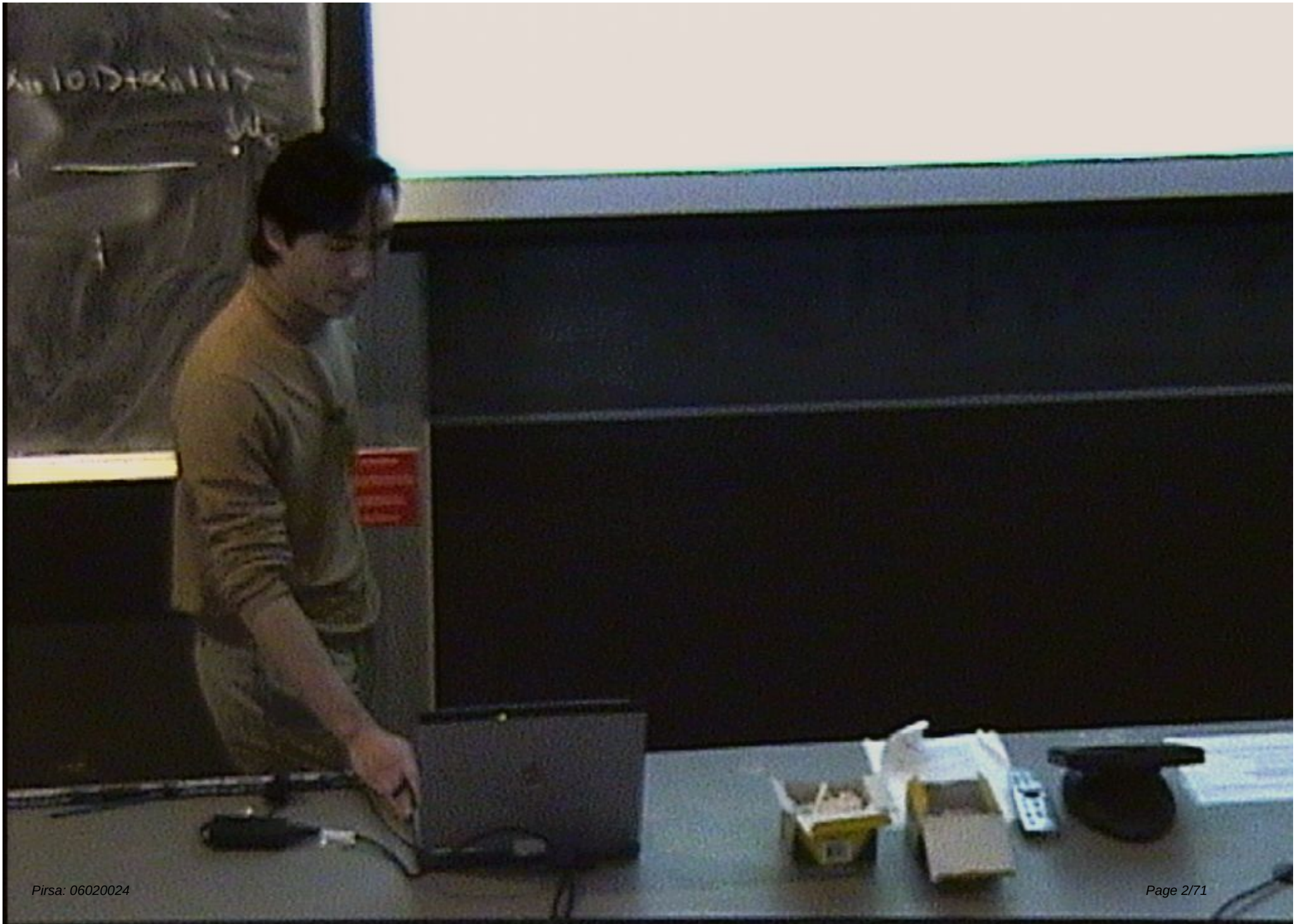


Title: Relation of adiabatic QC to other models

Date: Feb 11, 2006 10:45 AM

URL: <http://pirsa.org/06020024>

Abstract:



Workshop on Mathematical Aspects of the Quantum Adiabatic Approximation

Relation of Adiabatic Quantum Computing to Other Models

M. Stewart Siu, Stanford University

Reference:

1. Aharonov, Kempe, Regev, Van Dam, Landau & Lloyd, "Adiabatic Quantum Computation is Equivalent to Standard Quantum Computation", quant-ph/0405098
2. Kempe, Kitaev & Regev, "The Complexity of the Local Hamiltonian Problem", quant-ph/0406180
3. Oliveira & Terhal, "The Complexity of Quantum Spin Systems on a Two-Dimensional Square Lattice", quant-ph/0504050
4. Siu, "From Quantum Circuits to Adiabatic Algorithms", Phys. Rev. A 71, 062314 (2005), quant-ph/0409024

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Theme: Encoding a quantum circuit into a (smooth, time-dependent) Hamiltonian

OUTLINE

1. The Issue of Locality
2. Kitaev's History Approach (analog of Cook-Levin's proof of 3-SAT NP-completeness) and Its Adaption to AQC
3. Eigenvalue Gap and Stochastic Matrices
4. Connection to Holonomic Quantum Computing
5. Perturbation Theory Gadgets and Square Lattice Constructions

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1. The Issue of Locality

A straightforward way to map the circuit

Suppose $H|n\rangle = E_n|n\rangle \quad n = 0, 1, \dots$

$$(H + ?) U|n\rangle = E'_n U|n\rangle$$

Note that this implies

$$? U|n\rangle = [U, H]|n\rangle + (E'_n - E_n) U|n\rangle$$

which completely specifies "?"

Easy Case: $E'_n = E_n \Rightarrow H + ? = U H U^\dagger$

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We can use time-dependent
similarity transform

Easy Case: $E_n' = E_n \Rightarrow H + ? = U H U^\dagger$

We can use time-dependent
similarity transform

Given Quantum Gate U ,

Let $K = -i \log U$ and $\tilde{U}(t) = e^{iKt}$

$\tilde{U}(t) H \tilde{U}(t)^\dagger$ takes $|0\rangle$ to $U|0\rangle$ as $t \rightarrow 1$

Problem: If H is m -local and U is 2-local,
 $[U, H]$ can be $m+1$ -local.

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Problem : If H is m -local and U is 2-local,
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Can we avoid the problem

if we are willing to vary the gap?

