Title: Self-Testing of Quantum Circuits

Date: Nov 15, 2006 02:00 PM

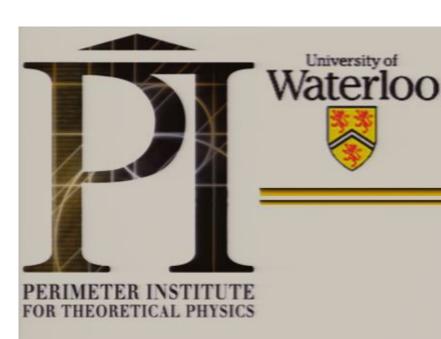
URL: http://pirsa.org/06110011

Abstract: I will explain how a quantum circuit together with measurement apparatuses and EPR sources can be fully verified without any reference to some other trusted set of quantum devices. Our main assumption is that the physical system we are working with consists of several identifiable sub-systems, on which we can apply some given gates locally.

To achieve our goal we define the notions of simulation and equivalence. The concept of simulation refers to producing the correct probabilities when measuring physical systems. The notion of equivalence is used to enable the efficient testing of the composition of quantum operations. Unlike simulation, which refers to measured quantities (i.e., probabilities of outcomes), equivalence relates mathematical objects like states, subspaces or gates.

Using these two concepts, we prove that if a system satisfies some simulation conditions, then it is equivalent to the one it is supposed to implement. In addition, with our formalism, we can show that these statements are robust, and the degree of robustness can be made explicit. Finally, we design a test for any quantum circuit whose complexity is linear in the number of gates and qubits, and polynomial in the required precision. Joint work with Frederic Magniez, Dominic Mayers and Harold Ollivier.

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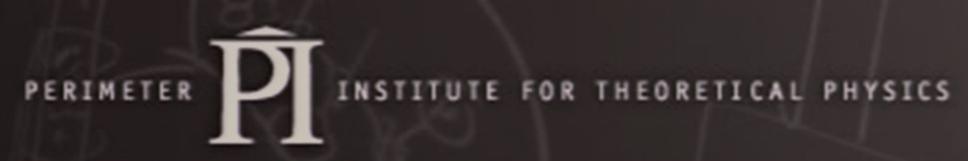
Self-Testing of Quantum Circuits

Michele Mosca

Canada Research Chair in Quantum Computation

Joint with F. Magniez, D. Mayers, H. Ollivier

Perimeter Institute Colloquium





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Why do we want to test?

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Why do we want to test?

What did we already know?

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What did we already know?

Current status.

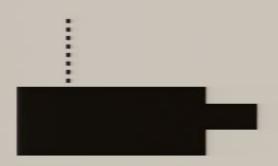
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Suppose you wish to buy a component for BB84 quantum cryptography, e.g. a source.

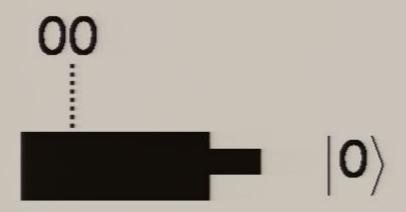
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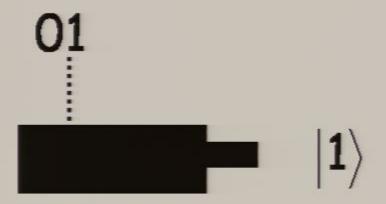
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Why should you trust this component?

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Why not?

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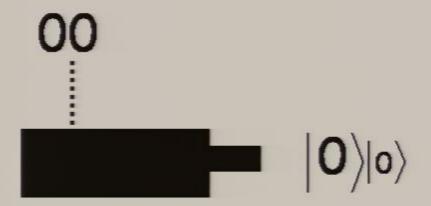
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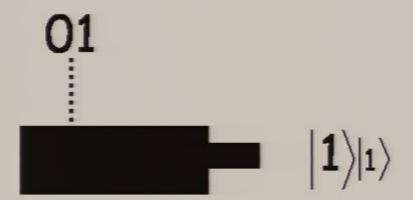
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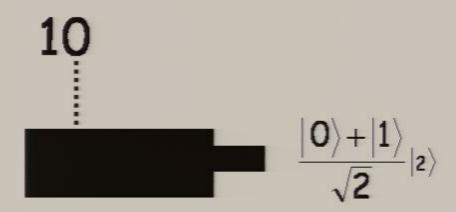
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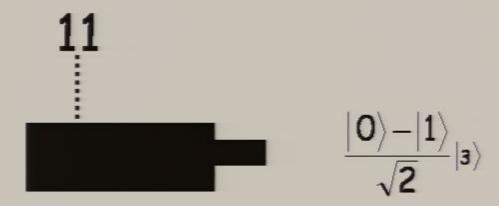
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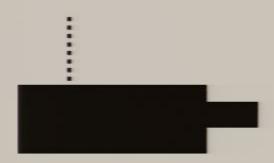


"side-channels"

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"side-channels"

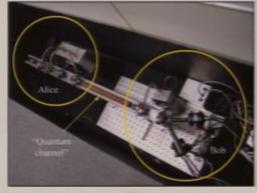
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http://www.research.ibm.com/journal/rd/481/smolin.html

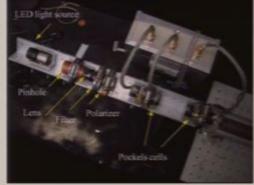
C. H. Bennett, F. Bessett, G. Brassard, L. Salvail, and J. Smolin, "Experimental Quantum Cryptography," J. Cryptol. 5, No. 1, 3–28 (1992).

Why not? What if what we really have is implethe following?

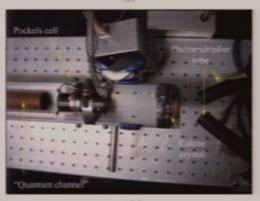








...



"side-channel

- 0

Figure 1

The apparatus used to perform the Page 30/154 cryptography experiment (a) The entire apparatus; (b) detailed view of Alice, (c) detailed view of Bob.



We need to make our assumptions and testing procedures explicit.

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We also don't want to rely on some other untrusted apparatus (e.g. in order to "just" do tomography).

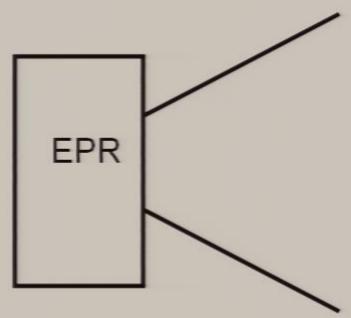
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Mayers and Yao devised a scheme for "self"testing **sources** for the purposes of QKD.

