Title: Interaction and information flow between quantum systems (Part 1B)

Date: Aug 30, 2007 04:30 PM

URL: http://pirsa.org/07080053

Abstract:

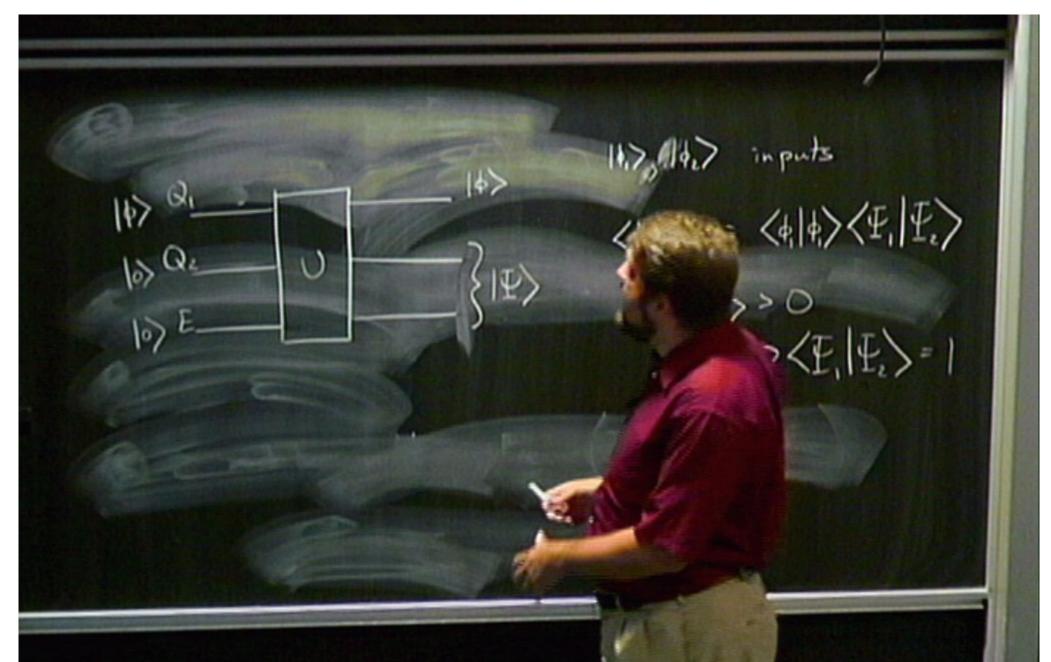
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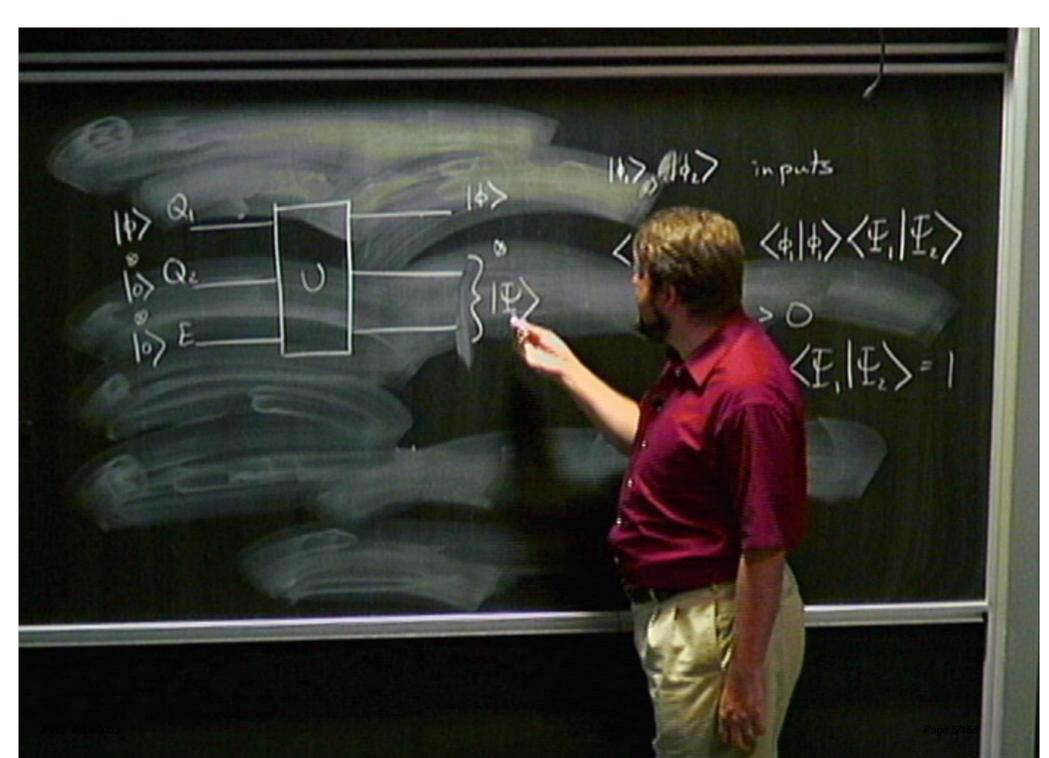
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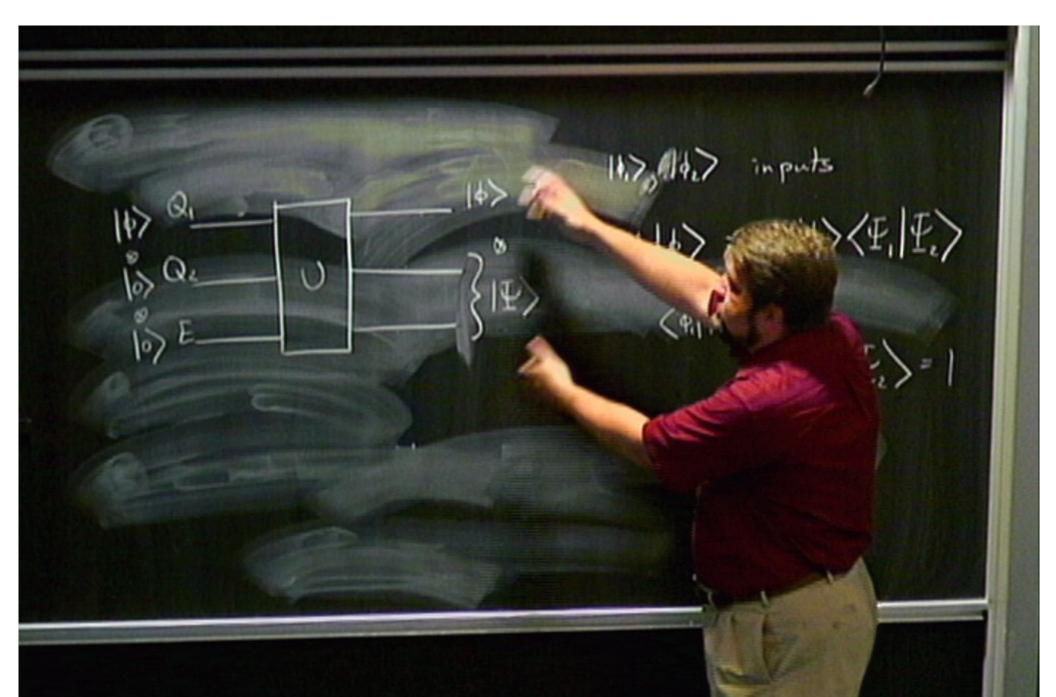
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Page 2/18

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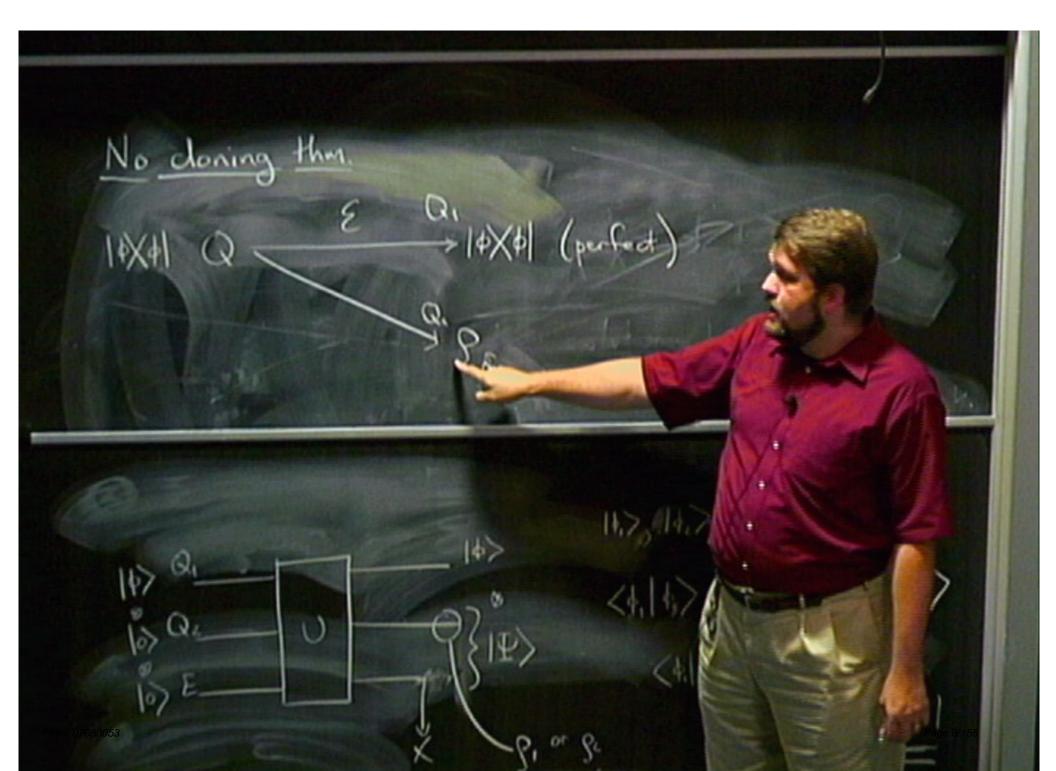


