

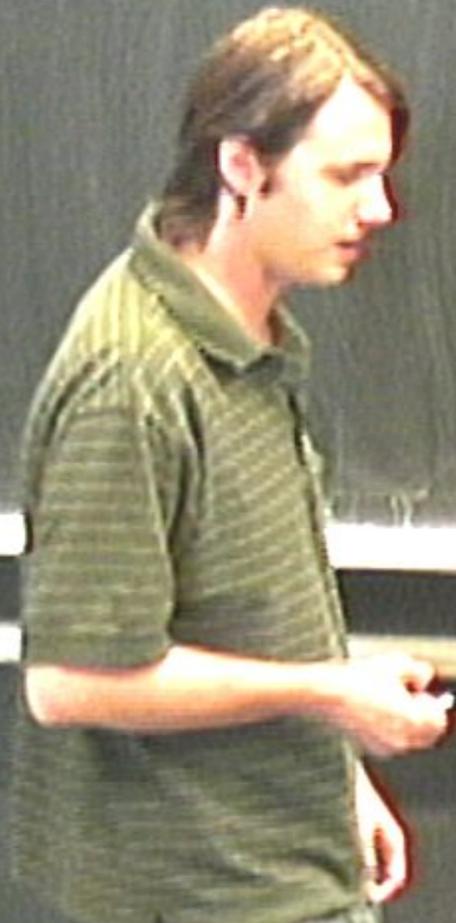
Title: Quantum communication with zero-capacity channels

Date: Sep 24, 2008 04:00 PM

URL: <http://pirsa.org/08090021>

Abstract: A quantum channel models a physical process in which noise is added to a quantum system via interaction with its environment. Protecting quantum systems from such noise can be viewed as an extension of the classical communication problem introduced by Shannon sixty years ago. A fundamental quantity of interest is the quantum capacity of a given channel, which measures the amount of quantum information which can be protected, in the limit of many transmissions over the channel. In this talk, I will show that certain pairs of channels, each with a capacity of zero, can have a strictly positive capacity when used together, implying that the quantum capacity does not completely characterize a channel's ability to transmit quantum information. As a corollary, I will show that a commonly used lower bound on the quantum capacity - the coherent information, or hashing bound - is an overly pessimistic benchmark against which to measure the performance of quantum error correction because the gap between this bound and the capacity can be arbitrarily large.

Jon Yard (LANL)  
Quantum Communication  
with 0-capacity channels  
w/ Greene Smith (IBM)  
Science Express (Aug '08)



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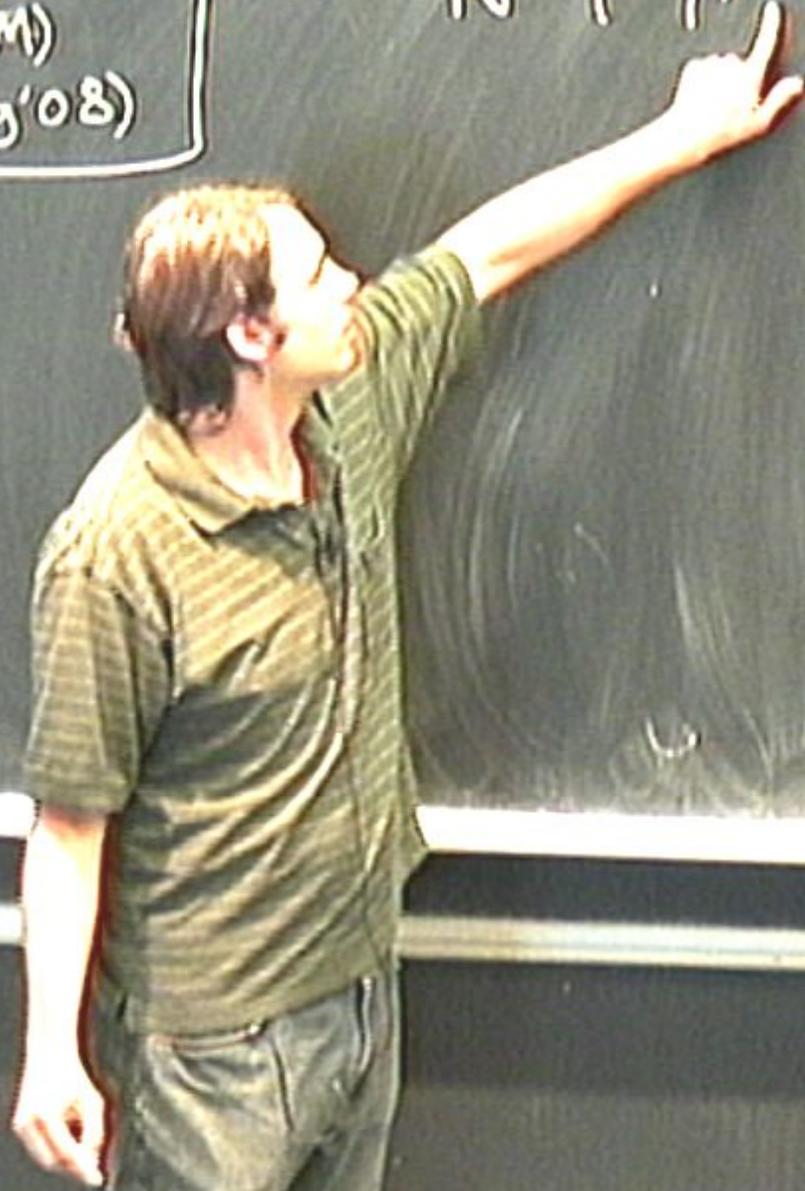
Shannon '48



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Shannon '48

$$N = p(y|x)$$



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Shannon '48

$$N = p(Y|X)$$

Alice

Bob

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Shannon '48

$$N = \rho(Y|X)$$

Alice

$N$

Bob

$N$

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Shannon '48

$$N = p(Y|X)$$

Alice

$N_1$

Bob

$(R, n, \epsilon)$   
code

$N$

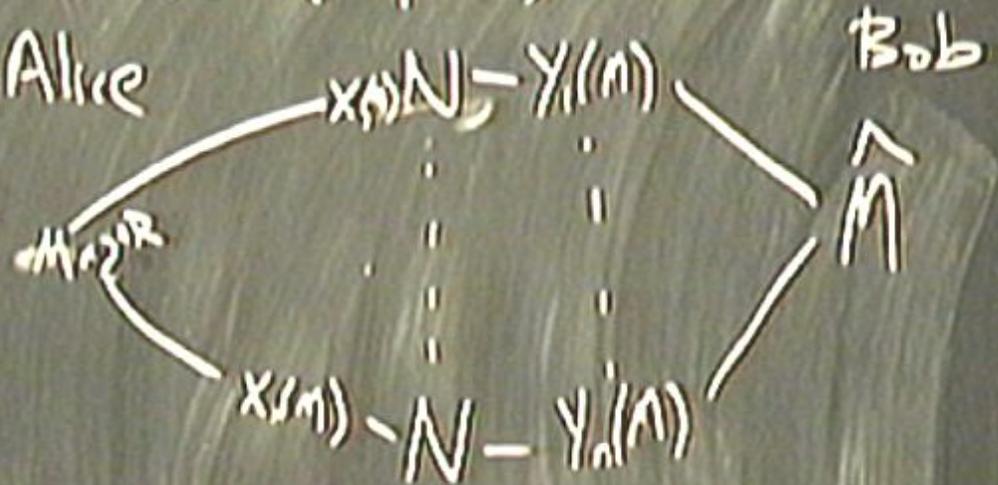
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$(R, n, \epsilon)$   
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Shannon '48

$$N = p(Y|X)$$

Alice

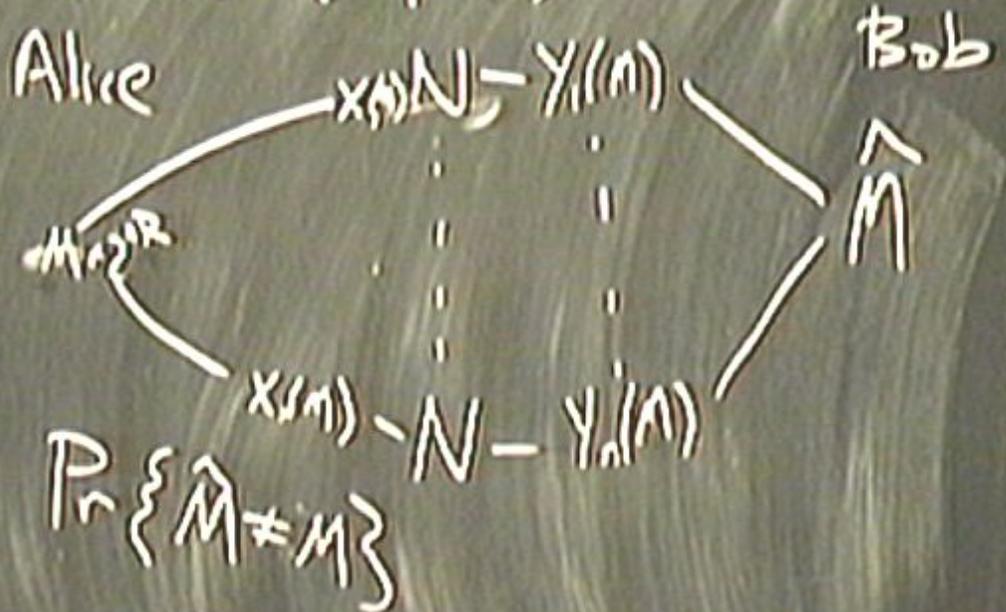


Bob

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Shannon '48

$$N = p(Y|X)$$



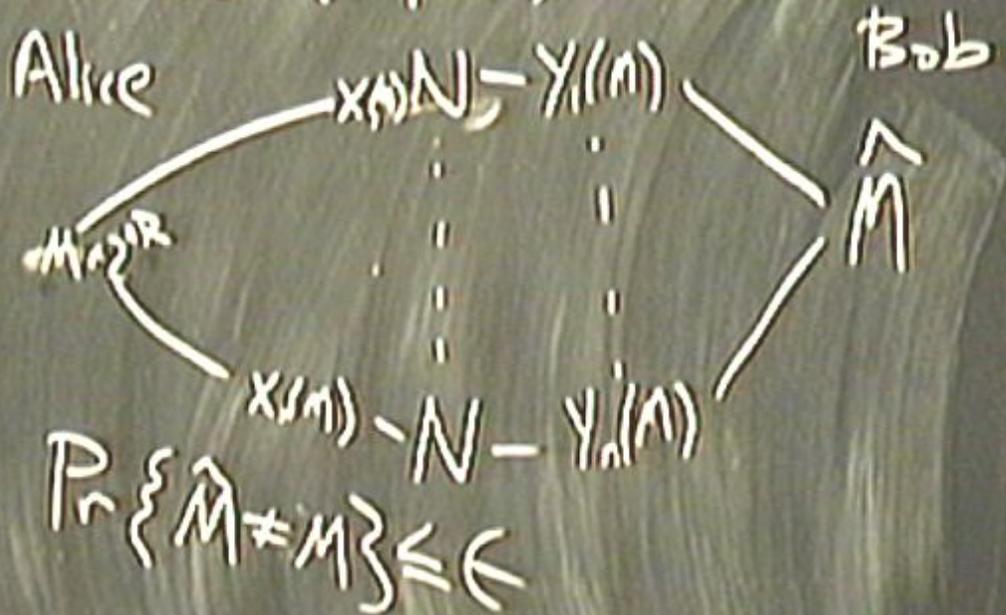
$(R, n, \epsilon)$   
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Shannon '48

