

Title: Additive Extensions of a Quantum Channel

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Abstract:

# Additive extensions of a quantum channel

Graeme Smith and John A Smolin

arXiv:0712.2471

# Additive Extensions of a Quantum Channel

- Quantum Channels
- Capacities and Additivity
- Known Bounds
- Additive and Degradable Extensions
- New Bounds on quantum and private capacity
- Relation to Symmetric Side Channels (SSW 2006)

# Shannon 1948

Mutual Information

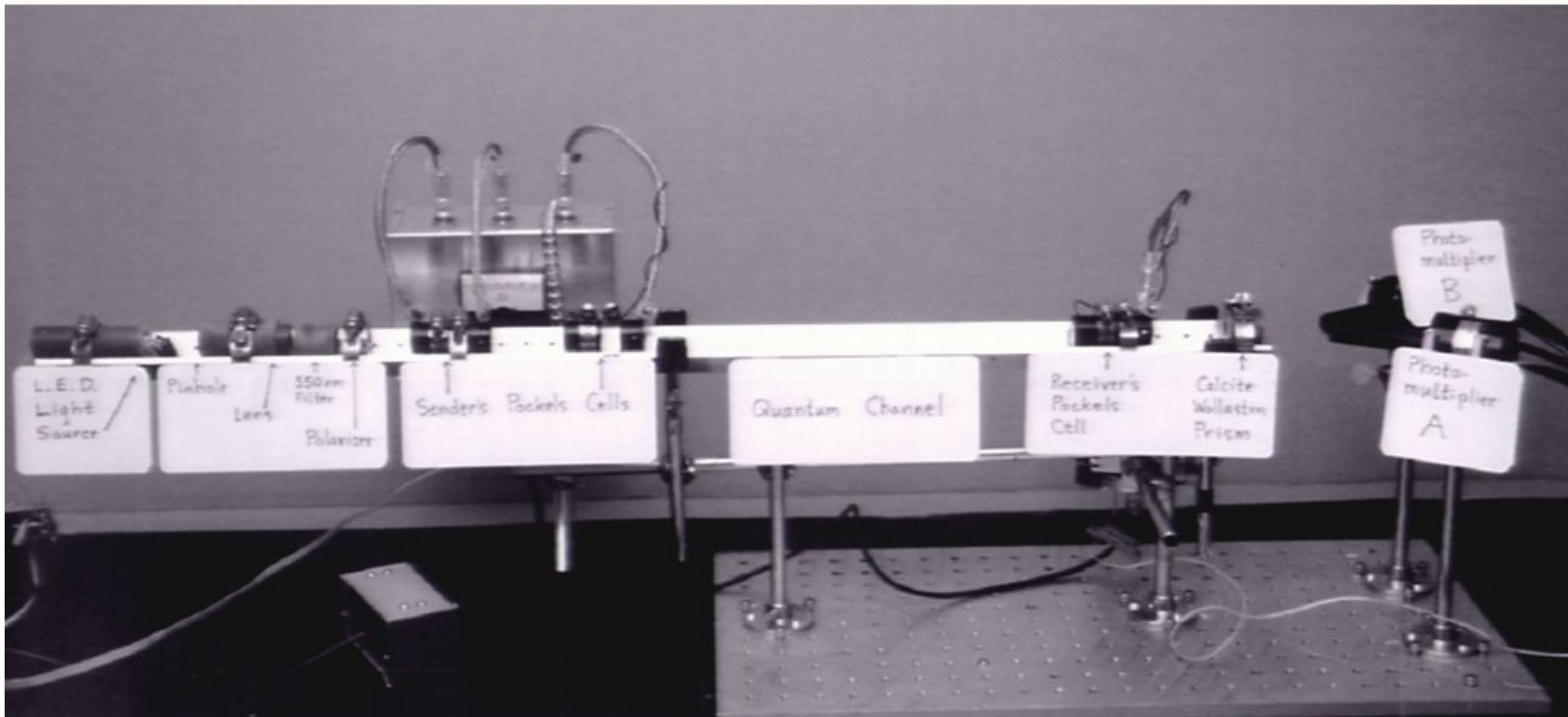
$$I(A:B) = H(A) + H(B) - H(A,B)$$

Channel maps  $X$  to  $B$

Classical capacity  $C = \text{Max}_X I(A:B)$

Remarkable that you can achieve this with random coding  
And that it is a single-letter formula

# A Quantum Channel



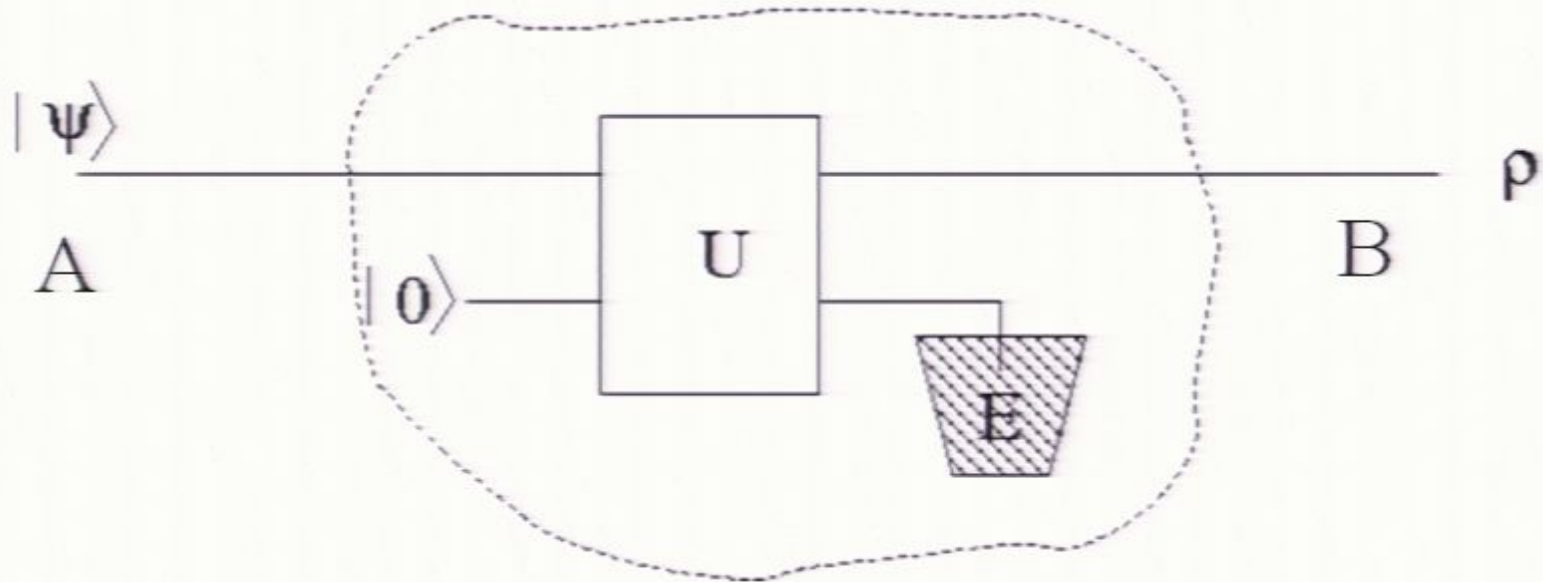
# Quantum channels

A quantum channel is represented by a trace-preserving, completely positive (TCP) map from density operators to density operators.

$$\rho \rightarrow \sum_i A_i \rho A_i^\dagger$$
$$\sum_i A_i^\dagger A_i = 1$$

# Isometry

$$A \dashrightarrow BE$$



# Coherent Information

Coherent Information :  $I^c(\rho_{AB}) = S(\rho_B) - S(\rho_{AB})$   
of a channel and state)  $I^c(\mathcal{N}, \rho) = I^c(I \otimes \mathcal{N}(\phi^{AB}))$ ,

where  $\text{Tr}_A |\phi^{AB}\rangle\langle\phi^{AB}| = \rho$

$$Q_1 = \max_{\rho} I^c(\mathcal{N}, \rho)$$



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Not the capacity

SS96

# Capacity Theorem

Have to regularize the one-shot capacity

$$\begin{aligned} \text{Quantum Capacity : } Q &= \lim_{n \rightarrow \infty} \frac{1}{n} \max_{\rho_n} I^c(\mathcal{N}^{\otimes n}, \phi_n) \\ &= \lim_{n \rightarrow \infty} \frac{1}{n} Q_1(\mathcal{N}^{\otimes n}) \\ Q_1 &= \max_{\rho} I^c(\mathcal{N}, \rho) \end{aligned}$$

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- \* Ugly
- \* Usually impossible to evaluate

# Additivity Questions Abound

Entanglement of formation  
Holevo classical capacity  
Minimum output entropy  
Output rank

Quantum capacity is simply not single-letter

Bad: Hard to calculate

Good: Lots to work on

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# Depolarizing channel

$$\mathcal{N}_p(\rho) = (1 - p)\rho + \frac{p}{3}X\rho X + \frac{p}{3}Y\rho Y + \frac{p}{3}Z\rho Z$$