Theorem (Haselgrove, Nielsen & Osborne):

$\forall$ density matrix ρ w/ eigenvalues $\rho_1 \leq \rho_2 \leq \dots$
 and $\text{tr}(\rho H) = \langle \psi | H | \psi \rangle$ 

$$\sum_{j=1}^{d-1} (E_j - E_0) \rho_{j+1} \leq (1 - |\langle \psi | E_0 \rangle|^2) E_{\text{tot}}$$

Suppose a n -qubit circuit generates
 $|\psi\rangle = \frac{1}{\sqrt{2}} (|000\dots\rangle + |111\dots\rangle)$ at some stage
 & we want $|\psi\rangle$ to be the ground state of
 k -local Hamiltonian H (where $k < n$),

$$\sum_{j=1}^n (L_j - L_0) \rho_{j+1} = \dots$$

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$$\text{tr}(\rho H) = \langle \psi | H | \psi \rangle \text{ since } k < n$$

$$\text{Inequality} \Rightarrow E_1 - E_0 \leq (1 - \langle \psi | E_0 \rangle^2) E_{\text{gap}}$$

K -local Hamiltonian $H = \sum_{i=1}^n H_i$ (where $K < n$),

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 $= 0$

$\Rightarrow$ Adiabatic evolution with local Hamiltonians cannot uniquely track evolution of the circuit.

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$\Rightarrow$ Adiabatic evolution with local Hamiltonians cannot uniquely track evolution of the circuit.

What can we do?

- 1) Track approximately
- 2) Add ancillary qubits
- 3) Allow degeneracy

2. The History Approach

Technique not developed for Adiabatic QC,
but for QMA-completeness.

Def. of QMA: $L \in \text{QMA}$ if $\exists$ Verifier circuit $U \ni$.

$$\forall x \in L \exists |y\rangle \ni P(U_x | y, |1\rangle) \geq \frac{2}{3} \quad \text{hand icon}$$

$$\forall x \notin L \forall |y\rangle \ni P(U_x | y, |1\rangle) \leq \frac{1}{3}$$

Def. of Local Hamiltonian Problem:

$$H = \sum \text{k-local hermitian operators on n-particles}$$

Is the smallest eigenvalue $> b$ or $< a$,
where $b - a = \text{poly}(n)^{-1}$

Like 3-SAT, the Local Hamiltonian is a set of constraints.



To prove any QMA problems can be written as LHP, follow Cook-Levin and constrain the history of verifier:

$$H = \left(\sum_{\text{all time step}} H_{\text{prop}} \right) + H_{\text{out}}$$

(For convenience, all terms are semi +ve definite)

$\uparrow$ $\uparrow$ $\uparrow$

$= 0$ iff each gate of verifier U is correctly applied

$= 0$ iff final state $= |1\rangle$

Initial state is left unconstrained. we don't know if $\exists |y\rangle \ni U|y\rangle \sim |1\rangle \otimes \dots$

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Use this for Adiabatic QC:

$$\hookrightarrow H = \left(\sum_{\text{all time step}} H_{\text{prop}} \right) + \cancel{H_{\text{out}}} + H_{\text{in}}$$

- Verifier circuit is replaced by the circuit we are mapping from.
- Instead of constraining the output, we constrain the input.
- The ground state is always the history of the circuit (Unless we keep H_{out} , in which case we'd have to measure the energy)

Let $|x_0\rangle \dots |x_L\rangle$ be the intermediate states of the quantum circuit. Make them orthogonal by adding clock ancillae.

The history state is

$$\frac{1}{\sqrt{L+1}} (|x_0\rangle + \dots + |x_L\rangle)$$

What H would have an equal superposition as ground state?

Observe that $\frac{1}{2} \begin{bmatrix} 1 & -1 \\ -1 & 1 \end{bmatrix} \begin{bmatrix} 1 \\ 1 \end{bmatrix} = 0$



we can stack these up:

$$\frac{1}{2} \begin{bmatrix} 1 & -1 & & \\ -1 & 2 & -1 & \\ & -1 & 2 & -1 \\ & & -1 & 1 \end{bmatrix} \begin{bmatrix} 1 \\ 1 \\ 1 \\ 1 \end{bmatrix} = 0$$

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$$\begin{bmatrix} 2 & -1 & & \\ & \ddots & \ddots & \\ & & 2 & -1 \\ & & -1 & 1 \end{bmatrix}$$

$$= \frac{1}{2} \sum_t |x_t\rangle \langle x_t| + |x_{t+1}\rangle \langle x_{t+1}| - |x_t\rangle \langle x_{t+1}| - |x_{t+1}\rangle \langle x_t|$$

etc.

★ The key to making the Hamiltonian k -local, where k is small, is finding good k -local approx. to the terms $|x_{t+1}\rangle \langle x_t|$ etc.

unary clock:

$$|t\rangle^{\text{clock}} = | \underbrace{111 \dots}_t 00 \dots \rangle$$

3-local example:

$$|x_{t+1}\rangle \langle x_t| \approx U_{t+1}^{\otimes} |1\rangle_{t+1}^{\text{clock}} \langle 0|_{t+1}^{\text{clock}} + J |01\rangle_{t,t+1}^{\text{clock}} \langle 01|_{t,t+1}^{\text{clock}}$$

for large J

The terms $|0\rangle_{t+1} \langle 0|_{t+1}$ etc.

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$$|\gamma_{t+1}\rangle \langle \gamma_t| \approx U_{t+1}^{\otimes 3} |1\rangle_{t+1}^{\text{clock}} \langle 0|_{t+1}^{\text{clock}} + J |0\rangle_{t,t+1}^{\text{clock}} \langle 0|_{t,t+1}^{\text{clock}}$$

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What is the gap of $H_{\text{in}} + \sum H_{\text{prop}}$?

Both H_{in} and $\sum H_{\text{prop}}$ have very simple spectra
if both have gaps $\geq \lambda$, then

$$\begin{aligned} &\text{gap of } H_{\text{in}} + \sum H_{\text{prop}} \\ &\leq \lambda \max \{ |\langle v | w \rangle|^2 \mid \begin{array}{l} v \in \text{kernel of } H_{\text{in}}, \\ w \in \text{kernel of } \sum H_{\text{prop}} \end{array} \} \end{aligned}$$

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$$\begin{bmatrix} 1 & & & \\ & -1 & & \\ & & 1 & \\ & & & -1 \end{bmatrix}$$

So $\sum H_{\text{prop}}$ in the basis $|\gamma_0\rangle, \dots, |\gamma_L\rangle$ looks like

$$\frac{1}{2} \begin{bmatrix} 1 & -1 & & \\ -1 & 2 & & \\ & & \ddots & \\ & & & 2 & -1 \\ & & & -1 & 1 \end{bmatrix}$$

$$= \frac{1}{2} \sum_t |\gamma_t\rangle \langle \gamma_t| + |\gamma_{t+1}\rangle \langle \gamma_{t+1}| - |\gamma_t\rangle \langle \gamma_{t+1}| - |\gamma_{t+1}\rangle \langle \gamma_t| \text{ etc.}$$