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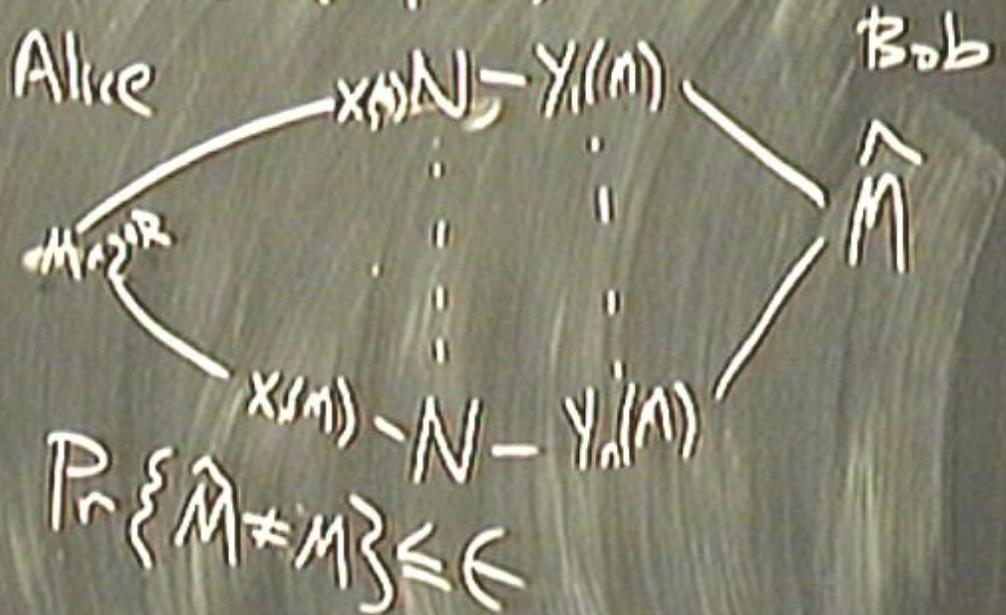


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$(R, n, \epsilon)$   
 code

Shannon '48

$$N = p(Y|X)$$



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Shannon '48

$$N = p(Y|X)$$

Alice

$$x^{(n)} - N - y^{(n)}$$

Bob

$M \in \mathcal{R}$

$\hat{M}$

$$x^{(n)} - N - y^{(n)}$$

$$P\{\hat{M} \neq M\} \leq \epsilon$$

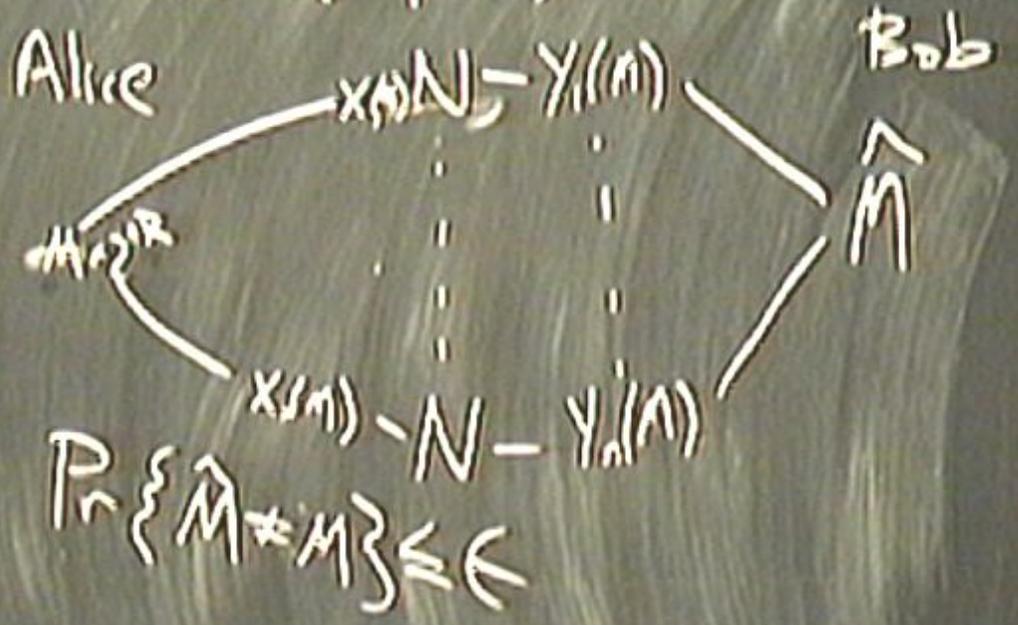
$$(R, n, \epsilon)$$

ach. in  $\exists$  seq  $\uparrow$  code

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Shannon '48

$$N = p(Y|X)$$



$$(R, n, \epsilon)$$

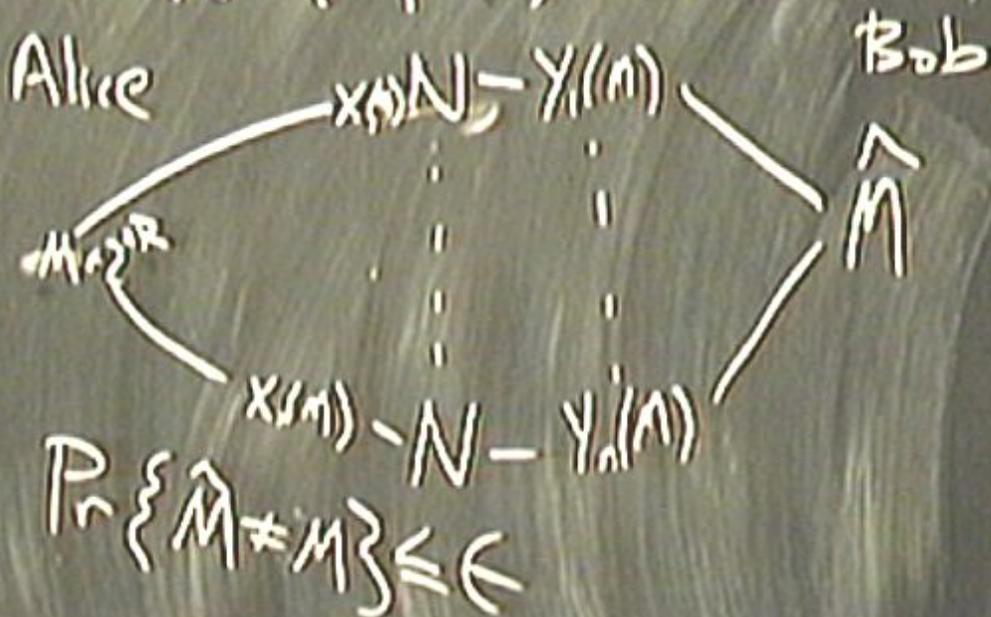
R ach. in } sy ↑ code  
 $\epsilon \rightarrow 0$

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$(R, n, \epsilon)$   
 R ach. in  $n$  sq  $\uparrow$  code  
 $\epsilon \rightarrow 0$

Shannon '48

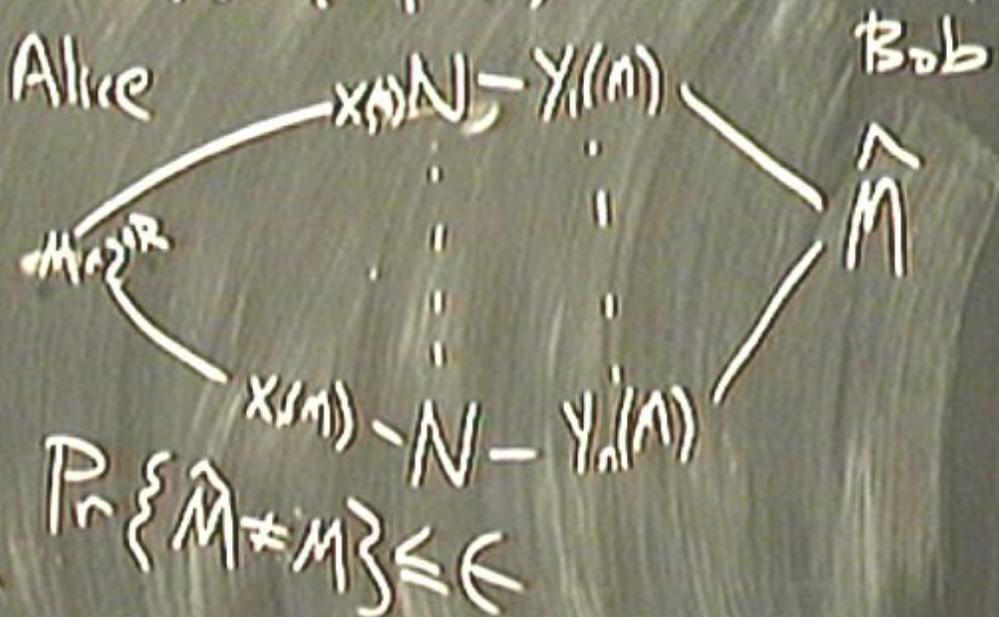
$$N = p(Y|X)$$



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Shannon '48

$$N = p(Y|X)$$



$$(R, n, \epsilon)$$

R ach. in } seq ↑ code  
 $\epsilon \rightarrow 0$

$$C(N) = \text{Sup. ach. rates } R$$

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Shannon '48

$$N = p(Y|X)$$

Alice

$$x^{(n)} - N - y^{(n)}$$

Bob

$M \in \mathcal{R}$

$\hat{M}$

$$(R, n, \epsilon)$$

R ach. if  $\exists$  seq  $\uparrow$  code  
 $\epsilon \rightarrow 0$

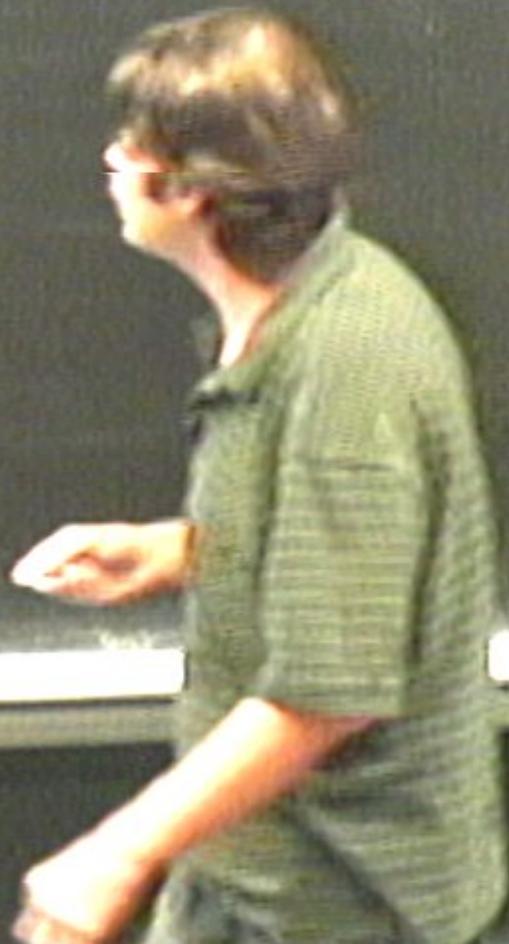
$$C(N) = \text{Sup. ach. rates } R$$

$$P_n \{ \hat{M} \neq M \} \leq \epsilon$$

Thm: Sh '48

$$C(N) = \max_{p(x)} I(X; Y)$$

$$I(X; Y) = H(X) + H(Y) - H(XY)$$



Thm: Sh '48

$$C(N) = \max_{p(n)}$$

$$I(X; Y)$$

$$H(X) + H(Y) - H(XY)$$

Thm: Sh '48

$$C(N) = \max_{p(x)} \underbrace{I(X; Y)}_{H(X) + H(Y) - H(XY)}$$

$$I(X; Y)$$

$$H(X) + H(Y) - H(XY)$$

Thm: Sh '48

$$C(N) = \max_{p(x,y)} \underbrace{I(X;Y)}_{H(X) + H(Y) - H(XY)}$$

$$C(N) = 0 \Leftrightarrow X \perp Y$$

Thm: Sh '48

$$C(N) = \max_{p(x)} \underbrace{I(X; Y)}_{H(X) + H(Y) - H(XY)}$$

$$C(N) = 0 \iff X \perp Y$$

$$C(N_1 \otimes N_2)$$

Thm: Sh '48

$$C(N) = \max_{p(x)} \underbrace{I(X; Y)}_{H(X) + H(Y) - H(XY)}$$

$$C(N) = 0 \iff X \perp Y$$

$$C(N_1 \otimes N_2) = C(N_1) + C(N_2)$$

Thm: Sh '48

$$C(N) = \max_{p(x)} \underbrace{I(X; Y)}_{H(X) + H(Y) - H(XY)}$$

$$C(N) = 0 \Leftrightarrow X \perp Y$$

$$C(N_1 \otimes N_2) = C(N_1) + C(N_2)$$

Thm: Sh '48

$$C(N) = \max_{p(x)} \underbrace{I(X; Y)}_{\text{Additive}} = H(X) + H(Y) - H(XY)$$

$$C(N) = 0 \iff X \perp Y \quad \leftarrow \text{Additive}$$

$$C(N_1 \otimes N_2) = C(N_1) + C(N_2)$$

Thm: Sh 148

$$C(N) = \left( \max_{p(x)} I(X; Y) \right) \left( H(X) + H(Y) - H(X, Y) \right) = C^{(1)}(N)$$

$$C(N) = 0 \iff X \perp Y \quad \text{--- Additive}$$

$$C(N_1 \otimes N_2) = C(N_1) + C(N_2)$$

Thm: Sh 148

$$C(N) = \left( \max_{p(x)} I(X; Y) \right) \left( H(X) + H(Y) - H(X, Y) \right)$$

