I+V)\rho]] = \min_{\rho(r)} \min_{\Omega \rightarrow \rho}$$

$\Omega$ : N-el. density matrix

$$\rho(r, r') = \int \Omega(r_1, r'_1, r_2, r'_2, \dots, r_N, r'_N) dr_2, \dots, dr_N$$

(DFT, DMRG)

# Density Functional Theory

N electrons in an ext. potential

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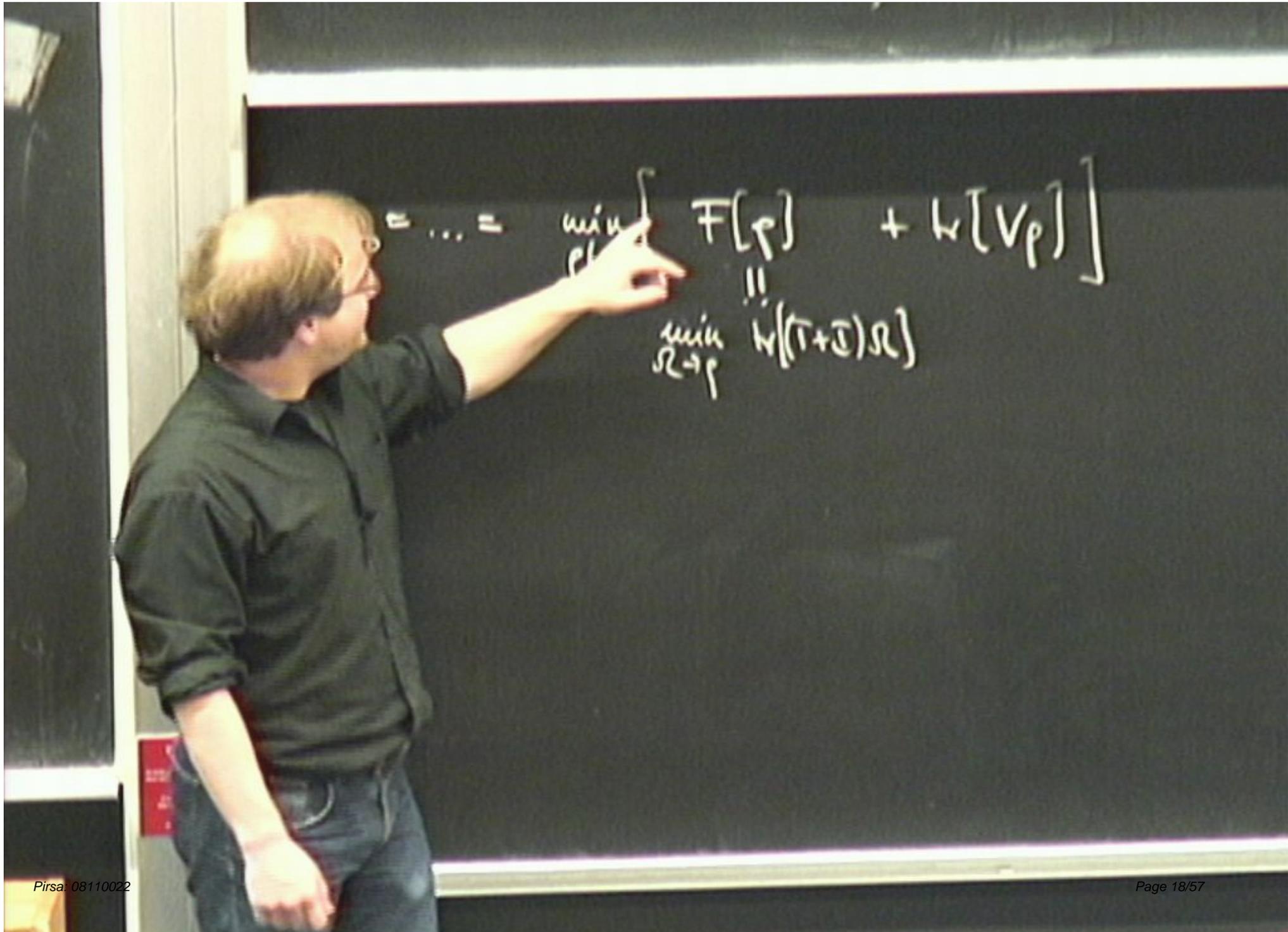
$$E_{gs} = \min_{\Omega} [\text{tr}[(T+I+V)\rho]] = \min_{\rho(r)} \min_{\Omega \rightarrow \rho} [H[(T+I)\rho] + \int [V\rho]]$$

$\Omega$   
↑  
antisym.  $N$ -el.  
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$$\rho(r, r') = \int \Omega(r_1, r'_1, r_2, r'_2, \dots, r_N, r'_N) dr_2, \dots, dr_N$$

(DFT, DMRC)

...



$$= \dots = \min_{\rho} [ F[\rho] + \kappa [V\rho] ]$$
$$!!$$
$$\min_{\Omega \rightarrow \rho} \kappa [(\Gamma + \mathcal{D})\Omega]$$

$$E_0 = \dots = \min_{p(r)} \left[ F(p) + h[V_p] \right]$$

$\min_{R \rightarrow p} h[(1+\sigma)R]$



CAUTION  
DO NOT TOUCH  
ELECTRICAL EQUIPMENT  
OR WIRING

$$E_0 = \dots = \min_{\rho(r)} \left[ \overset{\text{"universal functional"}}{F[\rho]} + h[V\rho] \right]$$