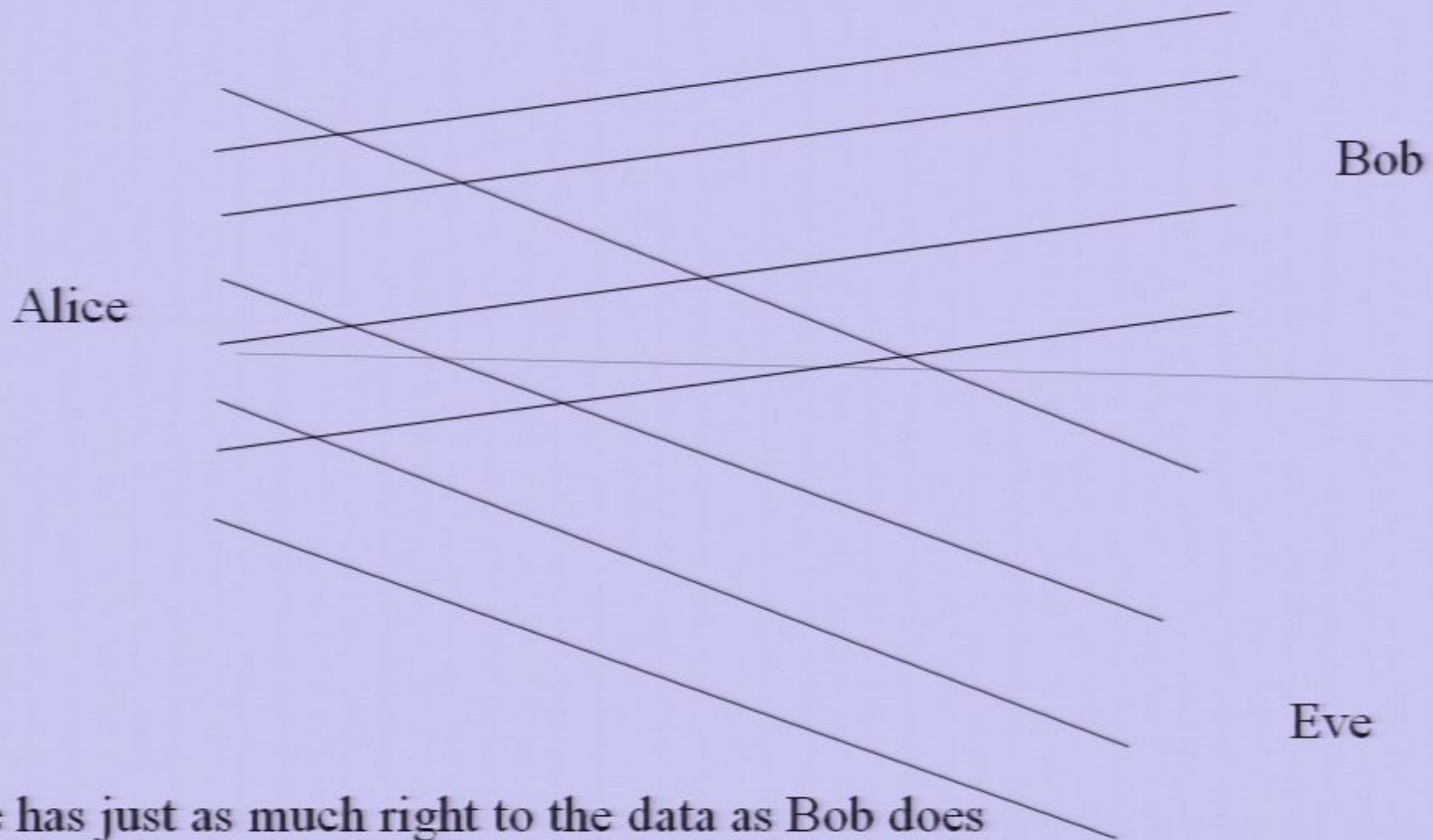
$$\mathcal{N}_x(\rho) = (1 - x)\rho + x\frac{I}{2} = (1 - x)\rho + \frac{x}{4}I\rho I + \frac{x}{4}X\rho X + \frac{x}{4}Y\rho Y + \frac{x}{4}Z\rho Z$$

$$p = \frac{3}{4}x$$

$$\mathcal{N}_p(U\rho U^\dagger) = U\mathcal{N}_p(\rho)U^\dagger$$



# 50% depolarizing probability



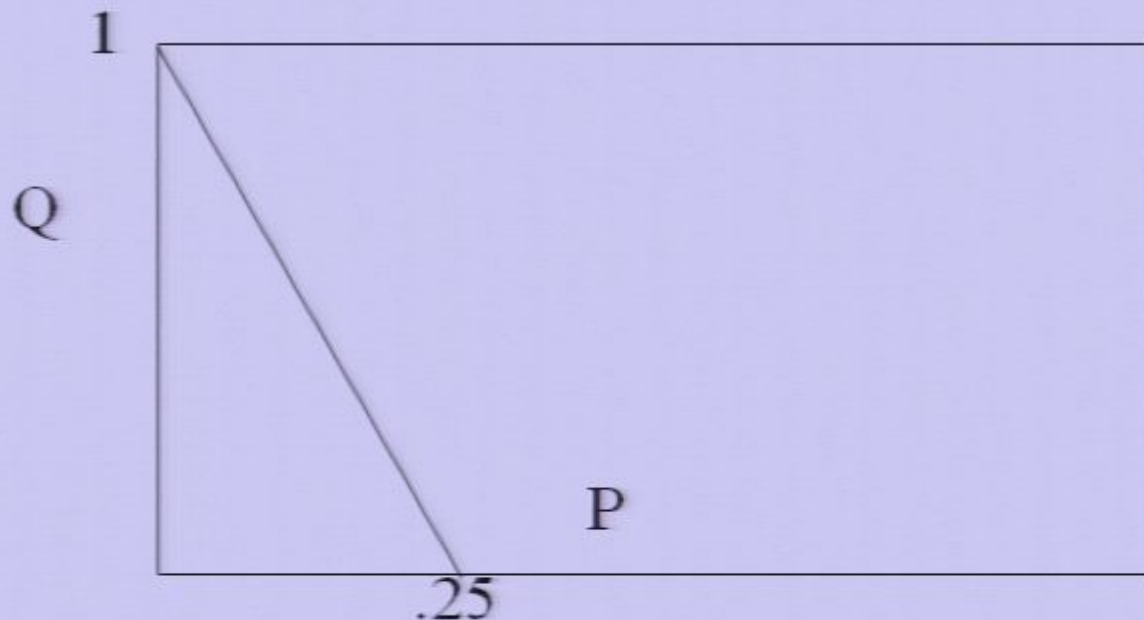
Eve has just as much right to the data as Bob does  
But quantum information cannot be cloned  $\langle \text{Dolly} \rangle \langle / \text{Dolly} \rangle$

# No-cloning bound

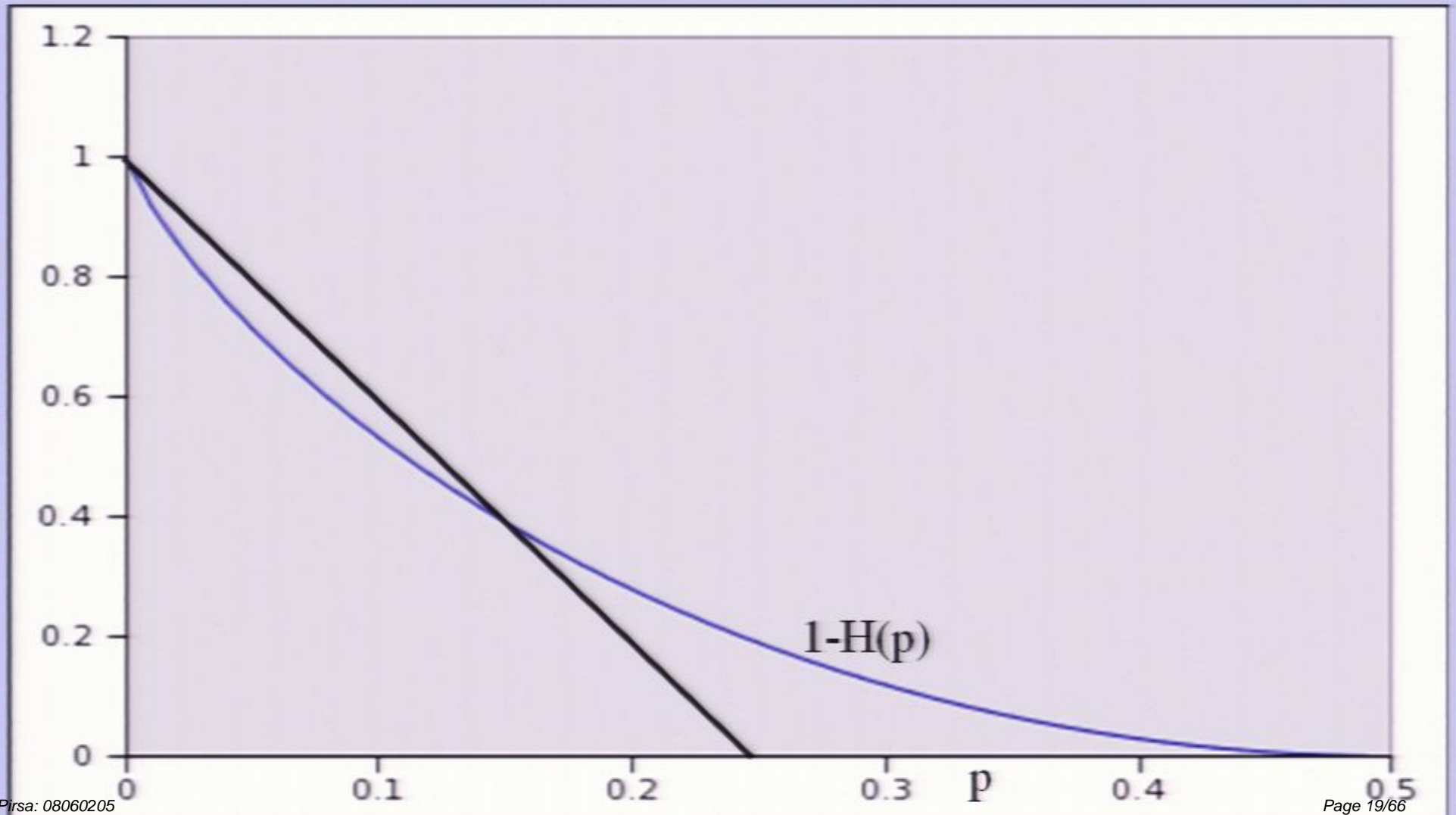
50% depolarizing probability corresponds to  $p=3/8$