$C(N) = C^{(1)}(N)$

$$C(N) = 0 \iff X \perp Y$$

$$C(N_1 \otimes N_2) = C(N_1) + C(N_2)$$

Thm: Sh 148

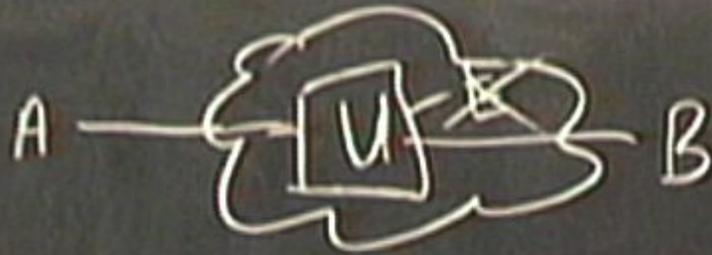
$$C(N) = \left( \max_{p(x)} \underbrace{I(X; Y)}_{C^{(1)}(N)} \right) \frac{H(X) + H(Y) - H(XY)}{C^{(1)}(N)}$$

$$C(N) = C^{(1)}(N)$$

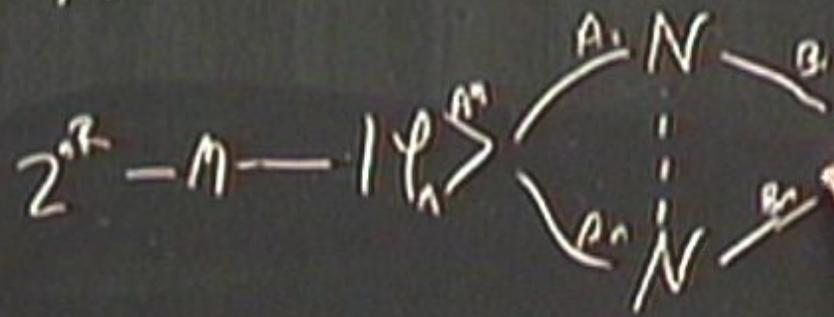
$$C(N) = 0 \iff X \perp Y \quad \leftarrow \text{Additive}$$

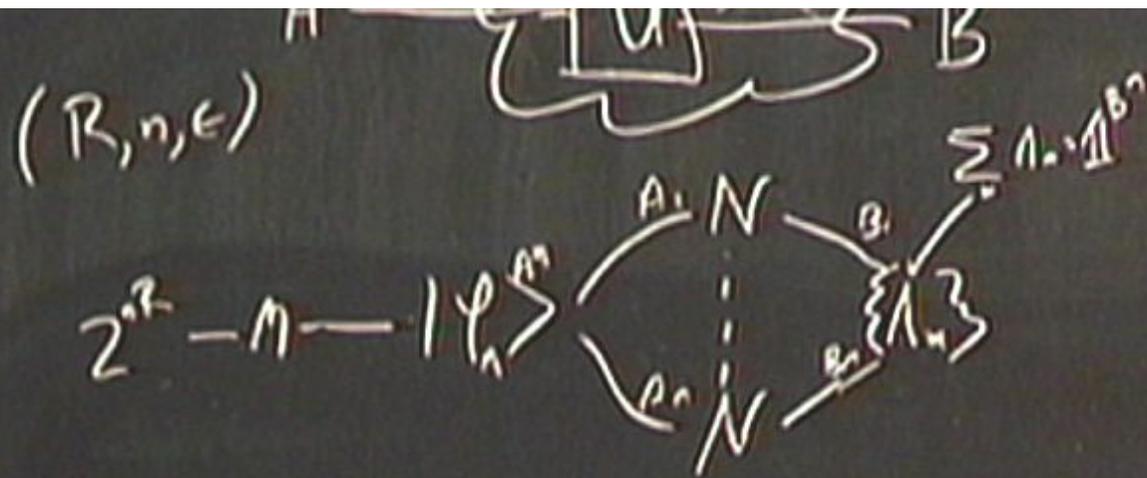
$$C(N_1 \otimes N_2) = C(N_1) + C(N_2)$$

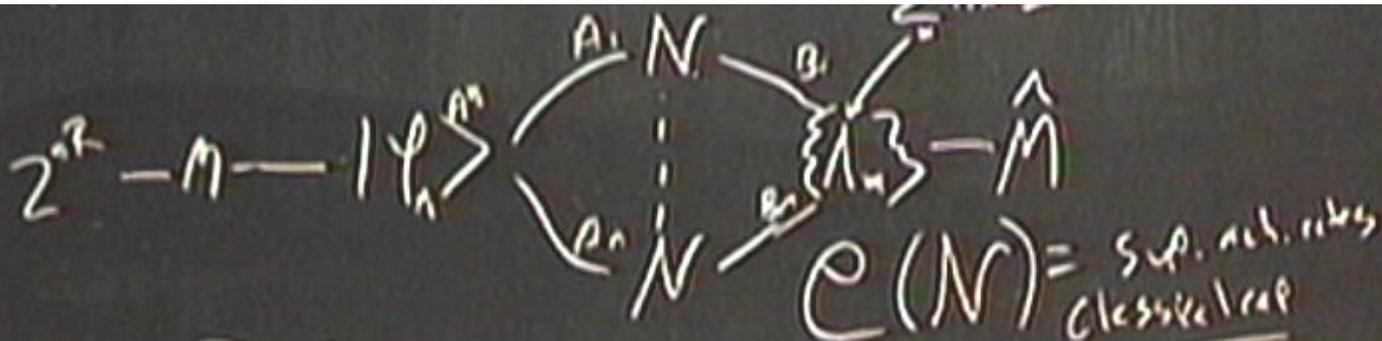
$N^{A \rightarrow B}$



$(R, n, \epsilon)$   $A \xrightarrow{\{U\}} B$

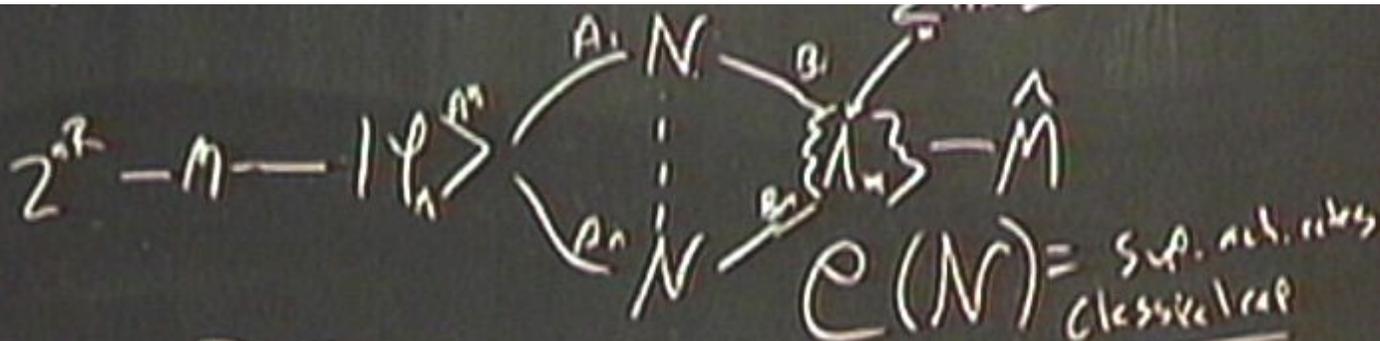






$$\Pr \{ \hat{M} \neq M \} \leq \epsilon$$

Thm (HSW).  
 $\mathcal{P}(N) = \lim_{n \rightarrow \infty} \dots$



$$\Pr\{\hat{M} \neq M\} \leq \epsilon$$

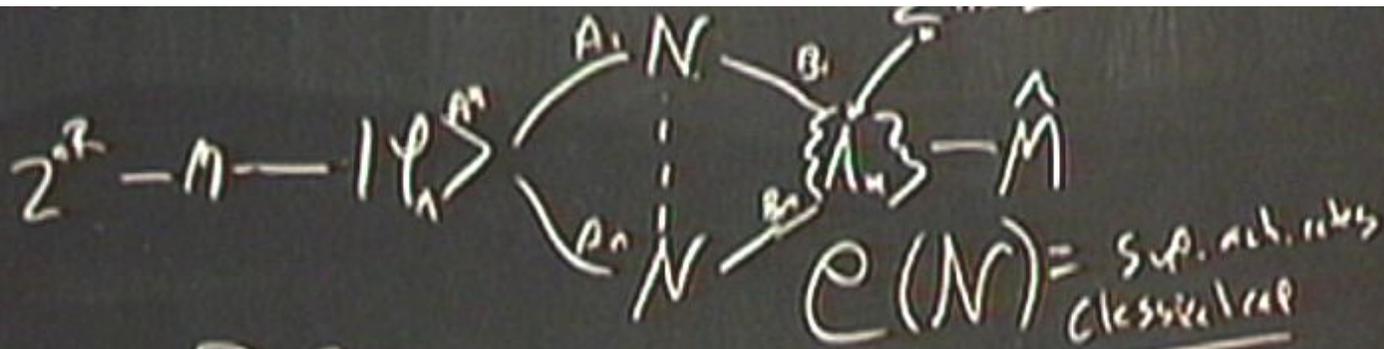
Thm (HSW).

$$\mathcal{C}(N) = \lim_{n \rightarrow \infty} \frac{1}{n} \mathcal{C}^{(n)}(N^{\otimes n})$$

$$\mathcal{C}^{(n)}(N) = \max_{\{P, H\}} \dots$$

$$I(X; B) =$$

$$= \chi(\xi \dots)$$



$$P_{\{\hat{M} \neq M\}} \leq \epsilon$$

Thm (HSW).

$$C(N) = \lim_{n \rightarrow \infty} \frac{1}{n} C(N^{\otimes n})$$

$$C^{(n)}(N) = \max_{\{P, P'\}} \dots$$

$$I(X; B) =$$

$$= \chi(\{P, N(\varphi_x)\})$$

$$P_r \{ \hat{M} \neq M \} \leq \epsilon$$

Thm (HSW).

