Title: Fast Scrambling

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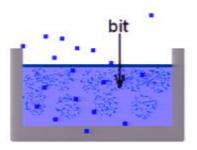
Abstract: TBA

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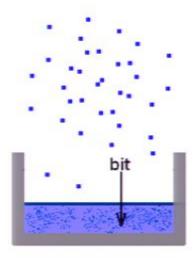
Fast Scrambling

 Patrick Hayden, John Preskill, 0708.4025

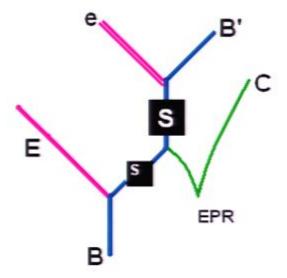
- Don Page, qc/9305007
- J. Lindesay, L. Susskind, World Sci 2005



time needed to recover bit = time to emit 1/2 the particles



time needed to recover bit
= scrambling time t*



How scrambled is SCRAMBLED?

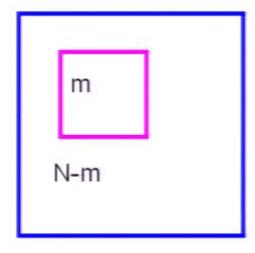
Haar scrambling

U chosen randomly with respect to Haar measure on Hilbert space of states.

Haar scambling is very inefficient. It takes non-polyomial steps.

But it is also OVERKILL.

Page scrambling (Don Page) is good enough.



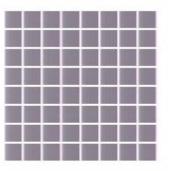
Page Scrambling means

$$S_{m} = log m + order (m/N)$$

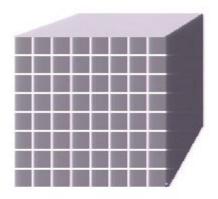
For all
$$m < N/2$$

but





$$d = 2$$



$$d = 3$$

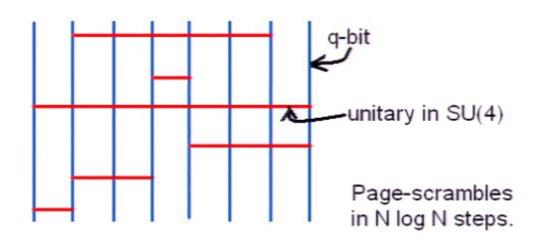
Typical scrambling time:

$$t_* = \beta N^{\frac{2}{d}}$$

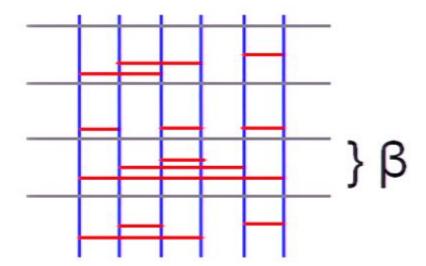
∃ systems that Page-scramble in time

$$t_* = \log N$$

Scrambler built from 2-bit unitary operations. Typically random but definite.



Parallel Processing



Time step = inverse temperature.

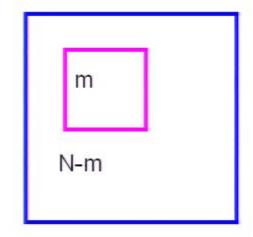
After m time-steps the 1st q-bit has directly or indirectly communicated with about 2^m others. The page-scrambling time is of order

That's fast!

$$t_* = \beta \log N$$

Hamiltonian Scramblers

$$|R\rangle = e^{iHt_*} |\Psi_0\rangle$$



$$P_{m} = TR_{N-m} \mid R > \langle R \mid$$

$$S_m = - tr \ P_m \log P_m$$

Page-thermal-scramblers satisfy

$$S_m = S_{thermal} + Order (m/N)$$

for
$$m < N/2$$

Fast Hamiltonian Scramblers?

Let the degrees of freedom be

$$X_1, X_2, X_3, X_4, \dots, X_N$$

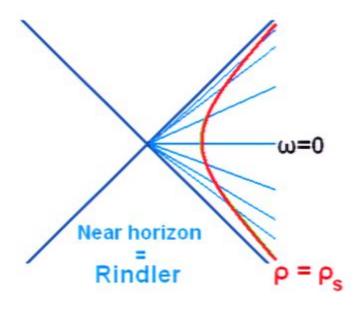
$$H = \sum_{i} F(X_i, X_j, X_k, X_l)$$

Clusters of fixed size.

Can
$$t * = \beta \log N$$
 ?