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$$\text{Combined effect} = \sum_{t \neq 0} |\gamma_t\rangle \langle \gamma_t|$$

Adiabatic Algorithm:

$$H_{\text{initial}} = H_{\text{in}} + \underbrace{|\gamma_{\text{clock}=0}\rangle \langle \gamma_{\text{clock}=0}|}_{\text{clock}} + H_{\text{clock}}$$

↑ make sure
clock states are
legal
↓

$$H_{\text{final}} = H_{\text{in}} + \sum H_{\text{prop}} + H_{\text{clock}}$$

In the basis $\{|\gamma_0\rangle \dots |\gamma_L\rangle\}$, this looks like

$$H_{\text{initial}} = \begin{bmatrix} 0 & 1 & & \\ & 1 & & \\ & & \ddots & \\ & & & 1 \end{bmatrix}$$

$$H_{\text{final}} = \frac{1}{2} \begin{bmatrix} 1 & -1 & & \\ -1 & 2 & & \\ & & \ddots & \\ & & & 2 & -1 \\ & & & -1 & 1 \end{bmatrix}$$

3. Eigenvalue Gap and Random Walk

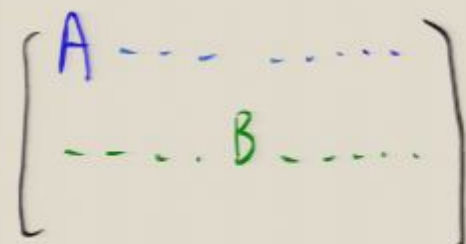
(cf. Spitzer's talk)

What is the gap of $H(s) = (1-s)H_{\text{initial}} + sH_{\text{final}}$?

Rough Idea:

Case I, $s \leq \frac{1}{3}$

Circle theorem:



$$H_{ii} - \sum_{j \neq i} |H_{ij}| \leq \text{eigenvalue}_{E_i} \leq H_{ii} + \sum_{j \neq i} |H_{ij}|$$

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$$\begin{bmatrix} A & \dots & \dots \\ \dots & B & \dots \end{bmatrix} \quad H_{ii} - \sum_{j \neq i} |H_{ij}| \leq \text{eigenvalue } E_i \leq H_{ii} + \sum_{j \neq i} |H_{ij}|$$

Since $H(s)$ close to $H_{\text{int}} = \begin{bmatrix} 0 & \dots & \dots \end{bmatrix}$

$$\begin{array}{l} A + \dots \leq \frac{1}{3} \\ B - \dots \geq \frac{2}{3} \end{array} \Rightarrow \text{gap} \geq \frac{1}{3}$$

Case II, $S > \frac{1}{3}$ (sketch only)

We want to use the property that the gap of stochastic matrix (positive entries, sum of rows/column = 1) can be related to its conductance (expansion parameter).

Define $G = 4I - H(S)$ and $P_{ij} = \frac{\alpha_j}{\mu \alpha_i} G_{ij}$, then P is stochastic.

Random walk on the line $x_0 \dots x_L$



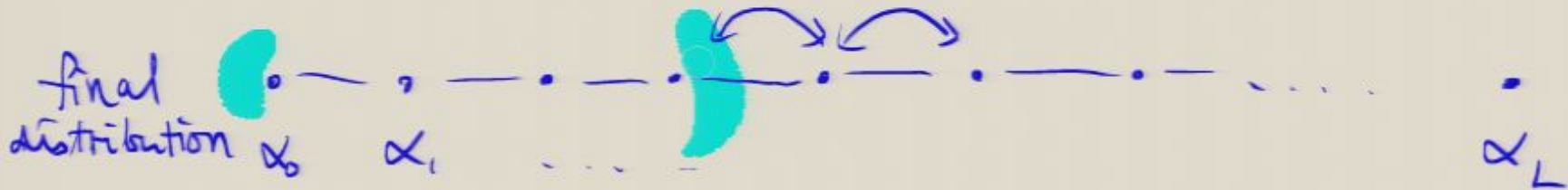
Limiting distribution is monotonic: $\alpha_0 \geq \alpha_1 \geq \dots \geq \alpha_L$

$$\text{Gap} \sim \text{conductance} = \min_B \left(\sum_{i \in B} \sum_{j \notin B} \pi_i P_{ij} \right) / \left(\sum_{i \in B} \pi_i \right) \leq \frac{1}{2}$$

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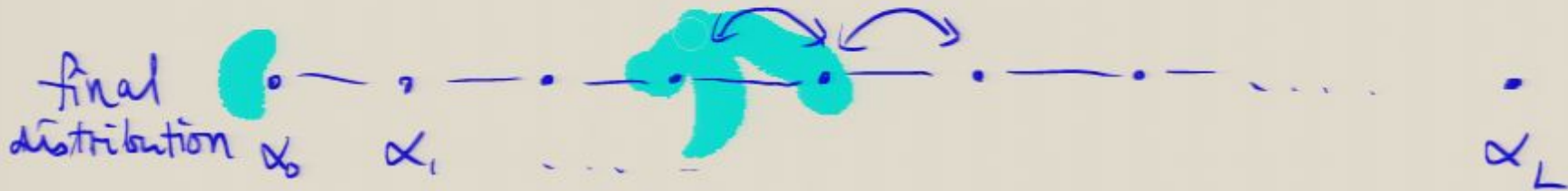
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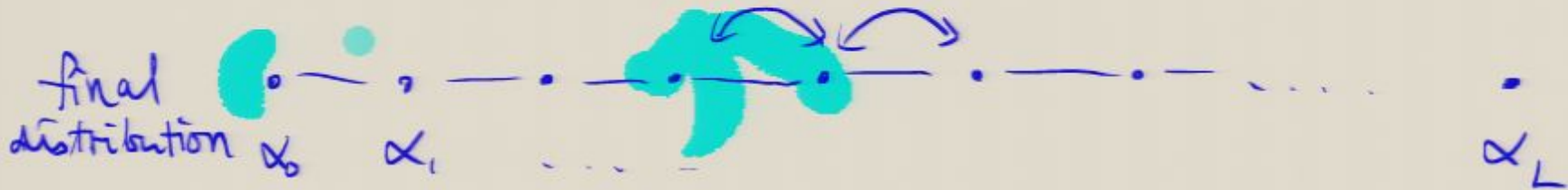
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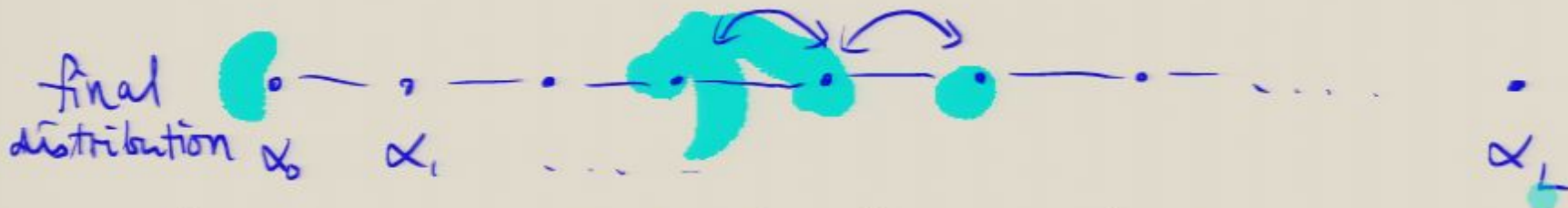
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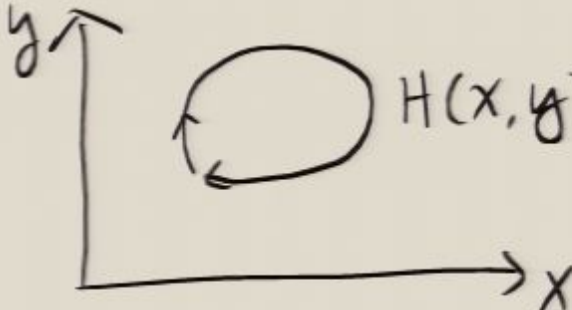
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4. Connection to Holonomic QC