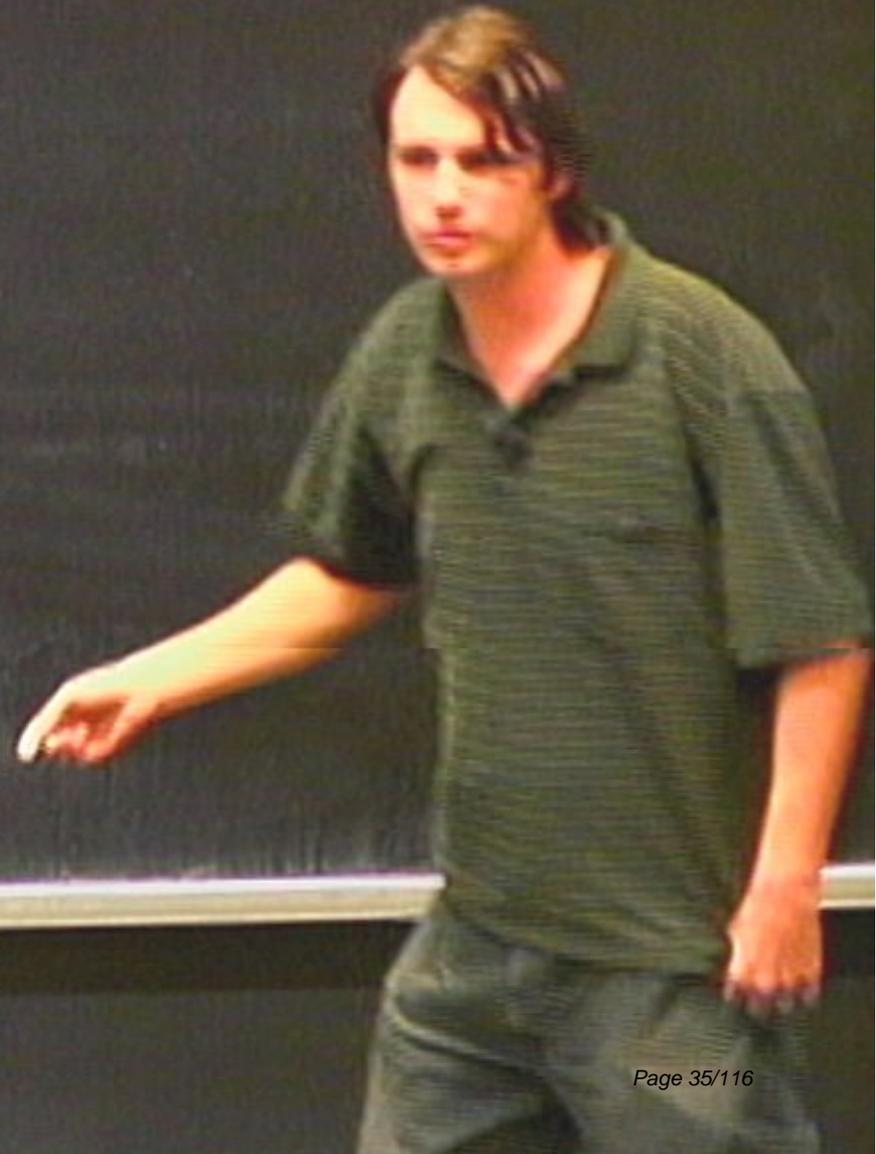
$$C(N) = \lim_{n \rightarrow \infty} \frac{1}{n} \log C(N^{ot n})$$

$$C^{(1)}(N) = \max_{\mathcal{E}_r, \mathcal{P}_x} \{ \dots \}$$

$$I(X; B) =$$

$$= \chi(\mathcal{E}_r, N(\mathcal{P}_x))$$

Thm. (Hastings):  $e^{(n)}(N) < e(N) \quad \exists N$



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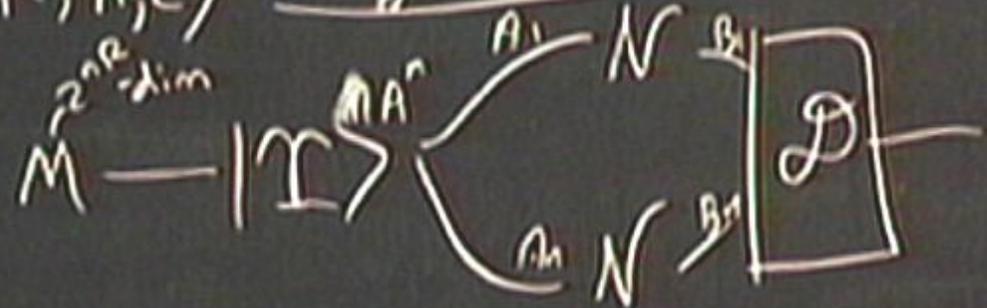
Thm: (Hastings):  $\underbrace{O^m(N)} < \underbrace{O(N)} \quad \exists N$

$(R, n, \epsilon)$  ent. generation

M

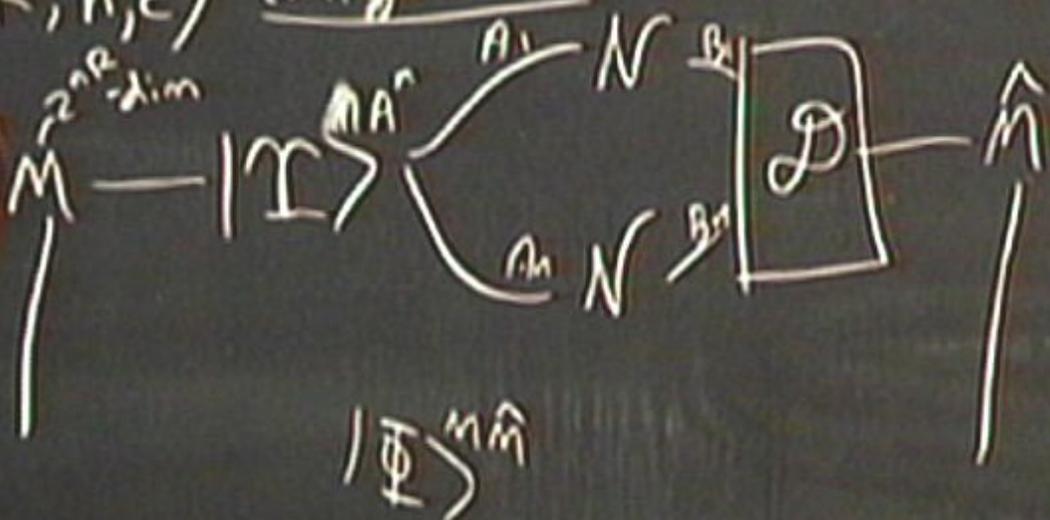
Thm: (Hastings):  $\underline{C^{\infty}(N)} < \underline{C(N)} \quad \exists N$

$(R, n, \epsilon)$  ent. generation



Thm: (Hofstadter):  $\mathcal{P}^{(n)}(N) < \mathcal{P}(N) \quad \exists N$

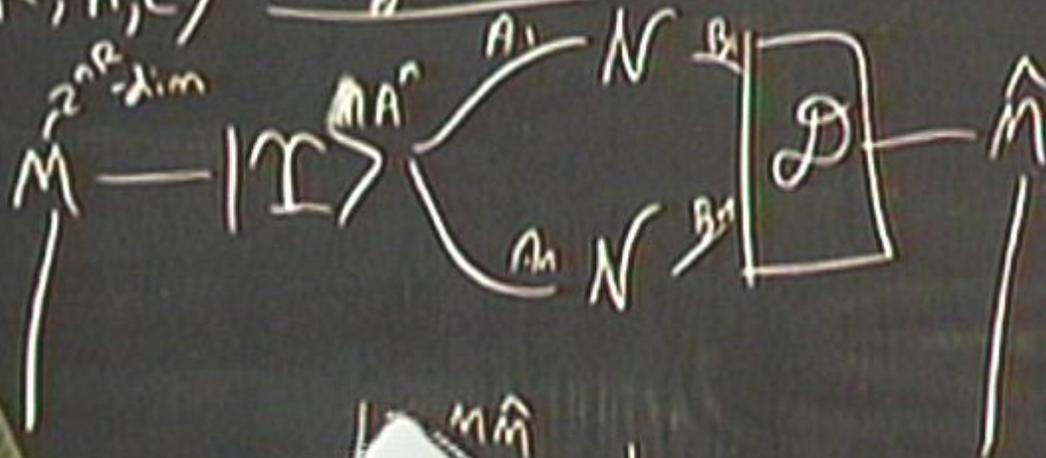
$(R, n, \epsilon)$  int. generation



$I \hat{I}$

Thm: (Hastings):  $\underline{C^{\infty}(N)} < \underline{C(N)} \quad \exists N$

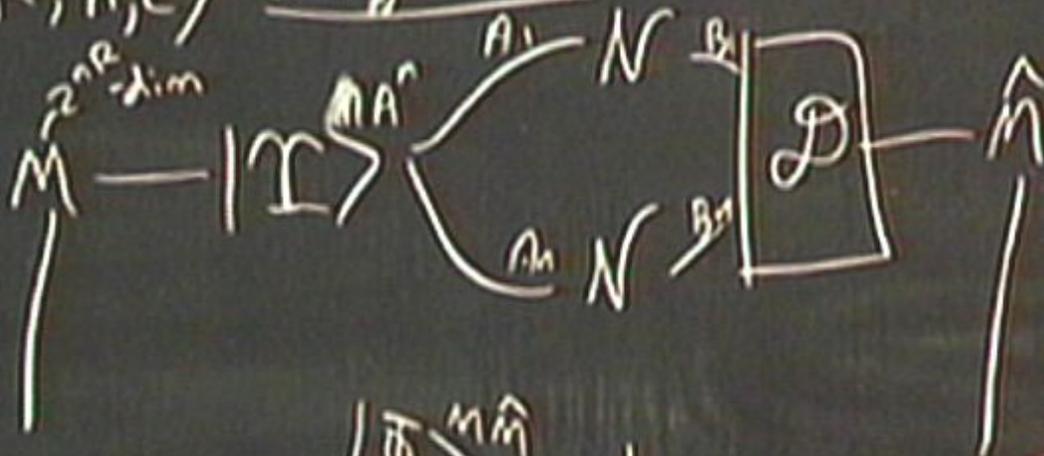
$(R, n, \epsilon)$  ent. generation



$\hat{M} = \frac{1}{\sqrt{2^n R}} \sum_{\mathbf{m}} |\mathbf{m}\rangle |\mathbf{m}\rangle$

Thm: (Hastings):  $\underline{C^{\infty}(N)} < \underline{C(N)} \quad \exists N$

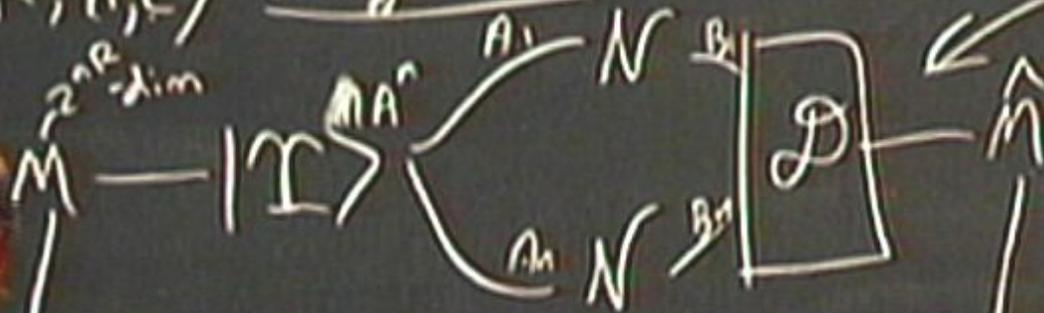
$(R, n, \epsilon)$  ent. generation



$$|\Phi\rangle^{\hat{m}\hat{n}} = \frac{1}{\sqrt{2nR}} \sum_{\hat{m}} |\hat{m}\rangle |\hat{n}\rangle$$