$$\min_{\Omega \rightarrow \rho} h[(T+V)\Omega]$$

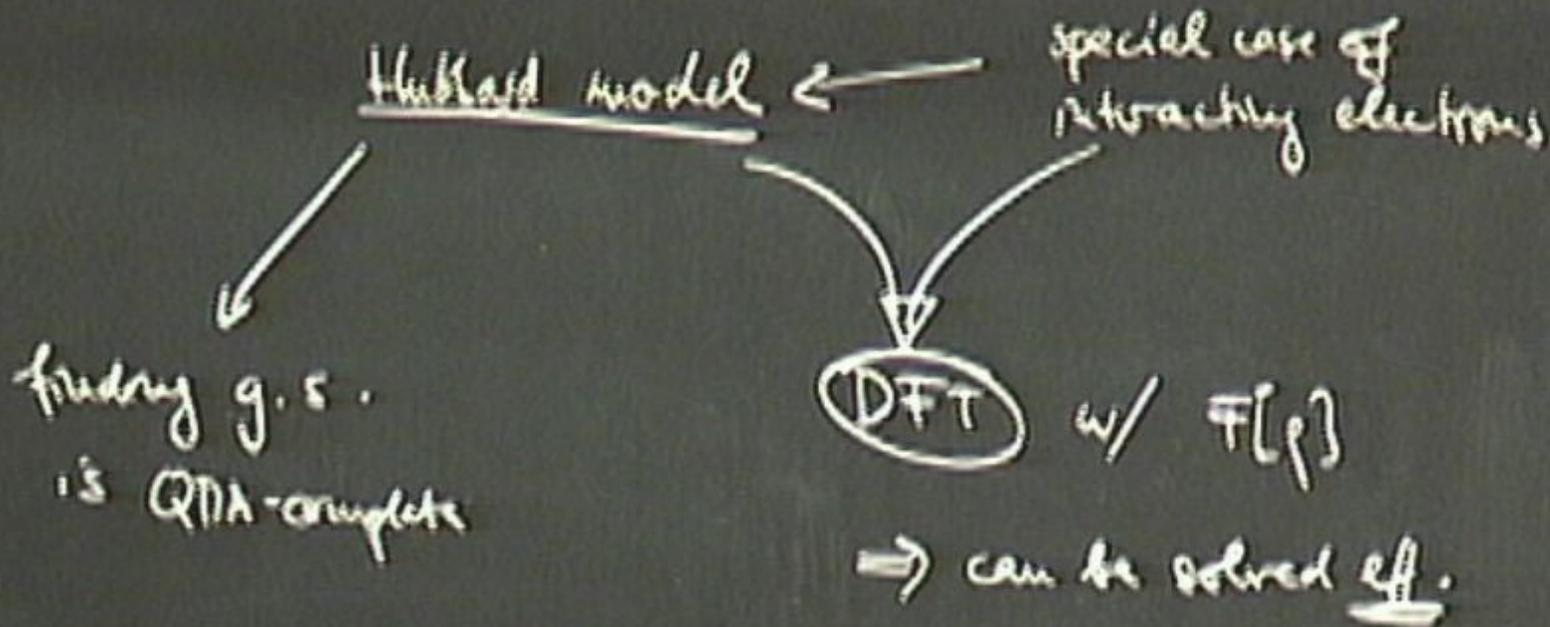
$$E_0 = \dots = \min_{\rho(r)} \left[ \overset{\text{"universal functional"}}{F[\rho]} + k[V\rho] \right]$$

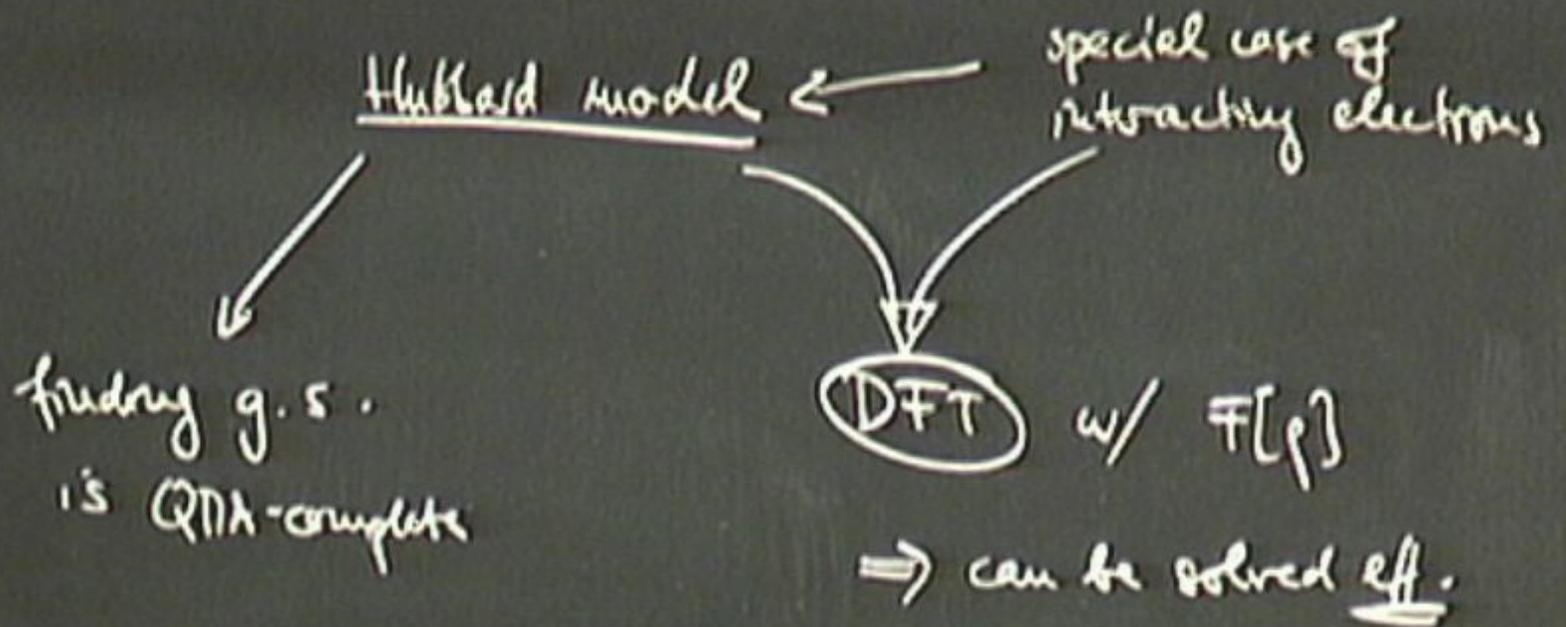
$$\min_{\Omega \rightarrow \rho} k[(T+V)\Omega]$$

Hubbard model

special case of  
interacting electrons

Frustrated g.s.  
is





Q. complexity:

P: can be comp in poly-time

NP: if answer is yes,  $\exists$  proof which can be checked eff.

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QMA problem: Given Loc. Ham.,

$$\text{is } E_0(H) < a \quad \left( \text{or above } b, \epsilon > \frac{1}{\text{poly}(N)} \right)$$

## Q. complexity:

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P: can be comp in poly-time

NP: if answer is yes,  $\exists$  proof which can be checked eff.

IQP: P w/ quantum

QMA -ess. Quantum NP

typ. problem: Given Loc. Ham.,

is  $E_0(H) < a$  (or above  $b$ ,  $b - a > \frac{1}{\text{poly}(N)}$ )

Proof: g.s.