Surprisingly, the approximate cloner of the previous slide is not optimal. Need ancilla to find optimal approximate cloning channel

Can achieve  $p=1/4$  with optimal approximate cloner

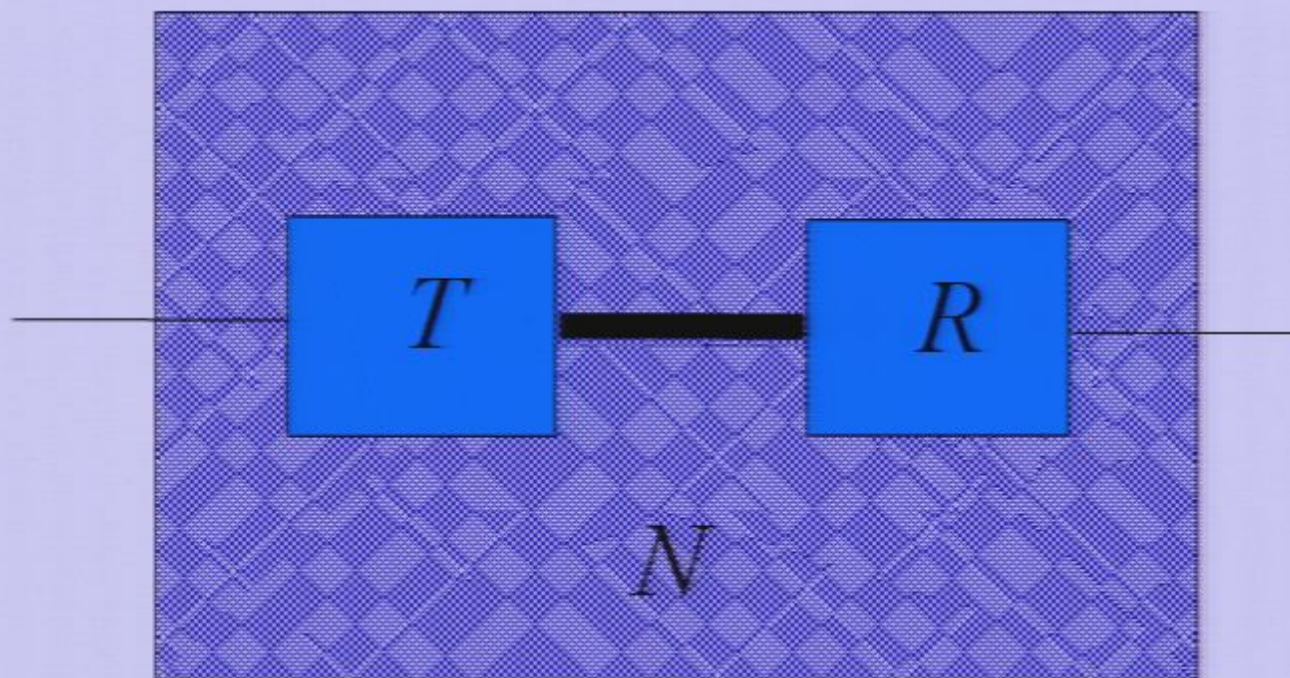


# Rains bound



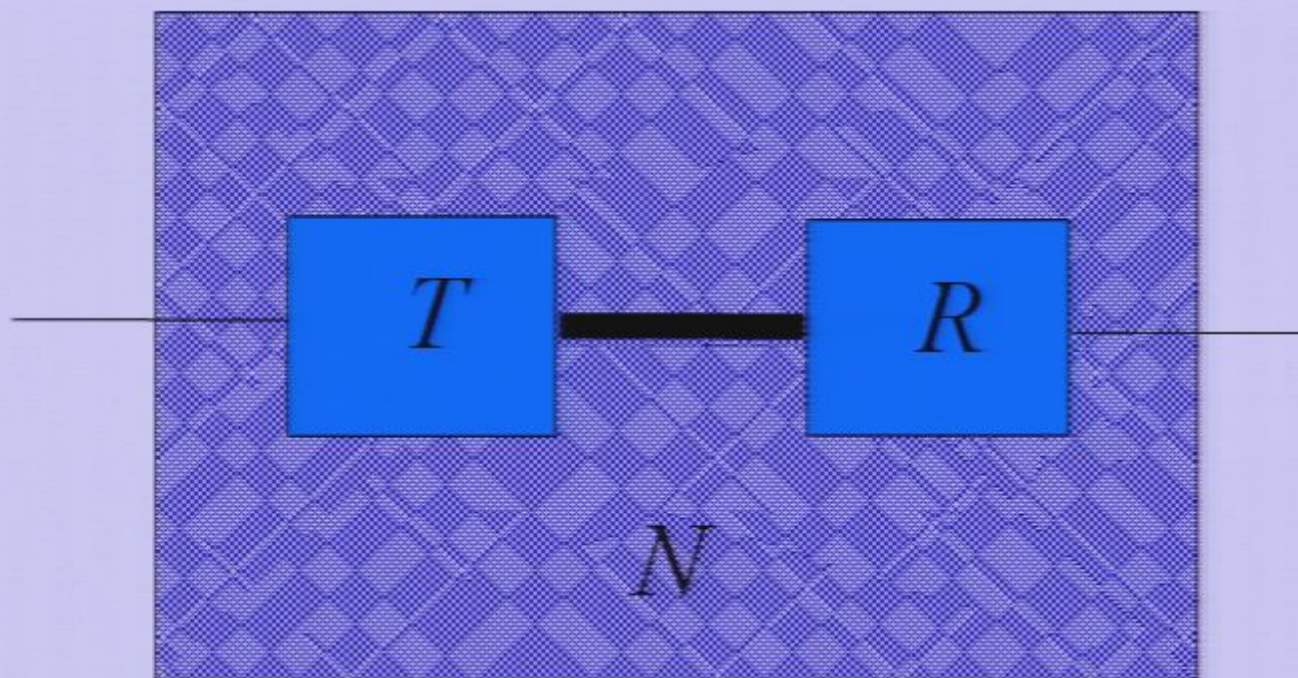
# Additive Extension

Definition:  $\mathcal{T}$  is an additive extension of a quantum channel  $\mathcal{N}$  if there is another channel  $\mathcal{R}$  such that  $\mathcal{N} = \mathcal{R} \circ \mathcal{T}$  and  $Q(\mathcal{T}) = Q_1(\mathcal{T})$ .