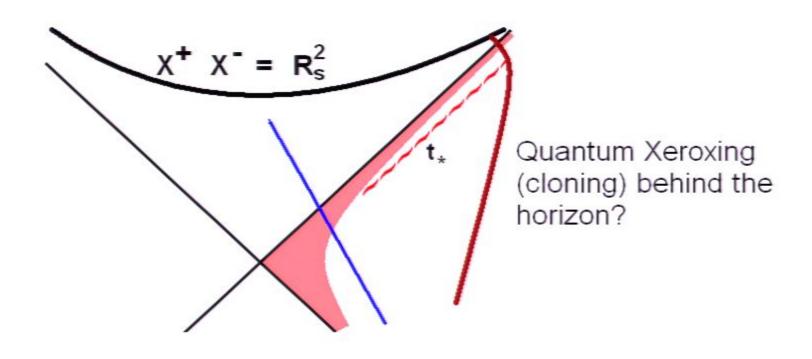
Black Holes are fast scramblers.

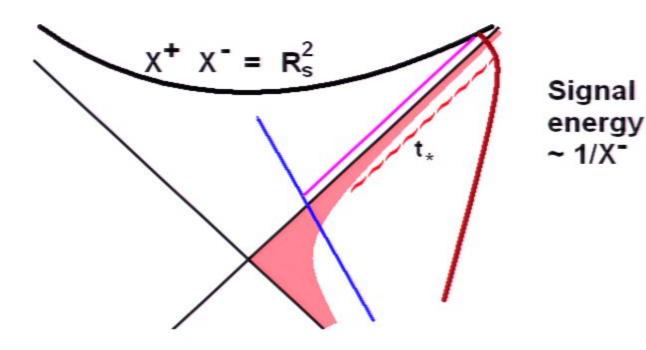
Hayden, Preskill Lindesay, LS



$$X^+ = \rho e^{\omega}$$

$$\omega = t/4MG = \beta t$$





$$M = \frac{R_s}{G} > \frac{1}{\chi}$$

$$X^+ > \frac{R_s^3}{G}$$

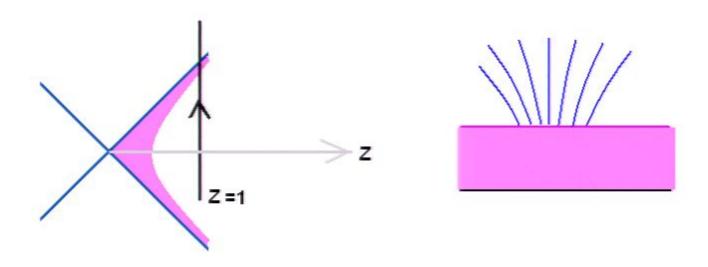
Recall $X^+ = \rho e^{\omega}$

Bound saturated when

$$\omega_* = \log \frac{R_s^3}{G\rho_0}$$

$$t_* = \beta \log \frac{R_s^3}{G\rho_0}$$

How fast does information (charge, mass, temperature) spread over the horizon? Let's examine an electric charge falling onto a horizon.