$H/QC =$  $H(x, y)$ moves adiabatically

Non-abelian Berry phase to implement U .

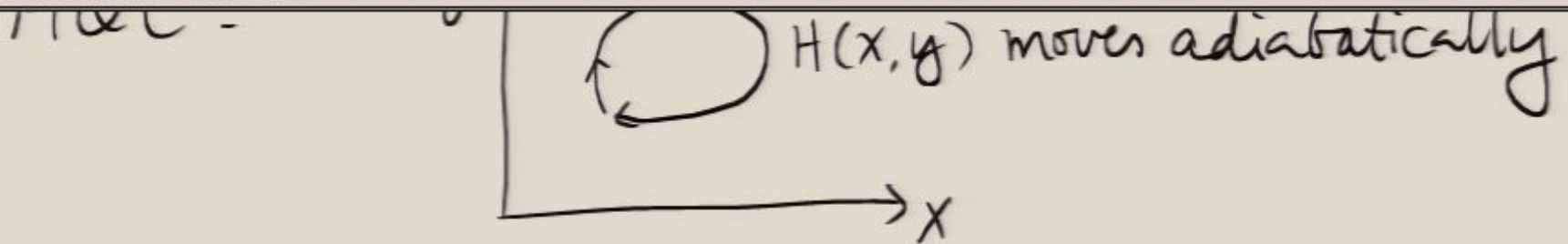
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What does it have to do with the history state?

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But we could have started with any H_{in} .



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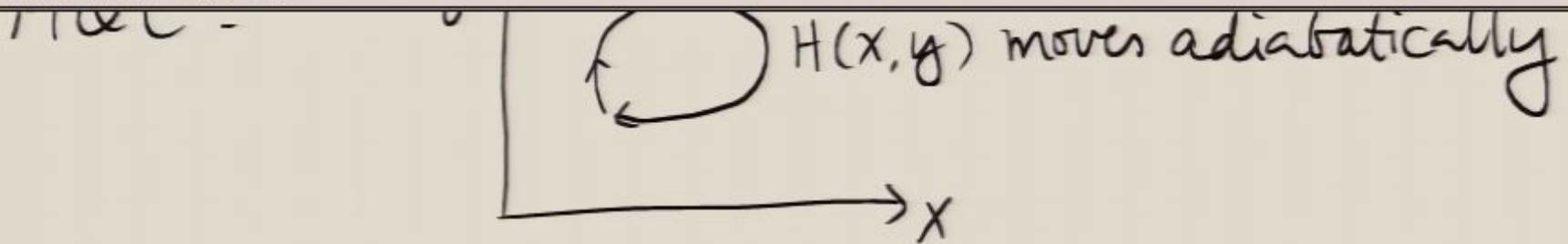
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Consider a one gate circuit mapped using the history approach. The initial and final states are:

$$|\psi_i\rangle = |0\rangle \otimes |0\rangle^{\text{clock}}$$

$$|\psi_f\rangle = \frac{1}{\sqrt{2}} (|0\rangle \otimes |0\rangle^{\text{clock}} + U|0\rangle \otimes |1\rangle^{\text{clock}})$$

The corresponding Hamiltonians are

$$H_{\text{initial}} = |1\rangle\langle 1|^{\text{clock}}$$

$$H_{\text{final}} = \frac{1}{2} (I \otimes |0\rangle\langle 0|^{\text{clock}} + I \otimes |1\rangle\langle 1|^{\text{clock}} - U \otimes |1\rangle\langle 0|^{\text{clock}} + U^\dagger \otimes |0\rangle\langle 1|^{\text{clock}})$$

Note that the same Hamiltonians could have taken

$$|\psi'_i\rangle = V|0\rangle \otimes |0\rangle^{\text{clock}}$$

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Idea:

$$H(t) = \frac{1}{2} (P(t) \otimes |0\rangle\langle 0|^{\text{clock}} + Q(t) |1\rangle\langle 1|^{\text{clock}} - R(t) U \otimes |1\rangle\langle 0|^{\text{clock}} - R(t) U^\dagger \otimes |0\rangle\langle 1|^{\text{clock}})$$

$$\text{The ground state} \propto R(t) |0\rangle |0\rangle^{\text{clock}} + P(t) U |0\rangle |1\rangle^{\text{clock}}$$

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$$H(t) = \frac{1}{2} \left(P(t) \otimes |0\rangle\langle 0| + Q(t) \otimes |1\rangle\langle 1| - R(t) U \otimes |1\rangle\langle 0|^{clock} - R(t) U^\dagger \otimes |0\rangle\langle 1|^{clock} \right)$$

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Note that we can even set $PQ = R^2$ and keep the gap constant!

This works for multiple gates as well.