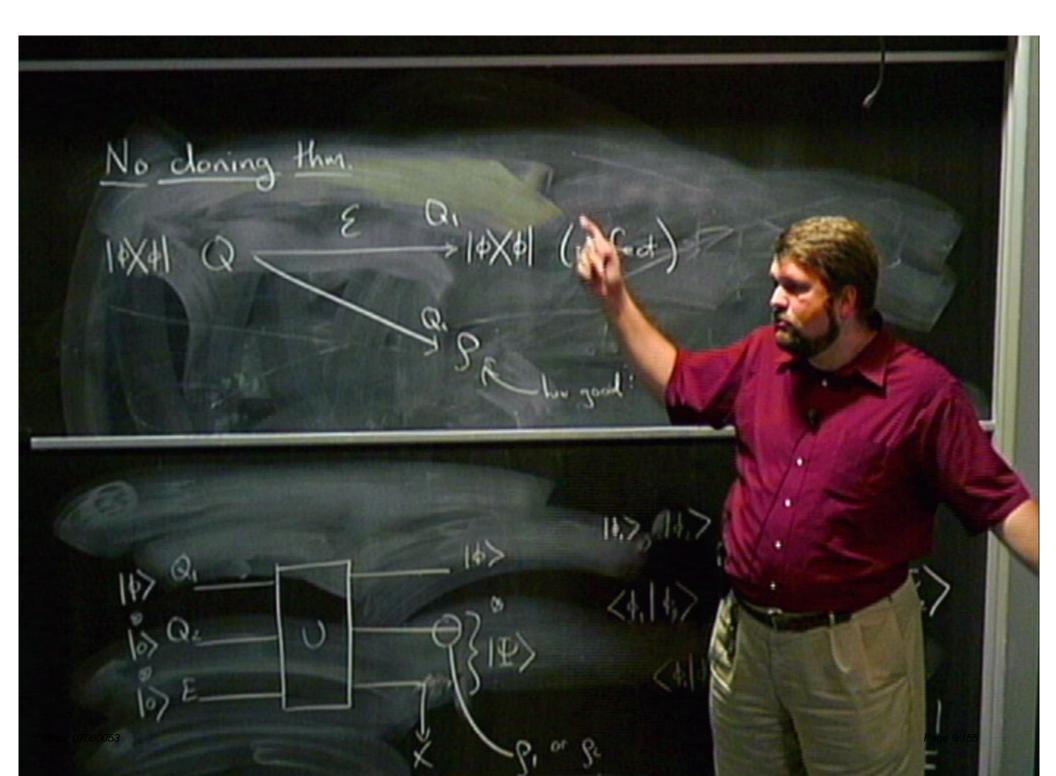


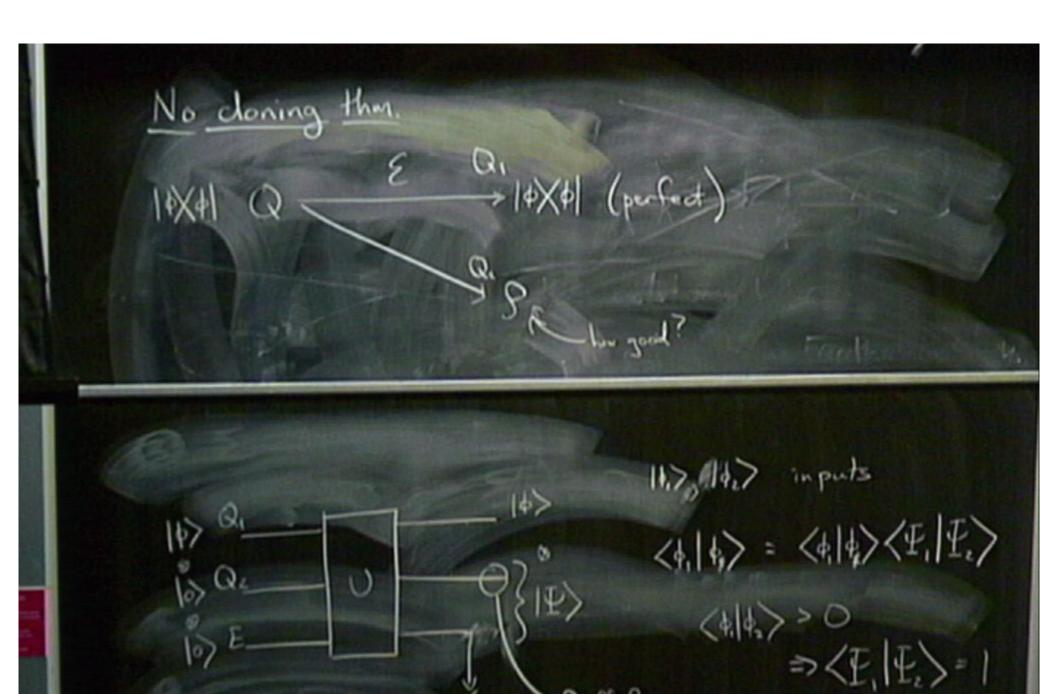


Page 6/15

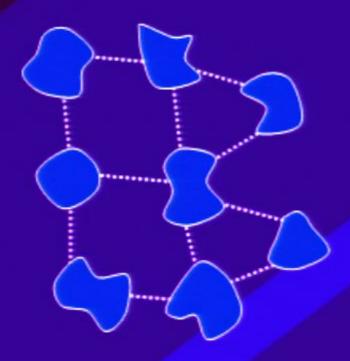
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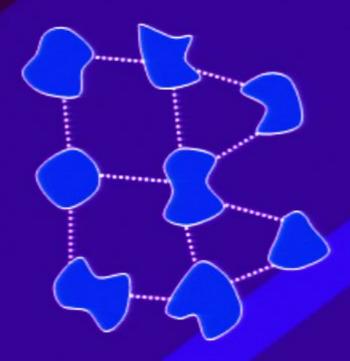
Page 10/155



Universe is divided into subsystems.

Subsystems interact and exchange information.

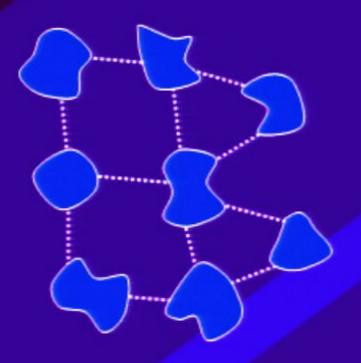
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Pirsa: 07080053 Page 12/15.

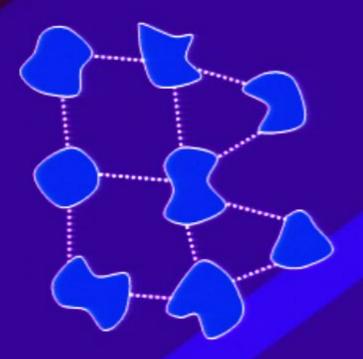


Universe is divided into subsystems.

Subsystems interact and exchange information.

Locality: Not all subsystems exchange information directly.

Pirsa: 07080053 Page 13/155

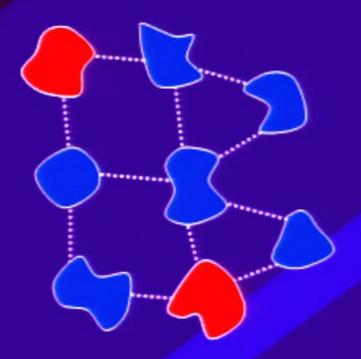


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Pirsa: 07080053 Page 14/155

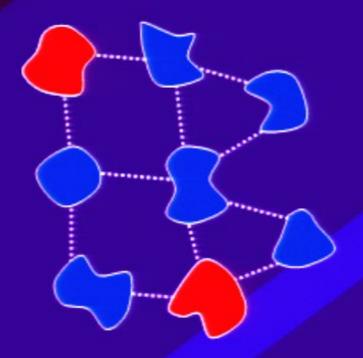


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Pirsa: 07080053 Page 15/155



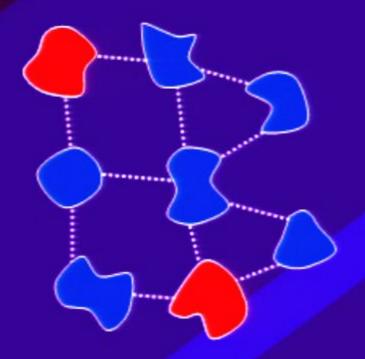
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What does quantum mechanics say about the rules of this web?