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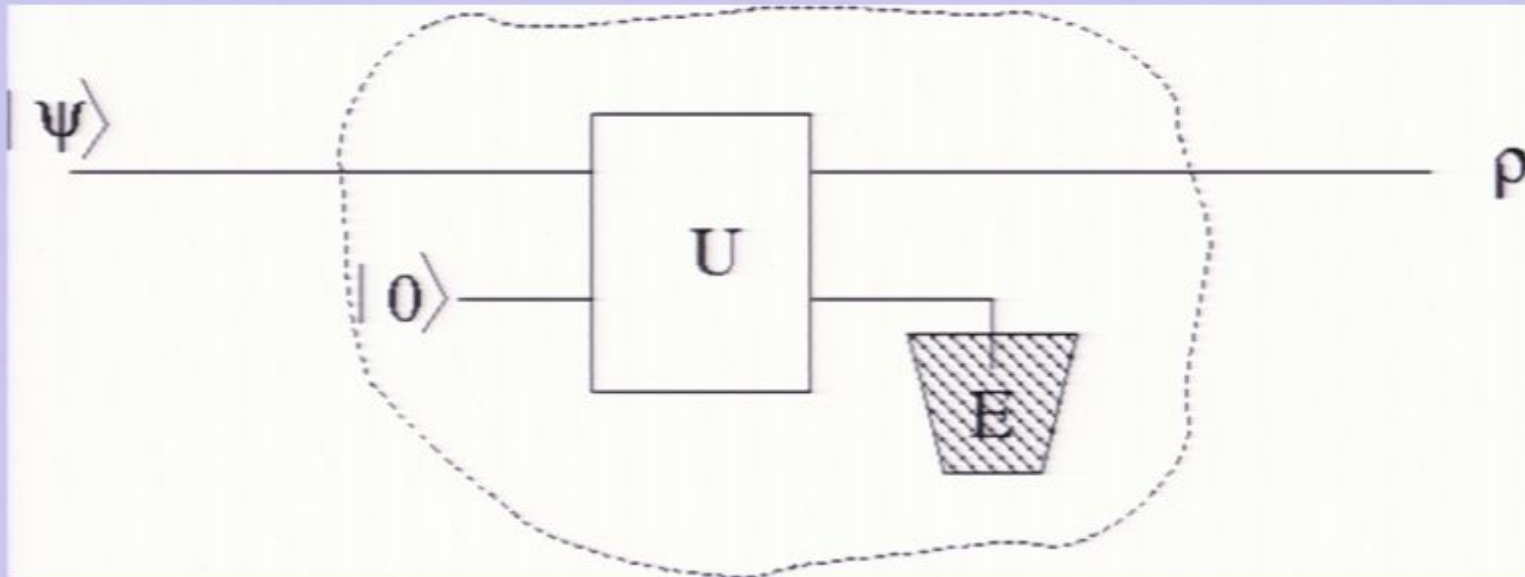
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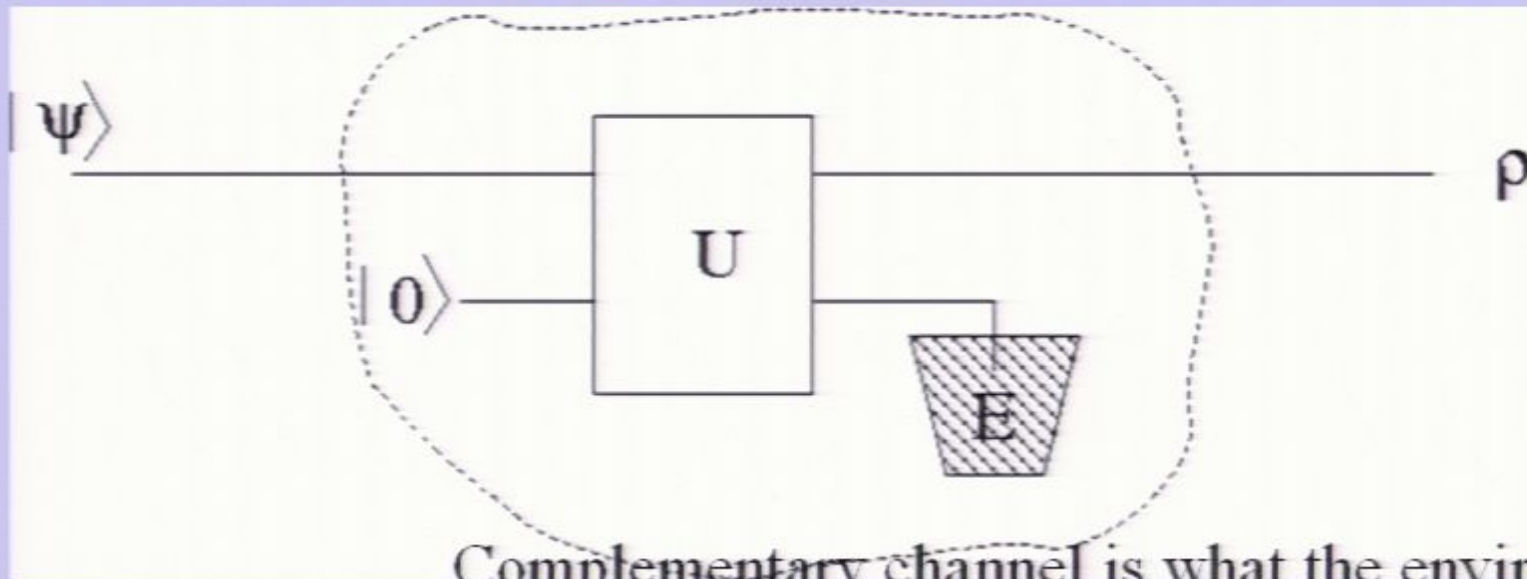
A degradable extension is an additive extension that is also degradable

definition: A channel  $\mathcal{N}$  with isometric extension  $U : A \rightarrow BE$  is called *degradable* if there is a degrading map  $\mathcal{D}$  such that  $\mathcal{D} \circ \mathcal{N} = \hat{\mathcal{N}}$ , where  $\hat{\mathcal{N}}(\rho) = \text{Tr}_B U \rho U^\dagger$ .  $\hat{\mathcal{N}}$  is called the *complementary channel* of  $\mathcal{N}$ .

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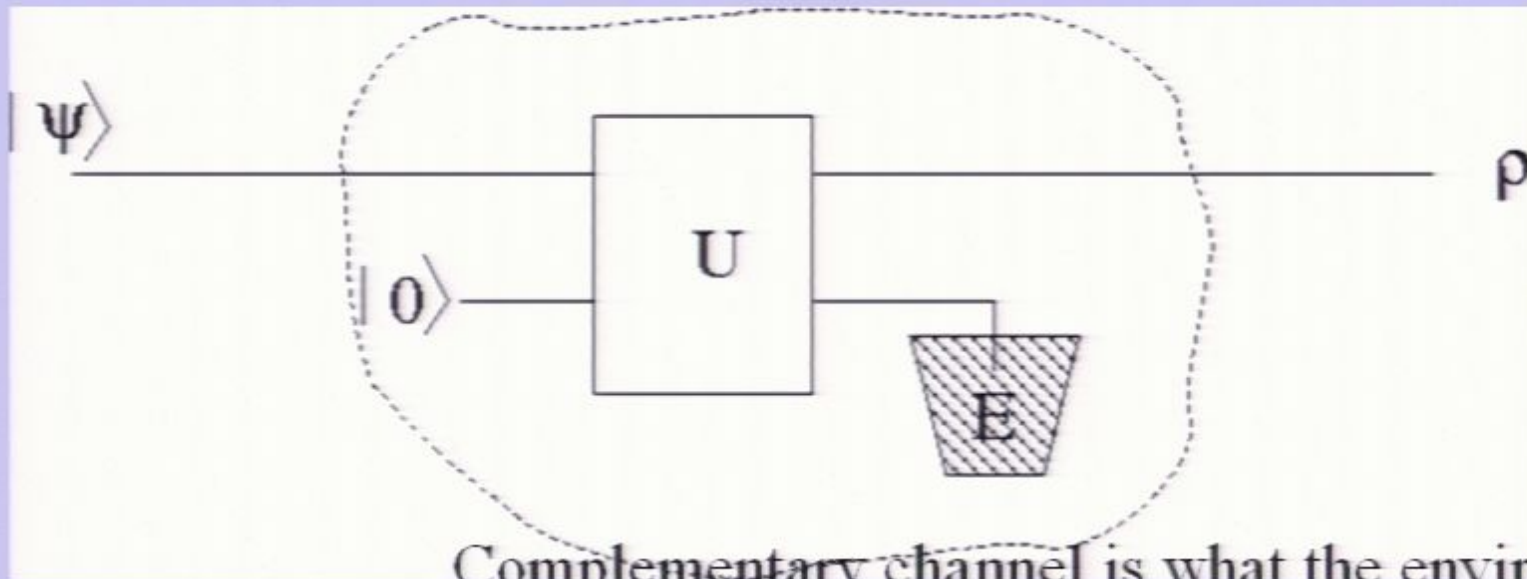
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Complementary channel is what the environment gets