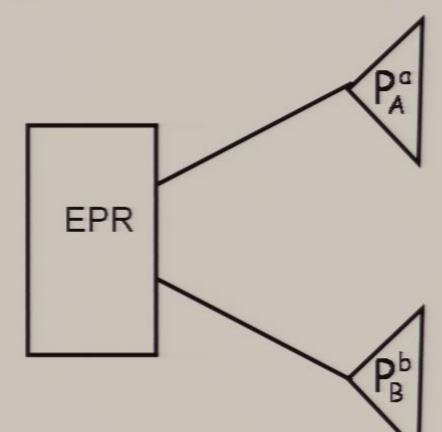
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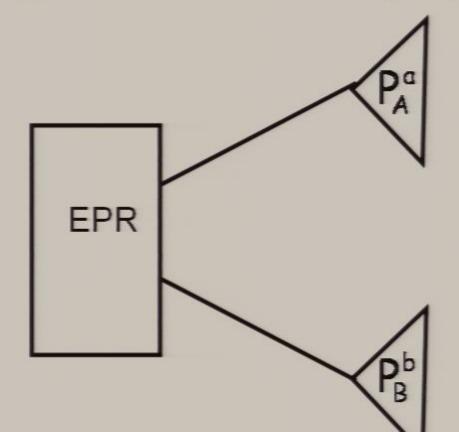
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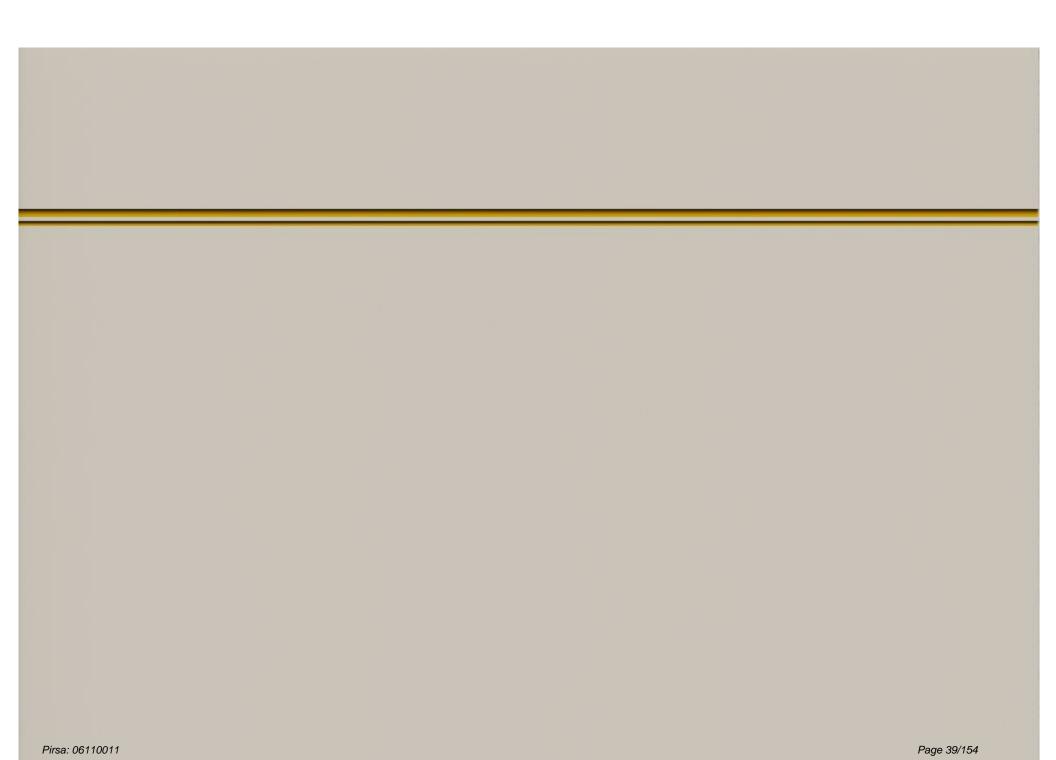


$$P^0 + P^{\pi/2} = I$$

$$P^{\pi/8} + P^{5\pi/8} = I$$

$$P^{\pi/4} + P^{3\pi/4} = I$$

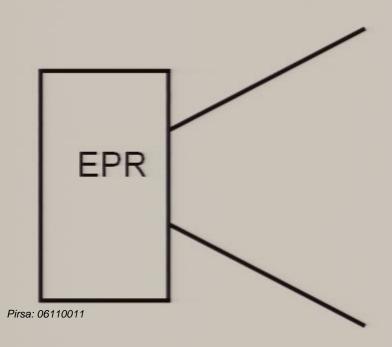
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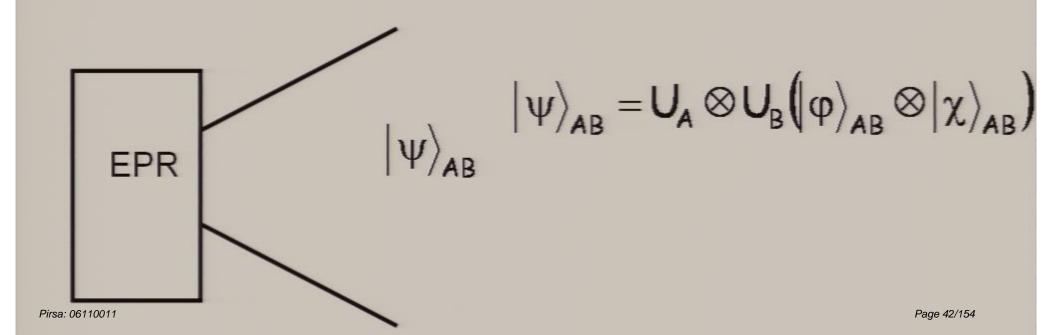
If the statistics are consistent with $|\phi\rangle = |00\rangle + |11\rangle$ then the output of the sources is locally unitarily equivalent to a state containing $|\phi\rangle$, and the projections are consistent with measuring the EPR pair.

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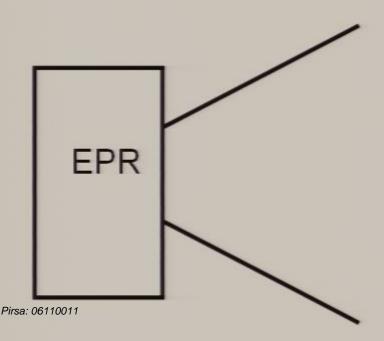


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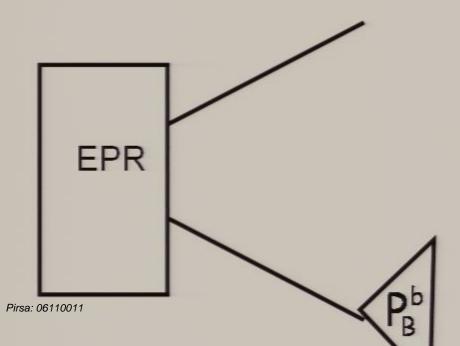


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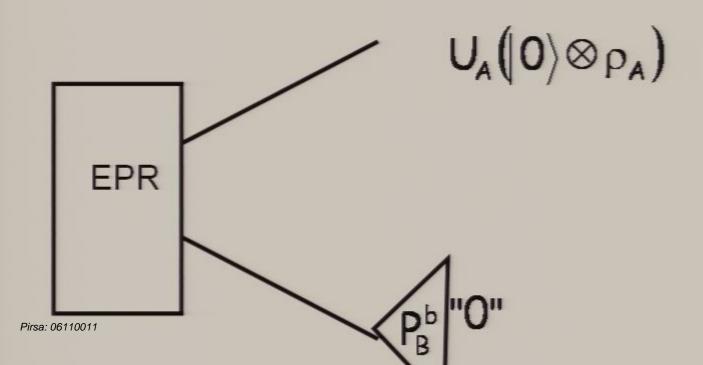
$$|\psi\rangle_{AB} = U_A \otimes U_B (|\phi\rangle_{AB} \otimes |\chi\rangle_{AB})$$