Let $|t\rangle^c$ denote the clock states at time step t ,
i.e. $|\chi_t\rangle = U_t \dots U_1 |0\rangle \otimes |t\rangle^c$

Suppose $H(s) = (1-s)H_i + sH_f$ takes

$$|\psi_i\rangle = |0\rangle \otimes |0\rangle^c$$

to $|\chi_f\rangle = \frac{1}{\sqrt{L+1}} \sum_{t=0}^L U_t \dots U_0 |0\rangle \otimes |t\rangle^c$

Now construct $H'(s) = (1-s)H'_i + sH'_f$ corresponding to circuit $U_L^+, U_{L-1}^+, \dots$. It would take

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Up to relabelling of the clock, $|\psi_f'\rangle = |\psi_f\rangle$

$\Rightarrow$ If we implement the evolution

$$H_i \rightarrow H_f (H_f') \rightarrow H_i'$$

We can take $|0\rangle$ to $U_L U_{L-1} \dots |0\rangle$, exactly the desired output.

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 "half" of Holonomic QC. This also tells us
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 there is much room for manipulation of
 gap size.

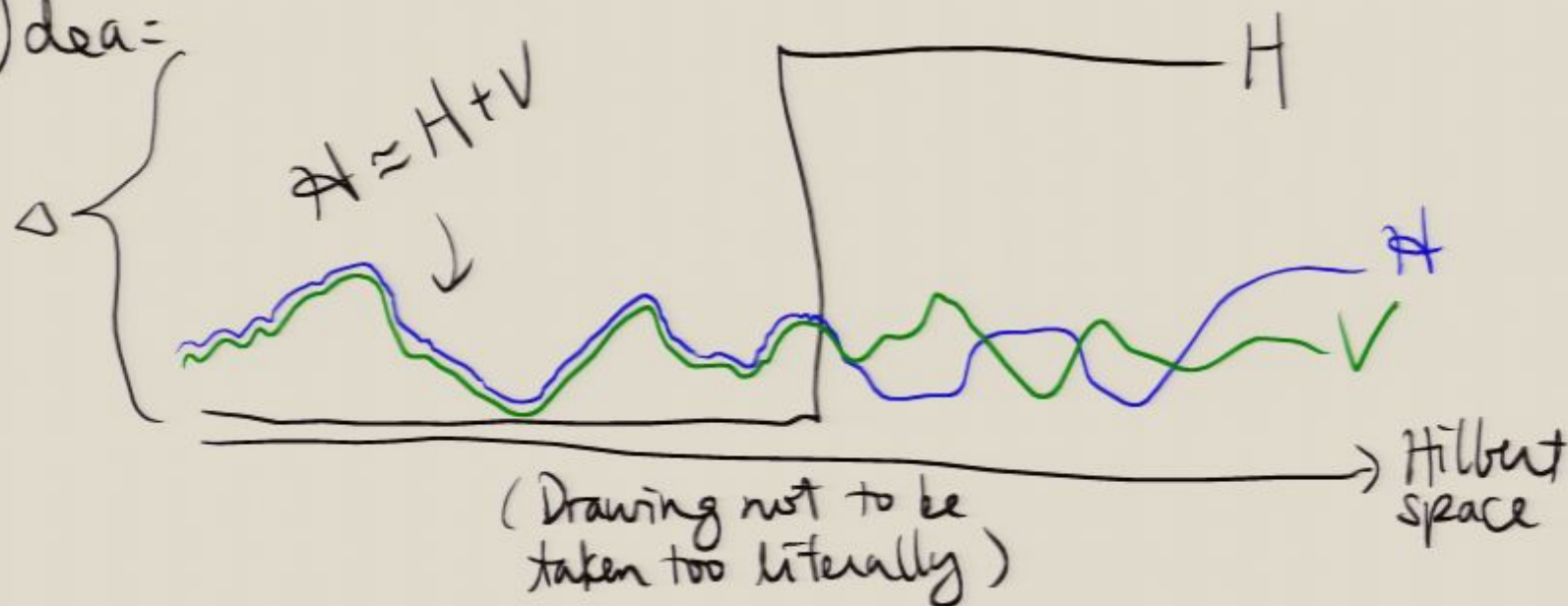
5. Perturbation Theory Gadgets

Interactions in Nature tends to be between nearest neighbors, i.e. they are 2-local and fit in a lattice. Can we implement AQC/HQC using 2-local Hamiltonians?

Yes, $\exists$ approximations that reduces nonlocality and preserves approximately both eigenvalues & eigenvectors.

Price we pay = Gap reduced by $\text{poly}(n)$ factor

Basic Idea =



Using ancillary qubits, we can often replace k -local Hamiltonian H by less-than- k -local Hamiltonians $H+V$, where H has a large gap Δ and ground state space $\mathcal{L}_-$ (write total Hilbert space $\mathcal{L} = \mathcal{L}_+ \oplus \mathcal{L}_-$) $\Rightarrow V|_{\mathcal{L}_-} \approx H$,

Theorem (roughly stated):

Let spectrum of $H \in [a, b]$. $\Delta/2 > a, b$

For $z \in [a, b]$, if $\Sigma_{--}(z)$ (defined below) is close to H ,
low energy eigenvalues and eigenvectors of $H+V$
are close to H .

$$\Sigma_{--} \equiv H_{--} + V_{--} + V_{-+} G_{++} V_{+-} + V_{-+} G_{++} V_{++} G_{++} V_{+-} + \dots$$

where "+-" denotes projections to L_+ , L_- etc.

$$G(z) = \frac{1}{z - H}$$

Remark: Why Σ_{--} and not just $H_{--} + V_{--}$?

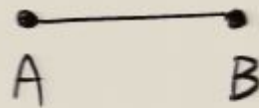
Because L_- , L_+ are eigenspaces of H ,

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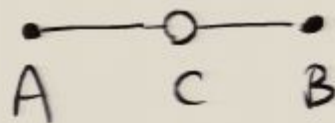
Remark: Why Σ_- and not just $H_- + V_-$?
Because L_- , L_+ are eigenspaces of H ,
not V

Simple example
of a gadget



H acts on AB

↓ approximation

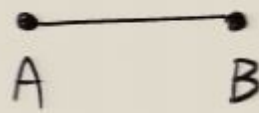


$H + V$, H acts on C ,

V acts on AC and CB

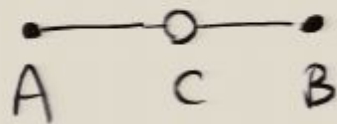
$$H = -(P_A + P_B)^2/2 \quad (\text{meaning } -(P_A \otimes I_B + I_A \otimes P_B)^2/2)$$