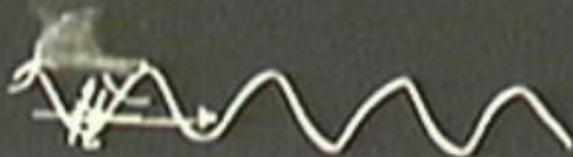
$$H_{\text{Hück}} = -t \sum_{\langle ij \rangle, s} a_{i,s}^\dagger a_{j,s} +$$



$$H_{\text{Hütt}} = -t \sum_{\langle i,j \rangle, s} a_{i,s}^\dagger a_{j,s} + \mu \sum_i n_{i,\uparrow} n_{i,\downarrow} + \sum_i \vec{B}_i \cdot \vec{\sigma}_i$$



$$H_{\text{Hubb}} = -t \sum_{\langle i,j \rangle, s} a_{i,s}^\dagger a_{j,s} + U \sum_i n_{i,\uparrow} n_{i,\downarrow} + \sum_i \vec{B}_i \cdot \vec{\sigma}_i$$



Finding g.s. energy of  $H_{\text{Hubb}}$  is QDA-complete.

$$H_{\text{Huff}} = -t \sum_{\langle i,j \rangle, s} a_{i,s}^\dagger a_{j,s} + \mu \sum_{i,j} u_{i,j} u_{i,j} + \sum_i \vec{B}_i \cdot \vec{\sigma}_i$$



$$\begin{array}{l} \sigma_i^+ \sigma_i^+ \sigma_i^2 \dots \\ \sigma_i^- \sigma_i^2 \dots \end{array}$$

Finding g.o. energy of H<sub>Huff</sub> is QDA-complete.



$$H_{\text{Huff}} = -t \sum_{\langle ij \rangle, s} a_{i,s}^\dagger a_{j,s} + \mu \sum_{i,j} u_{i,j} u_{i,j} + \sum_i \vec{B}_i \cdot \vec{\sigma}_i$$



$$\begin{matrix} \sigma^+ \sigma^+ \sigma^+ \dots \\ \sigma^- \sigma^- \dots \end{matrix}$$

Finding g.s. energy of  $H_{\text{Huff}}$  is QDA-complete.

- In QDA ✓
- hard for QDA ?

→ Show that a known QP-hard problem can be reduced to  $H_{\text{table}}$

→ Show that a known QP hard problem can be reduced to  $H_{\text{Harris}}$

$$H_{20} = \sum_{i,j} \lambda_{ij} A_{(ij)} \otimes B_{(ij)}, \quad A, B \text{ are } 2 \times 2 \text{ matrices}$$

→ Show that a known QP hard problem can be reduced to  $H_{\text{basis}}$

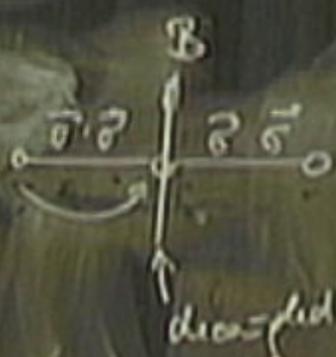
$$H_{\text{red}} = \sum_{(i,j)} \lambda_{ij} A_{(i,j)} \otimes B_{(i,j)}, \quad A, B \text{ are } 2 \times 2 \text{ matrices}$$

“Perturbation gadgets”:

→ Show that a known QMA-hard problem can be reduced to  $H_{\text{magic}}$

$H_{2D} = \sum_{i,j} \lambda_{ij} A_{(ij)} \otimes B_{(ij)}$ ,  $A, B$  are Pauli matrices

"Perturbation gadgets":



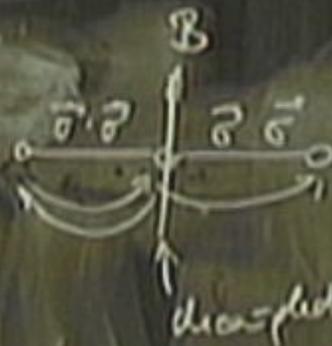
$H = B \cdot \sigma_{\text{control}} + \frac{h_{\text{weak}}}{\text{weak}}$



→ Show that a known QMA-hard problem can be reduced to  $H_{\text{local}}$

$$H_{\text{local}} = \sum_{(i,j)} \lambda_{ij} A_{(i,j)} \otimes B_{(i,j)}, \quad A, B \text{ are } \text{Pam} \otimes \text{U} \text{ matrices}$$

“Perturbation gadgets”:



$$H = B \cdot \underbrace{\sigma_{\text{control}}}_{\text{weak}} + \underbrace{\dots}_{\text{weak}}$$

→ Show that a known QMA-hard problem can be reduced to  $H_{\text{local}}$

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$$|\psi\rangle = \frac{1}{\sqrt{2}} (|xx\rangle + |yy\rangle) \otimes |zz\rangle$$

Permutation gadgets:



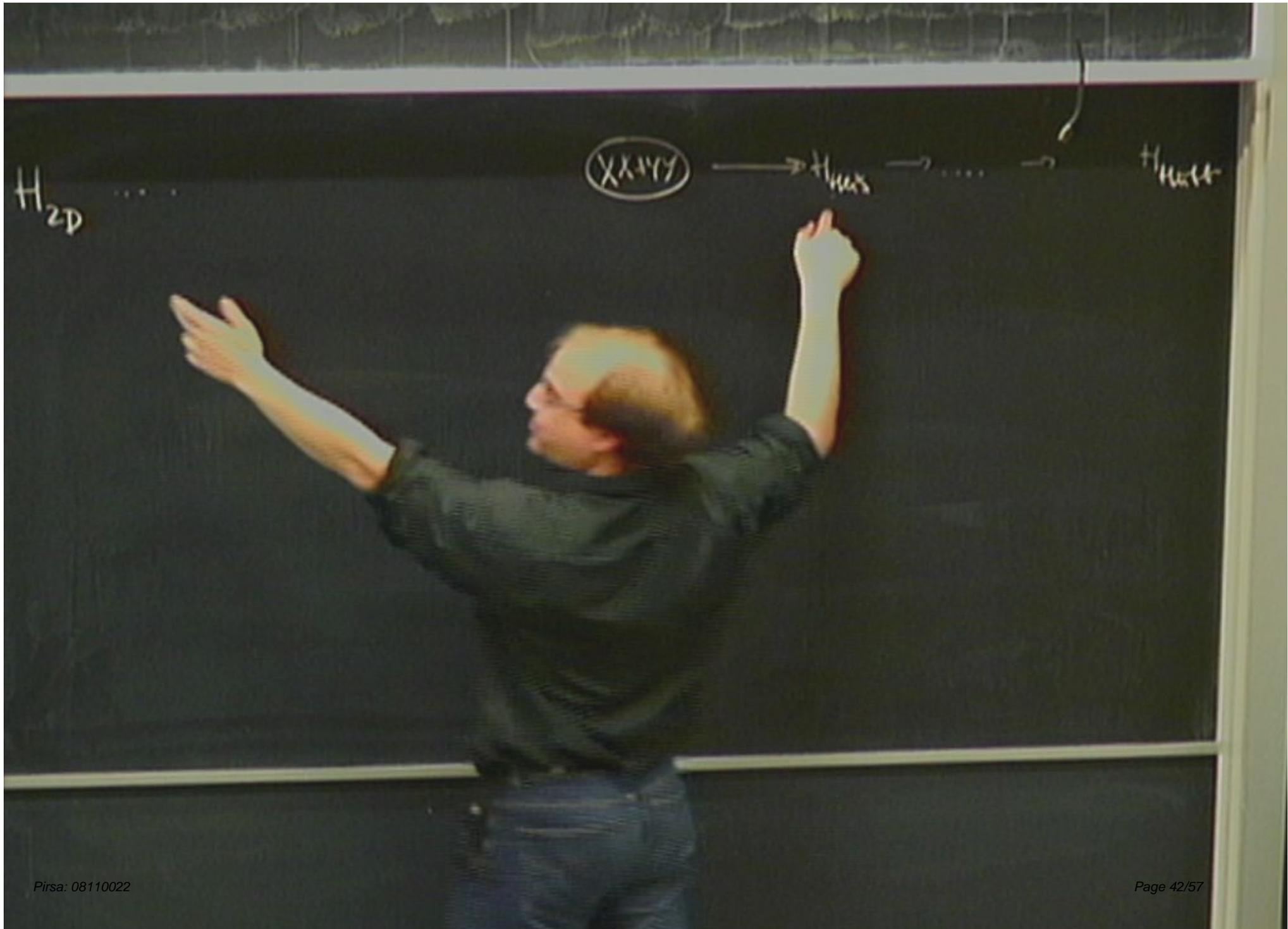
→ Show that a known QMA-hard problem can be reduced to  $H_{\text{table}}$

$$H_{20} = \sum_{i,j} \lambda_{ij} A_{(ij)} \otimes B_{(ij)}, \quad A, B \text{ are } \text{PancU} \text{ matrices}$$

$$|\psi\rangle = \frac{1}{\sqrt{3}} (|xx\rangle + |yy\rangle + |zz\rangle)$$

\* Restriction gadgets:





$H_{2D}$  ...

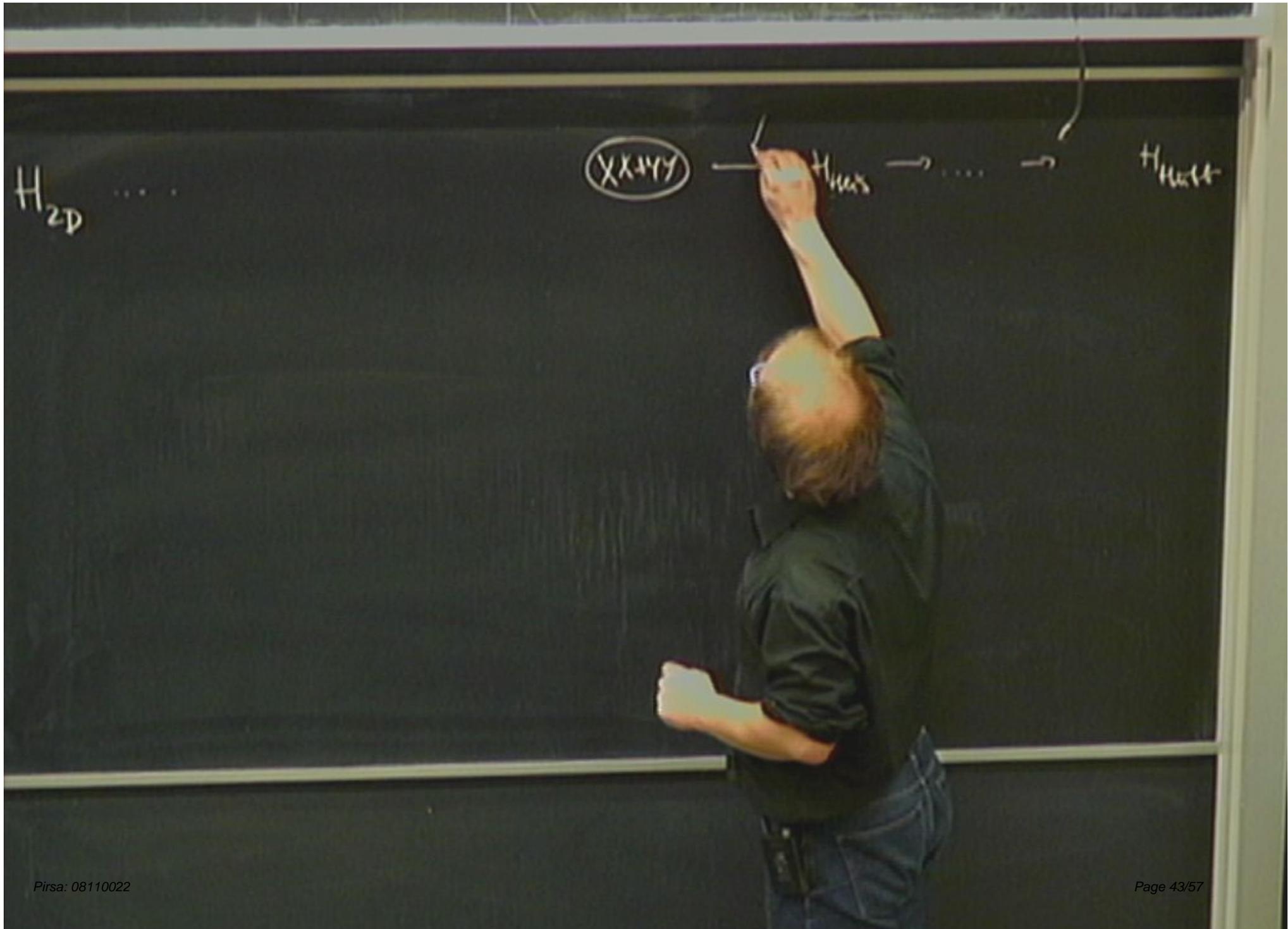
$(XY)$

→  $H_{2D}$

→ ...