Pirsa: 07080053 Page 16/155



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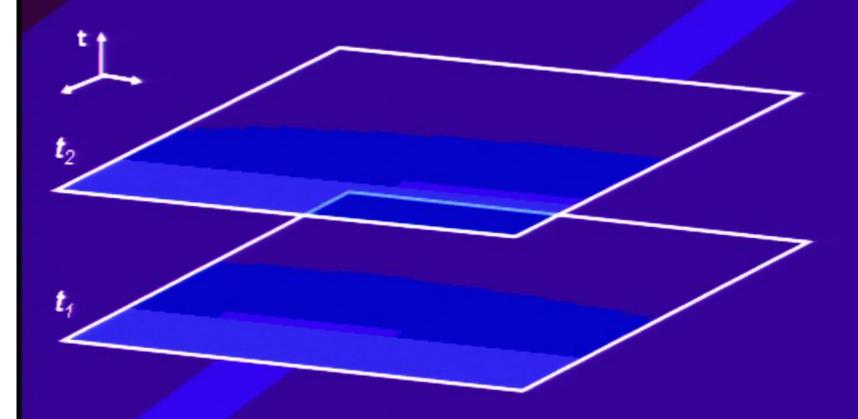
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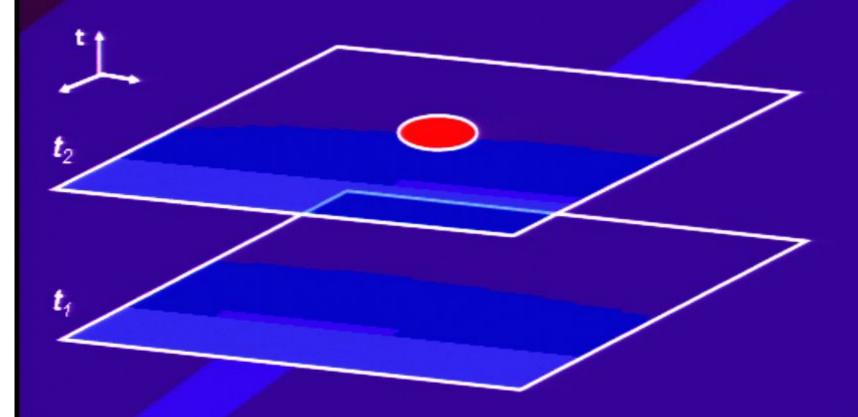
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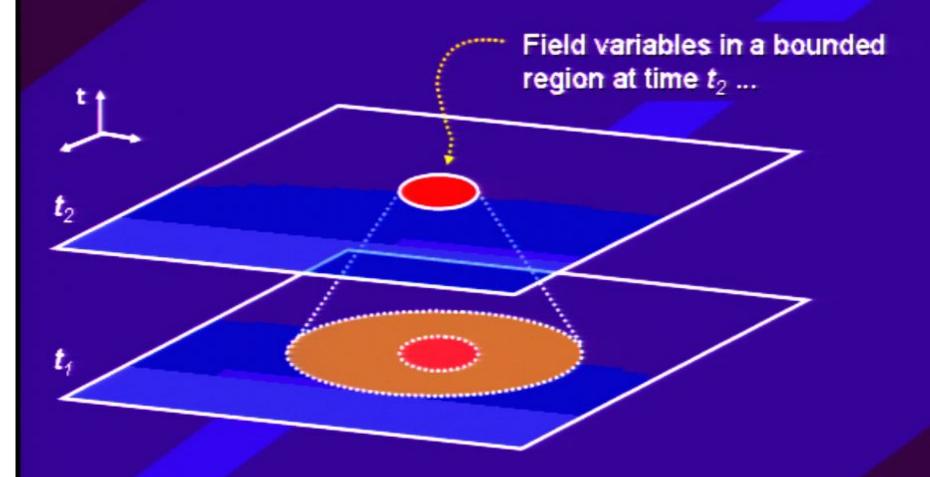
What does quantum mechanics say about locality?

Pirsa: 07080053 Page 17/155

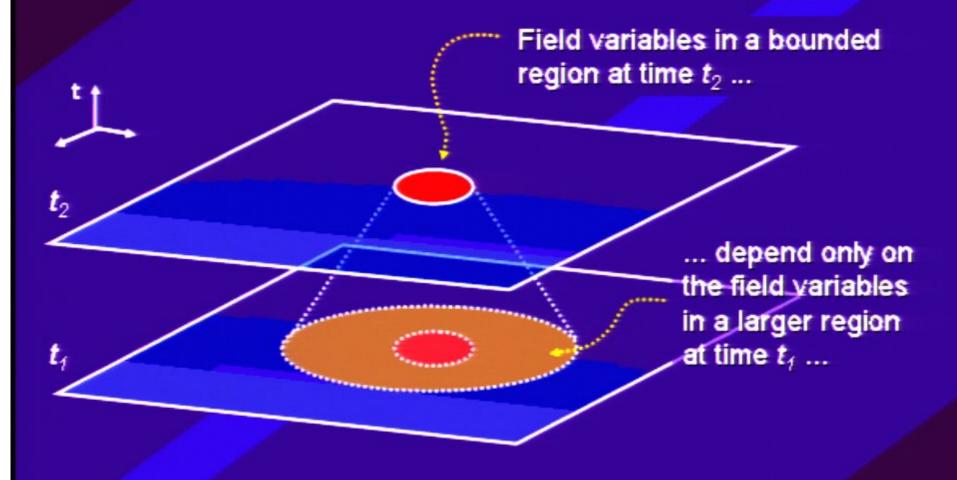
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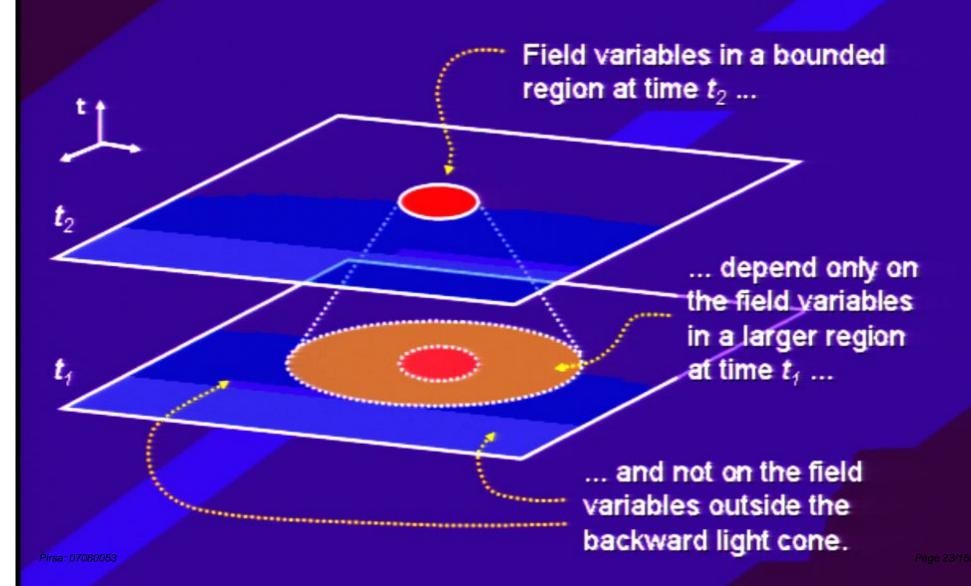






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A is the system of interest.

Pirsa: 07080053 Page 26/15.



A is the system of interest.

C is the distant "rest of the world"

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B is the rest of A's "neighborhood"

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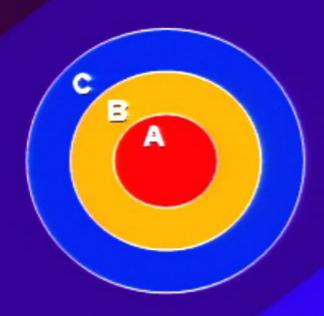


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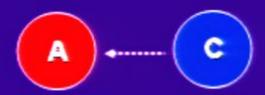
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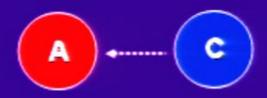
Locality: In one time step, there is no information transfer from C to A.

Pirsa: 07080053 Page 31/



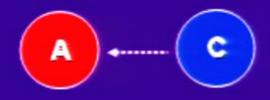
Pirsa: 07080053 Page 32/155

When does information "flow" from C to A?



Pirsa: 07080053 Page 33/15.

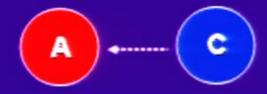
When does information "flow" from C to A?



 Information flows from C to A if the final state of A depends on the initial state of C.

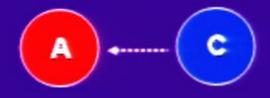
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When does information "flow" from C to A?



- Information flows from C to A if the final state of A depends on the initial state of C.
- Information does not flow from C to A if the final state of A does not depend on the initial state of C.





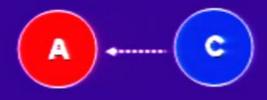
Note: We must consider all possible initial states of A and C.

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Information flow

When does information "flow" from C to A?



Note: We must consider all possible initial states of A and C.

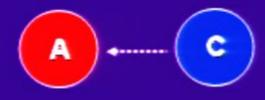
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Quantum difficulties!

Pirsa: 07080053 Page 37/155

Information flow

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Quantum difficulties!

Initial state of AC is not determined by the initial states of A and C separately – quantum entanglement.

Pirsa: 07080053 Page 39/15

Two classical bits.

Interaction: Controlled-NOT

C = control bit

T = target bit

Pirsa: 07080053 Page 40/155

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Pirsa: 07080053 Page 41/15.

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 $\begin{array}{c} CT \rightarrow CT \\ 0.0 \rightarrow 0.0 \\ 0.1 \rightarrow 0.1 \\ 1.0 \rightarrow 1.1 \\ 1.1 \rightarrow 1.0 \end{array}$

Two classical bits.

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Note: CNOT operation is reversible

 $\begin{array}{c} CT \rightarrow CT \\ 0.0 \rightarrow 0.0 \\ 0.1 \rightarrow 0.1 \\ 1.0 \rightarrow 1.1 \\ 1.1 \rightarrow 1.0 \end{array}$

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Final T state does depend on initial C state. There is information flow from C to T.

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Classical CNOT has one-way information flow from C to T.

Two qubits.

Interaction: quantum CNOT

C = control bit

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$$\begin{array}{c} |\text{CT}\rangle \rightarrow |\text{CT}\rangle \\ |0\ 0\rangle \rightarrow |0\ 0\rangle \\ |0\ 1\rangle \rightarrow |0\ 1\rangle \\ |1\ 0\rangle \rightarrow |1\ 1\rangle \\ |1\ 1\rangle \rightarrow |1\ 0\rangle \end{array}$$

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CNOT is unitary

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Two qubits.

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One-way information flow? No!



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C



|CT⟩ → |CT⟩

 $|0 0\rangle \rightarrow |0 0\rangle$

 $|0 1\rangle \rightarrow |0 1\rangle$

 $|10\rangle \rightarrow |11\rangle$

 $|11\rangle \rightarrow |10\rangle$

One-way information flow? No!

CNOT is unitary

Look at CNOT in a conjugate basis:

$$|+\rangle = \frac{1}{\sqrt{2}} (|\mathbf{0}\rangle + |\mathbf{1}\rangle)$$

$$\left|-\right\rangle = \frac{1}{\sqrt{2}} \left(\left| \mathbf{0} \right\rangle - \left| \mathbf{1} \right\rangle \right)$$

Two qubits.