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$$H = \Delta |1\rangle\langle 1|_c \quad (\text{on qubit } c)$$

$$V = \sqrt{\frac{\Delta}{2}} (P_A + P_B) \otimes (|1\rangle\langle 0| + |0\rangle\langle 1|)_c$$

From these we get

$$G_{++}(z) = \frac{|1\rangle\langle 1|_c}{z - \Delta}, \quad H_{--} + V_{--} = 0,$$

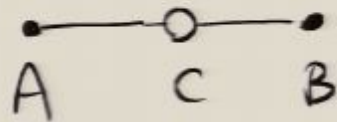
$$V_{+-} = \sqrt{\frac{\Delta}{2}} (P_A + P_B) \otimes |1\rangle\langle 0|_c$$

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$$\Sigma_{--} = \frac{\Delta}{2(z - \Delta)} (P_A + P_B)^2 \otimes |0\rangle\langle 0|_c + \mathcal{O}(\Delta^{-\frac{1}{2}})$$

we can choose Δ such that Σ_{--} is close to $\mathbb{I}$ on A, B .

Examples of other gadgets (changing locality/topology)

actions on A and C

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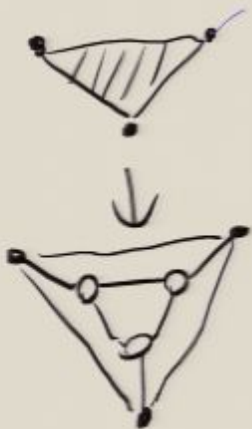
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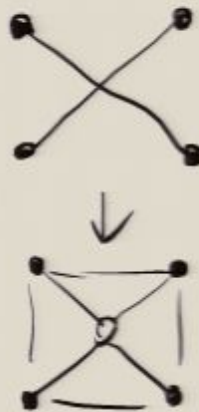
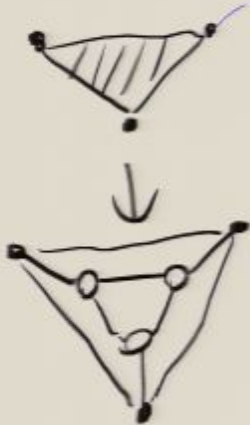


Gadgets can be used on the AQC/HQC

$$\Sigma_{--} = \frac{\Delta}{2(\beta - \Delta)} (P_A + P_B)^2 \otimes |0\rangle\langle 0|_c + \mathcal{O}(\Delta^{-\frac{1}{2}})$$

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Examples of other gadgets (changing locality/topology)



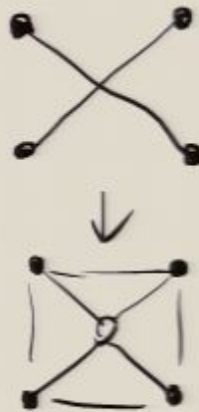
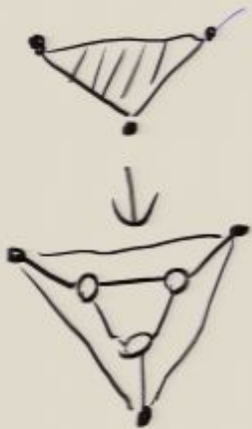
Gadgets can be used on the AQC/HQC

Hamiltonians so that they become 2-local.

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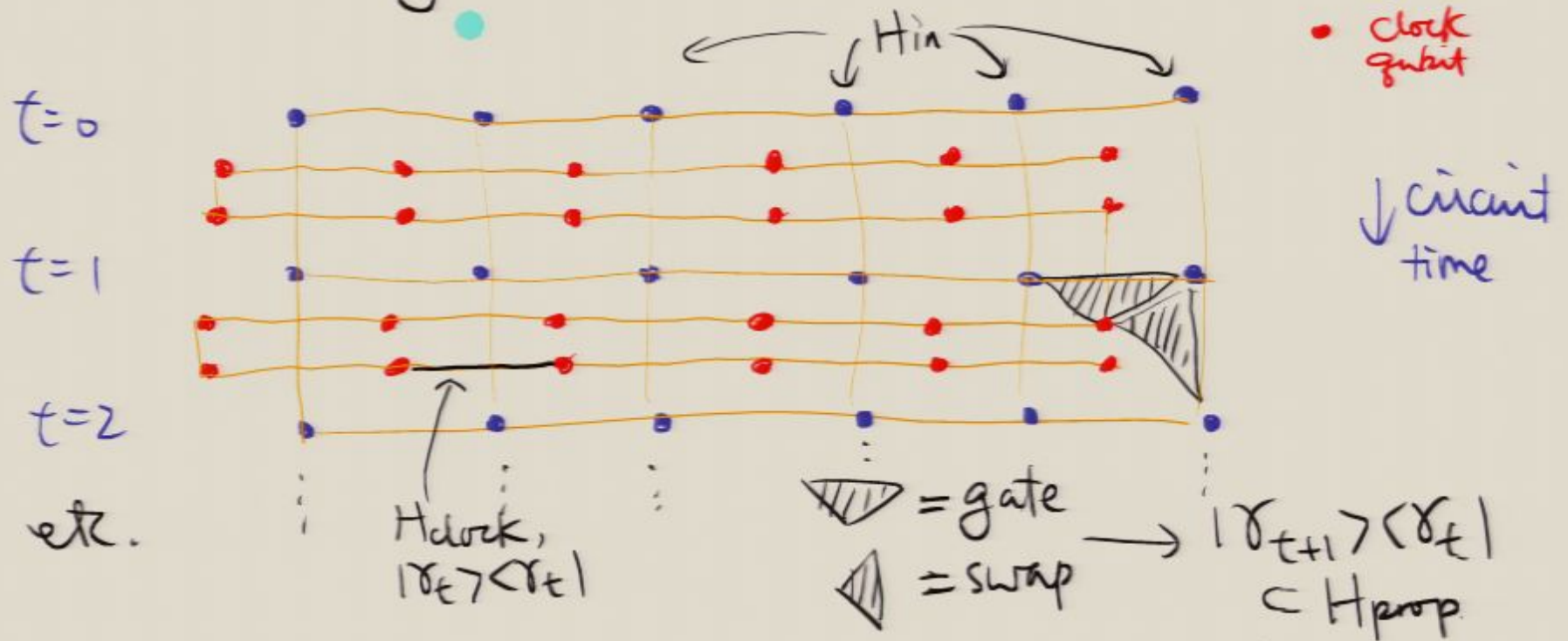
Examples of other gadgets (changing locality/topology)



Gadgets can be used on the AQC/HQC Hamiltonians so that they become 2-local.