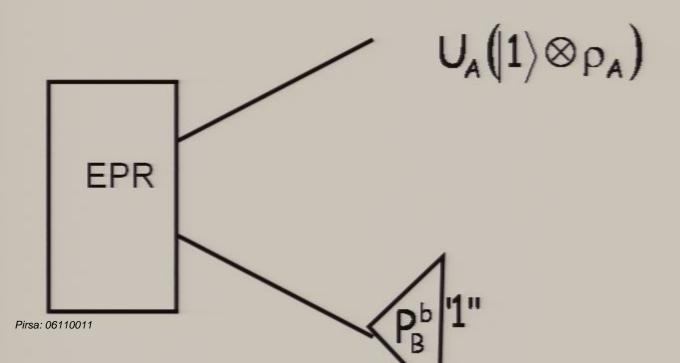
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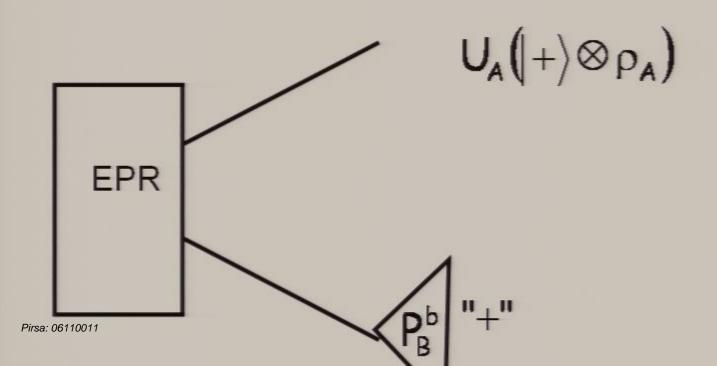
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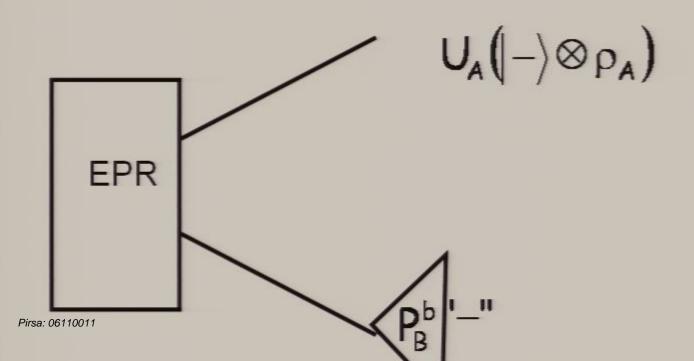
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The main assumptions of Mayers and Yao are

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- However, the results are not "robust". The results hold exactly if the statistics are satisfied exactly.
- Any realistic application will need to be robust.
- (Assuming robustness) This might be the only way,

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Suppose we are paying a lot of money to perform a large quantum computation, whose answer is not efficiently classically checkable.

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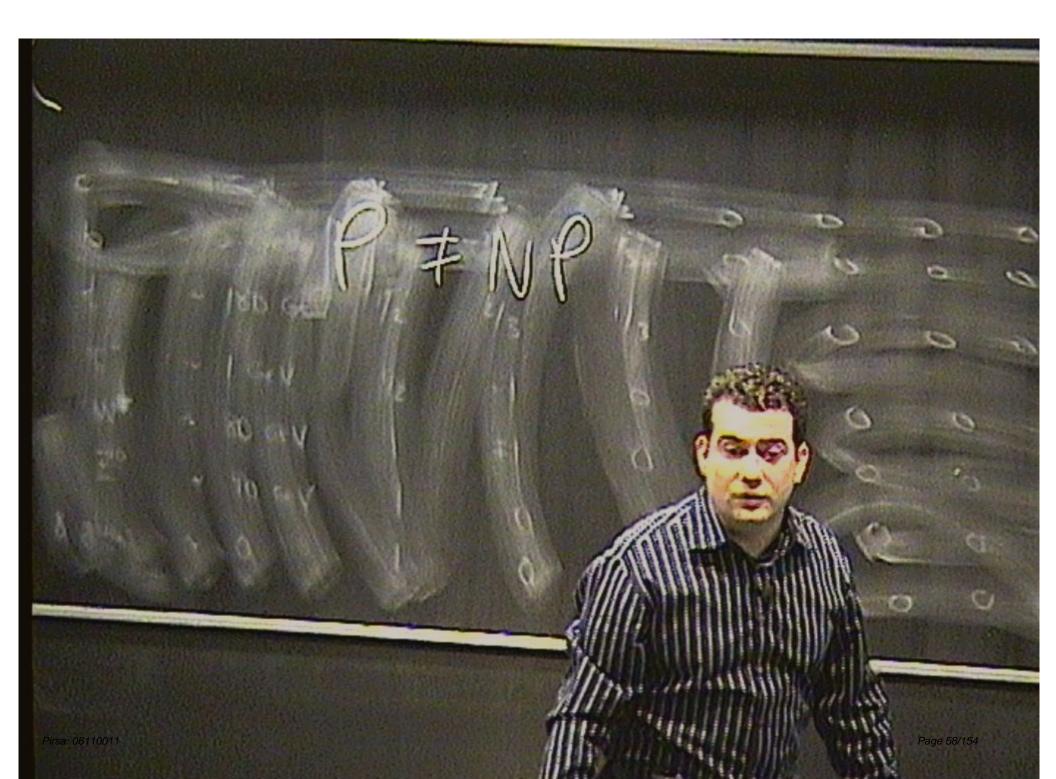
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Or, suppose we have proved one of the Clay Institute \$1M Millennium problem by a proof that needs to be run on a quantum computer.

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Van Dam, Magniez, M, Santha developed a series of self-tests for a universal and fault-tolerant set of quantum gates, with three additional assumptions:

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3) The ability to use the same gate more than once in the same experiment

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We wish to still have a "composable" technique for self-testing a large circuit; since we want it to be efficient.

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One step in that direction

Pirsa: 06110011 Page 69/154

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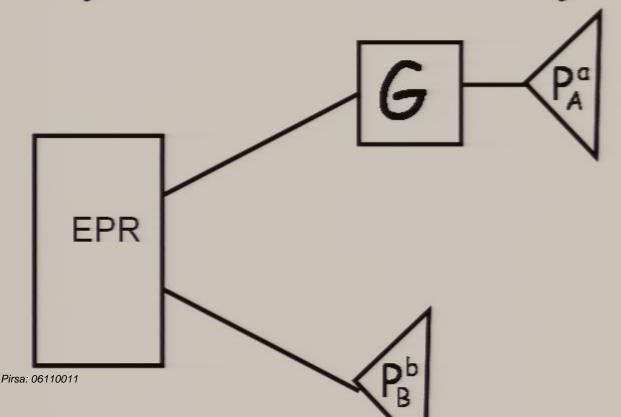
One step in that direction

We can combine the EPR self-test of Mayers-Yao with DMMS-style gate testing

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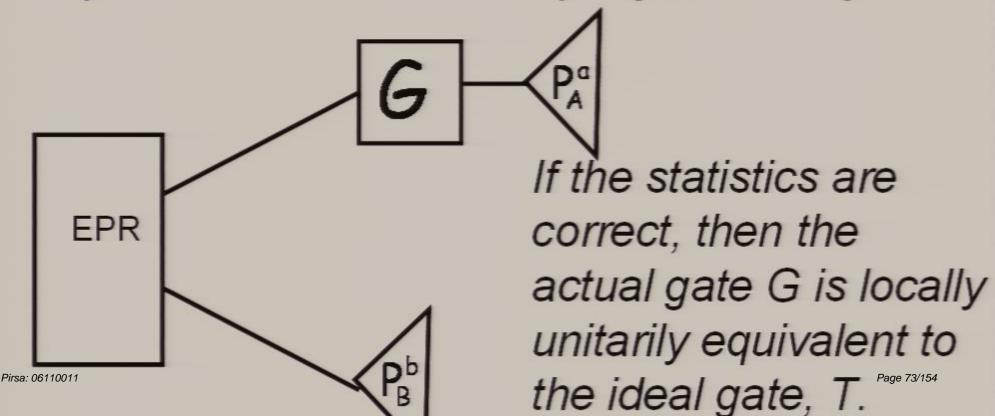
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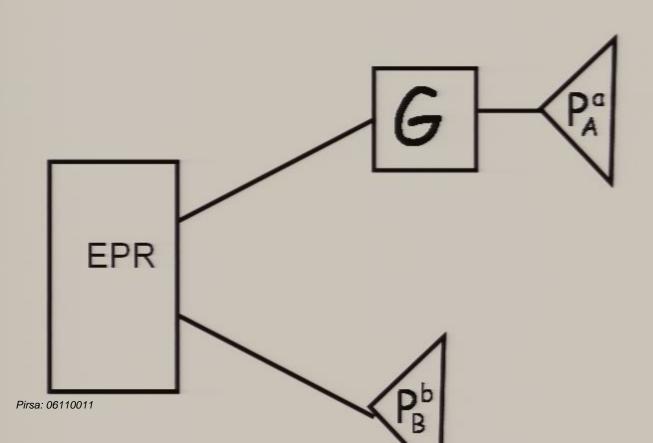