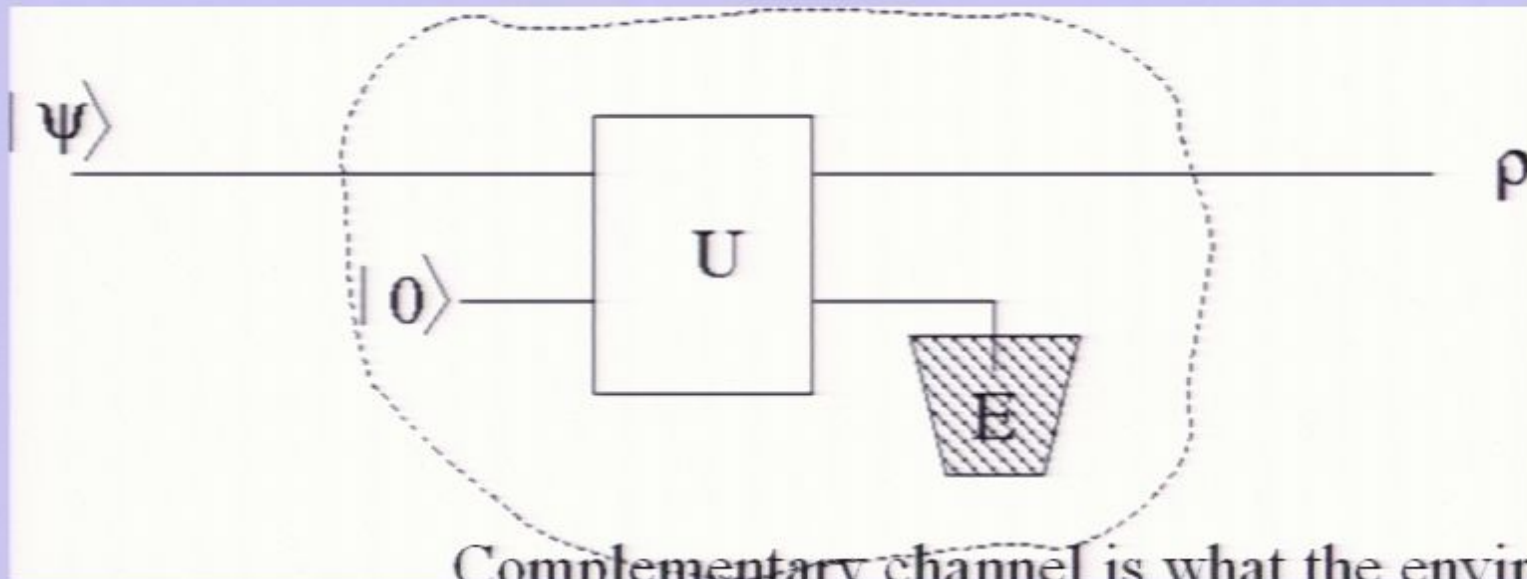
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Degradable channels have single-letter capacity Devetak-Shor 2003  
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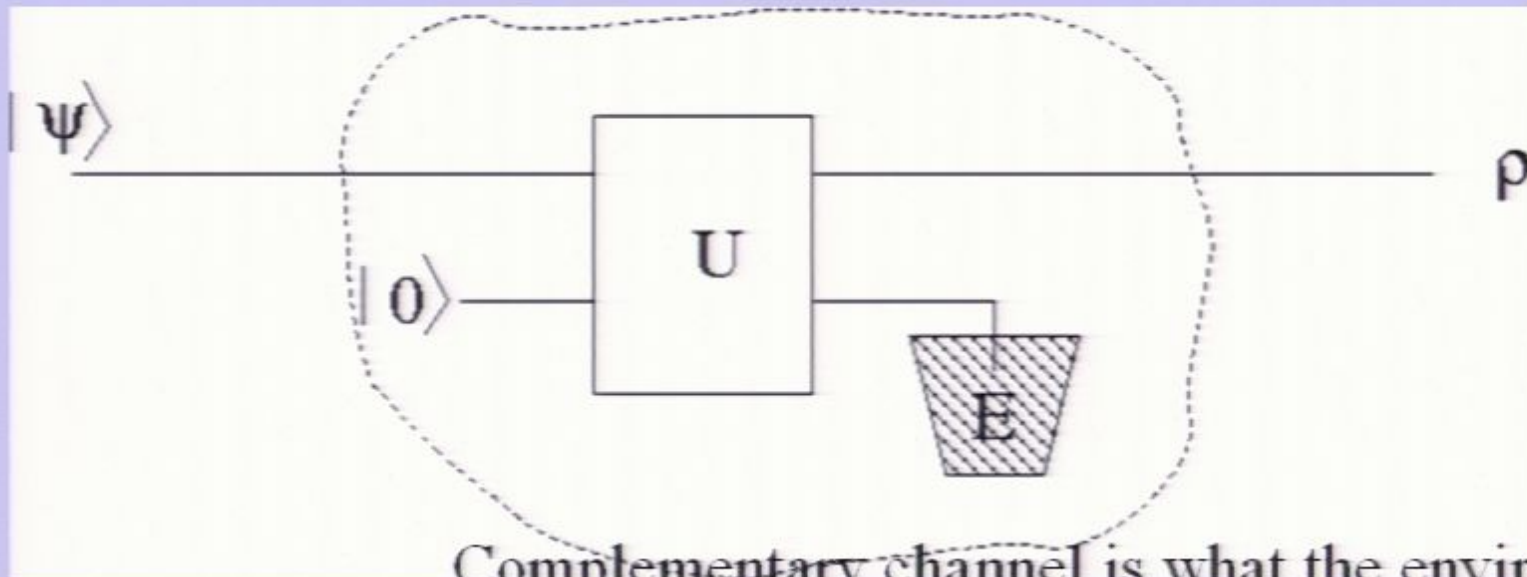
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Unfortunately, the depolarizing channel is not degradable

# Good Properties of Additive Extensions

- Simple to understand
- Additive (single-letter) bounds on  $Q$
- Includes all previous bounds on  $Q$
- Can get Rains bound without my having to understand his paper
- Also bounds private capacity when degradable
- Good convexity properties
- Fixes argument for no-cloning straight line
- Yields new, tighter bounds

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# Embarrassingly Simple Theorem

$Q(\mathcal{N}) \leq Q^{(1)}(\mathcal{T})$   $\mathcal{T}$  is an additive extension of  $\mathcal{N}$

$Q_p(\mathcal{N}) \leq Q^{(1)}(\mathcal{T})$  when  $\mathcal{T}$  is also degradable

Proof:

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$Q(\mathcal{N}) \leq Q^{(1)}(\mathcal{T})$   $\mathcal{T}$  is an additive extension of  $\mathcal{N}$

$C_p(\mathcal{N}) \leq Q^{(1)}(\mathcal{T})$  when  $\mathcal{T}$  is also degradable

**Proof:**

$Q^{(1)}(\mathcal{T}) = Q(\mathcal{T})$  by definition of additive extension

$Q(\mathcal{N}) \leq Q(\mathcal{T})$  can obtain  $\mathcal{N}$  from  $\mathcal{T}$  using  $\mathcal{R}$

similarly for  $C_p$  using  $C_p(\mathcal{T}) = Q^{(1)}(\mathcal{T})$  (Smith07)

# Flagged-Degradable Extensions

Lemma:

Suppose we have  $\mathcal{N} = \sum_i p_i \mathcal{N}_i$  with  $\mathcal{N}_i$  degradable  
 $\mathcal{T} = \sum_i p_i \mathcal{N}_i \otimes |i\rangle\langle i|$  is a degradable extension of  $\mathcal{N}$

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$\mathcal{N} = \mathcal{R} \circ \mathcal{T}$  : trace out flag system

$\mathcal{T}$  Degradable : read flag and apply  $\mathcal{D}_i$

# Convex combination bound

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$$\begin{aligned} Q_1(\mathcal{T}) &= S(\sum_i p_i \mathcal{N}_i(\phi) \otimes |i\rangle\langle i|) - S(\sum_i p_i \hat{\mathcal{N}}_i(\phi) \otimes |i\rangle\langle i|) \\ &= \sum_i p_i (S(\mathcal{N}_i(\phi)) - S(\hat{\mathcal{N}}_i(\phi))) \\ &\leq \sum_i p_i Q_1(\mathcal{N}_i) \end{aligned}$$

also  $Q(\mathcal{N}) < Q_1(\mathcal{T})$  QED

# No cloning as special case

To show off how general Additive extensions are

To let us use convexity results on no-cloning bounds

$\mathcal{N}$  Antidegradable :  $\mathcal{D} \circ \hat{\mathcal{N}} = \mathcal{N}$

Isometry of  $\mathcal{N}$  is  $U : A \rightarrow BE$

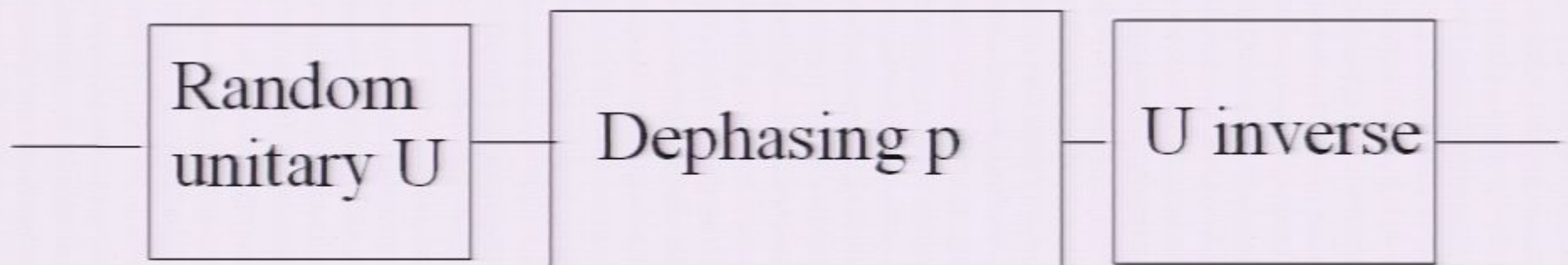
let  $V|\phi\rangle = \frac{1}{\sqrt{2}}U|\phi\rangle|0\rangle_{f_1}|1\rangle_{f_2} + \frac{1}{\sqrt{2}}\text{SWAP}_{BE}U|\phi\rangle|1\rangle_{f_1}|0\rangle_{f_2}$

$\mathcal{T}(\rho) = \text{Tr}_{E f_2} V \rho V^\dagger$  is a degradable extension of  $\mathcal{N}$