Thm: (Hasting):  $\mathcal{P}^{||}(N) < \mathcal{P}(N) \exists N$

$(R, n, \epsilon)$  ent. generation

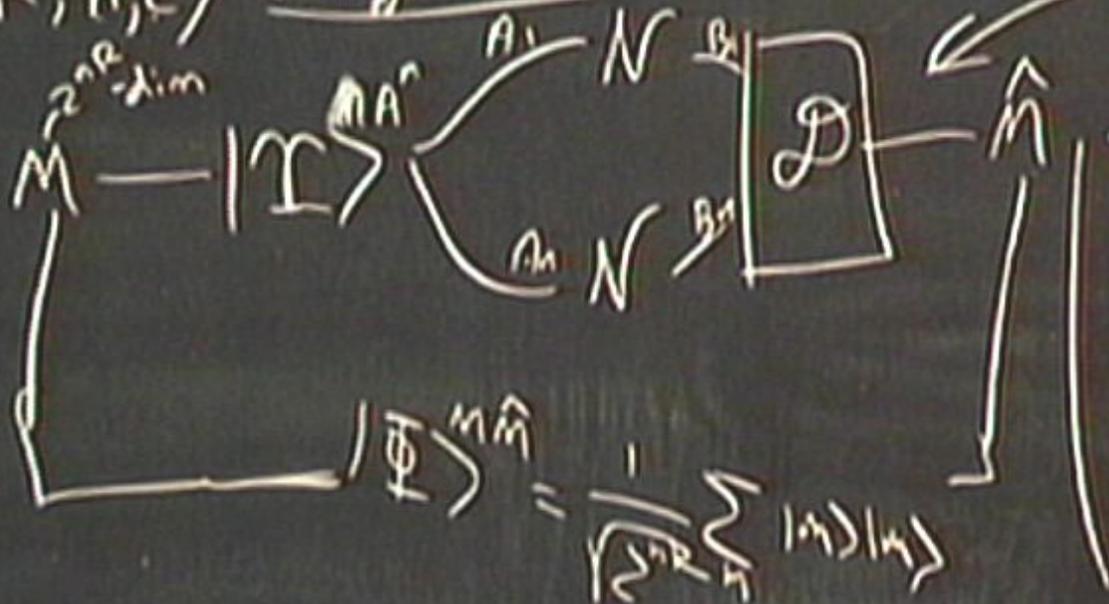


$$\langle \hat{\Phi} | \hat{\Phi} \rangle \geq 1 - \epsilon$$

$$|\hat{\Phi}\rangle = \frac{1}{\sqrt{2^n R}} \sum_{i,j} |i\rangle |j\rangle$$

Thm: (Hasting):  $\underline{C^{(1)}(N)} < \underline{C(N)} \quad \exists N$

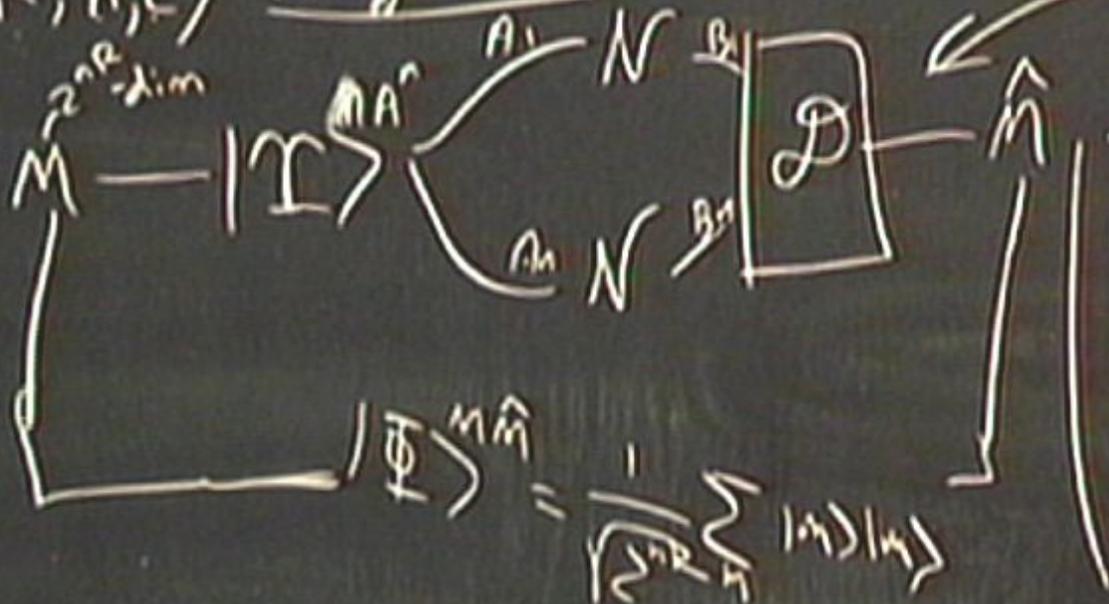
$(R, n, \epsilon)$  ent. generation



$$\langle \hat{\psi} | \psi \rangle^{m, n} \geq 1 - \epsilon$$

Thm: (Höding) :  $\underline{C}''(N) < \underline{C}(N) \exists N$

$(R, n, \epsilon)$  ent. generation

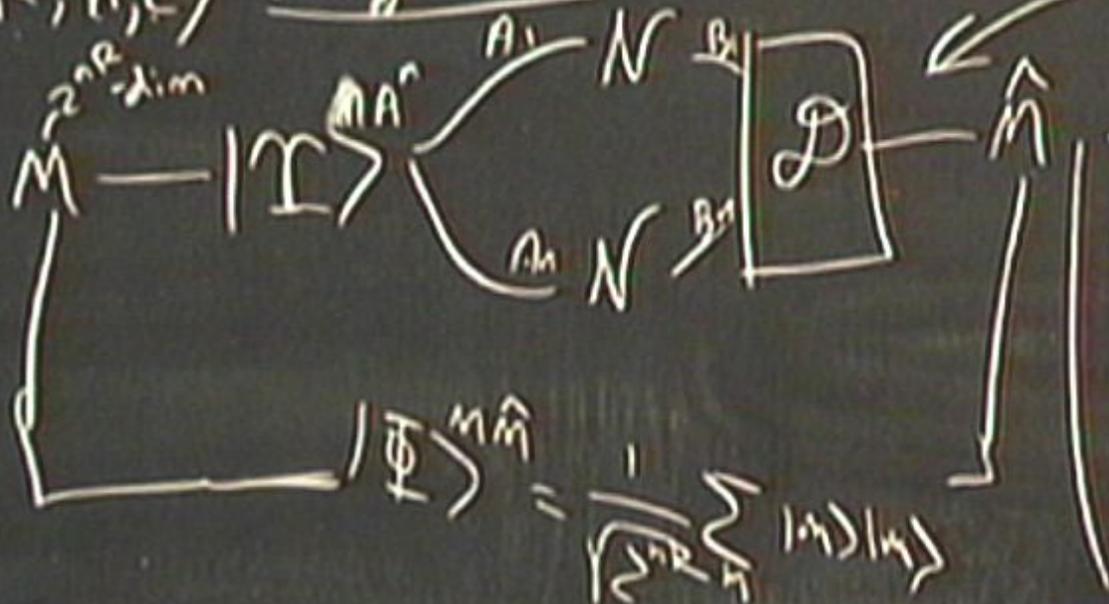


$$\langle \hat{\Phi} | \hat{\Phi} | \hat{\Phi} \rangle \geq 1 - \epsilon$$

$$Q(N) = \sup_{\text{ach. rels } R}$$

Thm: (Holtings):  $\underline{C}^{\infty}(N) < \underline{C}(N) \exists N$

$(R, n, \epsilon)$  ent. generation



$$\langle \hat{\Phi} | \hat{\Phi} | \hat{\Phi} \rangle \geq 1 - \epsilon$$

$$Q(N) = \sup_{\text{ach. rels } \mathbb{R}}$$

Thm LSD  $Q(N) = \lim_{n \rightarrow \infty} \frac{1}{n}$



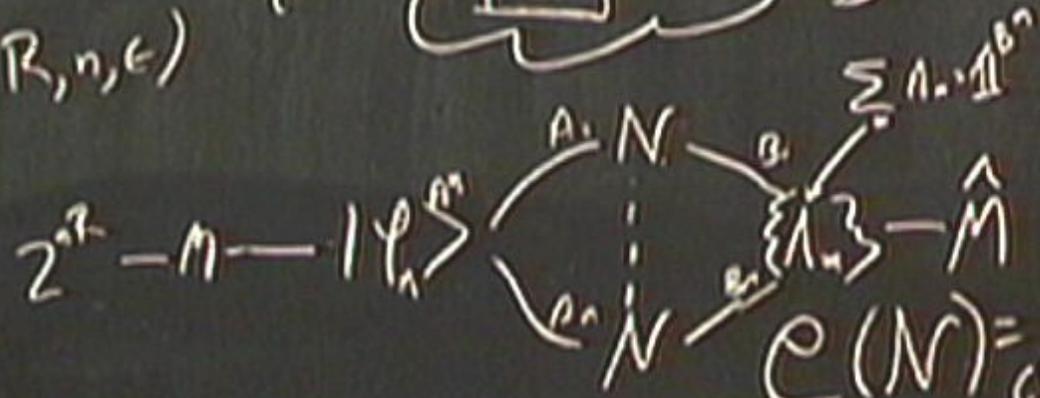
Thm LSD  $Q(W) = \lim_{n \rightarrow \infty} \frac{1}{n} Q^{(n)}(N^{(n)})$   
 $Q^{(n)}(N) = \max_{P_A} [H(B) - H(E)]$

$$\text{Thm LSD } Q(N) = \lim_{n \rightarrow \infty} \frac{1}{n} Q^{(n)}(N^{2^n})$$
$$Q^{(n)}(N) = \max_{P \in \mathcal{A}} [H(B) - H(E)]$$

$$P^{(n)}(N) < P(N) \quad \exists \Delta$$

$$N^A \rightarrow B$$

$$(R, n, \epsilon)$$



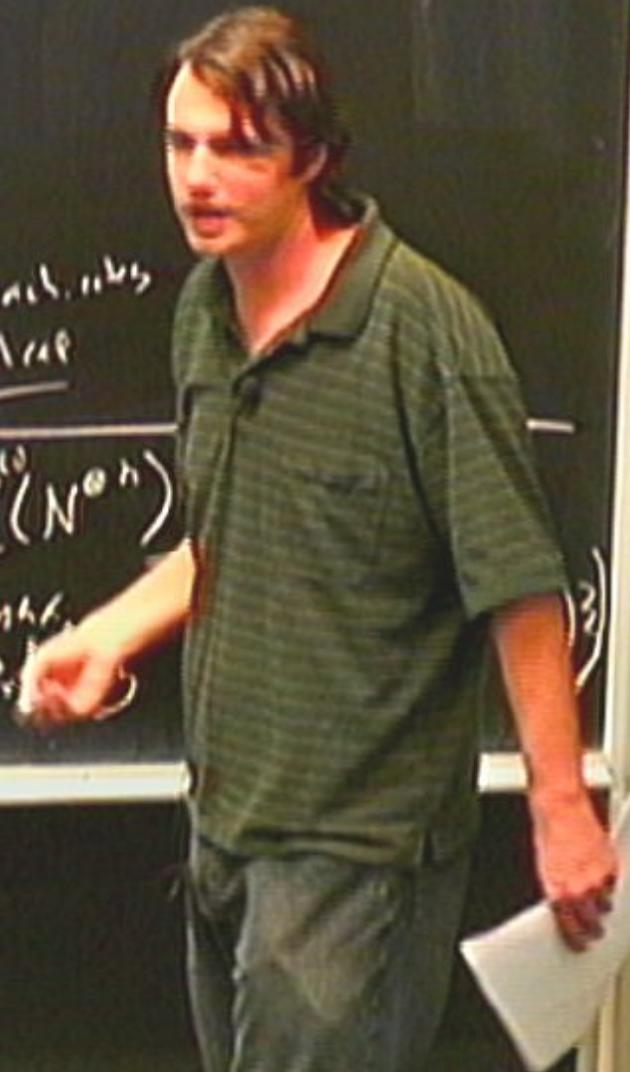
$$C(N) = \frac{\text{Sup. ach. rates}}{\text{classical cap}}$$

$$P_r \{ \hat{M} \neq M \} \leq \epsilon$$

Thm (HSW).

$$C(N) = \lim_{n \rightarrow \infty} \frac{1}{n} C(N^{\otimes n})$$

$$C^{(n)}(N) = \max_{\epsilon P, \dots}$$



$$\text{Thm LSD } Q(W) = \lim_{n \rightarrow \infty} \frac{1}{n} Q^{(n)}(N^n)$$
$$Q^{(n)}(N) = \max_{\rho_A} [H(B) - H(E)]$$

$$\text{Thm LSD } Q(N) = \lim_{n \rightarrow \infty} \frac{1}{n} Q^{(n)}(N^{2^n})$$

$$Q^{(n)}(N) = \max_{P \in \mathcal{A}} [H(B) - H(E)]$$

$$\text{Thm: SS '96 } Q(N) > Q^{(n)}(N) = 0$$

Thm LSD  $Q(N) = \lim_{n \rightarrow \infty} \frac{1}{n} Q^{(n)}(N^{(n)})$

$Q^{(n)}(N) = \max_{\rho_A} [H(B) - H(E)]$  ← relevant info.

Thm: SS '96  $Q(N) > Q^{(n)}(N) = 0$

Thm LSD  $Q(N) = \lim_{n \rightarrow \infty} \frac{1}{n} Q^{(n)}(N^{2^n})$

$Q^{(n)}(N) = \max_{P \in \mathcal{A}} [H(B) - H(E)]$  ← relevant info.

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relevant info.

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Thm. SS '96  $Q(N) > Q^{(n)}(N) = 0$

Thm LSD  $Q(N) = \lim_{n \rightarrow \infty} \frac{1}{n} Q^{(n)}(N^{2^n})$

$Q^{(n)}(N) = \max_{P \in \mathcal{A}} [H(B) - H(E)]$

relevant info.