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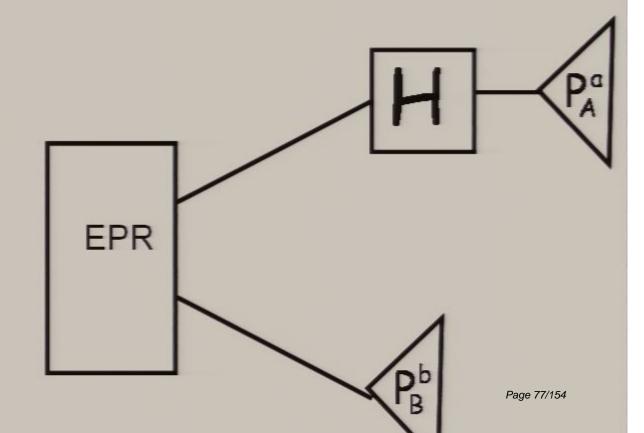
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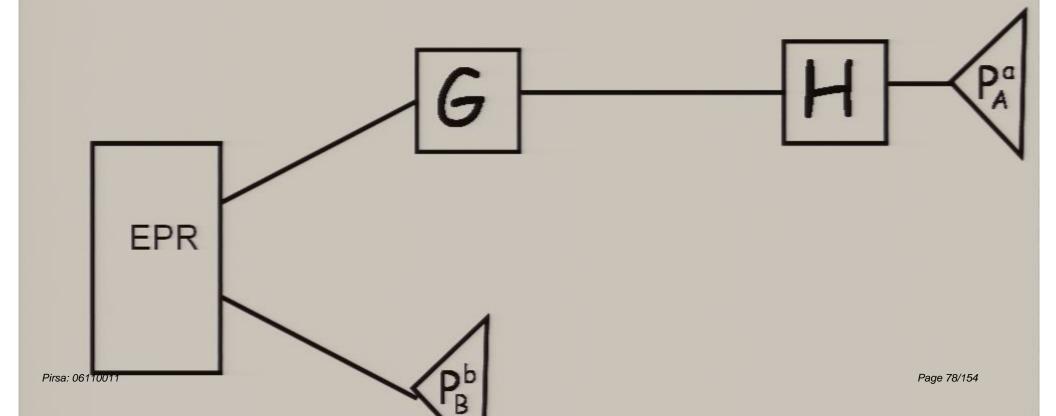


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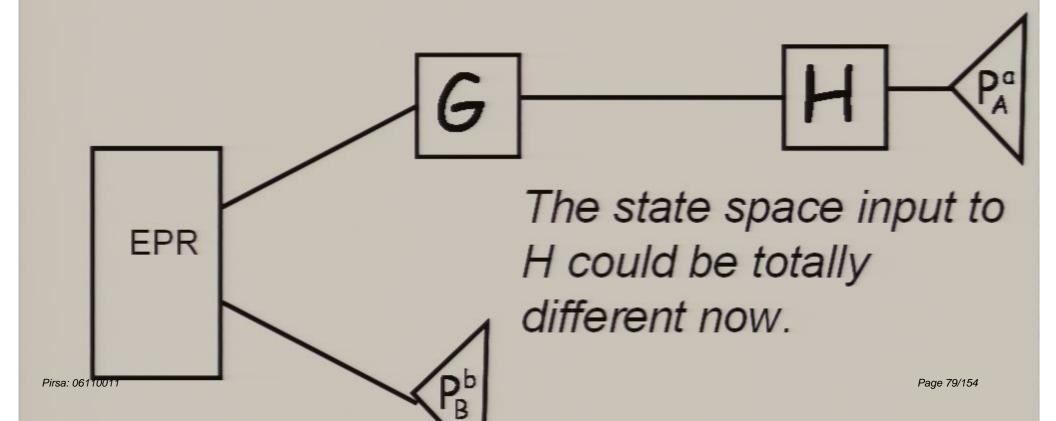


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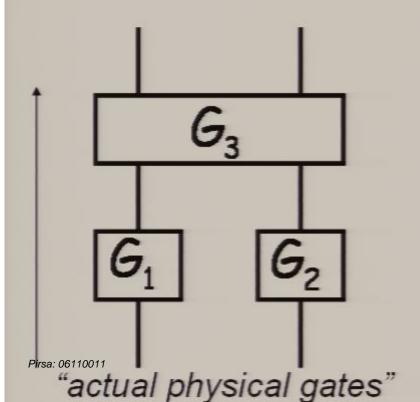
Goal

We ultimately wish to test the performance of an entire circuit (note that the circuits now flow up)

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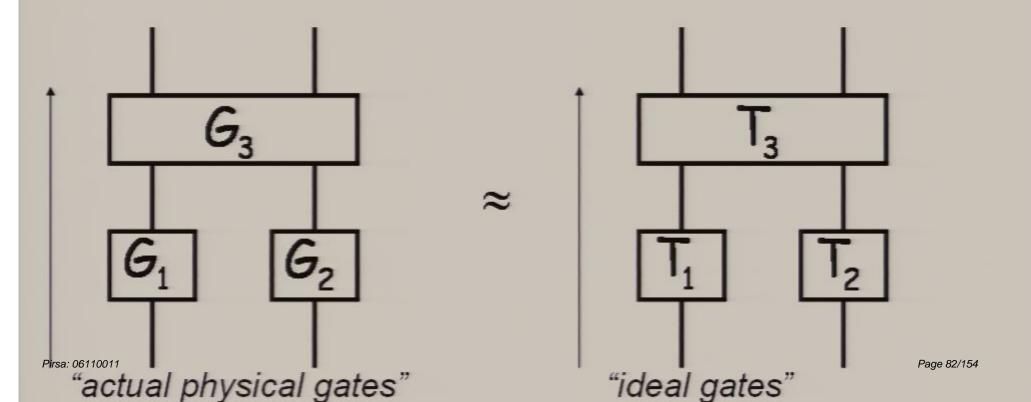
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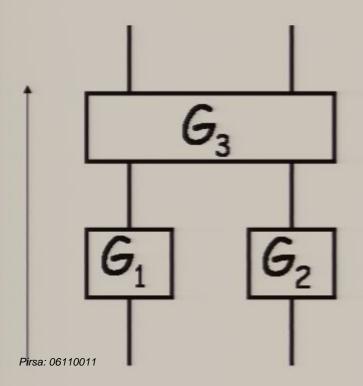
Goal