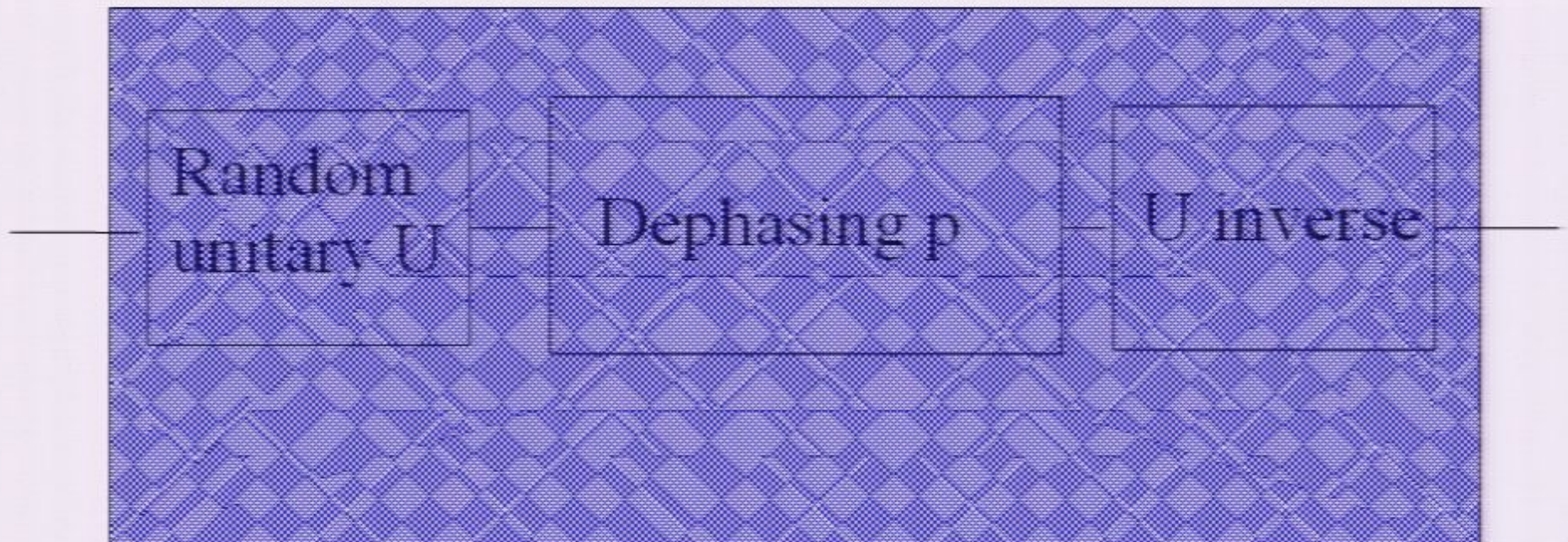
Degradable, antidegradable and zero capacity by symmetry

Can get  $N$  by applying antidegrading map when necessary

# Example: Rains bound

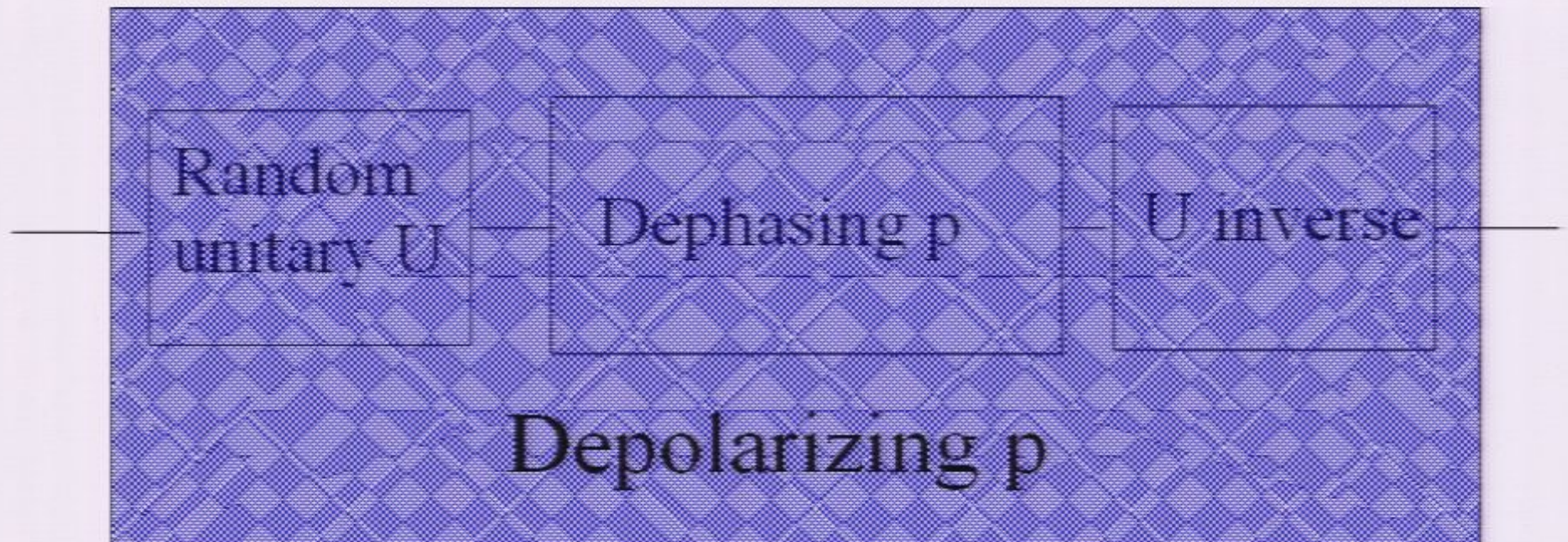


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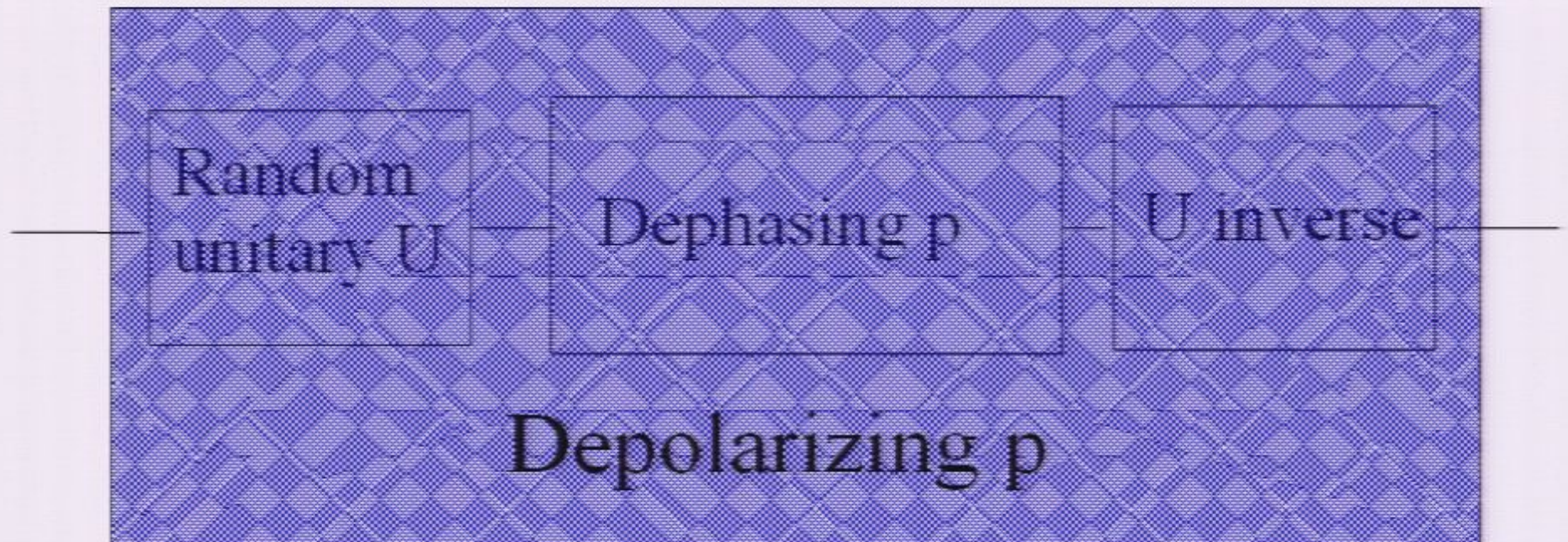