Thm: SS '96  $Q(N) > Q^{(n)}(N) = 0$

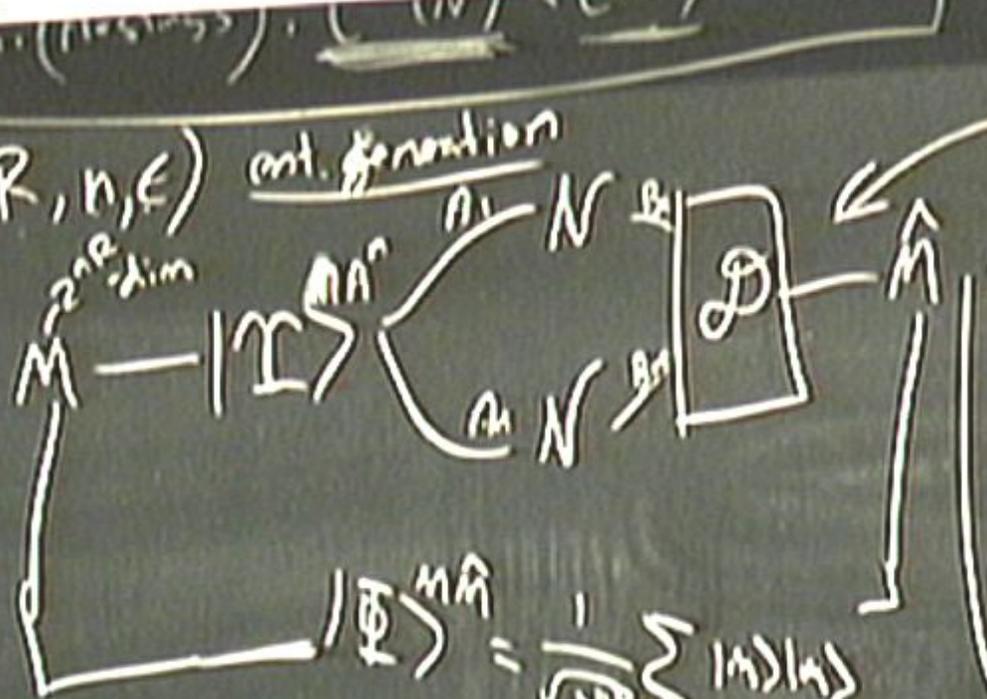


$$P_{\{ \hat{M} \neq M \}} \leq \epsilon$$

Thm. (HSW).  
 $C(N) = \lim_{n \rightarrow \infty} \frac{1}{n} \log C(N^{(n)})$   
 $C^{(n)}(N) = \max_{\mathcal{E}, \mathcal{P}} \{ \mathbb{I}(X; B) \}$   
 $= \chi(\mathcal{E}, N(\Psi_N^M))$

$(R, n, \epsilon)$

int. generation



$\langle \mathbb{I} | \hat{\Phi} | \mathbb{I} \rangle \geq 1 - \epsilon$

$Q(N) = \text{sup. ach. icht } \mathbb{R}$

Thm LSD  $Q(N) = \lim_{n \rightarrow \infty} \frac{1}{n} Q^{(1)}(N^{(n)})$

$$Q^{(1)}(N) = \max_{P_A} [H(B) - H(E)]$$

relevant  
info.

Thm. SS '96  $Q(N) > Q^{(1)}(N) = 0$

Thm LSD  $Q(N) = \lim_{n \rightarrow \infty} \frac{1}{n} Q^{(n)}(N^{2^n})$

$Q^{(n)}(N) = \max_{\rho \in \mathcal{A}} [H(B) - H(E)]$

relevant info.

Thm: SS '96  $Q(N) > Q^{(n)}(N) = 0$

Thm LSD  $Q(N) = \lim_{n \rightarrow \infty} \frac{1}{n} Q^{(1)}(N^{(n)})$

$Q^{(1)}(N) = \max_{p \in \mathcal{A}} [H(B) - H(E)]$

relevant info.

Thm: SS '96  $Q(N) > Q^{(1)}(N) = 0$

Thm LSD  $Q(N) = \lim_{n \rightarrow \infty} \frac{1}{n} Q^{(n)}(N^{(n)})$

$Q^{(n)}(N) = \max_{P_A} [H(B) - H(E)]$

relevant info.

Mc, Highton writes, Highton

Thm: SS '16  $Q(N) > Q^{(n)}(N) = 0$



Thm LSD  $Q(N) = \lim_{n \rightarrow \infty} \frac{1}{n} Q^{(n)}(N^{(n)})$

$Q^{(n)}(N) = \max_{\rho_A} [H(B) - H(E)]$

relevant info.

me, Highton  
writes, Highton

Thm: SS '76  $Q(N) > Q^{(n)}(N) = 0$

Then LSD  $Q(N) = \lim_{n \rightarrow \infty} \frac{1}{n} Q^{(n)}(N^{(n)})$

$Q^{(n)}(N) = \max_{p_A} [H(B) - H(E)]$

relevant info.

me, Highton write, Highton

Then SS 'T6  $Q(N) > Q^{(n)}(N) = 0$

$$\text{Thm LSR } Q(N) = \lim_{n \rightarrow \infty} \frac{1}{n} Q^{(n)}(N^{(n)})$$

$$\underline{Q^{(n)}(N)} = \max_{p \in \mathcal{A}} [H(B) - H(E)]$$

relevant info.

Mc, Highton  
writes, Highton

$$\text{Thm SS '96 } \underline{Q(N) > Q^{(n)}(N) = 0}$$

$$\text{Thm (SY '08)} \exists N_1, N_2 \text{ s.t. } Q(N_1, N_2)$$

$$\text{Thm LSR } Q(N) = \lim_{n \rightarrow \infty} \frac{1}{n} Q^{(n)}(N^{(n)})$$

$$\underline{Q^{(n)}(N)} = \max_{P^A} [H(B) - H(E)]$$

relevant  
info.

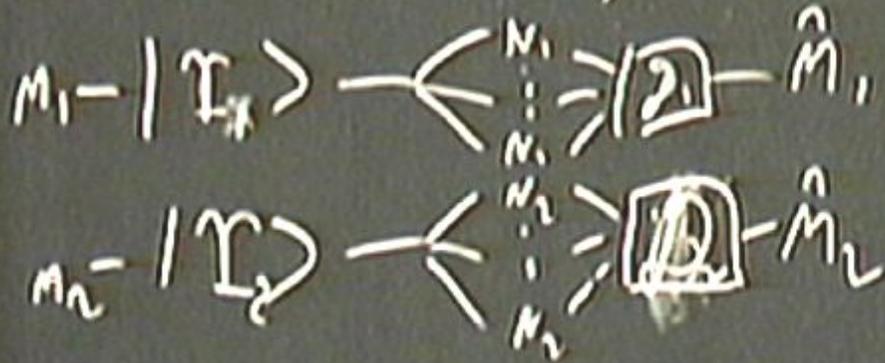
Mc, Hight  
write, Hight

$$\text{Thm SS '96 } \underline{Q(N) > Q^{(n)}(N) = 0}$$

$$\text{Thm (SY '08)} \exists N_1, N_2 \text{ s.t. } Q(N_1 \circ N_2) > Q(N_1) + Q(N_2) = 0$$

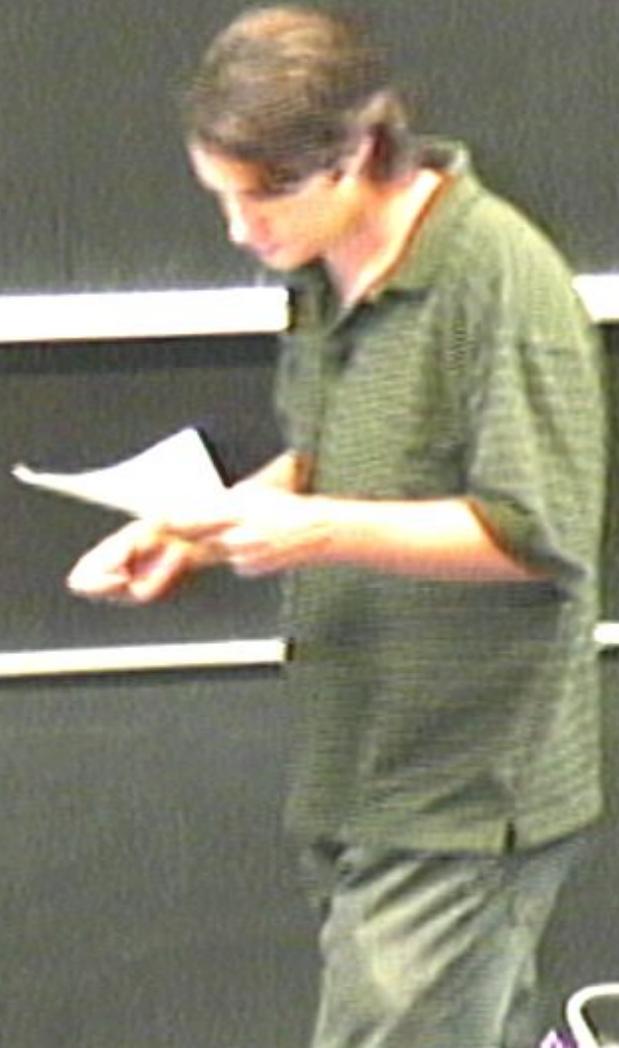
Thm 16  $Q(N) > Q'(N) = 0$

Thm (SY '08)  $\exists N_1, N_2$  s.t.  $Q(N_1 \circ N_2) > Q(N_1) + Q(N_2) = 0$



Thm  $\Rightarrow$  (6)  $Q(N) > Q(N) = 0$

Thm (SY '08)  $\exists N_1, N_2$  s.t.  $Q(N_1 \circ N_2) > Q(N_1) + Q(N_2) = 0$



$$\text{Thm } \Rightarrow \text{ (6) } Q(N) > Q(N) = 0$$

$$\text{Thm (SY '08)} \exists N_1, N_2 \text{ s.t. } Q(N_1 \circ N_2) > Q(N_1) + Q(N_2) = 0$$



Thm LSR  $Q(N) = \lim_{n \rightarrow \infty} \frac{1}{n} Q^{(n)}(N^{(n)})$

$Q^{(n)}(N) = \max_{P \in \mathcal{A}} [H(B) - H(E)]$  ← relevant info.

me, Highten write, handw.

Thm SS '06  $Q(N) > Q^{(n)}(N) = 0$

Thm (SY '08)  $\exists N_1, N_2$  s.t.  $Q(N_1 \circ N_2) > Q(N_1) + Q(N_2) = 0$



Thm LSR  $Q(N) = \lim_{n \rightarrow \infty} \frac{1}{n} Q^{(n)}(N)$

$Q^{(n)}(N) = \max_{P_A} [H(B) - H(E)]$

relevant info.

Mc, Hgdon  
wrote, Hordosh

Thm SS '96  $Q(N) > Q^{(n)}(N) = 0$

Thm (SY '08)  $\exists N_1, N_2$  s.t.  $Q(N_1 \circ N_2) > Q(N_1) + Q(N_2) = 0$



$Q_2(N) = 0$  when?

$Q_1(N) = 0$  when?

$Q(N)=0$  when?

1. Symmetric channels

$Q(N)=0$  when?

1. Symmetric channels

$$U \rightarrow (BE)_+$$

$Q_2(N) = 0$  when?

1. Symmetric channels  $U^A \rightarrow (BE)_+$   $Q_2^{(1)} = 0$   $H(B) = H(E)$   
 $Q_2 = 0$

2.

$Q_2(N) = 0$  when?

1. Symmetric channels  $U^A \rightarrow (BE)_+$   $Q_2^{(1)} = 0$   $H(B) = H(E)$   
 $Q_2 = 0$

2. Hrodocki channels

$Q_2(N) = 0$  when?

1. Symmetric channels  $U^A \rightarrow (BE)_+$   $Q_2^{(1)} = 0$   $H(B) = H(E)$   
 $Q_2 = 0$

2. Horodecki channels (PPT, entang. binding)

$Q(N)=0$  when?

1. Symmetric channels  $U^A \rightarrow (BE)_+$   $Q^{(1)}=0$   $H(B)=H(E)$   
 $Q_2=0$

2. Horodecki channels (PPT, entang. binding)

$$\rho^{AB} = (\mathbb{1}^A \otimes W)(\rho^{AB})$$

$Q_2(N) = 0$  when?