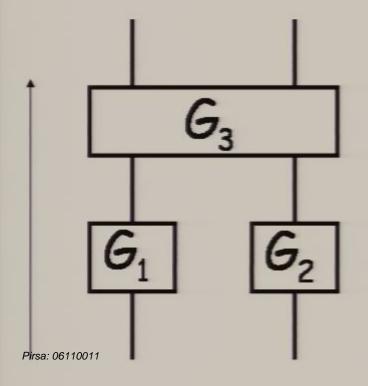
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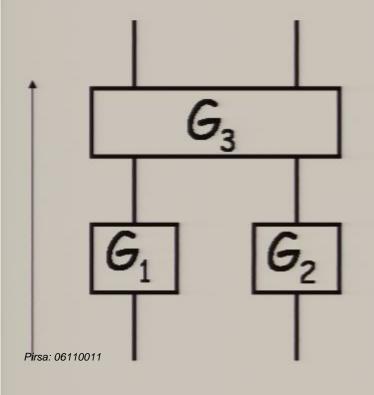
Suppose we wish to run the following circuit

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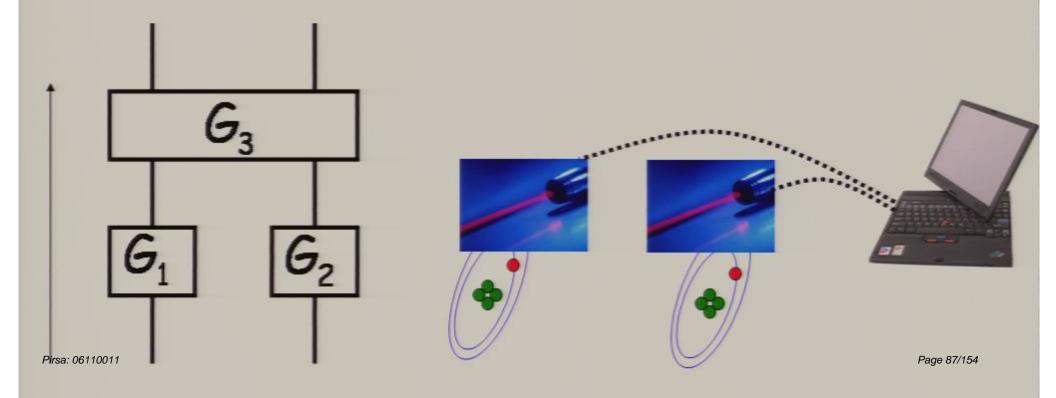


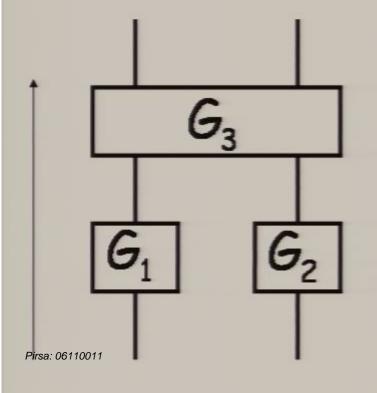








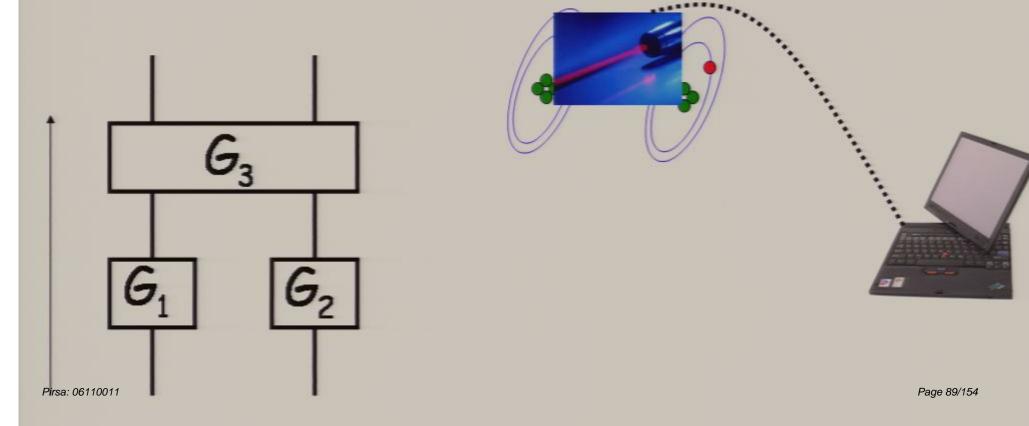


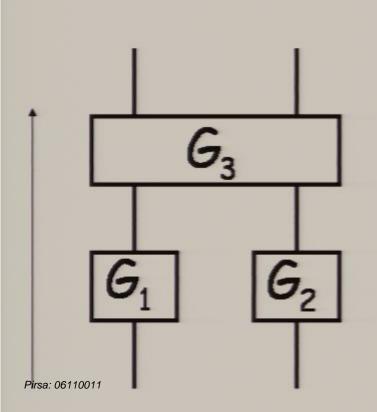








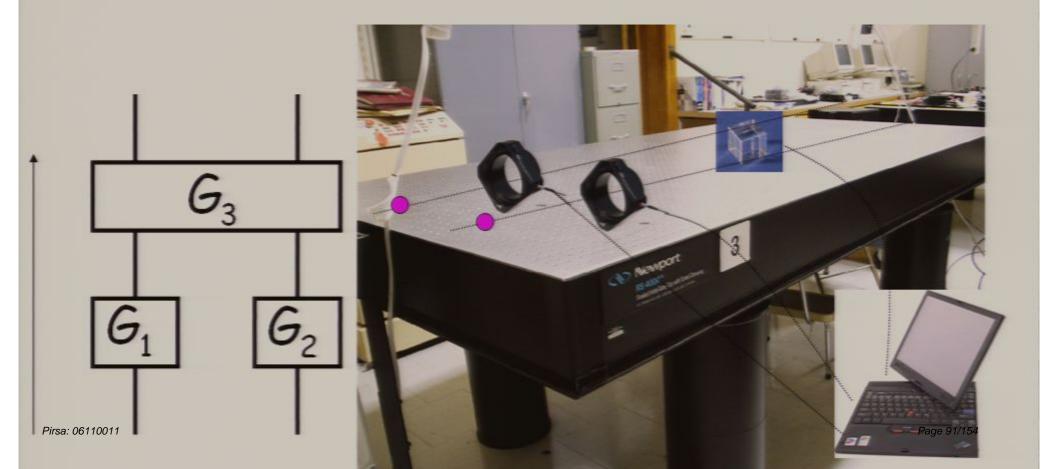








Example 2



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No finite set of tests will lead to a foolproof test. Why not?

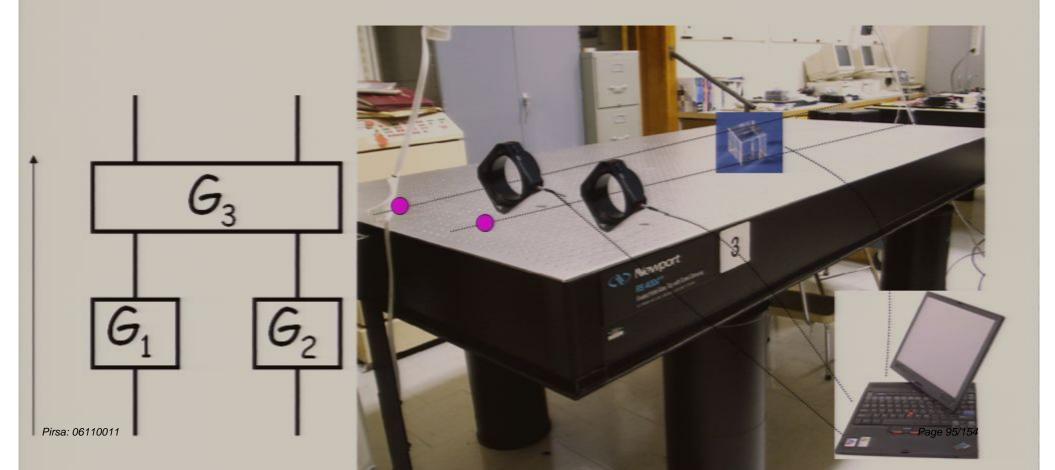
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No finite set of tests will lead to a foolproof test. Why not?