Interaction: quantum CNOT

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 $|\mathsf{CT}\rangle \to |\mathsf{CT}\rangle$

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$$|++\rangle \rightarrow |++\rangle$$

$$|+-\rangle \rightarrow |--\rangle$$

$$|-+\rangle \rightarrow |-+\rangle$$

$$|--\rangle \rightarrow |+-\rangle$$

Two qubits.

Interaction: quantum CNOT

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 $|CT\rangle \rightarrow |CT\rangle$

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$$|--\rangle \rightarrow |+-\rangle$$

In the conjugate basis, control and target qubits switch roles!

Two qubits.

Interaction: quantum CNOT

C = control bit

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 $|\text{CT}\rangle \rightarrow |\text{CT}\rangle$

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$$|-+\rangle \rightarrow |-+\rangle$$

$$|--\rangle \rightarrow |+-\rangle$$



Pirsa: 07080053 Page 55/158

 No unitary interaction leads to one-way information flow between quantum systems.

Pirsa: 07080053 Page 56/155



11,7 1/42 inputs < 1 | 4> = < 4 | 4> < 4 | 4> < 4 | 4> <

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- Quantum measurement

Pirsa: 07080053 Page 59/155

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Pirsa: 07080053 Page 60/155

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C = system of interest

Pirsa: 07080053 Page 61/155

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T = measuring apparatus

Pirsa: 07080053 Page 62/153

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Pirsa: 07080053 Page 63/15.

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We'd like to have information flow $C \rightarrow T$ only, so that we do not disturb the system. But any unitary interaction can make information flow either way.

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In general, we can have CP maps between states of different systems.

Pirsa: 07080053 Page 68/15

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Of particular interest: An initial state of a composite system leads to a final state of one subsystem.

E.g., Partial trace operation tr_R

Pirsa: 07080053 Page 69/158

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$$\hat{\rho}^{Q} = \mathcal{E}_{RQ}^{Q}(\rho^{RQ})$$

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E.g., Partial trace operation tr_R

$$\hat{\mathbf{p}}^{\mathcal{Q}} = \mathcal{E}_{RQ}^{\mathcal{Q}}(\mathbf{p}^{RQ})$$

We may also have maps with the final system larger than the initial system

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Global evolution map \mathcal{E}^{ABC} What does it mean to say "No information flows from C to A"?

Pirsa: 07080053 Page 74/15



Global evolution map \mathcal{E}^{ABC} What does it mean to say "No information flows from C to A"?

Locality #1

There exists a CP map giving final A states from initial AB states. That is,

 \mathcal{E}^{A}_{AR} exists



Global evolution map EABC

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Locality #2

If ABC is initially in a product pure state

 $|\alpha\rangle\otimes|\beta\rangle\otimes|\gamma\rangle$

then the final state of A does not depend on |y>



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Locality #3

Given any initial ABC state. Do the following:

- 1. Operation on C.
- 2. EABC

Final A state does not depend on choice of previous C operation.



Global evolution map EABC

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Good news: All three locality conditions for \mathcal{E}^{ABC} are equivalent.

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Can we have

- Global evolution of ABC unitary; and
- No information transfer from C to A?

Pirsa: 07080053 Page 80/15

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C _____

B _____

A —

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В

Yes, of course. Trivial cases:

Can we have

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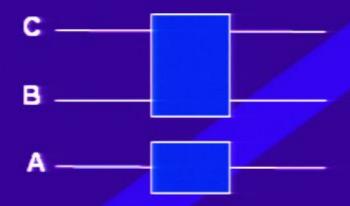
С В А

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A or C are isolated.

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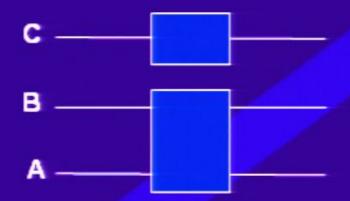
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Pirsa: 07080053 Page 84/15.

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Pirsa: 07080053 Page 85/15.

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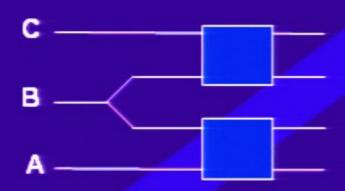
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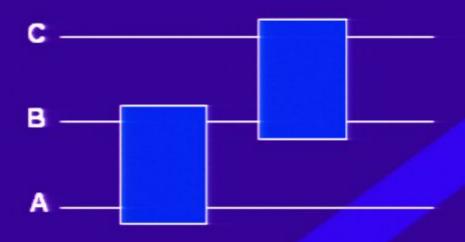
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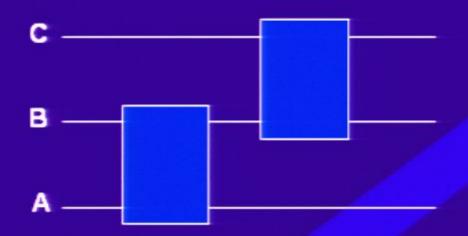
A and C interact separately with parts of a composite system B.

Pirsa: 07080053 Page 87/15.

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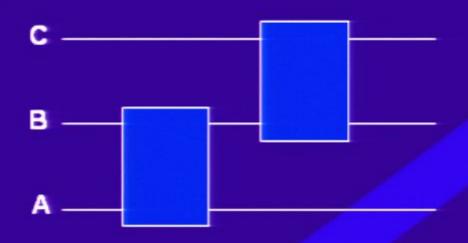


Pirsa: 07080053 Page 89/155



Points to note

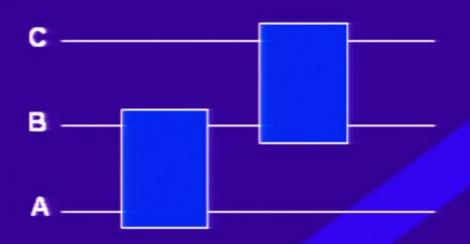
Pirsa: 07080053 Page 90/15



Points to note

 AB interaction followed by BC interaction.

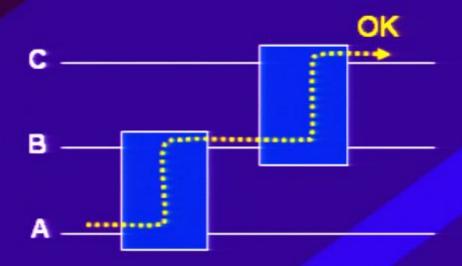
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Points to note

- AB interaction followed by BC interaction.
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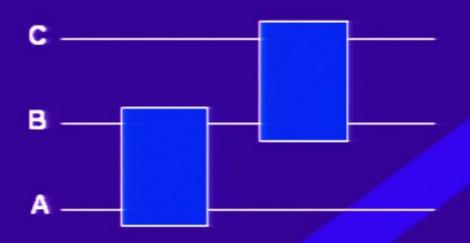
Pirsa: 07080053 Page 92/15.



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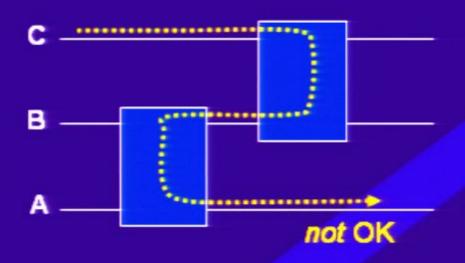
Pirsa: 07080053 Page 93/15.



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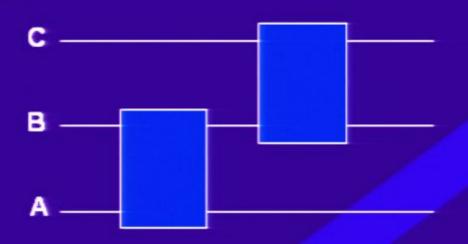
Pirsa: 07080053 Page 94/15.



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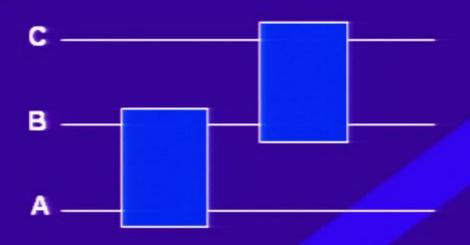
Pirsa: 07080053 Page 95/15



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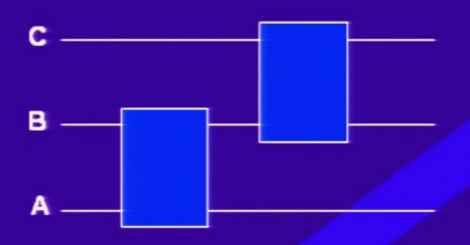
Pirea: 07080053



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- One-way information transfer: A → C but not C → A
- Previous examples can be converted to this general form

Pirsa: 07080053 Page 97/155



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- AB interaction followed by BC interaction.
- One-way information transfer: A → C but not C → A
- Previous examples can be converted to this general form

Remarkable fact: This is the only possibility!