1. Symmetric channels  $U^A \rightarrow (BE)_+$   $Q_2^{(1)} = 0$   $H(B) = H(E)$   
 $Q_2 = 0$

2. Horodecki channels (PPT, entang. binding)

$$\rho^{AB} = (\mathbb{1}^A \otimes W)(\rho^{AB}) \quad \rho^{AB} = \sum_{i,j} |i\rangle\langle j| \otimes \rho_{ij}^T$$

$Q_2(N) = 0$  when?

1. Symmetric channels  $U^A \rightarrow (BE)_+$   $Q_2^{(1)} = 0$   $H(B) = H(E)$   
 $Q_2 = 0$

2. Horodecki channels (PPT, entang. binding)

$$\rho^{AB} = (\mathbb{1}^A \otimes W)(\rho^{AB}) \quad \rho^{AB} = \sum_{i,j} \lambda_i \times \lambda_j |i\rangle\langle j|^T$$

$Q(N)=0$  when?

1. Symmetric channels  $U^{A \rightarrow (BE)}$   $Q^{(1)}=0$   $H(B)=H(E)$   
 $Q=0$

2. Horodecki channels (PPT, entang. binding)

$$\rho^{AB} = (\mathbb{1}^A \otimes W)(\rho^{AB}) \quad \rho^{AB} = \sum_{i,j} \lambda_{ij} x_{ij} |i\rangle\langle j|^T$$

$Q(N)=0$  when?

1. Symmetric channels  $U^A \rightarrow (BE)_+$   $Q^{(1)}=0$   $H(B)=H(E)$   
 $Q_2=0$

Product channels (PPT, entang. binding)

$$\rho^{AB} = (\mathbb{I}^A \otimes W)(\rho^{AB}) \quad \rho^{AB} = \sum_{i,j} \lambda_i \lambda_j^* |i\rangle\langle j| \otimes \rho_i$$

$Q(N)=0$  when?

→ 1. Symmetric channels  $U^{A \rightarrow (BE)}$   $Q^{(1)}=0$   $H(B)=H(E)$   
 $Q=0$

→ 2. Horodecki channels (PPT, entang. binding)

$$\rho^{AB} = (\mathbb{1}^A \otimes W)(\rho^{AB}) \quad \rho^{AB} = \sum_{i,j} |i\rangle\langle j| \otimes T_{ij}$$

$Q(N)=0$  when?

→ 1. Symmetric channels  $U^A \rightarrow (BE)_+$   $Q^{(1)}=0$   $H(B)=H(E)$   
 $Q_2=0$

→ 2. Horodecki channels (PPT, entang. binding)

$$\rho^{AB} = (\mathbb{1}^A \otimes W)(\rho^{AB}) \quad \rho^{AB} = \sum_{i,j} |i\rangle\langle j| \otimes \rho_{ij}^T$$

$$N^A \rightarrow B$$



$$(\mathbb{R}, n, \epsilon)$$



$$\|\psi_m^{E^A} - \psi_{m'}^{E^A}\|$$

$$Pr \{ \hat{M} \neq M \} \leq \epsilon$$

$$C(N) = \sup_{\text{classical test}}$$

Thm. (HSW).

$$C(N) = \lim_{n \rightarrow \infty} \frac{1}{n} C^{(n)}$$

$$C^{(n)}(N) =$$

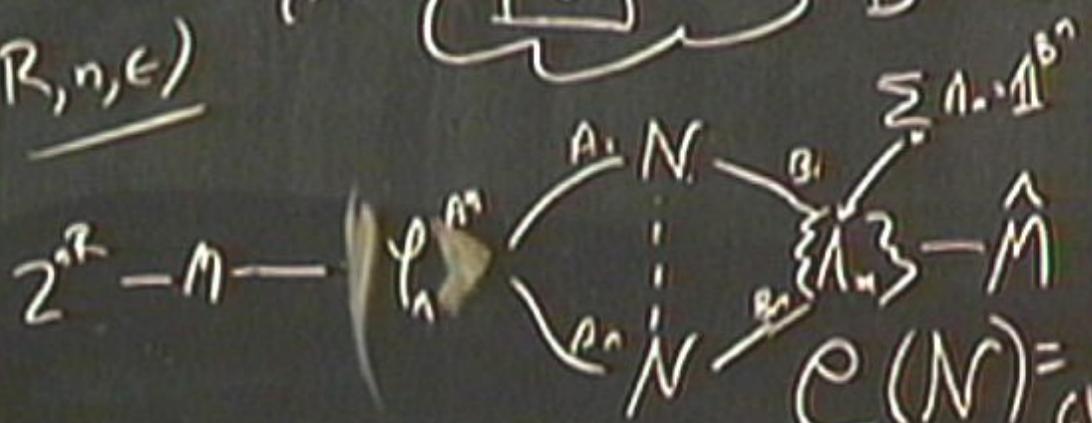
$$C^{(n)}(N) = \sup_{\text{classical test}} \dots$$

$N^A \rightarrow B$

$(R, n, \epsilon)$



$$\|\psi_m^{E^A} - \psi_{m'}^{E^A}\| \leq \epsilon$$



$C(N) = \text{sup. ach. rates, classical rep}$

$$P_r \{\hat{M} \neq M\} \leq \epsilon$$

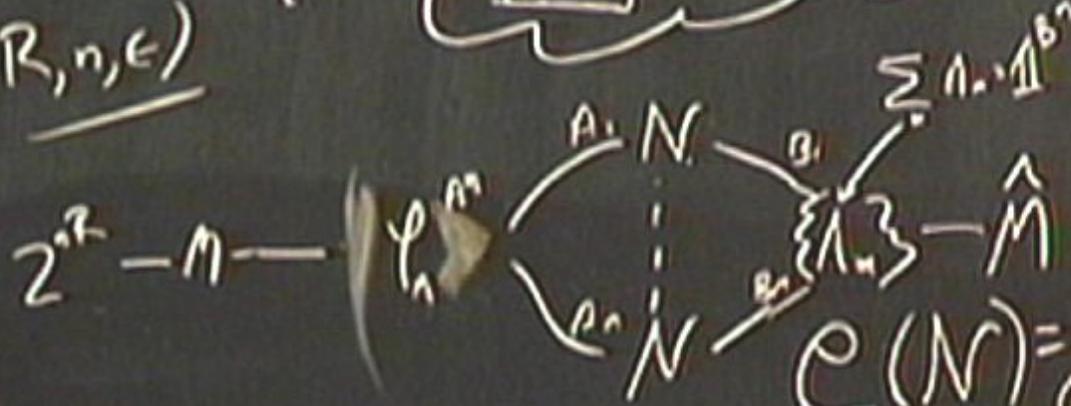
Thm (HSW).  
 $C(N) = \lim_{n \rightarrow \infty} \frac{1}{n} C(N^{\otimes n})$   
 $C^{(1)}(N) = \max_{\{P, P'\}} \{ I(X; B) - \chi(\epsilon_P, N^{\otimes n}) \}$

$$N^A \rightarrow B$$

$$(R, n, \epsilon)$$



$$\|\psi_m^{E_n} - \psi_{m'}^{E_n}\| \leq \epsilon$$



$$C(N) = \text{sup. act. rates, classical cap}$$

$$P_r \{ \hat{M} \neq M \} \leq \epsilon$$

Thm (HSW).

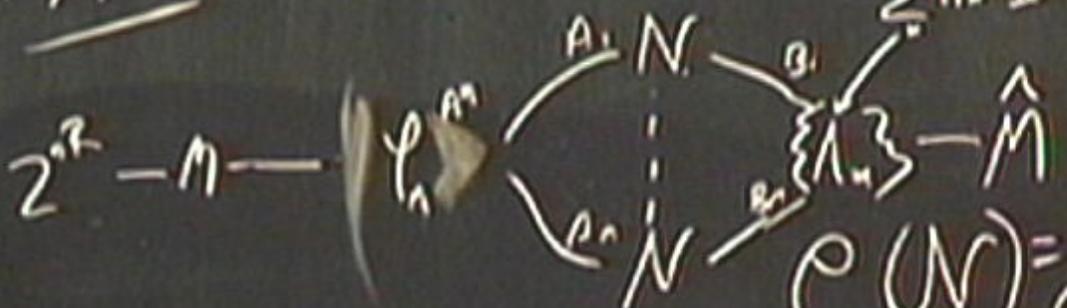
$$C(N) = \lim_{n \rightarrow \infty} \frac{1}{n} C(N^{\otimes n})$$

$$C^{(1)}(N) = \max_{\{P_A, P_B\}} I(X; B)$$

$$N^A \rightarrow B$$



$$(R, n, \epsilon)$$



$$\|\psi_m^{E^A} - \psi_{m'}^{E^A}\| \leq \epsilon$$

$$P(N) = \lim_{n \rightarrow \infty} P^{(n)}(N)$$

$$C(N) = \sup_{\text{classical ref}}$$

$$P_r \{ \hat{M} \neq M \} \leq \epsilon$$

Thm. (HSW).

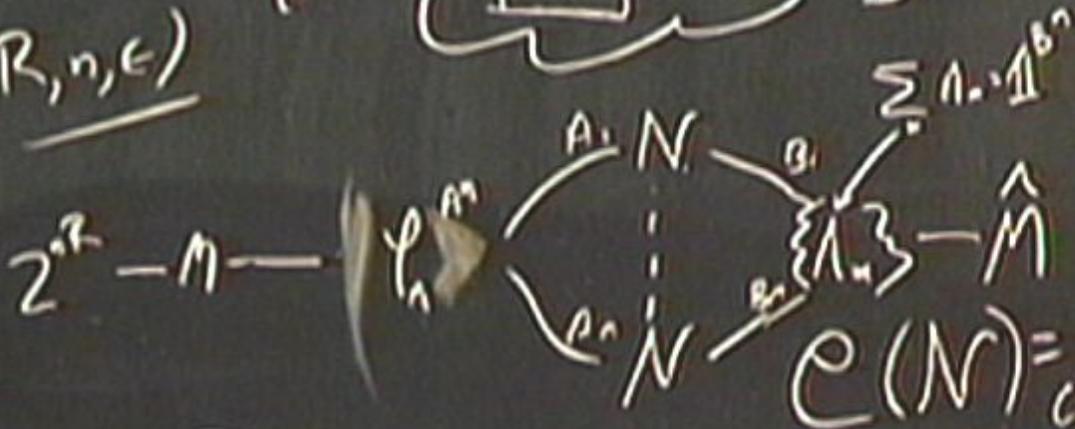
$$C(N) = \lim_{n \rightarrow \infty} \frac{1}{n} C(N^{\otimes n})$$

$$C^{(n)}(N) = \max_{\{P, P'\}} \dots$$

$$(\psi_x^A)^{\otimes n}$$

$N: A \rightarrow B$

$(R, n, \epsilon)$



$C(N) = \text{sup. act. rates}$   
classical rate

$$P_r \{ \hat{M} \neq M \} \leq \epsilon$$

Thm (HSW).

$$C(N) = \lim_{n \rightarrow \infty} \frac{1}{n} C(N^{(n)})$$

$$C^{(n)}(N) = \max_{\{P_x, P_y\}} \left[ I(X; B) - \chi(\epsilon_r, N(\varphi_x^n)) \right]$$