Square Lattice construction (Oliveira & Terhal)

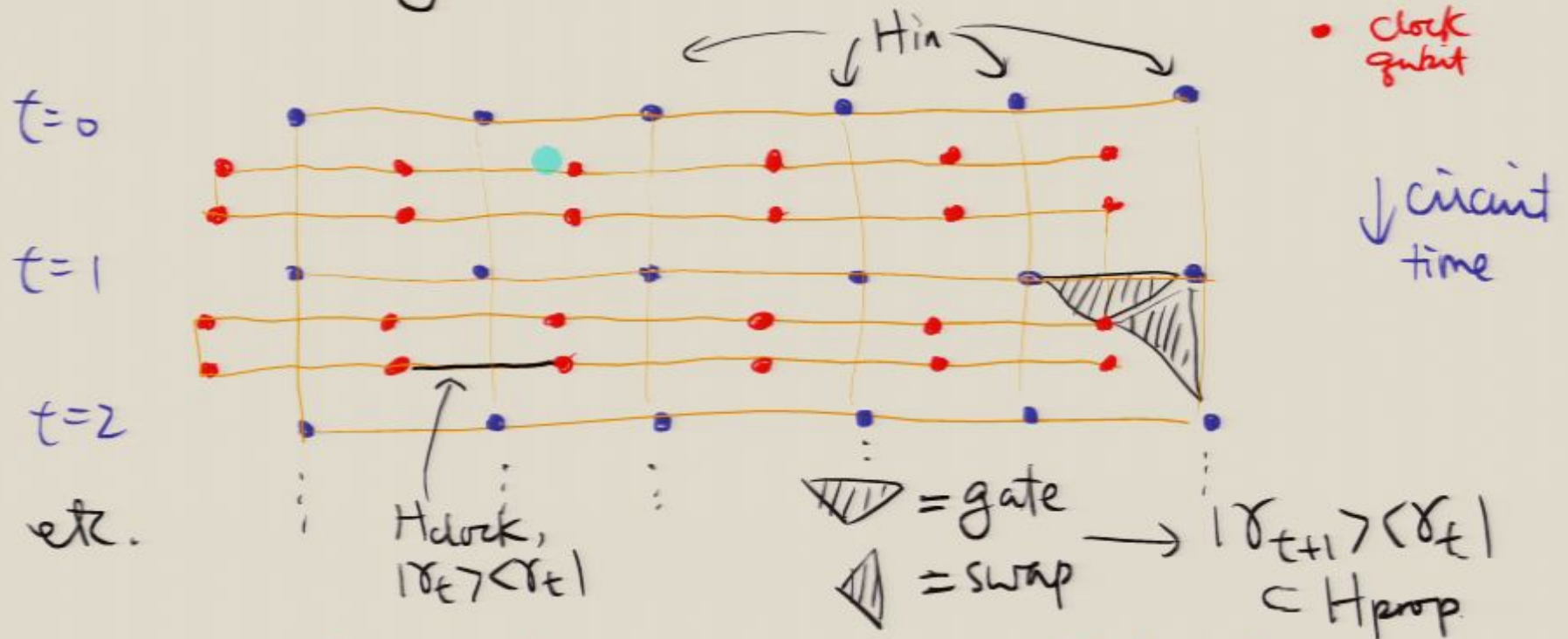
$H_f = H_{in} + \sum H_{prop}$ etc. can be implemented
3-locally on a square lattice of spins (qubits),
w/ one axis being circuit time.



Gadgets used on this can reduce to 2-local neighboring interaction

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Gadgets used on this can reduce to 2-local neighboring interaction

Adiabatic Algorithm:

$$H_{\text{initial}} = H_{\text{in}} + \underbrace{11}_{\substack{\text{clock} \\ t=0}} > < \underbrace{11}_{\substack{\text{clock} \\ t=0}} + H_{\text{clock}} \quad t \neq 0$$

↑ make sure
clock states are
legal
↓

$$H_{\text{final}} = H_{\text{in}} + \sum H_{\text{prop}} + H_{\text{clock}}$$

In the basis $\{|\sigma_0\rangle \dots |\sigma_L\rangle\}$, this looks like

$$H_{\text{initial}} = \begin{bmatrix} 0 & 1 & & \\ & 1 & \ddots & \\ & & \ddots & 1 \\ & & & 0 \end{bmatrix}$$

$$H_{\text{final}} = \frac{1}{2} \begin{bmatrix} 1 & -1 & & \\ -1 & 2 & & \\ & & \ddots & \\ & & & 2 & -1 \\ & & & -1 & 1 \end{bmatrix}$$

→ The key to making the Hamiltonian k -local, where k is small, is finding good k -local approx. to the terms $|\gamma_{t+1}\rangle\langle\gamma_t|$ etc.

$$\text{unary clock: } |t\rangle^{\text{clock}} = |\underbrace{111\dots}_t 000\dots\rangle$$

3-local example:

$$|\gamma_{t+1}\rangle\langle\gamma_t| \approx U_{t+1} \otimes |1\rangle_{t+1}^{\text{clock}} \langle 0|_{t+1}^{\text{clock}} + J |0\rangle_{t,t+1}^{\text{clock}} \langle 0|_{t,t+1}^{\text{clock}} \quad \text{for large } J$$