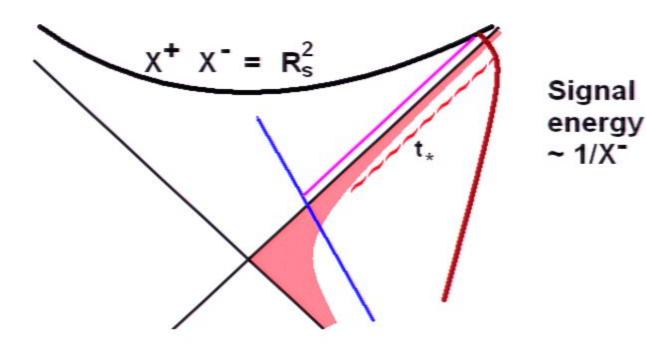


Radial electric field given by

$$E_{\rho} = \frac{Z-1}{[X^2 + Y^2 + (Z-1)^2]^{3/2}}$$

For large ω , on the stretched horizon this gives charge density

$$\sigma = \frac{\rho_0 e^{\omega}}{[X^2 + Y^2 + \rho_0^2 e^{2\omega}]^{3/2}}$$



$$M = \frac{R_s}{G} > \frac{1}{\chi}$$

$$X^+ > \frac{R_s^3}{G}$$

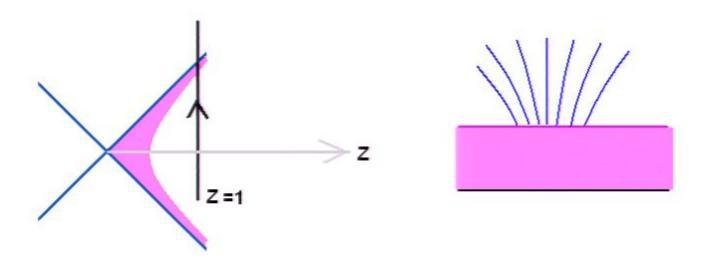
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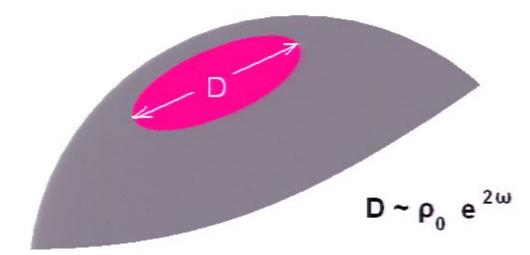


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Setting $D = R_s$ gives

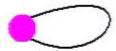
$$\omega_* = \log \frac{R_s}{\rho_0}$$

$$t_* = \beta \log \frac{R_s}{\rho_0}$$

$$\frac{\partial t_{*}T}{\partial \log S} \ge \hbar$$

Black holes are surely the fastest scramblers in nature. Consider a stellar mass black hole with a radius of about a kilometer and a temperature of .0000001degrees. Its scrambling time is .003 sec.

Concrete Hamiltonian? BFSS M(atrix) theory exactly describes a collection of N ten dimensional "D-particles" for large N.



x^a are quantum field operators for creating and anihilating a string attached to the DP. The string is "polarized" along the a axis where

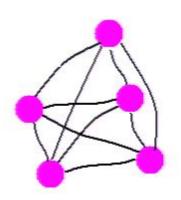
$$a = 1,...,9$$

x^a are the 9 spatial coordinates of the DP.

For slowly moving DP's

$$\mathcal{L} = \dot{X}^a \dot{X}^a$$

Now take an assembly of N D-particles



The "field operators" for strings connecting the m and n D-particles are (classical) matrices

X mn

$$\mathcal{L}_{kinetic} = \dot{X}_{mn}^{a} \dot{X}_{mn}^{a} = Tr \dot{X}^{a} \dot{X}^{a}$$

$$\mathcal{L}_{\text{interaction}} = \text{Tr} \{ X^a X^b X^a X^b - X^b X^a X^a X^b \}$$

Note: Every one of the 9N² matrix elements is coupled to every other one.

$$X_{mn}^{b}$$
 X_{nl}^{a} X_{lr}^{a} X_{rn}^{b}

Add energy to D-particle system and it becomes a 10-dimensional non-extremal black hole with a Rindler horizon.

ds =
$$f^{-1/2}$$
 [-(1- R_H^7 / r^7) dt^2 ] Horowitz, Strominger

Using 10-D general relativity we can compute the horizon radius in terms of the temperature **T**, and the charge, **N**.

Then using the (by now) standard (Rindler) formula

$$\omega_* = \log \frac{R_H}{\rho_0}$$

one finds

$$t_* = \beta \log \frac{N}{\beta}$$

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The dualities of String Theory require the M(atrix) Theory hamiltonian to be a fast scrambler.

In fact it has the correct properties:

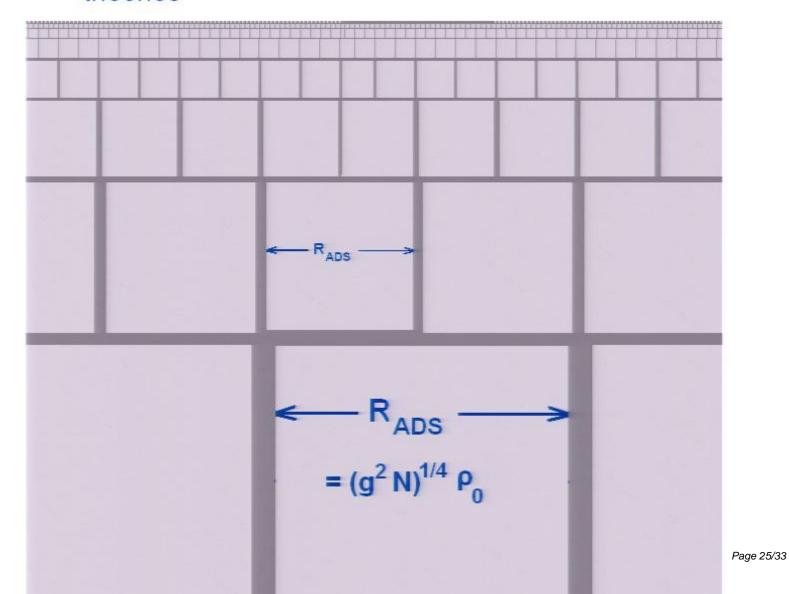
- The individual terms in H involve at most 4 operators.
- Every element is directly coupled to every other element.