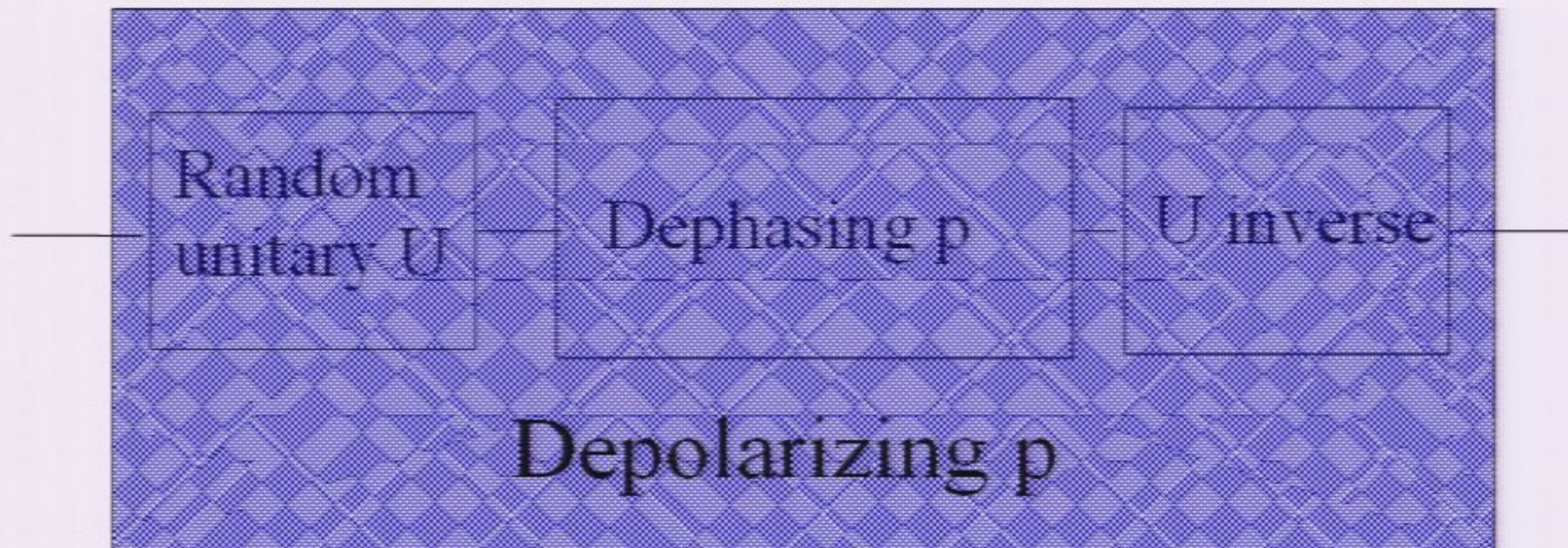
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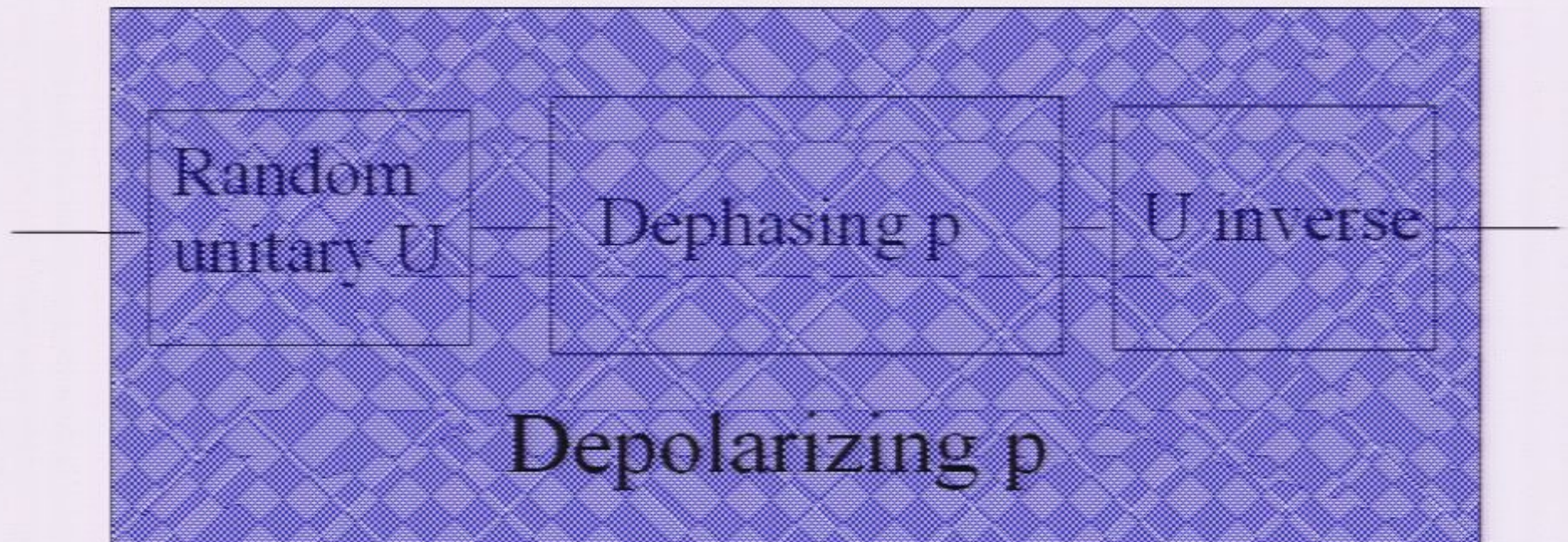


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also  $Q(\mathcal{N}) \leq Q_1(\mathcal{T})$  **QED**

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# General Degradable Qubit channel

$$A_+ = \begin{pmatrix} \cos(\frac{1}{2}(v - u)) & 0 \\ 0 & \cos(\frac{1}{2}(v + u)) \end{pmatrix}$$
$$A_- = \begin{pmatrix} 0 & \sin(\frac{1}{2}(v + u)) \\ \sin(\frac{1}{2}(v - u)) & 0 \end{pmatrix}$$

degradable when  $|\sin v| \leq |\cos u|$

Wolf Perez-Garcia 2007  
Cubitt, Ruskai, Smith 2008



# General degradable-depolarizing mixture

$$\mathcal{T}_{(u,v)}^{\text{dep}}(\rho) = \frac{1}{|\mathcal{C}|} \sum_{c \in \mathcal{C}} c^\dagger \mathcal{N}_{(u,v)}(c\rho c^\dagger) c \otimes |c\rangle\langle c|$$

$$\text{Tr}_f \mathcal{T}_{(u,v)}^{\text{dep}}(\rho) = \mathcal{N}_p(\rho)$$

$$p = 1 - \cos^2(u/2) \cos^2(v/2)$$

## Finally, some results

$$D(\mathcal{N}_p) \leq \min H \left[ \frac{1}{2}[1 + \sin u \sin v] \right] - H \left[ \frac{1}{2}[1 + \cos u \cos v] \right]$$

minimize over  $(u, v)$  s.t.  $\cos^2(u/2) \cos^2(v/2) = 1 - p$

$$D(\mathcal{N}_p) \leq \text{co} \left[ 1 - H(p), H\left(\frac{1-\gamma(p)}{2}\right) - H\left(\frac{\gamma(p)}{2}\right), 1 - 4p \right]$$

where  $\gamma(p) = 4\sqrt{1-p}(1 - \sqrt{1-p})$

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# Finally, some results

$$C(\mathcal{N}_p) \leq \min H \left[ \frac{1}{2} [1 + \sin u \sin v] \right] - H \left[ \frac{1}{2} [1 + \cos u \cos v] \right]$$