Pirsa: 07080053

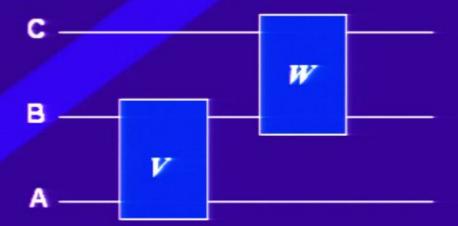
A decomposition theorem

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A decomposition theorem

Suppose system ABC evolves via unitary UABC, such that no information transfer is possible from C to A ("locality"). Then

$$U^{^{ABC}} = (1^{^{A}} \otimes W^{^{BC}})(V^{^{AB}} \otimes 1^{^{C}})$$

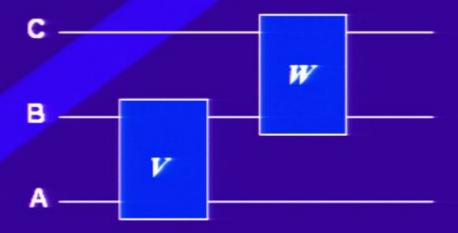


Pirsa: 07080053 Page 100/15

A decomposition theorem

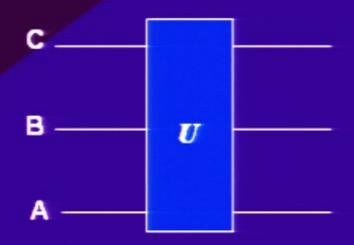
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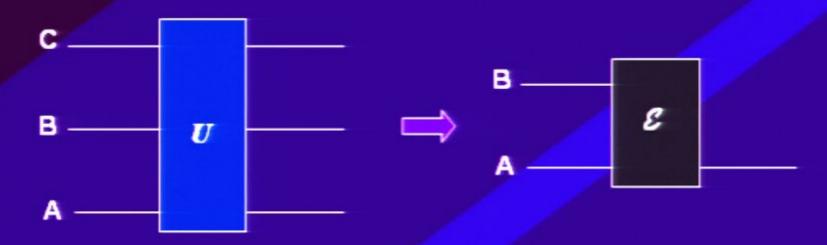


NB – We are not claiming that U actually happened this way

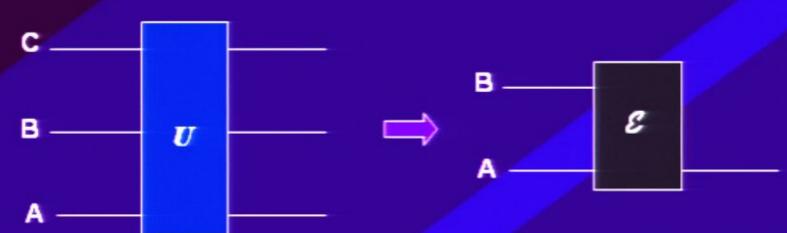
Pirsa: 07080053



Pirsa: 07080053 Page 103/15

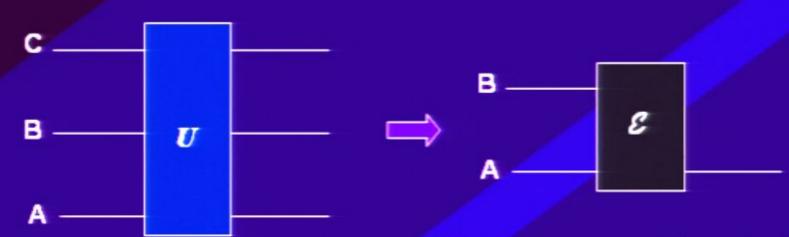


Pirsa: 07080053 Page 104/15



For every pure state output $|\psi\rangle$, construct the precursor subspace

$$S_{\psi} = \{ |\phi\rangle : \mathcal{E}(|\phi\rangle\langle\phi|) = \lambda |\psi\rangle\langle\psi| \}$$



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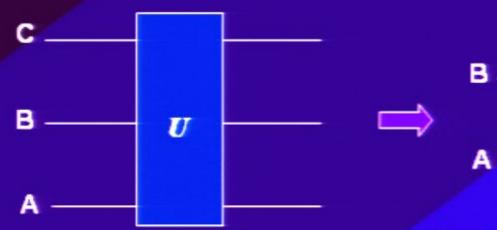


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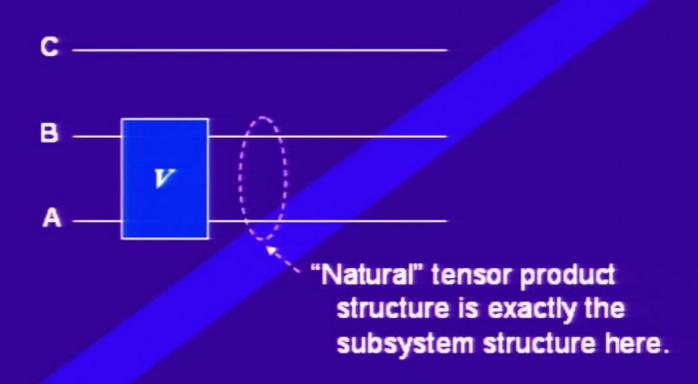
This is not the same as $\mathcal{A}_A \otimes \mathcal{A}_B - it$'s twisted by a unitary transformation \mathbf{V}^{\dagger}

Pirsa: 07080053 Page 109/15

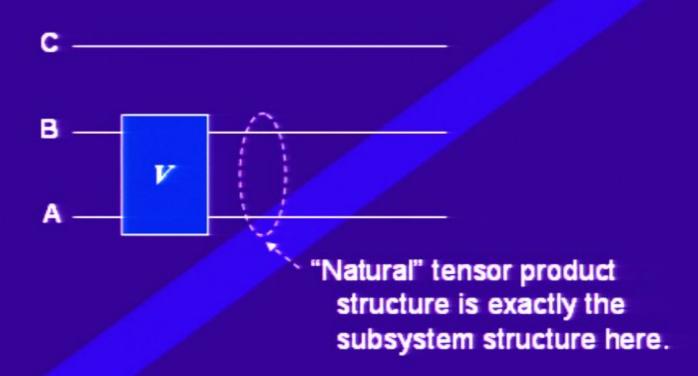
C _____



"Natural" tensor product structure is exactly the subsystem structure here.

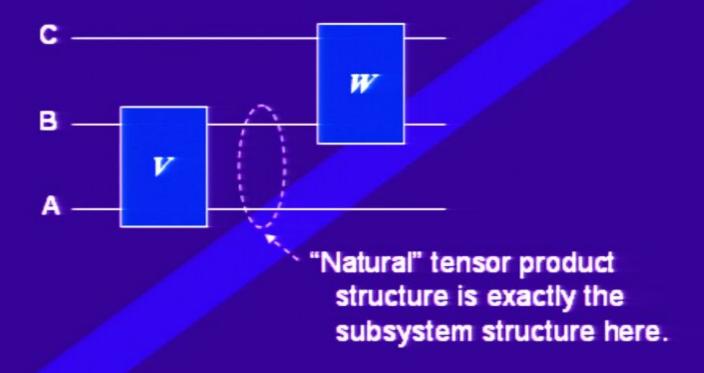


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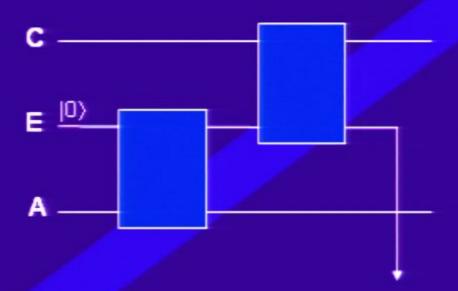


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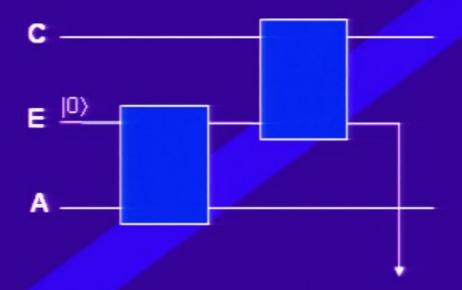
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Suppose \mathcal{E}^{AC} is a CP map such that no information is transferred from C to A. Then there is a unitary representation for \mathcal{E}^{AC} of the form



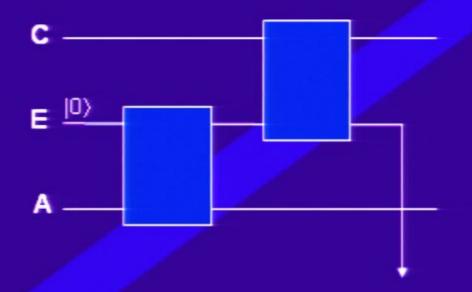
Pirsa: 07080053 Page 115/155

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Semicausal operations are semilocalizable Beckman et. al. (2001) Eggeling et al. (2002)

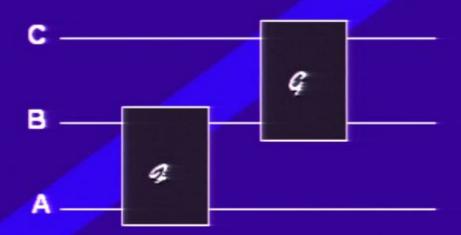
irsa: 07080053 Page 118/15.

Suppose ABC evolves according to a general CP map \mathcal{E} , and no information is transferred from C to A.

Pirsa: 07080053 Page 119/155

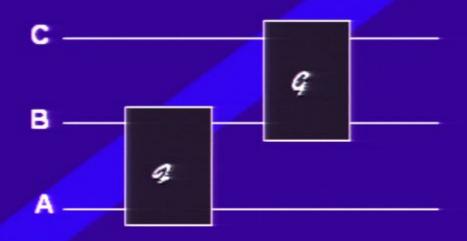
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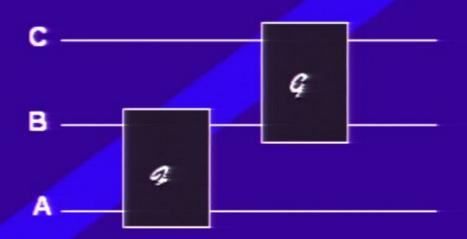


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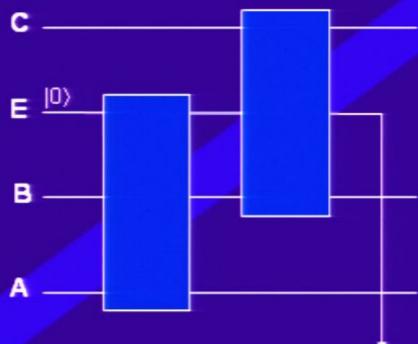
However

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Locality in general

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Suppose \mathcal{E}^{ABC} is a general CP map that is local – that is, no information can flow from C to A. Then the map has a unitary representation of the form:



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Reversible to unitary

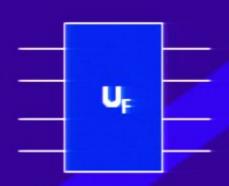
Reversible to unitary



Global update rule for N classical discrete variables:

$$(a,b,...,c) \rightarrow F(a,b,...,c)$$

Reversible map: F is 1-1 and onto.