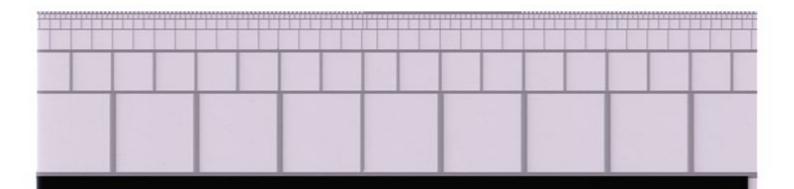
But as usual in String Theory, no direct proof exists.

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Scrambling in ADS/CFT

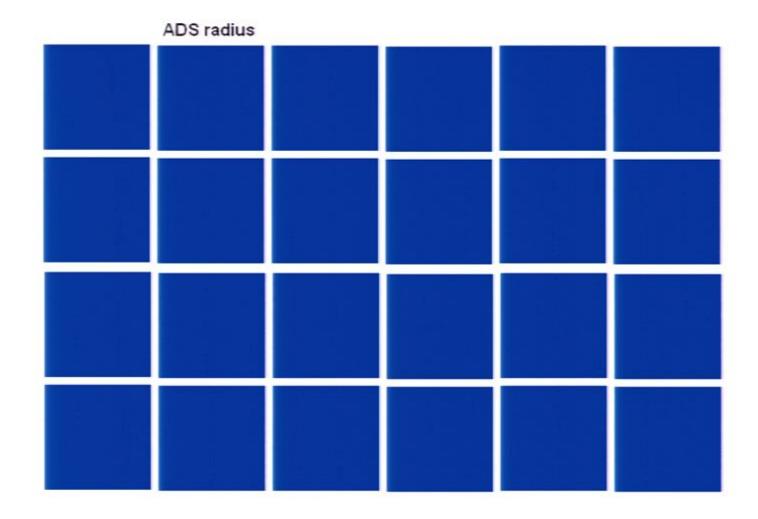
Each cutoff cell is described by N² degrees of freedom. ADS is a <u>lattice</u> of coupled matrix theories