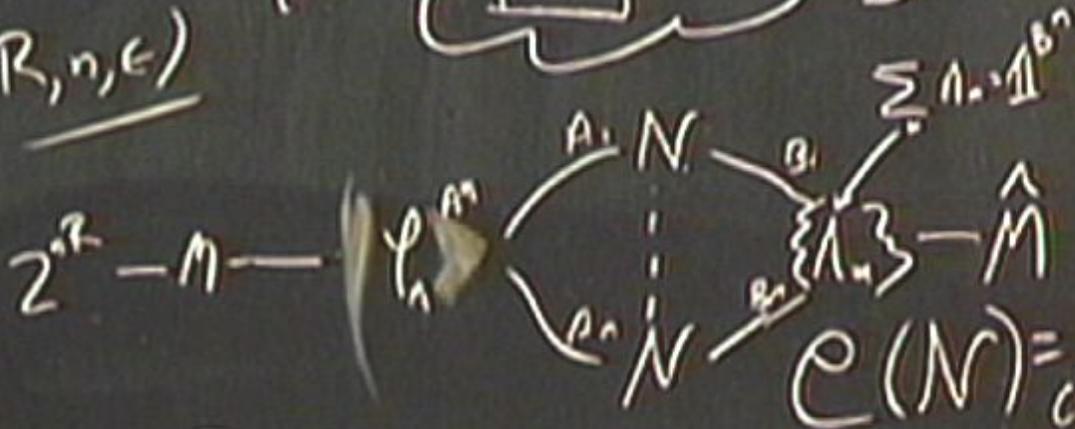
$$\| \varphi_m^{E^n} - \varphi_{m'}^{E^n} \| \leq \epsilon$$

$$P(N) = \lim_{n \rightarrow \infty} P^{(n)}(N)$$

$$\max_{\{P_x, P_y\}} [I(X; B) - I(X; E)]$$

$N: A \rightarrow B$

$(R, n, \epsilon)$



$$\|\psi_m^{E^A} - \psi_{m'}^{E^A}\| \leq \epsilon$$

$$P(N) = \lim_{n \rightarrow \infty} P^{(n)}(N)$$

$$\max_{\epsilon, P, \psi} [I(X; B) - I(X; E)]$$

$C(N) = \text{sup. act. rates}$   
classical rate

$$P_r \{ \hat{M} \neq M \} \leq \epsilon$$

Thm (HSW).

$$C(N) = \lim_{n \rightarrow \infty} \frac{1}{n} C^{(n)}(N^{\otimes n})$$

$$C^{(n)}(N) = \max_{\epsilon, P, \psi} \{ I(X; B) - \chi(\epsilon, N(\psi_x)) \}$$

$$I(X; B) - \chi(\epsilon, N(\psi_x))$$

$N: A \rightarrow B$



$(R, n, \epsilon)$



$$\|\Psi_m^{E^n} - \Psi_{m'}^{E^n}\| \leq \epsilon$$


---


$$P(N) = \lim_{n \rightarrow \infty} \frac{P^{(n)}(N^{(n)})}{n}$$


---


$$\max_{\{P_x, P_y\}} [I(X; B) - I(X; E)]$$

$C(N) = \text{sup. ach. rates classical cap}$

$$P_r \{ \hat{M} \neq M \} \leq \epsilon$$

Thm (HSW).

$$C(N) = \lim_{n \rightarrow \infty} \frac{1}{n} C(N^{(n)})$$

$$C^{(n)}(N) = \max_{\{P_x, P_y\}} [I(X; B) - \chi(\epsilon, N(\Psi_x^n))] - \chi(\epsilon, N(\Psi_x^n))$$

Thm:  $P_{\frac{1}{3}H^3}$  private Hor. channels exist

$$\exists N \quad Q(N) = 0, \quad P^{(1)}(N) > 0$$

Thm.  $P_{\frac{1}{3}H^3}$  provide Hor. channels exist

$$\exists N \quad Q(N) = 0, \quad \underline{P^{(1)}(N) > 0}$$

Thm.  $\mathcal{P}_{\frac{1}{3}}^H$  private Hor. channels exist

$$\exists N \ Q(N) = 0, \ \underline{\mathcal{P}^{(1)}(N)} > 0$$

Thm.  $\mathcal{P}_{\frac{1}{3}}^H$  :  $\exists N$  in 2 qubits st.  $\uparrow$

let unit ops.  $\left\{ |0\rangle\langle 0| \otimes \frac{11_2}{2}, |1\rangle\langle 1| \otimes \frac{11_2}{2} \right\}, \ I(X;B) - I(X;E) > .02$

Thm.  $\mathcal{Q}_{\frac{1}{3}}^H$  private Hor. channels exist

$$\exists N \ Q(N) = 0, \ P^{(1)}(N) > 0$$

Thm.  $\mathcal{P}_{\frac{1}{3}}^H$ :  $\exists N$  in 2 qubits st.  $\uparrow$

with unif. enst.  $\left\{ |0\rangle\langle 0| \otimes \frac{1}{2}, |1\rangle\langle 1| \otimes \frac{1}{2} \right\}, I(X;B) - I(X;E) > .02$

Thm.  $Q_{\frac{1}{3}}^H$  private Hor. channels exist

$$\exists N \quad Q(N) = 0, \quad \underline{P^{(1)}(N) > 0}$$

Thm.  $P_{\frac{1}{3}}^H$  :  $\exists N$  in 2 qubits st.  $\uparrow$

with unit ops.  $\left\{ |10\rangle\langle 01| \otimes \frac{11_2}{2}, |11\rangle\langle 11| \otimes \frac{11_2}{2} \right\}$ ,  $I(X;E) > .02$

Thm SY'08

given  $\{P_x, P_x^A\}$ ,  $N^A \xrightarrow{E} B$ , let  $\epsilon$

Thm.  $Q_{\frac{1}{3}}^H$  private Hor. channels exist

$$\exists N \text{ s.t. } Q(N) = 0, \underline{P^{(1)}(N)} > 0$$

Thm.  $P_{\frac{1}{3}}^H$ :  $\exists N$  in 2 qubits s.t.  $\uparrow$

with unif. ensembles  $\left\{ |10\rangle\langle 01| \otimes \frac{\mathbb{1}_2}{2}, |11\rangle\langle 11| \otimes \frac{\mathbb{1}_2}{2} \right\}$ ,  $I(X;B) - I(X;E) > .02$

Thm SY'08

given  $\{P_x, P_x^A\}$ ,  $N^A \xrightarrow{E} B$ , let  $\mathcal{E}$  be a 50% erasure channel

Thm.  $Q_{\frac{1}{3}}^H$  private Hor. channels exist

$$\exists N \ Q(N)=0, \ P^{(1)}(N) > 0$$

Thm.  $P_{\frac{1}{3}}^H$ :  $\exists N$  in 2 qubits st.  $\uparrow$

with unit ops.  $\left\{ |0\rangle\langle 0| \otimes \frac{I_2}{2}, |1\rangle\langle 1| \otimes \frac{I_2}{2} \right\}, I(x;B) - I(x;E) > .02$

Thm SY'08

given  $\{P_x, \rho_x^A\}$ ,  $N^A \xrightarrow{E} B$ , let  $E$  be a 50% erasure channel

with input dim  $\sim \sum_{x \in \mathcal{X}} \text{rank } \rho_x^A$

Thm. QH<sup>3</sup> private Hor. channels exist

$$\exists N \ Q(N) = 0, \quad \underline{Q^{(1)}(N) > 0}$$

Thm. PH<sup>3</sup>:  $\exists N$  on 2 qubits st.  $\uparrow$

Let unit ops.  $\left\{ |0\rangle\langle 0| \otimes \frac{I_2}{2}, |1\rangle\langle 1| \otimes \frac{I_2}{2} \right\}, \quad I(X;B) - I(X;E) > .02$

Thm SY'08

Given  $\{ \rho_x, \rho^A \}$ ,  $N^A \rightarrow B$ , let  $\mathcal{E}$  be a 50% erasure channel

with input dim  $\sum \text{rank } \rho_x^A$ . Then  $Q^{(1)}(N \circ \mathcal{E}) \geq \frac{1}{2} [I(X;B) - I(X;E)]$

Thm.  $Q_{\frac{1}{3}}^H$  private Hor. channels exist

$$\exists N \ Q(N) = 0, \quad \underline{Q^{(1)}(N) > 0}$$

Thm.  $P_{\frac{1}{3}}^H$ :  $\exists N$  in 2 qubits st.  $\uparrow$

with unit ops.  $\left\{ |10\rangle\langle 01| \otimes \frac{I_2}{2}, |11\rangle\langle 11| \otimes \frac{I_2}{2} \right\}, \quad I(X;B) - I(X;E) > .02$

Thm SY'08

given  $\{ \rho_A^x \}$ ,  $N^A \xrightarrow{E} B$ , let  $\mathcal{E}$  be a 50% erasure channel with input dim  $\rightarrow \sum \text{rank } \rho_A^x$ . Then  $Q^{(1)}(N \circ \mathcal{E}) \geq \frac{1}{2} [I(X;B) - I(X;E)]$

Thm. QH<sup>3</sup> private Her. channels exist

$$\exists N \ Q(N) = 0, \quad \underline{Q''(N) > 0}$$

Thm. PH<sup>3</sup>:  $\exists N$  on 2 qubits st.  $\uparrow$

with unit ops.  $\left\{ |0\rangle\langle 0| \otimes \frac{I_2}{2}, |1\rangle\langle 1| \otimes \frac{I_2}{2} \right\}, \quad I(X;B) - I(X;E) > .02$

Thm SY'08

given  $\{ \rho_x, \rho_A \}$ ,  $N^A \rightarrow B$ , let  $\mathcal{E}$  be a 50% erasure channel

with input dim  $\rightarrow \sum \text{rank } \rho_A$ . Then  $Q''(N \circ \mathcal{E}) \geq \frac{1}{2} [I(X;B) - I(X;E)]$





Thm:  $\mathbb{Q}_{\frac{1}{3}}^{\text{H}^3}$  private rec. channels exist

$$\exists N \ Q(N) = 0, \ \underline{P^{(1)}(N) > 0}$$

Thm:  $\mathbb{P}^{\text{H}^3}$ :  $\exists N$  on 2 qubits st.  $\uparrow$

$$\left\{ \begin{array}{l} 10 \times 01 \otimes \frac{11_2}{2}, \ 11 \times 01 \otimes \frac{11_2}{2} \end{array} \right\}, \ \underline{I(X;B) - I(X;E) > .02}$$

Y'08

in  $\{P_A, P_A^A\}$ ,  $N^A \xrightarrow{E}$ , let  $\mathcal{E}$  be a 50% private channel

Then  $Q^{(1)}(N \otimes \mathcal{E}) \geq \frac{1}{2} [I(X;B) - I(X;E)]$

Thm.  $Q_{\frac{1}{3}}^{\text{H}^3}$  private Hor. channels exist

$$\exists N \ Q(N) = 0, \quad \underline{Q^{(1)}(N) > 0}$$

Thm.  $P_{\frac{1}{3}}^{\text{H}^3}$ :  $\exists N$  in 2 qubits st.  $\uparrow$

with unit ops.  $\{ |10\rangle\langle 01| \otimes \frac{11_2}{2}, |11\rangle\langle 11| \otimes \frac{11_2}{2} \}$ ,  $\underline{I(X;B) - I(X;E) > .02}$

Thm SY'08

given  $\{ \rho_x, \rho^A \}$ ,  $N^A \xrightarrow{E} B$ , let  $E$  be a 50% private channel with input dim  $\leq \sum \text{rank } \rho^A$ . Then  $Q^{(1)}(N \otimes E) \geq \frac{1}{2} [I(X;B) - I(X;E)]$

Thm.  $Q_{\frac{1}{3}}^H$  private Hor. channels exist

$$\exists N \ Q(N) = 0, \quad \underline{P^{(1)}(N) > 0}$$

Thm.  $P_{\frac{1}{3}}^H$ :  $\exists N$  on 2 qubits st.  $\uparrow$

with unit ops.  $\left\{ \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 0 \\ 0 & i \end{pmatrix}, \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 0 & -1 \end{pmatrix} \right\}$ ,  $I(X;B) - I(X;E) > .02$

Thm SY'08

given  $\{P_x, P_A\}$ ,  $N^A \rightarrow B$ , let  $\mathcal{E}$  be a 50% erasure channel with input dim  $\rightarrow \sum \text{rank } P_A$ . Then  $Q^{(1)}(N \circ \mathcal{E}) \geq \frac{1}{2} [I(X;B) - I(X;E)]$

$$M \begin{cases} Q''(M) = 0 \\ Q'''(M) = 0 \end{cases}$$

$$M' \begin{cases} Q''(M) = 0 \end{cases}$$

$$Q''(M_0 + \epsilon) > Q''(M_0 - \epsilon)$$

$$\begin{array}{c}
 M_1 \quad M_{2k+1} \\
 \uparrow \\
 \underline{Q''(M) = 0} \\
 Q(M) \neq Q''(M \oplus M) > Q''(N \oplus \epsilon) > \underline{.01}
 \end{array}$$







$$M_k \quad 2k+1 \quad \dots$$

$$Q''(M) = 0 \quad \begin{matrix} \text{not} \\ \varepsilon^k \end{matrix}$$

$$Q(M) \neq Q''(M \circ M) > Q''(N \circ \varepsilon) > \textcircled{.01}$$

$$Q(M_k) > .01k \quad Q''(M_k) = 0$$