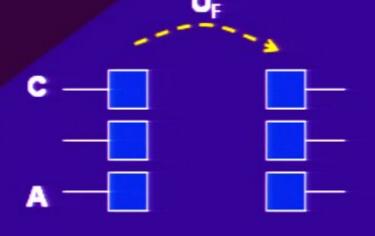
"Quantizing" the classical map

Product basis: | a,b,...,c >

Global unitary dynamics:

$$U_F | a,b,...,c \rangle = | F(a,b,...,c) \rangle$$

Quantum locality



Quantize previous problem (array of N finite quantum systems) with U_F

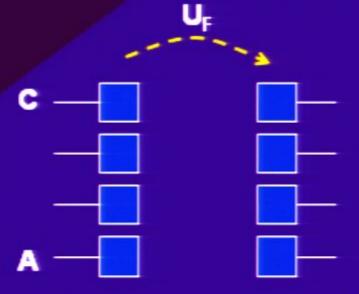
Is the global evolution U_F "local enough" to prevent information flow from C to A?

N = 3: Even though no information can flow from C to A in the classical case, it sometimes can in the quantum case.

Pirsa: 07080053 Page 127/15



Quantum locality



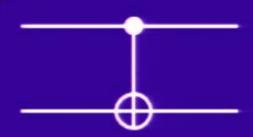
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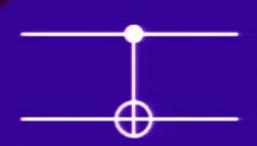
N = 3: Even though no information can flow from C to A in the classical case, it sometimes can in the quantum case.

N = 5: No information can flow even in the quantum case.

N = 4: No information can flow if the are qubits. We do not know the ageneral.

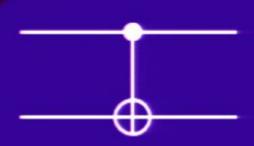


Pirsa: 07080053 Page 130/15



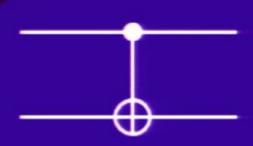
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Pirsa: 07080053 Page 131/15



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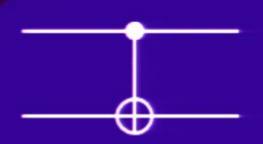
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What is the structure of information flow inside CNOT?



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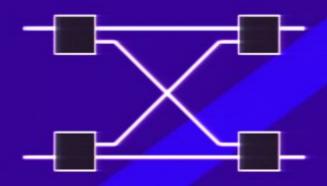


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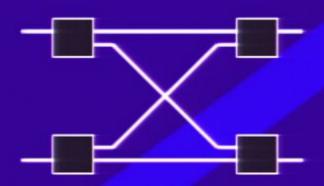


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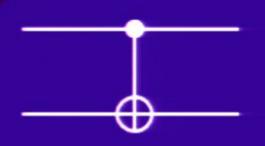
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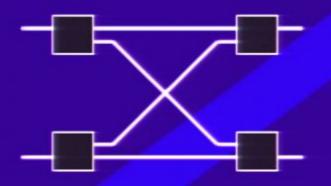
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Note that every classical gate can be modeled in this way. (Exchange copies!)



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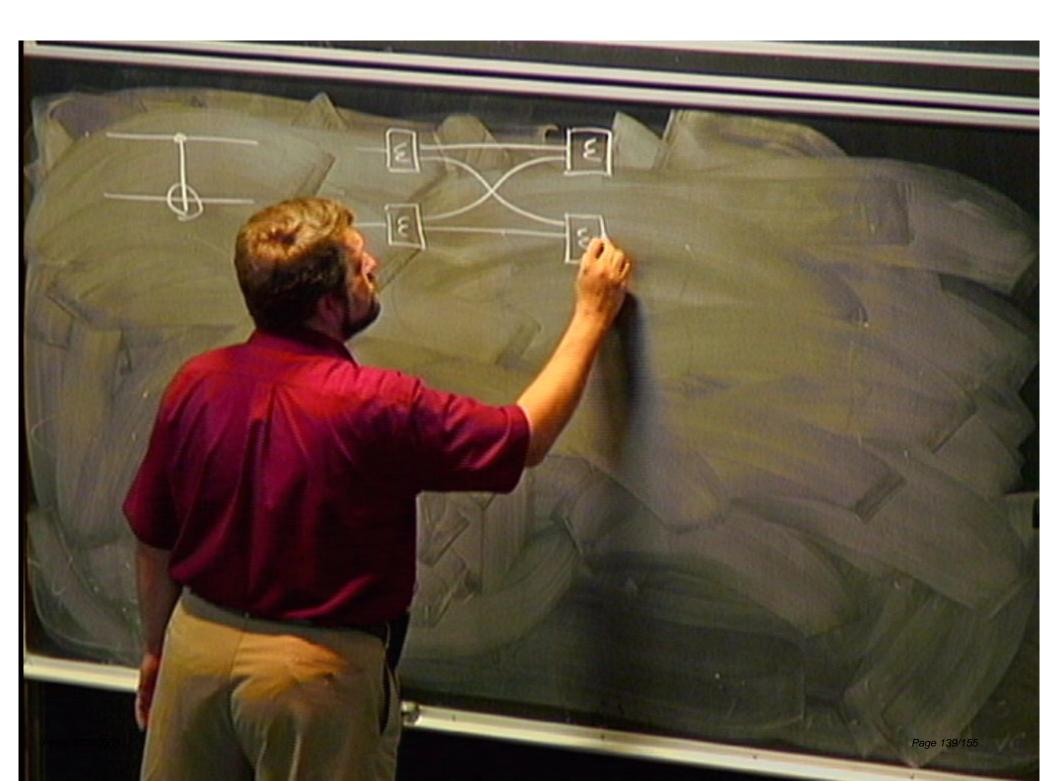
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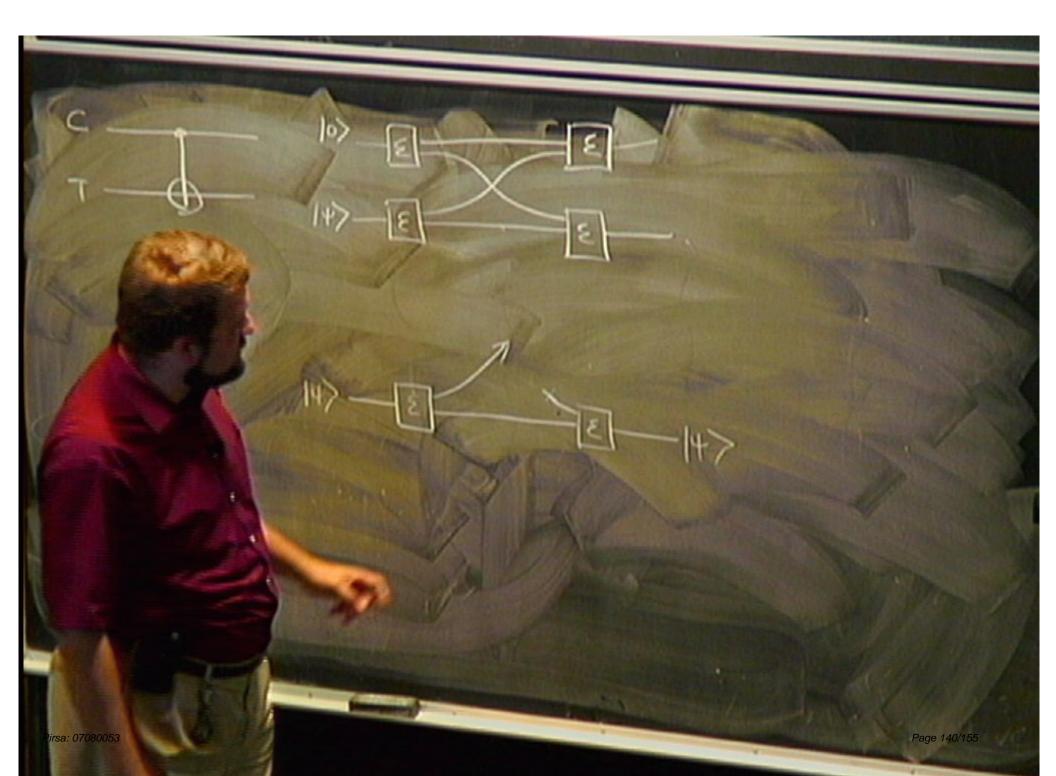