The gates can pass on the full (classical) history of their past to future gates in hidden degrees of freedom. Thus each gate knows the history of its input qubit(s), and can recognize when its history is no longer part of a test.

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Hint: Every circuit we would wish to run

needs to also be part of a test.

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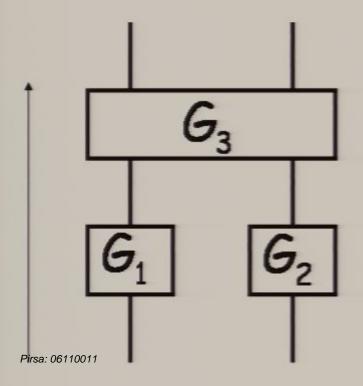
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Suppose we wish to run the following circuit

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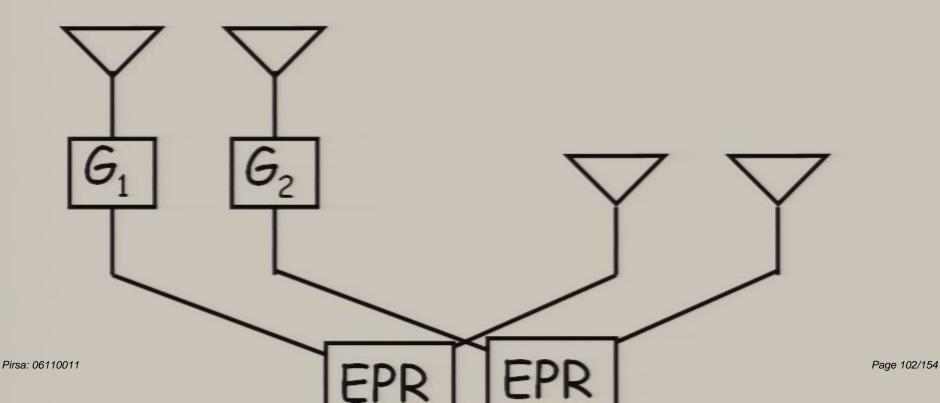




Verify the initial qubit sources.

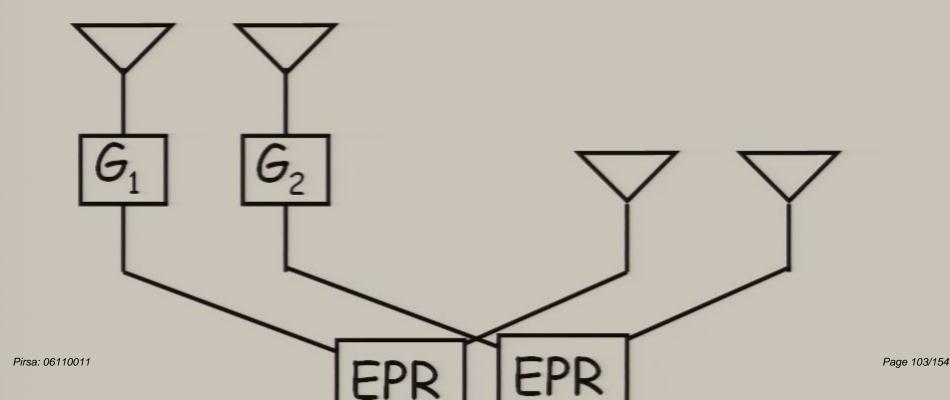


"tomography test"

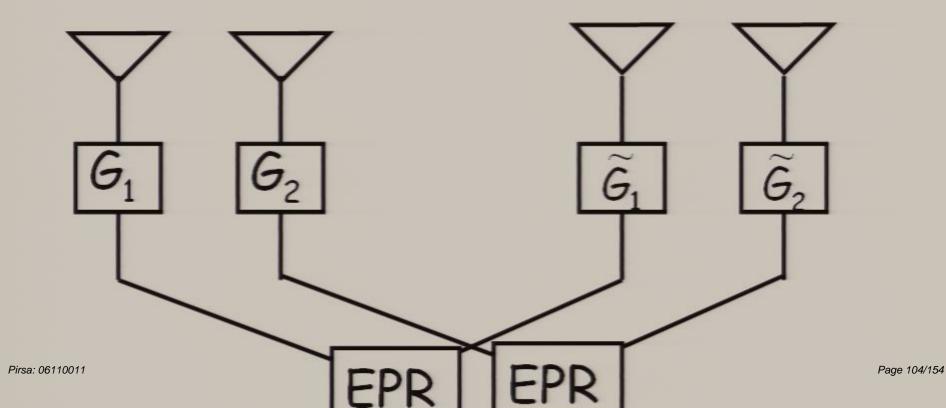


"tomography test"

We verify the behaviour of the gates G_1 and G_2

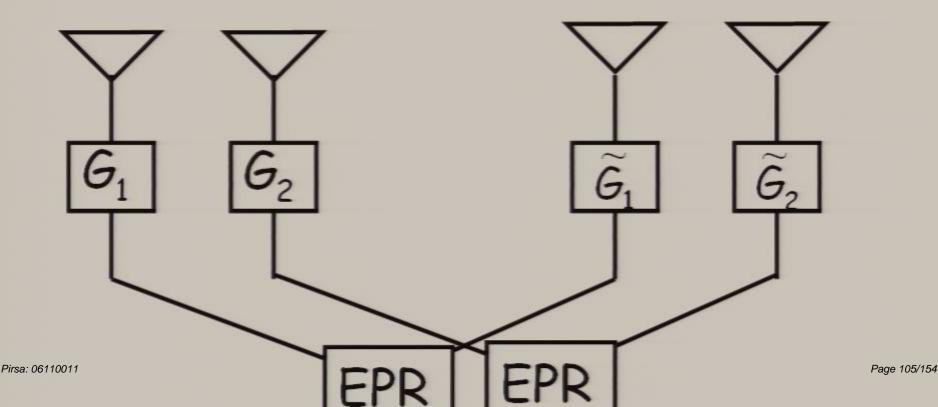


"conspiracy test"

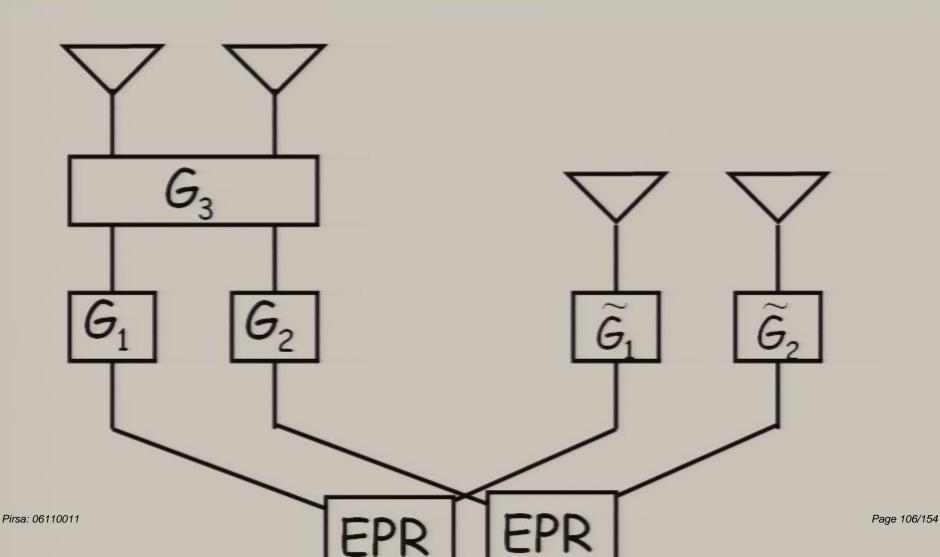


"conspiracy test"