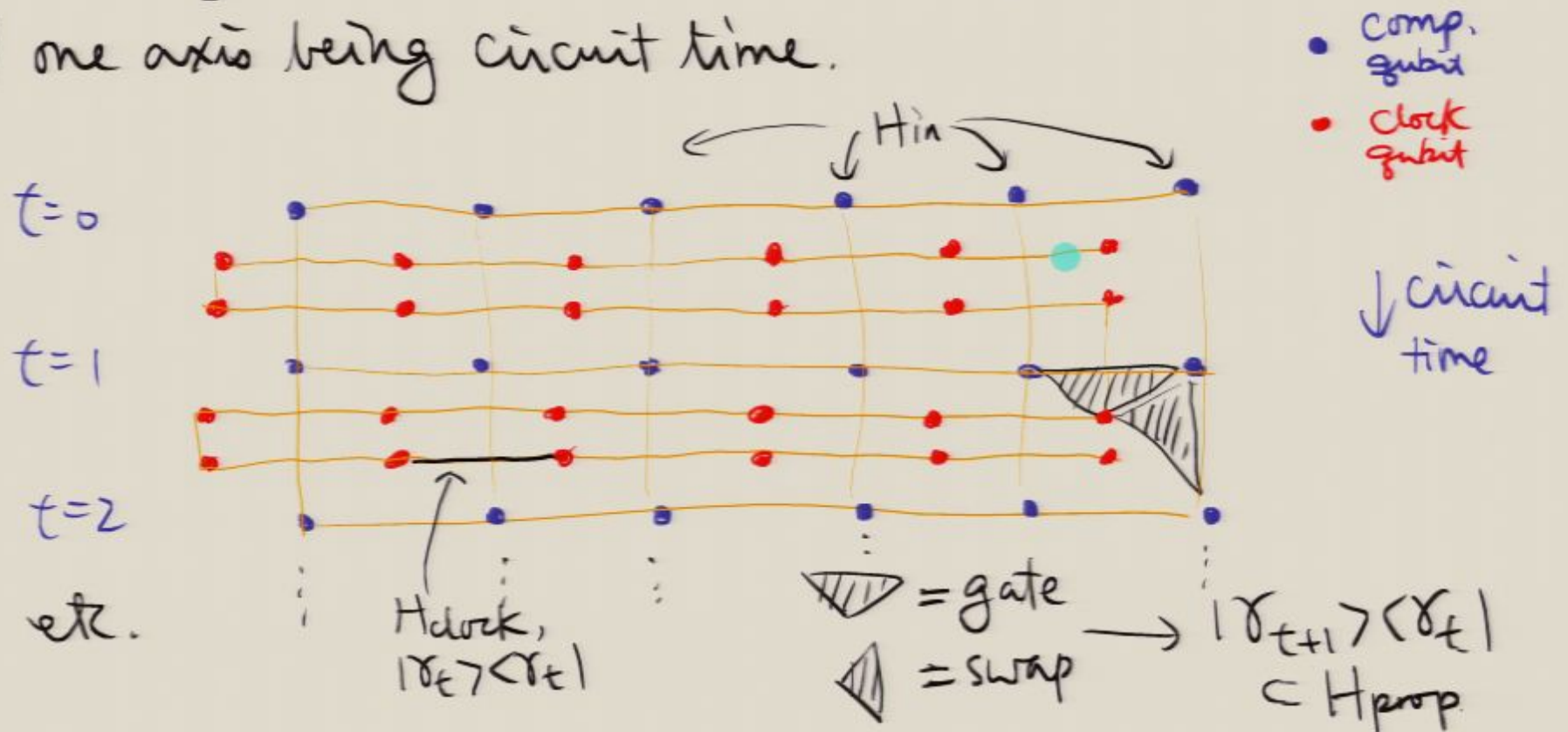
What is the gap of $H_{\text{in}} + \sum H_{\text{prop}}$?

Both H_{in} and $\sum H_{\text{prop}}$ have very simple spectra
if both have gaps $\geq \lambda$, then

gap of $H_{\text{in}} + \sum H_{\text{prop}}$

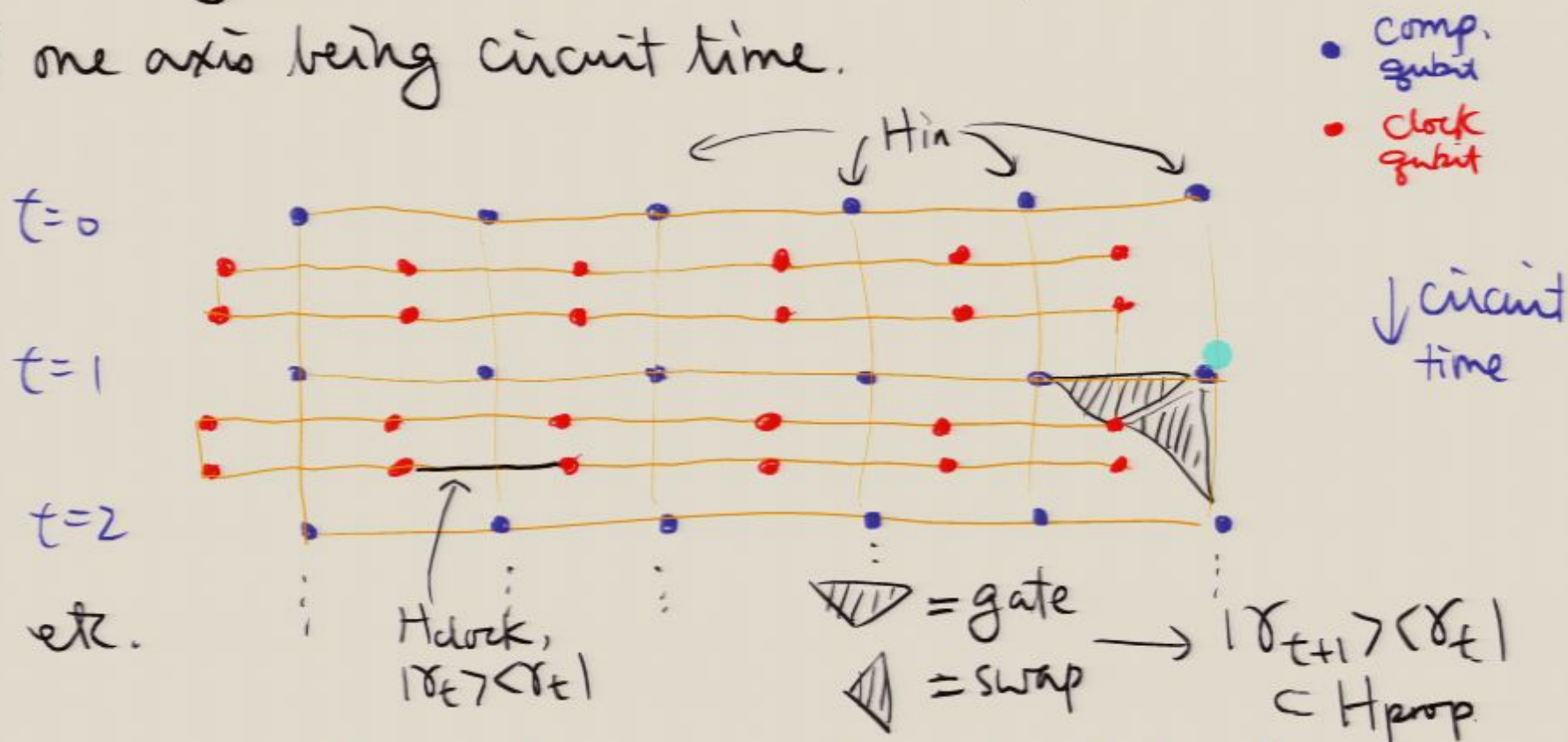
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