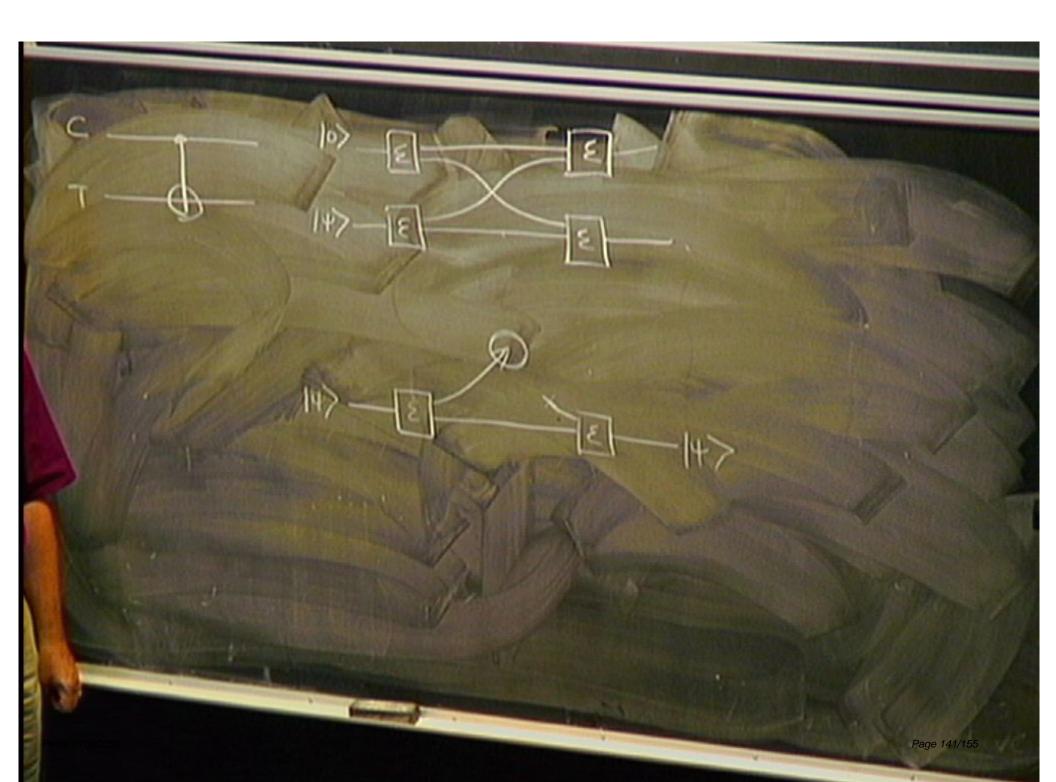
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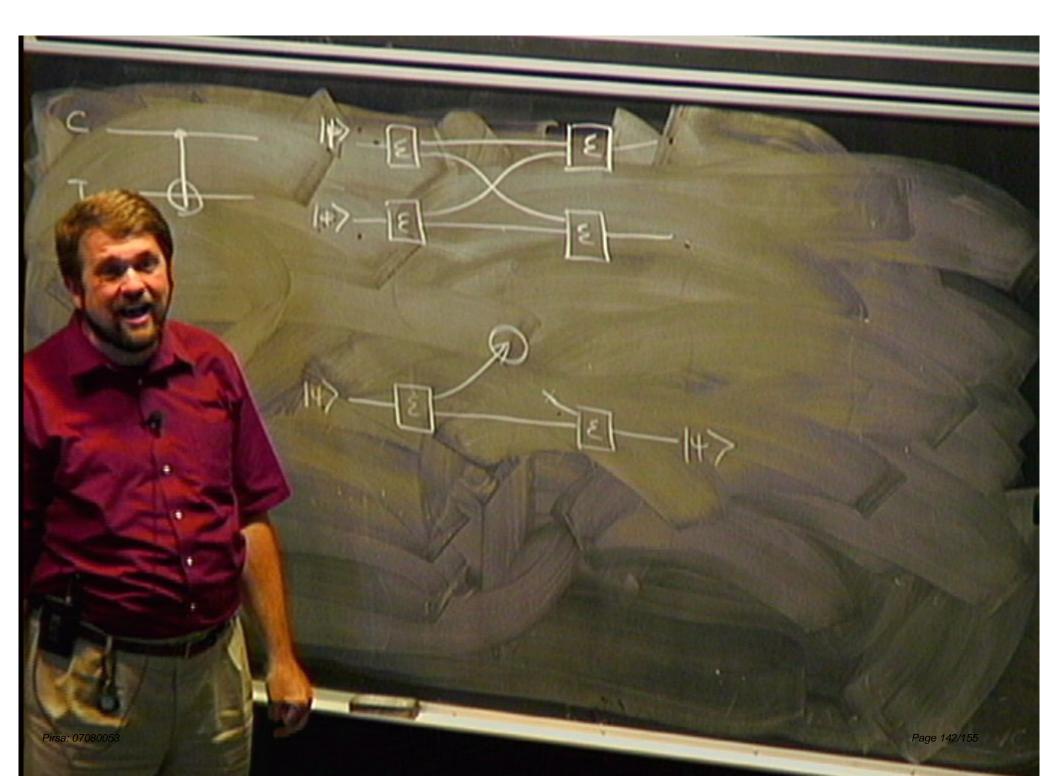
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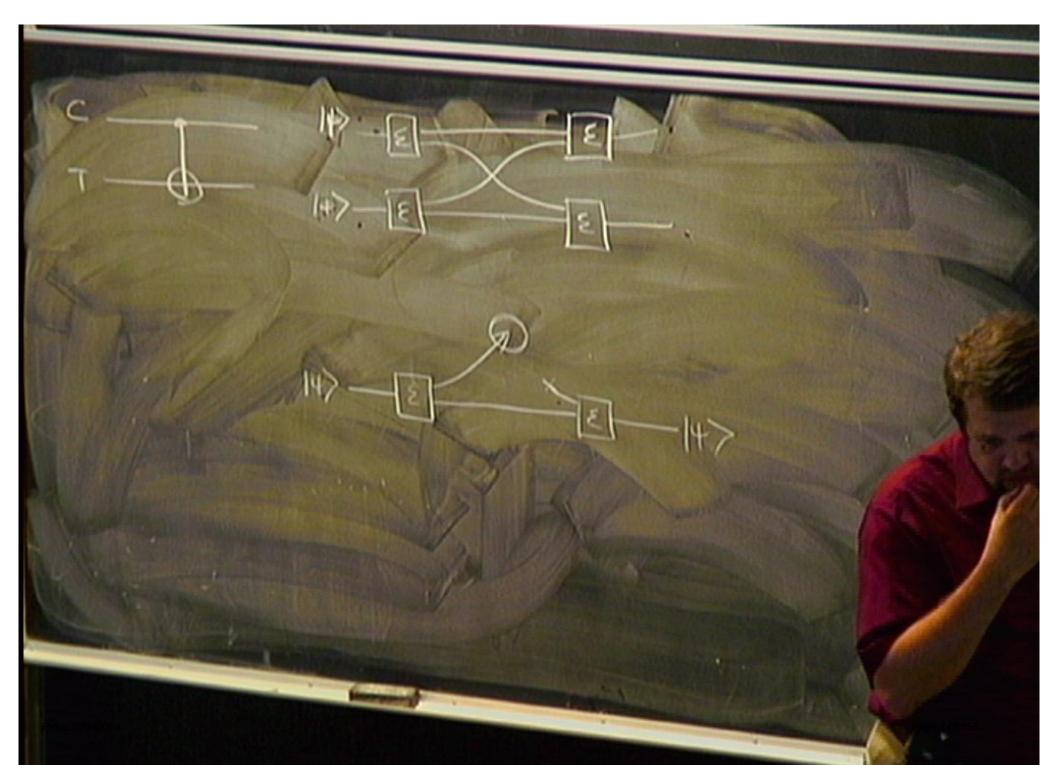
Can CNOT be modeled by local CP maps and simple information exchange?

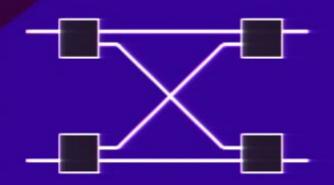








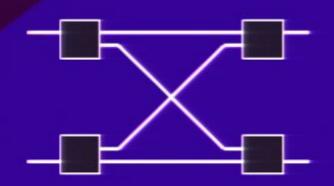




M. Nathanson: No entangling unitary twoqubit gate can be modeled by local CP maps and simple information exchange

(This holds only for qubit gates!)

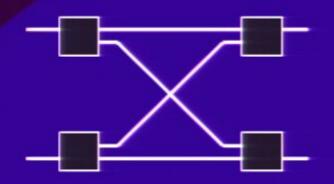
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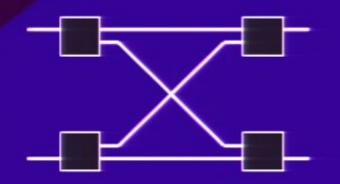
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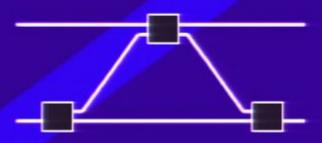
Here are two ways that you can model CNOT:



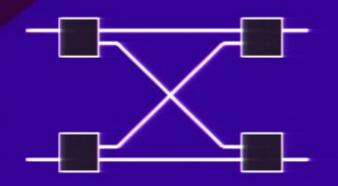
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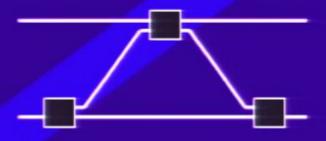
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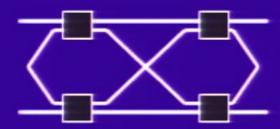