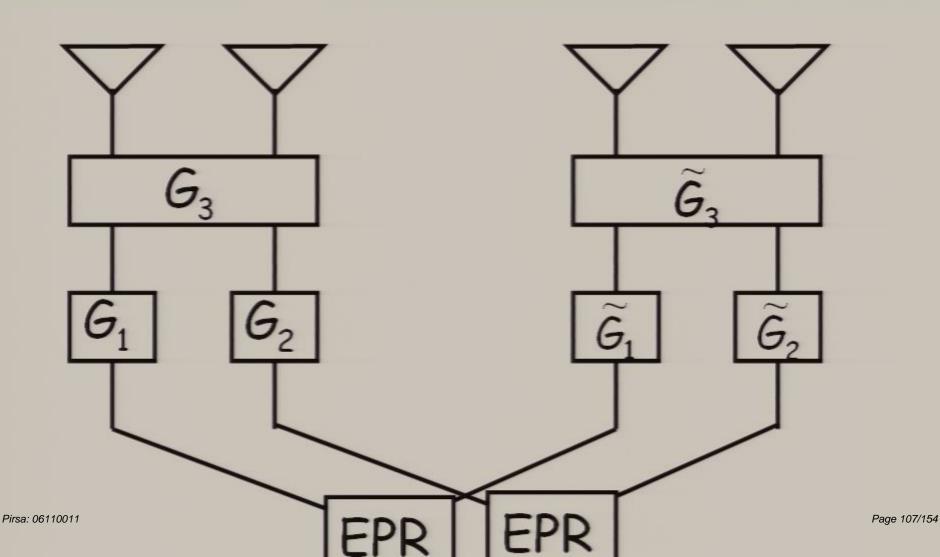
If $\widetilde{G}_1 = {G_1}^{\dagger}$, then this should recreate two EPR pairs.



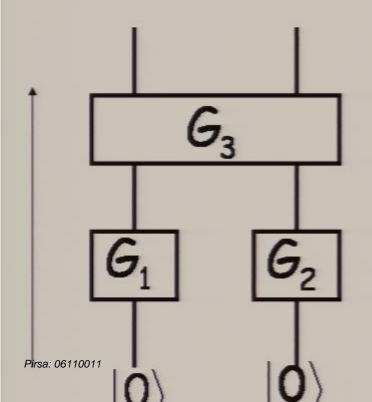
Tomography test



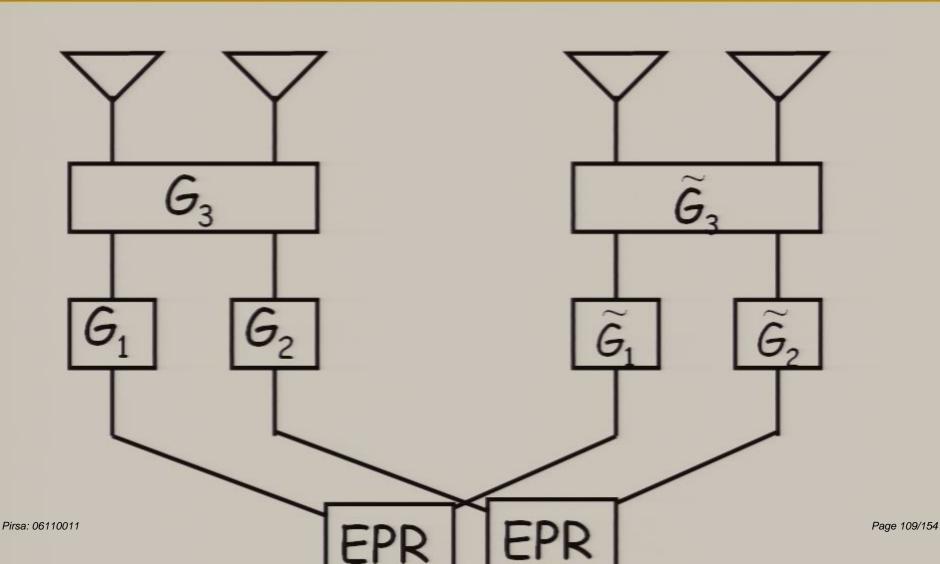
Conspiracy test



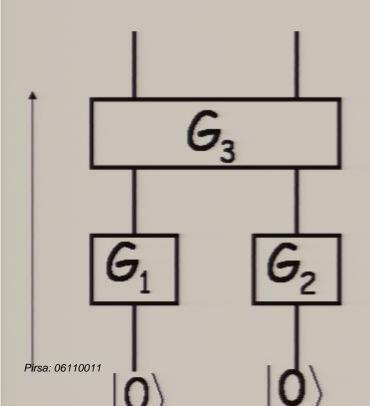
One last technical glitch



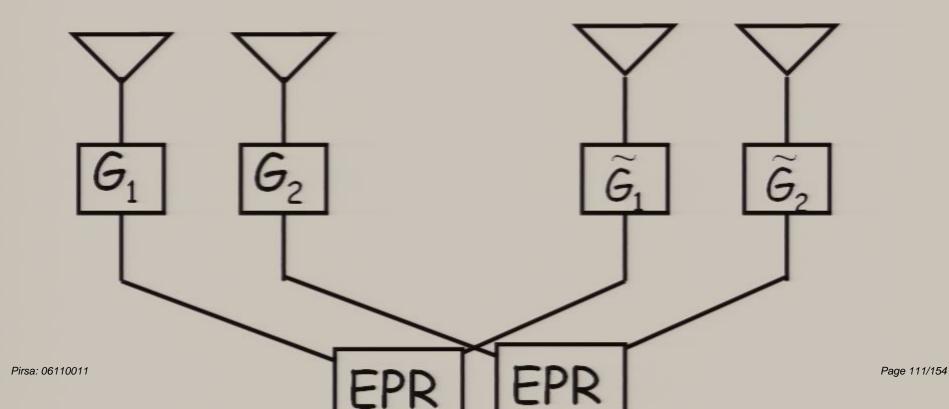
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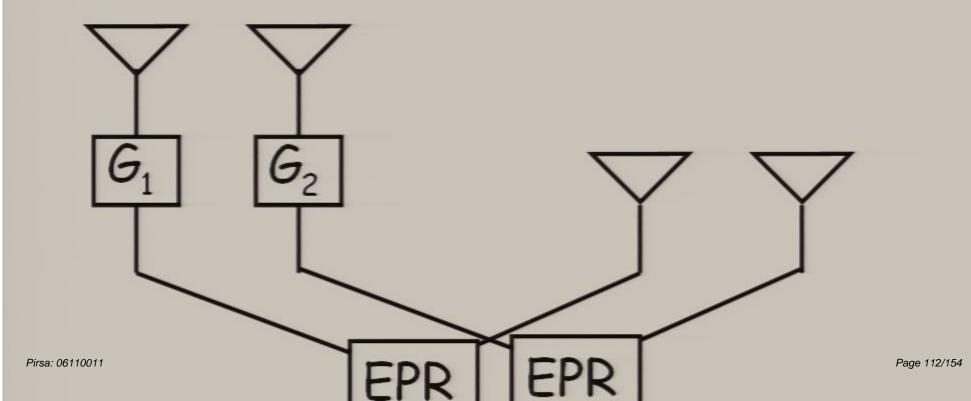