→

$H_{2D}$

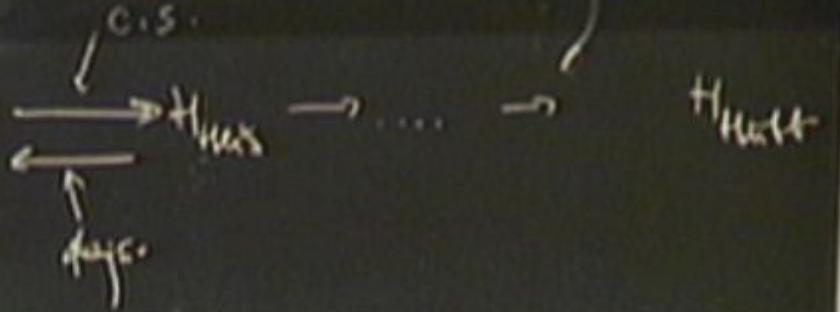


H<sub>2</sub>D ...

(XX+YY)

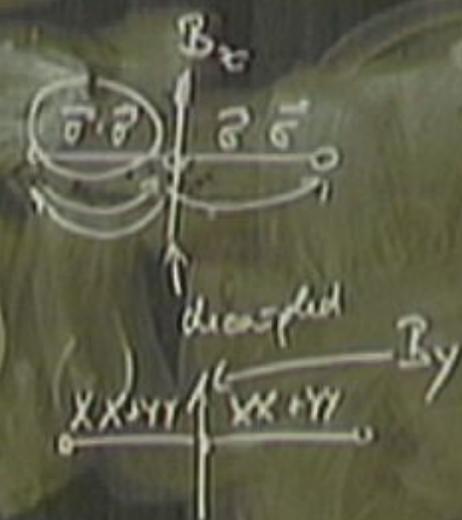
XX + ZZ

YY + ZZ



$$H_{2D} = \sum_{i,j} \lambda_{ij} A_{(ij)} \otimes B_{(ij)}, \quad A, B \text{ are rank 1 matrices}$$

Permutation gadgets:



$$\sigma \sigma = \begin{matrix} \boxed{XX+YY} & \boxed{YY} \\ \boxed{XX} & \boxed{YY} \end{matrix}$$

1015, 1107

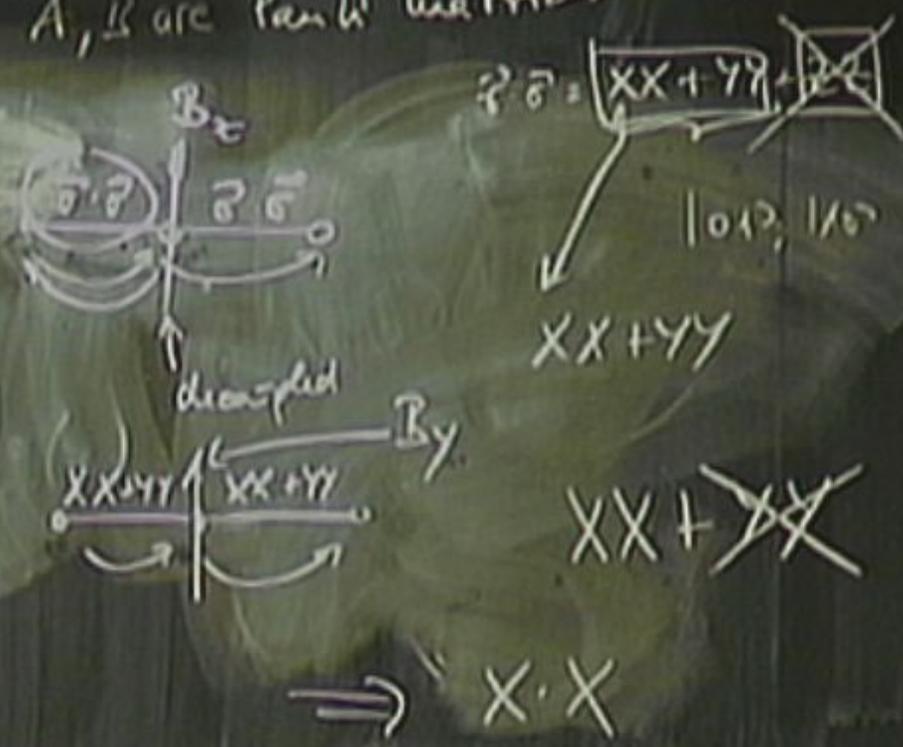
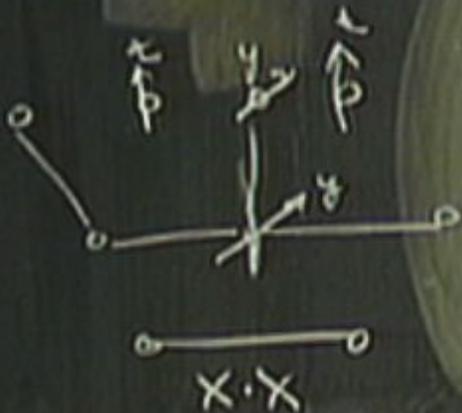
$$XX+YY$$

$$XX+YY$$



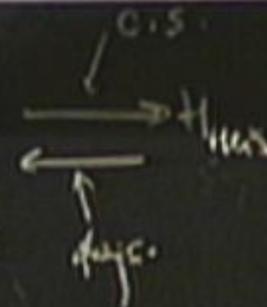
$$H_{2D} = \sum_{i,j} \lambda_{ij} A_{(ij)} \otimes B_{(ij)}, \quad A, B \text{ are Pauli matrices}$$

Permutation gadgets:

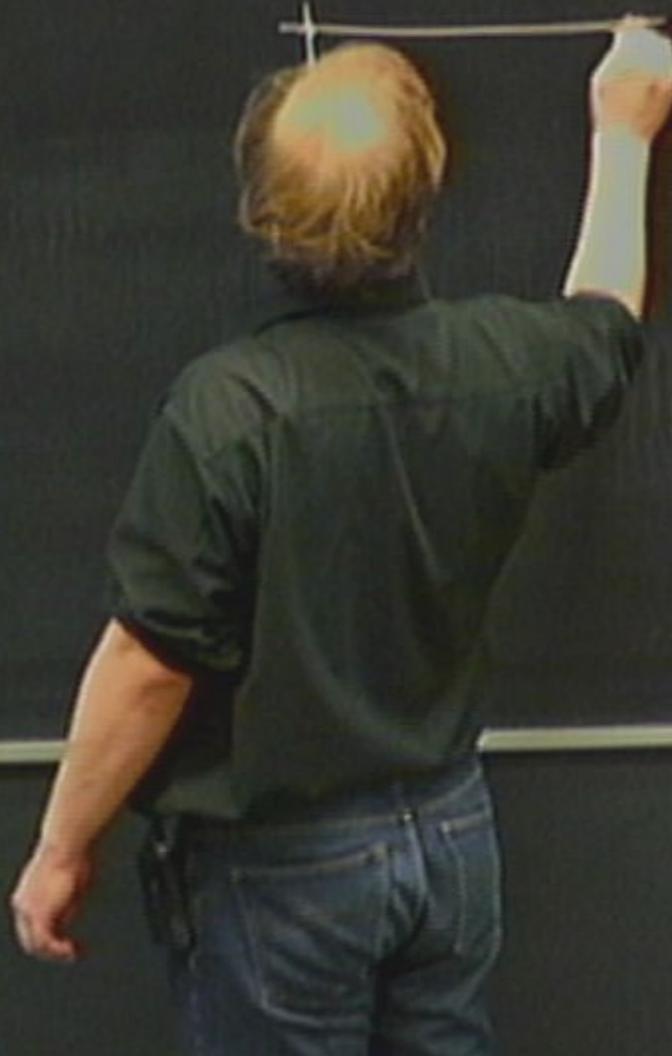
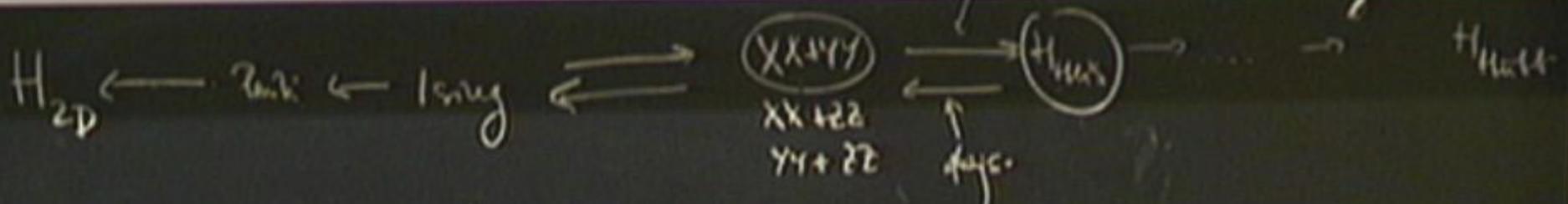


$H_{2D}$

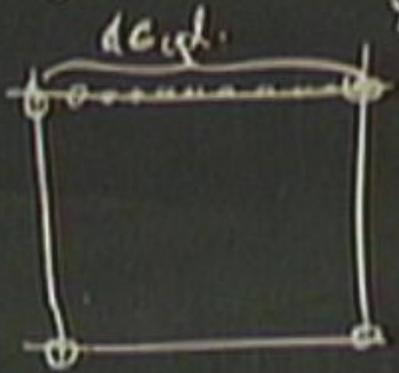
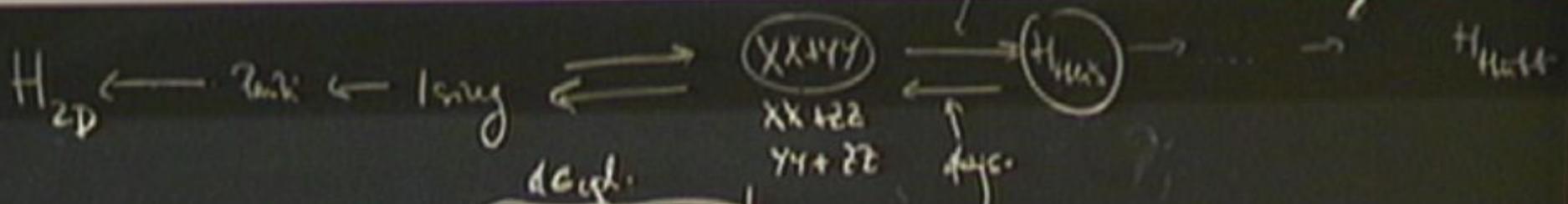
$\begin{pmatrix} XX+YY \\ XX+ZZ \\ YY+ZZ \end{pmatrix}$



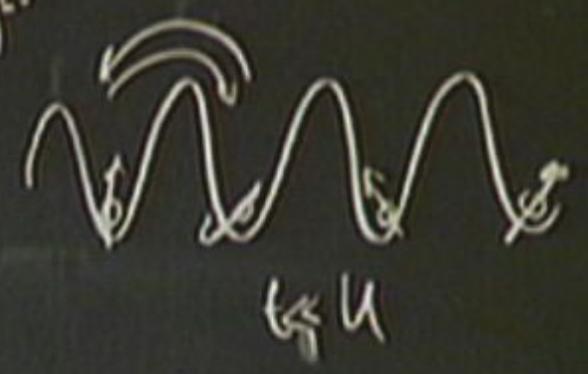
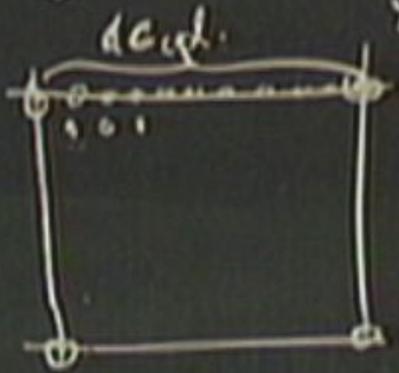
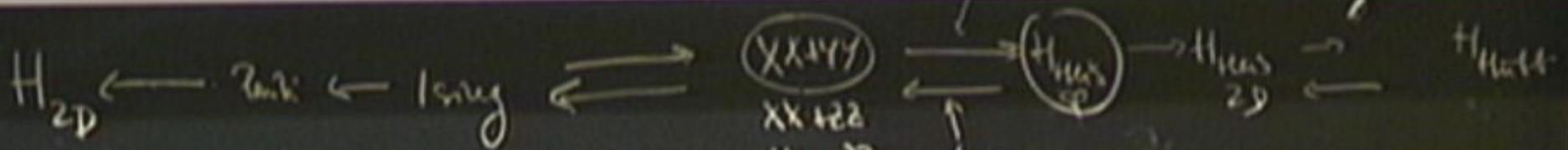
$-x \cdot x$     $-y \cdot y$