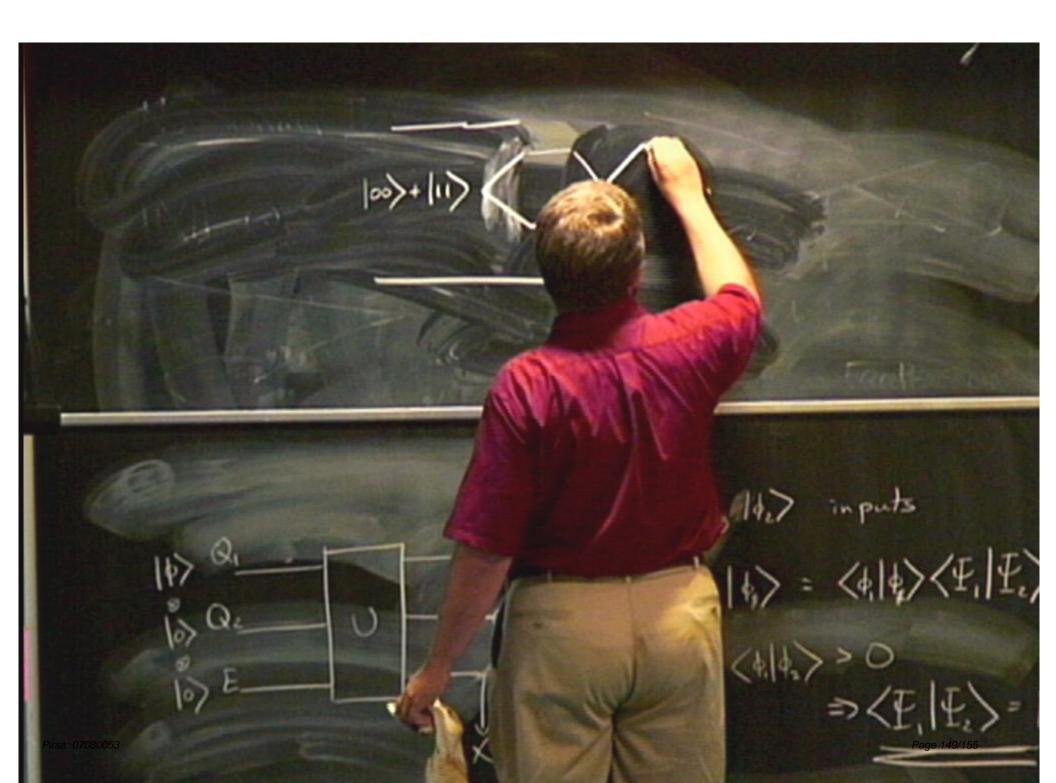
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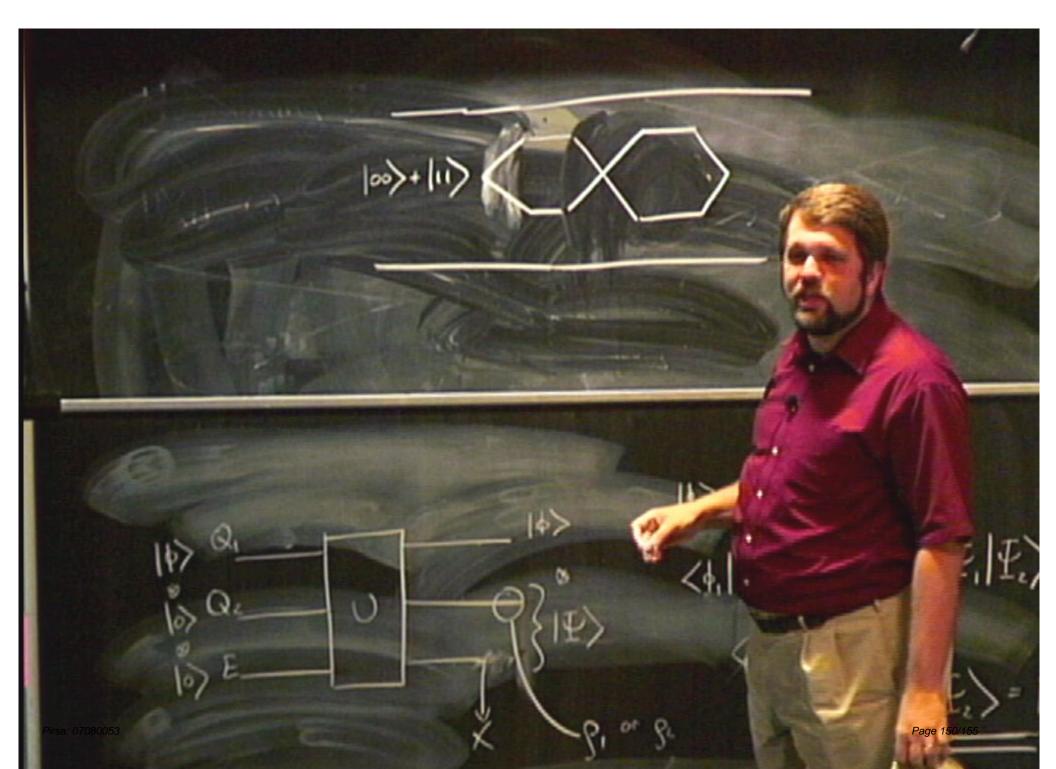
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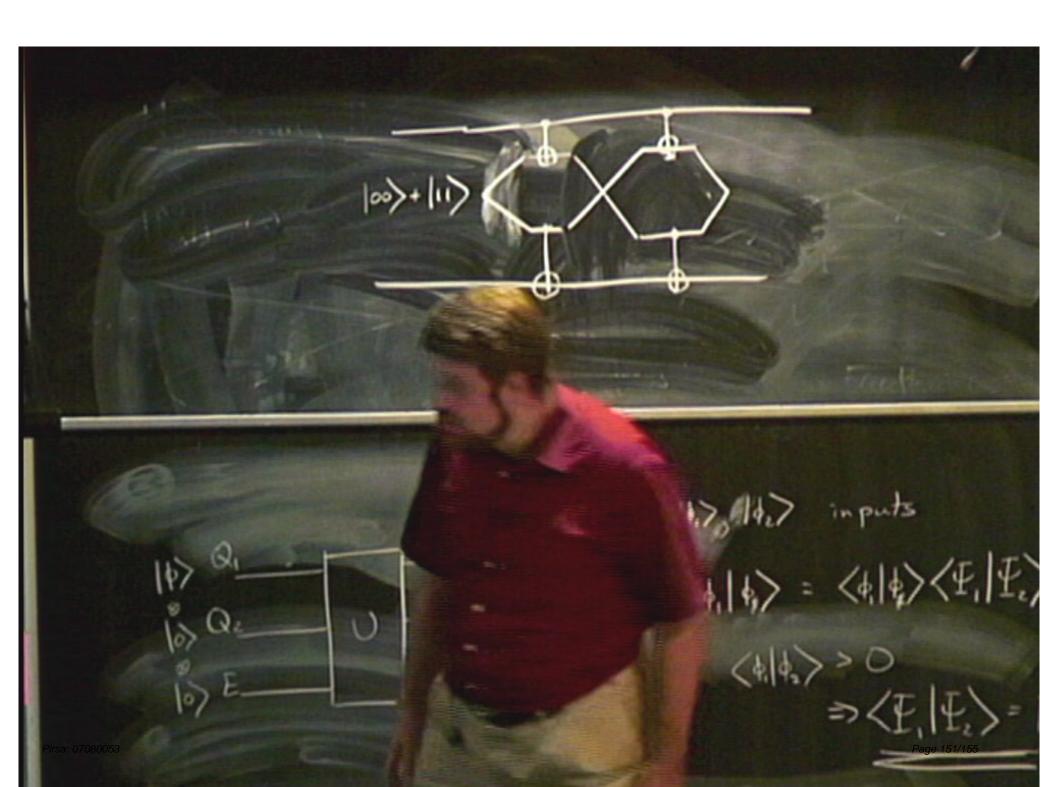
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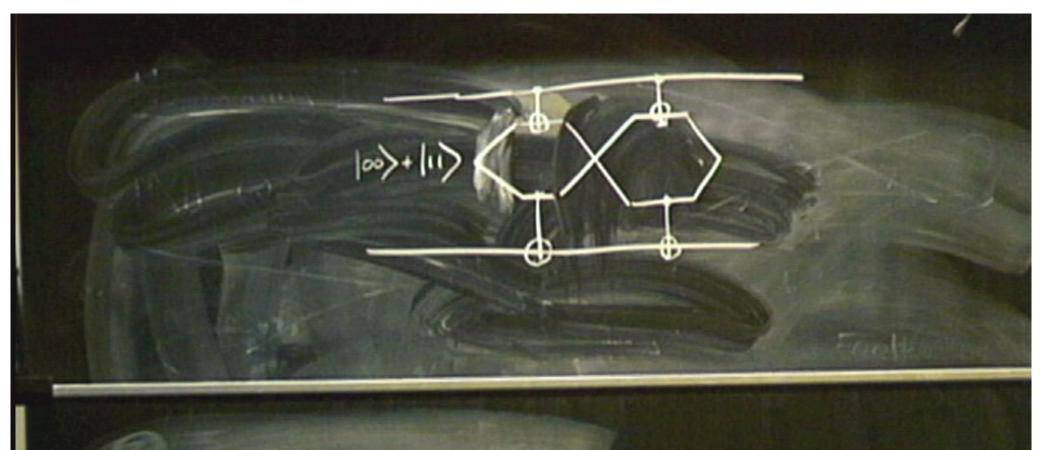


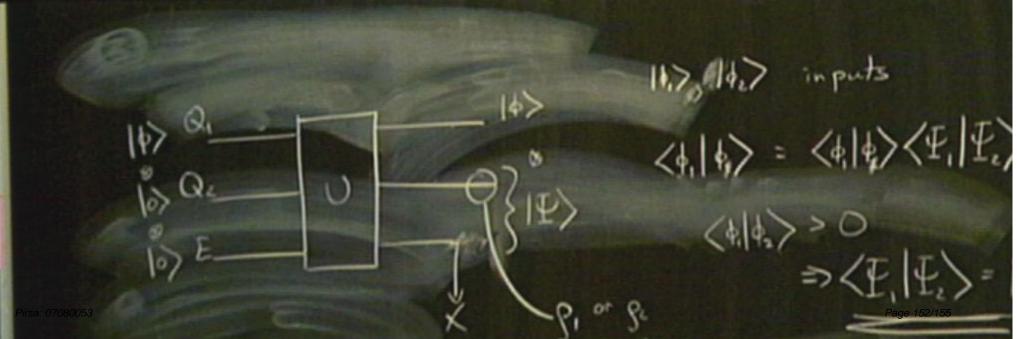


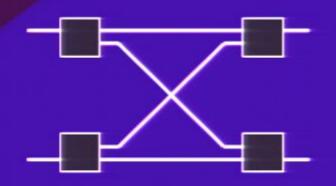








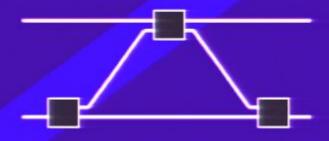


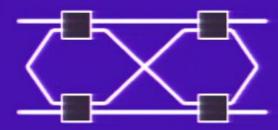


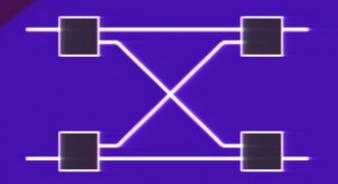
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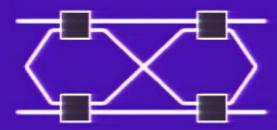


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What is the essential difference between these information flow patterns and simple information exchange?

References

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