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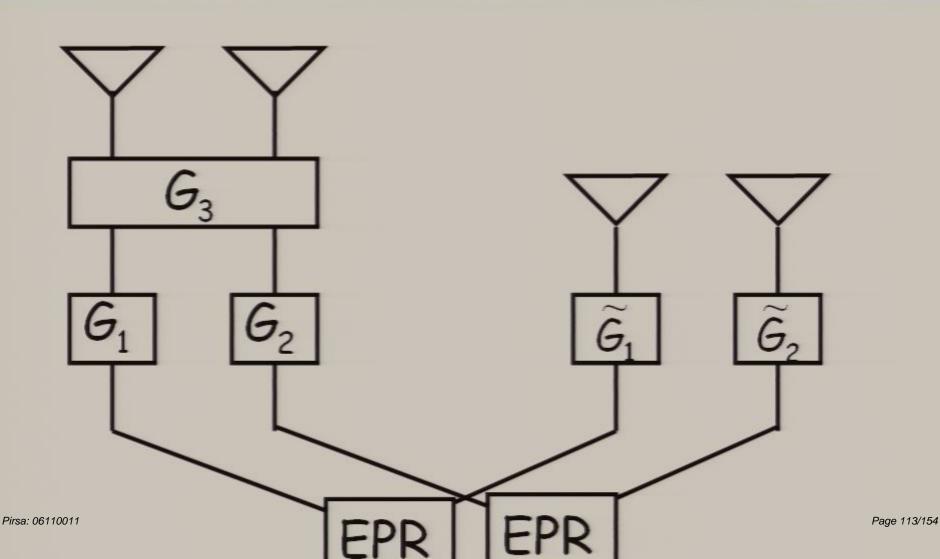


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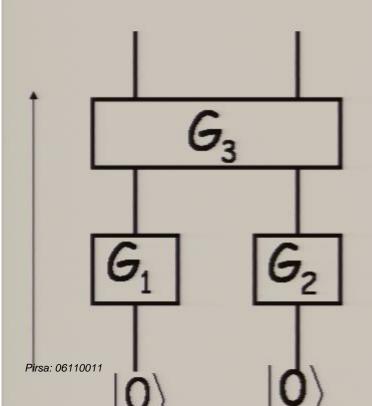
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Tomography test



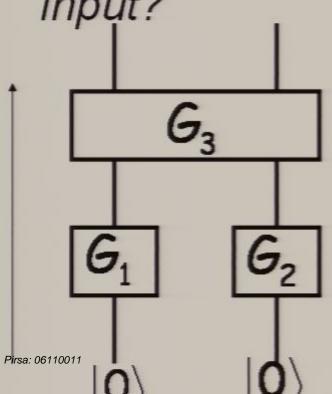
One last technical glitch



Preparing a specific input Pirsa: 06110011 Page 115/154

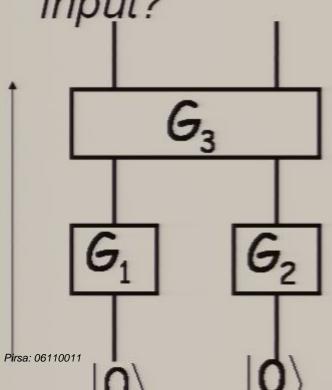
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This procedure verifies the behaviour of the circuit, but how do we run it on a specific input?



One last technical glitch

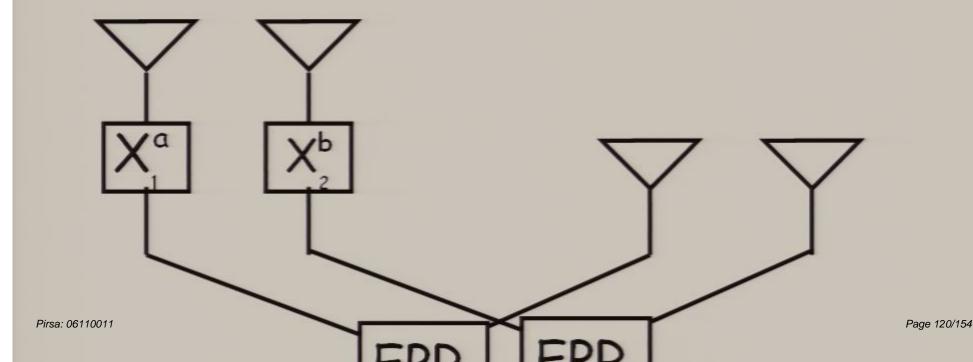
This procedure verifies the behaviour of the circuit, but how do we run it on a specific input?



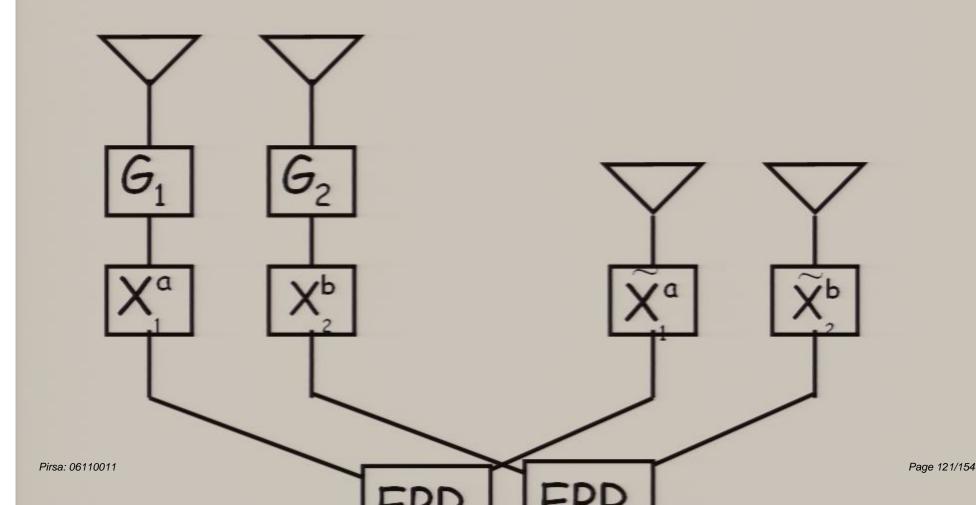
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Then proceed to test this modified circuit



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Then proceed to test this modified circuit Page 122/154 Pirsa: 06110011

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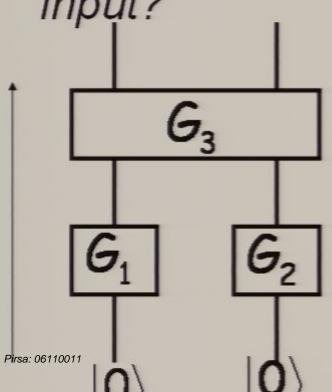
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We are merely saying that we do not have a procedure in the case of non-real coefficients.

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1 complex bit can be simulated by 2 real bits (see e.g. Rudolph and Grover quant-ph/0210187). But the two systems are not "equivalent" according to our notion of equivalence. E.g. inner products are not

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10)4(111)

10)4(11)

Let
$$T^1, T^2, \cdots, T^k \in U(2^n)$$
 (acting on a constant number of qubits each) $x \in \{0,1\}^n, \epsilon > 0, \gamma > 0$

If CircuitTest $(T^1, T^2, \cdots, T^k, x, \epsilon, \gamma)$ accepts, then with probability $1-O(\gamma)$ the outcome probability distribution of the circuit is at total variation distance $O((k+n)\epsilon^{1/8})$ from the distribution that comes from the measurement of $T^kT^{k-1}\cdots T^2T^1|x\rangle$ in the

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n complex tits

2 n+1 real Lits

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The main assumptions of Mayers and Yao are

- 1) Locality (i.e. measurements at A commute with those at B)
- Repeatability of experiments
- 3) Trusted classical apparatus
- However, the results are not "robust". The results hold exactly if the statistics are satisfied exactly.
- Any realistic application will need to be robust.
- (Assuming robustness) This might be the only way, was ing only these assumptions, to verifiably securely