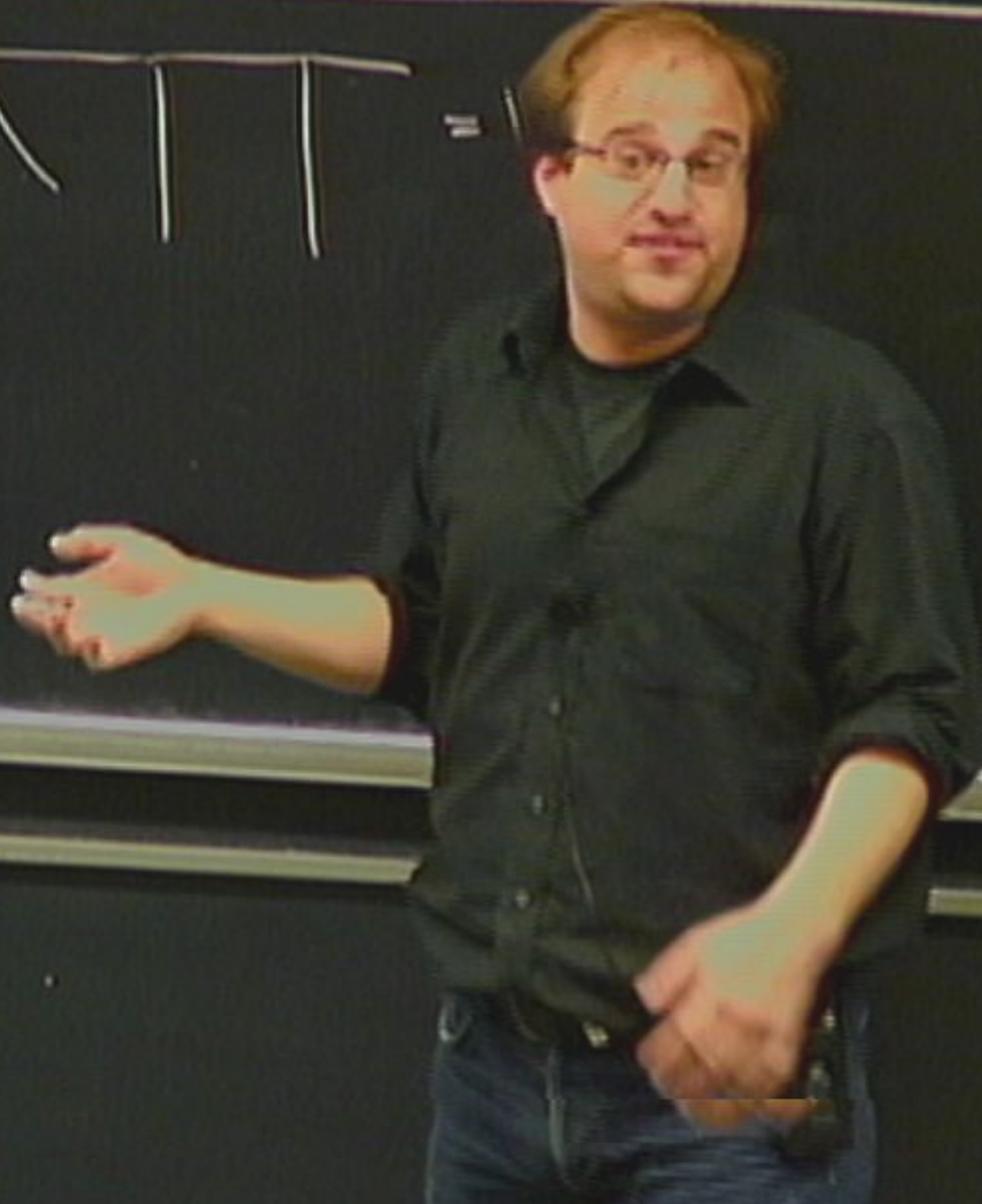
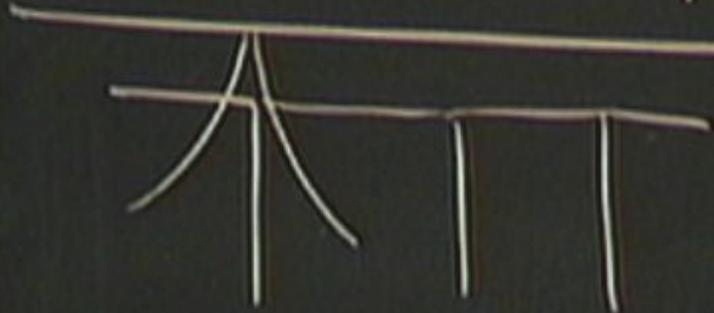
$-x \cdot x$     $-y \cdot y$



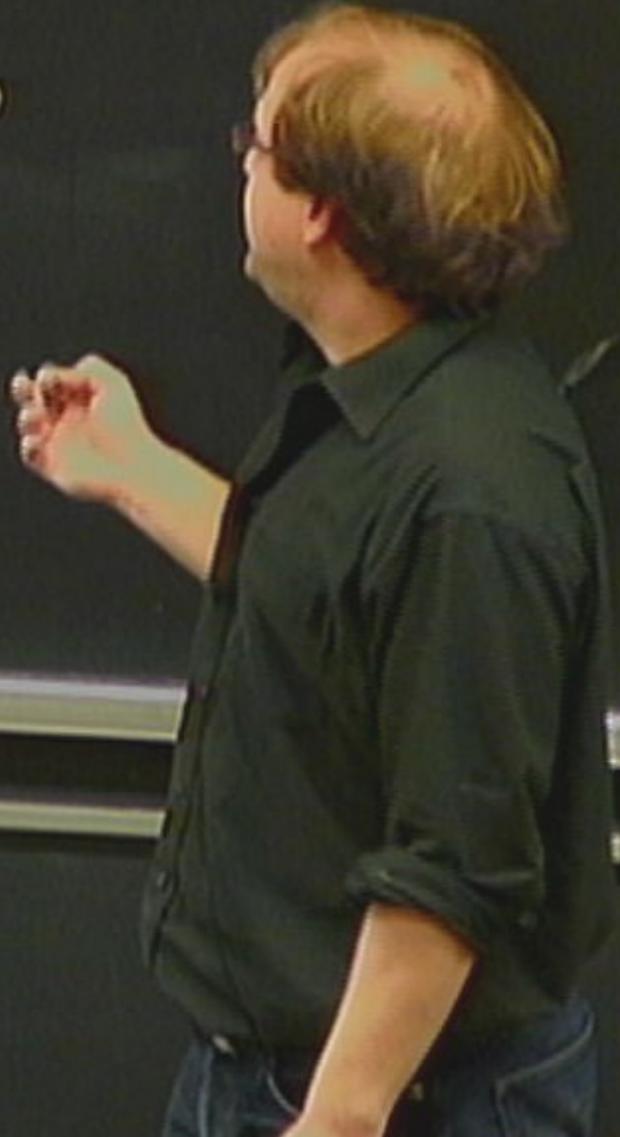
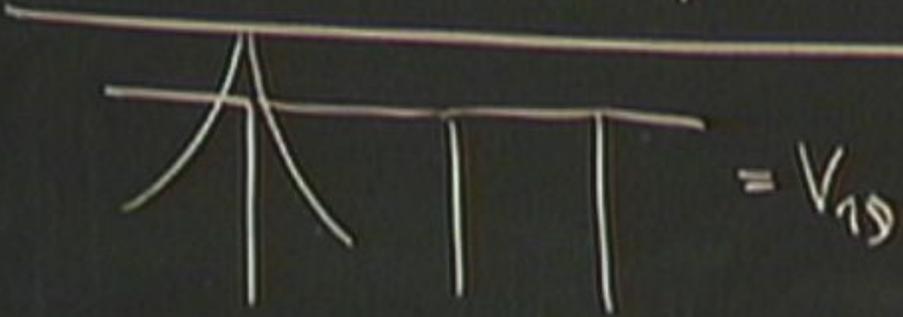
$-x \cdot x$     $-y \cdot y$