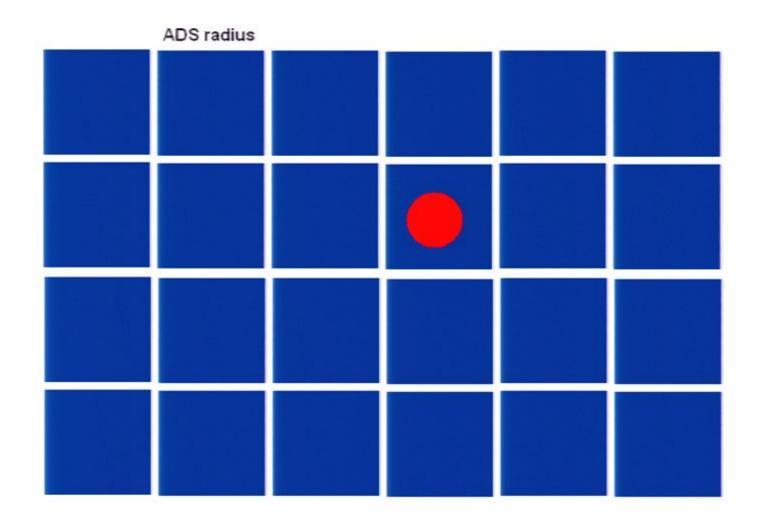


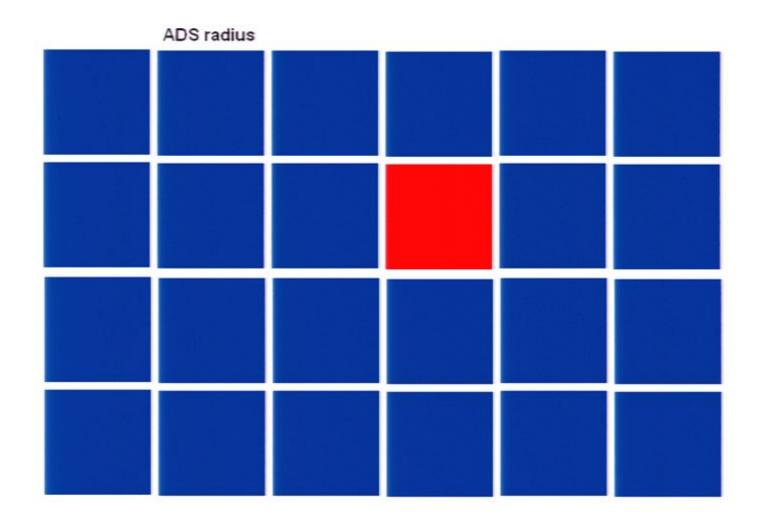
Using $t_* = \log R_{ADS}/\rho_0$ and $R_{ADS}=(g^2N)^{1/4}$ one finds that a single ADS volume is scrambled in time of order $t_* = \beta \log N$

On larger scales the spreading is slow characteristic of a 3+1 dimensional system.

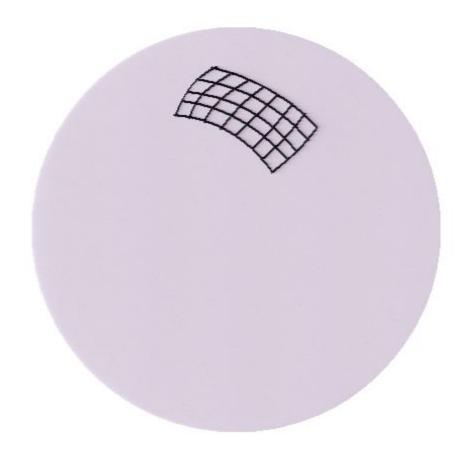


Looking down

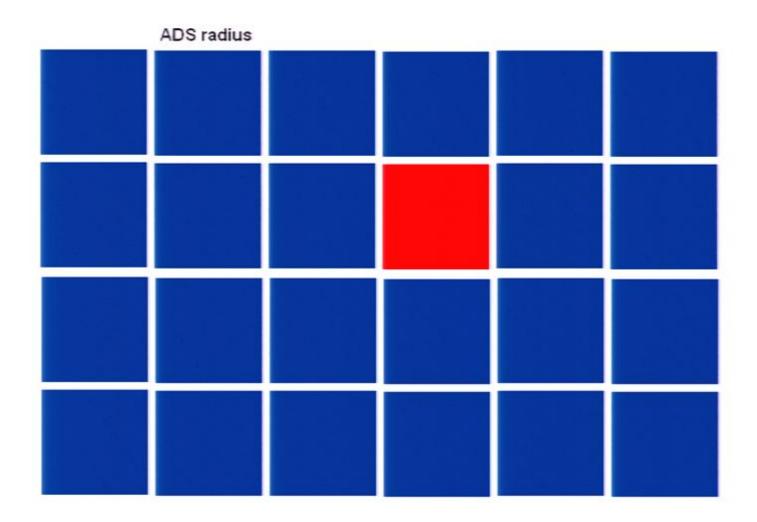




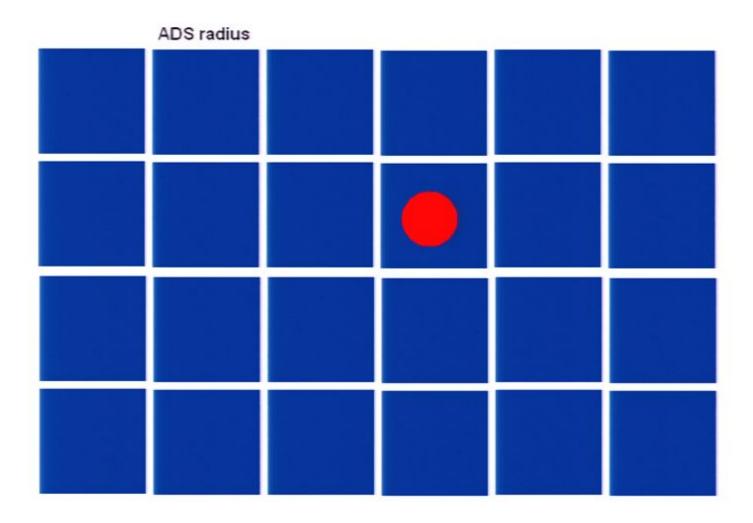
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There can be no sense in which the horizon is a two dimensional collection of "pixels" interacting locally. Every pixel is coupled to every other pixel to create an ultra fast information scrambler. In a definite sense it is infinite dimensional.



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