minimize over  $(u, v)$  s.t.  $\cos^2(u/2) \cos^2(v/2) = 1 - p$

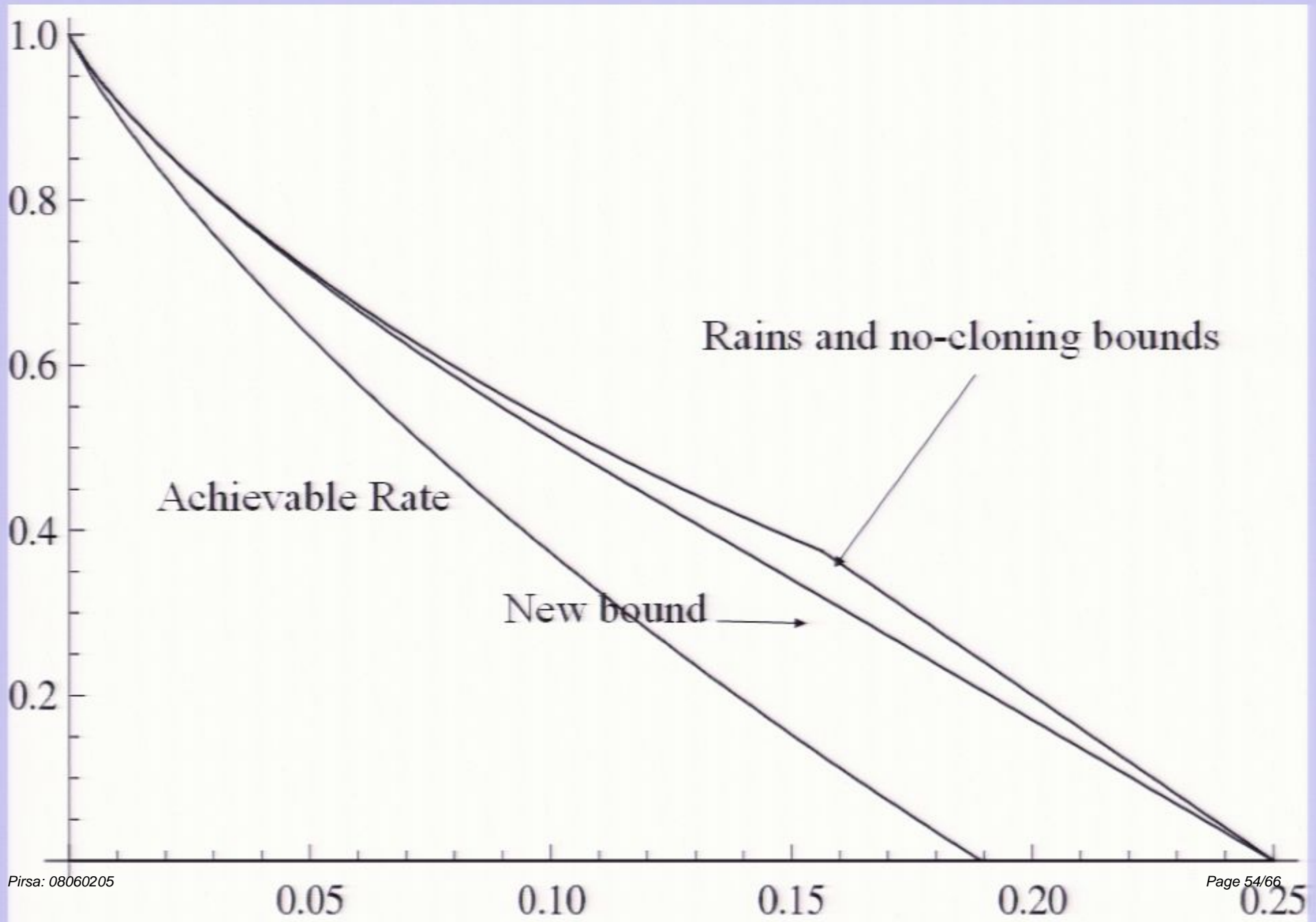
Rains

New bound

No cloning

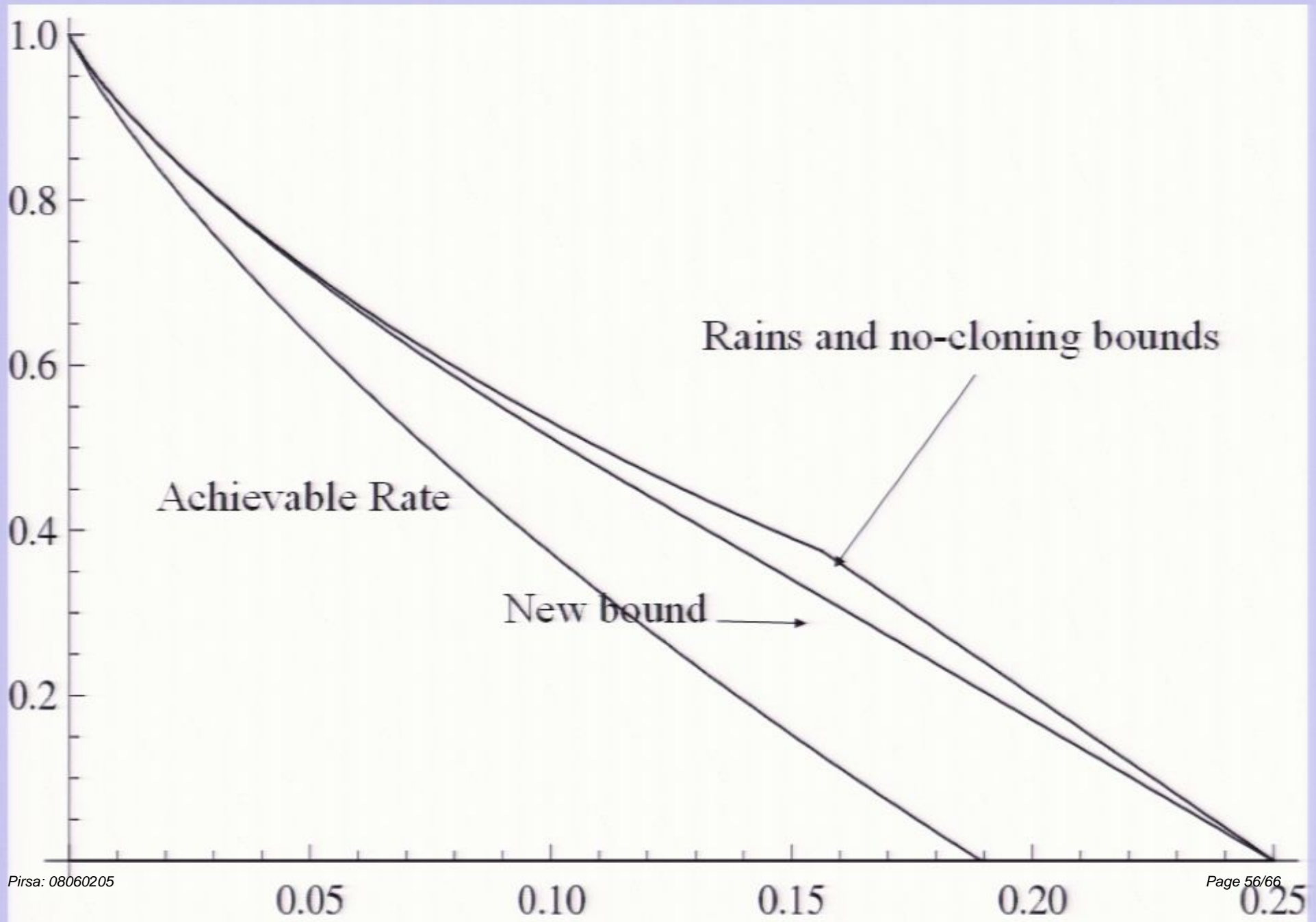
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# Another Result: BB84

$$\sqrt{q}^{\text{BB84}}(\rho) = (1 - q)^2 \rho + q(1 - q)X\rho X + q(1 - q)Z\rho Z + q^2 Y\rho Y$$





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Similar to depolarizing channel but with independent amplitude and phase noise

This channel has private capacity equal to that of BB84 with one-way postprocessing and bit error rate  $q$ .

$$C_p(\mathcal{N}_q^{\text{BB84}}) \leq H\left(\frac{1}{2} - 2q(1 - q)\right) - H(2q(1 - q))$$

# Amplitude Damping

$$\mathcal{N}_\gamma^{\text{ad}} : \quad A_0 = \begin{pmatrix} 1 & 0 \\ 0 & \sqrt{1-\gamma} \end{pmatrix}$$
$$A_1 = \begin{pmatrix} 0 & \sqrt{\gamma} \\ 0 & 0 \end{pmatrix}$$

degradable when  $\gamma \leq 1/2$

# BB84 extension

$$\tau_q^{\text{BB84}} = \frac{1}{2} \mathcal{N}_{\gamma(q)}^{\text{ad}}(\rho) \otimes |0\rangle\langle 0| + \frac{1}{2} Y \mathcal{N}_{\gamma(q)}^{\text{ad}}(Y \rho Y) Y \otimes |1\rangle\langle 1|$$
$$\gamma(q) = 4q(1 - q)$$

This is a degradable extension of the BB84 channel

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This is a degradable extension of the BB84 channel

Nice because it has just two terms

We can fully calculate its  $Q_1$

# SS Capacity

$$\text{symmetry : } V|(i, j)\rangle = \frac{1}{\sqrt{2}}(|i\rangle|j\rangle - |j\rangle|i\rangle)$$

$$\text{channel : } \mathcal{A}(\rho) = \text{Tr}_2 V \rho V^\dagger \quad Q(\mathcal{A}) = 0$$

$$Q_{\text{ss}}(\mathcal{N}) = \lim_{n \rightarrow \infty} \frac{1}{n} Q_1(\mathcal{N}^{\otimes n} \otimes \mathcal{A})$$

$$Q_{\text{ss}}(\mathcal{N}) = Q_1(\mathcal{N} \otimes \mathcal{A}) \quad \text{SSW2006}$$

Want to show  $\mathcal{N} \otimes \mathcal{A}$  is an additive extension of  $\mathcal{N}$   
with capacity  $Q(\mathcal{N} \otimes \mathcal{A}) = Q_{\text{ss}}(\mathcal{N})$

$$Q(\mathcal{N} \otimes \mathcal{A}) = \lim_{n \rightarrow \infty} \frac{1}{n} Q_1(\mathcal{N}^{\otimes n} \otimes \mathcal{A}^{\otimes n})$$



# A is an additive extension

$$\begin{aligned} Q(\mathcal{N} \otimes \mathcal{A}) &= \lim_{n \rightarrow \infty} \frac{1}{n} Q_1(\mathcal{N}^{\otimes n} \otimes \mathcal{A}^{\otimes n}) \\ &= \lim_{n \rightarrow \infty} \frac{1}{n} \left( Q_1(\mathcal{N}^{\otimes n} \otimes \mathcal{A}) + Q_1(\mathcal{A}^{\otimes(n-1)} \otimes \mathcal{A}) \right) \\ &= \lim_{n \rightarrow \infty} \frac{1}{n} (n Q_1(\mathcal{N} \otimes \mathcal{A})) = Q(\mathcal{N} \otimes \mathcal{A}) = Q_1(\mathcal{N} \otimes \mathcal{A}) \end{aligned}$$

End of slide show, click to exit.