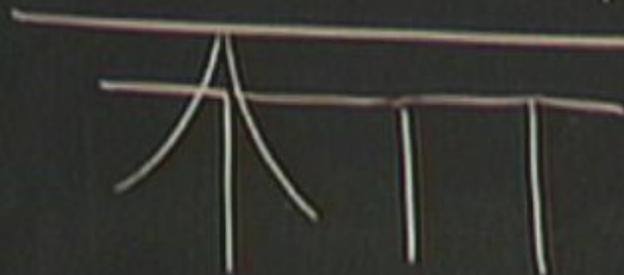


$$\min_{\Omega \rightarrow \rho} \chi[(T+\mathbb{D})\Omega]$$

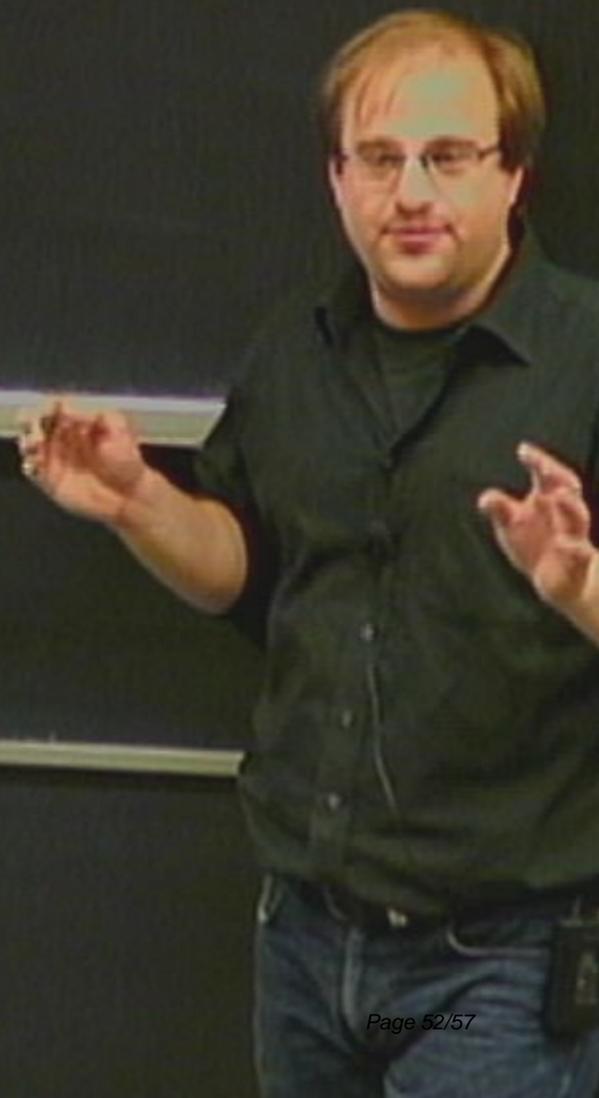


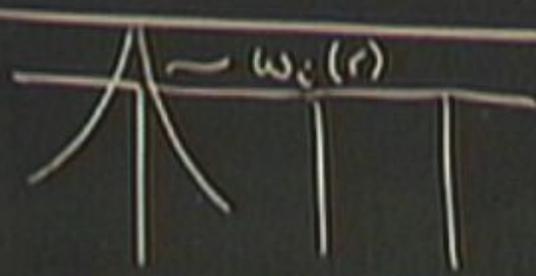
$$\min_{\Omega \rightarrow p} V[(T+D)\Omega]$$





$= V_{19} \rightsquigarrow V(x, y, z) = V(x) + V(y) + V(z)$

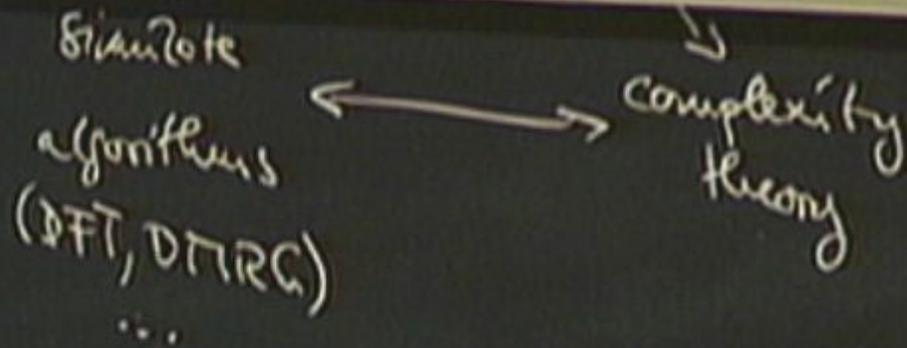



 $\sim w_i(r)$   
 $= V_{19} \rightsquigarrow V(x, y, z) = V(x) + V(y) + V(z)$

$$p(r) = \sum \underbrace{\lambda_{i, s, s'}}_{(s, s')} |sXs'| \cdot |w_i(r)|^2$$

$(s, s')$   
 $\dots$

$$P(r) = \sum_{\substack{\lambda_{i,s,s'} \\ \mathbb{R}^{4N}}} |sXs'| \cdot |w_i(r)|^2$$





$F[\rho]$  is convex!



$F(p)$  is convex!

$$E(p) = F(p) + h[V_p] \text{ is convex!}$$

$F(p)$  is convex!

$E(p) = F(p) + h[V(p)]$  is convex!

→ can be done efficiently!

Hubb. model QNA - complete.